

## Project 1--CDA 3101 (Spring 2014)

Worth: 100 points (10% of course grade)

Assigned: Friday, Jan 24, 2014

Due: 1:25 pm, Monday, Feb 24, 2014

### 1. Purpose

This project is intended to help you understand the instructions of a very simple assembly language and how to assemble programs into machine language.

### 2. Problem

This project has three parts. In the first part, you will write a program to take an assembly-language program and produce the corresponding machine language. In the second part, you will write a behavioral simulator for the resulting machine code. In the third part, you will write a short assembly-language program to multiply two numbers.

### 3. LC3101 Instruction-Set Architecture

For this project, you will be developing a simulator and assembler for the LC3101 (Little Computer, used in CDA 3101). The LC3101 is very simple, but it is general enough to solve complex problems. For this project, you will only need to know the instruction set and instruction format of the LC3101.

The LC3101 is an 8-register, 32-bit computer. All addresses are word-addresses (unlike MIPS which is byte-addressed). The LC3101 has 65536 words of memory. By assembly-language convention, register 0 will always contain 0 (i.e. the machine will not enforce this, but no assembly-language program should ever change register 0 from its initial value of 0).

There are 3 instruction formats (bit 0 is the least-significant bit). Bits 31-25 are unused for all instructions, and should always be 0.

R-type instructions (add, nand):

- bits 24-22: opcode
- bits 21-19: reg A
- bits 18-16: reg B
- bits 15-3: unused (should all be 0)
- bits 2-0: destReg

I-type instructions (lw, sw, beq):

bits 24-22: opcode  
 bits 21-19: reg A  
 bits 18-16: reg B  
 bits 15-0: offsetField (16-bit, range of -32768 to 32767)

O-type instructions (halt, noop):  
 bits 24-22: opcode  
 bits 21-0: unused (should all be 0)

Table 1: Description of Machine Instructions

Assembly language name for instruction	Opcode in binary (bits 24, 23, 22)	Action
add (R-type format)	000	add contents of regA with contents of regB, store results in destReg.
nand (R-type format)	001	nand contents of regA with contents of regB, store results in destReg.
lw (I-type format)	010	load regB from memory. Memory address is formed by adding offsetField with the contents of regA.
sw (I-type format)	011	store regB into memory. Memory address is formed by adding offsetField with the contents of regA.
beq (I-type format)	100	if the contents of regA and regB are the same, then branch to the address PC+1+offsetField, where PC is the address of the beq instruction.
cmov (R-type)	101	copy the value regA into destReg if the contents of regB != 0
halt (O-type format)	110	increment the PC (as with all instructions), then halt the machine (let the simulator notice that the machine halted).
noop (O-type format)	111	do nothing except increment PC.

-----  
-----  

#### 4. LC3101 Assembly Language and Assembler (40%)

The first part of this project is to write a program to take an assembly-language program and translate it into machine language. You will translate assembly-language names for instructions, such as beq, into their numeric equivalent (e.g. 100), and you will translate symbolic names for addresses into numeric values. The final output will be a series of 32-bit instructions (instruction bits 31-25 are always 0).

The format for a line of assembly code is:

```
label instruction field0 field1 field2 comments
```

The leftmost field on a line is the label field. Valid labels contain a maximum of 6 characters and can consist of letters and numbers (but must start with a letter). The label is optional (the white space following the label field is required). Labels make it much easier to write assembly-language programs, since otherwise you would need to modify all address fields each time you added a line to your assembly-language program!

After the optional label is white space which consists of any number of space or tab characters. The whitespace is followed by the instruction field, where the instruction can be any of the assembly-language instruction names listed in the above table. After more white space comes a series of fields. All fields are given as decimal numbers or labels. The number of fields depends on the instruction, and unused fields should be ignored (treat them like comments).

R-type instructions (add, nand) instructions require 3 fields: field0 is regA, field1 is regB, and field2 is destReg.

I-type instructions (lw, sw, beq) require 3 fields: field0 is regA, field1 is regB, and field2 is either a numeric value for offsetField or a symbolic address. Numeric offsetFields can be positive or negative; symbolic addresses are discussed below.

O-type instructions (noop and halt) require no fields.

