

復旦大學



使用状态机的方式实现 booth 乘法器

课程名称: 集成电路设计基础

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微电子学院

2022 年 1 月 3 日

使用状态机的方式实现booth乘法器

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Booth乘法器前端设计

Booth乘法器原理

课堂上讲解的Booth乘法器原理如图

Booth 算法:免去特殊操作(减)

X被乘数, Y乘数

$$XY = \sum_{i=0}^{N-1} \underline{(-Y_i + Y_{(i-1)})} 2^i X = \sum_{i=0}^{N-1} PP_i 2^i \quad (3-5)$$

部分积 $PP_i = (-Y_i + Y_{(i-1)})X$

对无符号数乘法:

$$D = A * B$$

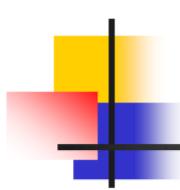
乘数 A 可表示为

$$A = \sum_{n=0}^{N-1} (A[n] * 2^n) \quad (3-10)$$

$$= \sum_{n=0}^{N-1} ((A[n] * B) * 2^n)$$

(3-12)

式(3-5)、式(3-12)表达形式是相同的,
无符号数与符号数的乘法就统一起来了



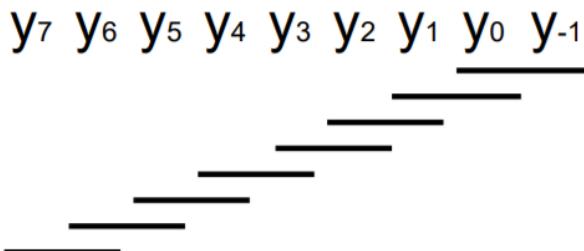
经典booth算法

- 经典的Booth算法（1951，A.D. Booth）主要是为了解决有符号运算中复杂的**符号修正问题**而研究的
 - 对于补码表示的两数相乘无需进行符号位**特殊操作**
 - 经典Booth算法通过每次在乘数中交叠地取两位(Y_{i+1}, Y_i)来产生部分积 $PP_i = (-Y_{i+1} + Y_i)X$, X 是被乘数

X被乘数,Y乘数

经典booth算法

$$PP_i = (-Y_{i+1} + Y_i)X$$



Booth算法取数操作

y_{i+1}	y_i	PP_i
0	0	+0
0	1	+X
1	0	-X
1	1	+0

表1.1：Booth算法规则

其中： $Y_{-1}=0$ 是附加考察位,帮助分析 Y_0
 X 是被乘数

Booth乘法器设计思路

Booth乘法器不会加快计算速度，只是单纯解决了有符号运算中复杂的符号修正问题，所以需要进行32次乘加操作。而这32次乘加操作我们可以在一个周期内完成，但这样会造成**关键路径延时过高**；我们也可以分成数级流水线，通过流水线的方式打断关键路径，但是可能会**面积和功耗升高，达不到要求**；所以我们选择使用**状态机**的方式完成Booth乘法器的设计，虽然会降低功耗和面积，也不会有很长时间的关键路径，但是**会有很长的输入输出延时**，但是Spec中并没有对此做出要求，所以我们认为这是一个不重要的参数，**无论输入输出延时多长都可以接受**。如下图：

