

A Completeness Theorem for Probabilistic Regular Expressions

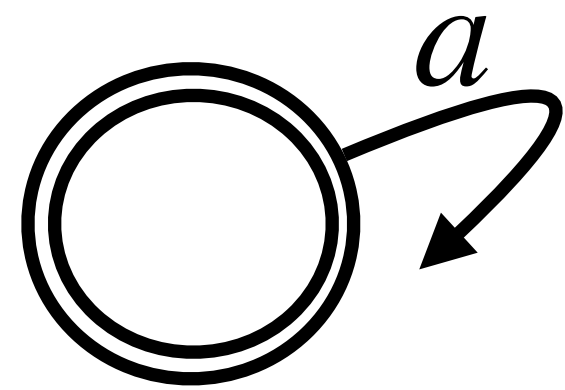
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Regular expressions

a^*

(Non)deterministic
finite automata



Regular languages

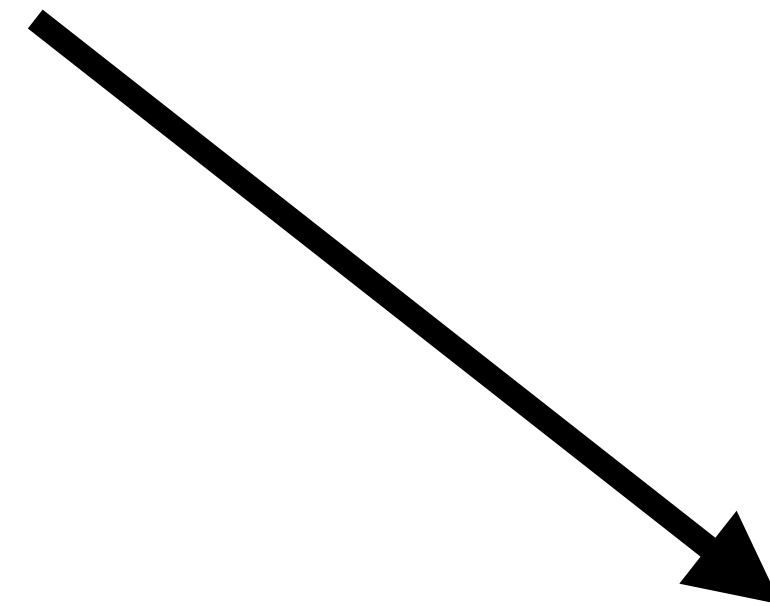
$\{\varepsilon, a, aa, aaa, \dots\} \in A^* \rightarrow 2$

This talk:

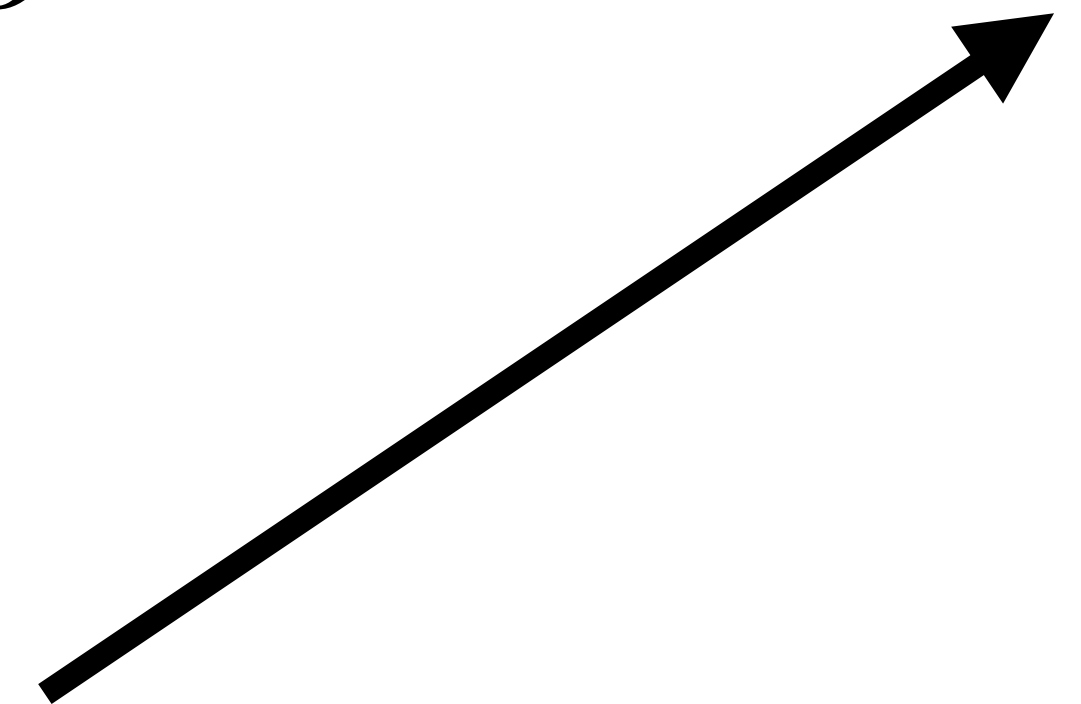
Probabilistic Regular Expressions (PRE)

