Project 1--CDA 3101 (Spring 2014)

Worth: 100 points (10% of course grade)

Assigned: Friday, Jan 24, 2014

bits 24-22: opcode

bits 21-19: reg A bits 18-16: reg B

bits 15-0: offsetField (16-bit, range of -32768 to 32767)

O-type instructions (halt, noop):

bits 24-22: opcode

bits 21-0: unused (should all be 0)

Table 1: Description of Machine Instructions

Assembly language Opcode in binary Action name for instruction (bits 24, 23, 22)

add (R-type format) 000 add contents of regA with contents of regB, store results in destReg.

nand (R-type format) 001 nand contents of regA with contents of regB, store results in destReg.

sw (I-type format) 011 store regB into memory. Memory address is formed by adding offsetField with the contents of regA.

beq (I-type format)

100

if the contents of regA and regB are the same, then branch to the address PC+1+offsetField, where PC is the address of the beg instruction.

cmov (R-type) 101 copy the value regA into destReg if the contents of regB != 0

halt (O-type format) 110 increment the PC (as with all instructions), then halt the machine (let the simulator notice that the machine halted).

noop (O-type format) 111 do nothing except increment PC.

4. LC3101 Assembly Language and Assembler (40%)

The first part of this project is to write a program to take an assembly-language program and translate it into machine language. You will

translate assembly-language names for instructions, such as beq, into their

numeric equivalent (e.g. 100), and you will translate symbolic names for addresses into numeric values. The final output will be a series of 32-bit

instructions (instruction bits 31-25 are always 0).

The format for a line of assembly code is:

label instruction field0 field1 field2 comments

The leftmost field on a line is the label field. Valid labels contain a maximum of 6 characters and can consist of letters and numbers (but must start

with a letter). The label is optional (the white space following the

field is required). Labels make it much easier to write assembly-language

programs, since otherwise you would need to modify all address fields each time

you added a line to your assembly-language program!

After the optional label is white space which consists of any number of space

or tab characters. The writespace is followed by the instruction field, where the instruction can be any of the assembly-language instruction names

listed in the above table. After more white space comes a series of fields.

All fields are given as decimal numbers or labels. The number of fields depends on the instruction, and unused fields should be ignored (treat them

like comments).

R-type instructions (add, nand) instructions require 3 fields: field0 is regA, field1 is regB, and field2 is destReg.

I-type instructions (lw, sw, beq) require 3 fields: field0 is regA, field1

is regB , and field2 is either a numeric value for offsetField or a symbolic

address. Numeric offsetFields can be positive or negative; symbolic addresses are discussed below.

O-type instructions (noop and halt) require no fields.

Symbolic addresses refer to labels. For lw or sw instructions, the assembler

should compute offsetField to be equal to the address of the label. This could

be used with a zero base register to refer to the label, or could be used with

a non-zero base register to index into an array starting at the label. For beg

instructions, the assembler should translate the label into the numeric offsetField needed to branch to that label.

After the last used field comes more white space, then any comments. The comment field ends at the end of a line. Comments are vital to creating understandable assembly-language programs, because the instructions themselves

are rather cryptic.

In addition to LC3101 instructions, an assembly-language program may contain

directions for the assembler. The only assembler directive we will use is .fill

(note the leading period). .fill tells the assembler to put a number into the

place where the instruction would normally be stored. .fill instructions use

one field, which can be either a numeric value or a symbolic address. For

example, ".fill 32" puts the value 32 where the instruction would normally be

stored. .fill with a symbolic address will store the address of the label.

In the example below, ".fill start" will store the value 2, because the label

"start" is at address 2.

The assembler should make two passes over the assembly-language program. In the

first pass, it will calculate the address for every symbolic label. Assume

that the first instruction is at address 0. In the second pass, it will generate a machine-language instruction (in decimal) for each line of assembly

language. For example, here is an assembly-language program (that counts down

from 5, stopping when it hits 0).

lw 1 five load reg1 with 5 (uses symbolic address) lw 1 2 3 load reg2 with -1 (uses numeric address) 2 start add 1 1 decrement reg1 1 beq 0

neg1 .fill -1
stAddr .fill start

Hints: the example assembly-language program above is a good case to include

in your test suite, though you'll need to write more test cases to get full

credit. Remember to create some test cases that test the ability of an assembler to check for the errors in Section 4.2.

