

# Discrete Math Problem Set 10

Will Krzastek

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1. Construct the explicit functions requested below.

(a) A bijection from  $\{x \in \mathbb{R} \mid -1 < x < 1\}$  to  $\{x \in \mathbb{R} \mid -\pi < x < \pi\}$ .

**Proof.** Let  $A := \{x \in \mathbb{R} \mid -1 < x < 1\}$  and let  $B := \{x \in \mathbb{R} \mid -\pi < x < \pi\}$ . Let  $f : A \rightarrow B$ . We will define  $f$  as  $f(a) = a \cdot \pi$  where  $a \in A$ . Now, we will show that  $f$  is a bijection.

*Injective:*

Let  $a, b \in A$ . Assume  $f(a) = f(b)$ . Thus by the definition of  $f$ ,  $a\pi = b\pi$ . So,  $a = b$ . Thus,  $f$  is injective.

*Surjective:*

Let  $c \in B$ . Let  $d := \frac{c}{\pi}$ . By the definition of  $B$ , we know that  $-\pi < c < \pi$ , thus  $-1 < d < 1$ . So,  $d \in A$ . Thus, we obtain  $f(d) = d\pi = \frac{c}{\pi}\pi = c$ . Because we chose an arbitrary element of  $B$ , we know that every element of  $B$  has an element of the domain which maps to it.

Thus,  $f$  is surjective.

Thus,  $f$  is a bijection.

■

(b) A surjection from  $\mathbb{N}$  to  $\{p \mid p \text{ is prime}\}$ .

**Proof.** Let  $P := \{p \mid p \text{ is prime}\}$ .

Let  $f : \mathbb{N} \rightarrow P$  be defined as  $f(n) := p$  such that  $|\{q \in P \mid q \leq p\}| = n$ . We will show that  $f$  is a surjection.

Let  $k \in P$  where  $k \in \mathbb{N}$ . This means that  $k$  is prime by definition of  $P$ . Observe,  $k$  is finite. This lets us know that there are finitely many prime numbers less than  $k$ . As such,  $|\{q \in P \mid q \leq k\}|$  is finite, which allows us to say that  $|\{q \in P \mid q \leq k\}| = m$  for some  $m \in \mathbb{N}$ . Then by the definition of  $f$ , we know that  $f(m) = k$ . Because we showed this for an arbitrary element of  $P$ , we know that every element of  $P$  has an element of the domain which maps to it.

Therefore,  $f$  is surjective.

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- (c) An injection from  $X$  to  $\mathbb{P}(X)$  for every set  $X$ .

**Proof.** Let  $X$  be a set.

Let  $f : X \rightarrow \mathbb{P}(X)$  be defined as  $f(x) := \{x\}$ . We will show that  $f$  is injective.

Let  $a, b \in X$ . Assume  $f(a) = f(b)$ .

This means that  $\{a\} = \{b\}$ . By definition, we know  $\forall c (c \in \{a\} \iff c \in \{b\})$ .

Therefore, all elements of  $a$  and  $b$  are the same. So,  $a = b$ .

Therefore,  $f$  is injective. ■

- (d) A surjection from  $\mathbb{P}(X)$  to  $X$  for every set  $X \neq \emptyset$ .

**Proof.** Let  $X$  be a set where  $X \neq \emptyset$ . Let  $x \in X$ .

Let  $f : \mathbb{P}(X) \rightarrow X$  be defined

$$f(a) = \begin{cases} \cup a & |a| = 1 \\ x & |a| \neq 1 \end{cases}$$

We will show that  $f$  is surjective.

Let  $y \in X$ . By the definition of a subset, we know that  $\{y\} \subseteq X$ , which means that  $\{y\} \in \mathbb{P}(X)$  by the definition of a power set. We also know that  $|\{y\}| = 1$  because there is only one element. Therefore by the definition of  $f$ , we know that  $f(\{y\}) = \cup\{y\} = y$ .

Because we showed this for an arbitrary element of  $X$ , we know that  $f$  is surjective for every set  $X$  where  $X \neq \emptyset$ . ■

2. Let  $A$  be an arbitrary finite set of cardinality  $|A| = n$ , where  $n \in \mathbb{N}$ . How many finite strings over  $A$  are there?

**Proof.** Let  $A$  be a set, where  $|A| = n$  for some  $n \in \mathbb{N}$ . Let  $s := \{z \mid (\exists c \in \mathbb{N})(z = \{s \mid s : c \rightarrow A\})\}$ . Observe,  $\cup s = \bigcup_{i \in \mathbb{N}} \{s \mid s : i \rightarrow A\}$ . Let's define  $B \in s$  for some  $m \in \mathbb{N}$ . So,  $B = \{s \mid s : m \rightarrow A\}$ . Thus,  $|B| = |A|^m = n^m$  by *Theorem 6.10*. So,  $|B|$  is finite thus  $B$  is countable. This means that  $\cup s$  is a countable collection of countable sets. Therefore by *Theorem 8.8*, we know that  $|\cup s|$  is countable, meaning  $|\cup s| \leq |\mathbb{N}|$ .