以32 bit×32 bit有符号乘法器设计为例，进行半定制芯片设计流程训练。

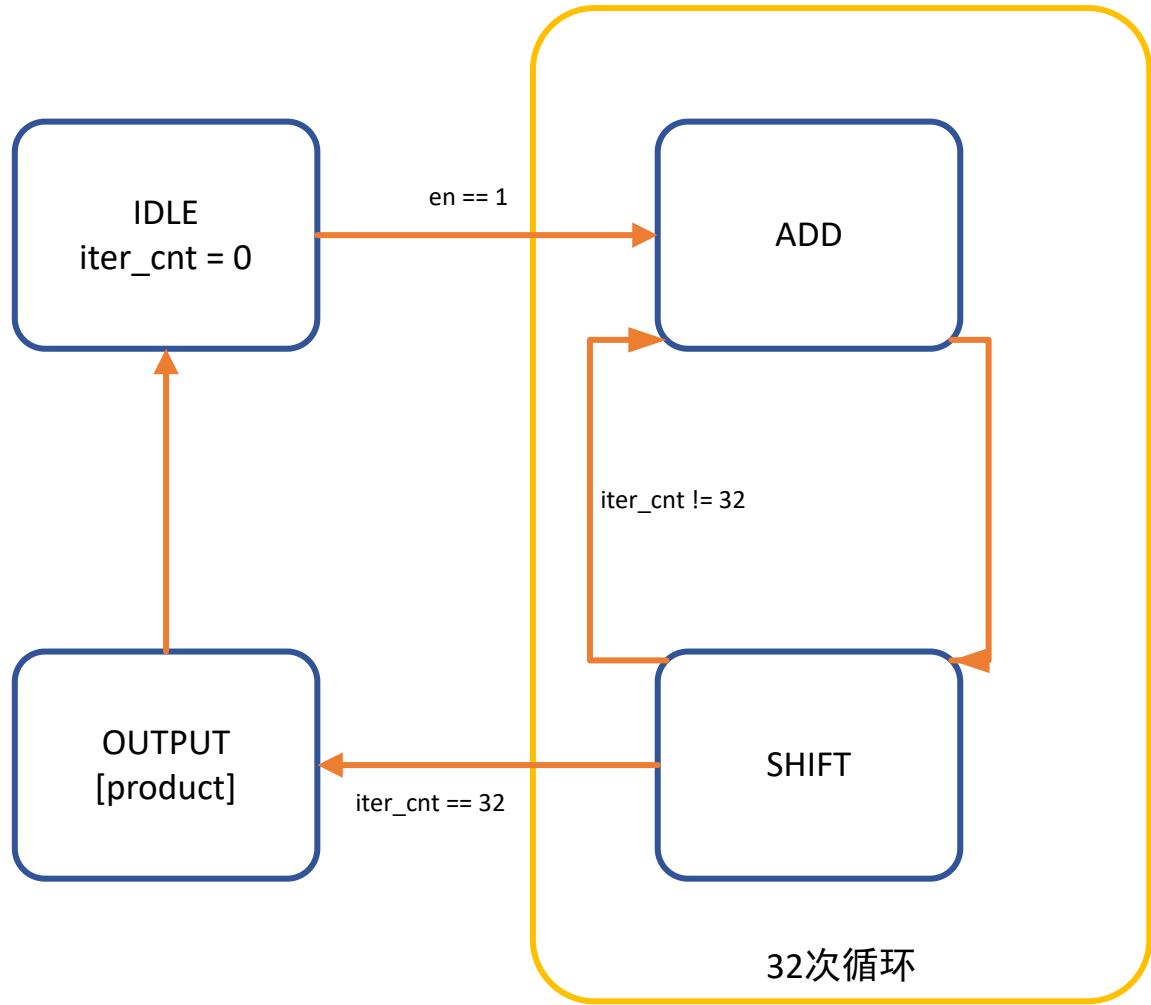
一、参考目标

SMIC 180nm工艺下，在典型情况下，逻辑综合（DC）报告数据：关键路径延迟小于3.6ns，功耗小于32.7mw，面积小于0.2mm²

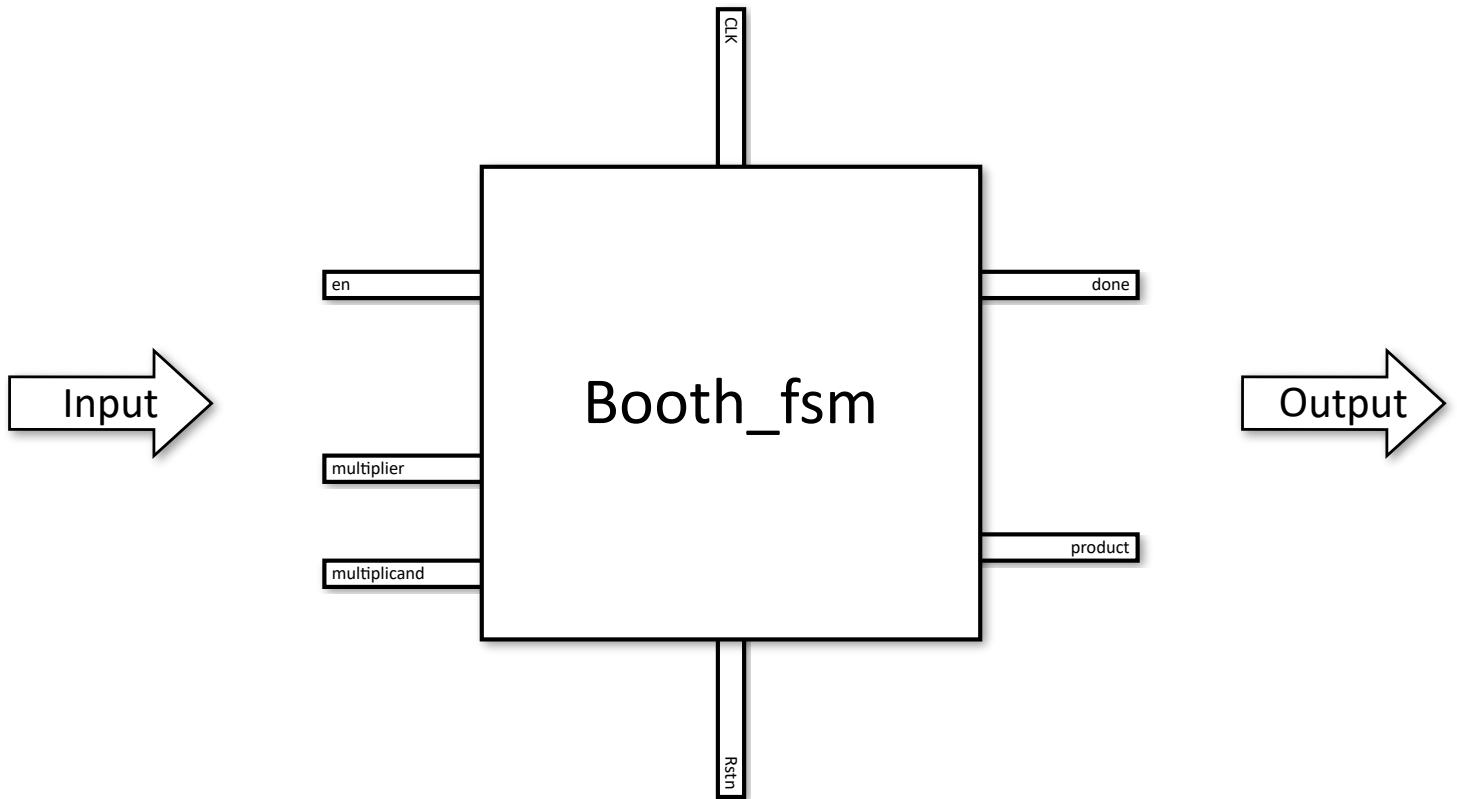
Booth乘法器状态机

通过Booth算法原理，我们发现Booth算法完成一项部分积的计算需要一次加法和移位操作（移位操作指乘2），如果要完成整个乘法计算，我们需要进行32次加法和移位。

状态机图片如下



所以我们的Booth乘法器的Top模块应当如下图



Booth乘法器部分积生成

在状态机的IDLE阶段生成部分积，分别保存在 `a_reg` 和 `s_reg`。 `a_reg` 保存正的部分积，相当于做了加法； `s_reg` 保存负的部分积，相当于做了减法。两个寄存器的宽度为66位，我们将部分积保存在高位，部分积压缩后，进行右移，最终可以发现我们进行的部分积为最低位的部分积。

```

.....
reg [2*DATAWIDTH+1:0] a_reg,s_reg,p_reg,sum_reg; // computational values.
.....
assign multiplier_neg = -{multiplier[DATAWIDTH-1],multiplier};
//取补码，相当于每一位取反，然后加1
.....
always @(posedge clk or negedge rstn) begin
.....
case (current_state)
IDLE : begin
    a_reg <= {multiplier[DATAWIDTH-1],multiplier,{(DATAWIDTH+1){1'b0}}};
    s_reg <= {multiplier_neg,{(DATAWIDTH+1){1'b0}}};
    p_reg <= {{(DATAWIDTH+1){1'b0}},multiplicand,1'b0};
    iter_cnt <= 0;
    done <= 1'b0;
end
endcase
.....
end
.....

```

Booth乘法器部分积压缩

因为我们采用状态机方式设计，所以，每个状态机循环进行一次部分积压缩，即使用加法器对其应用加法，并将结果进行移位。因为我们将部分积保存在了高位，所以每次压缩完对结果右移；这样做，相当于将部分积保存在低位，每次压缩完，将部分积寄存器左移。

```
.....
adder_66 adder1(
    .A(data_in1),
    .B(data_in2),
    .S(sum),
    .c66()
);
.....
assign data_in1 =      p_reg;
assign data_in2 =      (p_reg[1:0]==2'b01) ? a_reg:// + multiplier
                           (p_reg[1:0]==2'b10) ? s_reg:// - multiplier
                           0;// nop

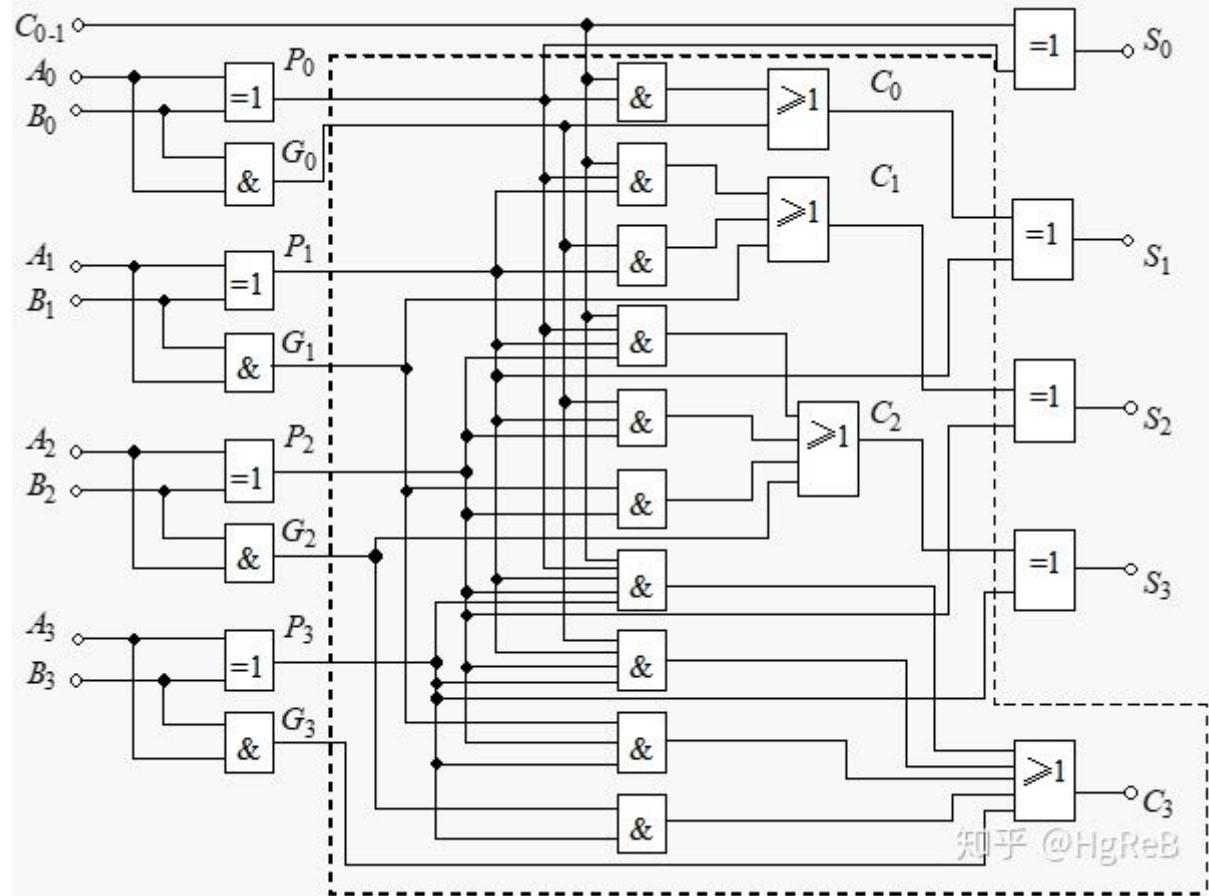
// algorithm implemenation details.
always @(posedge clk or negedge rstn) begin
.....
    case (current_state)
        IDLE : begin
.....
            end
        ADD  : begin
            sum_reg <= sum;
            iter_cnt <= iter_cnt + 1;
        end
        SHIFT : begin
            p_reg <= {sum_reg[2*DATAWIDTH+1],sum_reg[2*DATAWIDTH+1:1]}; // right shift
        end
        OUTPUT : begin
            product <= p_reg[2*DATAWIDTH:1];
            done <= 1'b1;
        end
    endcase
end
end
```

另外，我们的加法器为66位加法器，我们66位加法器的实现方式是将64位超前进位加法器和两个1bit全加器级联而成，下面是我们64bit超前进位加法器的设计。

64位超前进位加法器设计

64位超前进位加法器由4位超前进位加法器级联而成。

4位超前进位加法器电路图如下



代码如下

```

//一位全加器
module adder(X,Y,Cin,F,Cout);

    input X,Y,Cin;
    output F,Cout;

    assign F = X ^ Y ^ Cin;
    assign Cout = (X ^ Y) & Cin | X & Y;
endmodule

///四位CLA部件
module CLA(c0,c1,c2,c3,c4,p1,p2,p3,p4,g1,g2,g3,g4);

    input c0,g1,g2,g3,g4,p1,p2,p3,p4;
    output c1,c2,c3,c4;

    assign c1 = g1 ^ (p1 & c0),
           c2 = g2 ^ (p2 & g1) ^ (p2 & p1 & c0),
           c3 = g3 ^ (p3 & g2) ^ (p3 & p2 & g1) ^ (p3 & p2 & p1 & c0),
           c4 = g4 ^ (p4 & g3) ^ (p4 & p3 & g2) ^ (p4 & p3 & p2 & g1)
           ^ (p4 & p3 & p2 & p1 & c0);

endmodule

//四位超前进位加法器
module adder_4(x,y,c0,c4,F,Gm,Pm);
    input [4:1] x;
    input [4:1] y;
    input c0;
    output c4,Gm,Pm;
    output [4:1] F;

    wire p1,p2,p3,p4,g1,g2,g3,g4;
    wire c1,c2,c3;
    adder adder1(
        .X(x[1]),
        .Y(y[1]),
        .Cin(c0),
        .F(F[1]),
        .Cout()
    );
    adder adder2(
        .X(x[2]),
        .Y(y[2]),
        .Cin(c1),
        .F(F[2]),
        .Cout()
    );
    adder adder3(

