

# Behavioural Metrics

## Compositionality of the Kantorovich Lifting and an Application to Up-To Techniques

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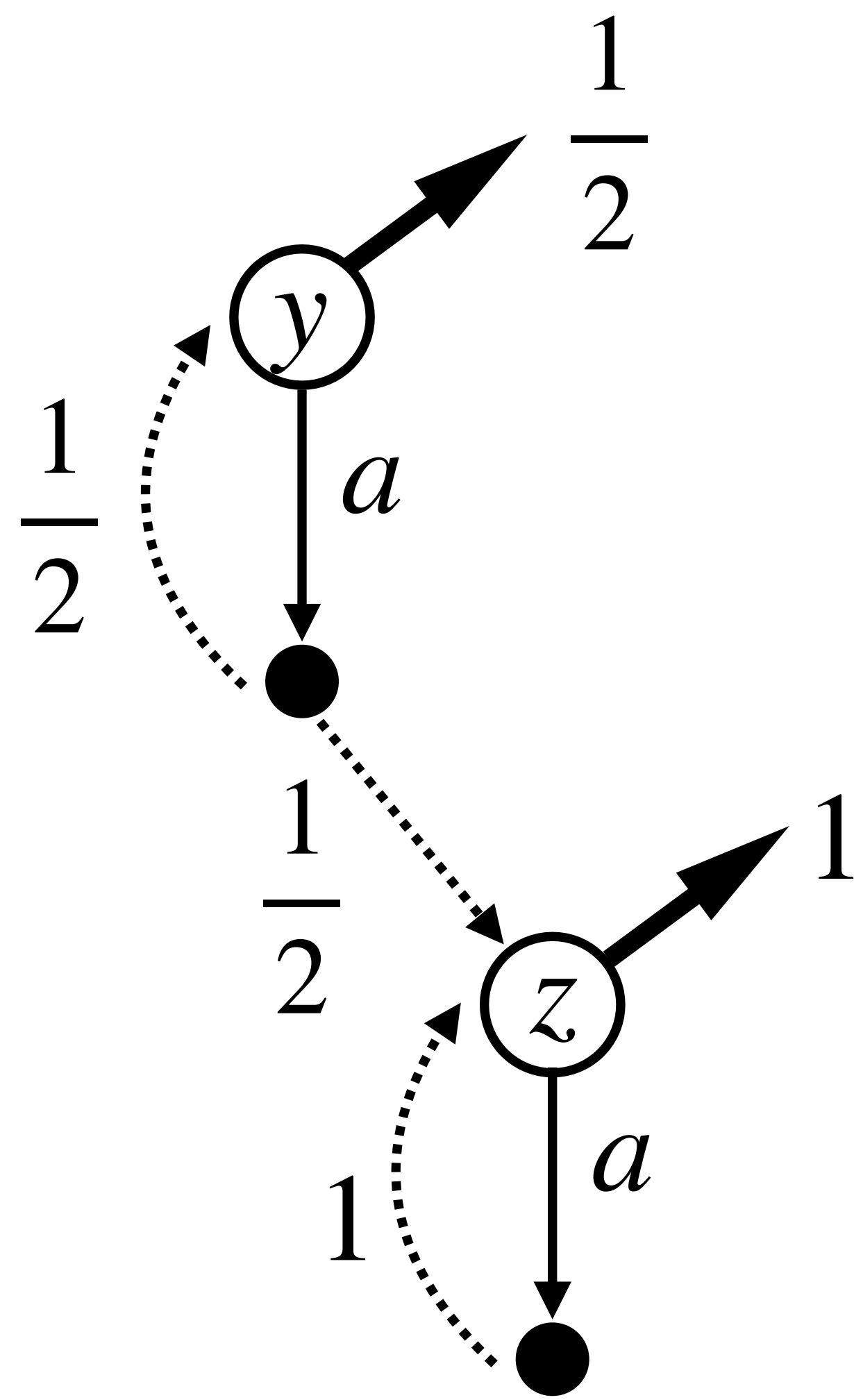
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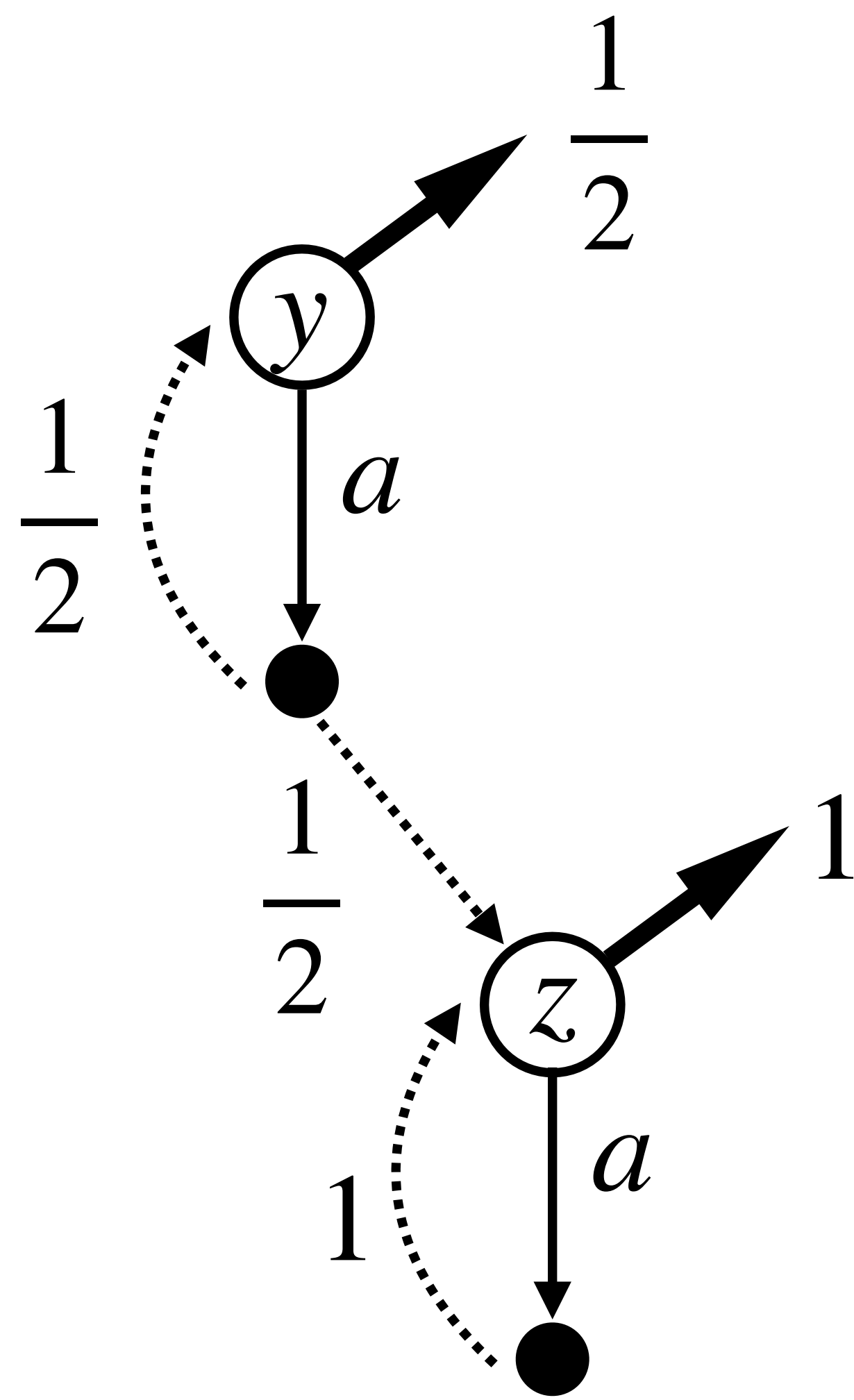
**We give a sound technique for checking bounds on these distances**

# Rabin probabilistic automata



$$(X, X \rightarrow [0,1] \times \mathcal{D}(X)^A)$$

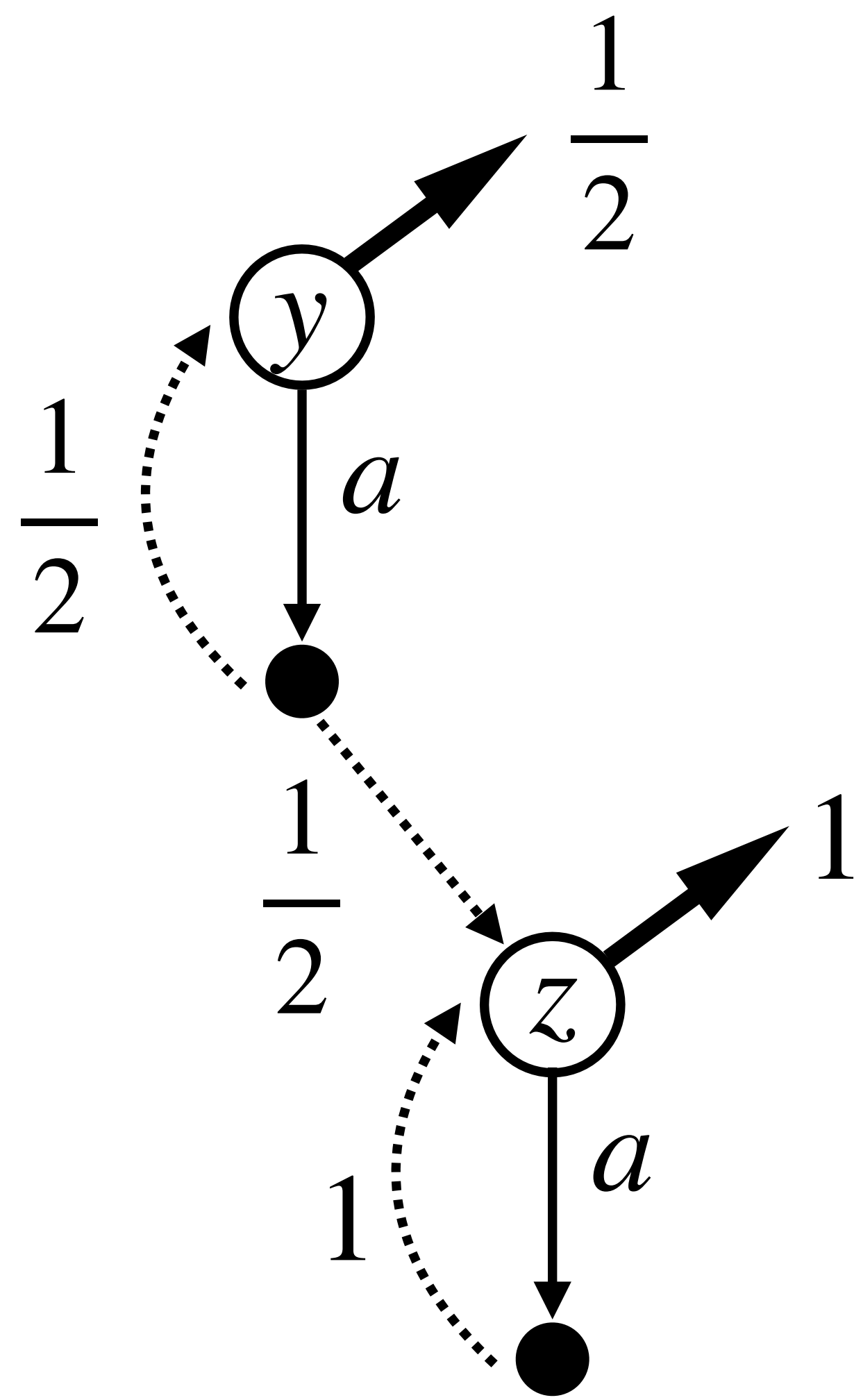
# Rabin probabilistic automata



Set of states

$(X, X \rightarrow [0,1] \times \mathcal{D}(X)^A)$

# Rabin probabilistic automata



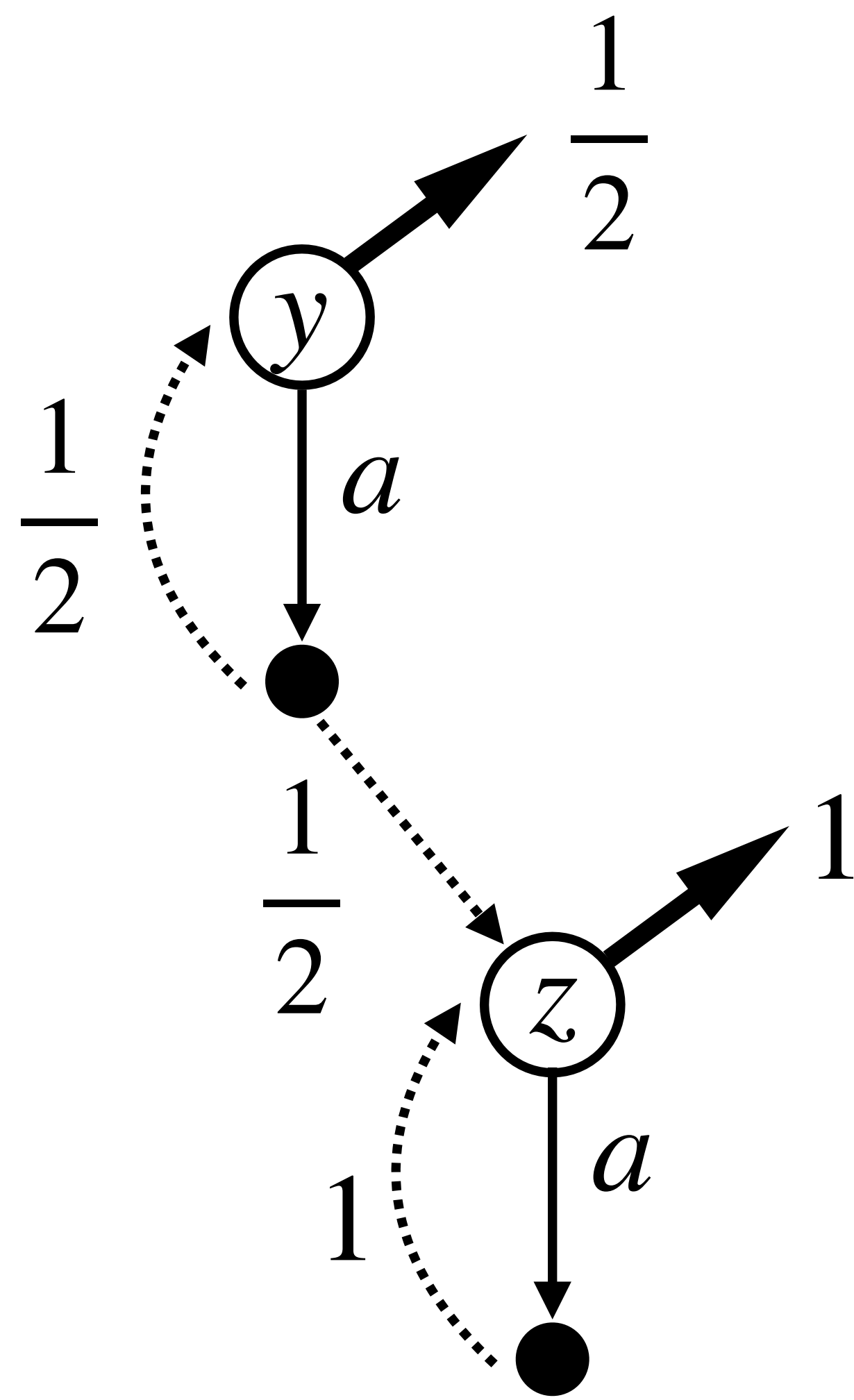
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Acceptance weight



