# Why Frame Pointers?

### C Code

```
f(1);
...;
g(2)
```

### Assembly





# Why Frame Pointers?

### Assembly

```
addiu $29,$29,-16  # space for 4 args
li $4,1  # argument setup for f
jal f  # call f

# addiu $29,$29,16  no more pop stack
...  # code for ...
addiu $29,$29,-16  # space for 4 args
li $4,2  # argument setup for g
jal g  # call g

# addiu $29,$29,16  no more pop stack
```

Save initial \$sp in the frame pointer





# Optimization

### Calculate max stack adjustment, change \$sp once

### Assembly

```
entry: addiu $29,$29,-24 # allocate space once sw $31,20($29) # save return address ...
li $4,1 # argument setup for f jal f # call f ... # code for ...
li $4,2 # argument setup for g jal g # call g
```

Stack pointer must be double-word aligned.





# More Argument Passing

What if I want to pass a half-word (short) or a byte (char)?

• Use a full word (least significant bits)

If the argument is stored on the stack (argument 5 and greater):

- Memory is big-endian
- Therefore uses higher addresses

Example: reading fifth argument of type char:

lbu \$8,19(\$29)





### More Stacks...

#### Stack frame:

Region of stack allocated by function

If a function doesn't call another one:

- Called a "leaf function"
- Doesn't save return address on stack
- Stack only used for local variables





# Leaf Function With Local Storage

#### C Code

```
void storage (int i)
{
  int x[16];
  ...
  return;
}
```

Large array, need to save in memory.





# Leaf Function With Local Storage





# Stacks: Summary

- Stack pointer points to top of stack
- Stack grows downward in memory
- Stack pointer is double-word aligned
- Stack frame is at least 24 bytes
- Registers 16-23 are saved by callee
- Argument passing: register 4-7, space on stack
- Return values in register 2-3
- Local variables, saved registers, return address saved on stack





# Memory Layout

0x7ffffffc

Stack **Dynamic Data** Static data Code  $0 \times 00400000$ Reserved

**Simulator Init SP:** 

0x7fffae50