```

```

.X(x[3]),
.Y(y[3]),
.Cin(c2),
.F(F[3]),
.Cout()
);

adder adder4(
    .X(x[4]),
    .Y(y[4]),
    .Cin(c3),
    .F(F[4]),
    .Cout()
);
);

CLA CLA(
    .c0(c0),
    .c1(c1),
    .c2(c2),
    .c3(c3),
    .c4(c4),
    .p1(p1),
    .p2(p2),
    .p3(p3),
    .p4(p4),
    .g1(g1),
    .g2(g2),
    .g3(g3),
    .g4(g4)
);
;

assign p1 = x[1] ^ y[1],
      p2 = x[2] ^ y[2],
      p3 = x[3] ^ y[3],
      p4 = x[4] ^ y[4];

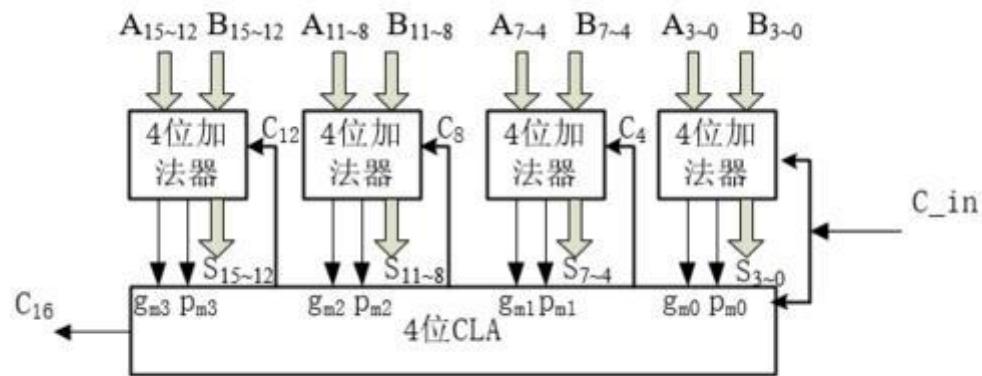
assign g1 = x[1] & y[1],
      g2 = x[2] & y[2],
      g3 = x[3] & y[3],
      g4 = x[4] & y[4];

assign Pm = p1 & p2 & p3 & p4,
      Gm = g4 ^ (p4 & g3) ^ (p4 & p3 & g2) ^ (p4 & p3 & p2 & g1);

endmodule

```

通过复用CLA电路可以级联完成16位超前进位加法器，如图



代码如下

```

//16位加法器
module adder_16(A,B,c0,c16,S,px,gx);
    input [16:1] A;
    input [16:1] B;
    input c0;
    output c16,gx,px;
    output [16:1] S;

    wire c4,c8,c12;
    wire Pm1,Gm1,Pm2,Gm2,Pm3,Gm3,Pm4,Gm4;

    adder_4 adder1(
        .x(A[4:1]),
        .y(B[4:1]),
        .c0(c0),
        .c4(),
        .F(S[4:1]),
        .Gm(Gm1),
        .Pm(Pm1)
    );
    adder_4 adder2(
        .x(A[8:5]),
        .y(B[8:5]),
        .c0(c4),
        .c4(),
        .F(S[8:5]),
        .Gm(Gm2),
        .Pm(Pm2)
    );
    adder_4 adder3(
        .x(A[12:9]),
        .y(B[12:9]),
        .c0(c8),
        .c4(),
        .F(S[12:9]),
        .Gm(Gm3),
        .Pm(Pm3)
    );
    adder_4 adder4(
        .x(A[16:13]),
        .y(B[16:13]),
        .c0(c12),
        .c4(),
        .F(S[16:13]),
        .Gm(Gm4),
        .Pm(Pm4)
    );

```

```
CLA CLA(
    .c0(c0),
    .c1(c4),
    .c2(c8),
    .c3(c12),
    .c4(c16),
    .p1(Pm1),
    .p2(Pm2),
    .p3(Pm3),
    .p4(Pm4),
    .g1(Gm1),
    .g2(Gm2),
    .g3(Gm3),
    .g4(Gm4)
);
assign px = Pm1 & Pm2 & Pm3 & Pm4,
      gx = Gm4 ^ (Pm4 & Gm3) ^ (Pm4 & Pm3 & Gm2) ^ (Pm4 & Pm3 & Pm2 & Gm1);
endmodule
```

最终16位加法器级联完成64位超前进位加法器。代码如下

```
//64位并行进位加法器顶层模块
module adder_64(A,B,c0,c64,S,px,gx);
    input [64:1] A;
    input [64:1] B;
    input c0;
    output [64:1] S;
    output c64,px,gx;

    wire px1,gx1,px2,gx2,px3,gx3,px4,gx4;
    wire c16,c32,c48,c64;

    adder_16 adder1(
        .A(A[16:1]),
        .B(B[16:1]),
        .c0(c0),
        .c16(),
        .S(S[16:1]),
        .px(px1),
        .gx(gx1)
    );
    adder_16 adder2(
        .A(A[32:17]),
        .B(B[32:17]),
        .c0(c16),
        .c16(),
        .S(S[32:17]),
        .px(px2),
        .gx(gx2)
    );
    adder_16 adder3(
        .A(A[48:33]),
        .B(B[48:33]),
        .c0(c32),
        .c16(),
        .S(S[48:33]),
        .px(px3),
        .gx(gx3)
    );
    adder_16 adder4(
        .A(A[64:49]),
        .B(B[64:49]),
        .c0(c48),
        .c16(),
        .S(S[64:49]),
        .px(px4),
        .gx(gx4)
    );

```

```

CLA CLA(
    .c0(c0),
    .c1(c16),
    .c2(c32),
    .c3(c48),
    .c4(c64),
    .p1(px1),
    .p2(px2),
    .p3(px3),
    .p4(px4),
    .g1(gx1),
    .g2(gx2),
    .g3(gx3),
    .g4(gx4)
);
assign px = px1 & px2 & px3 & px4,
      gx = gx4 ^ (px4 & gx3) ^ (px4 & px3 & gx2) ^ (px4 & px3 & px2 & gx1);
endmodule

```

Booth乘法器前仿验证

乘法器验证我们分为前仿、和后仿两部分。在前仿和后仿我们均使用自动化的验证脚本。前仿我们在windows平台，使用modelsim完成。

testbench分为两部分，分别为66位加法器的testbench，和booth乘法器的testbench。

无论是加法器的testbench，还是乘法器的testbench，我们均选择随机生成 $16'bFFFF$ ，即 65535 组操作数，并作为激励输入，并自动验证结果是否正确

加法器testbench如下

```

`timescale 1ns/1ps

// Basic exhaustive self checking test bench.
`define TEST_WIDTH 66
module adder_tb;

reg clk;
reg rstn;
integer src1;
integer src2;
reg [`TEST_WIDTH-1:0] A;
reg [`TEST_WIDTH-1:0] B;

//输入 : 要定义有符号和符号, 输出: 无要求
wire signed [`TEST_WIDTH-1:0] data_out;
wire signed [`TEST_WIDTH-1:0] data_in1;
wire signed [`TEST_WIDTH-1:0] data_in2;

reg signed [`TEST_WIDTH-1:0] adder_ref;
reg signed [`TEST_WIDTH-1:0] data_out_reg;
assign data_in1 = A[`TEST_WIDTH-1:0];
assign data_in2 = B[`TEST_WIDTH-1:0];

adder_66 adder
(
    .A(data_in1),
    .B(data_in2),
    .S(data_out),
    .c66()
);

always #1 clk = ~clk;

integer num_good;
integer i,j;
initial begin
    clk = 1;
    rstn = 1;
    #2 rstn = 0; #2 rstn = 1;
    num_good = 0;
    for(i = 0; i < 32'hFFFF;i=i+1) begin
        A={$random}%(66'h3_FFFF_FFFF_FFFF_FFFF);//
        B={$random}%(66'h3_FFFF_FFFF_FFFF_FFFF);//
        if (data_out_reg !== adder_ref)

```

```

        $display("A = %d B = %d data_out =%d adder_ref =%d",data_in1,data_in2,data_out_reg,adder_ref);
    else
        num_good = num_good + 1;
    @(posedge clk);
end

$display("sim done. num good = %d",num_good);
$finish;

end

always @(posedge clk or negedge rstn) begin
    adder_ref      <=  data_in1+data_in2;
    data_out_reg   <=  data_out;
end

initial begin
    // $fsdbDumpvars();
    // $fsdbDumpMDA();
    // $dumpvars();
    // if(finish === 1)
    //     $finish;
end
endmodule

```

加法器仿真结果如下

```

# Compile of lookahead_adder.v was successful.
# Compile of boothmul_tb.v was successful.
# Compile of lookahead_adder_tb.v was successful.
# Compile of booth_fsm.v was successful.
# 4 compiles, 0 failed with no errors.
VSIM 5> vsim -gui work.adder_tb
# End time: 12:58:51 on Jan 03, 2022, Elapsed time: 0:05:48
# Errors: 0, Warnings: 8
# vsim -gui work.adder_tb
# Start time: 12:58:51 on Jan 03, 2022
# Loading work.adder_tb
# Loading work.adder_66
# Loading work.adder_64
# Loading work.adder_16
# Loading work.adder_4
# Loading work.adder
# Loading work.CLA
VSIM 6> run -all
# sim done. num good =      65535
# ** Note: $finish      : C:/Users/admin/Desktop/booth/tb/lookahead_adder_tb.v(59)
#      Time: 131072 ns  Iteration: 1  Instance: /adder_tb
# 1
# Break in Module adder_tb at C:/Users/admin/Desktop/booth/tb/lookahead_adder_tb.v line 59

VSIM 7>