4.4. Assembler Hints

Since offsetField is a 2's complement number, it can only store numbers ranging

from -32768 to 32767. For symbolic addresses, your assembler will compute

offsetField so that the instruction refers to the correct label.

Remember that offsetField is only an 16-bit 2's complement number. Since most.

machines you run your assembler on have 32-bit or 64-bit integers, you will

have to truncate all but the lowest 16 bits for negative values of offsetField.

5. Behavioral Simulator (40%)

The second part of this assignment is to write a program that can simulate any

legal LC3101 machine-code program. The input for this part will be the machine-code file that you created with your assembler. With a program name

run as follows:

simulate program.mc > output

This directs all print statements to the file "output".

As with the assembler, you will write a suite of test cases to validate the LC3101 simulator.

The test cases for the simulator part of this project will be short assembly-language programs that, after being assembled into machine code, serve as input to a simulator. You will submit your suite of test cases together with your s

program halts. You may assume that the two input numbers are at most 15 bits

and are positive; this ensures that the (positive) result fits in an ${\tt LC3101}$

word. See the algorithm on page 252 of the textbook for how to multiply. Remember that shifting left by one bit is the same as adding the number to

itself. Given the LC3101 instruction set, it's easiest to modify the algorithm so that you avoid the right shift. Submit a version of the program ${}^{\prime}$

that computes (32766 * 10383).

Your multiplication program must be reasonably efficient—it must be at most

50 lines long and execute at most 1000 instructions for any valid numbers (this

is several times longer and slower than the solution). To achieve this, you

must use a loop and shift algorithm to perform the multiplication; algorithms

such as successive addition (e.g. multiplying 5 * 6 by adding 5 six times)

will take too long.

7. Grading and Formatting

We will grade primarily on functionality, including error handling, correctly

assembling and simulating all instructions, input and output format, method of

executing your program, correctly multiplying, and comprehensiveness of the

test suites.

The best way to debug your program is to generate your own test cases, figure

out the correct answers, and compare your program's output to the correct answers. This is also one of the best ways to learn the concepts in the project.

The student suite of test cases for the assembler and simulator parts of

project will be graded according to how thoroughly they test an LC3101 assembler or simulator. We will judge thoroughness of at most suite by how

well it exposes potentially bugs in an assembler or simulator.

For the assembler test suite, we will use each test case as input to a set $\ensuremath{\mathsf{S}}$

of buggy assemblers. A test case exposes a buggy assembler by causing it to generate a different answer from a correct assembler. The test suite is graded based on how many of the buggy assemblers were exposed by at least one test case. Tat is known as "mutation testing" in the research literature on automated testing.

For the simulator test suite, we will correctly assemble each test case, then use it as input to a set of buggy simulators. A test case exposes a buggy simulator by causing it to generate a different answer from a correct simulator. The test suite is graded based on how many of the buggy

simulators were exposed by at least one test case.

8. Turning in the Project

Submit you files through blackboard.

Each part should be archived in a .tar or .zip file to help with grading.

Here are the files you should submit for each project part:

- 1) assembler (part 1a)
 - a. C/C++ program for your assembler
- b. suite of test cases (each test case is an assembly-language $\operatorname{program}$

in a separate file)

- 2) simulator (part 1s)
 - a. C/C++ program for your simulator
- $\ensuremath{\text{b.}}$ suite of test cases (each test case is an assembly-language program

in a separate file)

- 3) multiplication (part 1m)
 - a. assembly program for multiplication

Your assembler and simulator must each be in a single ${\tt C}$ or ${\tt C++}$ file. We will compile

your program on linprog using "gcc program.c -lm" (or g++), so your program

should not require additional compiler flags or libraries.

The official time of submission for your project will be the time the last file

is sent. If you send in anything after the due date, your project will be considered late (and will use up your late days or will receive a zero).