Now, we would like to prove that  $\cup s$  is an infinite set. Towards a contradiction, assume  $\cup s$  is finite. This means that for some  $l \in \mathbb{N}$ ,  $|\cup s| = l$ . So, there exists some surjection  $f : l \rightarrow \cup s$ .

Let  $a \in A$ . Let  $h := f(0) \cup f(1) \cup \dots \cup f(l-1) \cup a$ . So,  $|h| = \sum_{i=0}^{l-1} |f(i)| + 1$ . So by definition,

$$(\forall k \in l)(|f(k)|) \leq \sum_{i=0}^{l-1} |f(i)| < \sum_{i=0}^{l-1} |f(i)| + 1 = |h|. \text{ This tells us that } (\forall k \in l)(|f(k)| \neq |h|).$$

Therefore,  $(\forall k \in l)(f(k) \neq h)$ . So,  $h$  is a finite string over  $A$  where  $h \in \cup s$ , but  $f$  does

not map to  $h$ . Therefore  $f$  is not a surjection. However, we assumed  $f$  was a surjection  $\nmid$ . Therefore,  $(\forall l \in \mathbb{N})(\cup s \neq l)$ , so  $\cup s$  is infinite.

Therefore, we know that  $\cup s$  is both countable and infinite.

Therefore by definition, we know  $|\cup s| = |\mathbb{N}| = \aleph_0$ . So, there are  $\aleph_0$  finite strings over  $A$ .

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3. Imagine that, one day, you encounter a library. At the entrance of this library is an enormous tome  $\mathcal{B}$  listing *all* of the *possible* sentences in the English language, indexed by natural numbers.

Walking past the entrance, you see that the library has rows of bookshelves numbered  $0, 1, 2, \dots$ , so that there is exactly one row of books for each  $n \in \mathbb{N}$ . A great owl—perched on the pedestal that supports  $\mathcal{B}$ —informs you that, for each  $n \in \mathbb{N}$ , the  $n^{\text{th}}$  row of bookshelves contains *all* of the books that could possibly ever be that begin with the  $n^{\text{th}}$  sentence in  $\mathcal{B}$ . With the sound of granite scraping against marble, the doors to the library close behind you. The owl makes you the following proposal: you will be free to go *if and only if* you can read every book in the library *in a countable amount of time*.

Will you be set free? The owl demands a proof to justify your answer.

**Proof.** Observe a few things. There are countably many sentences in  $\mathcal{B}$  because they are numbered.  $\mathcal{B}$  is a set of cardinality  $\mathbb{N}$  where  $n \in \mathbb{N}$  is a sentence at the  $n^{\text{th}}$  position of  $\mathcal{B}$ . We know that every book contains an infinitely countable amount of sentences, because they are indexed. Thus, we will represent the amount of sentences in each book with  $|\mathbb{N}|$ . Let  $S$  be  $S := \{s \mid s : \mathbb{N} \rightarrow \mathbb{N}\}$  where  $S$  is the set of all books of infinite length in the library. Thus,  $S \subseteq$  “all of the books in the library”.

Towards a contradiction, assume  $|\mathbb{N}| \geq |S|$ . This means there exists some  $f : \mathbb{N} \rightarrow S$  such that  $f$  is surjective. Let  $g \in S$ . Consider,  $g : \mathbb{N} \rightarrow \mathbb{N}$  where:

$$g(i) := \begin{cases} 1 & f(i)(i) \neq 1 \\ 2 & f(i)(i) = 1 \end{cases}$$

Because  $f$  is surjective,  $(\exists t \in \mathbb{N})(f(t) = g)$ . So,  $(\forall i \in \mathbb{N})(f(t)(i) = g(i))$ . However, observe  $f(t)(t) = 1 \Leftrightarrow g(t) = 2 \neq 1$ , and  $f(t)(t) \neq 1 \Leftrightarrow g(t) = 1$ . This means that  $f(t)(t) \neq g(t)$ , so  $(\exists i \in \mathbb{N})f(t)(i) \neq g(i)$ .

Therefore,  $|\mathbb{N}| < |S|$ . Because  $|S| > |\mathbb{N}|$ , we know  $|S| > \aleph_0$ , which means by definition that  $S$  is uncountably infinite. So I will sadly *not* be set free and the owl will own me forever :/

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