```

乘法器testbench如下

```

`timescale 1ns/1ps

// Basic exhaustive self checking test bench.
`define TEST_WIDTH 32
module booth_fsm_tb;

reg clk;
reg rstn;
reg en;
integer multiplier1;
integer multiplicand1;
reg [`TEST_WIDTH-1:0] multiplier;
reg [`TEST_WIDTH-1:0] multiplicand;
wire done;
reg finish;

//输入 : 要定义有符号和符号, 输出: 无要求
wire signed [2*`TEST_WIDTH-1:0] product;
wire signed [`TEST_WIDTH-1:0] m1_in;
wire signed [`TEST_WIDTH-1:0] m2_in;

reg signed [2*`TEST_WIDTH-1:0] product_ref;
reg [2*`TEST_WIDTH-1:0] product_ref_u;
assign m1_in = multiplier[`TEST_WIDTH-1:0];
assign m2_in = multiplicand[`TEST_WIDTH-1:0];

booth_fsm #( .DATAWIDTH(`TEST_WIDTH) ) booth
(
    .clk(clk),
    .rstn(rstn),
    .en(en),
    .multiplier(multiplier),
    .multiplicand(multiplicand),
    .done(done),
    .product(product)
);
always #1 clk = ~clk;

integer num_good;
integer i,j;
initial begin
    clk = 1;
    en = 0;
    rstn = 1;
    finish = 0;
    #2 rstn = 0; #2 rstn = 1;

    num_good = 0;
    multiplier=0;

```

```

multiplicand=0;
#9;

for(i = 0; i < 32'hFFFF;i=i+1) begin

    en = 1;

    wait (done == 0);
    wait (done == 1);
    multiplier={$random}%(32'hFFFF_FFFF);
    multiplicand={$random}%(32'hFFFF_FFFF);
    if (product_ref !== product)
        $display("multiplier = %d multiplicand = %d proudct =%d product_ref =%d",m1_in,m2_in,product_ref,product);
    else
        num_good = num_good + 1;
    @(posedge clk);
end
finish = 1;
$display("sim done. num good = %d",num_good);
$finish;

end

always @(posedge clk or negedge rstn) begin
    product_ref<=m1_in*m2_in;
    product_ref_u<=m1_in*m2_in;
end

initial begin
    // $fsdbDumpvars();
    // $fsdbDumpMDA();
    // $dumpvars();
    // if(finish === 1)
    //     $finish;
end
endmodule

```

乘法器仿真结果如下：

```
transcript
# Compile of boothmul_tb.v was successful.
# Compile of lookahead_adder_tb.v was successful.
# Compile of booth_fsm.v was successful.
# 4 compiles, 0 failed with no errors.
# Compile of lookahead_adder.v was successful.
# Compile of boothmul_tb.v was successful.
# Compile of lookahead_adder_tb.v was successful.
# Compile of booth_fsm.v was successful.
# 4 compiles, 0 failed with no errors.
ModelSim> vsim -gui work.booth_fsm_tb
# vsim -gui work.booth_fsm_tb
# Start time: 12:50:24 on Jan 03, 2022
# Loading work.booth_fsm_tb
# Loading work.booth_fsm
# Loading work.adder_66
# Loading work.adder_64
# Loading work.adder_16
# Loading work.adder_4
# Loading work.adder
# Loading work.CLA
VSIM 2> run -all
# sim done. num good =      65535
# ** Note: $finish    : C:/Users/admin/Desktop/booth/tb/boothmul_tb.v(71)
#   Time: 8650634 ns  Iteration: 1  Instance: /booth_fsm_tb
# 1
# Break in Module booth_fsm_tb at C:/Users/admin/Desktop/booth/tb/boothmul_tb.v line 71

```

Booth乘法器后端设计

本章节采用 SMIC180 纳米工艺库，对Booth乘法器进行版图综合设计，。整个流程大致如下：首先使用 Design Compiler 对所编写的 Verilog 文件进行逻辑综合，并添加相应约束。之后对输出的网表文件进行时序验证和形式验证。接着用 ICC 进行版图综合，并进行 LVS/DRC 检查。最后进行后仿，验证设计是否符合要求。

DC综合

指定工艺库

在进行DC综合时，需要先进行工艺库的指定，这里指定 smic180 typical 工艺库，工作电压为1.0V。所用的tcl命令如下：

```

#####
# Library Setup Variables
#####

# For the following variables, use a blank space to separate multiple entries.
# Example: set TARGET_LIBRARY_FILES "lib1.db lib2.db lib3.db"
# Additional search path to be added to the default search path
set ADDITIONAL_SEARCH_PATH      "/home/IC-smic180/digital/sc/synopsys \
                                    /home/IC-smic180/digital/io/synopsys " ;

# Target technology logical libraries
set TARGET_LIBRARY_FILES         "typical.db ";

# Extra link logical libraries not included in TARGET_LIBRARY_FILES
set ADDITIONAL_LINK_LIB_FILES   "SP018N_V1p0_typ.db ";

# List of max min library pairs "max1 min1 max2 min2 max3 min3"...
set MIN_LIBRARY_FILES           "slow.db                  fast.db \
                                    SP018N_V1p0_min.db    SP018N_V1p0_max.db"          \ ;
                                ;

puts "RM-Info: Completed script [info script]\n"

```

制定约束条件

在进行逻辑综合时，选择约束条件如下：

预计达到的时钟频率为 200 MHz。采用 create_clock 对寄存器间的时钟频率进行约束。
对 clock_latency, clock_uncertainty, clock_transition 按照经验值进行设定。
将输入延时和输出延时按照经验值设定。
其他约束

约束代码如下

```

*****  

#set design optimization constraints  

*****  

create_clock -name clock_main -period 5 [get_ports "clk"]  
  

set_clock_latency      0.05 [get_clocks "clock_main"]  

set_clock_uncertainty  1.5 -setup [get_clocks "clock_main"]  

set_clock_uncertainty  0.05 -hold [get_clocks "clock_main"]  

set_clock_transition   0.05 [get_clocks "clock_main"]  
  

#clock_main  

set_input_delay -max 1           -clock clock_main [remove_from_collection [all_inputs] [get_port  

set_input_delay -min 0.1        -clock clock_main [remove_from_collection [all_inputs] [get_port  

set_output_delay -max 1          -clock clock_main [all_outputs]  

set_output_delay -min 0.1        -clock clock_main [all_outputs]  
  

*****  

#set design rule constraints  

*****  
  

set_ideal_network [get_ports "rstn"]  

set_ideal_network [get_ports "clk"]  
  

#set_driving_cell -lib_cell "PSBI8N" -from_pin "A" -pin "P" [all_inputs]  
  

group_path -name REGOUT       -to    [all_outputs]  

group_path -name REGIN        -from  [remove_from_collection [all_inputs] [get_ports "clk"]]  

group_path -name FEEDTHROUGH -from  [remove_from_collection [all_inputs] [get_ports "clk"]] -to [
```

综合以及结果写出

逻辑综合部分完成会写出网表文件，sdc约束文件，时序报告，面积报告，功耗报告以及violation报告。

各报告详细信息如下

Violation

无violation

报告如下：

```
*****
Report : constraint
    -all_violators
    -verbose
Design : booth_fsm
Version: K-2015.06
Date   : Mon Jan  3 14:29:51 2022
*****
```

1

Power

总功耗为 4.1157mW , 以动态功耗为主, 大多数功耗为寄存器的Internal Power。符合功耗要求不高于 32.7mW 。

报告如下。

Report : power

-analysis_effort low
-verbose

Design : booth_fsm

Version: K-2015.06

Date : Mon Jan 3 14:29:51 2022

Library(s) Used:

typical (File: /home/IC/smic180/digital/sc/synopsys/typical.db)

Operating Conditions: typical Library: typical

Wire Load Model Mode: top

Global Operating Voltage = 1.8

Power-specific unit information :

Voltage Units = 1V

Capacitance Units = 1.000000pf

Time Units = 1ns

Dynamic Power Units = 1mW (derived from V,C,T units)

Leakage Power Units = 1pW

Cell Internal Power = 3.7915 mW (92%)

Net Switching Power = 324.1418 uW (8%)

Total Dynamic Power = 4.1156 mW (100%)

Cell Leakage Power = 101.9960 nW

Power Group	Internal Power	Switching Power	Leakage Power	Total Power	(%)	Att

io_pad	0.0000	0.0000	0.0000	0.0000	(0.00%)	
memory	0.0000	0.0000	0.0000	0.0000	(0.00%)	
black_box	0.0000	0.0000	0.0000	0.0000	(0.00%)	
clock_network	0.0000	0.0000	0.0000	0.0000	(0.00%)	
register	3.5466	6.8441e-02	5.8482e+04	3.6151	(87.84%)	
sequential	0.0000	0.0000	0.0000	0.0000	(0.00%)	
combinational	0.2449	0.2557	4.3514e+04	0.5006	(12.16%)	

Total	3.7915 mW	0.3241 mW	1.0200e+05 pW	4.1157 mW		

Area

因为使用状态机电路，面积远远低于要求，面积为 0.0186mm²，符合要求，不高于 0.2mm²，报告如下：

```
*****
Report : area
Design : booth_fsm
Version: K-2015.06
Date   : Mon Jan  3 14:29:51 2022
*****
```

Library(s) Used:

typical (File: /home/IC-smic180/digital/sc/synopsys/typical.db)

Number of ports: 297
Number of nets: 1429
Number of cells: 1200
Number of combinational cells: 902
Number of sequential cells: 296
Number of macros/black boxes: 0
Number of buf/inv: 120
Number of references: 43

Combinational area: 13448.635328
Buf/Inv area: 804.988811
Noncombinational area: 18644.471741
Macro/Black Box area: 0.000000
Net Interconnect area: undefined (No wire load specified)