Symbolic addresses refer to labels. For lw or sw instructions, the assembler should compute offsetField to be equal to the address of the label. This could be used with a zero base register to refer to 8 675.34 Tlabel, or could be used with a non-zero base register to index into an array starting at the label. For beq instructions, 8 675.34 T assembler should translate the label into the numeric offsetField needed to branch to thatTlabel.

After t5.34 Tlast used field comes more white space, 8 675.34 n any commen8 67s. The commen8 67 field ends at the end ofTaTline. Commen8 67s are vital to creating understandable assembly-language programs, because the instructions themselves are rather cryptic.

In addition to LC3101 instructions, an assembly-language program may contain directions for the assembler. The only assembler directive we will use is .fill (note t5.34 Tleading period). .fill tells the assembler totpa number into the place where t5.34 Tinstruction would normally be stored. .fill instructions use one field, which can be eitherTaTnumeric value orTaTsymbolic address. For example, ".fill 32" puts t5.34 Tvalue 32 where t5.34 Tinstruction would normally be stored. .fill with a symbolic address will store the address ofTthe label. In the example below, ".fill start" will store the value 2, because the label "start"34 Tis at address 2.

The assembler should make two passes over 8 675.34 Tassembly program. In the first pass,34 Tit will calculate the address forTevery symbolic label. Assume that the first instruction is at address 0. In the second pass,34 Tit will generate a machine-language instruction (in decimal) for each line of assembly language. For example, here is an34 Tassembly language program (thatTcoun8 67s down from 5, stopping when it hi8 67s 0).

	lw	0	1	five	load reg1 with 5 (uses symbolic address)
	lw	1	2	3	load reg2 with -1 (uses numeric address)
start	add	1	2	1	decremen8 67 reg1
	beq	0	1	2	goto end ofTprogram when 67 reg1==0
	beq	0	0	start	

```
negl .fill -1
stAddr .fill start          will contain the address of start (2)
```

And here is the corresponding machine language:

```
(address 0): 8454151 (hex 0x810007)
(address 1): 9043971 (hex 0x8a0003)
(address 2): 655361 (hex 0xa0001)
(address 3): 16842754 (hex 0x1010002)
(address 4): 16842749 (hex 0x100fffd)
(address 5): 29360128 (hex 0x1c00000)
(address 6): 25165824 (hex 0x1800000)
(address 7): 5 (hex 0x5)
(address 8): -1 (hex 0xffffffff)
(address 9): 2 (hex 0x2)
```

Be sure you understand how the above assembly-language program got translated to machine language.

Since your programs will always start at address 0, your program should only output the contents, not the addresses.

```
8454151
9043971
655361
16842754
16842749
29360128
25165824
5
-1
2
```

#### 4.1. Running Your Assembler

Write your program to take two command-line arguments. The first argument is the file name where the assembly-language program is stored, and the second argument is the file name where the output (the machine-code) is written. For example, with a program name of "assemble", an assembly-language program in "program.as", the following would generate a machine-code file "program.mc":

```
assemble program.as program.mc
```

Note that the format for running the command must use command-line arguments for the file names (rather than standard input and standard output). Your program should store only the list of decimal numbers in the machine-code

file, one instruction per line. Any deviation from this format (e.g. extra spaces or empty lines) will render your machine-code file ungradable. Any other output that you want the program to generate (e.g. debugging output) can be printed to standard output.

#### 4.2. Error Checking

Your assembler should catch the following errors in the assembly-language program: use of undefined labels, duplicate labels, offsetFields that don't fit in 16 bits, and unrecognized opcodes. Your assembler should exit(1) if it detects an error and exit(0) if it finishes without detecting any errors. Your assembler should NOT catch simulation-time errors, i.e. errors that would occur at the time the assembly-language program executes (e.g. branching to address -1, infinite loops, etc.).

#### 4.3. Test Cases

An integral (and graded) part of writing your assembler will be to write a suite of test cases to validate any LC3101 assembler. This is common practice in the real world--software companies maintain a suite of test cases for their programs and use this suite to check the program's correctness after a change. Writing a comprehensive suite of test cases will deepen your understanding of the project specification and your program, and it will help you a lot as you debug your program.