# Rabin probabilistic automata



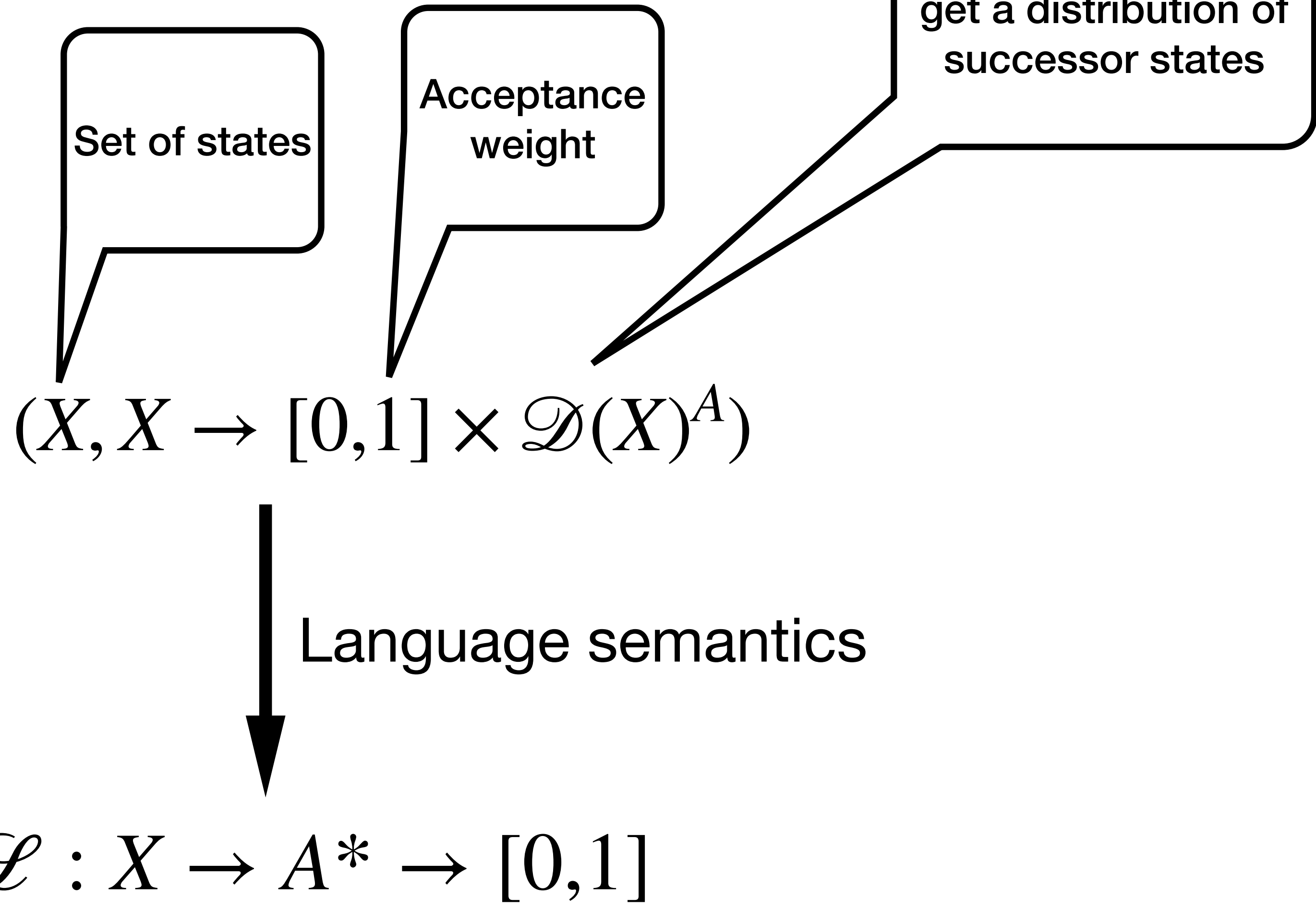
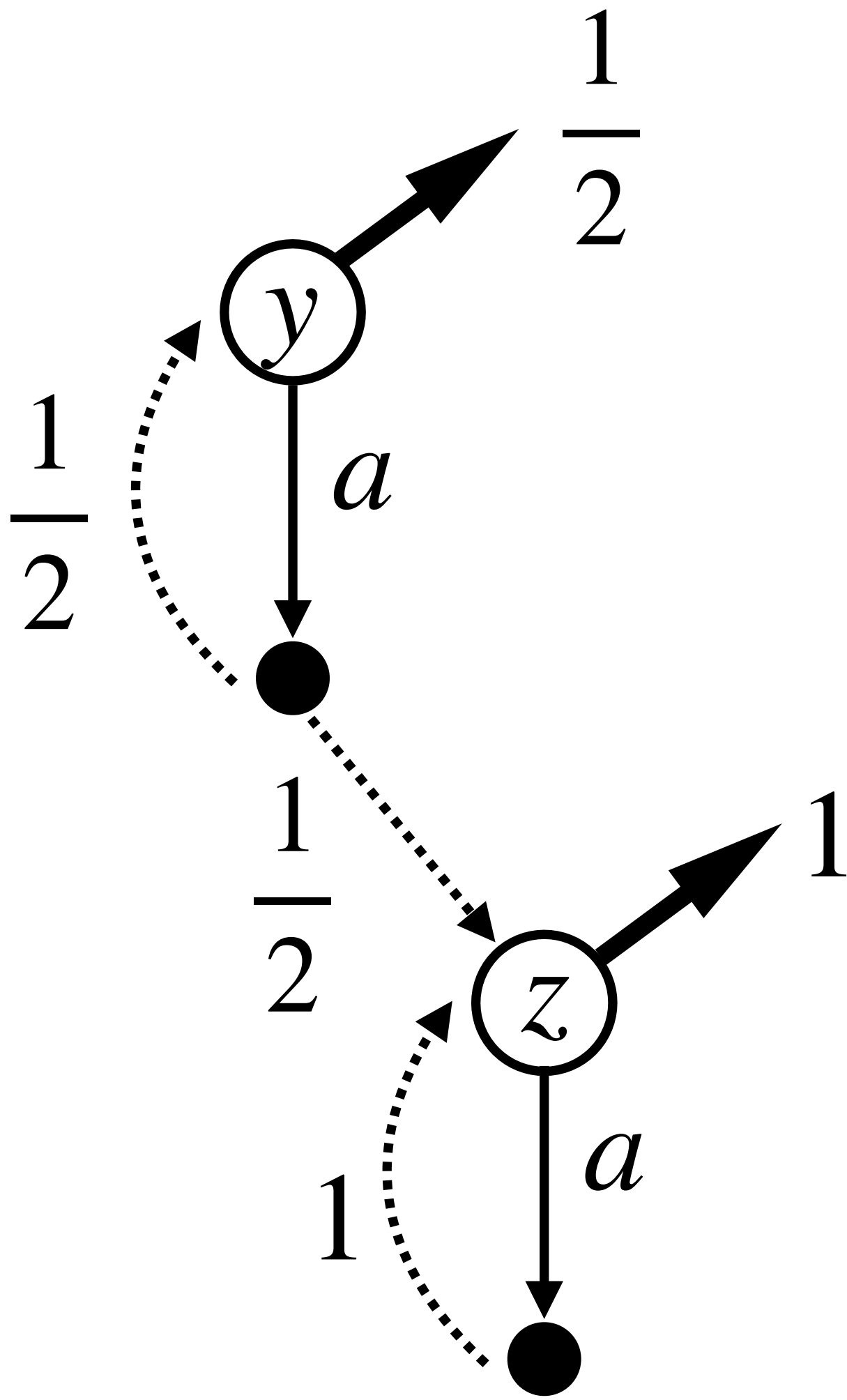
Set of states

$(X, X \rightarrow [0,1] \times \mathcal{D}(X)^A)$

Acceptance weight

Given a letter  $a$ , we get a distribution of successor states

# Rabin probabilistic automata



# The framework of Universal Coalgebra

$$X \rightarrow 2 \times X^A$$



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$$\left. \begin{array}{l} X \rightarrow 2 \times X^A \\ X \rightarrow \mathcal{P}(1 + A \times X) \\ X \rightarrow [0,1] \times \mathcal{D}(X)^A \end{array} \right\} X \rightarrow FX$$



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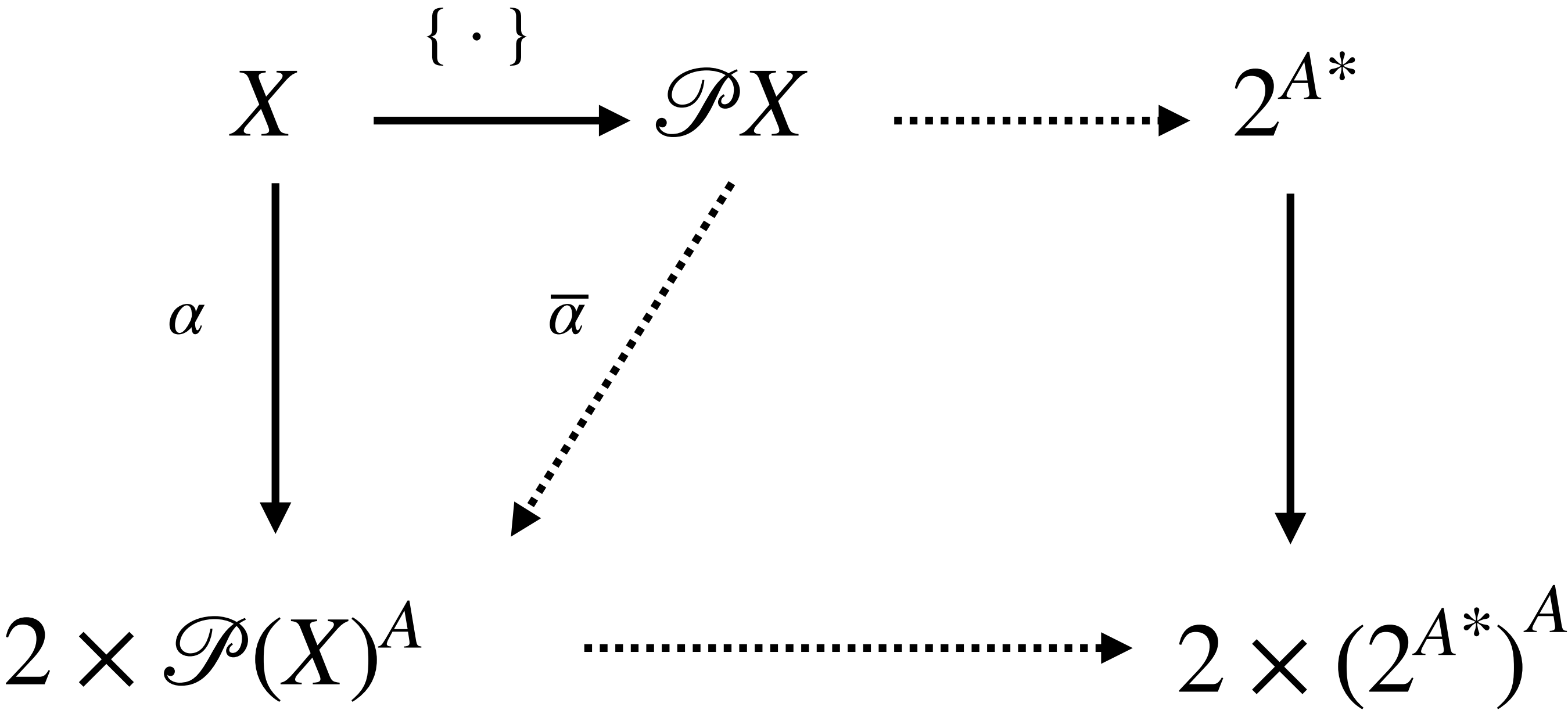
Endofunctor  $F$  describes one-step dynamics of the system

$$\begin{array}{ccc} X & \cdots \rightarrow & \Omega \\ \downarrow & & \downarrow \\ FX & \cdots \rightarrow & F\Omega \end{array}$$

Final coalgebra: canonical domain for branching-time semantics. Exists under mild size constraints.

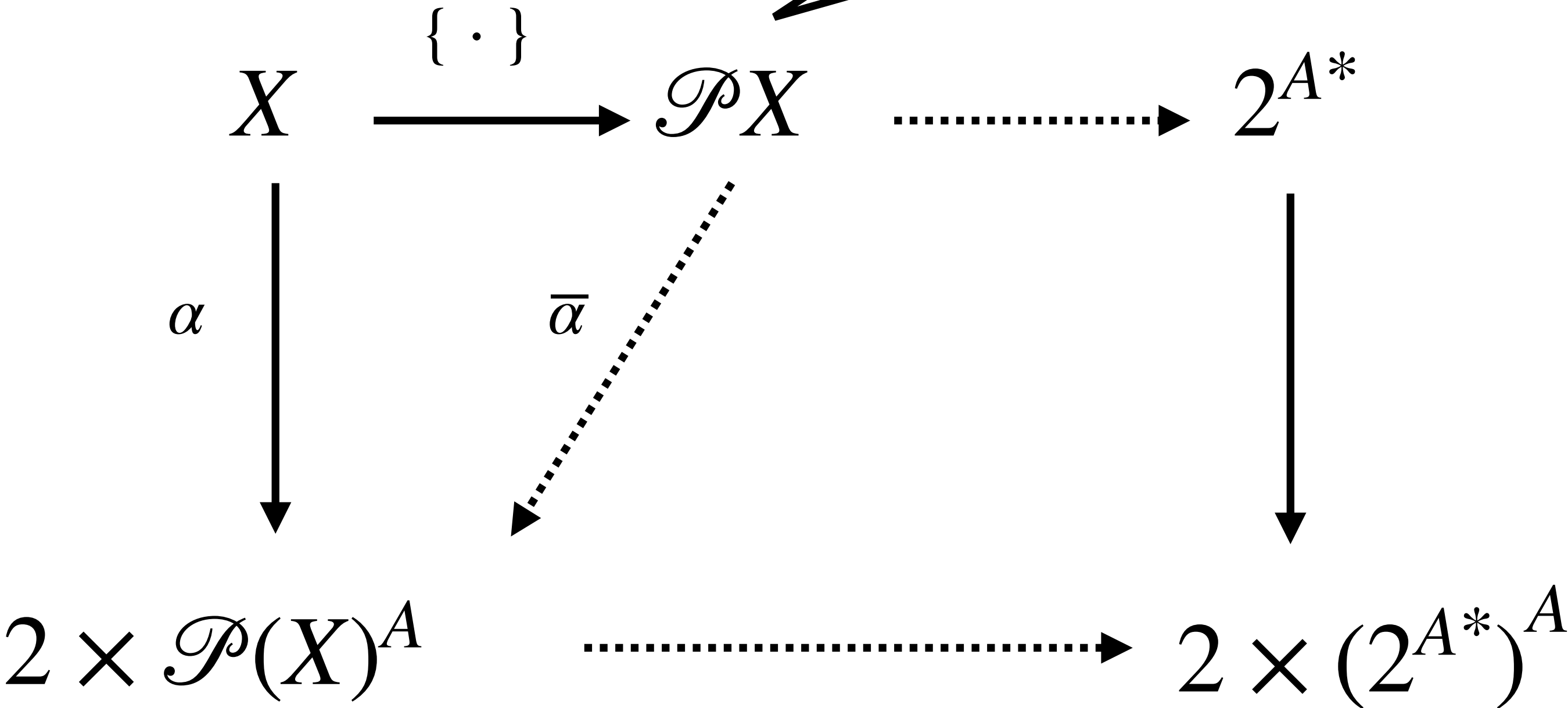


# NFA determinisation



# NFA determinisation

State space has extra structure of a semilattice



# Generalized determinisation

$$\begin{array}{ccccc} X & \xrightarrow{\eta} & TX & \cdots\cdots\cdots & \Omega \\ \downarrow \alpha & & \swarrow \bar{\alpha} & & \downarrow \\ FTX & \cdots\cdots\cdots & & \cdots\cdots\cdots & F\Omega \end{array}$$



# Generalized determinisation

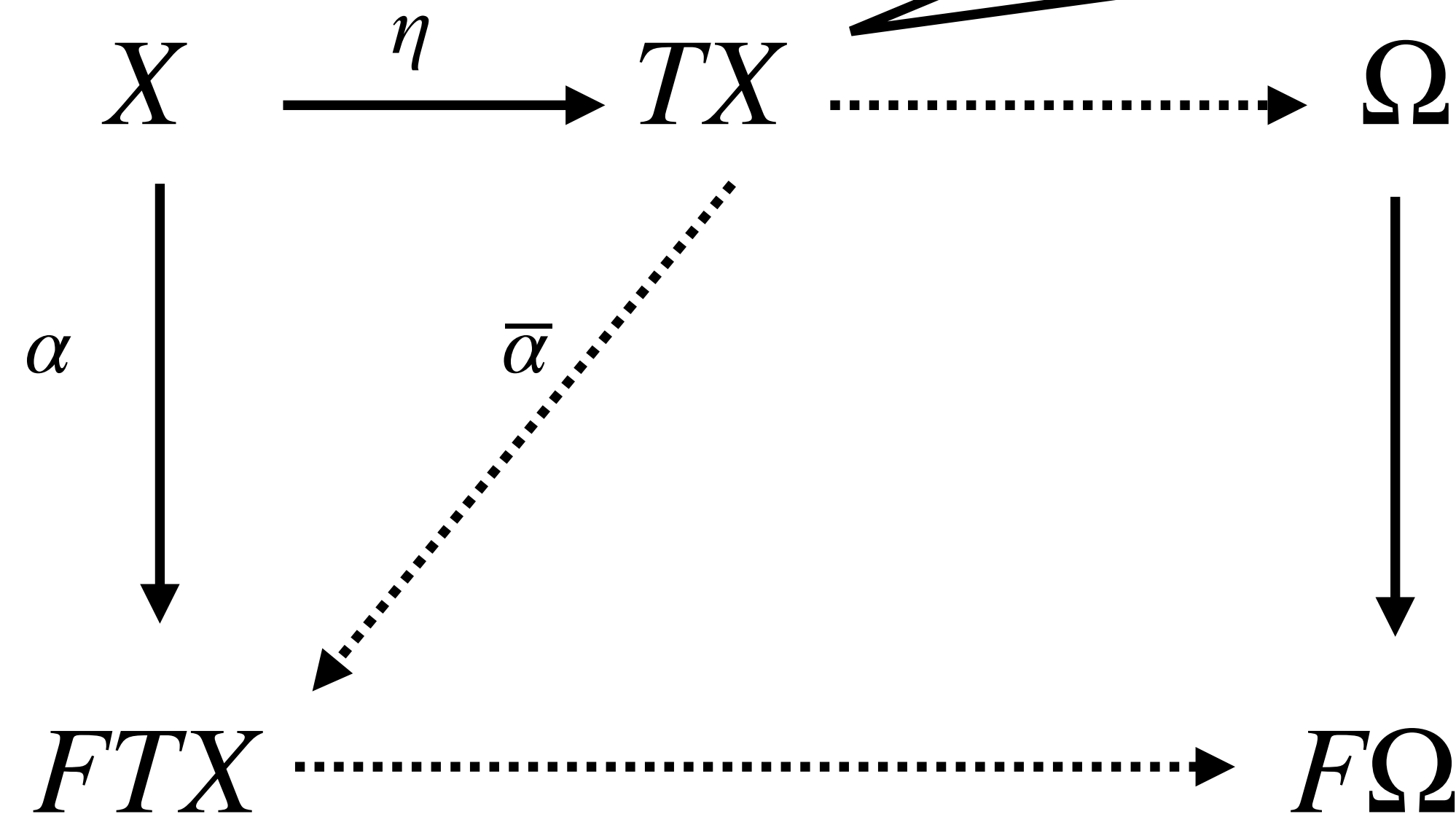
$$\begin{array}{ccccc} X & \xrightarrow{\eta} & TX & \cdots \longrightarrow & \Omega \\ \alpha \downarrow & & \searrow \bar{\alpha} & & \downarrow \\ FTX & \cdots \longrightarrow & F\Omega & & \end{array}$$

$T$  is a monad, that interacts with  $F$  via a distributive law  $\lambda : TF \Rightarrow FT$

# Generalized determinisation

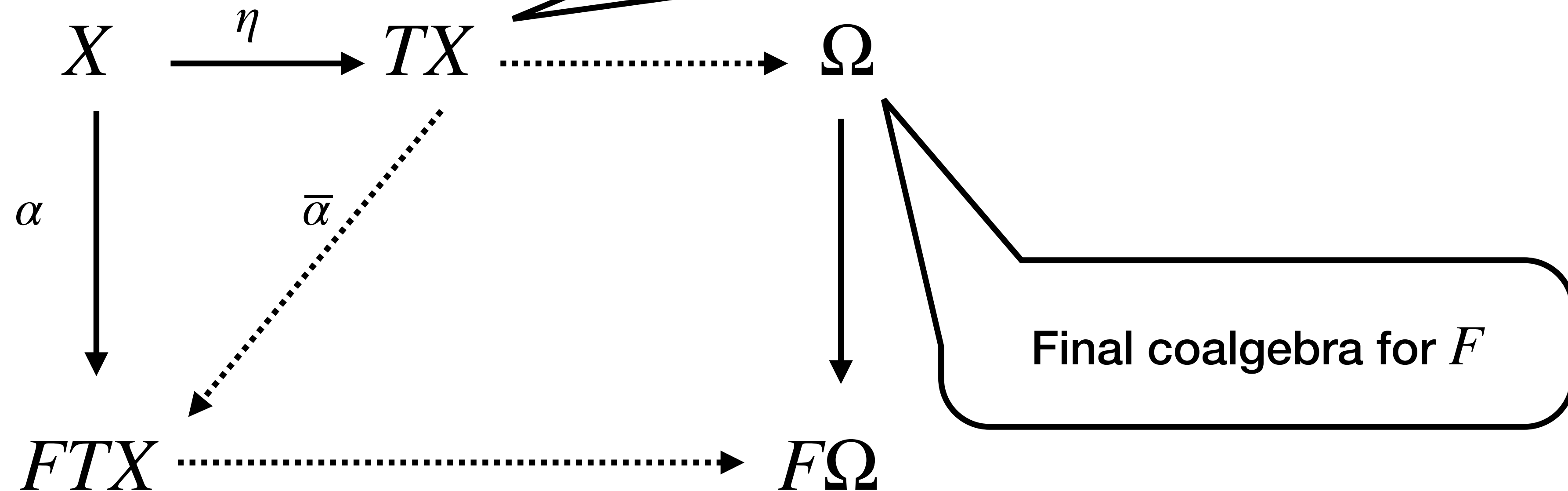
Extra structure of an algebra  
for a monad:

$$\mu_X : TTX \rightarrow TX$$



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# Generalized determinisation

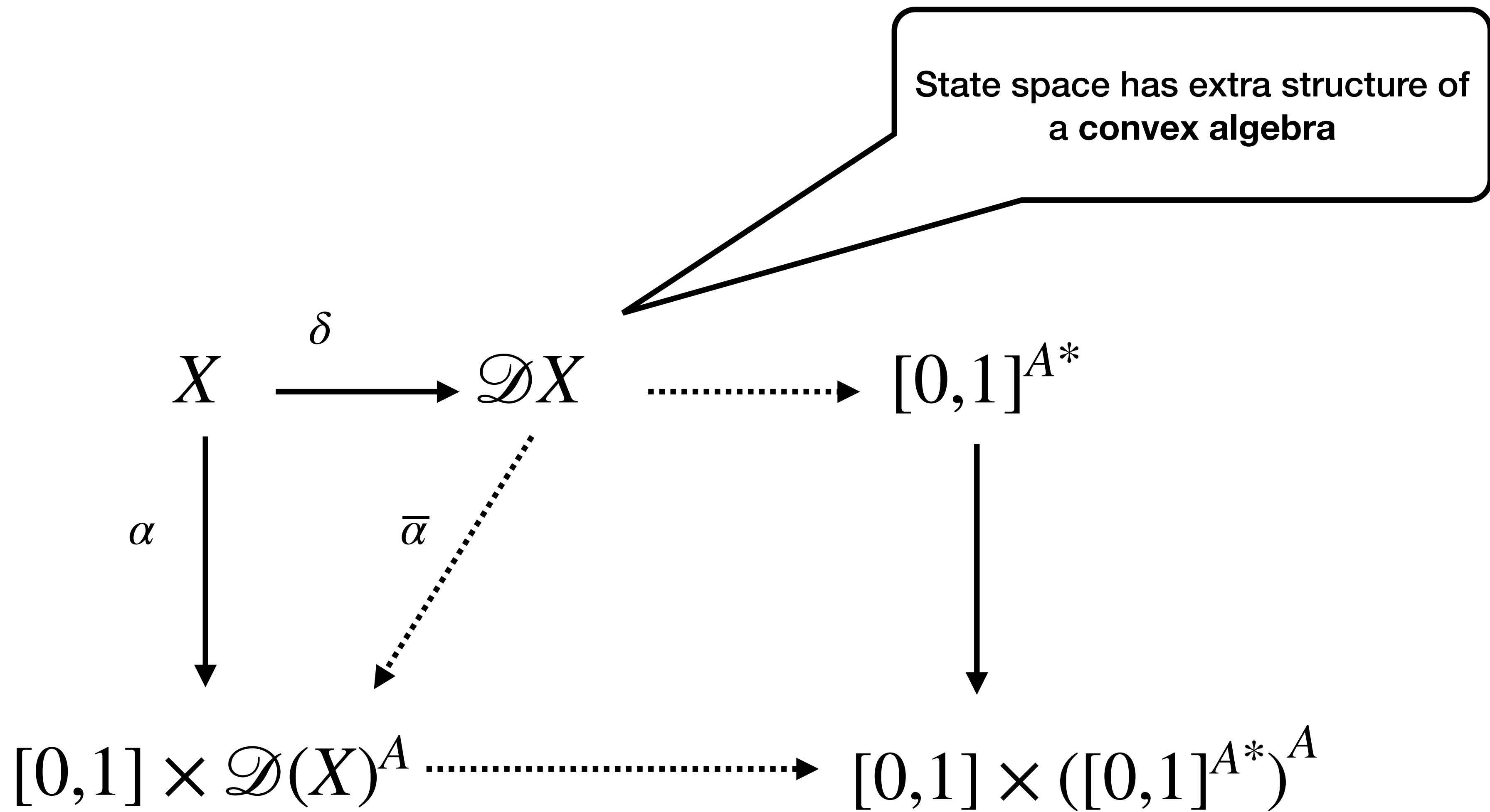


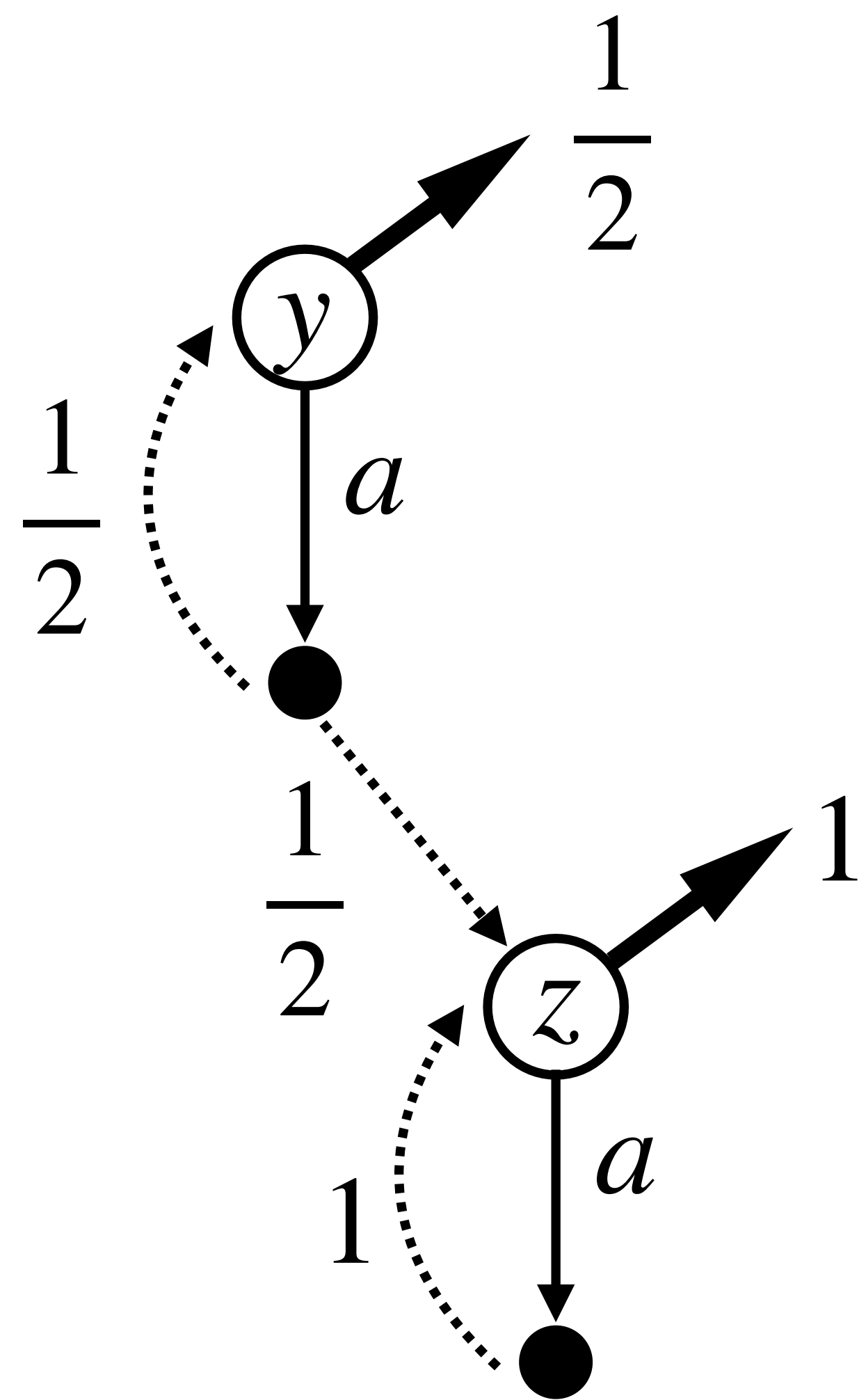
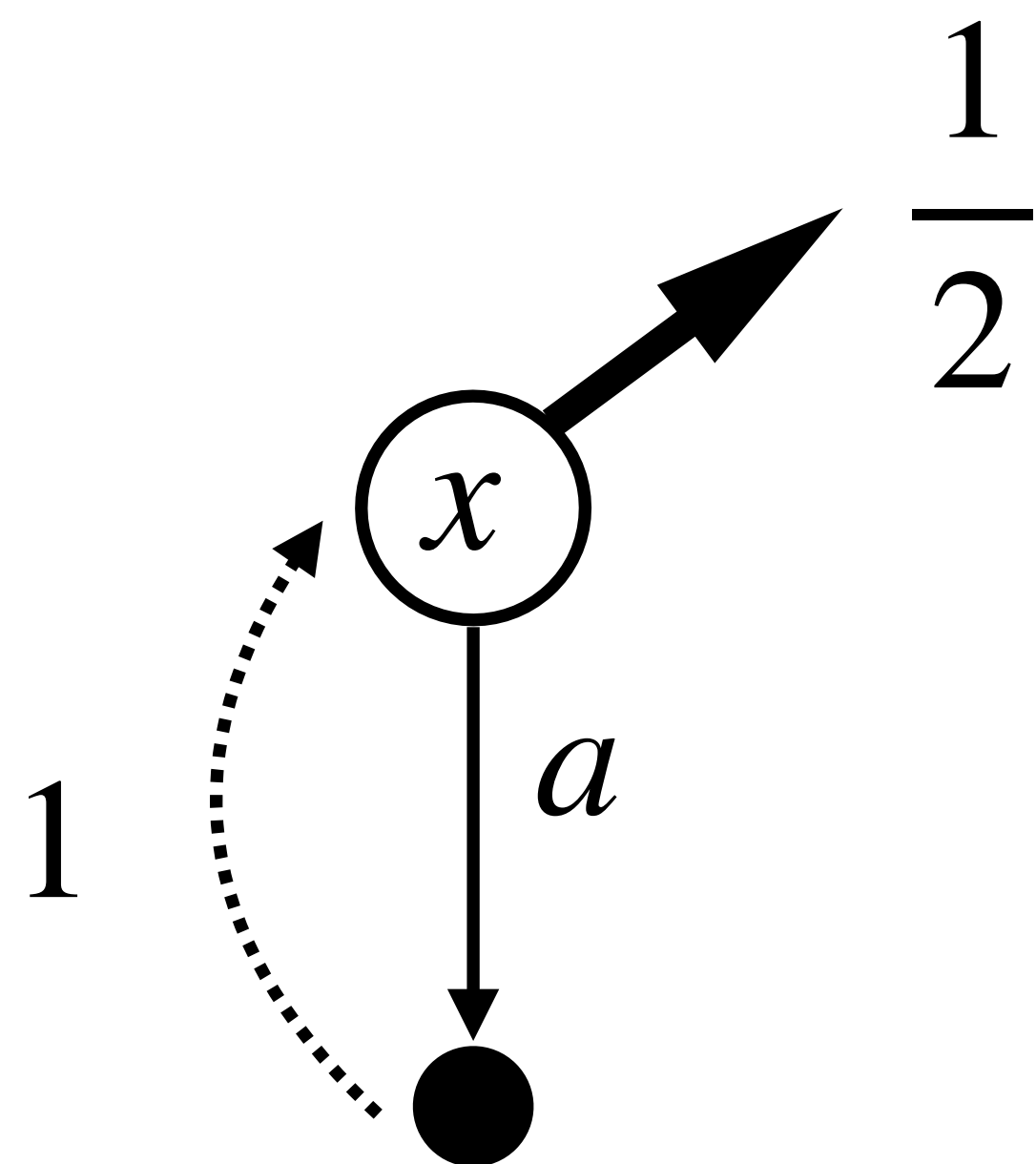
$T$  is a monad, that interacts with  $F$  via a distributive law  $\lambda : TF \Rightarrow FT$