Total cell area: 32093.107069
Total area: undefined

Hierarchical area distribution

Hierarchical cell	Global cell area			Local cell area		
	Absolute Total	Percent Total	Combi-national	Noncombi-national	Black-boxes	Design
booth_fsm	32093.1071	100.0	10491.4657	18644.4717	0.0000	booth_fsm
add_x_4	1410.3936	4.4	1410.3936	0.0000	0.0000	booth_fsm
sub_x_2	1546.7760	4.8	1546.7760	0.0000	0.0000	booth_fsm
Total			13448.6353	18644.4717	0.0000	

Timing

关键路径延时为 3.36ns , 低于要求的 3.6ns , 经过检查发现, 关键路径为66bits加法器, 符合预期。报告如下:

Report : timing
-path full
-delay max
-nets
-max_paths 1
-transition_time

Design : booth_fsm

Version: K-2015.06

Date : Mon Jan 3 14:29:51 2022

Operating Conditions: typical Library: typical
Wire Load Model Mode: top

Startpoint: p_reg_reg_1_
 (rising edge-triggered flip-flop clocked by clock_main)
Endpoint: sum_reg_reg_65_
 (rising edge-triggered flip-flop clocked by clock_main)
Path Group: clock_main
Path Type: max

Attributes:

d - dont_touch
u - dont_use
mo - map_only
so - size_only
i - ideal_net or ideal_network
inf - infeasible path

Point	Fanout	Trans	Incr	Path	Attributes
<hr/>					
clock clock_main (rise edge)			0.00	0.00	
clock network delay (ideal)			0.05	0.05	
p_reg_reg_1_/CK (DFFSX1)		0.05	0.00	0.05 r	i
p_reg_reg_1_/Q (DFFSX1)		0.18	0.42	0.47 r	
n1125 (net)	2		0.00	0.47 r	
U1272/Y (NOR2X4)		0.10	0.07	0.55 f	
n791 (net)	16		0.00	0.55 f	
U1280/Y (NAND2XL)		0.07	0.06	0.61 r	
n736 (net)	1		0.00	0.61 r	
U1281/Y (NAND2XL)		0.07	0.06	0.67 f	
n1011 (net)	2		0.00	0.67 f	
U1282/Y (NAND2XL)		0.16	0.11	0.78 r	
n1010 (net)	4		0.00	0.78 r	
U1283/Y (NAND2XL)		0.06	0.05	0.82 f	
n740 (net)	1		0.00	0.82 f	
U805/Y (AOI2BB2XL)		0.28	0.18	1.01 r	

n939 (net)	3	0.00	1.01	r
U1287/Y (INVXL)		0.07	0.04	1.05 f
n742 (net)	1	0.00	1.05	f
U1290/Y (OAI21XL)		0.24	0.17	1.22 r
n935 (net)	3	0.00	1.22	r
U1294/Y (NAND2XL)		0.07	0.05	1.27 f
n752 (net)	1	0.00	1.27	f
U806/Y (AOI21XL)		0.22	0.15	1.42 r
n1053 (net)	2	0.00	1.42	r
U810/Y (AOI2BB2XL)		0.24	0.17	1.59 r
n958 (net)	2	0.00	1.59	r
U718/Y (NAND2X1)		0.09	0.06	1.65 f
n951 (net)	2	0.00	1.65	f
U717/Y (AOI22X1)		0.24	0.15	1.80 r
n962 (net)	3	0.00	1.80	r
U753/Y (INVXL)		0.07	0.04	1.85 f
n782 (net)	1	0.00	1.85	f
U1310/Y (OAI21X1)		0.21	0.15	2.00 r
n930 (net)	2	0.00	2.00	r
U727/Y (NOR2X2)		0.07	0.05	2.05 f
n929 (net)	3	0.00	2.05	f
U820/Y (AOI21X1)		0.15	0.11	2.16 r
n917 (net)	2	0.00	2.16	r
U1257/Y (NOR2X1)		0.09	0.06	2.22 f
n916 (net)	3	0.00	2.22	f
U748/Y (NOR2X1)		0.13	0.10	2.32 r
n908 (net)	2	0.00	2.32	r
U1316/Y (OR4X2)		0.05	0.10	2.42 r
n808 (net)	1	0.00	2.42	r
U824/Y (AOI21X1)		0.11	0.07	2.49 f
n965 (net)	3	0.00	2.49	f
U843/Y (AOI31XL)		0.15	0.11	2.60 r
n842 (net)	1	0.00	2.60	r
U842/Y (OAI22XL)		0.13	0.10	2.69 f
n896 (net)	3	0.00	2.69	f
U902/Y (NOR4BXL)		0.10	0.14	2.84 f
n849 (net)	1	0.00	2.84	f
U900/Y (AOI211XL)		0.31	0.20	3.03 r
n865 (net)	3	0.00	3.03	r
U851/Y (NOR2XL)		0.10	0.07	3.11 f
n870 (net)	2	0.00	3.11	f
U864/Y (NOR2XL)		0.15	0.11	3.22 r
n878 (net)	2	0.00	3.22	r
U879/Y (INVXL)		0.05	0.04	3.25 f
n881 (net)	1	0.00	3.25	f
U912/Y (NAND4XL)		0.18	0.07	3.32 r
n885 (net)	1	0.00	3.32	r
U909/Y (NAND4XL)		0.09	0.07	3.39 f
n647 (net)	1	0.00	3.39	f
sum_reg_reg_65_/D (DFFSX1)		0.09	0.00	3.39 f
data arrival time				3.39

```

clock clock_main (rise edge)           5.00      5.00
clock network delay (ideal)          0.05      5.05
clock uncertainty                   -1.50     3.55
sum_reg_reg_65_/CK (DFFSX1)         0.00      3.55 r
library setup time                  -0.16     3.39
data required time                  3.39
-----
data required time                  3.39
data arrival time                  -3.39
-----
slack (MET)                         0.00

```

1

总结

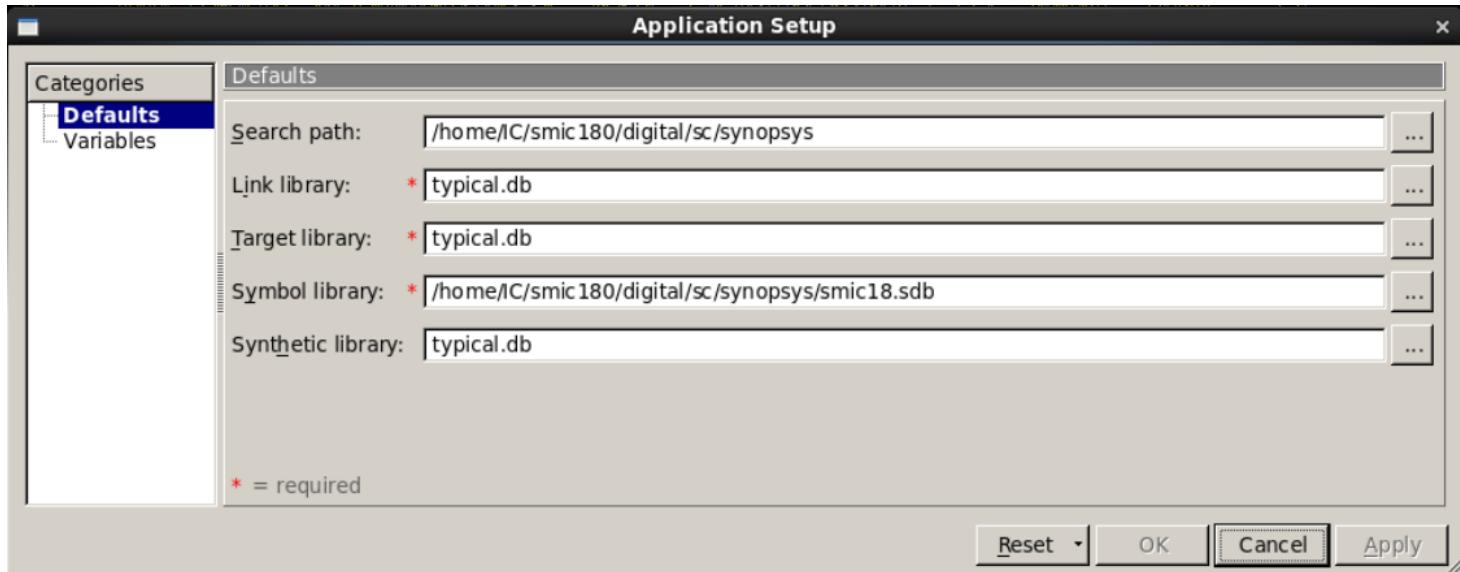
通过分析报告发现，因为采用了状态机的方式完成本设计，面积功耗远低于预期，为预期的 10% 左右，另外关键路径为加法器结果的最高（需要进位传递）符合预期，另外，因为我们设定了很高的时钟不确定度，导致裕度为0，但其实关键路径延时达到了Spec设定的要求。