9. Code Fragment for Assembler

The focus of this class is machine organization, not ${\tt C}$ programming skills. To

"build" your computer, however, you will be doing a lot of C programming. To

help you, here is a fragment of the C program for the assembler. This shows

how to specify command-line arguments to the program (via argc and argv), how

to parse the assembly-language file, etc.. This fragment is provided strictly

to help you, though it may take a bit for you to understand and use the file.

```
You may also choose to not use this fragment.
/* Assembler code fragment for LC3101 */
#include <stdlib.h>
#include <stdio.h>
#include <string.h>
#define MAXIINELENGTH 1000
int readAndParse(FILE *, char *, char *, char *, char *, char *);
int isNumber(char *);
int
main(int argc, char *argv[])
    char *inFileString, *outFileString;
    FILE *inFilePtr, *outFilePtr;
    char label[MAXLINELENGTH], opcode[MAXLINELENGTH],
arg0[MAXLINELENGTH],
            arg1[MAXLINELENGTH], arg2[MAXLINELENGTH];
    if (argc != 3) {
        printf("error: usage: %s <assembly-code-file> <machine-code-</pre>
file>\n",
            arqv[0]);
        exit(1);
    }
    inFileString = argv[1];
    outFileString = argv[2];
    inFilePtr = fopen(inFileString, "r");
    if (inFilePtr == NULL) {
        printf("error in opening %s\n", inFileString);
        exit(1);
    outFilePtr = fopen(outFileString, "w");
    if (outFilePtr == NULL) {
        printf("error in opening %s\n", outFileString);
        exit(1);
    }
    /* here is an example for how to use readAndParse to read a line from
        inFilePtr */
    if (! readAndParse(inFilePtr, label, opcode, arg0, arg1, arg2) ) {
        /* reached end of file */
    }
    /* this is how to rewind the file ptr so that you start reading from
the
        beginning of the file */
    rewind(inFilePtr);
```

```
/* after doing a readAndParse, you may want to do the following to
test the
        opcode */
    if (!strcmp(opcode, "add")) {
        /* do whatever you need to do for opcode "add" */
    }
    return(0);
}
 * Read and parse a line of the assembly-language file. Fields are
returned
 * in label, opcode, arg0, arg1, arg2 (these strings must have memory
already
 * allocated to them).
 * Return values:
       0 if reached end of file
       1 if all went well
 * exit(1) if line is too long.
 */
int
readAndParse(FILE *inFilePtr, char *label, char *opcode, char *arg0,
    char *arg1, char *arg2)
{
    char line[MAXLINELENGTH];
    char *ptr = line;
    /* delete prior values */
    label[0] = opcode[0] = arg0[0] = arg1[0] = arg2[0] = '\0';
    /* read the line from the assembly-language file */
    if (fgets(line, MAXLINELENGTH, inFilePtr) == NULL) {
     /* reached end of file */
        return(0);
    /* check for line too long (by looking for a \n) */
    if (strchr(line, '\n') == NULL) {
        /* line too long */
     printf("error: line too long\n");
     exit(1);
    /* is there a label? */
    ptr = line;
    if (sscanf(ptr, "%[^{tn}]", label)) {
     /* successfully read label; advance pointer over the label */
        ptr += strlen(label);
    }
    /*
```

```
* Parse the rest of the line. Would be nice to have real regular
              * expressions, but scanf will suffice.
           sscanf(ptr, "%*[\t\n ]%[^\t\n ]%*[\t\n ]%[^\t\n ]%[^\t\
]%*[\t\n ]%[^\t\n ]",
                      opcode, arg0, arg1, arg2);
           return(1);
}
int
isNumber(char *string)
           /* return 1 if string is a number */
           int i;
           return( (sscanf(string, "%d", &i)) == 1);
}
10. Code Fragment for Simulator
Here is some C code that may help you write the simulator. Again, you
should
take this merely as a hint. You may have to re-code this to make it do
exactly
what you want, but this should help you get started. Remember not to
change stateStruct or printState.
/* instruction-level simulator for LC3101 */
#include <stdio.h>
#include <string.h>
#define NUMMEMORY 65536 /* maximum number of words in memory */
#define NUMREGS 8 /* number of machine registers */
#define MAXLINELENGTH 1000
typedef struct stateStruct {
           int pc;
           int mem[NUMMEMORY];
           int reg[NUMREGS];
           int numMemory;
} stateType;
void printState(stateType *);
int
main(int argc, char *argv[])
           char line[MAXLINELENGTH];
           stateType state;
           FILE *filePtr;
           if (argc != 2) {
```

```
printf("error: usage: %s <machine-code file>\n", argv[0]);
     exit(1);
    filePtr = fopen(argv[1], "r");
    if (filePtr == NULL) {
     printf("error: can't open file %s", argv[1]);
     perror("fopen");
     exit(1);
    /* read in the entire machine-code file into memory */
    for (state.numMemory = 0; fgets(line, MAXLINELENGTH, filePtr) !=
NULL;
     state.numMemory++) {
     if (sscanf(line, "%d", state.mem+state.numMemory) != 1) {
          printf("error in reading address %d\n", state.numMemory);
          exit(1);
     printf("memory[%d]=%d\n", state.numMemory,
state.mem[state.numMemory]);
    return(0);
}
void
printState(stateType *statePtr)
    int i;
    printf("\n@@@\nstate:\n");
    printf("\tpc %d\n", statePtr->pc);
    printf("\tmemory:\n");
     for (i=0; i<statePtr->numMemory; i++) {
          printf("\t\tmem[ %d ] %d\n", i, statePtr->mem[i]);
    printf("\tregisters:\n");
     for (i=0; i<NUMREGS; i++) {</pre>
         printf("\t\treg[ %d ] %d\n", i, statePtr->reg[i]);
    printf("end state\n");
}
11. Programming Tips
Here are a few programming tips for writing C/C++ programs to manipulate
bits:
1) To indicate a hexadecimal constant in, precede the number by 0x. For
example, 27 decimal is 0x1b in hexadecimal.
```