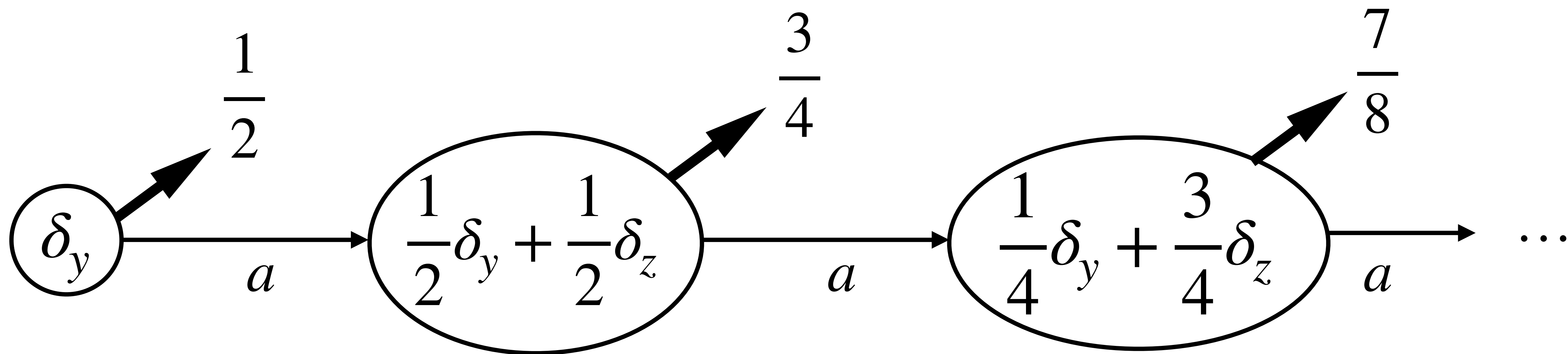
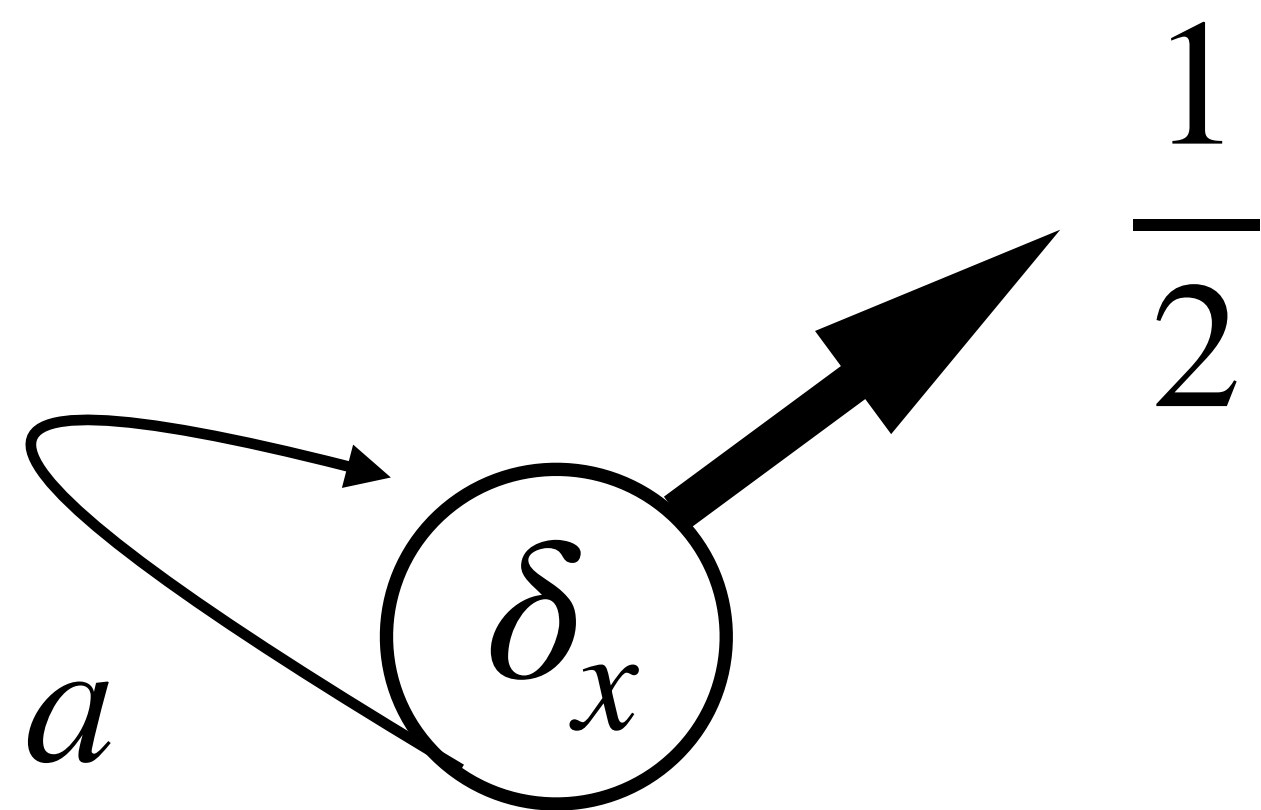
**Back to the example of Rabin  
automata**

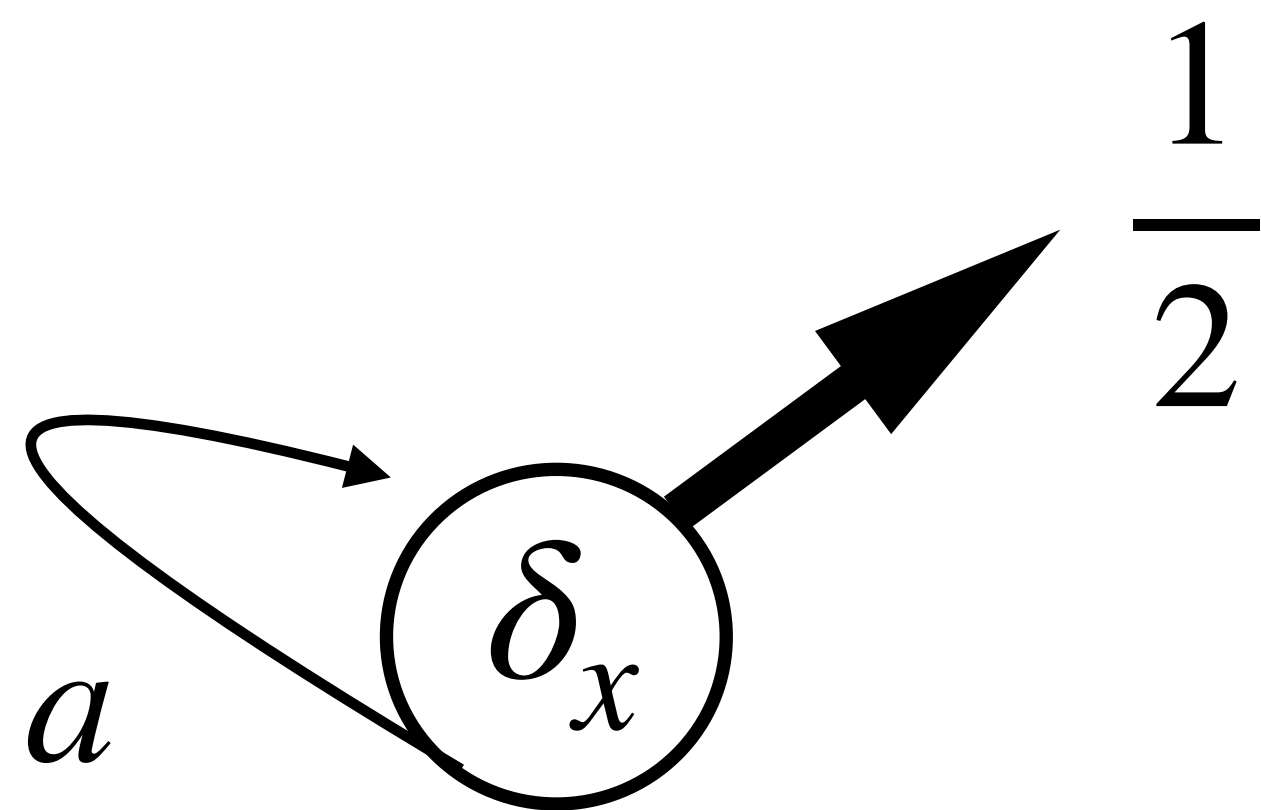
$$\begin{array}{ccccc}
 X & \xrightarrow{\delta} & \mathcal{D}X & \cdots\cdots\cdots\rightarrow & [0,1]^{A*} \\
 \downarrow \alpha & & \searrow \bar{\alpha} & & \downarrow \\
 [0,1] \times \mathcal{D}(X)^A & \cdots\cdots\cdots\rightarrow & & & [0,1] \times ([0,1]^{A*})^A
 \end{array}$$



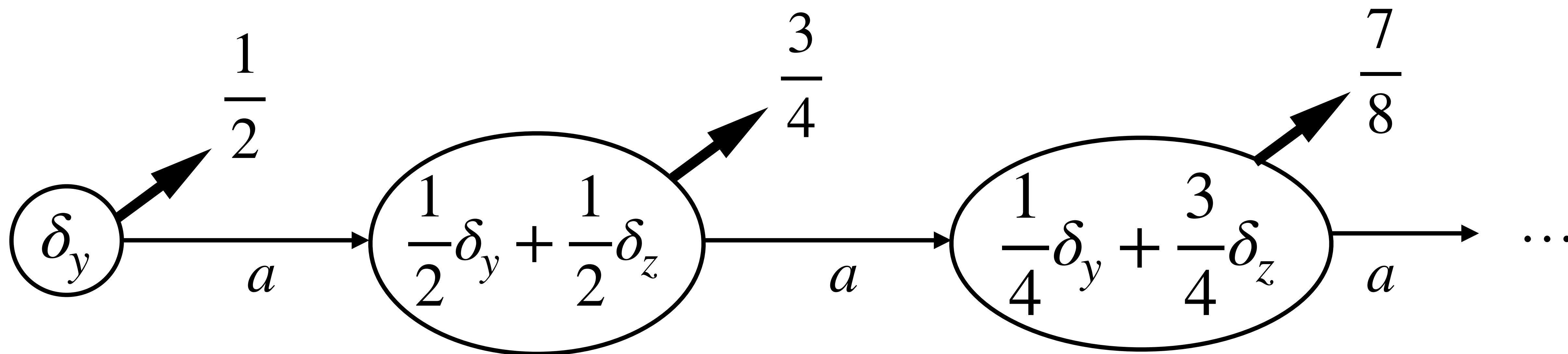


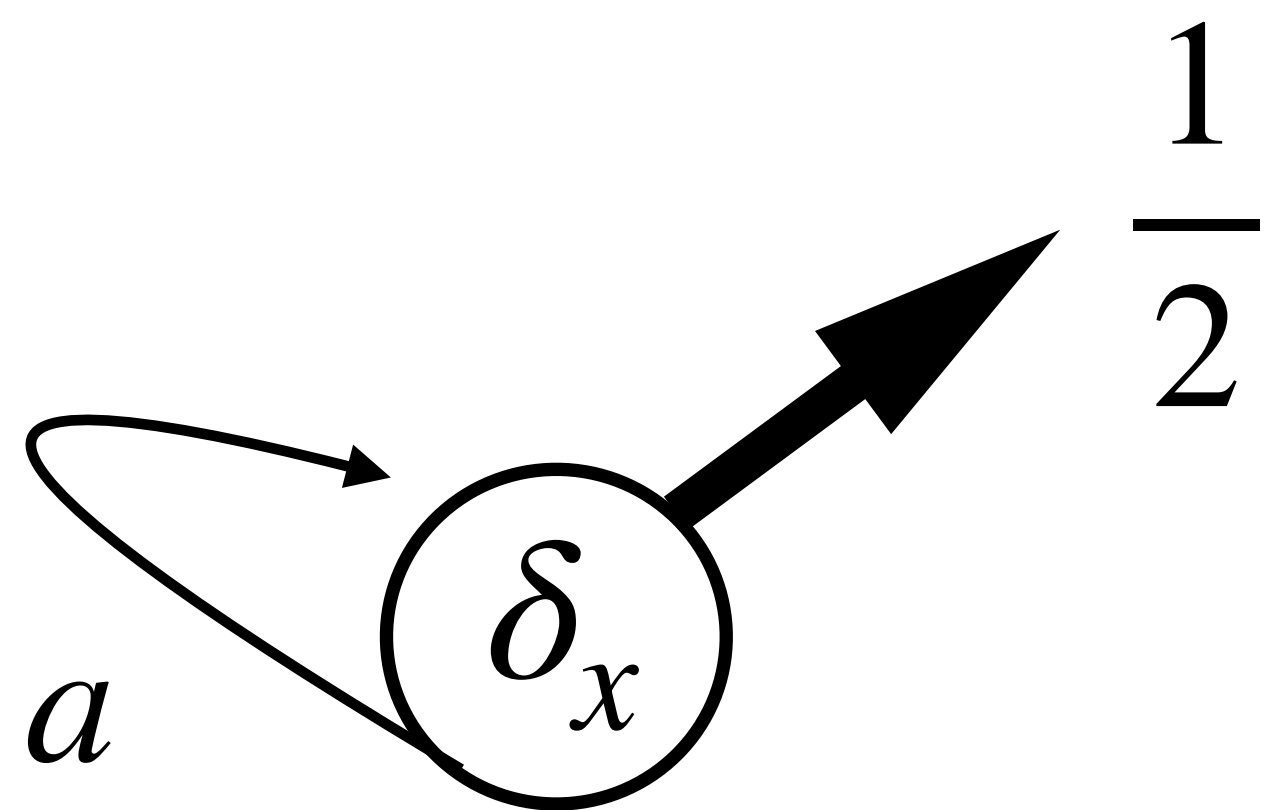




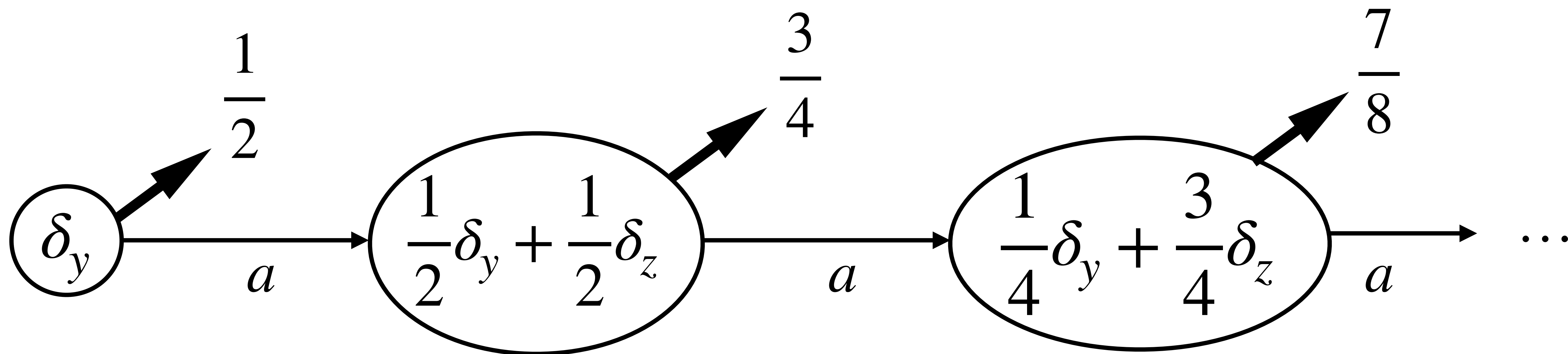


**Each state assigns a weight  
(expected payoff) to a word**





	$\epsilon$	$a$	$aa$	$\dots$
$\delta_x$	$\frac{1}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	$\dots$
$\delta_y$	$\frac{1}{2}$	$\frac{3}{4}$	$\frac{7}{8}$	$\dots$



**Behavioural distances via lifting**



$F: \mathbf{Set} \rightarrow \mathbf{Set}$

Endofunctor describing one-  
step behaviour

$F: \mathbf{Set} \rightarrow \mathbf{Set}$

$d$  is a pseudometric

Endofunctor describing one-  
step behaviour

$$\overline{F}(X, d: X \times X \rightarrow [0,1]) = (FX, d^F: FX \times FX \rightarrow [0,1])$$

Given a pseudometric  $d$  on the set of states, we can make a **new** one:  $\text{beh}(d)$

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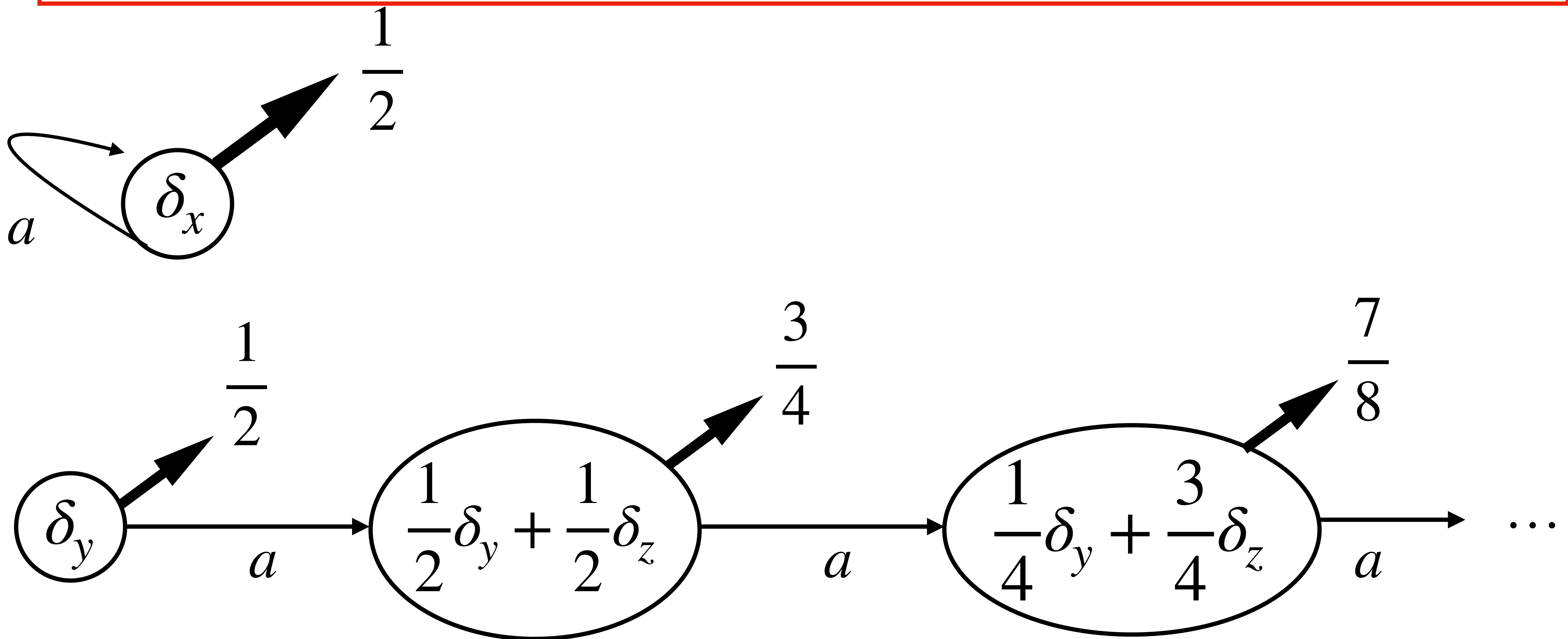
**Least** fixpoint  $\mu \text{beh}$  gives as the behavioural distance



$$F = [0,1] \times X^A$$

$$d^F(\langle o_1, t_1 \rangle, \langle o_1, t_2 \rangle) = \max\{ |o_1 - o_2|, \max_{a \in A} d(t_1(a), t_2(a)) \}$$

Is it true that  $d(\delta_x, \delta_y) \leq \frac{1}{2}$  and  $d(\delta_x, \delta_z) \leq \frac{1}{2}$ ?