ICC版图设计

进行了DC的逻辑综合之后，我们要开始版图设计，版图设计工具为ICC。在使用ICC时，需要将DC生成的网表文件和SDC约束文件进行导入。

工艺库选定

在进行 DC 综合时，已经对工艺库进行了选定，在进行 ICC 设计时同样选择SMIC180的typical工艺库。

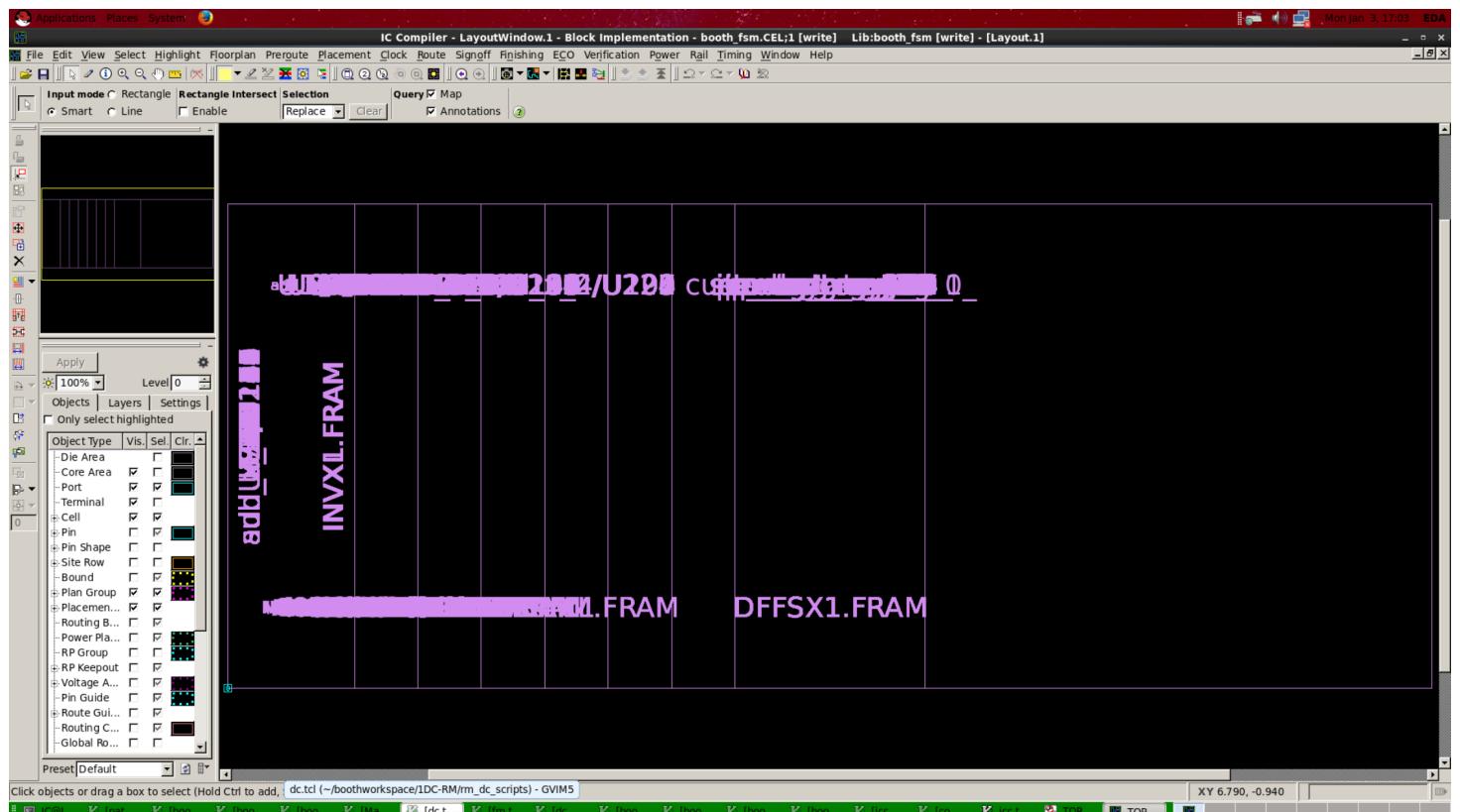


ICC初始化

初始化部分代码如下所示，读入设计文件和sdc文件，并初始化

```
create_mw_lib -technology /home/IC-smic180/digital/sc/apollo/tf-smic18_6lm.tf -mw_reference_library {/home/IC-smic180/digital/sc/apollo/sm1c18} -bus_naming_style {[%d]} -open boc  
import_designs -format verilog -top booth_fsm -cel booth_fsm{/home/IC/boothworkspace/1DC-RM/icc  
uniquify_fp_mw_cel  
current_design booth_fsm  
link  
set_tlu_plus_files -max_tluplus /home/IC-smic180/itf_tluplus/sm1clog018_6lm_cell_max.tluplus \  
-min_tluplus /home/IC-smic180/itf_tluplus/sm1clog018_6lm_cell_min.tluplus \  
-tech2itf_map /home/IC-smic180/itf_tluplus/6lm.map  
read_sdc -version Latest "/home/IC/boothworkspace/1DC-RM/icc3/booth_fsm.mapped.sdc"
```

初始化完成如下图所示：



电源网络

接下来开始电源网络的搭建

电源网络搭建的命令如下

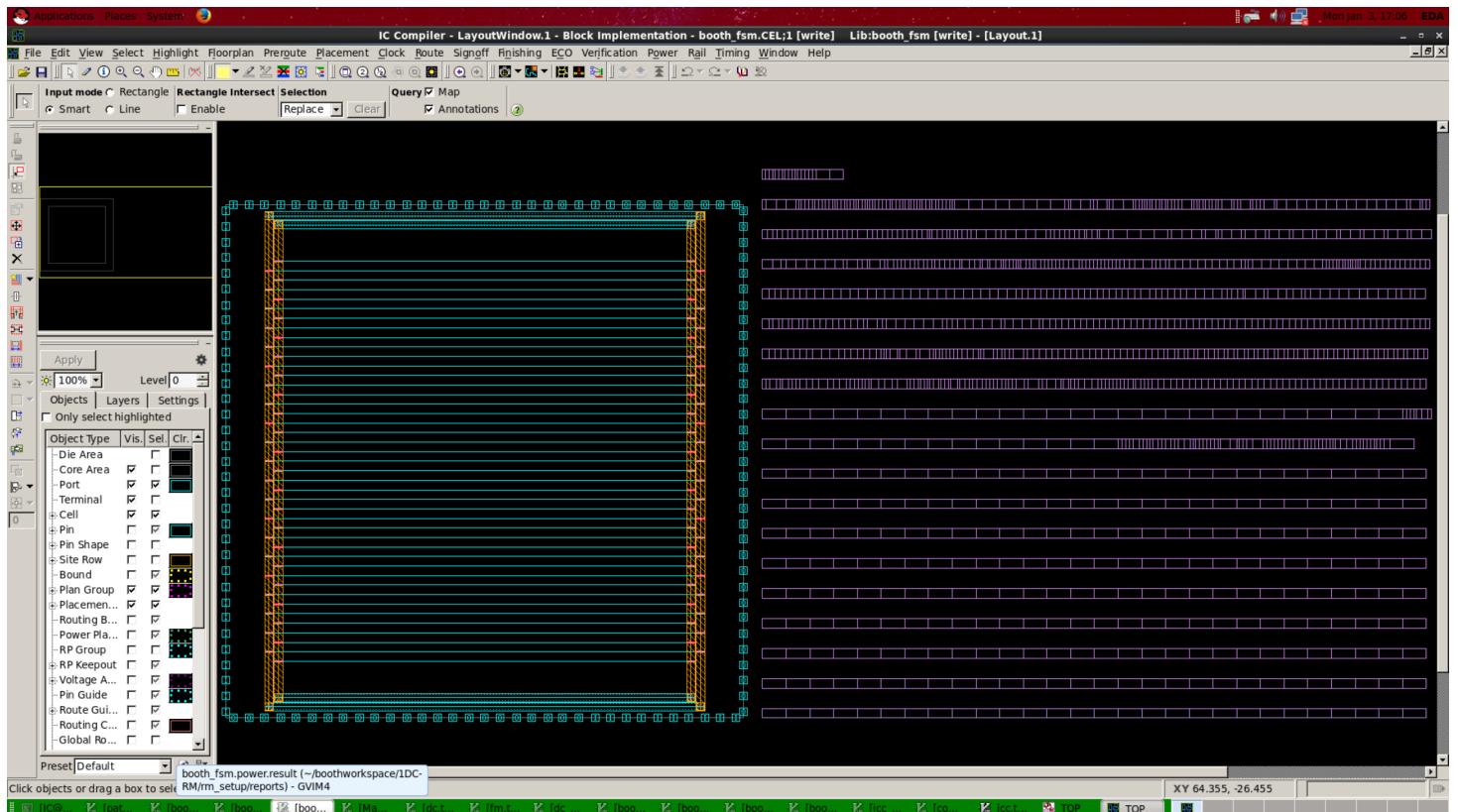
```

derive_pg_connection -power_net {VDD} -ground_net {VSS} -power_pin {VDD} -ground_pin {VSS}
derive_pg_connection -power_net {VDD} -ground_net {VSS} -tie
create_rectangular_rings -nets {VDD VSS} -left_segment_layer METAL6 \
                           -left_segment_width 4.5 \
                           -right_segment_layer METAL6 \
                           -right_segment_width 4.5 \
                           -bottom_offset 17 \
                           -bottom_segment_width 4.5 -top_offset 17 \
                           -top_segment_width 4.5 \
                           -offsets absolute
preroute_standard_cells -nets {VDD VSS} -connect horizontal \
                        -fill_empty_rows \
                        -port_filter_mode off \
                        -cell_master_filter_mode off \
                        -cell_instance_filter_mode off \
                        -voltage_area_filter_mode off \
                        -route_type {P/G Std. Cell Pin Conn}

verify_pg_nets -std_cell_pin_connection ignore -macro_pin_connection all -pad_pin_connection al