$e, f \in \text{Exp} ::= 0 \mid 1 \mid a \in A \mid e + f \mid e; f \mid e^*$

do e with probability p or do f with probability $1 - p$



$e, f \in \text{Exp} ::= 0 \mid 1 \mid a \in A \mid e \oplus_p f \mid e; f \mid e^{[p]}$



do e with probability p and then start again
or terminate with probability $1 - p$

$$\text{Sem} : A^* \rightarrow [0,1]$$

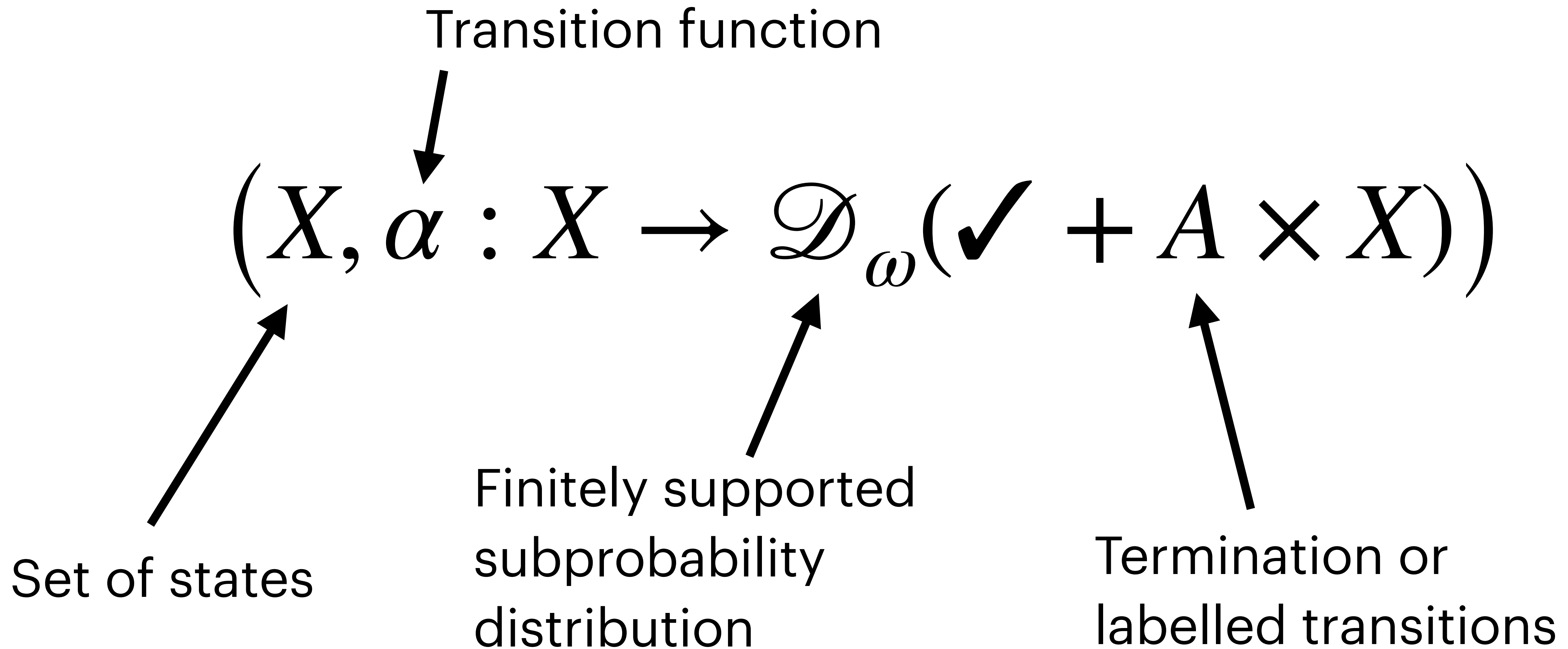
$$\text{Sem} \left(a^{\left[\frac{1}{2} \right]} \right) (w) = \begin{cases} \frac{1}{2} \times \left(\frac{1}{2} \right)^n & w = a^n \\ 0 & \text{otherwise} \end{cases}$$