Induction proof  
principle

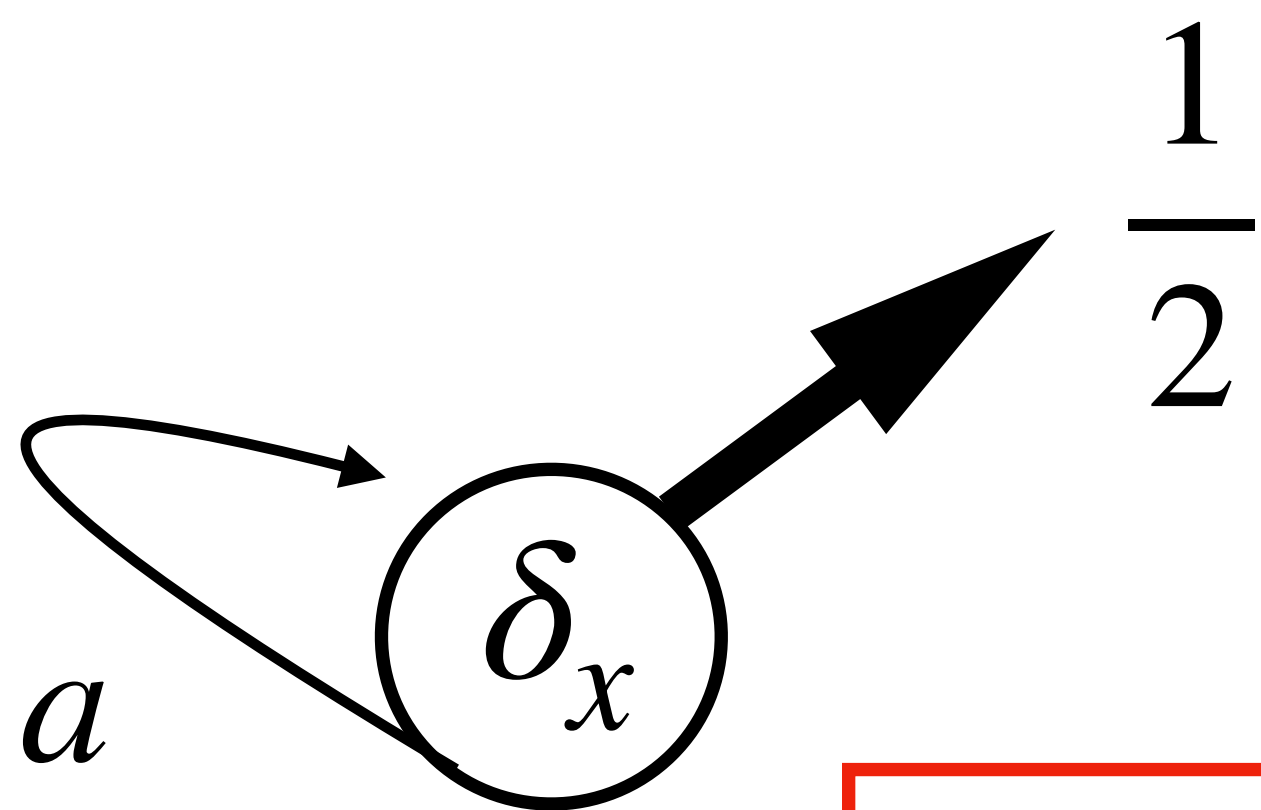
$$\text{beh}(d) \leq d$$

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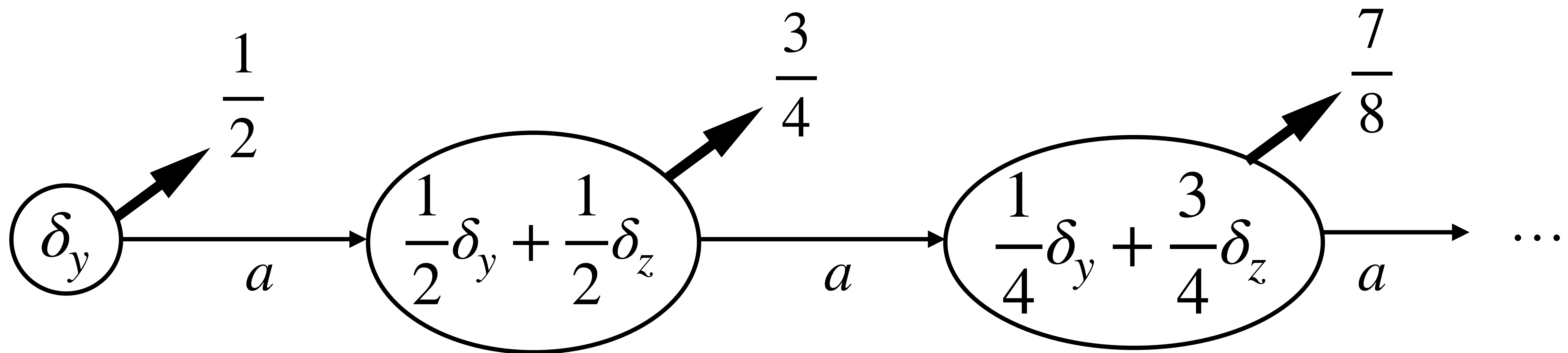
$$\mu\text{beh} \leq d$$

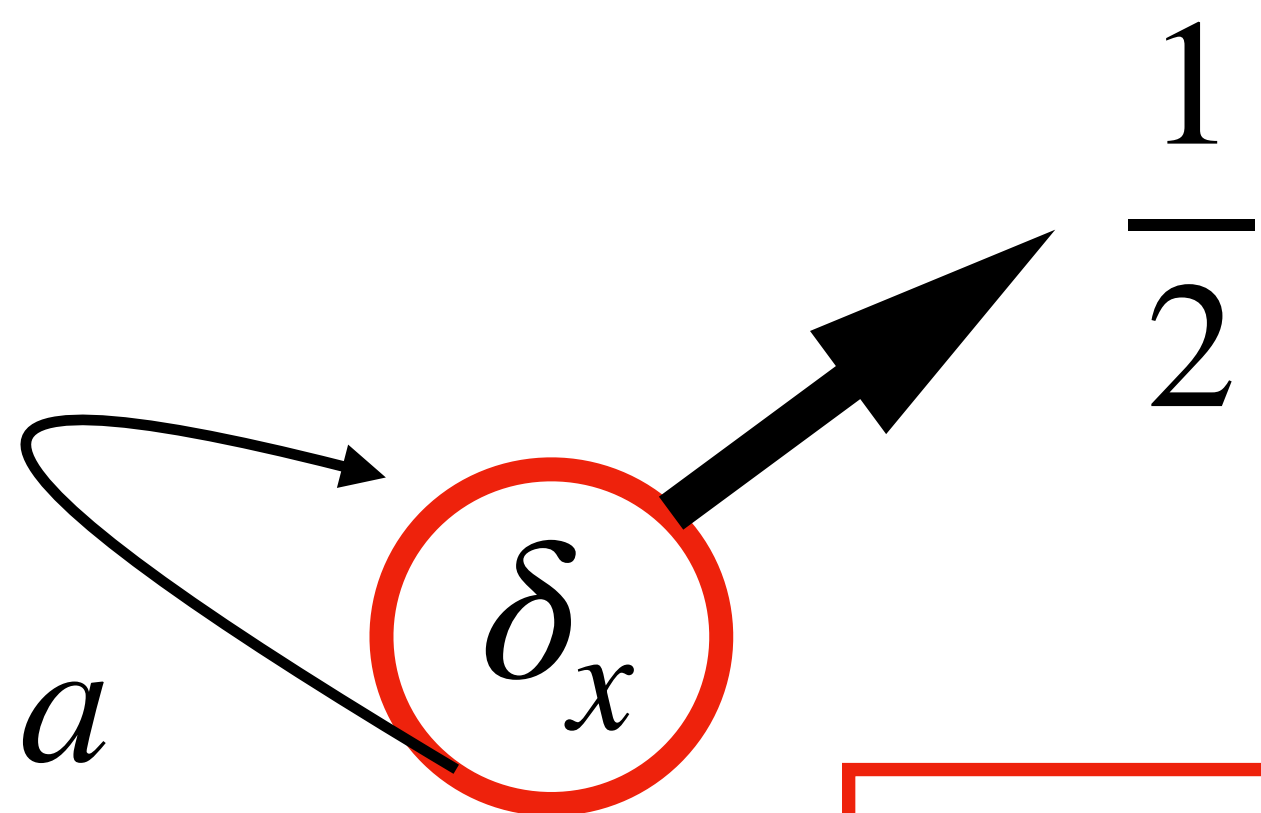
Witness of the upper  
bound

$$d(\delta_x, \delta_y) = \frac{1}{2}, d(\delta_x, \delta_z) = \frac{1}{2}, d(p, q) = 1$$

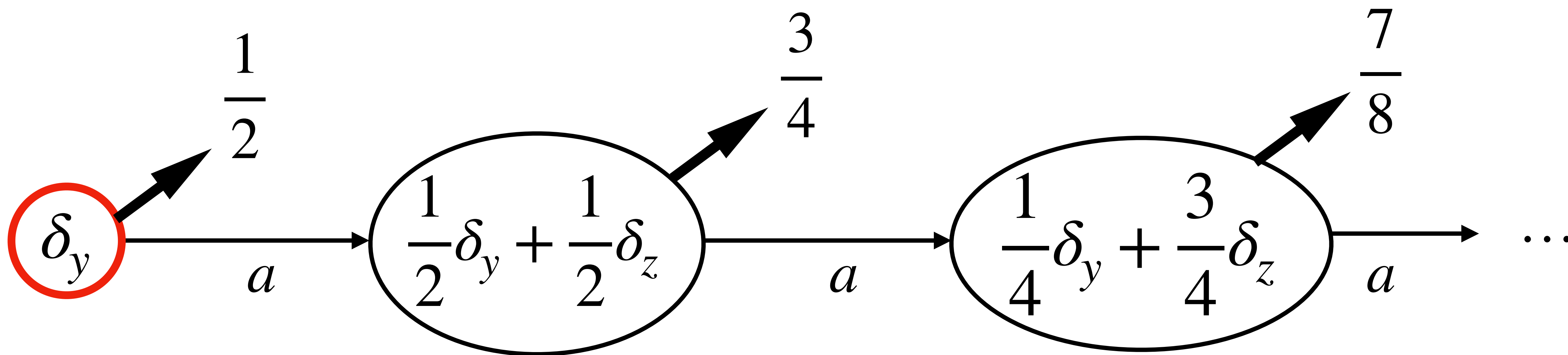


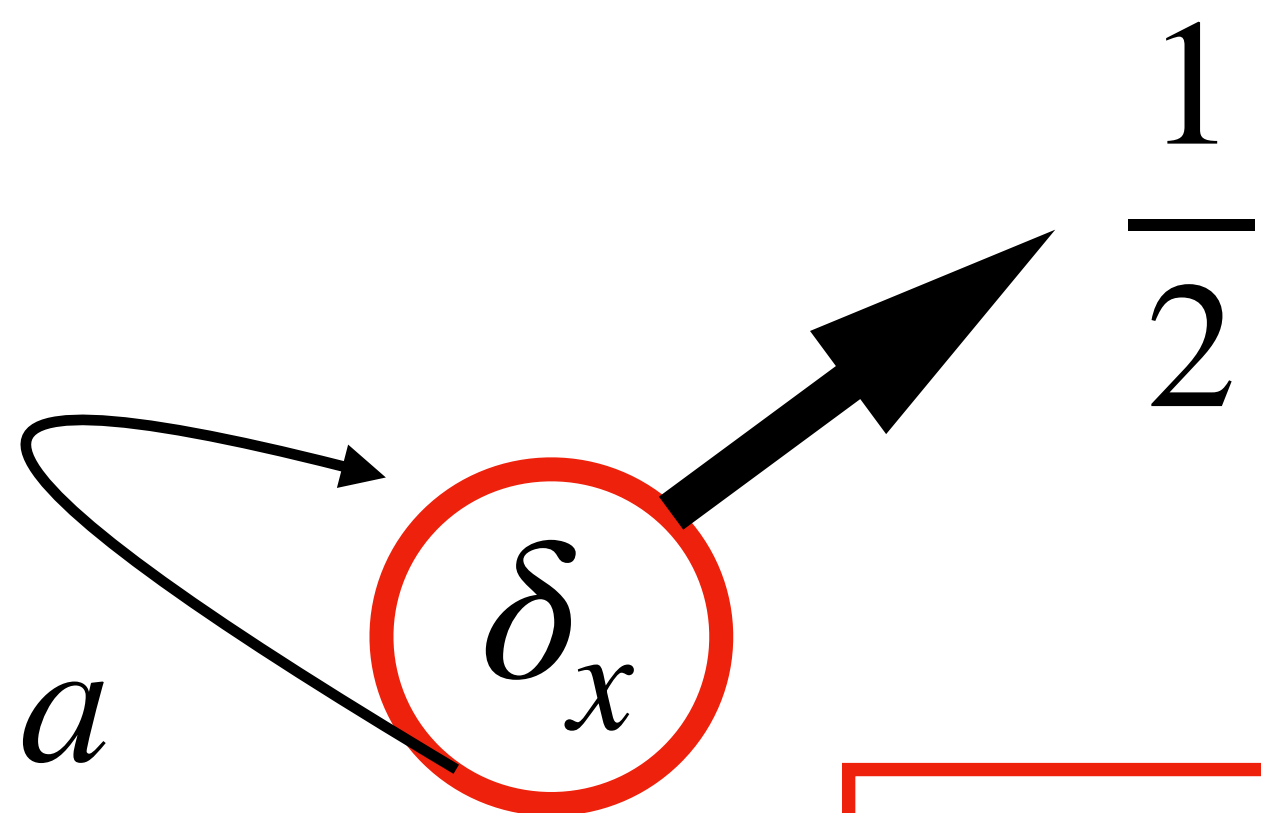
$$\text{beh}(d)(\delta_x, \delta_y) = \max\left\{\frac{1}{2} - \frac{1}{2}, d\left(\delta_x, \frac{1}{2}\delta_y + \frac{1}{2}\delta_z\right)\right\}$$