The test cases for the assembler part of this project will be short assembly-language programs that serve as input to an assembler. You will submit your suite of test cases together with your assembler, and we will grade your test suite according to how thoroughly it exercises an assembler. Each test case may be at most 50 lines long, and your test suite may contain up to 20 test cases. These limits are much larger than needed for full credit (the solution test suite is composed of 5 test cases, each < 10 lines long). See Section 7 for how your test suite will be graded.

Hints: the example assembly-language program above is a good case to include in your test suite, though you'll need to write more test cases to get full

As with the assembler, you will write a suite of test cases to validate the LC3101 simulator.

The test cases for the simulator part of this project will be short assembly-language programs that, after being assembled into machine code, serve as input to a simulator. You will submit your suite of test cases together with your simulator, and we will grade your test suite according to how thoroughly it exercises an LC3101 simulator. Each test case may execute at most 200 instructions on a correct simulator, and your test suite may contain up to 20 test cases. These limits are much larger than needed for full credit (the solution test suite is composed of a couple test cases, each executing less than 40 instructions). See Section 7 for how your test suite will be graded.

## 5.2. Simulator Hints

Be careful how you handle `offsetField` for `lw`, `sw`, and `beq`. Remember that it's a 2's complement 16-bit number, so you need to convert a negative `offsetField` to a negative 32-bit integer on the Sun workstations (by sign extending it). To do this, use the following function.

```
int
convertNum(int num)
{
    /* convert a 16-bit number into a 32-bit Sun integer */
    if (num & (1<<15) ) {
```



program halts. You may assume that the two input numbers are at most 15 bits and are positive; this ensures that the (positive) result fits in an LC3101 word. See the algorithm on page 252 of the textbook for how to multiply. Remember that shifting left by one bit is the same as adding the number to itself. Given the LC3101 instruction set, it's easiest to modify the algorithm so that you avoid the right shift. Submit a version of the program that computes  $(32766 * 10383)$ .

Your multiplication program must be reasonably efficient--it must be at most 50 lines long and execute at most 1000 instructions for any valid numbers (this is several times longer and slower than the solution). To achieve this, you must use a loop and shift algorithm to perform the multiplication; algorithms such as successive addition (e.g. multiplying  $5 * 6$  by adding 5 six times) will take too long.

## 7. Grading and Formatting

We will grade primarily on functionality, including error handling, correctly assembling and simulating all instructions, input and output format, method of executing your program, correctly multiplying, and comprehensiveness of the test suites.

The best way to debug your program is to generate your own test cases, figure out the correct answers, and compare your program's output to the correct answers. This is also one of the best ways to learn the concepts in the project.

For the simulator test suite, we will correctly assemble each test case, then use it as input to a set of buggy simulators. A test case exposes a buggy simulator by causing it to generate a different answer from a correct simulator. The test suite is graded based on how many of the buggy simulators were exposed by at least one test case.

## 8. Turning in the Project

Submit your files through blackboard.  
Each part should be archived in a .tar or .zip file to help with grading.

Here are the files you should submit for each project part:

- 1) assembler (part 1a)
    - a. C/C++ program for your assembler
    - b. suite of test cases (each test case is an assembly-language program in a separate file)
  - 2) simulator (part 1s)
    - a. C/C++ program for your simulator
    - b. suite of test cases (each test case is an assembly-language program in a separate file)
  - 3) multiplication (part 1m)
    - a. assembly program for multiplication
- Submit your files to the directory /cs/5.024/67Tm/linprog so your files are visible to the grader. Your assembler and simulator must each be in a single C or C++ file. We will compile your program on linprog using "gcc program.c -lm" (or g++), so your program must link with the math library.