The word *aaa* is accepted with probability $\frac{1}{16}$

Contribution:
**Complete axiomatisation of probabilistic
language equivalence of PREs**

In the case of RE, problem studied by Salomaa, Krob, Boffa and Kozen

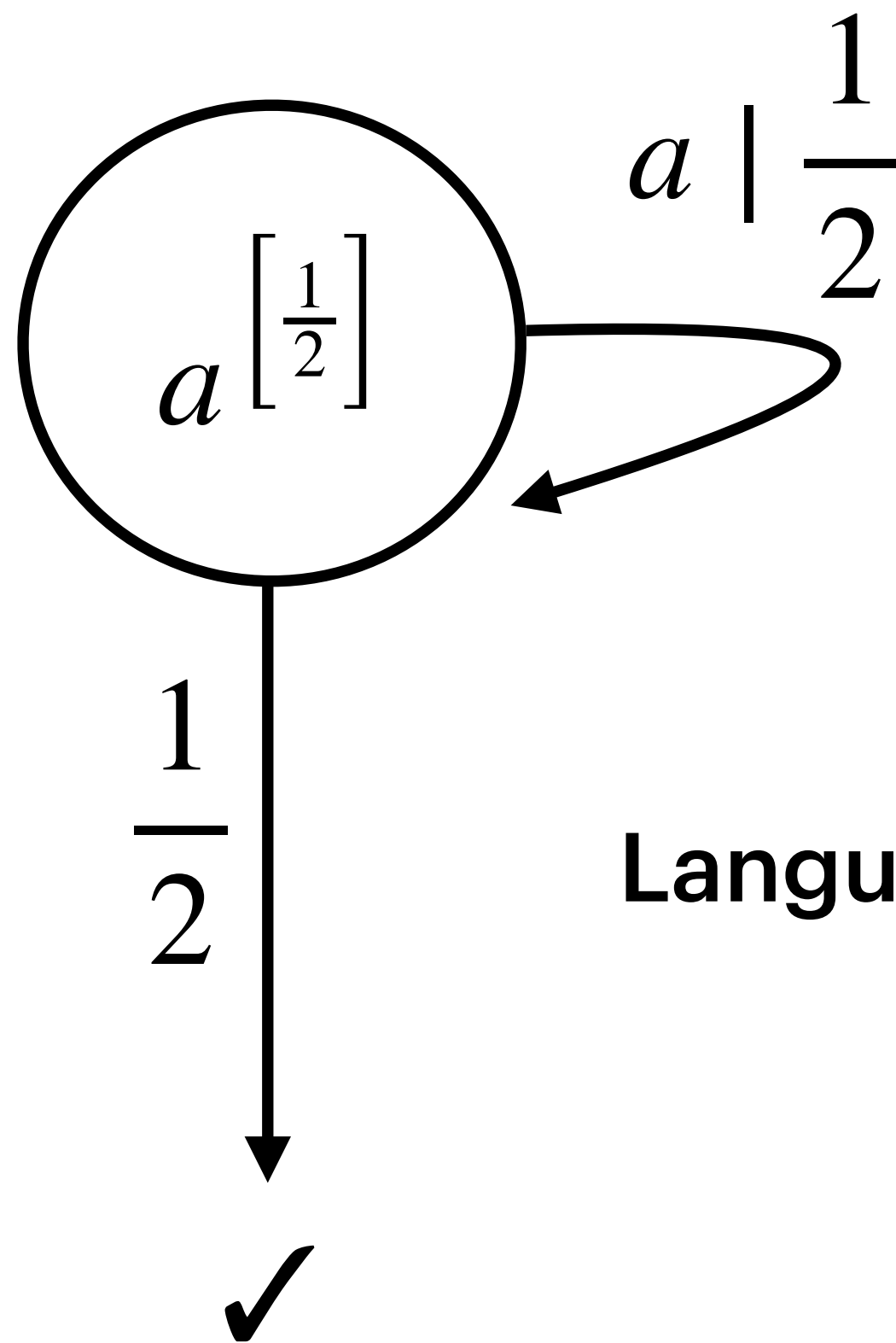
Generative Probabilistic Transition Systems