2) The value of the expression (a >> b) is the number "a" shifted right

by "b"

bits. Neither a nor b are changed. E.g. (25 >> 2) is 6. Note that 25 is 11001 in

binary, and 6 is 110 in binary.

3) The value of the expression (a << b) is the number "a" shifted left by "b"

bits. Neither a nor b are changed. E.g. (25 << 2) is 100. Note that 25 is 11001

in binary, and 100 is 1100100 in binary.

4) To find the value of the expression (a & b), perform a logical AND on each

bit of a and b (i.e. bit 31 of a ANDED with bit 31 of b, bit 30 of a ANDED with

bit 30 of b, etc.). E.g. (25 & 11) is 9, since:

11001 (binary)

& 01011 (binary)

- = 01001 (binary), which is 9 decimal.
- 5) To find the value of the expression (a \mid b), perform a logical OR on each bit

of a and b (i.e. bit 31 of a ORED with bit 31 of b, bit 30 of a ORED with bit 30 $\,$

of b, etc.). E.g. (25 | 11) is 27, since:

11001 (binary) & 01011 (binary)

a 01011 (201121)

- = 11011 (binary), which is 27 decimal.
- 6) ~a is the bit-wise complement of a (a is not changed).

Use these operations to create and manipulate machine-code. E.g. to look at bit

- 3 of the variable a, you might do: (a>>3) & 0x1. To look at bits (bits 15-12) of
- a 16-bit word, you could do: (a>>12) & 0xF. To put a 6 into bits 5-3 and a 3 $\,$

into bits 2-1, you could do: $(6 << 3) \mid (3 << 1)$. If you're not sure what an operation is doing, print some intermediate results to help you debug.