You may also choose to not use this fragment.

```
/* Assembler code fragment for LC3101 */
```

```
#include <stdlib.h>
```

```
#include <stdio.h>
```

```
#include <string.h>
```

```

    /* after doing a readAndParse, you may want to do the following to
test the
    opcode */
    if (!strcmp(opcode, "add")) {
        /* do whatever you need to do for opcode "add" */
    }

    return(0);
}

/*
 * Read and parse a line of the assembly-language file. Fields are
returned
 * in label, opcode, arg0, arg1, arg2 (these strings must have memory
already
 * allocated to them).
 *
 * Return values:
 *     0 if reached end of file
 *     1 if all went well
 *
 * exit(1) if line is too long.
 */
int
readAndParse(FILE *inFilePtr, char *label, char *opcode, char *arg0,
             char *arg1, char *arg2)
{
    char line[MAXLINELENGTH];
    char *ptr = line;

    /* delete prior values */
    label[0] = opcode[0] = arg0[0] = arg1[0] = arg2[0] = '\0';

    /* read the line from the assembly-language file */
    if (fgets(line, MAXLINELENGTH, inFilePtr) == NULL) {
        /* reached end of file */
        return(0);
    }

    /* check for line too long (by looking for a \n) */
    if (strchr(line, '\n') == NULL) {
        /* line too long */
        printf("error: line too long\n");
        exit(1);
    }

    /* is there a label? */
    ptr = line;
    if (sscanf(ptr, "%[^\\t\\n ]", label)) {
        /* successfully read label; advance pointer over the label */
        ptr += strlen(label);
    }

    /*

```

```

        * Parse the rest of the line.  Would be nice to have real regular
        * expressions, but scanf will suffice.
        */
        sscanf(ptr, "%*[\t\n ]%[^\\t\\n ]%*[\t\n ]%[^\\t\\n ]%*[\t\n ]%[^\\t\\n
]%^*[\t\n ]%[^\\t\\n ]",
               opcode, arg0, arg1, arg2);
        return(1);
    }

int
isNumber(char *string)
{
    /* return 1 if string is a number */
    int i;
    return( (sscanf(string, "%d", &i)) == 1);
}

```

## 10. Code Fragment for Simulator

Here is some C code that may help you write the simulator. Again, you should take this merely as a hint. You may have to re-code this to make it do exactly what you want, but this should help you get started. Remember not to change stateStruct or printState.

```

/* instruction-level simulator for LC3101 */

#include <stdio.h>
#include <string.h>

#define NUMMEMORY 65536 /* maximum number of words in memory */
#define NUMREGS 8 /* number of machine registers */
#define MAXLINELENGTH 1000

typedef struct stateStruct {
    int pc;
    int mem[NUMMEMORY];
    int reg[NUMREGS];
    int numMemory;
} stateType;

void printState(stateType *);

int
main(int argc, char *argv[])
{
    char line[MAXLINELENGTH];
    stateType state;
    FILE *filePtr;

    if (argc != 2) {

```



bits. Neither a nor b are changed. E.g. (25 >> 2) is 6. Note that 25 is 11001 in binary, and 6 is 110 in binary.

3) The value of the expression (a << b) is the number "a" shifted left by "b" bits. Neither a nor b are changed. E.g. (25 << 2) is 100. Note that 25 is 11001 in binary, and 100 is 1100100 in binary.

4) To find the value of the expression (a & b), perform a logical AND on each bit of a and b (i.e. bit 31 of a ANDED with bit 31 of b, bit 30 of a ANDED with bit 30 of b, etc.). E.g. (25 & 11) is 9, since:

```
      11001 (binary)
      & 01011 (binary)
      -----
      = 01001 (binary), which is 9 decimal.
```

5) To find the value of the expression (a | b), perform a logical OR on each bit of a and b (i.e. bit 31 of a ORED with bit 31 of b, bit 30 of a ORED with bit 30 of b, etc.). E.g. (25 | 11) is 27, since:

```
      11001 (binary)
      & 01011 (binary)
      -----
      = 11011 (binary), which is 27 decimal.
```

6) ~a is the bit-wise complement of a (a is not changed).

Use these operations to create and manipulate machine-code. E.g. to look at bit 3 of the variable a, you might do: (a>>3) & 0x1. To look at bits (bits 15-12) of a 16-bit word, you could do: (a>>12) & 0xF. To put a 6 into bits 5-3 and a 3 into bits 2-1, you could do: (6<<3) | (3<<1). If you're not sure what an operation is doing, print some intermediate results to help you debug.