```

搭建好电源网络如下图所示



时钟树

接下来构建时钟树，并对原件进行布局

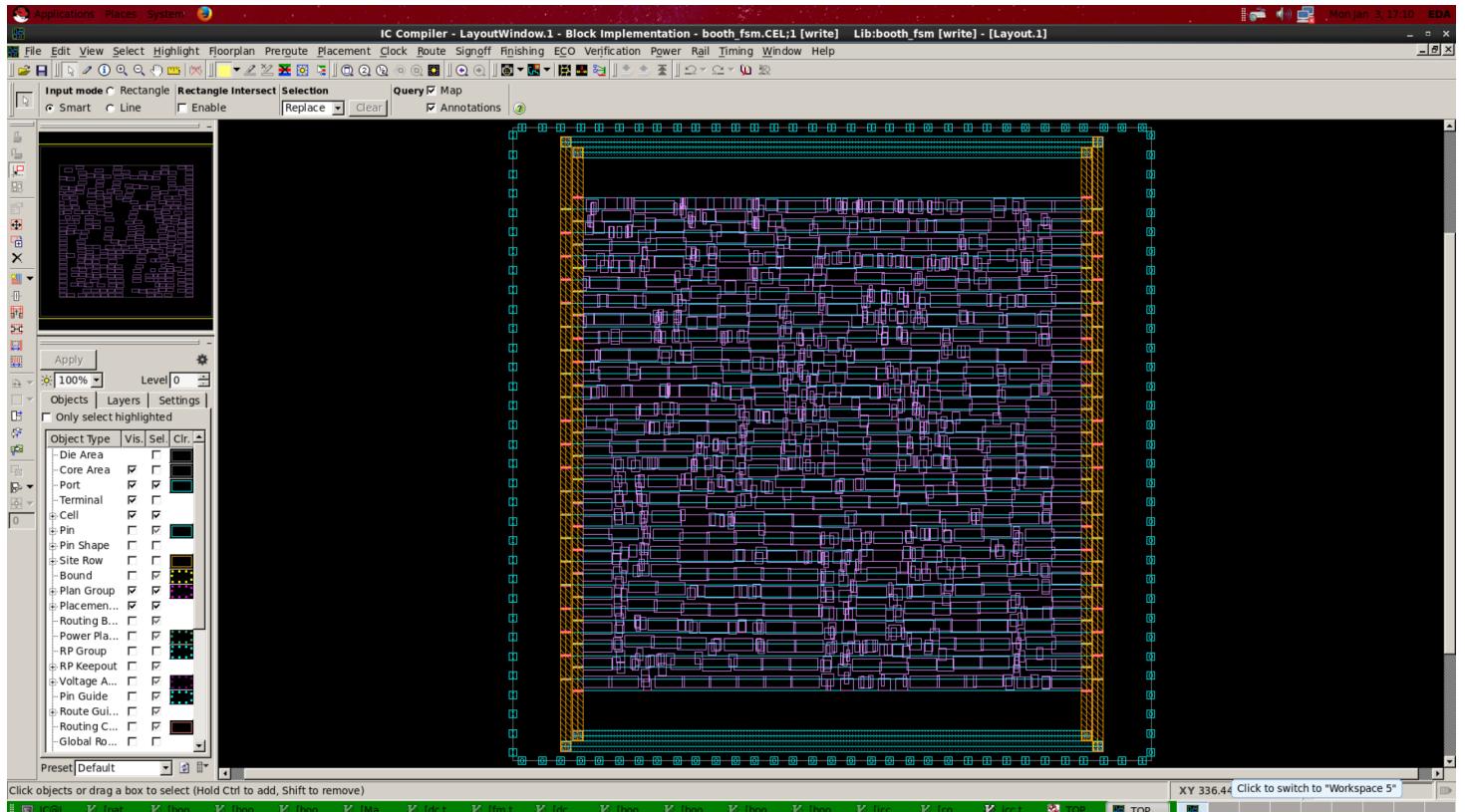
时钟树构建命令如下

```

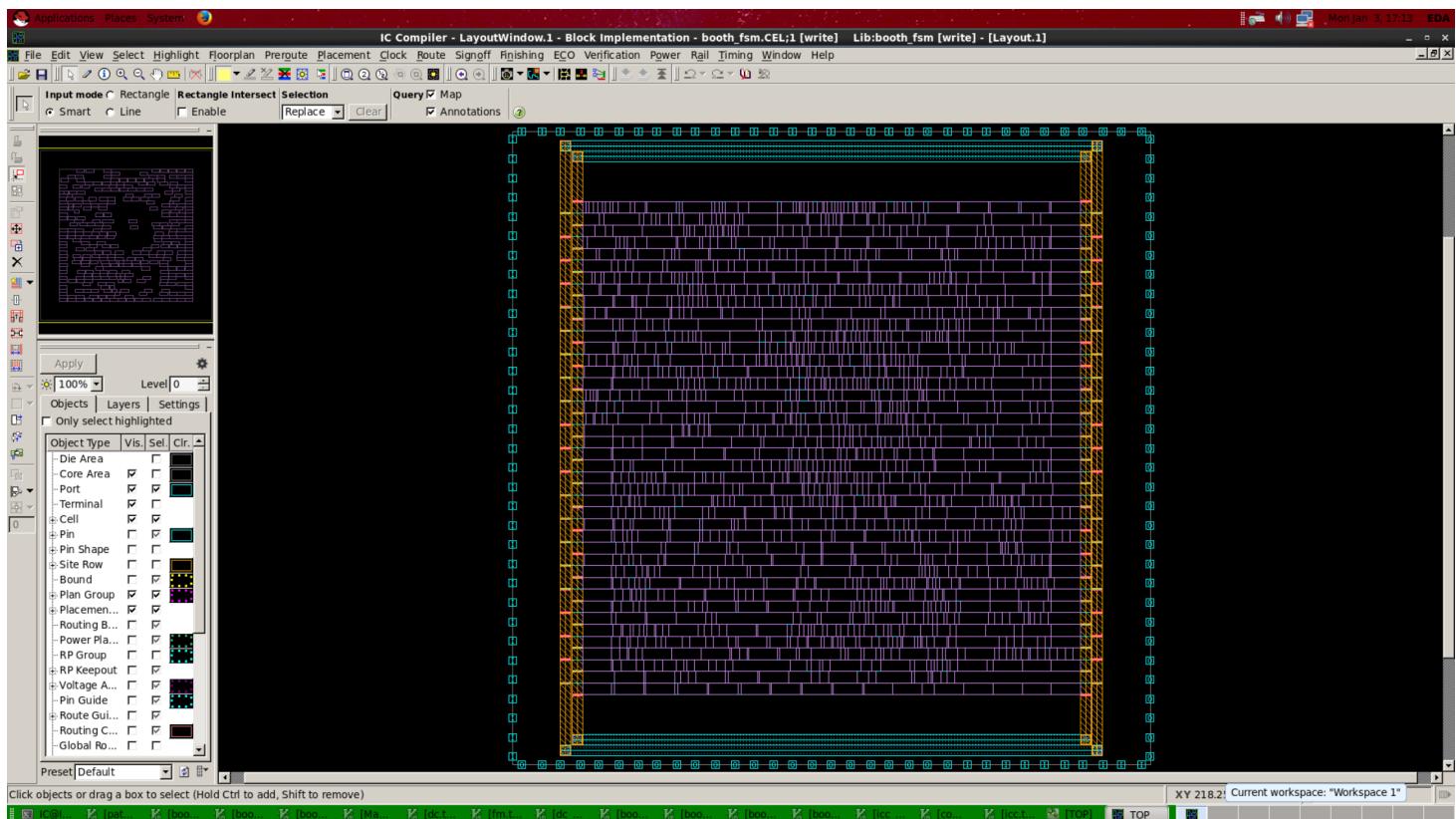
set_ideal_network [all_fanout -flat -clock_tree]
create_placement -congestion -timing_driven
place_opt -effort high -congestion
remove_routing_rules -all
define_routing_rule -default_reference_rule -multiplier_width 2 -multiplier_spacing 2 {clk_c
set_clock_tree_options -clock_tree clock_main -routing_rule clk_dsdw \
                        -use_default_routing_for_sinks 1 -layer_list_for_sinks {METAL3 METAL4 ME
                        -use_leaf_routing_rule_for_sinks 1 -max_transition 0.500 -leaf_max_trans
                        -use_leaf_max_transition_on_exceptions FALSE -use_leaf_max_transition_or
                        -max_capacitance 0.600 -max_fanout 2000 -max_rc_scale_factor 0.000 \
                        -target_early_delay 0.000 -target_skew 0.000 -buffer_relocation TRUE \
                        -gate_sizing FALSE -buffer_sizing TRUE -gate_relocation TRUE \
                        -layer_list {METAL3 METAL4 METAL5 METAL6}
set_clock_tree_references -references {typical/CLKBUFX1 typical/CLKINVXL typical/CLKINVX8 typi
                                         typical/CLKINVX3 typical/CLKINVX20 typical/CLKINVX2 typ
                                         typical/CLKBUFX12 typical/CLKINVX12 typical/CLKBUFLX ty
                                         typical/CLKBUFX4 typical/CLKBUFX3 typical/CLKBUFX20 typ
                                         typical/CLKBUFX16 typical/CLKINVX1 } -clock_trees {cloc
clock_opt -only_cts -no_clock_route

```

构建完理想时钟树，并按照时钟树放置元件，如下图所示



接下来进行布局优化，将元件放置于合适的位置，如下图所示



时钟树报告如下

```
*****
Report : clock tree
Design : booth_fsm
Version: L-2016.03-SP1
Date   : Mon Jan  3 17:12:07 2022
*****
```

Information: Float pin scale factor for the 'max' operating condition of scenario 'default' is s

```
===== Clock Tree Summary =====
Clock          Sinks      CTBuffers ClkCells  Skew      LongestPath TotalDRC  BufferArea
-----
clock_main      295        6         6    0.0263    0.3057      0     123.0768
1
```