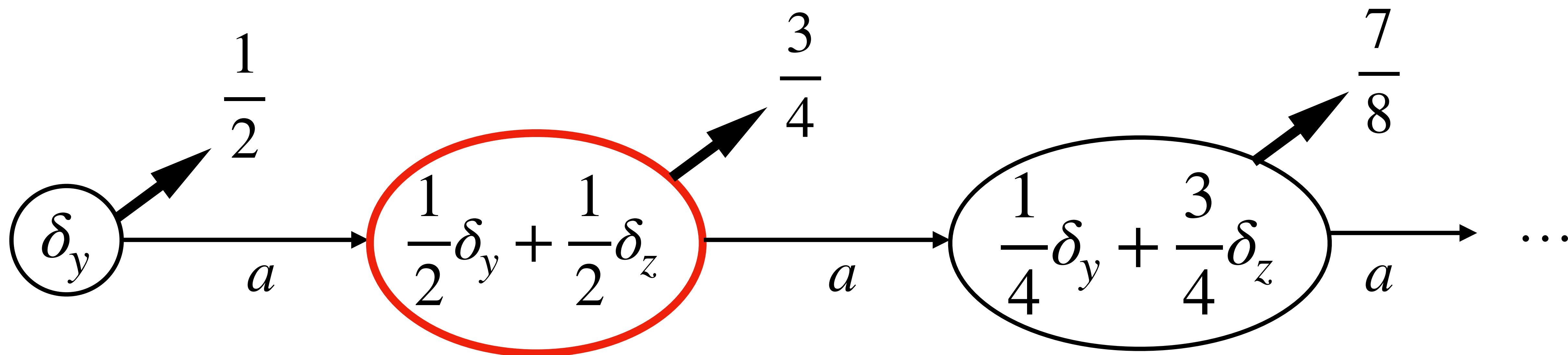


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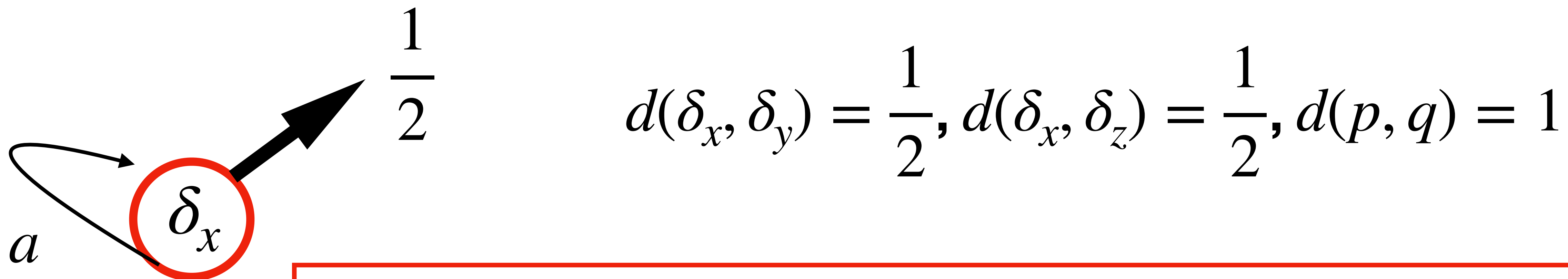




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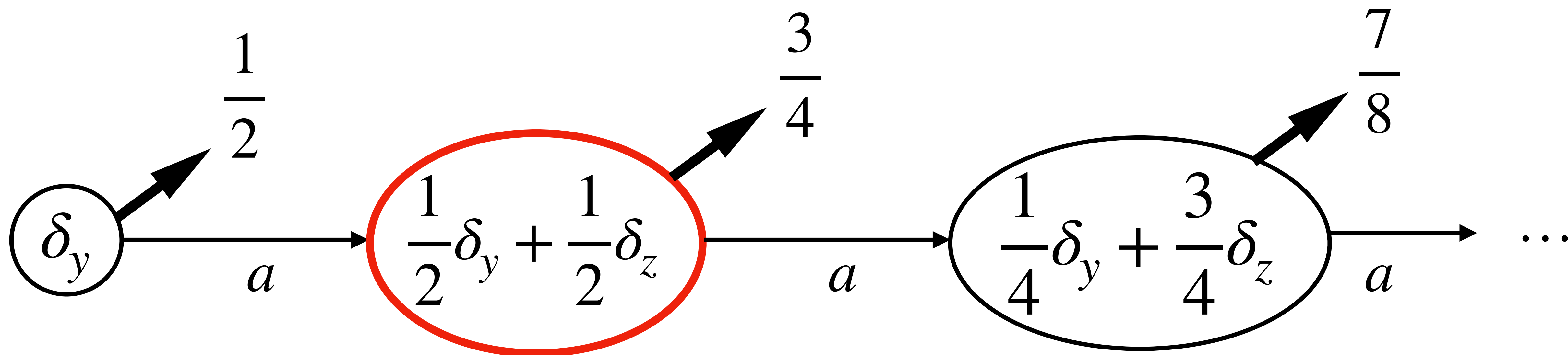






$$d(\delta_x, \delta_y) = \frac{1}{2}, d(\delta_x, \delta_z) = \frac{1}{2}, d(p, q) = 1$$

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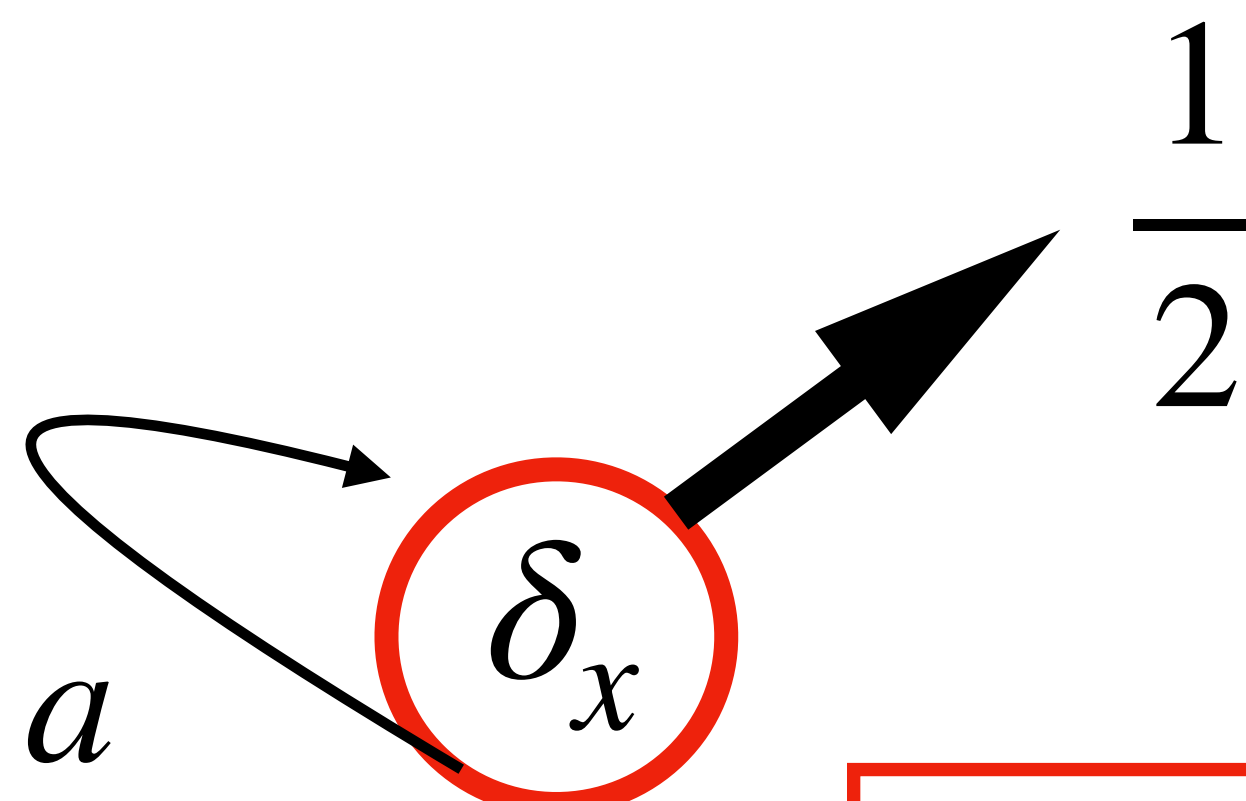
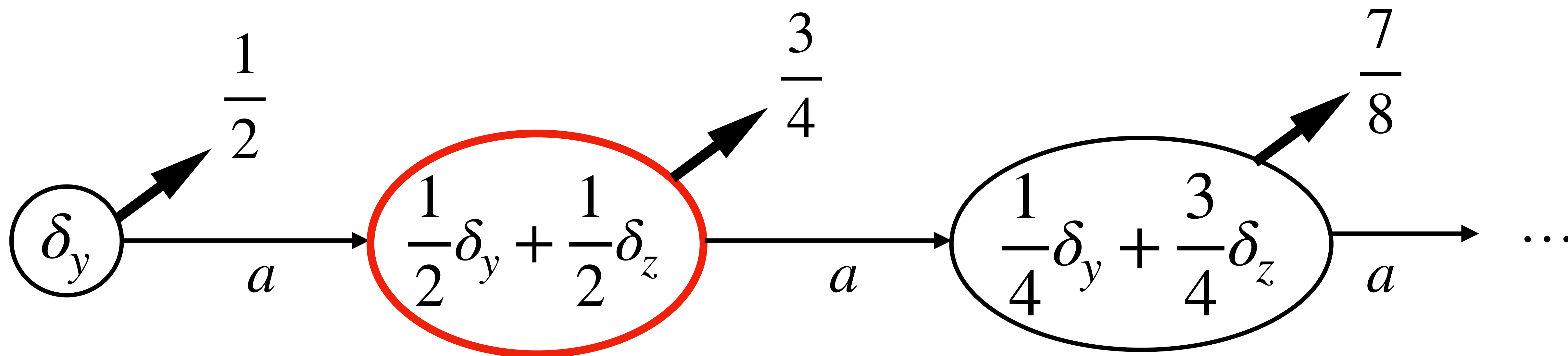


Diagram showing a node  $\delta_x$  (circled in red) with a self-loop labeled  $a$  and an outgoing arrow labeled  $\frac{1}{2}$ .

$$d(\delta_x, \delta_y) = \frac{1}{2}, d(\delta_x, \delta_z) = \frac{1}{2}, d(p, q) = 1$$

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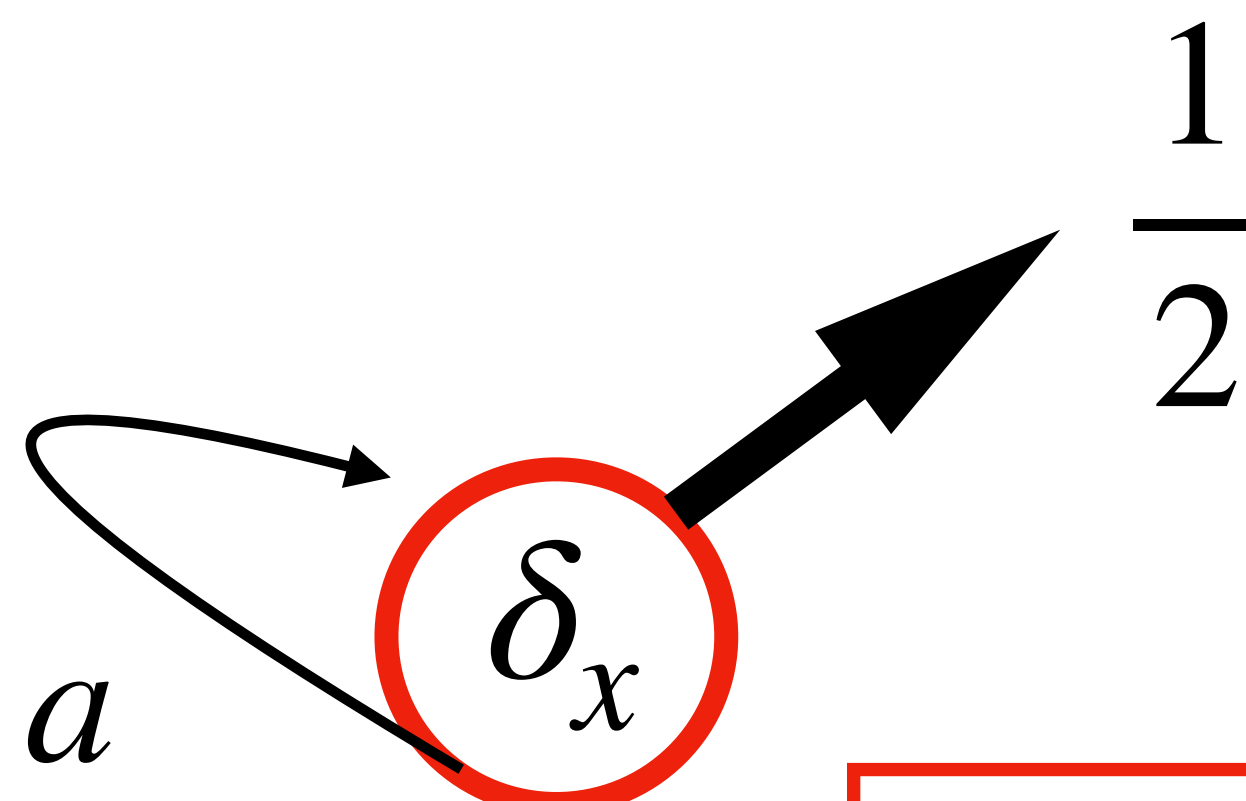
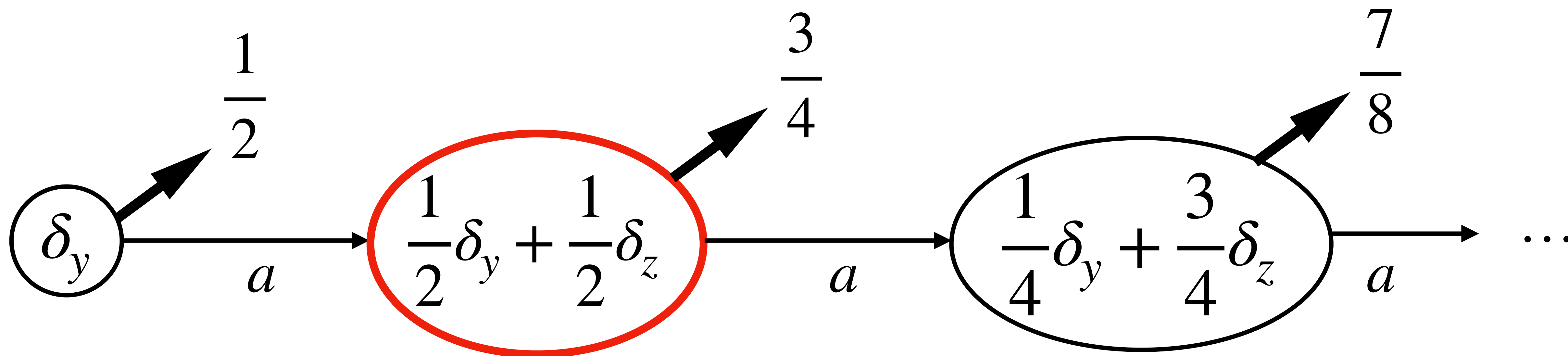


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$$d(\delta_x, \delta_y) \leq \frac{1}{2} \textbf{ and } d(\delta_x, \delta_z) \leq \frac{1}{2}$$

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$$d(\delta_x, p\delta_y + (1 - p)\delta_z) \leq p\frac{1}{2} + (1 - p)\frac{1}{2} = \frac{1}{2}$$

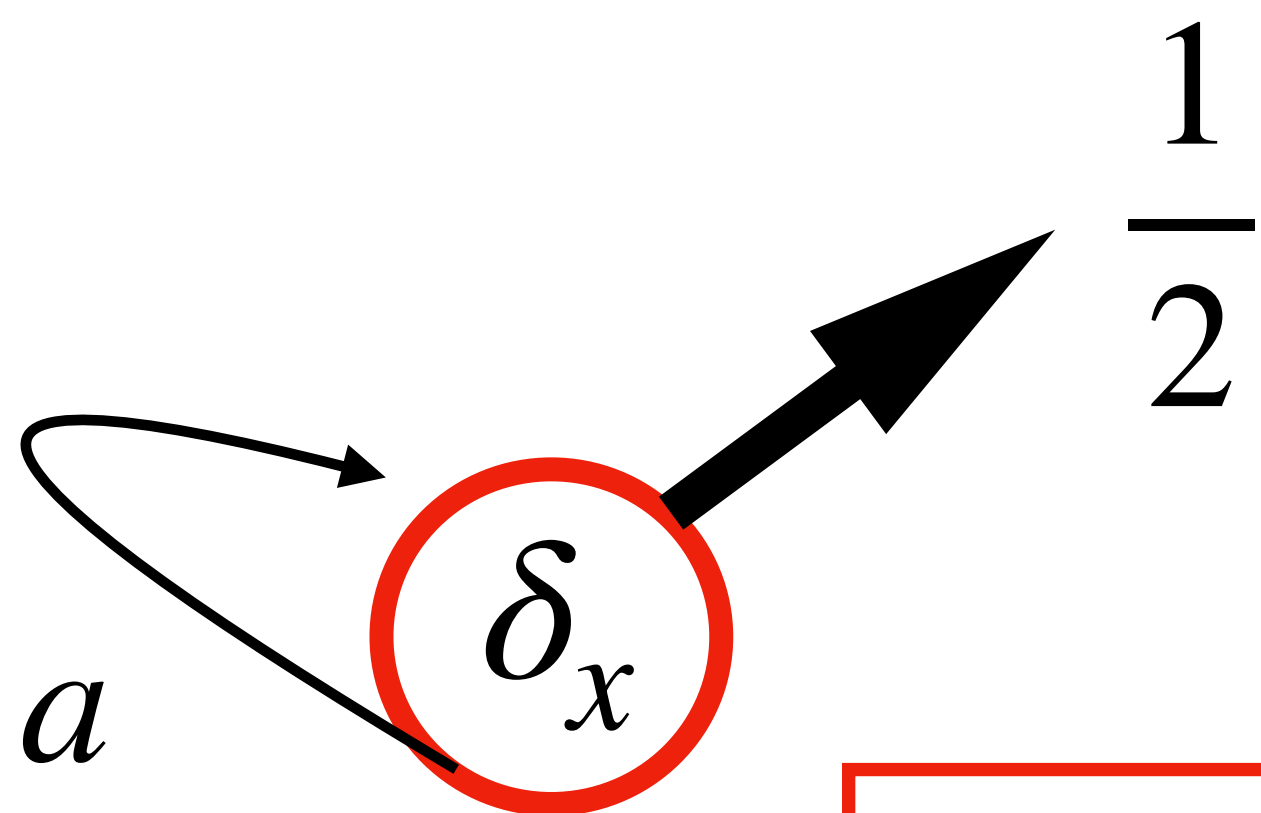
$$\frac{\text{beh}(u(d)) \leq d}{\mu\text{beh} \leq d}$$