## 12. Example Run of Simulator

```
memory[0]=8454151
memory[1]=9043971
memory[2]=655361
memory[3]=16842754
memory[4]=16842749
memory[5]=29360128
memory[6]=25165824
```

```
memory[7]=5
memory[8]=-1
memory[9]=2
```

```
@@@
```

```
state:
```

```
pc 0
```

```
memory:
```

```
mem[ 0 ] 8454151
mem[ 1 ] 9043971
mem[ 2 ] 655361
mem[ 3 ] 16842754
mem[ 4 ] 16842749
mem[ 5 ] 29360128
mem[ 6 ] 25165824
mem[ 7 ] 5
mem[ 8 ] -1
mem[ 9 ] 2
```

```
registers:
```

```
reg[ 0 ] 0
reg[ 1 ] 0
reg[ 2 ] 0
reg[ 3 ] 0
reg[ 4 ] 0
reg[ 5 ] 0
reg[ 6 ] 0
reg[ 7 ] 0
```

```
end state
```

```
@@@
```

```
state:
```

```
pc 1
```

```
memory:
```

```
mem[ 0 ] 8454151
mem[ 1 ] 9043971
mem[ 2 ] 655361
mem[ 3 ] 16842754
mem[ 4 ] 16842749
mem[ 5 ] 29360128
mem[ 6 ] 25165824
mem[ 7 ] 5
mem[ 8 ] -1
mem[ 9 ] 2
```

```
registers:
```

```
reg[ 0 ] 0
reg[ 1 ] 5
reg[ 2 ] 0
reg[ 3 ] 0
reg[ 4 ] 0
reg[ 5 ] 0
reg[ 6 ] 0
reg[ 7 ] 0
```

```
end state
```



```

@@@
state:
    pc 2
    memory:
        mem[ 0 ] 8454151
        mem[ 1 ] 9043971
        mem[ 2 ] 655361
        mem[ 3 ] 16842754
        mem[ 4 ] 16842749
        mem[ 5 ] 29360128
        mem[ 6 ] 25165824
        mem[ 7 ] 5
        mem[ 8 ] -1
        mem[ 9 ] 2
    registers:
        reg[ 0 ] 0
        reg[ 1 ] 5
        reg[ 2 ] -1
        reg[ 3 ] 0
        reg[ 4 ] 0
        reg[ 5 ] 0
        reg[ 6 ] 0
        reg[ 7 ] 0
end state

```

```

@@@
state:
    pc 3
    memory:
        mem[ 0 ] 8454151
        mem[ 1 ] 9043971
        mem[ 2 ] 655361
        mem[ 3 ] 16842754
        mem[ 4 ] 16842749
        mem[ 5 ] 29360128
        mem[ 6 ] 25165824
        mem[ 7 ] 5
        mem[ 8 ] -1
        mem[ 9 ] 2
    registers:
        reg[ 0 ] 0
        reg[ 1 ] 4
        reg[ 2 ] -1
        reg[ 3 ] 0
        reg[ 4 ] 0
        reg[ 5 ] 0
        reg[ 6 ] 0
        reg[ 7 ] 0
end state

```

```

@@@
state:
    pc 4

```

```

memory:
    mem[ 0 ] 8454151
    mem[ 1 ] 9043971
    mem[ 2 ] 655361
    mem[ 3 ] 16842754
    mem[ 4 ] 16842749
    mem[ 5 ] 29360128
    mem[ 6 ] 25165824
    mem[ 7 ] 5
    mem[ 8 ] -1
    mem[ 9 ] 2
registers:
    reg[ 0 ] 0
    reg[ 1 ] 4
    reg[ 2 ] -1
    reg[ 3 ] 0
    reg[ 4 ] 0
    reg[ 5 ] 0
    reg[ 6 ] 0
    reg[ 7 ] 0
end state

```

@@@

```

state:
    pc 2
    memory:
        mem[ 0 ] 8454151
        mem[ 1 ] 9043971
        mem[ 2 ] 655361
        mem[ 3 ] 16842754
        mem[ 4 ] 16842749
        mem[ 5 ] 29360128
        mem[ 6 ] 25165824
        mem[ 7 ] 5
        mem[ 8 ] -1
        mem[ 9 ] 2
    registers:
        reg[ 0 ] 0
        reg[ 1 ] 4
        reg[ 2 ] -1
        reg[ 3 ] 0
        reg[ 4 ] 0
        reg[ 5 ] 0
        reg[ 6 ] 0
        reg[ 7 ] 0
end state

