

# Lecture 11)

## 11.1) Critical Paths and Scheduling.

Last time we stated Thm.

$\rho$ -regular vels  $\Rightarrow$  ell in long run

no ell in long run  $\Leftrightarrow \exists$  weights  
(non-negative, not all 0)  
~~properties~~ st the Contact

graph is free b-load.

The contact graph is the ~~graph~~  
~~spanning~~ set of ~~vertices~~

as edges on  $S^2$  between edges

exactly dist  $\theta$  and for

$$\theta = 2\pi$$

Thm  $\text{Conf}(N, \rho)$

$$= \text{Conf}(N; \frac{\theta}{2\pi})$$

This also gives an  
algorithm for sampling.

critical / maximal configurations

$\rho$ -regular

update...

conflating...

11.2) Examples with pictures.

$N=3$  } full probn  
 $N=4$

$N = \dots$  max flows.

$N = 10$  - local max?

$N = 12$  - local max?

higher critical probn.  
Jittering?

11.3)

Local maxes and standard  
entropy.

Answer

~~False~~: all 'jump' states, have  
equal probability.  
of equal prob

Related to the evolution of  
p flow. Also the model...  
not always true...?

~~ex. cancelled~~

# Boltzmann Measure

Discrete Setting:

$N$  particles, each in one of  $k$  states.  
Each state has an energy  $\epsilon_i, i \in \{1, \dots, k\}$

System closed so:

$$E = \sum_{i=1}^N \epsilon_i \quad n_i = \# \text{ particles in state } \epsilon_i$$

$$N = \sum_{i=1}^k n_i$$

With # of ways to place particles  $\{n_1, \dots, n_k\}$

is the multiplicity  $\Omega$ .

$$\frac{N!}{n_1! n_2! \dots n_k!}$$

$\Omega \approx$   
the optimal partition dominates  
for  $N$  large

determine partition  
that maximizes  $W$ .

replace with  $\log W$  [log function]

$$\log W = (N \ln N - N) - \sum_{i=1}^K (N_i \ln N_i - N_i) + \alpha$$

$$L_{\text{with}} \left[ f = \ln(w) + \alpha(N - \sum_{i=1}^K N_i) + \beta(E - \sum_{i=1}^K E_i N_i) \right] \text{ const. } \dots$$

$$0 = \frac{\partial f}{\partial N_i} = - \ln N_i - (\alpha + \beta E_i)$$

$$\frac{\partial}{\partial N_i} (N_i \ln N_i + \alpha N_i + \beta E_i N_i)$$

$$-1 + \ln N - (\alpha + \beta E_i)$$

$$- \ln N_i - (\alpha + \beta E_i) = 0$$

$$N_i = e^{-\alpha - \beta E_i}$$

$$\overline{\overline{\frac{\partial \ln Z}{\partial \beta}}} = \frac{Z}{\frac{\partial \ln Z}{\partial \beta}} = \frac{Z}{\sum_k E_k e^{-\beta E_k}} =$$

$$\overline{E} = \frac{1}{Z} \sum_k E_k e^{-\beta E_k}$$

$$\Pr(\text{state } i) = \frac{e^{-\beta E_i}}{Z}$$

$$F = -T \ln Z = \text{Free energy}$$

$$\beta = \frac{1}{T} \cdot \frac{1}{k_B}$$

$$Z = \sum_k e^{-\beta E_k}$$

$$\frac{1}{Z} \sum_k E_k e^{-\beta E_k} = \overline{E}$$

Lemma:  $-\ln Z$  convex:

Show:  $-\frac{\partial}{\partial \beta} \ln Z$  monotone.

why?  $\frac{\partial^2}{\partial \beta^2} \ln Z > 0$

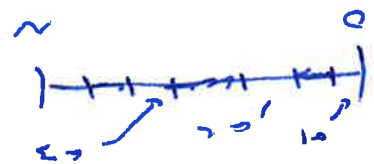
→ with (1) variance of  $E > 0$ .

$\frac{1}{\beta} \ln Z := \text{Free energy}$ .