Closure under  
algebraic structure

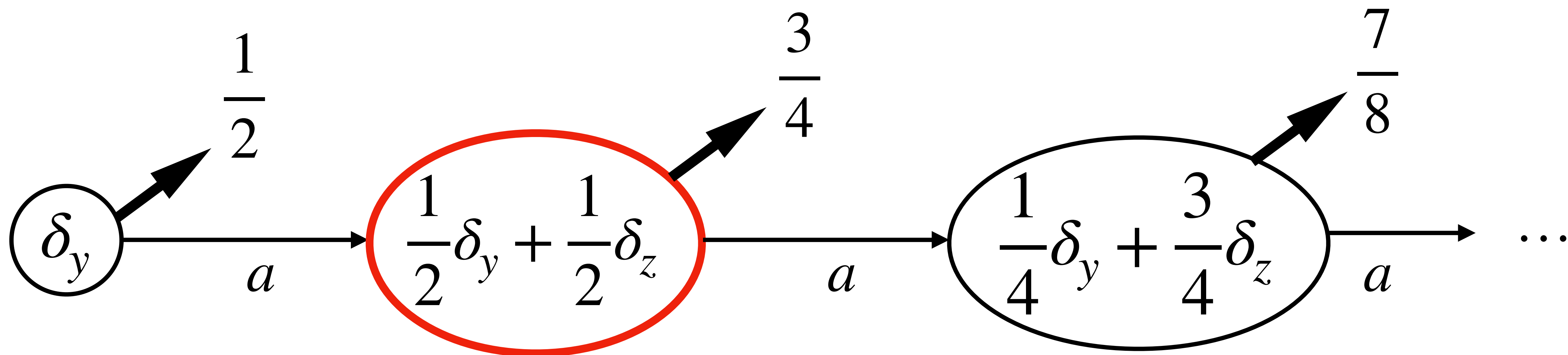
$$\text{beh}(u(d)) \leq d$$

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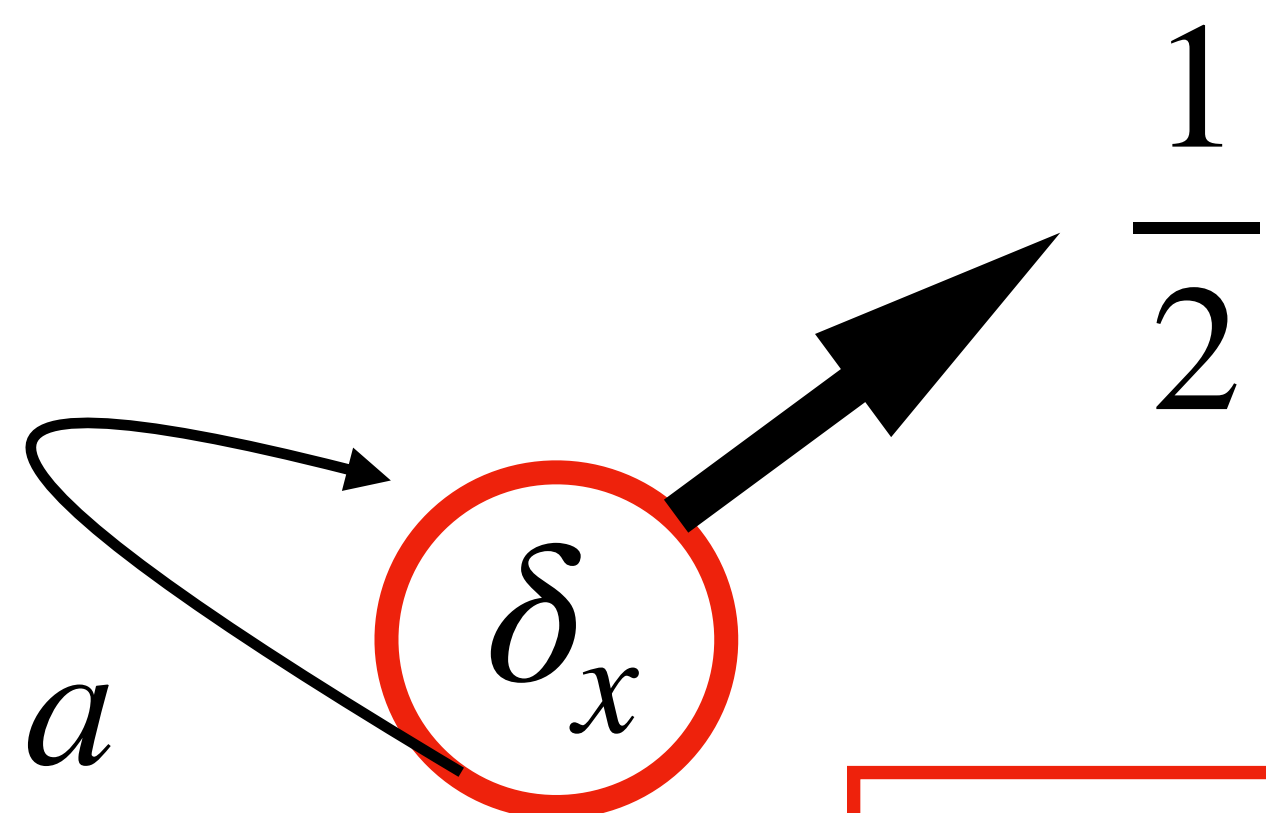
$$\mu\text{beh} \leq d$$



$$\text{beh}(u(d))(\delta_x, \delta_y) = \max\left\{\frac{1}{2} - \frac{1}{2}, d\left(\delta_x, \frac{1}{2}\delta_y + \frac{1}{2}\delta_z\right)\right\} \leq \frac{1}{2}$$

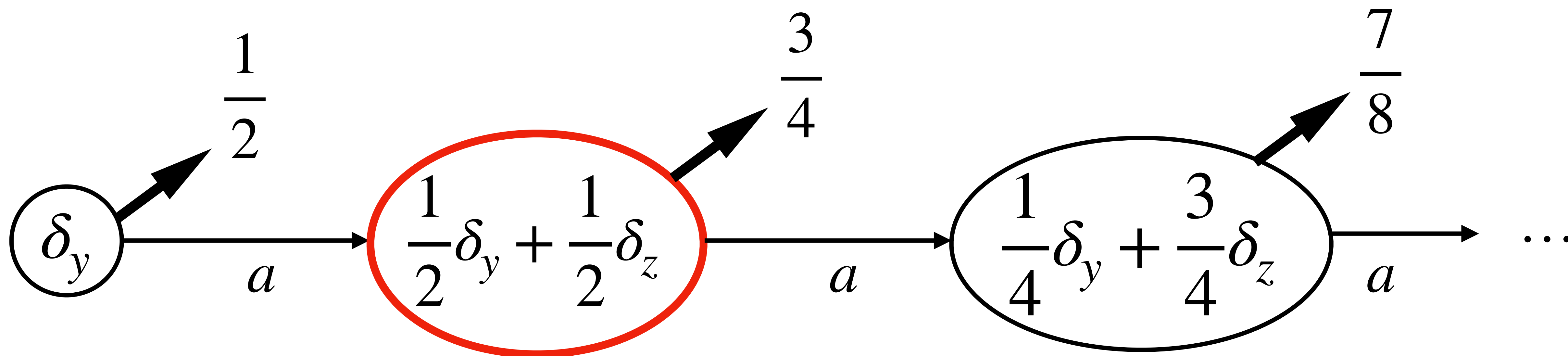






Now it works!

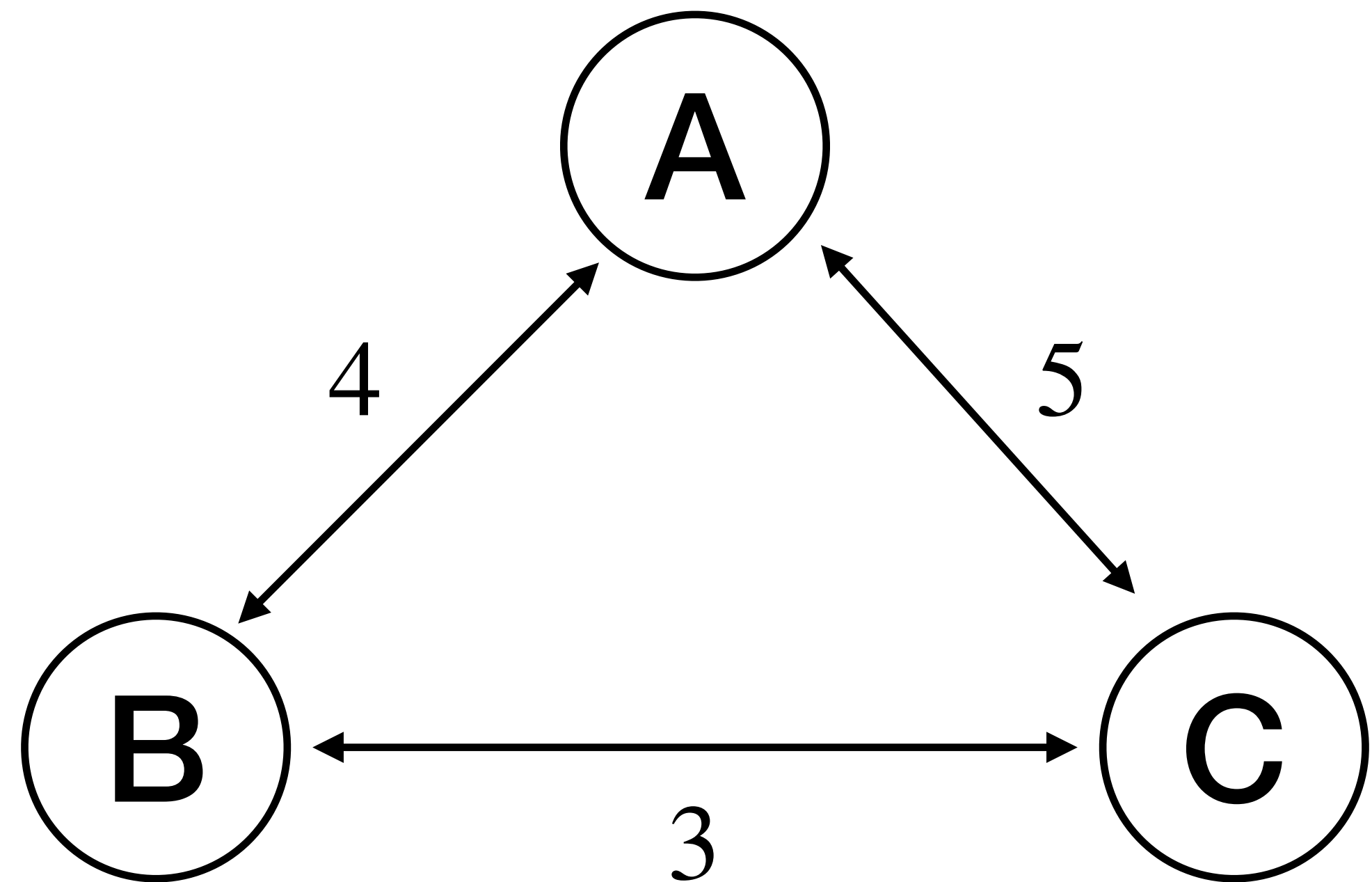
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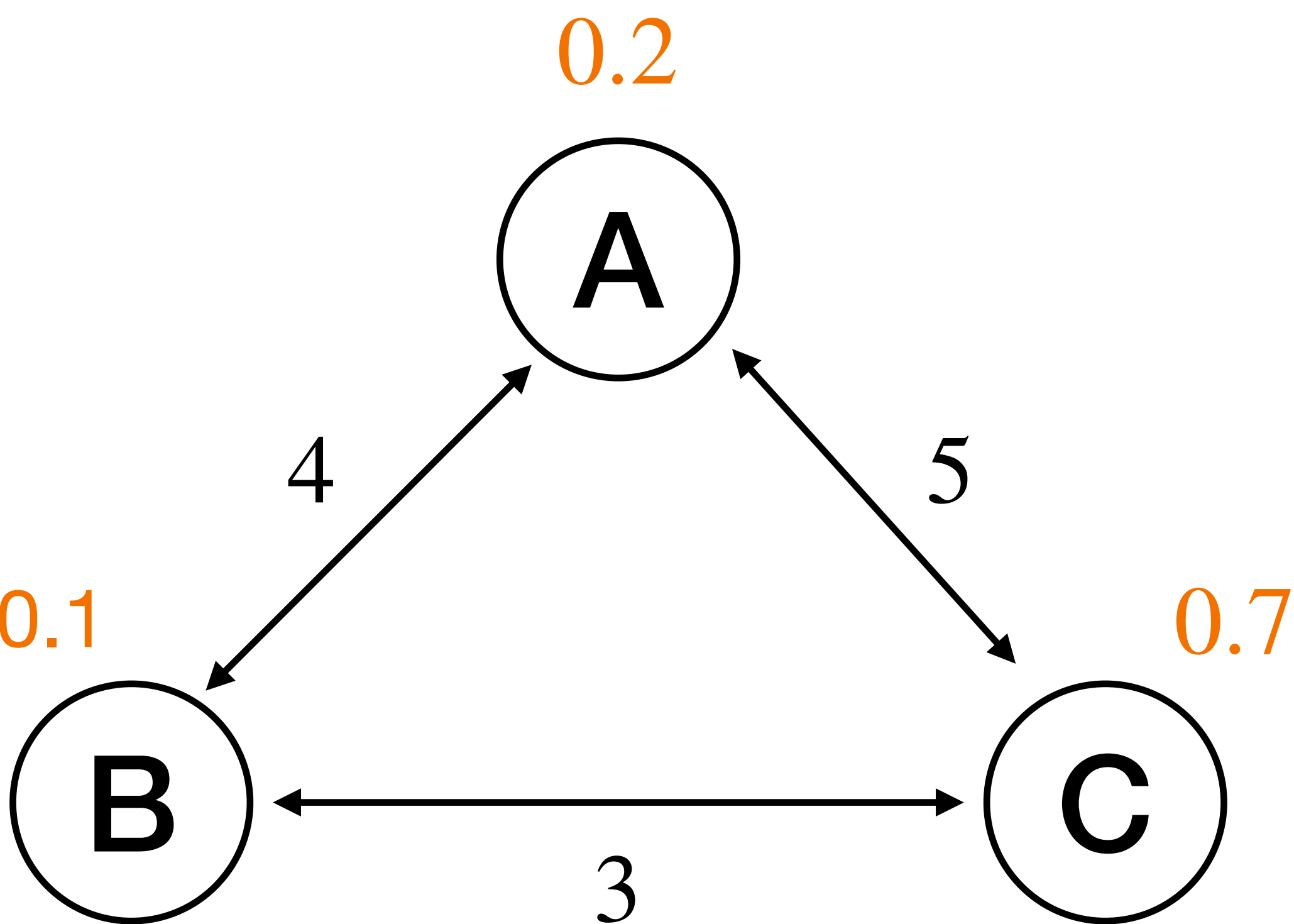
If we can lift  $\lambda : TF \Rightarrow FT$  into  $\lambda : \overline{TF} \Rightarrow \overline{FT}$ ,  
then the up-to technique is **sound**