```

@@@

```

state:
    pc 3
    memory:
        mem[ 0 ] 8454151
        mem[ 1 ] 9043971
        mem[ 2 ] 655361

```

```

        mem[ 3 ] 16842754
        mem[ 4 ] 16842749
        mem[ 5 ] 29360128
        mem[ 6 ] 25165824
        mem[ 7 ] 5
        mem[ 8 ] -1
        mem[ 9 ] 2
    registers:
        reg[ 0 ] 0
        reg[ 1 ] 3
        reg[ 2 ] -1
        reg[ 3 ] 0
        reg[ 4 ] 0
        reg[ 5 ] 0
        reg[ 6 ] 0
        reg[ 7 ] 0
end state

```

```

@@@
state:
    pc 4
    memory:
        mem[ 0 ] 8454151
        mem[ 1 ] 9043971
        mem[ 2 ] 655361
        mem[ 3 ] 16842754
        mem[ 4 ] 16842749
        mem[ 5 ] 29360128
        mem[ 6 ] 25165824
        mem[ 7 ] 5
        mem[ 8 ] -1
        mem[ 9 ] 2
    registers:
        reg[ 0 ] 0
        reg[ 1 ] 3
        reg[ 2 ] -1
        reg[ 3 ] 0
        reg[ 4 ] 0
        reg[ 5 ] 0
        reg[ 6 ] 0
        reg[ 7 ] 0
end state

```

```

@@@
state:
    pc 2
    memory:
        mem[ 0 ] 8454151
        mem[ 1 ] 9043971
        mem[ 2 ] 655361
        mem[ 3 ] 16842754
        mem[ 4 ] 16842749
        mem[ 5 ] 29360128
        mem[ 6 ] 25165824

```

```

        mem[ 7 ] 5
        mem[ 8 ] -1
        mem[ 9 ] 2
    registers:
        reg[ 0 ] 0
        reg[ 1 ] 3
        reg[ 2 ] -1
        reg[ 3 ] 0
        reg[ 4 ] 0
        reg[ 5 ] 0
        reg[ 6 ] 0
        reg[ 7 ] 0
end state

```

@@@

state:

pc 3

memory:

```

        mem[ 0 ] 8454151
        mem[ 1 ] 9043971
        mem[ 2 ] 655361
        mem[ 3 ] 16842754
        mem[ 4 ] 16842749
        mem[ 5 ] 29360128
        mem[ 6 ] 25165824
        mem[ 7 ] 5
        mem[ 8 ] -1
        mem[ 9 ] 2

```

registers:

```

        reg[ 0 ] 0
        reg[ 1 ] 2
        reg[ 2 ] -1
        reg[ 3 ] 0
        reg[ 4 ] 0
        reg[ 5 ] 0
        reg[ 6 ] 0
        reg[ 7 ] 0

```

end state

@@@

state:

pc 4

memory:

```

        mem[ 0 ] 8454151
        mem[ 1 ] 9043971
        mem[ 2 ] 655361
        mem[ 3 ] 16842754
        mem[ 4 ] 16842749
        mem[ 5 ] 29360128
        mem[ 6 ] 25165824
        mem[ 7 ] 5
        mem[ 8 ] -1
        mem[ 9 ] 2