布线优化

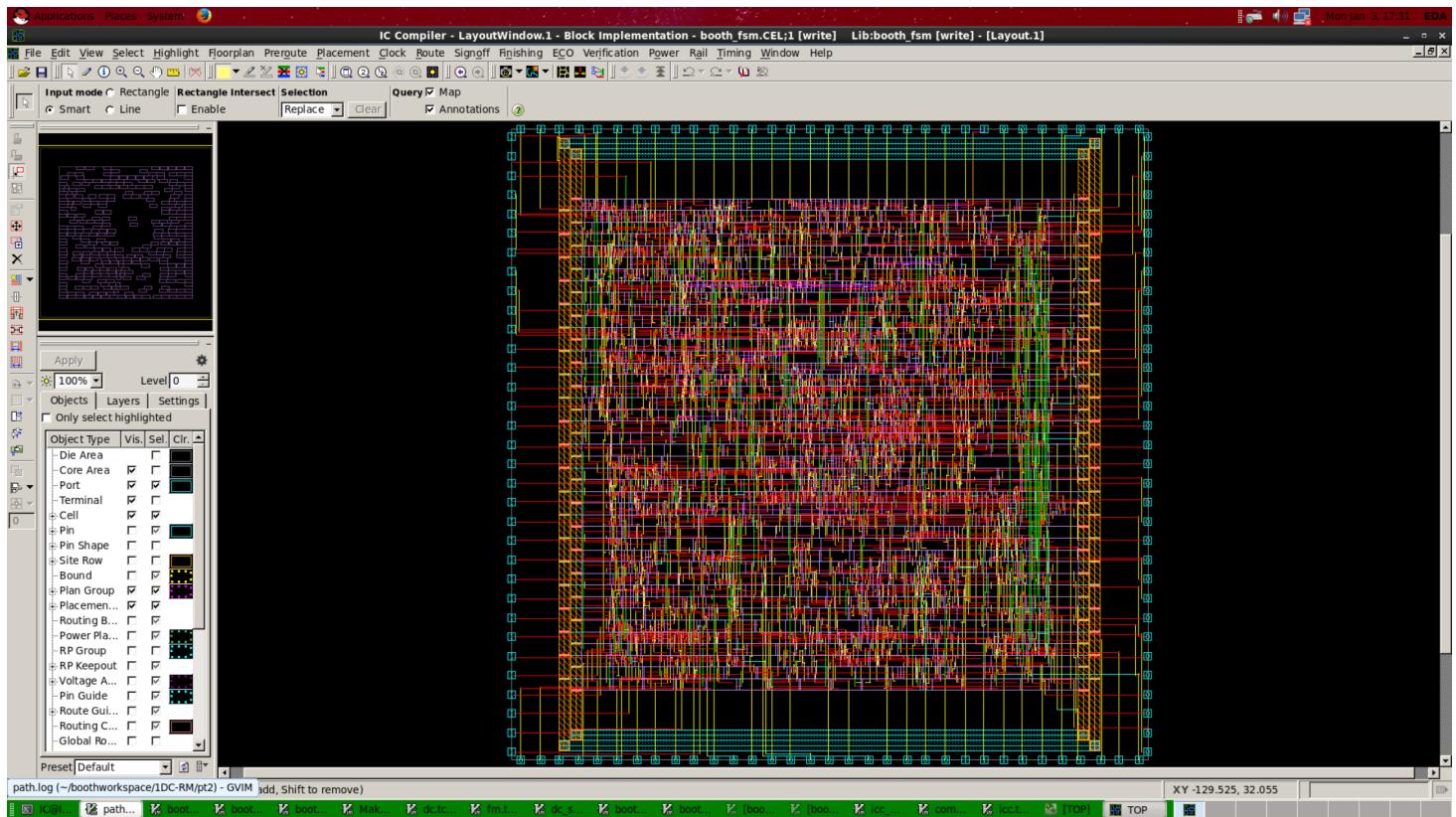
布局完成之后，要进行布线，并进行 hold_up time 修复，命令如下：

```

remove_clock_latency clock_main
remove_ideal_network [all_fanout -flat -clock_tree]
set_propagated_clock clock_main
set_clock_uncertainty -setup 0.01 clock_main
set_clock_uncertainty -hold 0.005 clock_main
set_fix_hold [all_clocks]
clock_opt -fix_hold_all_clocks -no_clock_route
derive_pg_connection -power_net {VDD} -power_pin {VDD} -ground_net {VSS} -ground_pin {VSS}
derive_pg_connection -power_net {VDD} -ground_net {VSS} -tie
route_zrt_auto
report_constraint -all_violators

```

结果如下所示



优化后的violation检查

violation说明没有冲突。

```
icc_shell> report_constraints -all_violators
```

```
*****
Report : constraint
    -all_violators
Design : booth_fsm
Version: L-2016.03-SP1
Date   : Mon Jan  3 17:30:20 2022
*****
```

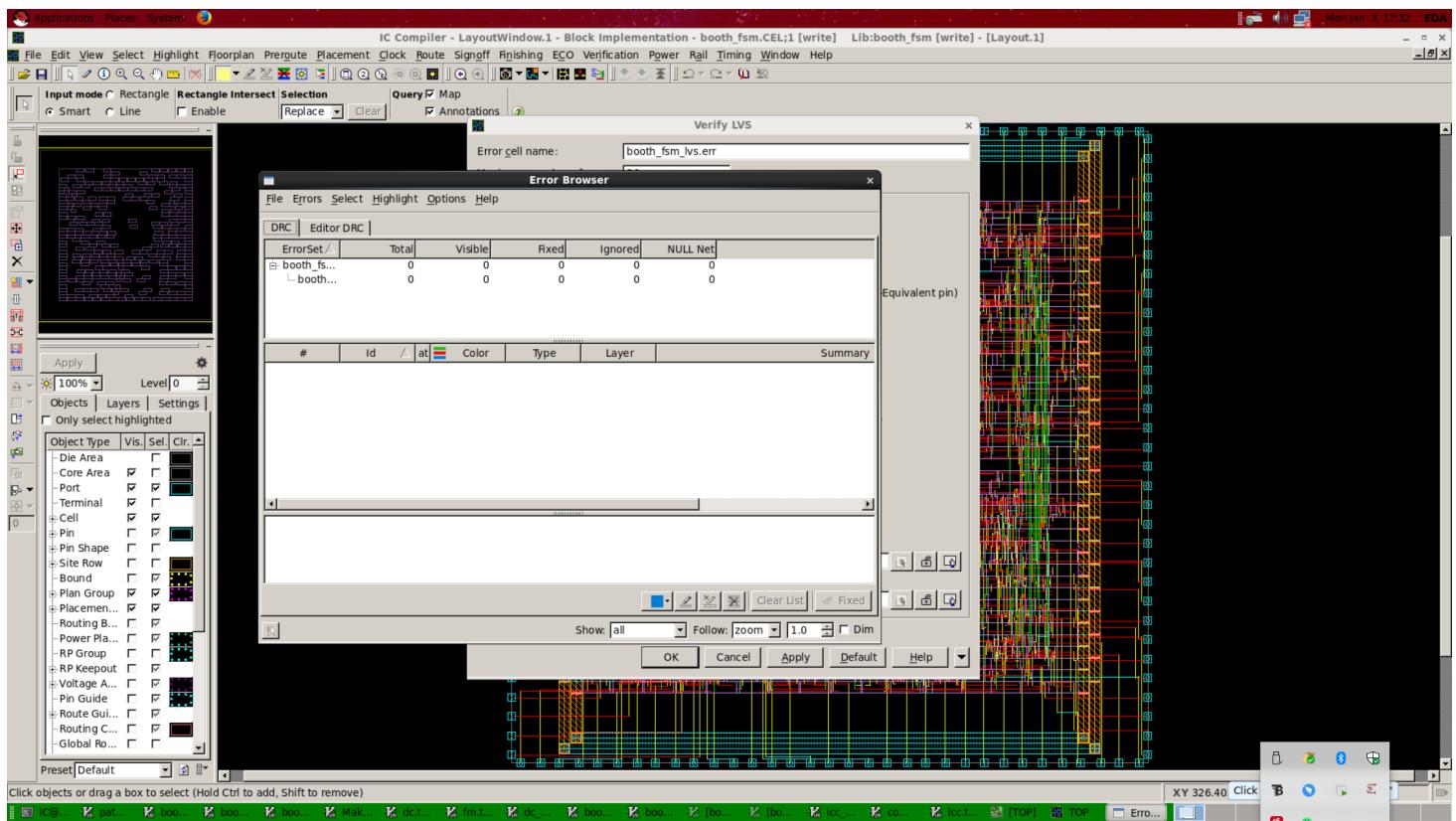
Parasitic source : LPE
 Parasitic mode : RealRC
 Extraction mode : MIN_MAX
 Extraction derating : 25/25/25

This design has no violated constraints.

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优化后的LVS/DRC检查

通过ICC的LVS/DRC检查工具检查没有问题。



后仿

后仿使用VCS进行仿真，同样使用前仿的仿真脚本，即随机生成65535个 testcase去计算；仿真结果如下

```
make[1]: Leaving directory `/home/IC/boothworkspace/1DC-RM/booth_vcs/csrc'
Chronologic VCS simulator copyright 1991-2014
Contains Synopsys proprietary information.
Compiler version I-2014.03_Full64; Runtime version I-2014.03_Full64; Jan 3 17:47 2022
sim done. num good = 65535
$finish called from file "./boothmul_tb.v", line 71.
$finish at simulation time 8650642000
V C S   S i m u l a t i o n   R e p o r t
```

总结

经过了前端和后端设计，最终我们完成了我们这个Booth乘法器的设计。

本设计在SMIC180nm工艺库下进行设计，并综合、布局、布线，最终完成这个booth乘法器，本booth乘法器采用状态机的方式进行计算，计算延时为66个cycles，面积和功耗均为要求的10%，虽然性能比较差，但是面积和功耗消耗非常小。