12. Example Run of Simulator

memory[0] = 8454151

memory[1] = 9043971

memory[2] = 655361

memory[3]=16842754

memory[4]=16842749

memory[5] = 29360128

memory[6] = 25165824

```
memory[7]=5
memory[8]=-1
memory[9]=2
999
state:
     pc 0
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[ 2 ] 655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[6]25165824
           mem[ 7 ] 5
           mem[8]-1
           mem[ 9 ] 2
     registers:
           reg[ 0 ] 0
           reg[ 1 ] 0
           reg[ 2 ] 0
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
000
state:
     pc 1
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[2]655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[ 8 ] -1
           mem[ 9 ] 2
     registers:
           reg[ 0 ] 0
           reg[ 1 ] 5
           reg[ 2 ] 0
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
```

```
@ @ @
state:
     pc 2
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[2]655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[8]-1
           mem[ 9 ] 2
     registers:
           reg[ 0 ] 0
           reg[ 1 ] 5
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
999
state:
     рс 3
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[2]655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[8]-1
           mem[ 9 ] 2
     registers:
           reg[ 0 ] 0
           reg[ 1 ] 4
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
@ @ @
state:
     pc 4
```

```
memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[2]655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[8]-1
           mem[ 9 ] 2
     registers:
           reg[ 0 ] 0
           reg[ 1 ] 4
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
999
state:
     pc 2
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[ 2 ] 655361
           mem[3]16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[ 8 ] -1
           mem[ 9 ] 2
     registers:
           reg[ 0 ] 0
           reg[ 1 ] 4
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
999
state:
     pc 3
     memory:
           mem[ 0 ] 8454151
           mem[ 1 ] 9043971
           mem[ 2 ] 655361
```

```
mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[ 8 ] -1
           mem[ 9 ] 2
     registers:
           reg[ 0 ] 0
           reg[ 1 ] 3
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
999
state:
     pc 4
     memory:
           mem[ 0 ] 8454151
           mem[ 1 ] 9043971
           mem[2]655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[ 8 ] -1
           mem[ 9 ] 2
     registers:
           reg[ 0 ] 0
           reg[ 1 ] 3
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
999
state:
     pc 2
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[ 2 ] 655361
           mem[3]16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
```

```
mem[ 7 ] 5
           mem[ 8 ] -1
           mem[ 9 ] 2
      registers:
           reg[ 0 ] 0
           reg[ 1 ] 3
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
000
state:
     рс 3
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[ 2 ] 655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[ 8 ] -1
           mem[ 9 ] 2
      registers:
           reg[ 0 ] 0
           reg[ 1 ] 2
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
999
state:
     pc 4
     memory:
           mem[ 0 ] 8454151
           mem[ 1 ] 9043971
           mem[ 2 ] 655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[7]5
           mem[ 8 ] -1
           mem[ 9 ] 2
      registers:
```

```
reg[ 0 ] 0
           reg[ 1 ] 2
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
000
state:
     pc 2
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[ 2 ] 655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[ 8 ] -1
           mem[ 9 ] 2
      registers:
           reg[ 0 ] 0
           reg[ 1 ] 2
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
999
state:
     рс 3
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[ 2 ] 655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[7]5
           mem[8]-1
           mem[ 9 ] 2
      registers:
           reg[ 0 ] 0
           reg[ 1 ] 1
           reg[ 2 ] -1
           reg[ 3 ] 0
```

```
reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
000
state:
     pc 4
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[2]655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[ 8 ] -1
           mem[ 9 ] 2
      registers:
           reg[ 0 ] 0
           reg[ 1 ] 1
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
999
state:
     pc 2
     memory:
           mem[ 0 ] 8454151
           mem[ 1 ] 9043971
           mem[2]655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[7]5
           mem[ 8 ] -1
           mem[ 9 ] 2
      registers:
           reg[ 0 ] 0
           reg[ 1 ] 1
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
```

```
end state
999
state:
     pc 3
     memory:
           mem[0] 8454151
           mem[ 1 ] 9043971
           mem[2]655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[ 8 ] -1
           mem[ 9 ] 2
      registers:
            reg[ 0 ] 0
            reg[ 1 ] 0
           reg[ 2 ] -1
            reg[ 3 ] 0
            reg[ 4 ] 0
            reg[ 5 ] 0
           reg[ 6 ] 0
            reg[ 7 ] 0
end state
<u>a</u> a a
state:
     pc 6
     memory:
           mem[ 0 ] 8454151
           mem[ 1 ] 9043971
           mem[ 2 ] 655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[7]5
           mem[ 8 ] -1
           mem[ 9 ] 2
      registers:
           reg[ 0 ] 0
            reg[ 1 ] 0
            reg[ 2 ] -1
            reg[ 3 ] 0
            reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
machine halted
total of 17 instructions executed
final state of machine:
```

```
000
state:
     pc 7
     memory:
           mem[ 0 ] 8454151
           mem[ 1 ] 9043971
           mem[ 2 ] 655361
           mem[ 3 ] 16842754
           mem[ 4 ] 16842749
           mem[ 5 ] 29360128
           mem[ 6 ] 25165824
           mem[ 7 ] 5
           mem[ 8 ] -1
           mem[ 9 ] 2
     registers:
           reg[ 0 ] 0
           reg[ 1 ] 0
           reg[ 2 ] -1
           reg[ 3 ] 0
           reg[ 4 ] 0
           reg[ 5 ] 0
           reg[ 6 ] 0
           reg[ 7 ] 0
end state
```