Operational semantics

Transition system whose states are expressions (similar to Antimirov derivatives)

$$(\text{Exp}, \partial : \text{Exp} \rightarrow \mathcal{D}_\omega(\checkmark + A \times \text{Exp}))$$



$$\text{Sem}(e) := \text{PLan}_{(\text{Exp}, \partial)}(e)$$

Language semantics obtained through a construction from coalgebra

The axioms

Probabilistic choice

$$e \oplus_p e \equiv e$$

$$e \oplus_1 f \equiv e$$

$$e \oplus_p f \equiv f \oplus_{1-p} e$$

$$(e \oplus_p f) \oplus_q g \equiv e \oplus_{pq} (f \oplus_{\frac{(1-p)q}{1-pq}} g)$$

Can be generalised to an n-ary subprobabilistic choice and becomes Positive Convex Algebra

Sequential composition

$$1; e \equiv e \equiv e; 1$$

$$0; e \equiv 0 \equiv e; 0$$

$$e; (f; g) \equiv (e; f); g$$

$$e; (f \oplus_p g) \equiv e; g \oplus_p e; g$$

$$(e \oplus_p f); g \equiv e; g \oplus_p f; g$$

Loops

$$e^{[p]} \equiv e; e^{[p]} \oplus_p 1$$

$$(e \oplus_p 1)^{[q]} \equiv e \left[\frac{pq}{1 - (1-p)q} \right]$$

$$1^{[1]} \equiv 0$$

$$g \equiv eg \oplus_p f \quad e \text{ accepts the empty word with probability } 0$$

$$g \equiv e^{[p]}; f$$

Soundness: If $e \equiv f$ then $\text{Sem}(e) = \text{Sem}(f)$ ✓

Completeness: If $\text{Sem}(e) = \text{Sem}(f)$ then $e \equiv f$



Equivalent to establishing a certain universal property

Using a series of very hard results in convex algebra

(Due to Sokolova, Woracek, Milius)

Milius, S. (2018). Proper Functors and Fixed Points for Finite Behaviour. *Logical Methods in Computer Science*

Sokolova, A. and Harald Woracek (2015). Congruences of convex algebras. *Journal of pure and applied algebra*