```

registers:

```
        reg[ 0 ] 0
        reg[ 1 ] 2
        reg[ 2 ] -1
        reg[ 3 ] 0
        reg[ 4 ] 0
        reg[ 5 ] 0
        reg[ 6 ] 0
        reg[ 7 ] 0
end state
```

```
@@@
state:
    pc 2
    memory:
        mem[ 0 ] 8454151
        mem[ 1 ] 9043971
        mem[ 2 ] 655361
        mem[ 3 ] 16842754
        mem[ 4 ] 16842749
        mem[ 5 ] 29360128
        mem[ 6 ] 25165824
        mem[ 7 ] 5
        mem[ 8 ] -1
        mem[ 9 ] 2
    registers:
        reg[ 0 ] 0
        reg[ 1 ] 2
        reg[ 2 ] -1
        reg[ 3 ] 0
        reg[ 4 ] 0
        reg[ 5 ] 0
        reg[ 6 ] 0
        reg[ 7 ] 0
end state
```

```
@@@
state:
    pc 3
    memory:
        mem[ 0 ] 8454151
        mem[ 1 ] 9043971
        mem[ 2 ] 655361
        mem[ 3 ] 16842754
        mem[ 4 ] 16842749
        mem[ 5 ] 29360128
        mem[ 6 ] 25165824
        mem[ 7 ] 5
        mem[ 8 ] -1
        mem[ 9 ] 2
    registers:
        reg[ 0 ] 0
        reg[ 1 ] 1
        reg[ 2 ] -1
        reg[ 3 ] 0
```

```
        reg[ 4 ] 0
        reg[ 5 ] 0
        reg[ 6 ] 0
        reg[ 7 ] 0
end state
```

@@@

state:

pc 4

memory:

```
mem[ 0 ] 8454151
mem[ 1 ] 9043971
mem[ 2 ] 655361
mem[ 3 ] 16842754
mem[ 4 ] 16842749
mem[ 5 ] 29360128
mem[ 6 ] 25165824
mem[ 7 ] 5
mem[ 8 ] -1
mem[ 9 ] 2
```

registers:

```
reg[ 0 ] 0
reg[ 1 ] 1
reg[ 2 ] -1
reg[ 3 ] 0
reg[ 4 ] 0
reg[ 5 ] 0
reg[ 6 ] 0
reg[ 7 ] 0
```

end state

@@@

state:

pc 2

memory:

```
mem[ 0 ] 8454151
mem[ 1 ] 9043971
mem[ 2 ] 655361
mem[ 3 ] 16842754
mem[ 4 ] 16842749
mem[ 5 ] 29360128
mem[ 6 ] 25165824
mem[ 7 ] 5
mem[ 8 ] -1
mem[ 9 ] 2
```

registers:

```
reg[ 0 ] 0
reg[ 1 ] 1
reg[ 2 ] -1
reg[ 3 ] 0
reg[ 4 ] 0
reg[ 5 ] 0
reg[ 6 ] 0
reg[ 7 ] 0
```

end state

@@@

state:

pc 3

memory:

mem[ 0 ]	8454151
mem[ 1 ]	9043971
mem[ 2 ]	655361
mem[ 3 ]	16842754
mem[ 4 ]	16842749
mem[ 5 ]	29360128
mem[ 6 ]	25165824
mem[ 7 ]	5
mem[ 8 ]	-1
mem[ 9 ]	2

registers:

reg[ 0 ]	0
reg[ 1 ]	0
reg[ 2 ]	-1
reg[ 3 ]	0
reg[ 4 ]	0
reg[ 5 ]	0
reg[ 6 ]	0
reg[ 7 ]	0

end state

@@@

state:

pc 6

memory:

mem[ 0 ]	8454151
mem[ 1 ]	9043971
mem[ 2 ]	655361
mem[ 3 ]	16842754
mem[ 4 ]	16842749
mem[ 5 ]	29360128
mem[ 6 ]	25165824
mem[ 7 ]	5
mem[ 8 ]	-1
mem[ 9 ]	2

registers:

reg[ 0 ]	0
reg[ 1 ]	0
reg[ 2 ]	-1
reg[ 3 ]	0
reg[ 4 ]	0
reg[ 5 ]	0
reg[ 6 ]	0
reg[ 7 ]	0

end state

machine halted

total of 17 instructions executed

final state of machine:

```
@@@
state:
  pc 7
  memory:
    mem[ 0 ] 8454151
    mem[ 1 ] 9043971
    mem[ 2 ] 655361
    mem[ 3 ] 16842754
    mem[ 4 ] 16842749
    mem[ 5 ] 29360128
    mem[ 6 ] 25165824
    mem[ 7 ] 5
    mem[ 8 ] -1
    mem[ 9 ] 2
  registers:
    reg[ 0 ] 0
    reg[ 1 ] 0
    reg[ 2 ] -1
    reg[ 3 ] 0
    reg[ 4 ] 0
    reg[ 5 ] 0
    reg[ 6 ] 0
    reg[ 7 ] 0
end state
```