# Kantorovich lifting

Going from distances on  $X$  to  $\mathcal{D}X$

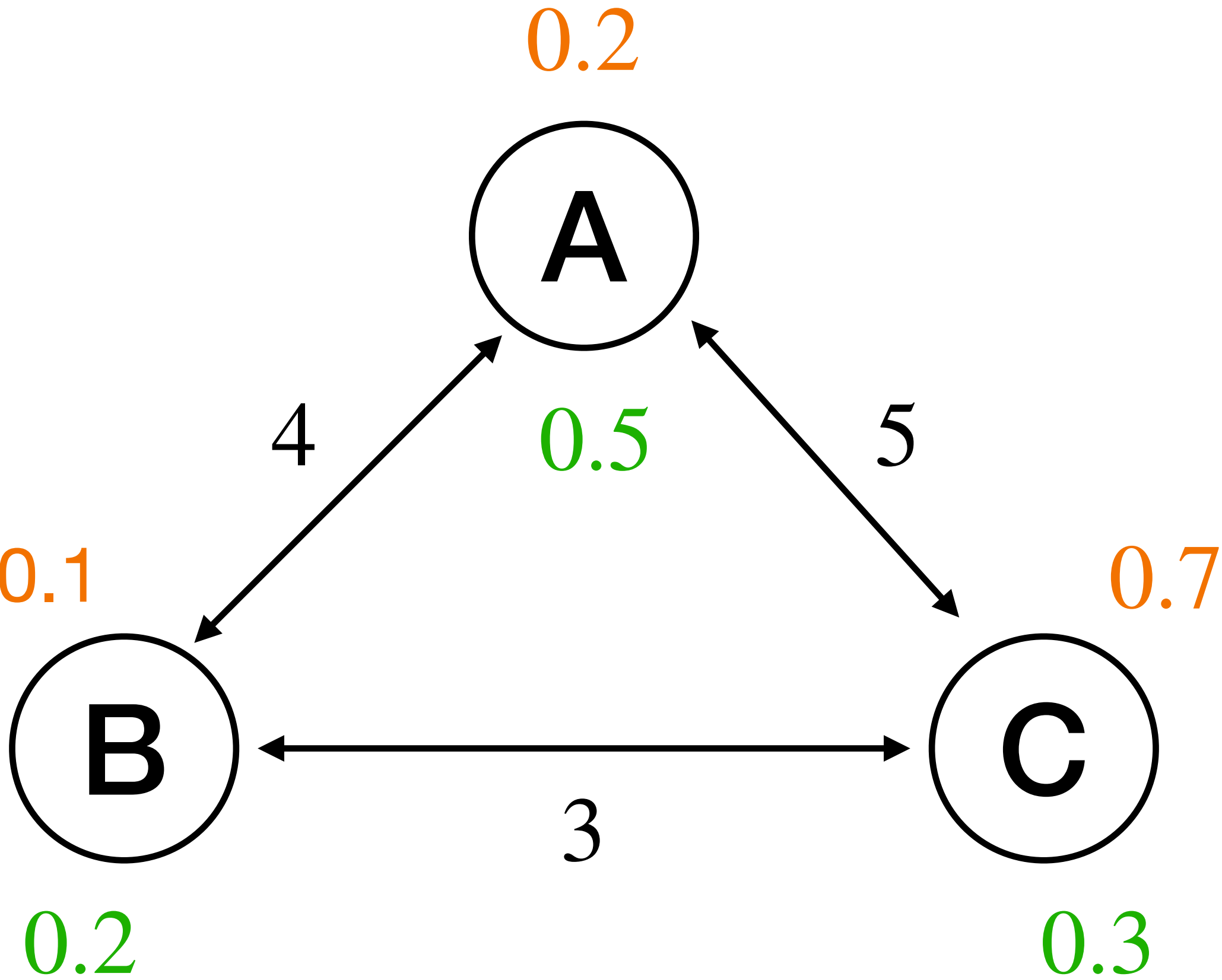


Going from distances on  $X$  to  $\mathcal{D}X$



Supply

Going from distances on  $X$  to  $\mathcal{D}X$



Supply

Demand

$$[0,1]^{X\times X} \xrightarrow{\gamma_X} \mathcal{P}([0,1]^X) \xrightarrow{\mathcal{P}(\mathbb{E}_X)} \mathcal{P}([0,1]^{\mathcal{D}X}) \xrightarrow{\alpha_X} [0,1]^{\mathcal{D}X\times\mathcal{D}X}$$

$$[0,1]^{X \times X} \xrightarrow{\gamma_X} \mathcal{P}([0,1]^X) \xrightarrow{\mathcal{P}(\mathbb{E}_X)} \mathcal{P}([0,1]^{\mathcal{D}X}) \xrightarrow{\alpha_X} [0,1]^{\mathcal{D}X \times \mathcal{D}X}$$

Generate price plans:  
nonexpansive functions  
 $(X, d) \rightarrow ([0,1], d_e)$



$$[0,1]^{X \times X} \xrightarrow{\gamma_X} \mathcal{P}([0,1]^X) \xrightarrow{\mathcal{P}(\mathbb{E}_X)} \mathcal{P}([0,1]^{\mathcal{D}X}) \xrightarrow{\alpha_X} [0,1]^{\mathcal{D}X \times \mathcal{D}X}$$

Generate price plans:  
nonexpansive functions  
 $(X, d) \rightarrow ([0,1], d_e)$

Evaluate the expected  
value for each plan

$$[0,1]^{X \times X} \xrightarrow{\gamma_X} \mathcal{P}([0,1]^X) \xrightarrow{\mathcal{P}(\mathbb{E}_X)} \mathcal{P}([0,1]^{\mathcal{D}X}) \xrightarrow{\alpha_X} [0,1]^{\mathcal{D}X \times \mathcal{D}X}$$

Use the plan that  
maximises the profit

Generate price plans:  
nonexpansive functions  
 $(X, d) \rightarrow ([0,1], d_e)$

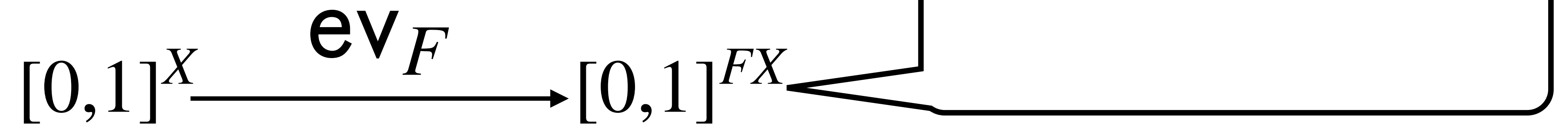
Evaluate the expected  
value for each plan

$$[0,1]^X \xrightarrow{\mathbb{E}_X} [0,1]^{\mathcal{D}X}$$

$$[0,1]^X \xrightarrow{\mathbf{ev}_F} [0,1]^{FX}$$

Goncharov, S., Hofmann, S., Nora, P., Schröder, L., Wild, P. (2023) “Kantorovich Functors and Characteristic Logics for Behavioural Distances”

Beohar, H., Gurke, S., König, B., Messing, K., Forster, J., Schröder, L. and Wild, P. (2024). Expressive Quantale-Valued Logics for Coalgebras: An Adjunction-Based Approach.

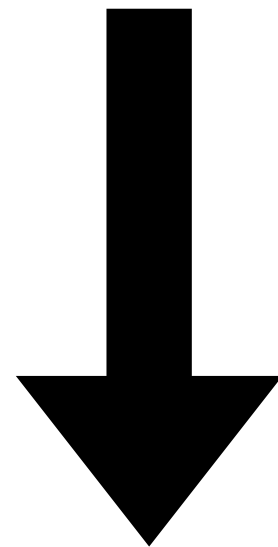


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$$[0,1]^X \xrightarrow{\mathbf{ev}_F} [0,1]^{FX}$$

Test one-step behaviour



$$[0,1]^X \xrightarrow{\Lambda_X^F} \mathcal{P}([0,1]^{FX})$$

We can have multiple  
evaluation maps

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$$d: X \times X \rightarrow [0,1]$$

$$0 \leq d(x, x)$$

$$d(x, y) = d(y, x)$$

$$d(x, z) \leq d(x, y) + d(y, z)$$

Complete lattice

$$(\mathcal{V}, \sqcup)$$

Monoid

$$(\mathcal{V}, \otimes, k)$$



Complete lattice

$$(\mathcal{V}, \sqcup)$$

$$([0,1], \oplus)$$

Monoid

$$(\mathcal{V}, \otimes, k)$$

Complete lattice

$$(\mathcal{V}, \sqcup)$$

Monoid

$$(\mathcal{V}, \otimes, k)$$

$$([0,1], \oplus)$$

$$(\{0,1\}, \wedge)$$

Complete lattice

$$(\mathcal{V}, \sqcup)$$

Monoid

$$(\mathcal{V}, \otimes, k)$$

$$([0,1], \oplus)$$

$$(\{0,1\}, \wedge)$$

$$([0,\infty], \max)$$

$$d\colon X\times X\rightarrow\mathcal{V}$$

$$k\sqsubseteq d(x,x)$$

$$d(x,y)=d(y,x)$$

$$d(x,z)\sqsubseteq d(x,y)\otimes d(y,z)$$

$$d: X \times X \rightarrow \mathcal{V}$$

$$k \sqsubseteq d(x, x)$$

~~$$d(x, y) = d(y, x)$$~~

$$d(x, z) \sqsubseteq d(x, y) \otimes d(y, z)$$

$$d: X \times X \rightarrow \mathcal{V}$$

$$k \sqsupseteq d(x, x)$$

~~$$d(x, y) = d(y, x)$$~~

$$d(x, z) \sqsupseteq d(x, y) \otimes d(y, z)$$

Directed metrics

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For Kantorovich, we know that  $\overline{T} \overline{F} \sqsubseteq \overline{TF}$  for free

# Compositionality

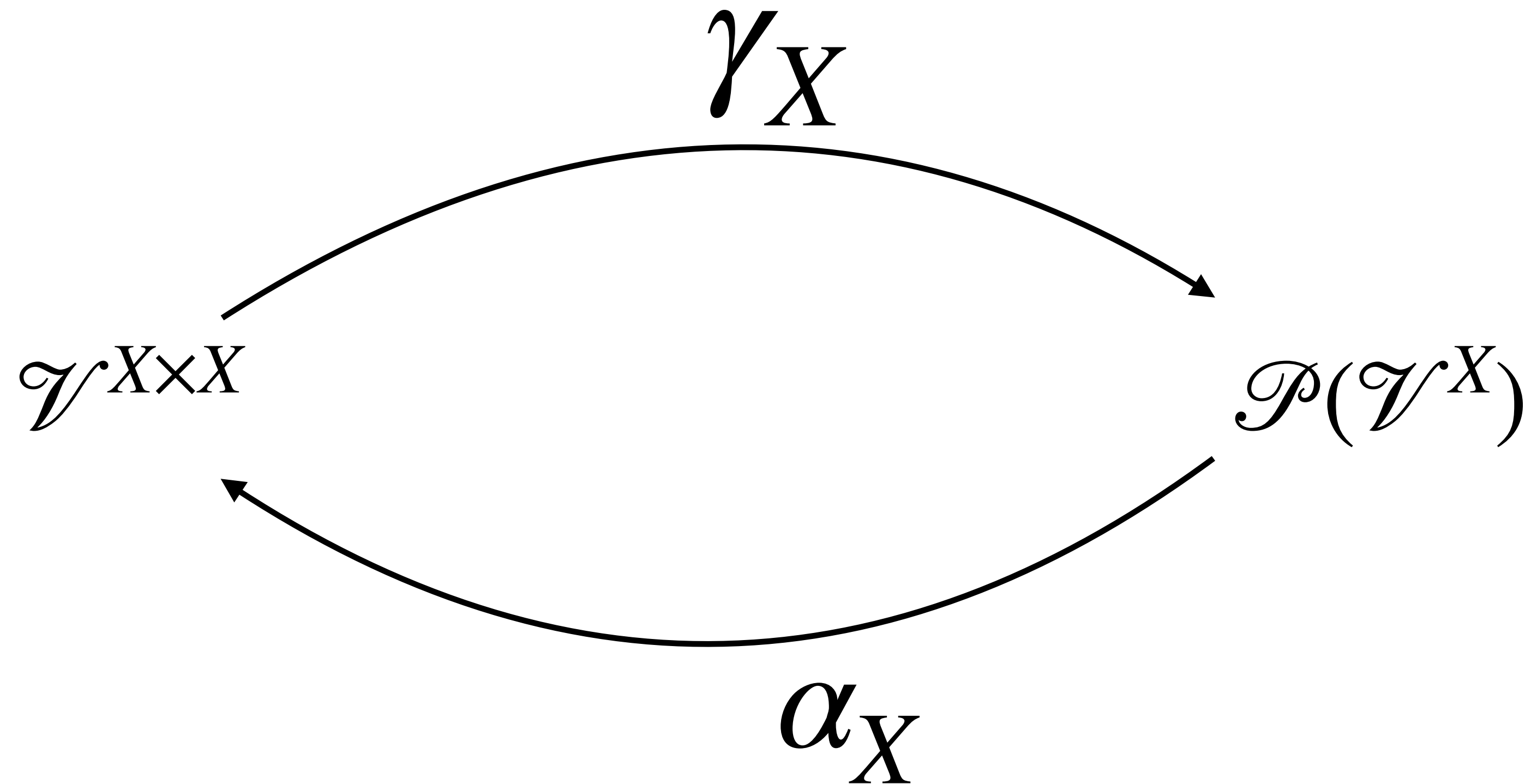
$$\overline{F} \overline{T} = \overline{FT}$$

# Sound up-to technique for $FT$ -coalgebras

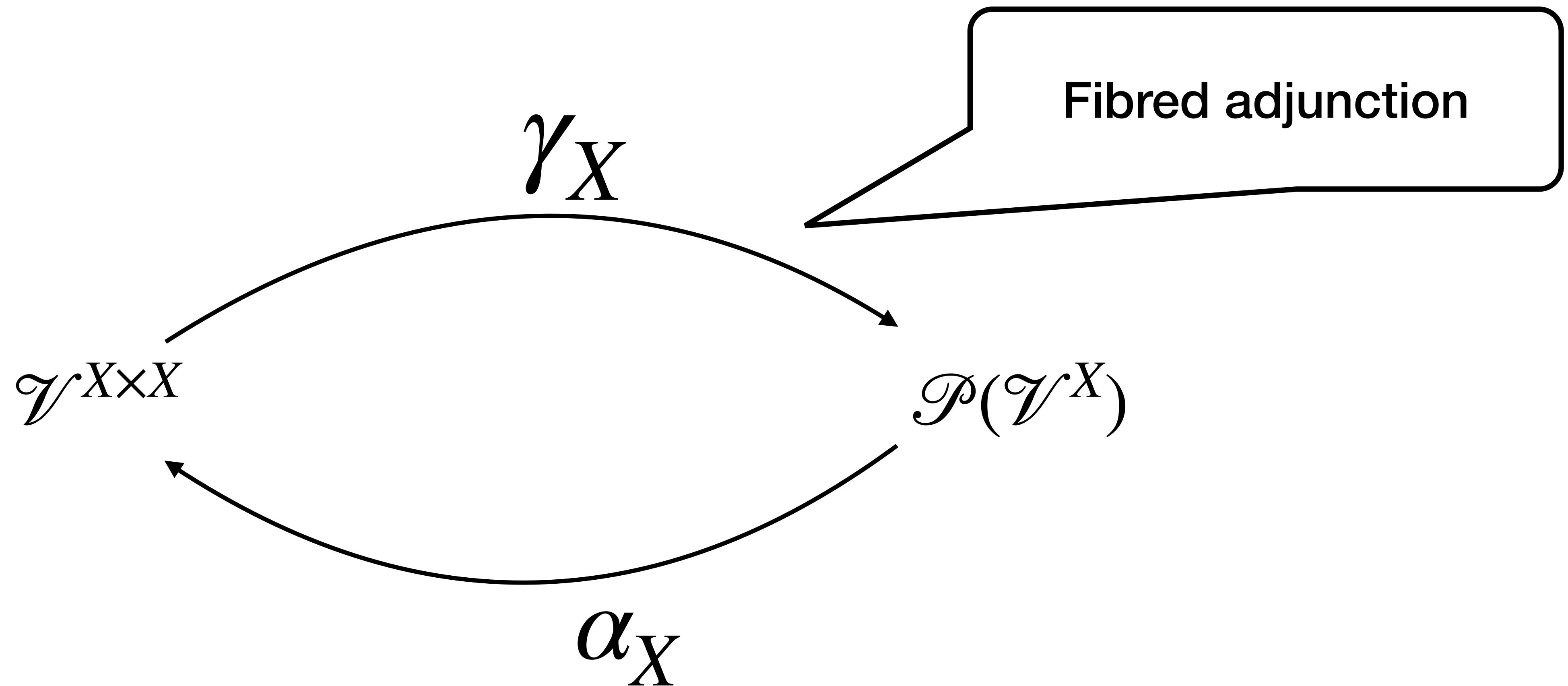
$$F = C_B \mid \text{Id} \mid \prod_{i \in I} F_i \mid F_1 + F_2$$

To make that work, we construct appropriate sets of evaluation maps that interact well with the distributive law

# Core results of the result come from fibred category theory



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# Quantale-valued behavioural distances for a wide class of transition systems

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**Kantorovich lifting is flexible and  
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coproducts**