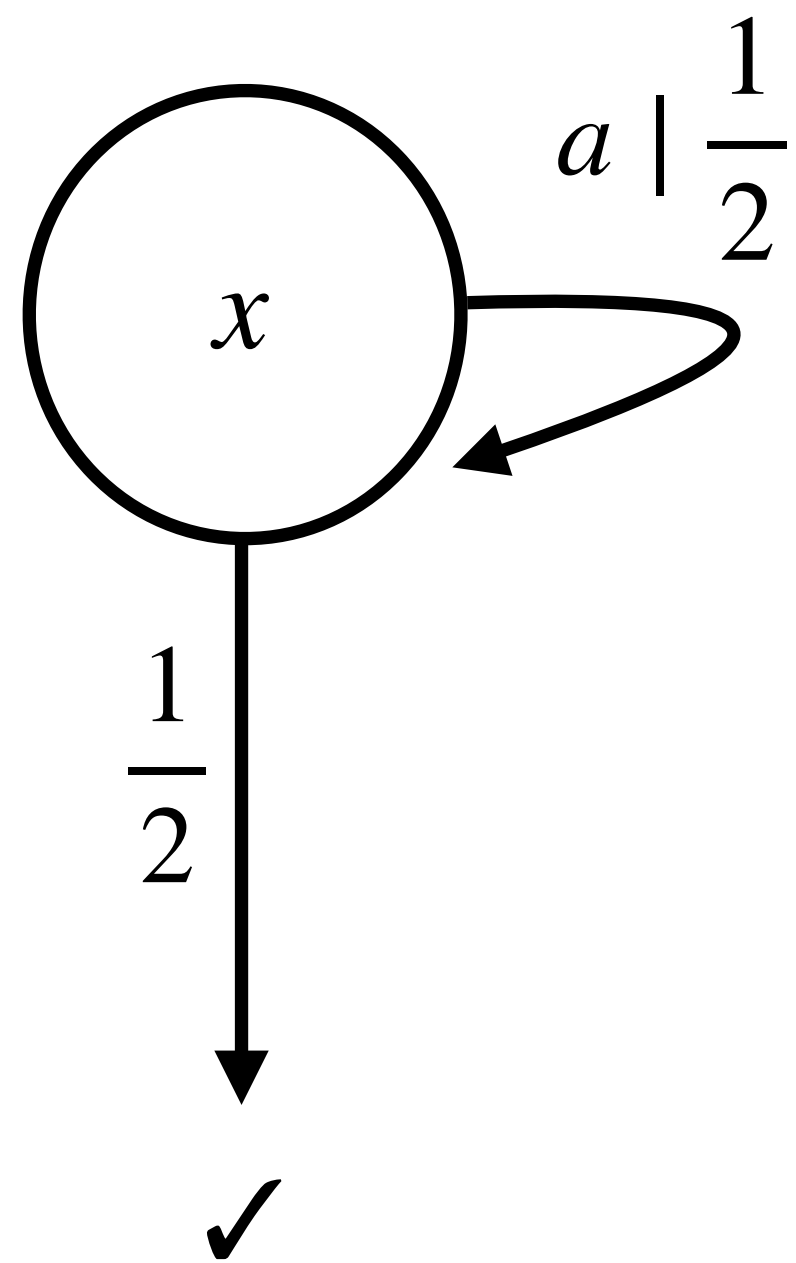
Sokolova, A. and Woracek, H. (2018). Proper Semirings and Proper Convex Functors. *Foundations of Software Science and Computation Structures (FoSSaCS)*

...and through a reduction which is the most technical bit of this work

It suffices to show:

**Every finite-state automaton can be uniquely converted to an expression
(up to the axioms of \equiv)**

Automaton



System of equations

$$x \equiv ax \oplus_{\frac{1}{2}} 1$$

Solution

$$x \rightarrow a^{[\frac{1}{2}]}; 1$$

Thank you!

- Generalisation of classic automata theory through heavy tools from coalgebra and convex algebra
- Previous work by Silva and Sokolova relied on extending the complete axiomatisation modulo bismilarity of a probabilistic process calculus
- Future work: quantitative axiomatisations, algebraic Kozen-style axiomatisation



Check out the full paper :)