$T \rightarrow \infty \quad P_i(t) \rightarrow \frac{1}{K}$

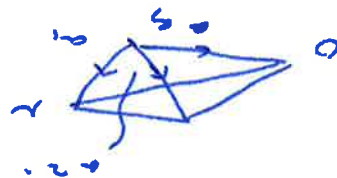
Hard disc model.

input



$$\{a_1, a_2, a_3\} = \{a_1, a_2, a_3\}$$

$$0 \leq a_1, a_2, a_3 \leq N-3$$



$$\frac{i}{\epsilon(N-3)} = \frac{\epsilon(N-3)}{\epsilon(N-3)}$$

input

$$\frac{i}{\epsilon(N-3)} = \frac{\epsilon(N-3)}{\epsilon(N-3)}$$

Give energy  $E_0$  for particle.

$$Z = e^{-\beta E_0} \text{ } \xrightarrow{\text{differentiate}}$$

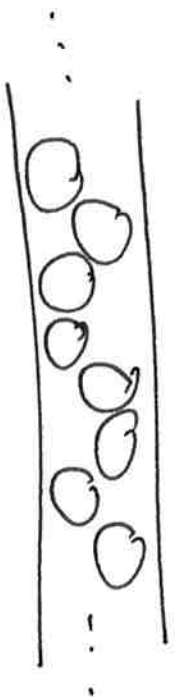
$$F_N(Z) = \sum_{k=0}^N Z^k \text{Vol}(\Omega_{N,k})$$

$$= \sum_{k=0}^N \binom{N-k}{k} Z^k$$

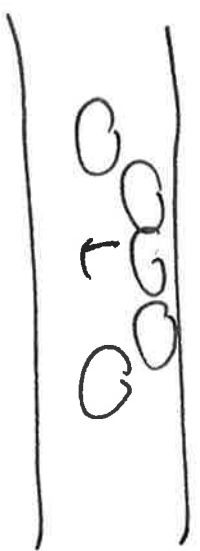


11.4) quasi-1D model with open ends.

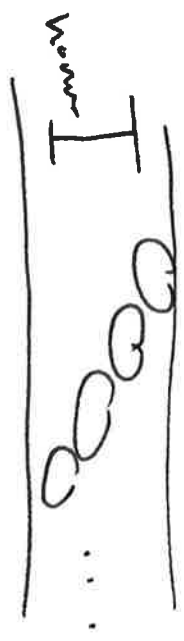
~~width~~ narrow channel: constant width.



— Density  
local maximal density.



centrally saddle



length is  $l$ .

$$\text{disks } n=1 \text{ channel } Z < \text{disk} < Z + \sqrt{3}.$$

Combinatorics: how many <sup>for</sup> maxima t?

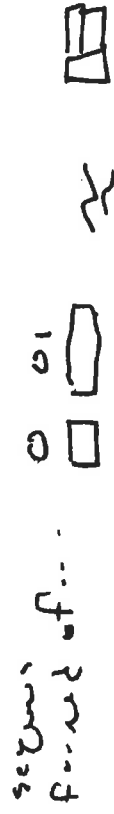
The length is determined by # of adjacent  
and opposite max.



(11) (0)

Counting is # of binary words with no  
adjacent 1's

This is eq to a tiling domino problem. stars and bars.



This satisfies the fibrec...

version. (node up/down symmetric.)

~~the contents of~~  $t$  of length  $l$



given exactly  $k$  is in a word of length  $l$ ,

$\Rightarrow$  inserting  $k$  is into a copy of  $l-k$

OS.  $0, 0, 0, \dots, 0$

$l-k$  into slots.

So for each  $k \Rightarrow \binom{2 \cdot t + 1}{k}$  class of nodes

So if  $\log_2 n \cdot f \in (1)^j$  is 2

and  $\in \mathcal{O}^*$  is  $L$

we have

$$t \cdot 2 = k \cdot 2 + (2 - k) L + 2$$

+2  
under  
rule.

we can also compute the max of rule nodes.

$$\frac{\binom{n-k+1}{k}}{k} \rightarrow \frac{5}{10} - \frac{\sqrt{5}}{10}$$

And then is "university" direct policy level.

00000...

in last line as well...

000000000000000000

}

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Ex. law of the & capital  
cur.

→ typical capital losses in

$$\frac{1}{2} \rightarrow \frac{55}{10} \quad (15) \quad (20)$$

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(Compare capitals ... ) → ( $\frac{1}{2}$ )

prints values of control pts  
flow to just after ...  
probably minimal size control  
section could be used 0 pages

==>

? complete the values!

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also.. for 1578 N

So why, in, why are: