

# Lecture 14: Joins!

# Announcements: Two Hints

- You may want to do *Trigger activity* for project 2.
  - We've noticed those who do it have less trouble with project!
  - Seems like we're good here ☺ Exciting for us!
- We posted an activity for you to do on your own... it may overlap heavily with a ps #3 problem... (*this is not necessary but helpful*).
  - The solutions will **not** be posted.
- Sorry the Google lecture was not recorded! Last minute thing...

# 1. Nested Loop Joins

# What you will learn about in this section

1. RECAP: Joins
2. Nested Loop Join (NLJ)
3. Block Nested Loop Join (BNLJ)
4. Index Nested Loop Join (INLJ)

# RECAP: Joins

# Joins: Example

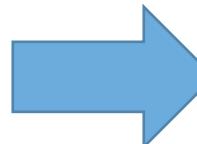
$R \bowtie S$

```
SELECT R.A, B, C, D  
FROM   R, S  
WHERE  R.A = S.A
```

Example: Returns all pairs of tuples  $r \in R, s \in S$  such that  $r.A = s.A$

R	A	B	C
1	0	1	
2	3	4	
2	5	2	
3	1	1	

S	A	D
3	7	
2	2	
2	3	



A	B	C	D
2	3	4	2

# Joins: Example

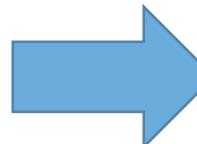
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	A	B	C	D
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2	3	4	3	

# Joins: Example

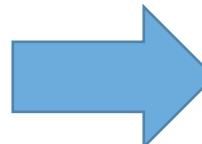
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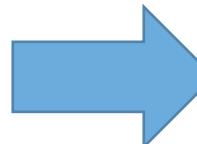
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	A	B	C	D
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2	3	4	3	
2	5	2	2	
2	5	2	3	

# Joins: Example

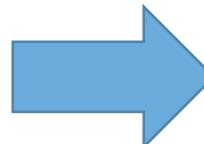
 $R \bowtie S$ 

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SELECT R.A, B, C, D  
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Example: Returns all pairs of tuples  $r \in R, s \in S$  such that  $r.A = s.A$

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3	7	
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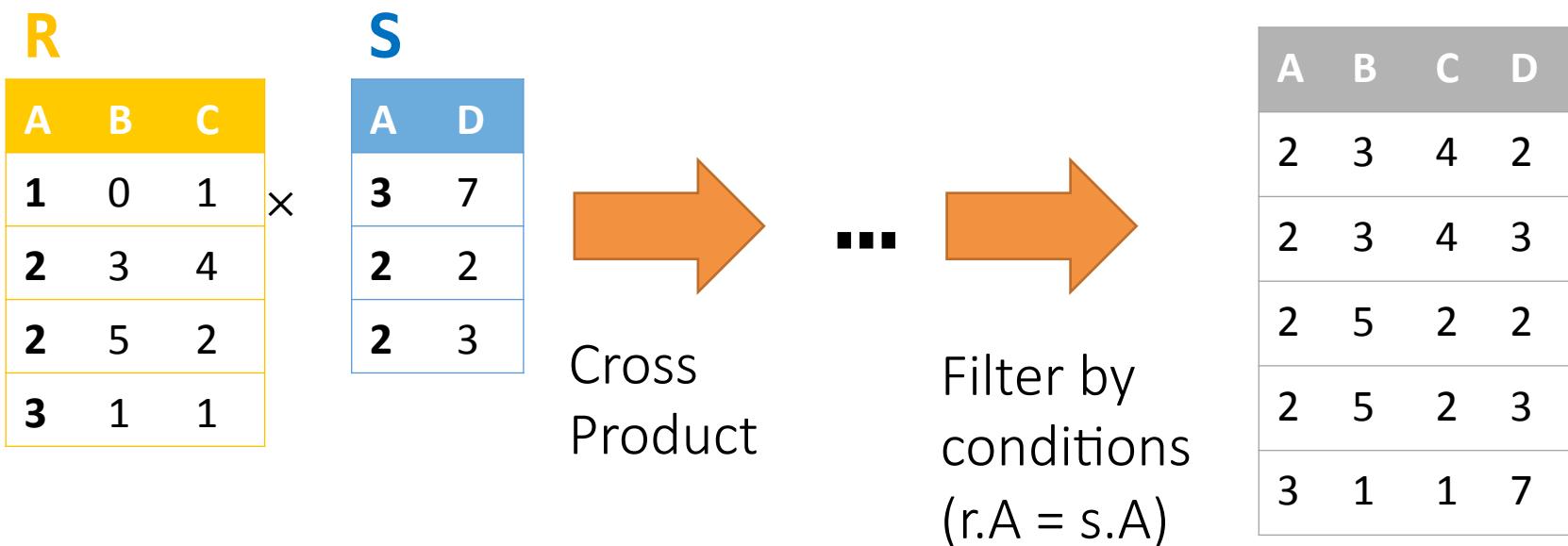
	A	B	C	D
2	3	4	2	
2	3	4	3	
2	5	2	2	
2	5	2	3	
3	1	1	7	

# Semantically: A Subset of the Cross Product

 $R \bowtie S$ 

```
SELECT R.A, B, C, D
FROM   R, S
WHERE  R.A = S.A
```

Example: Returns all pairs of tuples  $r \in R, s \in S$  such that  $r.A = s.A$



Can we actually implement a join in this way?

# Notes

- We write  $R \bowtie S$  to mean *join R and S by returning all tuple pairs where all shared attributes are equal*
- We write  $R \bowtie S \text{ on } A$  to mean *join R and S by returning all tuple pairs where attribute(s) A are equal*
- For simplicity, we'll consider joins on **two tables** and with **equality constraints** ("equijoins")

However joins *can* merge > 2 tables, and some algorithms do support non-equality constraints!

# Nested Loop Joins

# Notes

- We are again considering “IO aware” algorithms:  
***care about disk IO***
- Given a relation R, let:
  - $T(R)$  = # of tuples in R
  - $P(R)$  = # of pages in R
- Note also that we omit ceilings in calculations...  
good exercise to put back in!

Recall that we read / write  
entire pages with disk IO

# Nested Loop Join (NLJ)

Compute  $R \bowtie S$  on  $A$ :

```
for r in R:  
    for s in S:  
        if r[A] == s[A]:  
            yield (r,s)
```

# Nested Loop Join (NLJ)

Compute  $R \bowtie S$  on  $A$ :

```
for r in R:  
    for s in S:  
        if r[A] == s[A]:  
            yield (r,s)
```

Cost:

$P(R)$

**1. Loop over the tuples in  $R$**

Note that our IO cost is based on the number of *pages* loaded, not the number of tuples!

# Nested Loop Join (NLJ)

```
Compute R $\bowtie$ S on A:  
for r in R:  
    for s in S:  
        if r[A] == s[A]:  
            yield (r,s)
```

Cost:

$$P(R) + T(R)*P(S)$$

1. Loop over the tuples in R
2. For every tuple in R, loop over all the tuples in S

Have to read *all of S* from disk for *every tuple in R!*

# Nested Loop Join (NLJ)

```
Compute R  $\bowtie$  S on A:  
for r in R:  
    for s in S:  
        if r[A] == s[A]:  
            yield (r, s)
```

Cost:

$$P(R) + T(R)*P(S)$$

1. Loop over the tuples in R
2. For every tuple in R, loop over all the tuples in S
- 3. Check against join conditions**

Note that NLJ can handle things other than equality constraints... just check in the *if* statement!

# Nested Loop Join (NLJ)

Compute  $R \bowtie S$  on  $A$ :

```
for r in R:
    for s in S:
        if r[A] == s[A]:
            yield (r, s)
```

What would  $OUT$  be if our join condition is trivial (if TRUE)? cross product

$OUT$  could be bigger than  $P(R)*P(S)$ ... but usually not that bad

Cost:

$$P(R) + T(R)*P(S) + OUT$$

1. Loop over the tuples in  $R$
2. For every tuple in  $R$ , loop over all the tuples in  $S$
3. Check against join conditions
4. Write out (to page, then when page full, to disk)

# Nested Loop Join (NLJ)

Compute  $R \bowtie S$  on  $A$ :

```
for r in R:  
    for s in S:  
        if r[A] == s[A]:  
            yield (r,s)
```

Cost:

$$P(R) + T(R)*P(S) + OUT$$

*What if  $R$  ("outer") and  $S$  ("inner") switched?*



$$P(S) + T(S)*P(R) + OUT$$

Outer vs. inner selection makes a huge difference-  
DBMS needs to know which relation is smaller!

# IO-Aware Approach

# Block Nested Loop Join (BNLJ)

Compute  $R \bowtie S \text{ on } A$ :

```
for each  $B-1$  pages pr of R:  
    for page ps of S:  
        for each tuple r in pr:  
            for each tuple s in ps:  
                if  $r[A] == s[A]$ :  
                    yield (r,s)
```

Given  $B+1$  pages of memory

Cost:

$P(R)$

1. Load in  $B-1$  pages of R at a time (leaving 1 page each free for S & output)

*Note: There could be some speedup here due to the fact that we're reading in multiple pages sequentially however we'll ignore this here!*

# Block Nested Loop Join (BNLJ)

Compute  $R \bowtie S$  on  $A$ :

```
for each  $B-1$  pages  $pr$  of  $R$ :
    for page  $ps$  of  $S$ :
        for each tuple  $r$  in  $pr$ :
            for each tuple  $s$  in  $ps$ :
                if  $r[A] == s[A]$ :
                    yield  $(r, s)$ 
```

Given  $B+1$  pages of memory

Cost:

$$P(R) + P(R)/B - 1 \cdot P(S)$$

1. Load in  $B-1$  pages of  $R$  at a time (leaving 1 page each free for  $S$  & output)
2. **For each  $(B-1)$ -page segment of  $R$ , load each page of  $S$**

Note: Faster to iterate over the *smaller* relation first!

# Block Nested Loop Join (BNLJ)

Compute  $R \bowtie S$  on  $A$ :

```
for each  $B-1$  pages  $pr$  of  $R$ :
    for page  $ps$  of  $S$ :
        for each tuple  $r$  in  $pr$ :
            for each tuple  $s$  in  $ps$ :
                if  $r[A] == s[A]$ :
                    yield  $(r, s)$ 
```

Given  $B+1$  pages of memory

Cost:

$$P(R) + P(R)/(B-1) P(S)$$

1. Load in  $B-1$  pages of  $R$  at a time (leaving 1 page each free for  $S$  & output)
2. For each  $(B-1)$ -page segment of  $R$ , load each page of  $S$
- 3. Check against the join conditions**

BNLJ can also handle non-equality constraints

# Block Nested Loop Join (BNLJ)

Compute  $R \bowtie S$  on  $A$ :

```

for each  $B-1$  pages pr of R:
    for page ps of S:
        for each tuple r in pr:
            for each tuple s in ps:
                if  $r[A] == s[A]$ :
                    yield (r,s)
    
```

Again,  $OUT$  could be bigger than  $P(R)*P(S)...$  but usually not that bad

Given  $B+1$  pages of memory

Cost:

$$P(R) + P(R)/(B-1) P(S) + OUT$$

1. Load in  $B-1$  pages of R at a time (leaving 1 page each free for S & output)
2. For each  $(B-1)$ -page segment of R, load each page of S
3. Check against the join conditions
4. Write out

# BNLJ vs. NLJ: Benefits of IO Aware

- In BNLJ, by loading larger chunks of R, we minimize the number of full *disk reads* of S
  - We only read all of S from disk for ***every (B-1)-page segment of R!***
  - Still the full cross-product, but more done only *in memory*

NLJ

$$P(R) + T(R)*P(S) + OUT$$



BNLJ

$$P(R) + P(R)/B-1 P(S) + OUT$$

BNLJ is faster by roughly  $(B-1)T(R)/P(R)$  !

# BNLJ vs. NLJ: Benefits of IO Aware

- Example:
  - R: 500 pages
  - S: 1000 pages
  - 100 tuples / page
  - We have 12 pages of memory ( $B = 11$ ) 1 for output
- NLJ: Cost =  $500 + 50,000 * 1000 = 50 \text{ Million IOs} \approx \underline{140 \text{ hours}}$
- BNLJ: Cost =  $500 + 500 * 1000 / 10 = 50 \text{ Thousand IOs} \approx \underline{0.14 \text{ hours}}$

*Ignoring OUT here...*

A very real difference from a small  
change in the algorithm!

# Smarter than Cross-Products

# Smarter than Cross-Products: From Quadratic to Nearly Linear

- All joins that compute the *full cross-product* have some **quadratic** term
  - For example we saw:

$$\text{NLJ} \quad P(R) + \textcolor{red}{T(R)P(S)} + OUT$$

$$\text{BNLJ} \quad P(R) + \textcolor{green}{P(R)/B-1} \textcolor{black}{P(S)} + OUT$$

amortizing  
the overhead  
of the read

- Now we'll see some (nearly) linear joins:
  - $\sim O(P(R) + P(S) + OUT)$ , where again ***OUT*** could be quadratic but is usually better

We get this gain by *taking advantage of structure*- moving to equality constraints ("equijoin") only!

# Index Nested Loop Join (INLJ)

Cost:

Compute  $R \bowtie S$  on  $A$ :

Given index  $\text{idx}$  on  $S.A$ :

for  $r$  in  $R$ :

s in  $\text{idx}(r[A])$ :  
probing the index

yield  $r, s$

$$P(R) + T(R)*L + OUT$$

where  $L$  is the IO cost to access all the distinct values in the index; assuming these fit on one page,  $L \sim 3$  is good est.

$L=4$  for huge database

→ We can use an **index** (e.g. B+ Tree) to *avoid doing the full cross-product!*

# Joins: A Cage Match

*Message: It's all about the memory!*

# Today's Lecture

1. Sort-Merge Join (SMJ)
2. Hash Join (HJ)
3. The Cage Match: SMJ vs. HJ

# 1. Sort-Merge Join (SMJ)



# What you will learn about in this section

1. Sort-Merge Join
2. “Backup” & Total Cost
3. Optimizations
4. ACTIVITY: Sequential Flooding

# Sort Merge Join (SMJ): Basic Procedure

To compute  $R \bowtie S$  on  $A$ :

1. Sort  $R, S$  on  $A$  using ***external merge sort***
2. ***Scan*** sorted files and “merge”
3. [May need to “backup”- see next subsection]

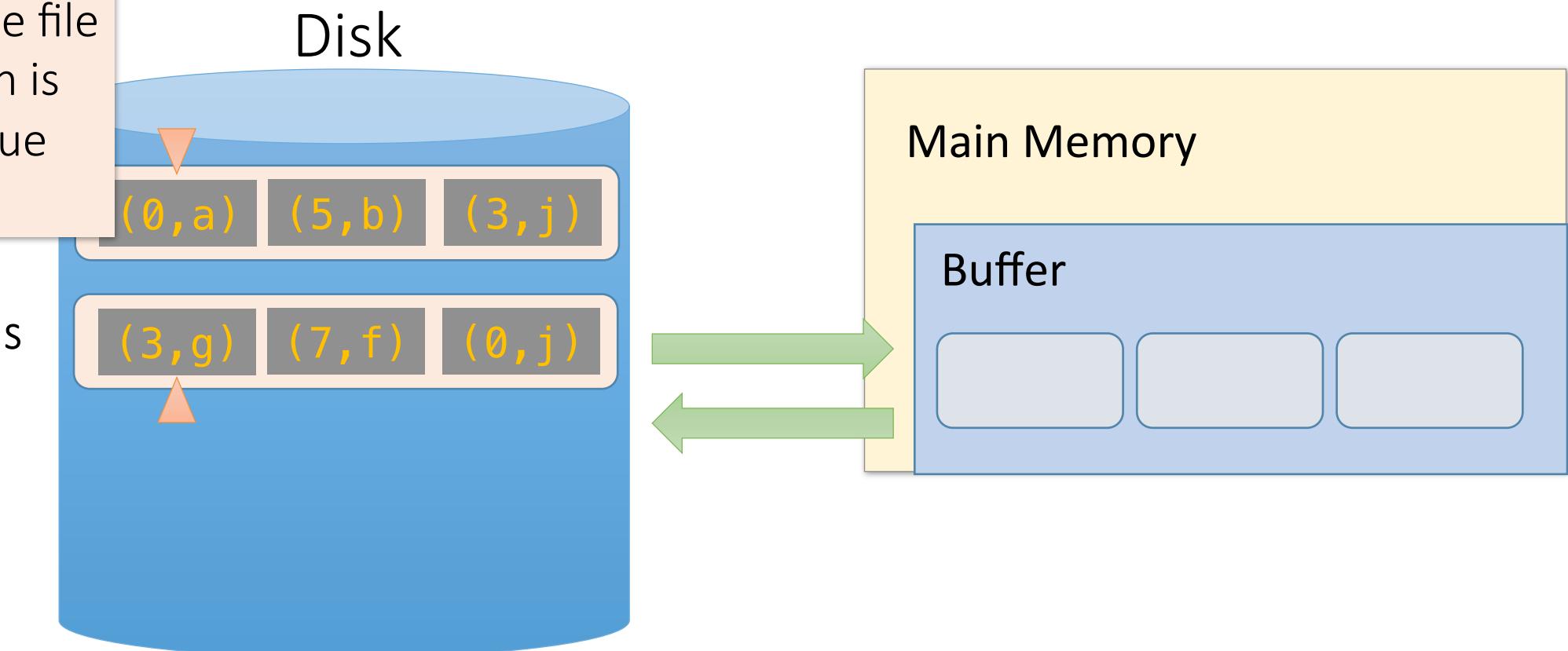
Note that we are only considering equality join conditions here

Note that if  $R, S$  are already sorted on  $A$ ,  
SMJ will be awesome!

# SMJ Example: $R \bowtie S$ on $A$ with 3 page buffer

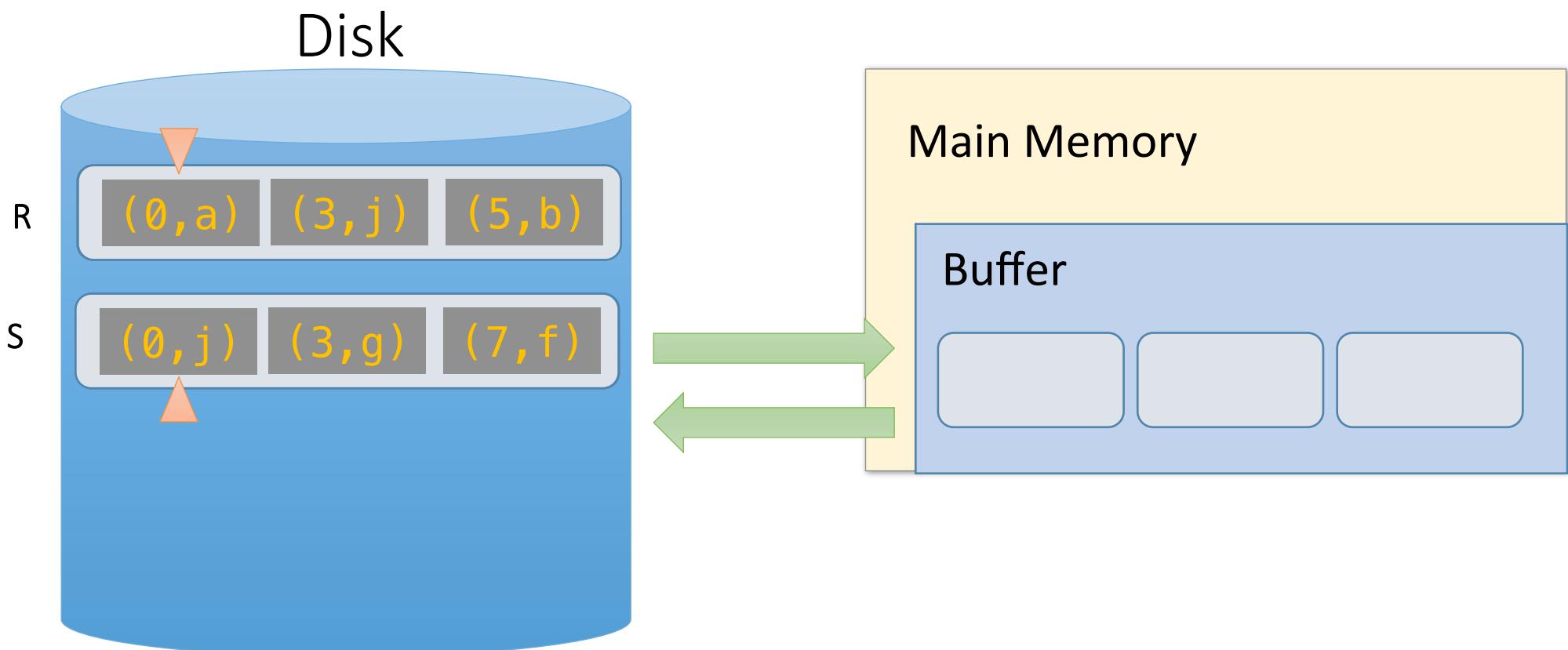
- For simplicity: Let each page be **one tuple**, and let the first value be A

We show the file HEAD, which is the next value to be read!



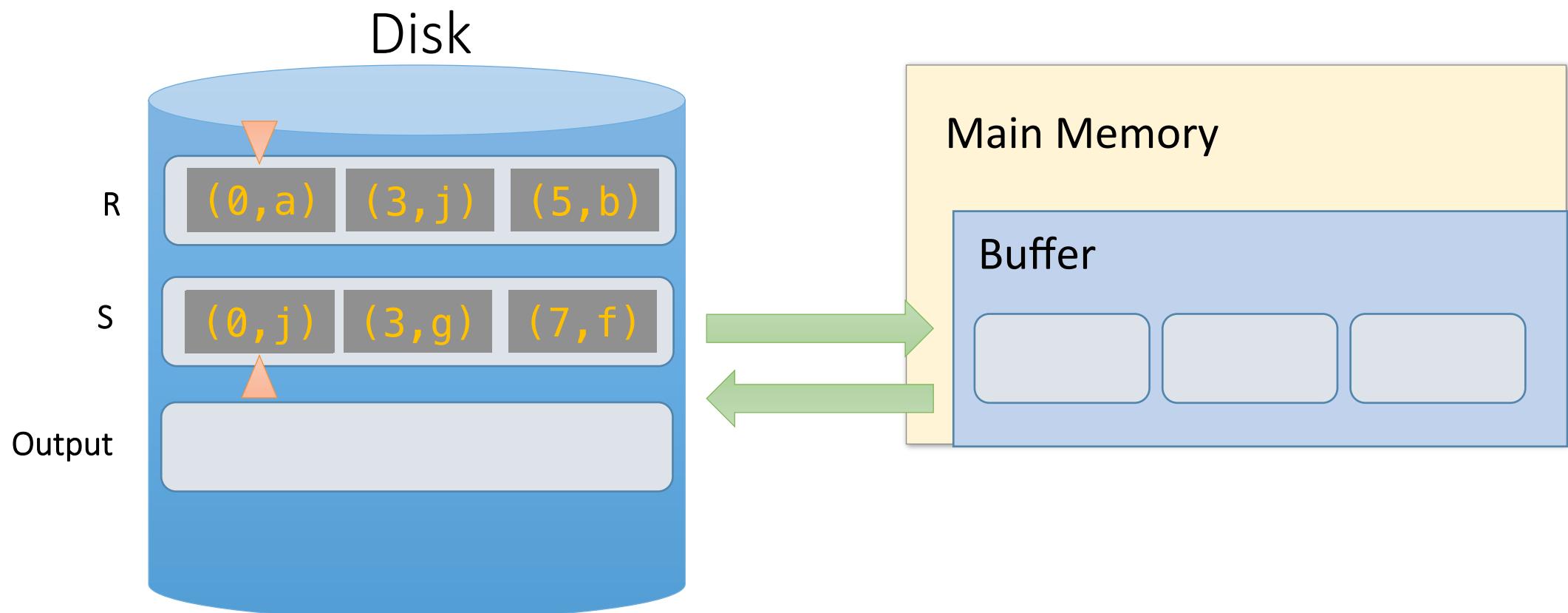
# SMJ Example: $R \bowtie S \text{ on } A$ with 3 page buffer

1. Sort the relations R, S on the join key (first value)



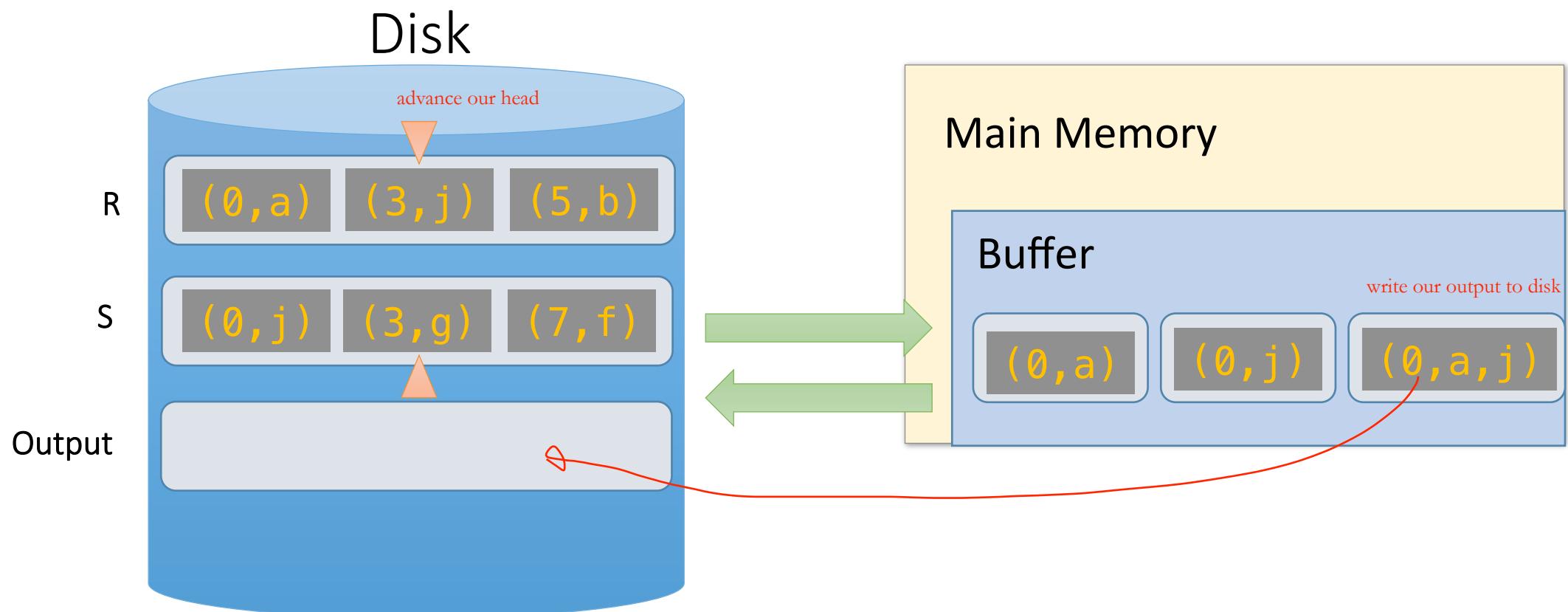
# SMJ Example: $R \bowtie S$ on $A$ with 3 page buffer

2. Scan and “merge” on join key!



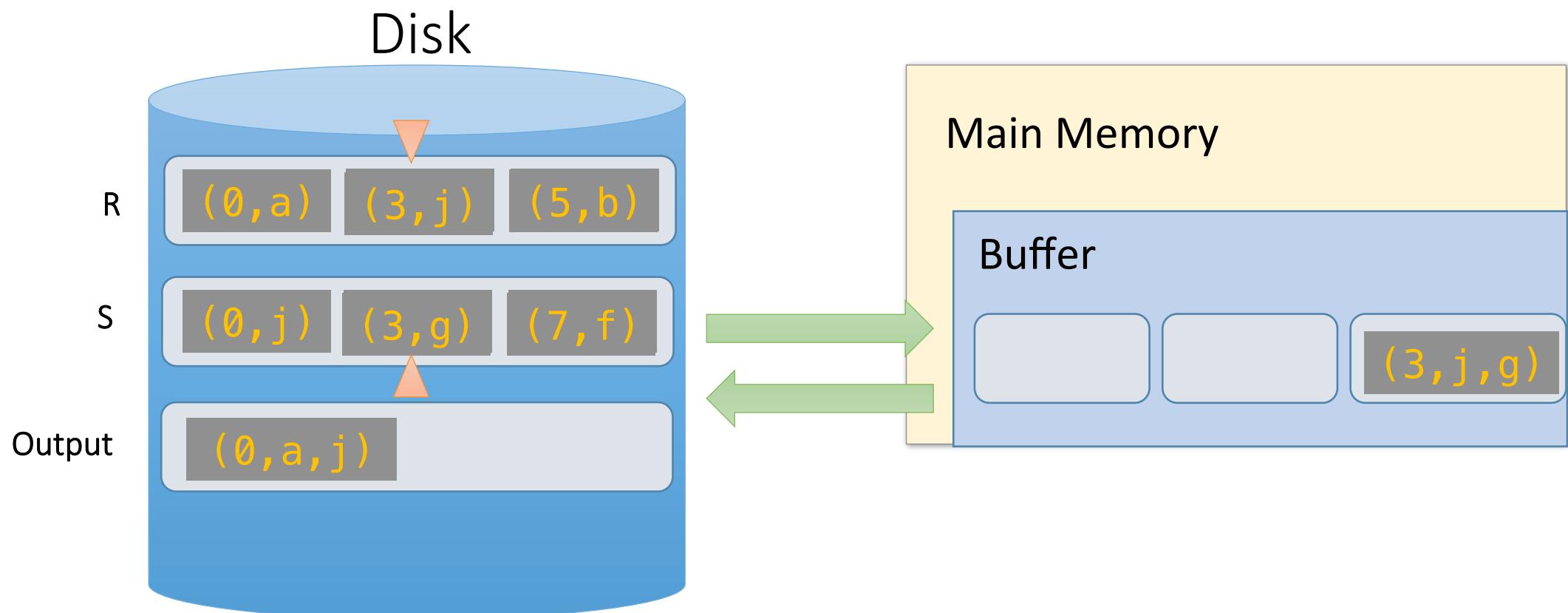
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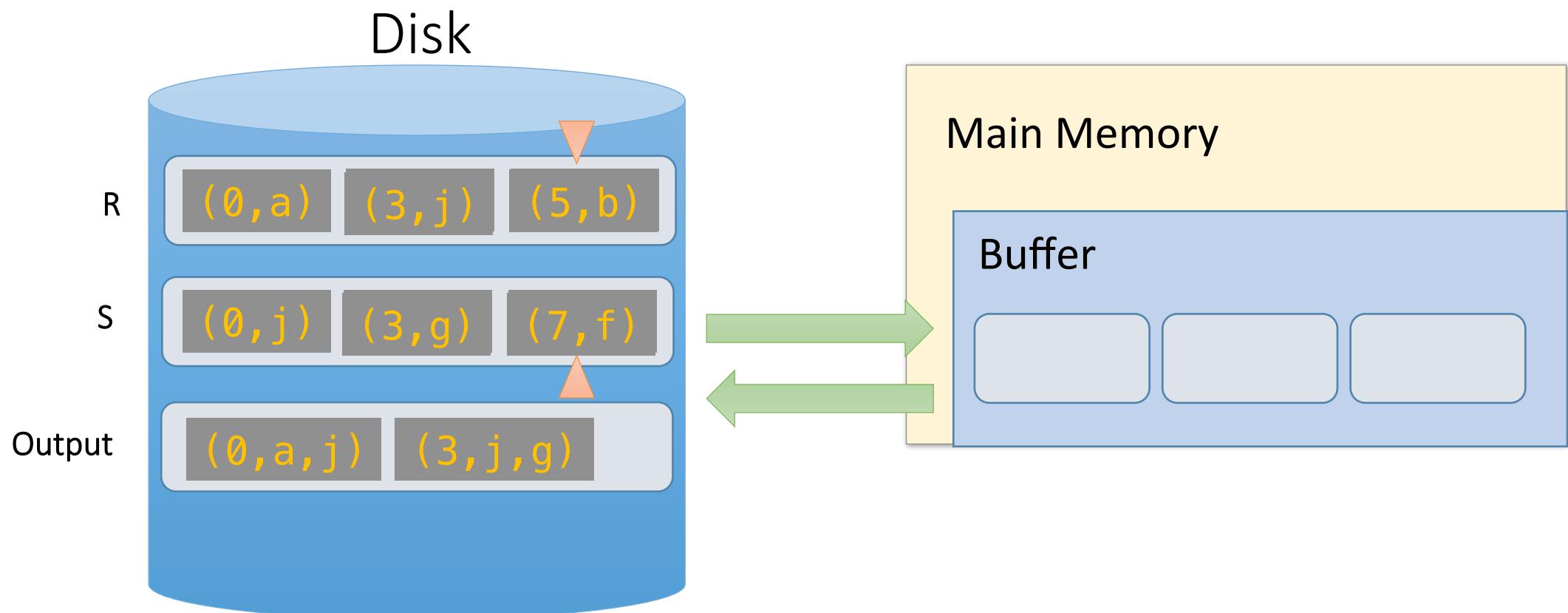
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2. Scan and “merge” on join key!



# SMJ Example: $R \bowtie S$ on $A$ with 3 page buffer

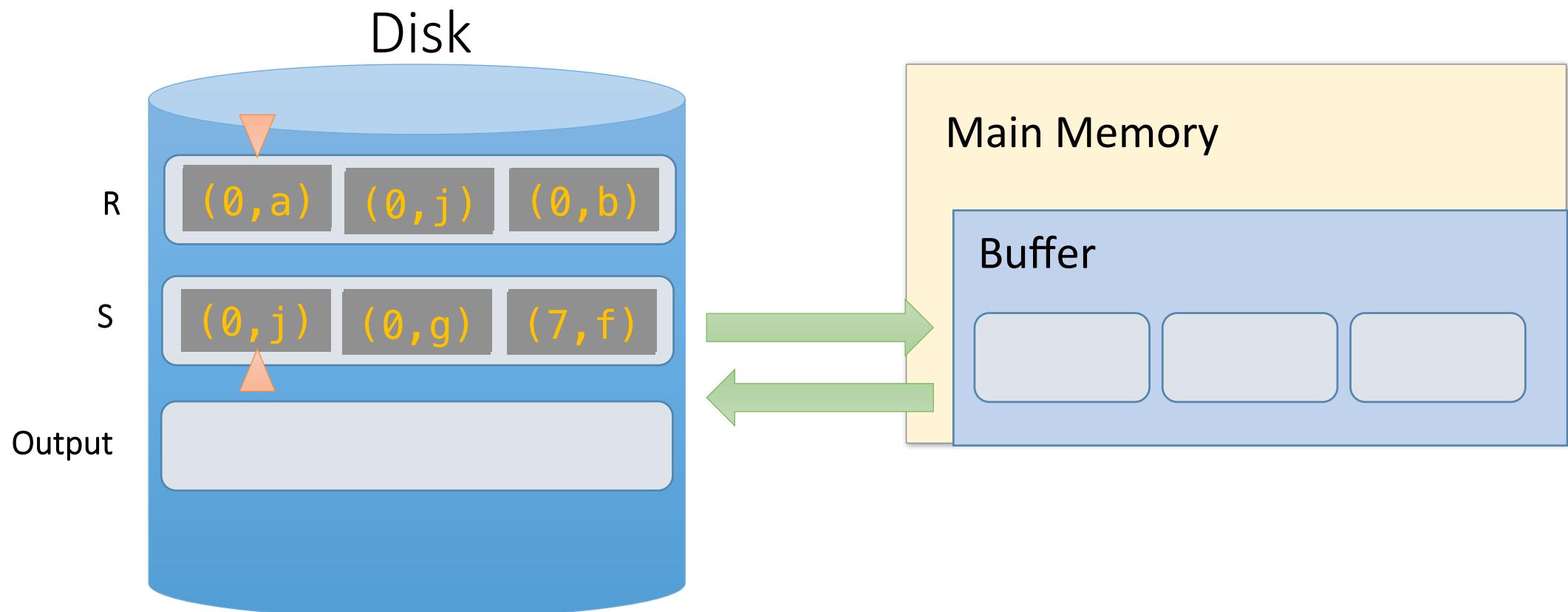
2. Done!



What happens with duplicate join keys?

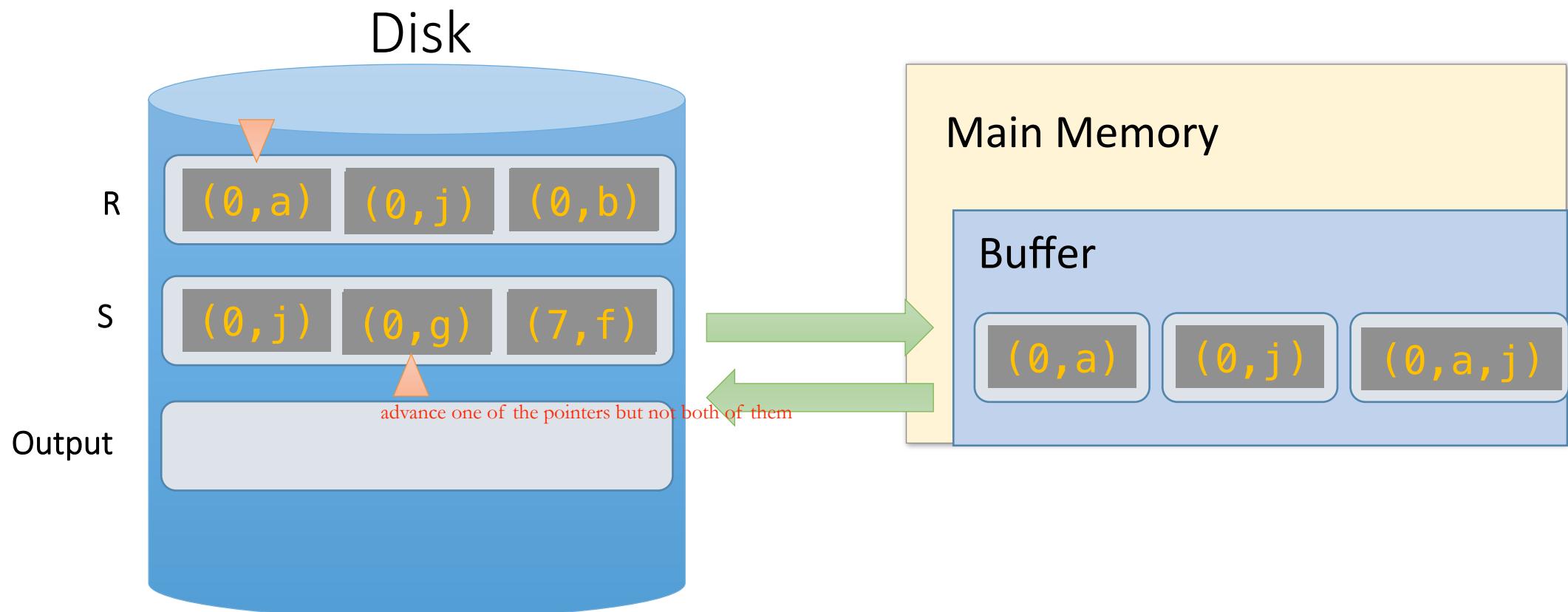
# Multiple tuples with Same Join Key: “Backup”

1. Start with sorted relations, and begin scan / merge...



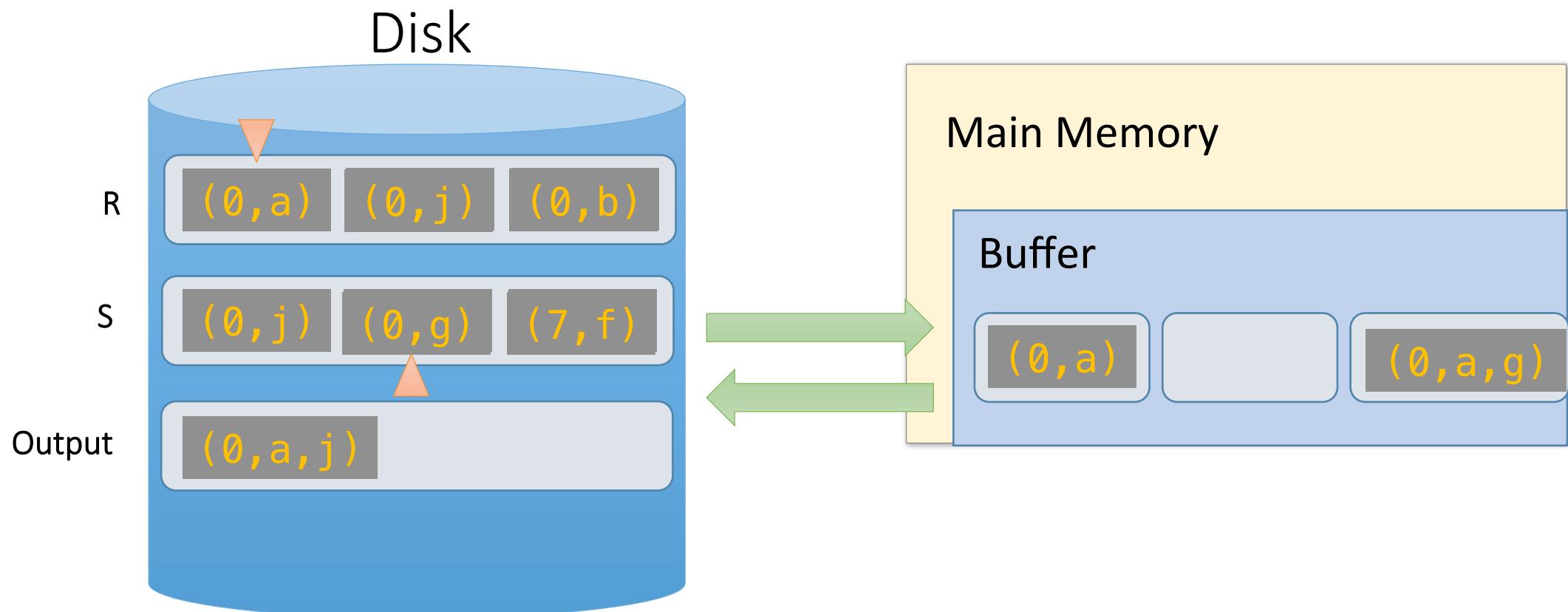
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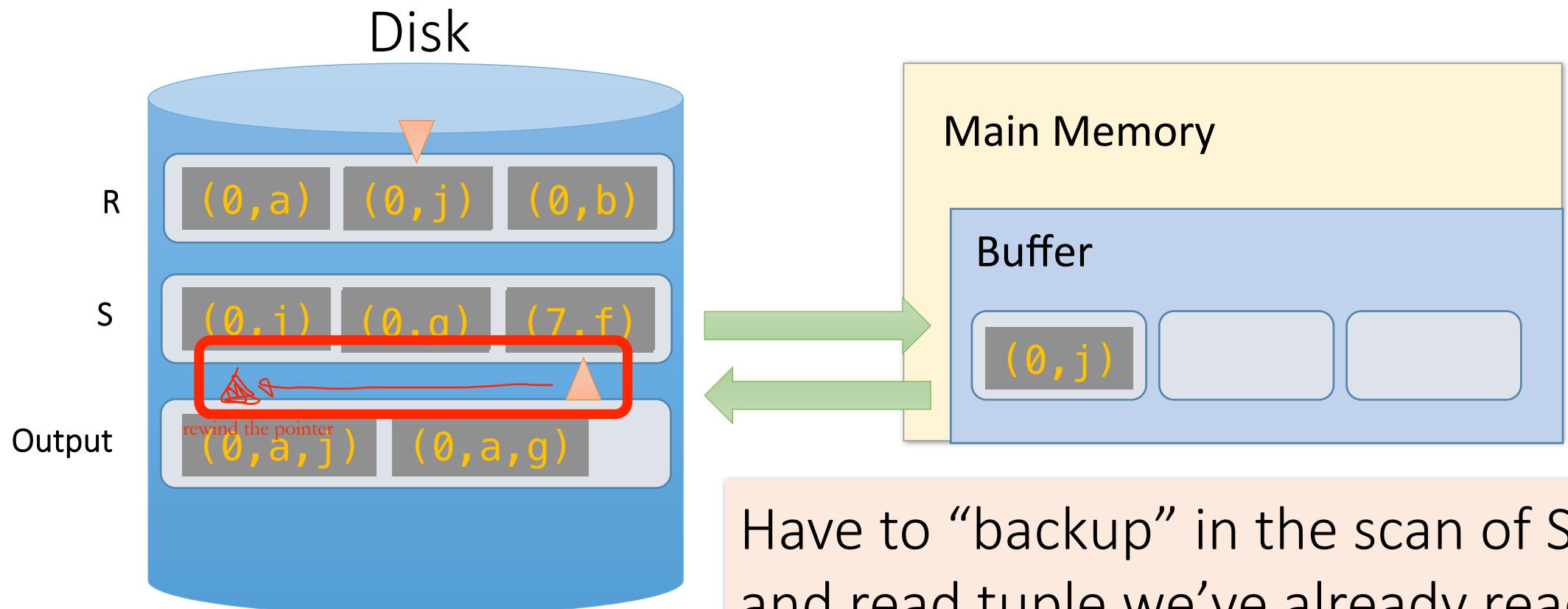
# Multiple tuples with Same Join Key: “Backup”

1. Start with sorted relations, and begin scan / merge...



# Multiple tuples with Same Join Key: “Backup”

1. Start with sorted relations, and begin scan / merge...



# Backup

- At best, no backup → scan takes  $P(R) + P(S)$  reads
  - For ex: if no duplicate values in join attribute
- At worst (e.g. full backup each time), scan could take  $P(R) * P(S)$  reads!
  - For ex: if *all* duplicate values in join attribute, i.e. all tuples in R and S have the same value for the join attribute
  - Roughly: For each page of R, we'll have to *back up* and read each page of S...
- Often not that bad however, plus we can:
  - Leave more data in buffer (for larger buffers)
  - Can “zig-zag” (see animation)

# SMJ: Total cost

- Cost of SMJ is **cost of sorting R and S...**
- Plus the **cost of scanning**:  $\sim P(R) + P(S)$ 
  - Because of *backup*: in worst case  $P(R)*P(S)$ ; but this would be very unlikely
- Plus the **cost of writing out**:  $\sim P(R) + P(S)$  but in worst case  $T(R)*T(S)$

$\sim \text{Sort}(P(R)) + \text{Sort}(P(S))$   
 $+ P(R) + P(S) + \text{OUT}$

Recall:  $\text{Sort}(N) \approx 2N(\log \downarrow B \mathbf{N/2(B+1)} \downarrow + 1)$   
Note: *this is using repacking, where we estimate that we can create initial runs of length  $\sim 2(B+1)$*

# SMJ vs. BNLJ: Steel Cage Match

- If we have 100 buffer pages,  $P(R) = 1000$  pages and  $P(S) = 500$  pages:
  - Sort both in two passes:  $2 * 2 * 1000 + 2 * 2 * 500 = 6,000 \text{ IOs}$
  - Merge phase  $1000 + 500 = 1,500 \text{ IOs}$
  - = 7,500 IOs + OUT

What is BNLJ?

- $500 + 1000 * \lceil 500 / 98 \rceil = \underline{6,500 \text{ IOs + OUT}}$
- But, if we have 35 buffer pages?
  - Sort Merge has same behavior (still 2 passes)
  - BNLJ? 15,500 IOs + OUT!



SMJ is ~ linear vs. BNLJ is quadratic...  
But it's all about the memory.

# A Simple Optimization: Merges Merged!

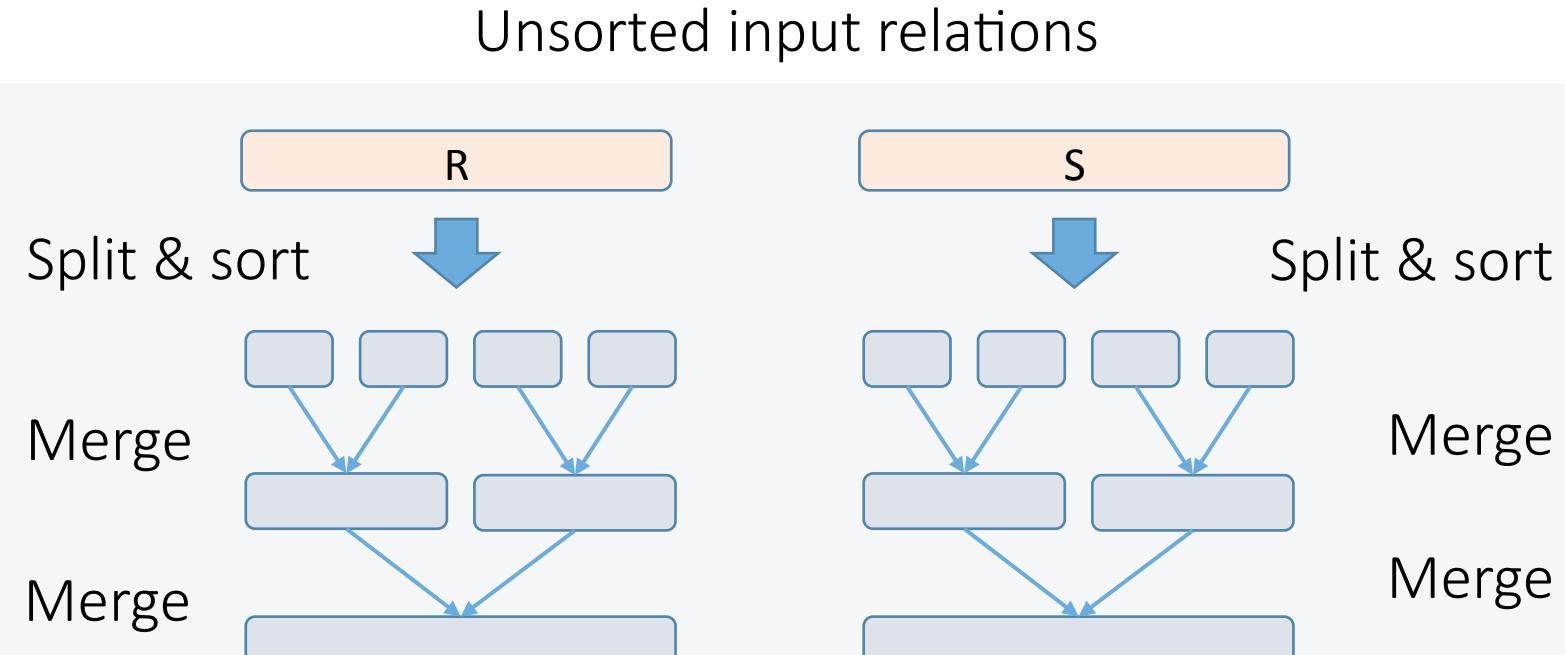
Given  $B+1$  buffer pages

- SMJ is composed of a ***sort phase*** and a ***merge phase***
- During the ***sort phase***, run passes of external merge sort on R and S
  - Suppose at some point, R and S have  $\leq B$  (sorted) runs in total
    - We could do two merges (for each of R & S) at this point, complete the sort phase, and start the merge phase...
    - OR, we could combine them: do **one** B-way merge and complete the join!

# Un-Optimized SMJ

Given  $B+1$  buffer pages

## Sort Phase (Ext. Merge Sort)



## Merge / Join Phase

Joined output  
file created!

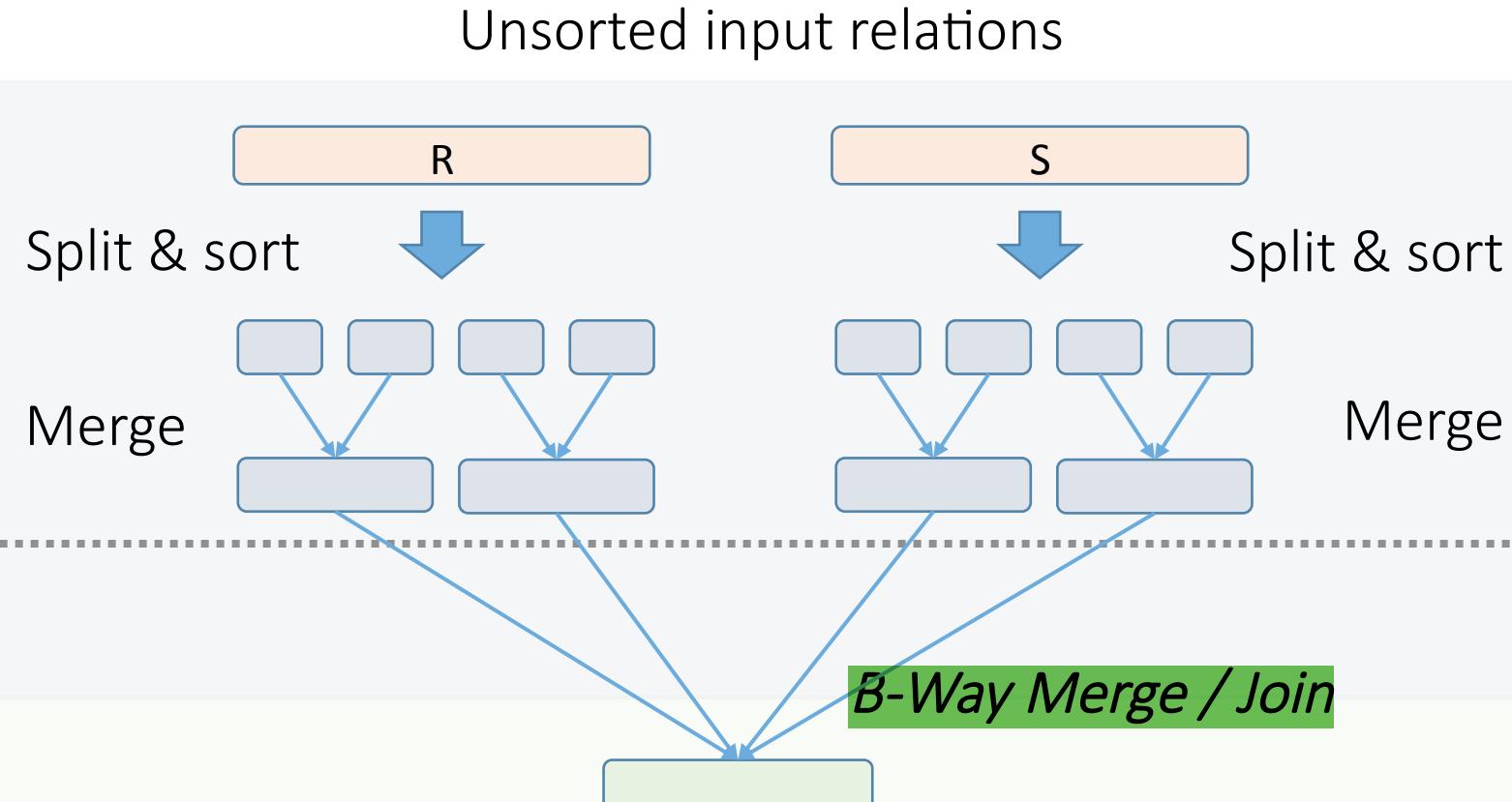
# Simple SMJ Optimization

Given  $B+1$  buffer pages

Sort Phase  
(Ext. Merge Sort)

$\leq B$  total runs

Merge / Join Phase



Joined output  
file created!

# Simple SMJ Optimization

Given  $B+1$  buffer pages

- Now, on this last pass, we only do  $P(R) + P(S)$  IOs to complete the join!
- If we can initially split R and S into **B total runs each of length approx.  $\leq 2(B+1)$** , *assuming repacking lets us create initial runs of  $\sim 2(B+1)$* - then we only need  **$3(P(R) + P(S)) + OUT$**  for SMJ!
  - 2 R/W per page to sort runs in memory, 1 R per page to B-way merge / join!
- How much memory for this to happen?
  - $P(R)+P(S)/B \leq 2(B+1) \Rightarrow \sim P(R)+P(S) \leq 2B^2$
  - **Thus,  $\max\{P(R), P(S)\} \leq B^2$  is an approximate sufficient condition**

See Lecture 15,  
Slide 13-14 – to  
clarify this slide.

If the larger of R,S has  $\leq B^2$  pages, then SMJ costs  
 $3(P(R)+P(S)) + OUT$ !

# Takeaway points from SMJ

If input already sorted on join key, skip the sorts.

- SMJ is basically linear.
- Nasty but unlikely case: Many duplicate join keys.

SMJ needs to sort **both** relations

- If  $\max \{ P(R), P(S) \} < B^2$  then cost is  $3(P(R)+P(S)) + OUT$

# 4. Hash Join (HJ)



# What you will learn about in this section

1. Hash Join
2. Memory requirements

# Recall: Hashing

- **Magic of hashing:**
  - A hash function  $h_B$  maps into  $[0, B-1]$
  - And maps nearly uniformly
- A hash **collision** is when  $x \neq y$  but  $h_B(x) = h_B(y)$ 
  - Note however that it will never occur that  $x = y$  but  $h_B(x) \neq h_B(y)$
- We hash on an attribute  $A$ , so our hash function is  $h_B(t)$  has the form  $h_B(t.A)$ .
  - **Collisions** may be more frequent.

# Recall: Mad Hash Collisions



Say something here to justify this slide's existence? [TODO]

# Hash Join: High-level procedure

To compute  $R \bowtie S$  on  $A$ :

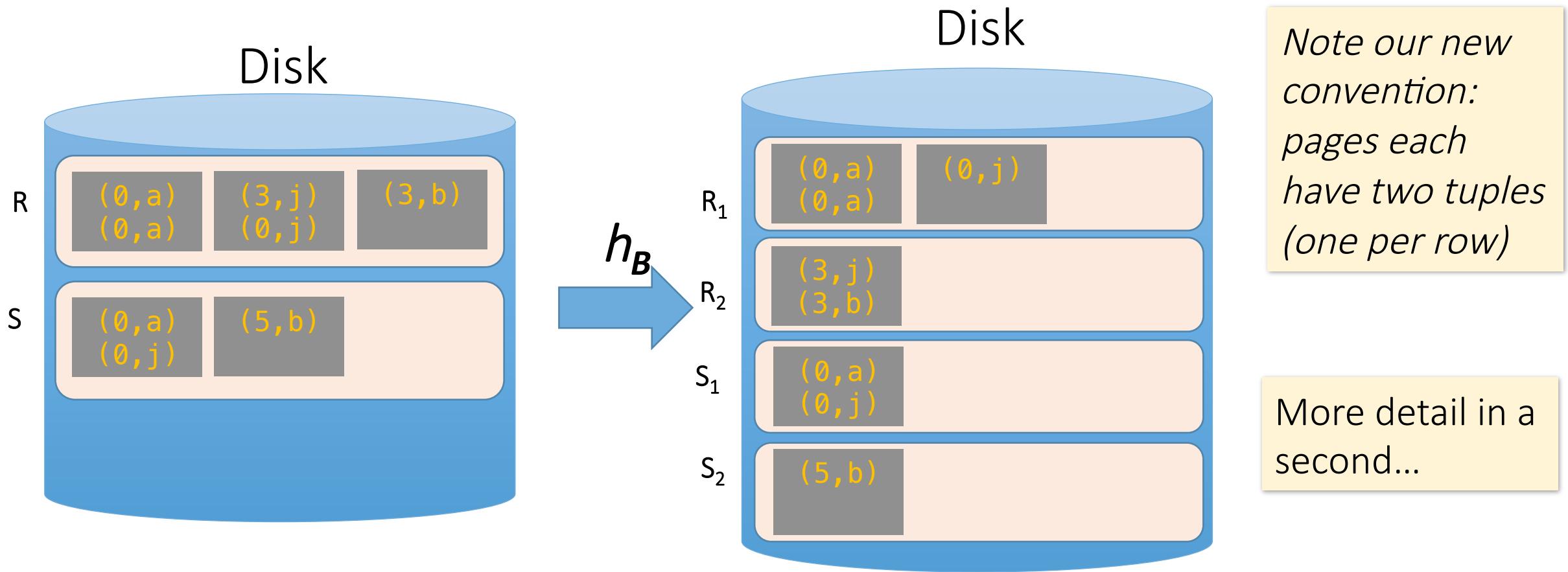
Note again that we are only considering equality constraints here

1. **Partition Phase:** Using one (shared) hash function  $h_B$ , partition  $R$  and  $S$  into  $B$  buckets
2. **Matching Phase:** Take pairs of buckets whose tuples have the same values for  $h$ , and join these
  1. Use BNJ here; or hash again → either way, operating on small partitions so fast!

We *decompose* the problem using  $h_B$ , then complete the join

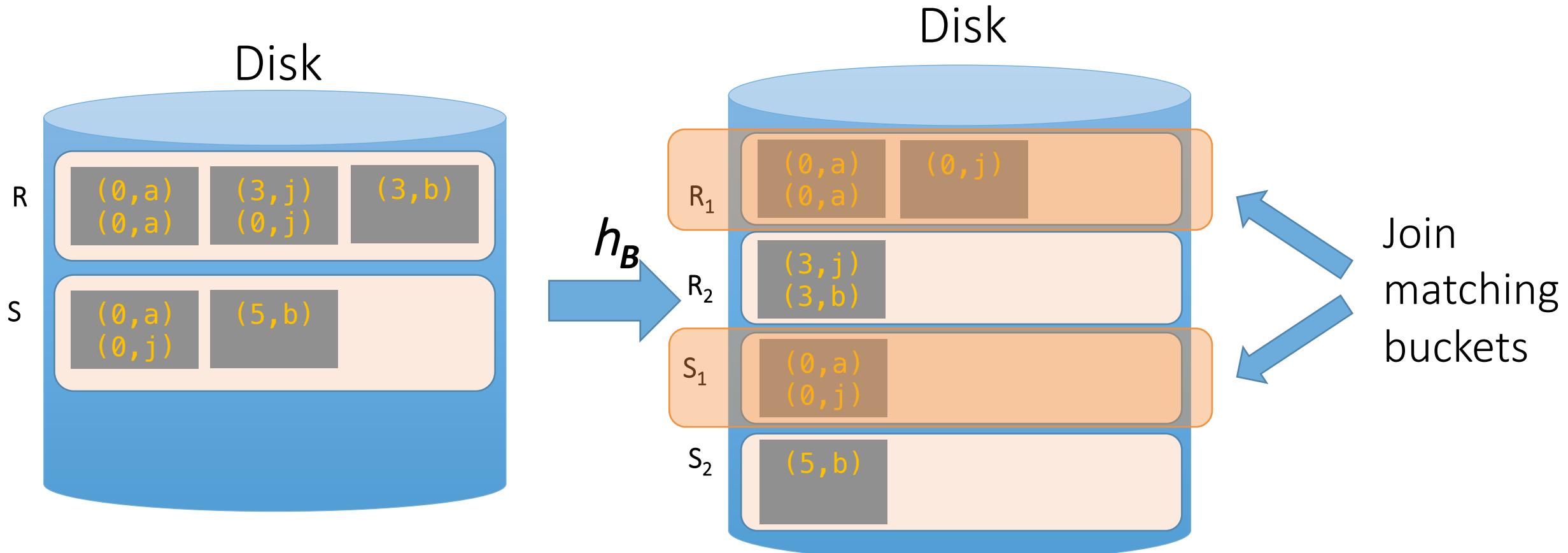
# Hash Join: High-level procedure

**1. Partition Phase:** Using one (shared) hash function  $h_B$ , partition R and S into  $B$  buckets



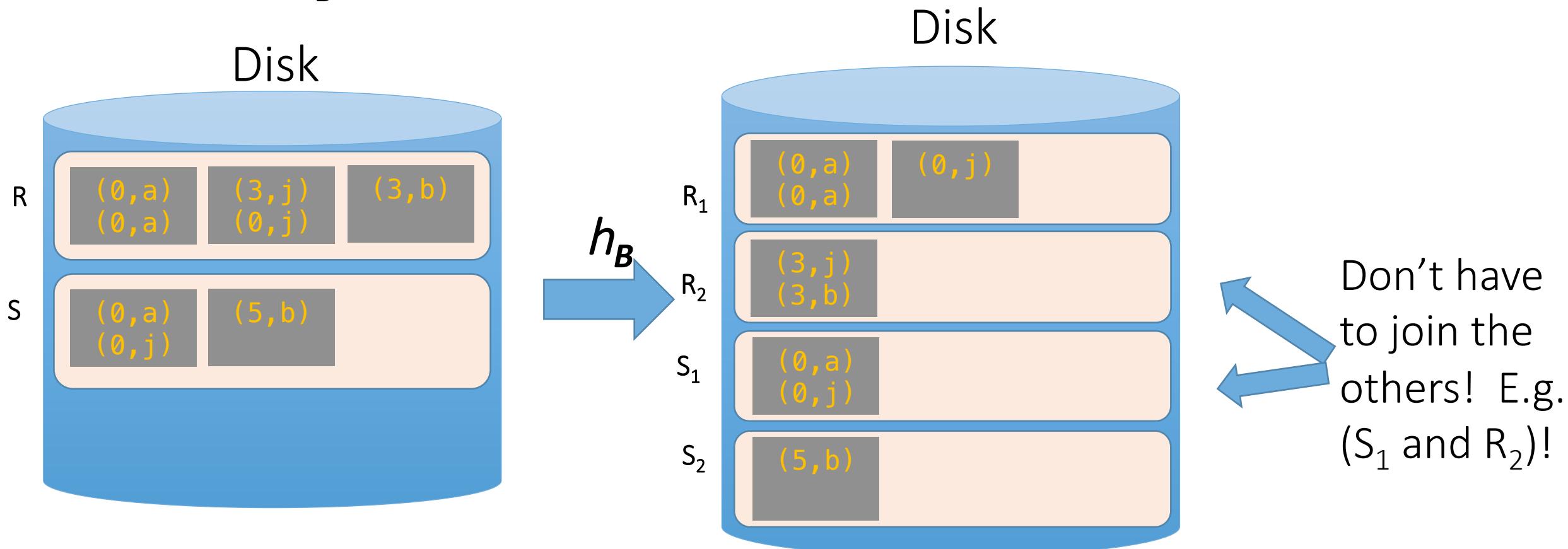
# Hash Join: High-level procedure

**2. Matching Phase:** Take pairs of buckets whose tuples have the same values for  $h_B$ , and join these



# Hash Join: High-level procedure

**2. Matching Phase:** Take pairs of buckets whose tuples have the same values for  $h_B$ , and join these



# Hash Join Phase 1: Partitioning

**Goal:** For each relation, partition relation into **buckets** such that if  $h_B(t.A) = h_B(t'.A)$  they are in the same bucket

Given  $B+1$  buffer pages, we partition into  $B$  buckets:

- We use  $B$  buffer pages for output (one for each bucket), and 1 for input
  - The “dual” of sorting.
  - For each tuple  $t$  in input, copy to buffer page for  $h_B(t.A)$
  - When page fills up, flush to disk.

# How big are the resulting buckets?

Given  $B+1$  buffer pages

- Given **N input pages, we partition into B buckets:**
  - → Ideally our buckets are each of size  $\sim N/B$  pages
- What happens if there are **hash collisions?**
  - Buckets could be  $> N/B$
  - **We'll do several passes...**
- What happens if there are **duplicate join keys?**
  - Nothing we can do here... could have some **skew** in size of the buckets

# How big *do we want* the resulting buckets?

- Ideally, our buckets would be of size  $\leq B-1$  pages
  - 1 for input page, 1 for output page,  $B-1$  for each bucket

- Recall: If we want to join a bucket from R and one from S, we can do BNLJ in linear time if for *one of them (wlog say R)*,  $P(R) \leq B-1$ !

- And more generally, being able to fit bucket in memory is advantageous

- We can keep partitioning buckets that are  $> B-1$  pages, until they are  $\leq B-1$  pages
  - Using a new hash key which will split them...

Given  $B+1$  buffer pages

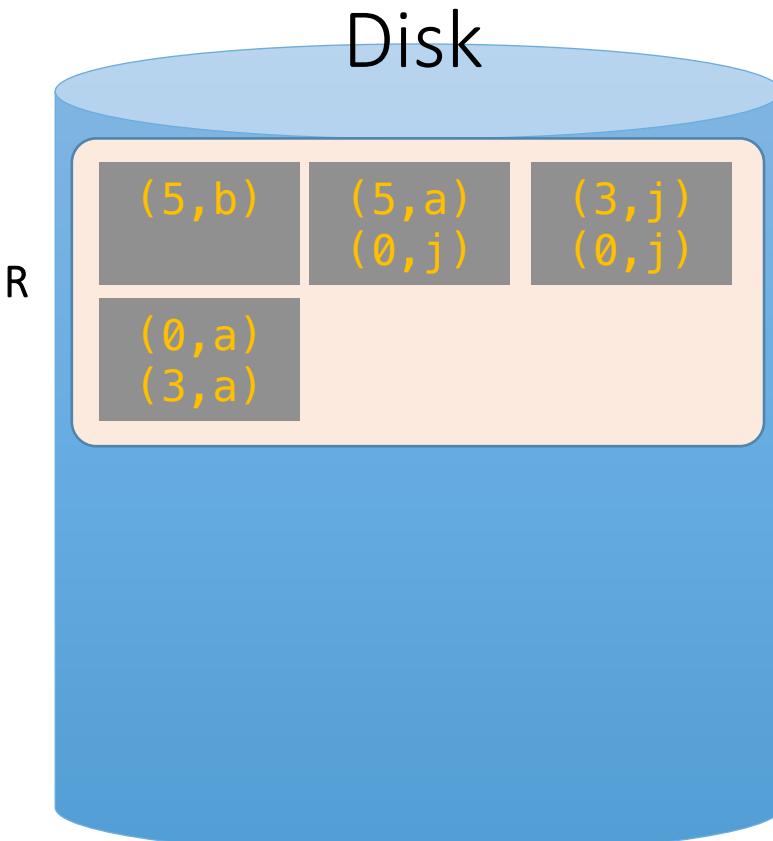
Recall for BNLJ:  
 $P(R) + P(R)P(S)/B-1$

We'll call each of these  
 a "pass" again...

# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

We partition into  $B = 2$  buckets **using hash function  $h_2$**  so that we can have one buffer page for each partition (and one for input)



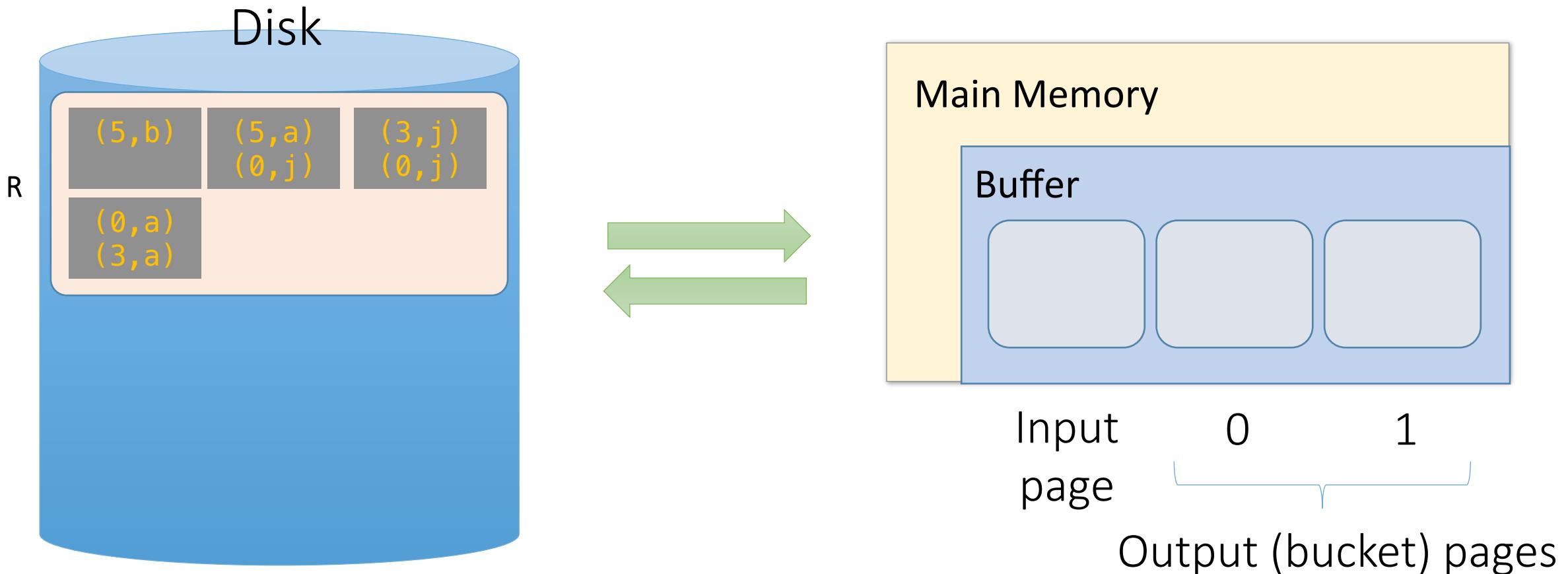
For simplicity, we'll look at partitioning one of the two relations- we just do the same for the other relation!

Recall: our goal will be to get  $B = 2$  buckets of size  $\leq B-1 \rightarrow 1$  page each

# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

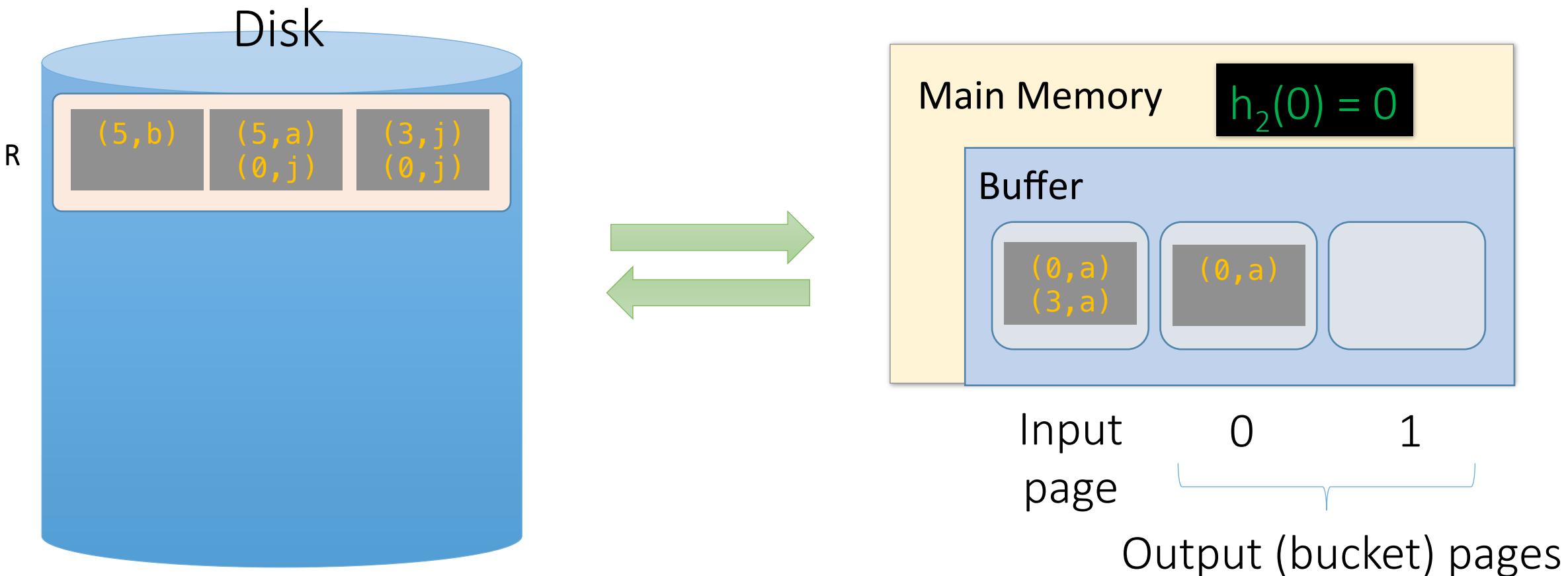
1. We read pages from R into the “input” page of the buffer...



# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

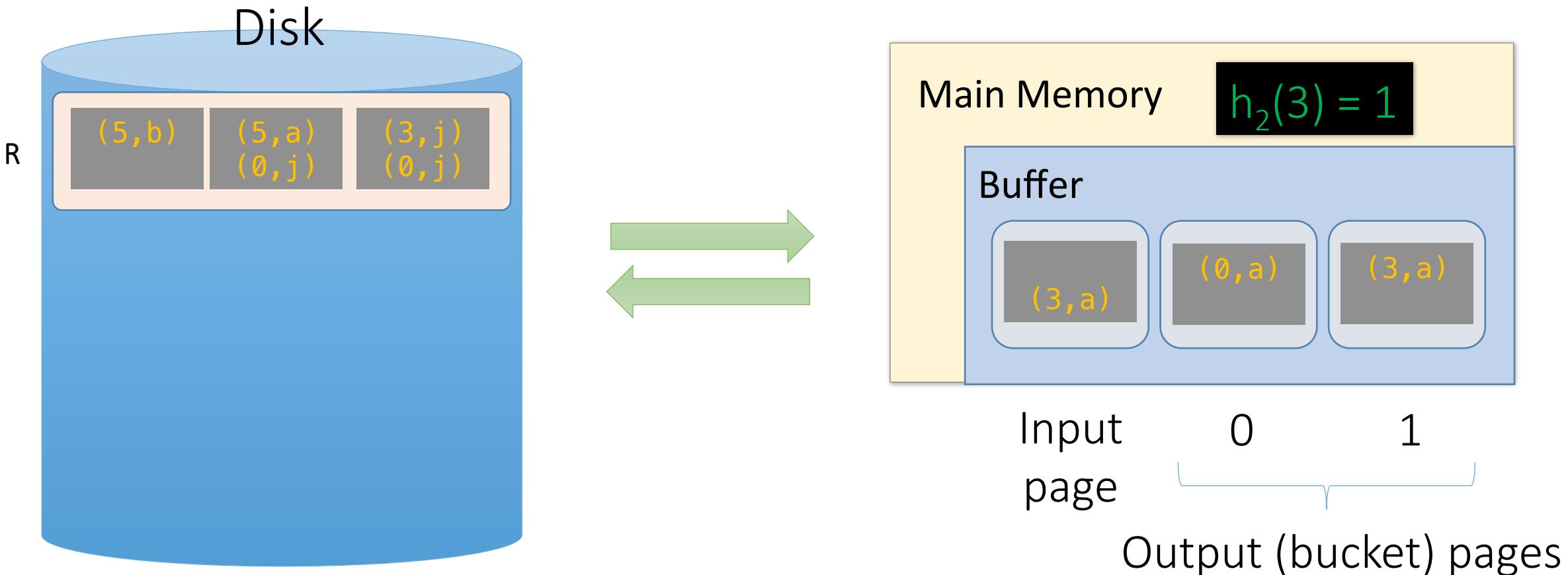
2. Then we use **hash function  $h_2$**  to sort into the buckets, which each have one page in the buffer



# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

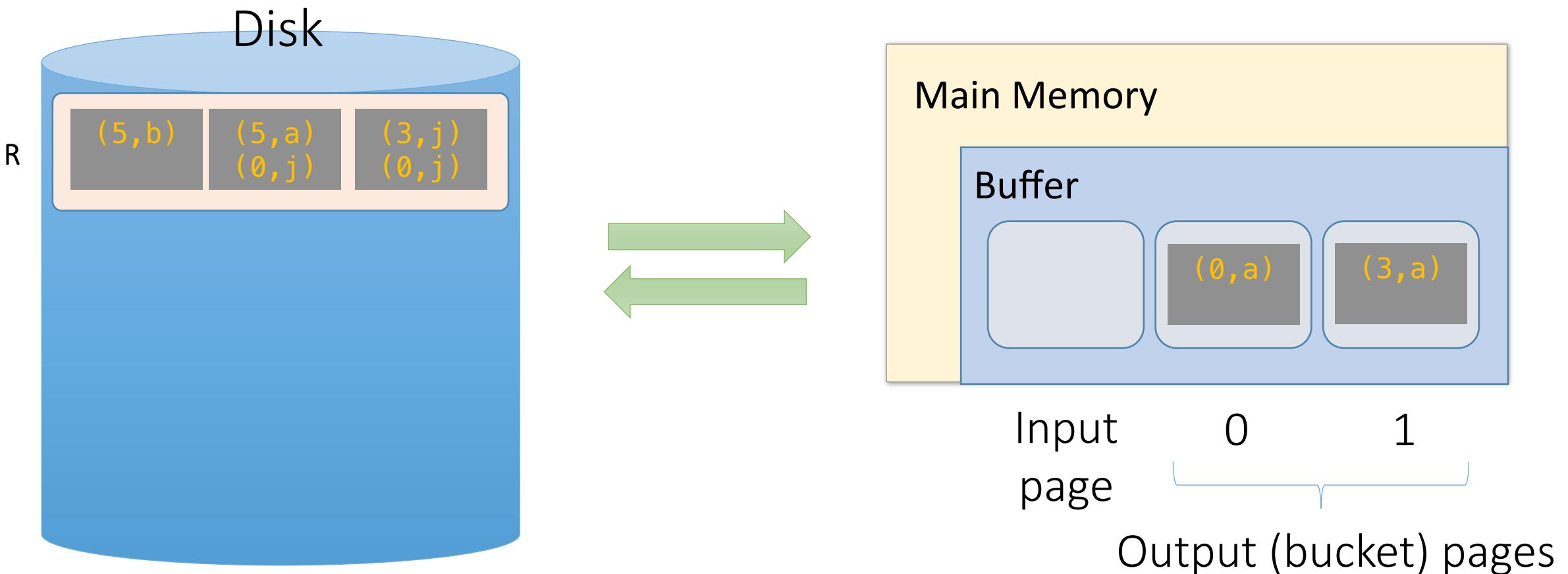
2. Then we use **hash function  $h_2$**  to sort into the buckets, which each have one page in the buffer



# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

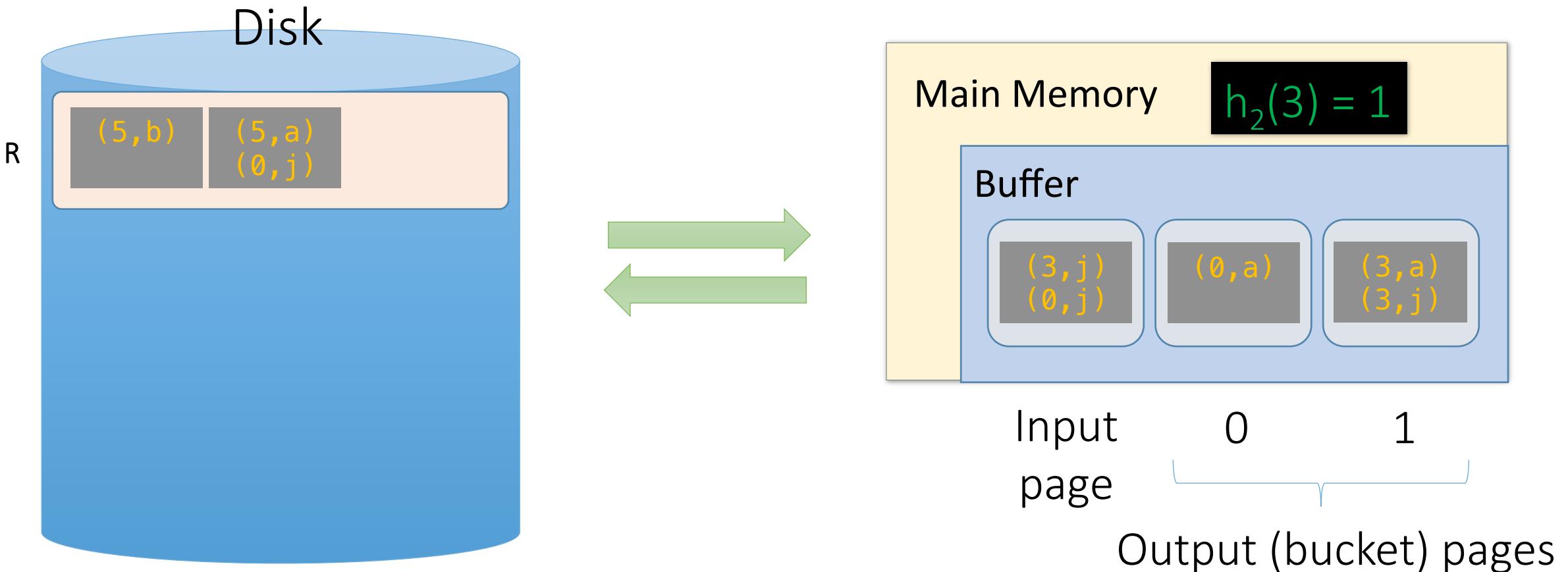
3. We repeat until the buffer bucket pages are full...



# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

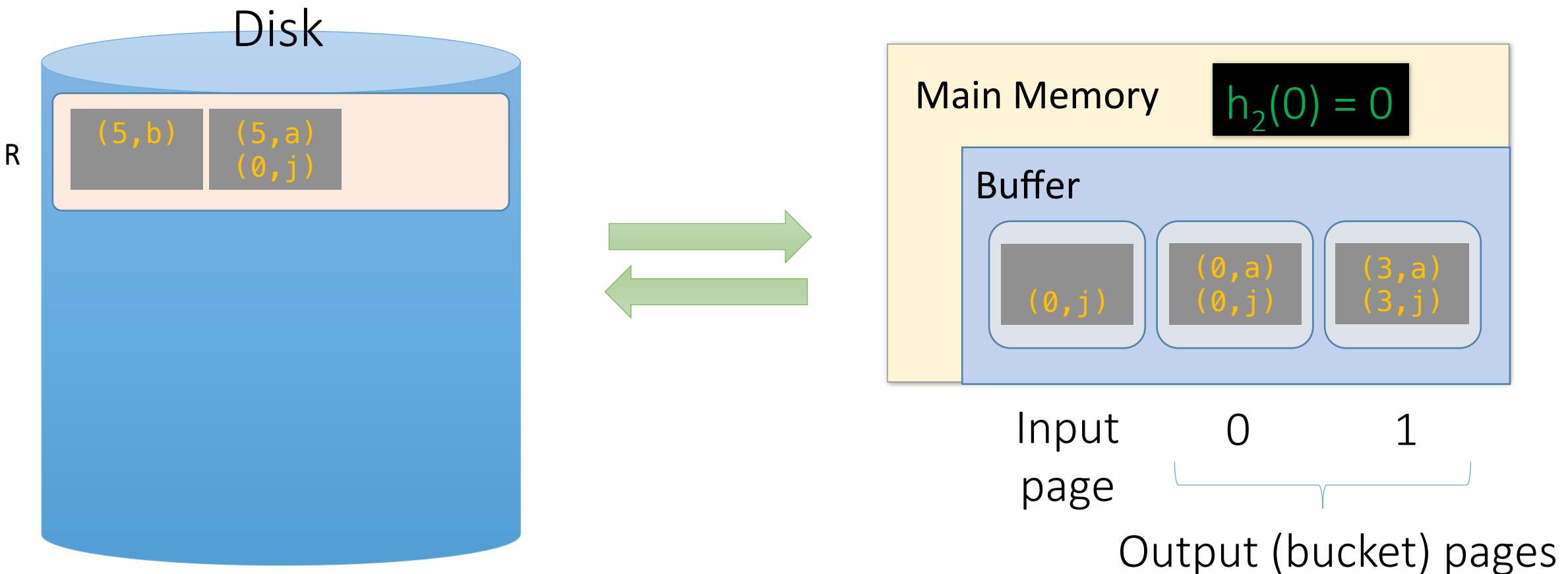
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# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

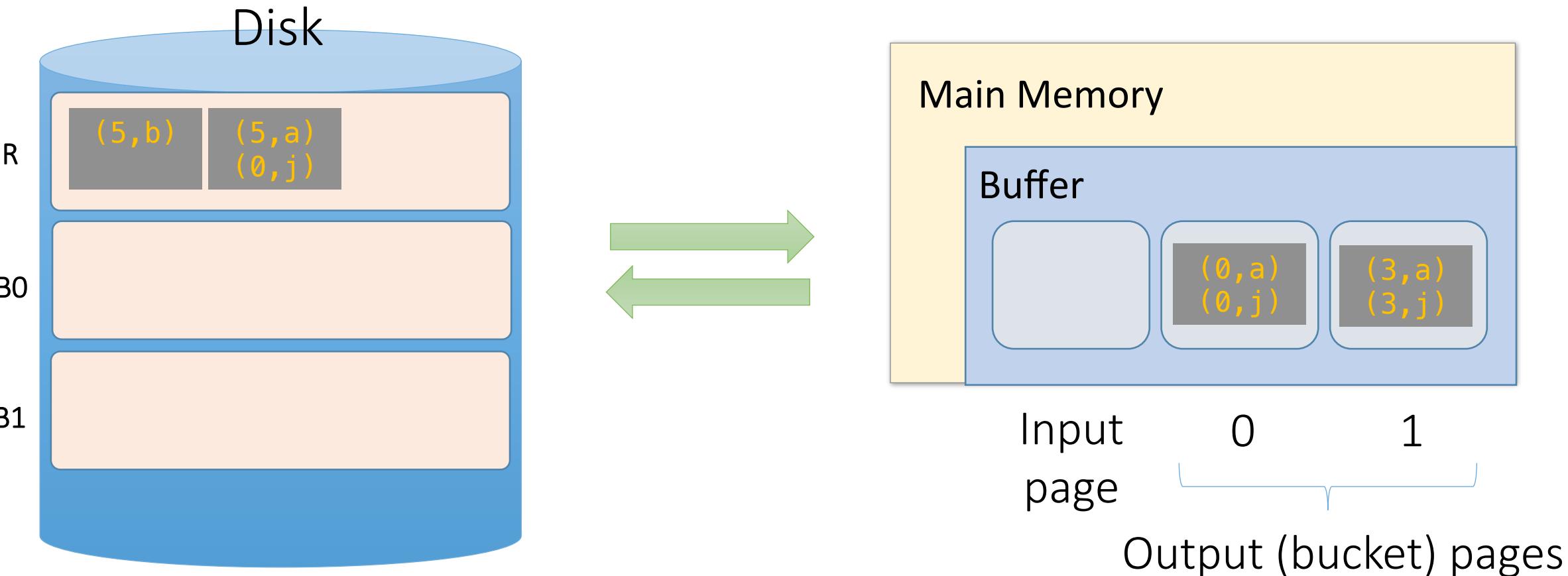
3. We repeat until the buffer bucket pages are full...



# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

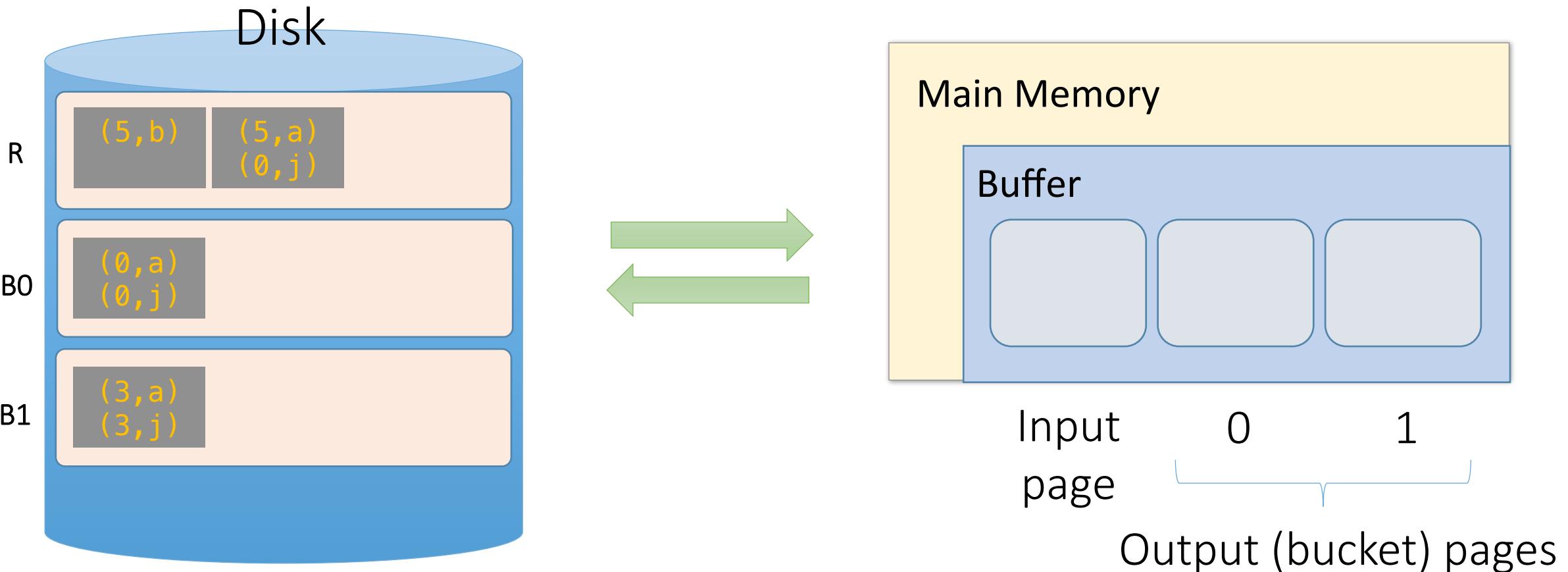
3. We repeat until the buffer bucket pages are full... then flush to disk



# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

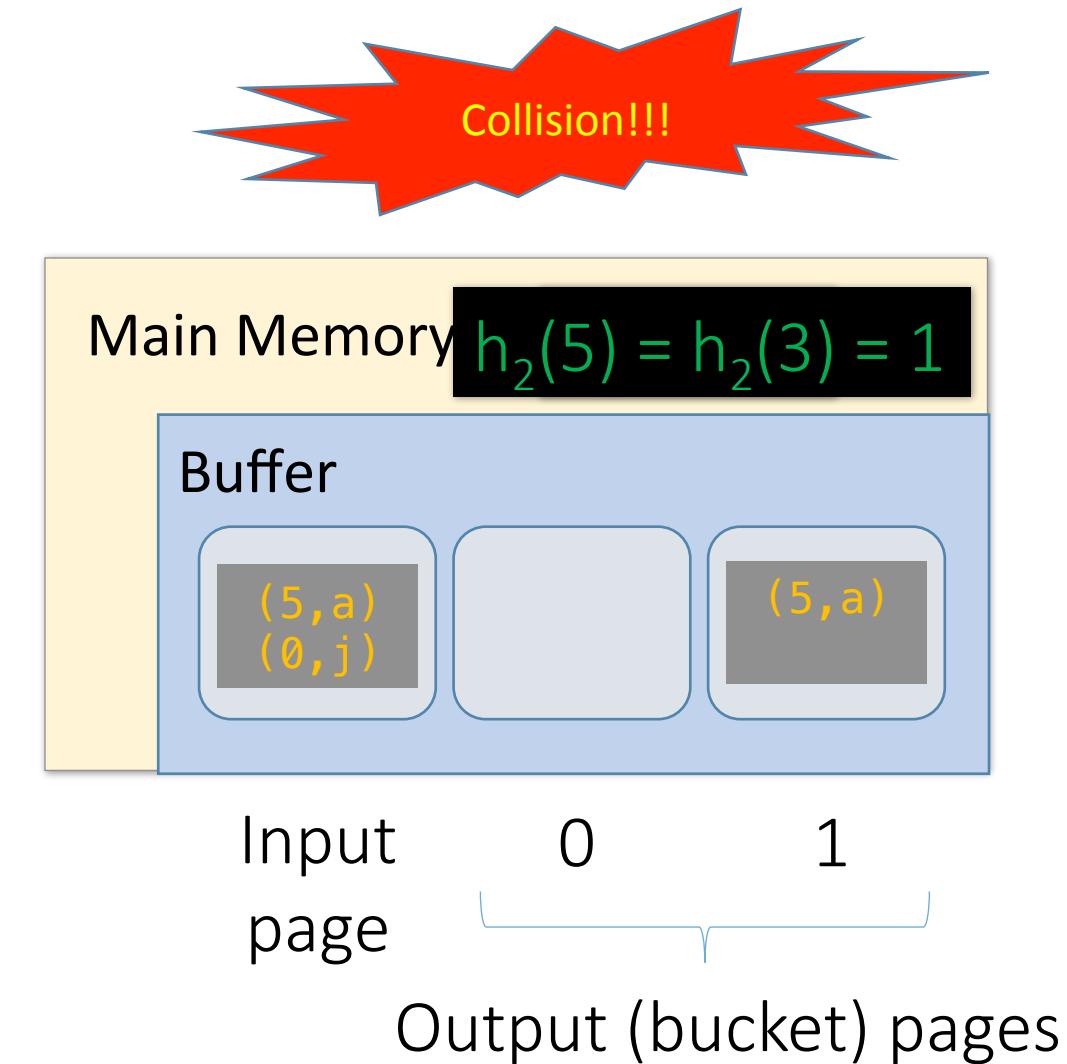
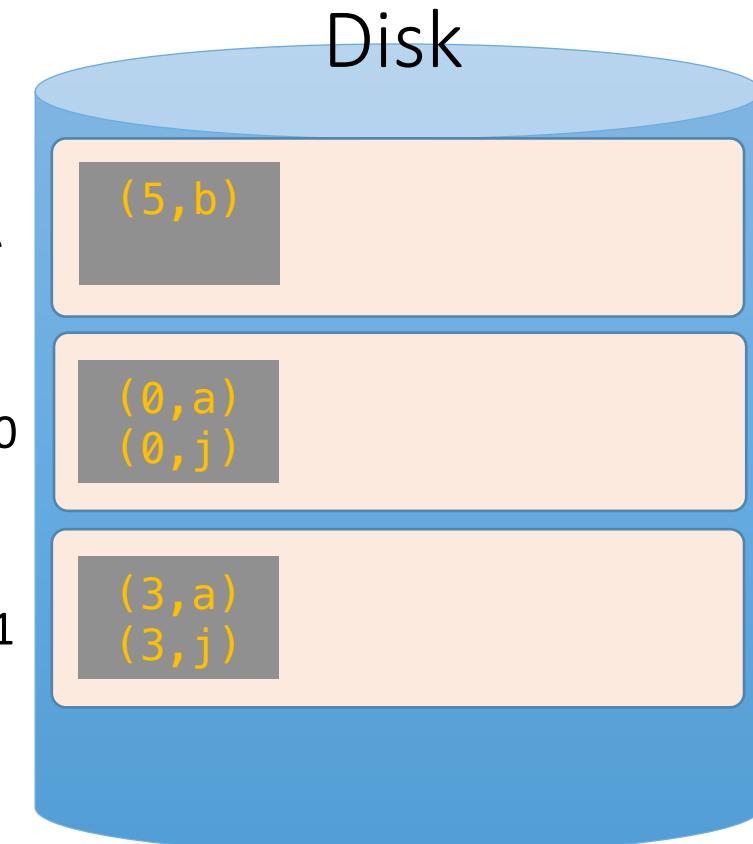
3. We repeat until the buffer bucket pages are full... then flush to disk



# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

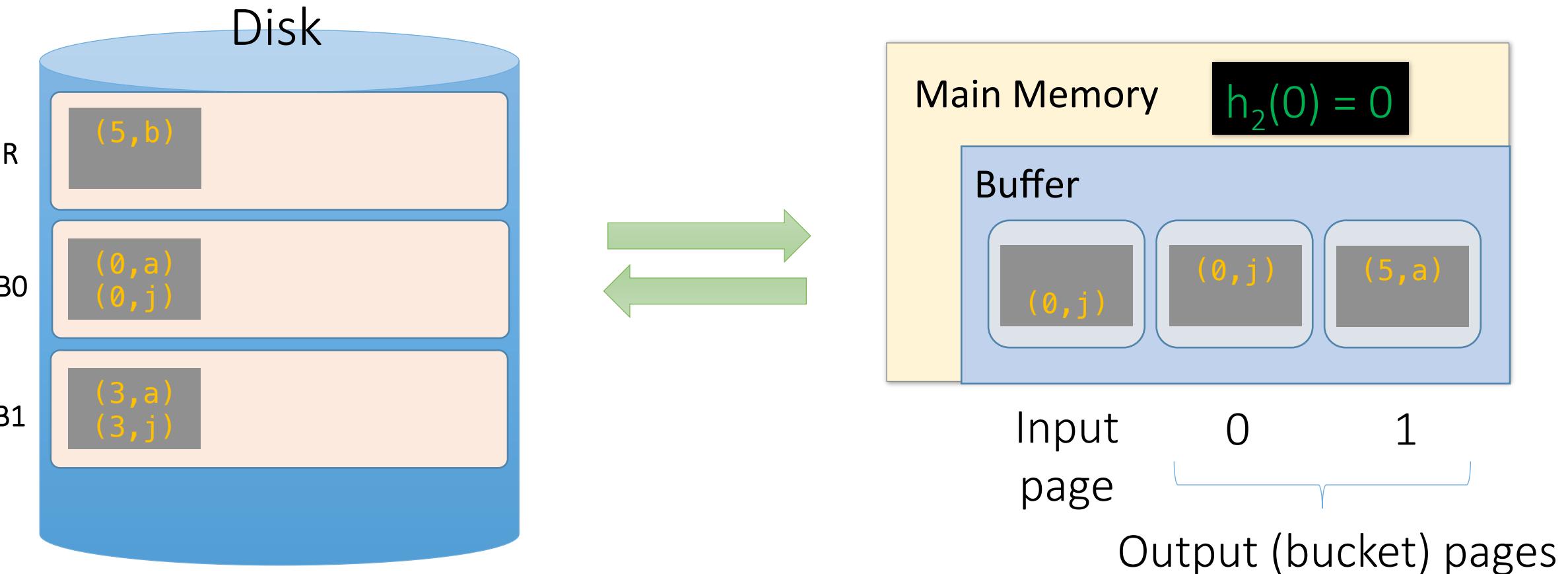
Note that collisions can occur!



# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

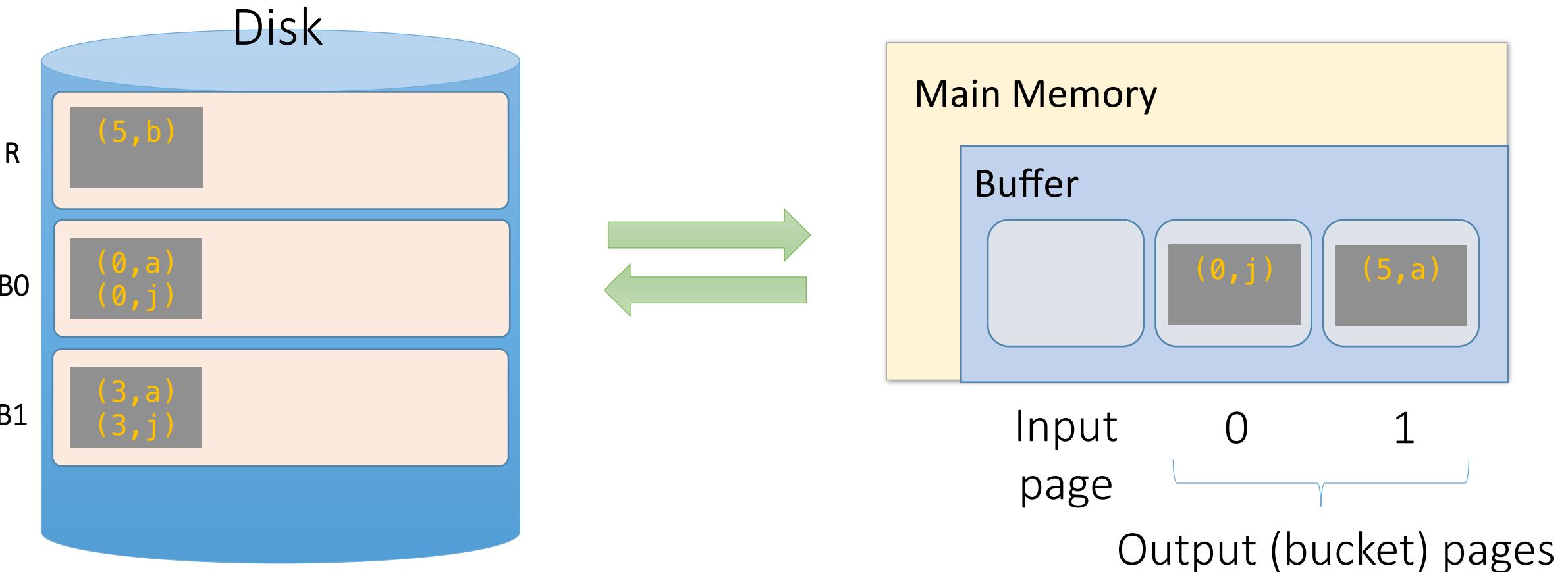
Finish this pass...



# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

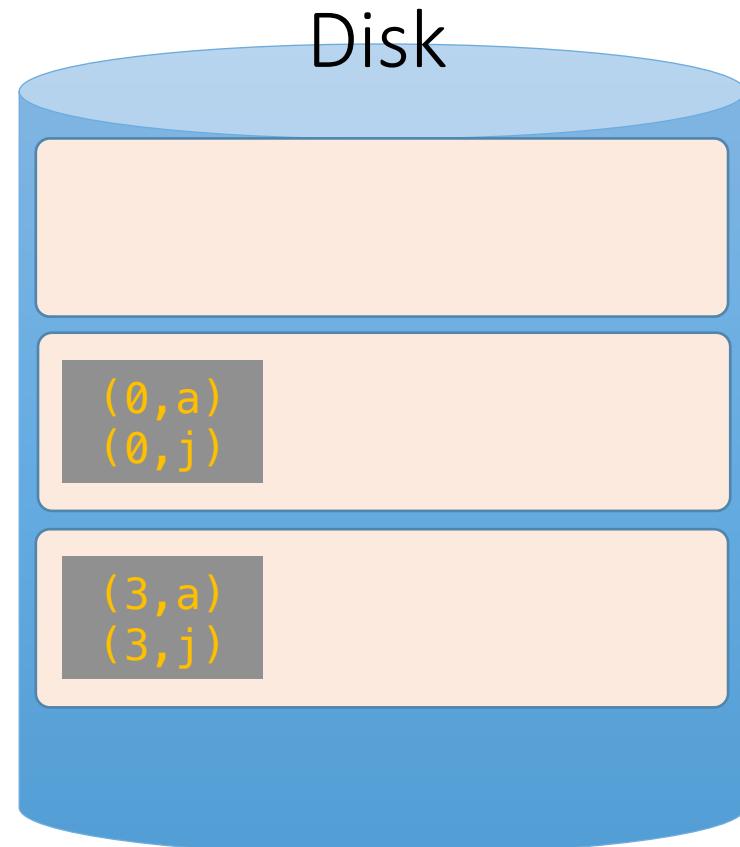
Finish this pass...



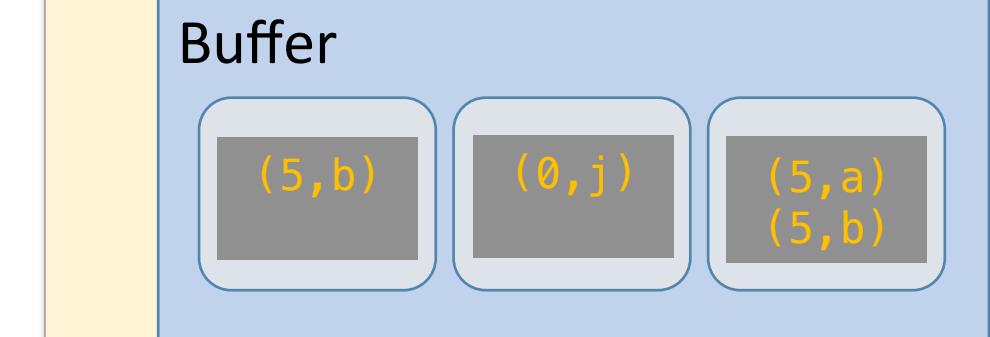
# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages

Finish this pass...



Main Memory  $h_2(5) = h_2(3) = 1$



Input page

0      1

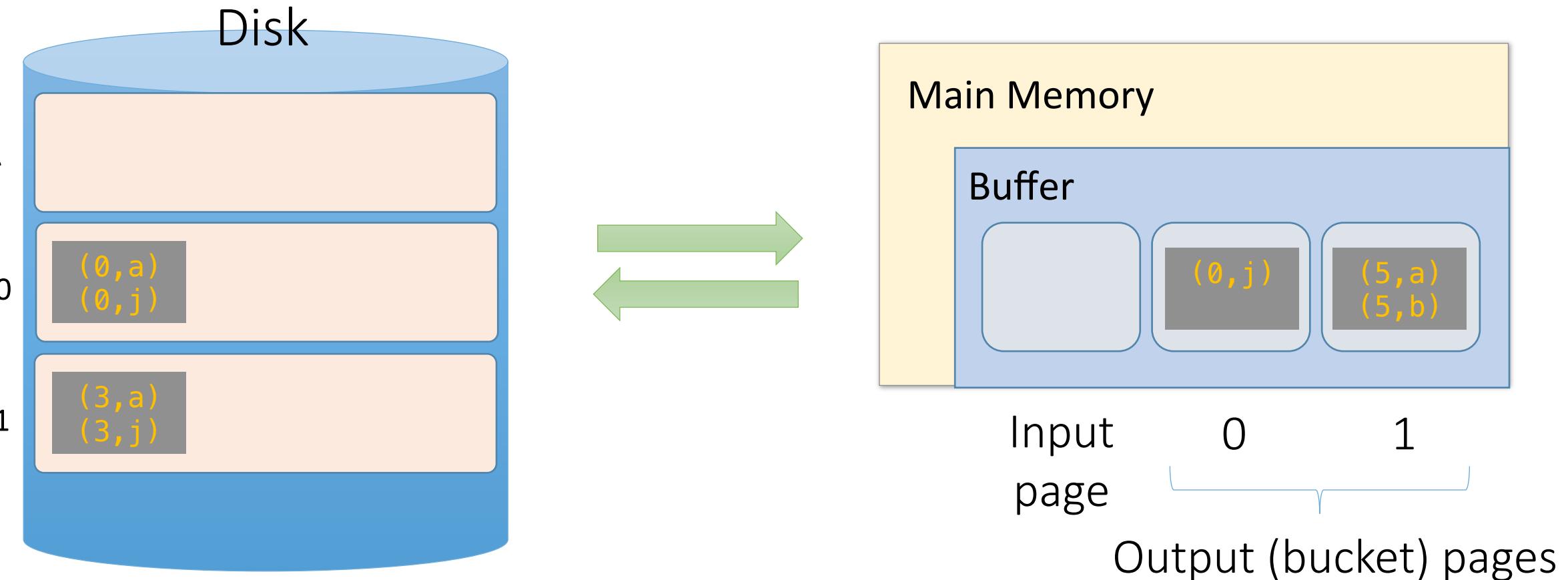
Output (bucket) pages



# Hash Join Phase 1: Partitioning

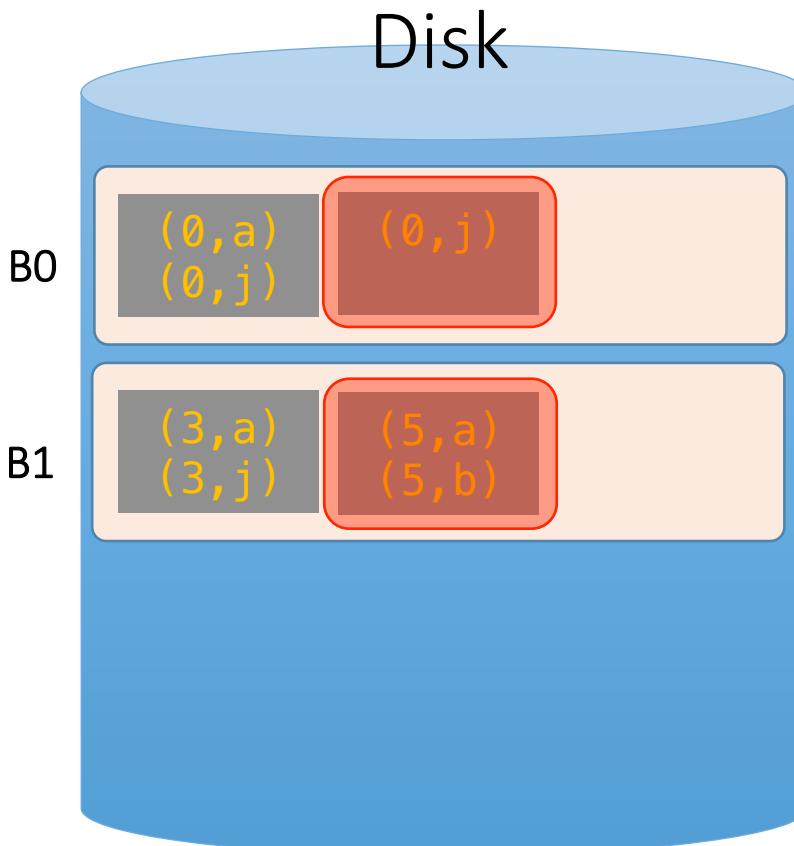
Given  $B+1 = 3$  buffer pages

Finish this pass...



# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages



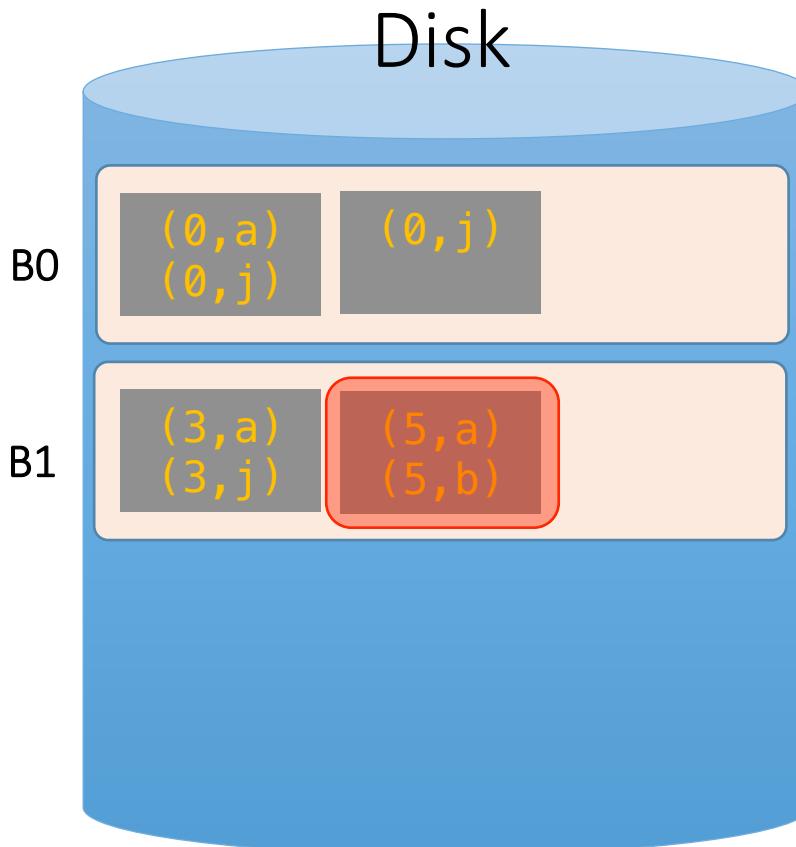
We wanted buckets of size  $B-1 = 1$ ...  
*however we got larger ones due to:*

(1) Duplicate join keys

(2) Hash collisions

# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages



To take care of larger buckets caused by (2) hash collisions, we can just do another pass!

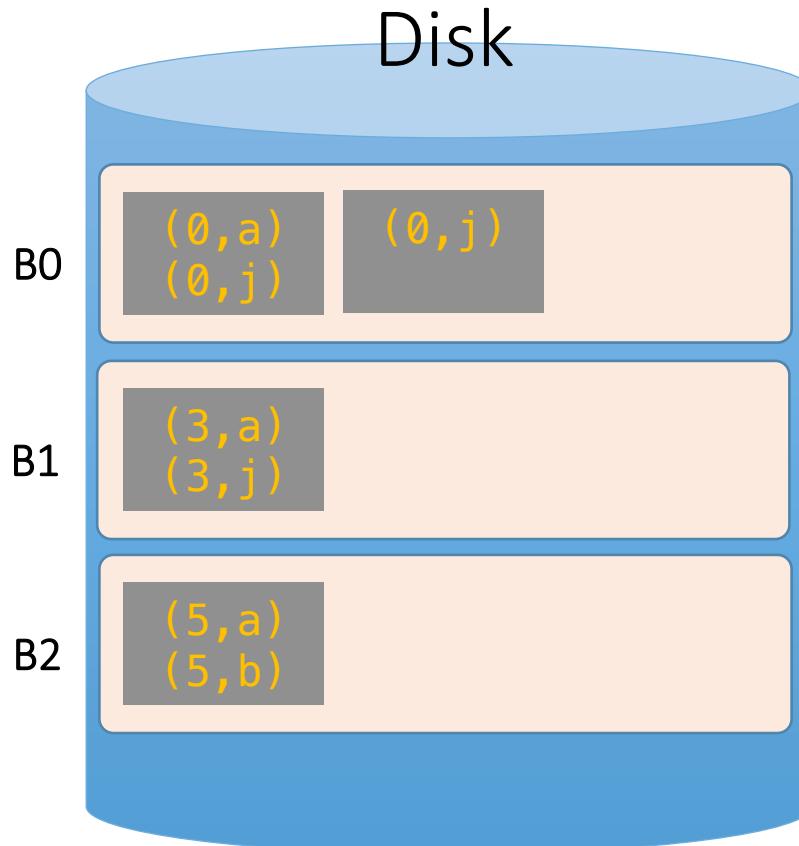
What hash function should we use?

Do another pass with a different hash function,  $h'_2$ , ideally such that:

$$h'_2(3) \neq h'_2(5)$$

# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages



To take care of larger buckets caused by (2) hash collisions, we can just do another pass!

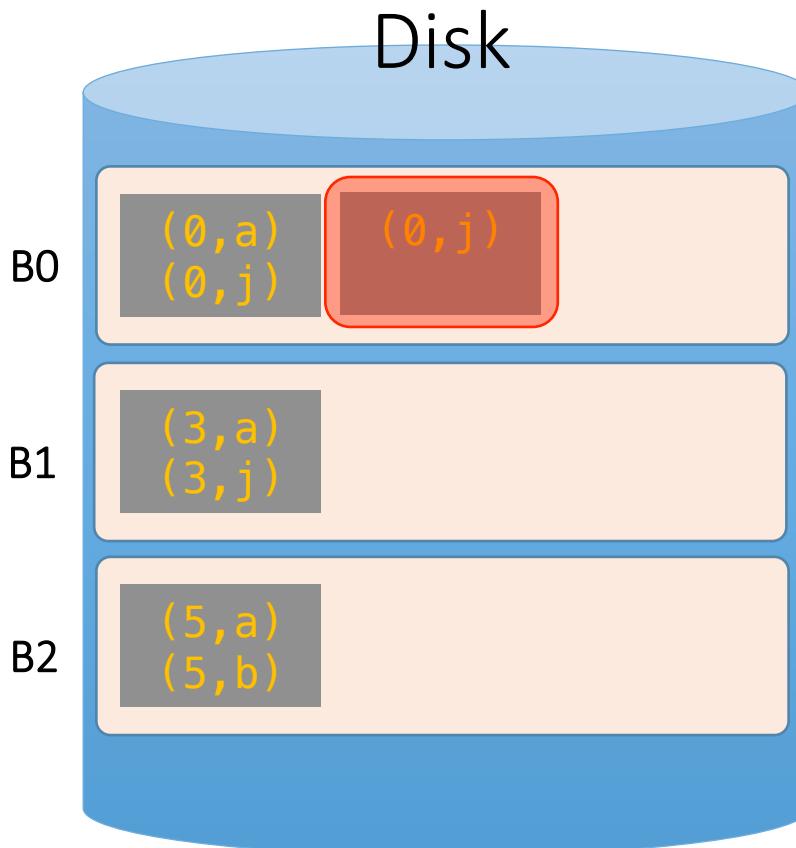
What hash function should we use?

Do another pass with a different hash function,  $h'_2$ , ideally such that:

$$h'_2(3) \neq h'_2(5)$$

# Hash Join Phase 1: Partitioning

Given  $B+1 = 3$  buffer pages



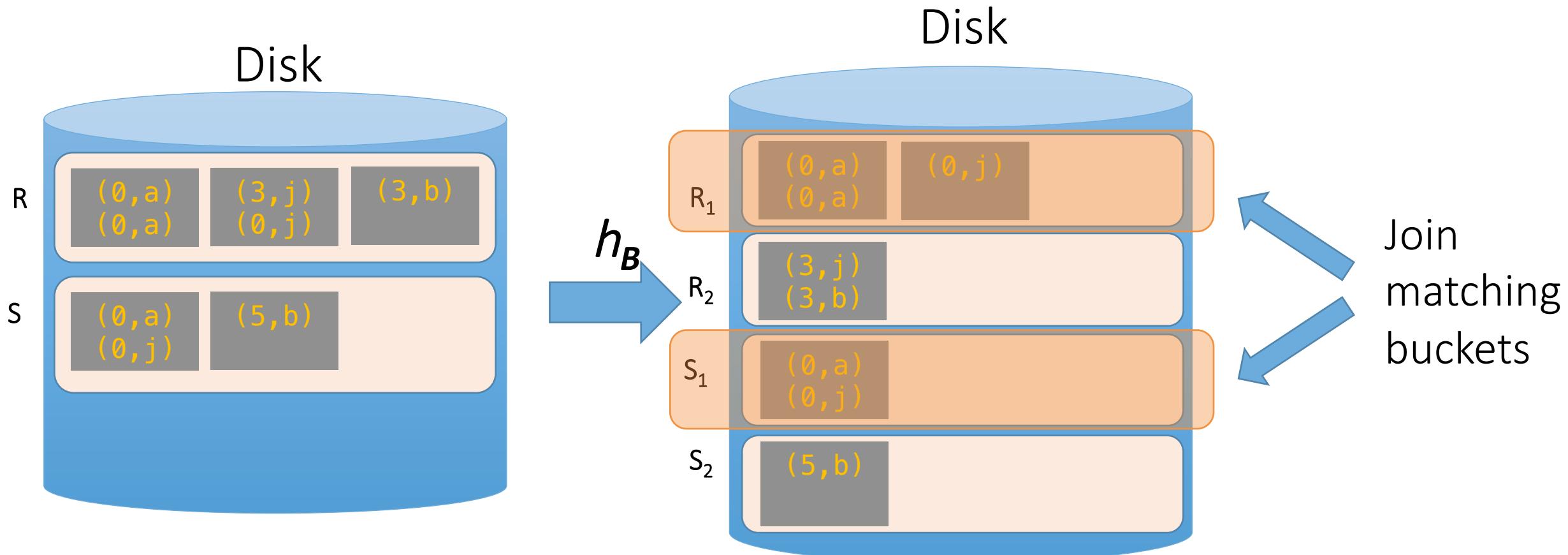
What about duplicate join keys?  
Unfortunately this is a problem... but  
usually not a huge one.

We call this unevenness  
in the bucket size skew

Now that we have partitioned R and S...

# Hash Join Phase 2: Matching

- Now, we just join pairs of buckets from R and S that have the same hash value to complete the join!



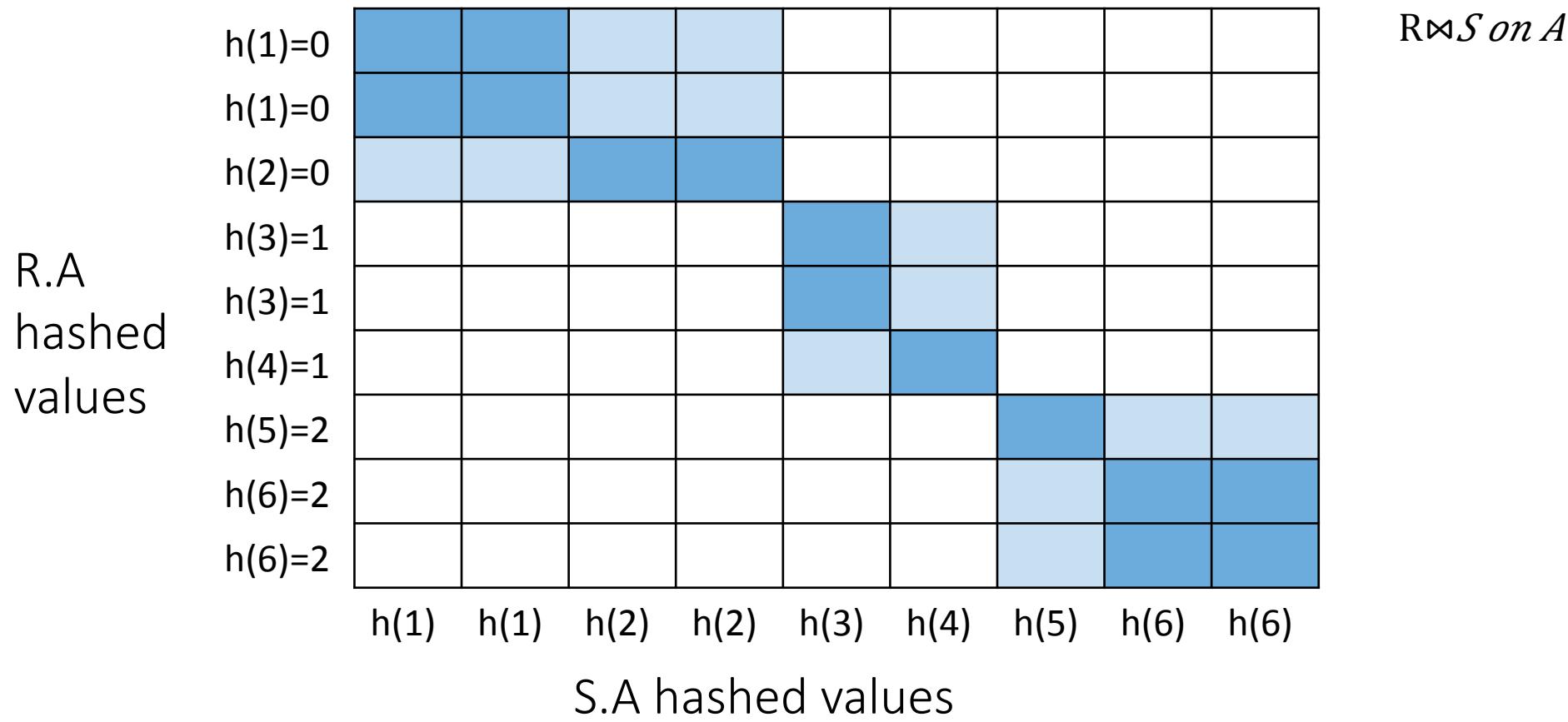
# Hash Join Phase 2: Matching

- Note that since  $x = y \rightarrow h(x) = h(y)$ , we only need to consider pairs of buckets (one from R, one from S) that have the same hash function value
- If our buckets are  $\sim B-1$  pages, can join each such pair using BNLJ *in linear time*; recall (with  $P(R) = B-1$ ):

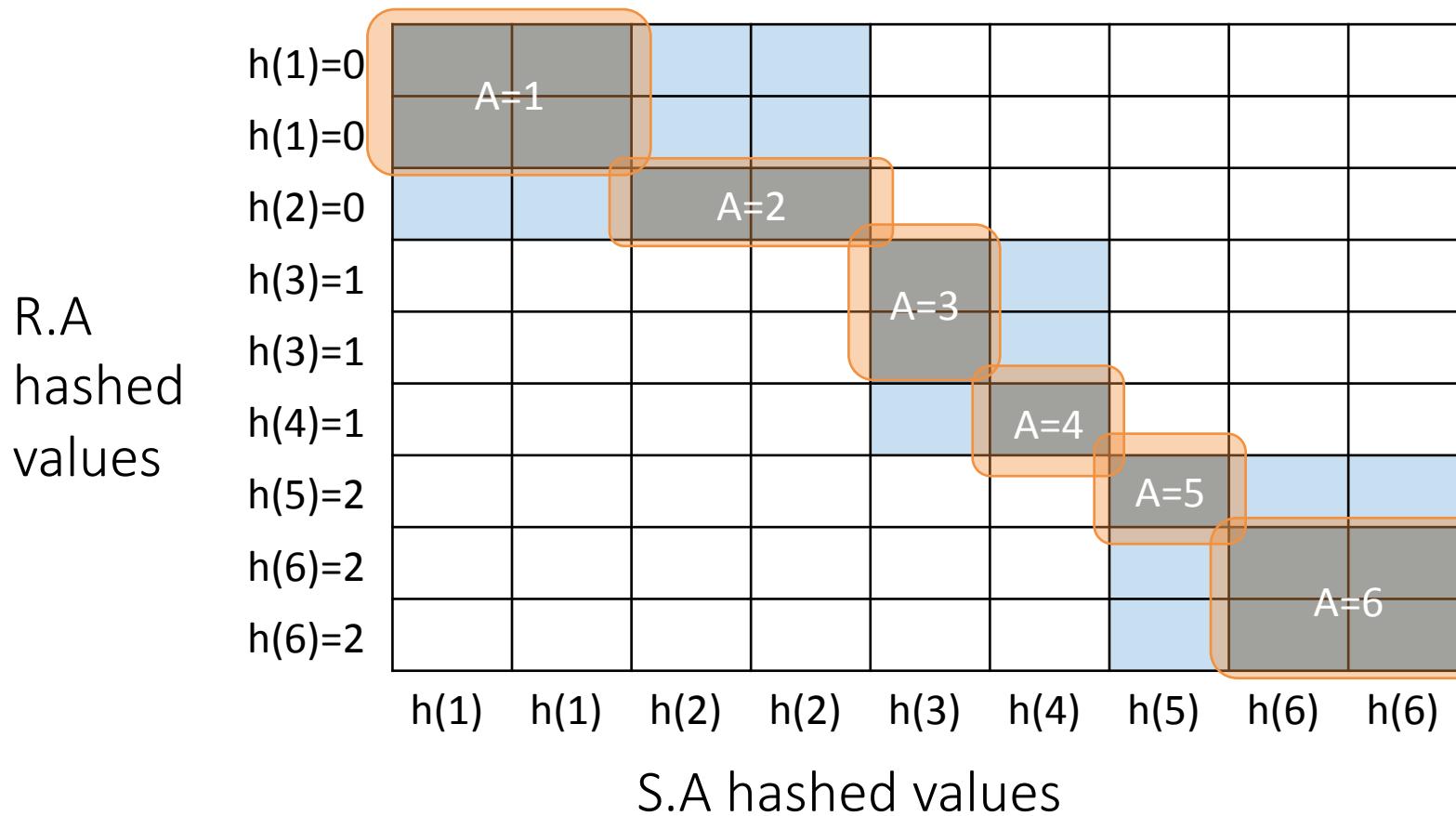
BNLJ Cost:  $P(R) + P(R)P(S)/B-1 = P(R) + (B-1)P(S)/B-1 = P(R) + P(S)$

Joining the pairs of buckets is linear!  
(As long as smaller bucket  $\leq B-1$  pages)

# Hash Join Phase 2: Matching



# Hash Join Phase 2: Matching

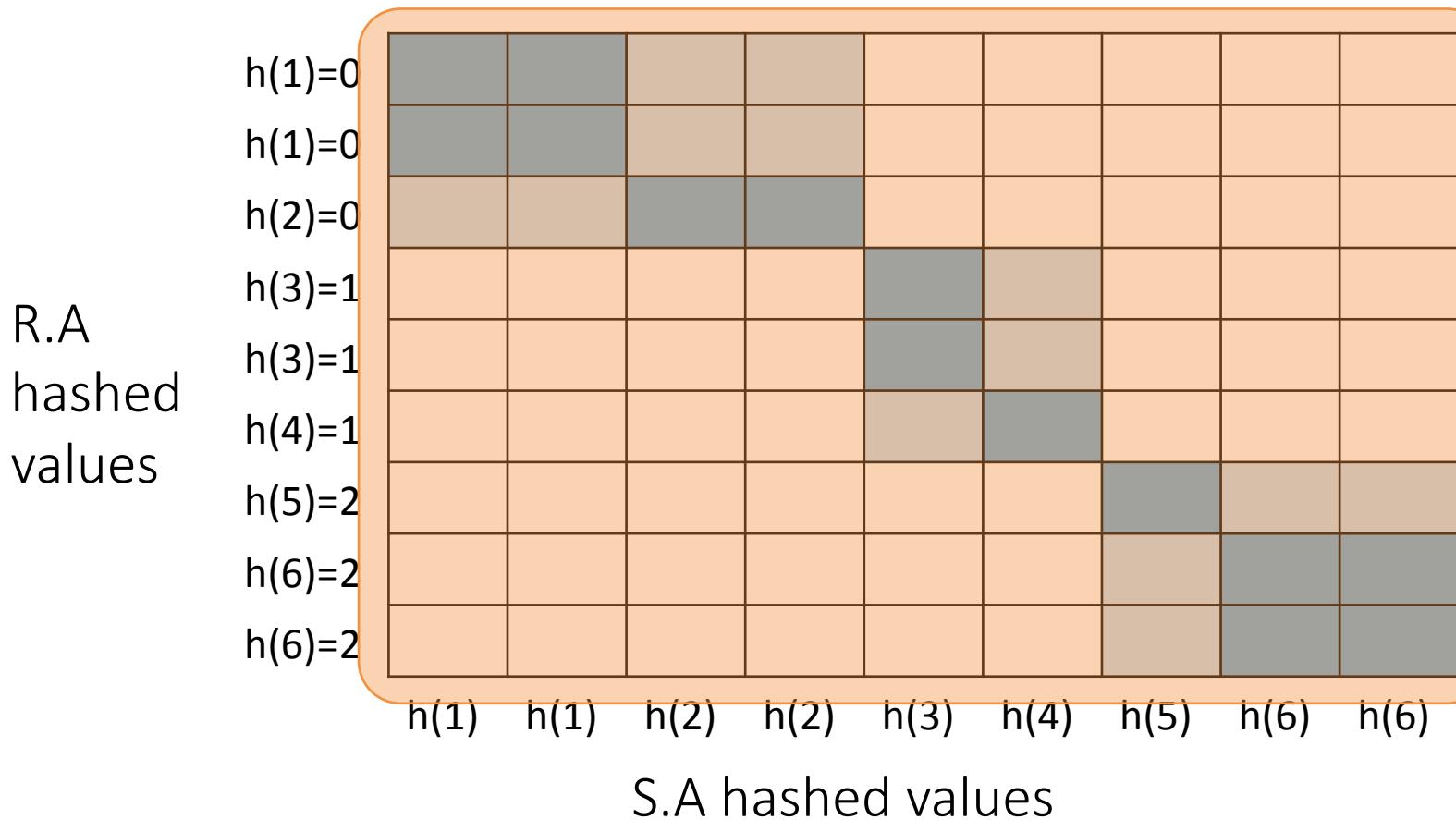


$R \bowtie S \text{ on } A$

To perform the join, we ideally just need to explore the dark blue regions

= the tuples with same values of the join key A

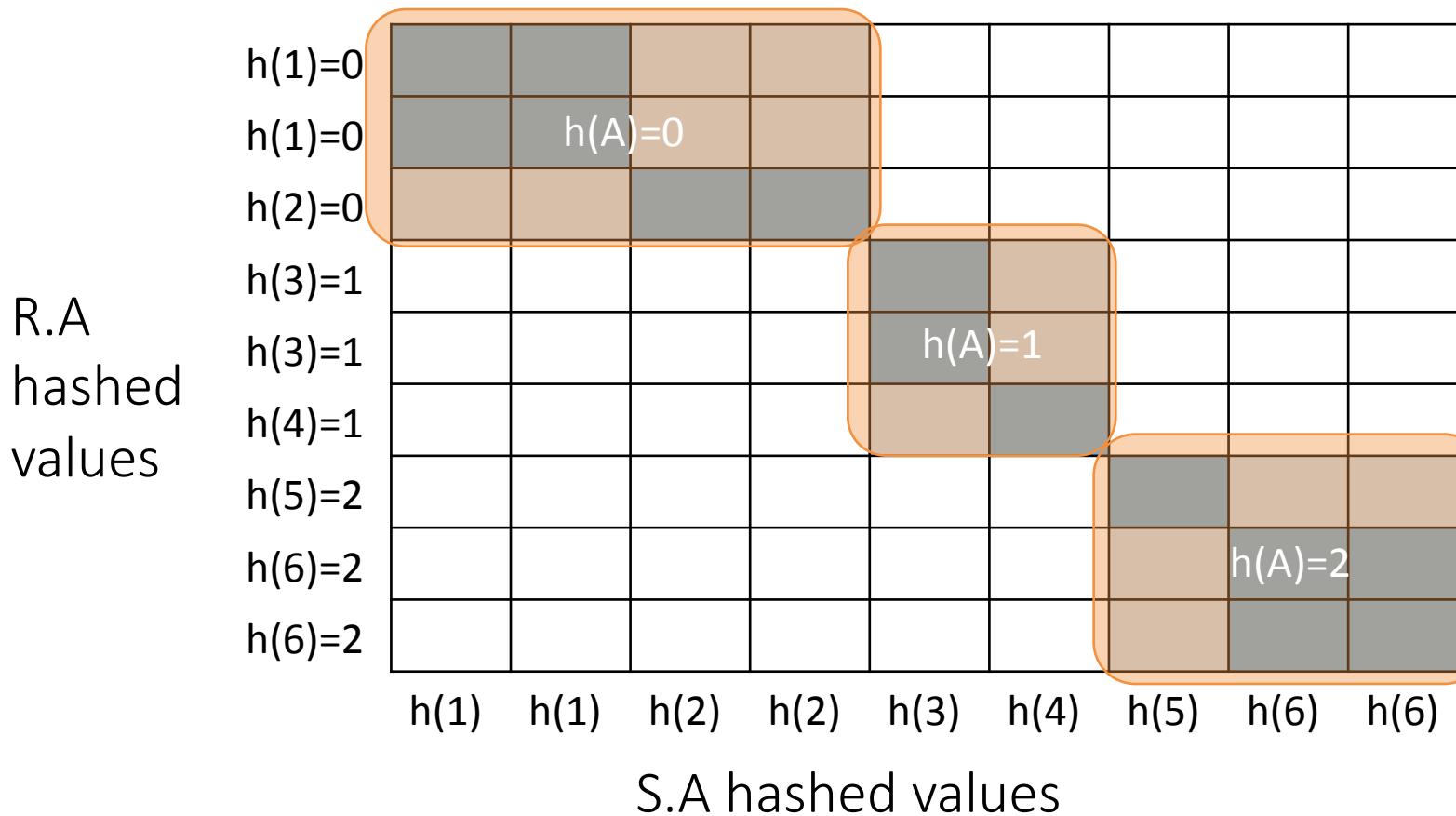
# Hash Join Phase 2: Matching



$R \bowtie S \text{ on } A$

With a join algorithm like BNLJ that doesn't take advantage of equijoin structure, we'd have to explore this ***whole grid!***

# Hash Join Phase 2: Matching



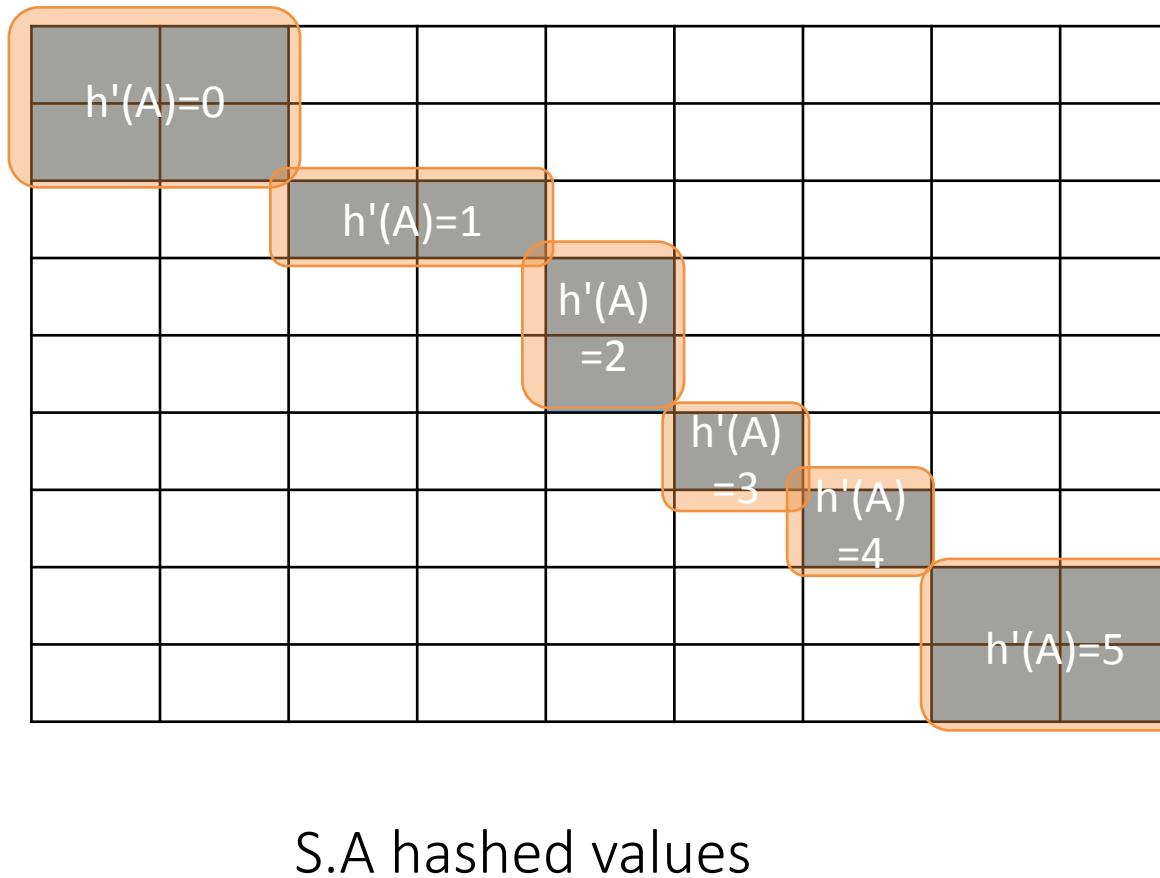
$R \bowtie S \text{ on } A$

With HJ, we only explore the *blue* regions  
= the tuples with same values of  $h(A)$ !

We can apply BNLJ to each of these regions

# Hash Join Phase 2: Matching

R.A  
hashed  
values



$R \bowtie S \text{ on } A$

An alternative to applying BNLJ:

We could also hash again, and keep doing passes in memory to reduce further!

# How much memory do we need for HJ?

- Given  $B+1$  buffer pages + WLOG: Assume  $P(R) \leq P(S)$
- Suppose (reasonably) that we can partition into  $B$  buckets in 2 passes:
  - For  $R$ , we get  $B$  buckets of size  $\sim P(R)/B$
  - To join these buckets in linear time, we need these buckets to fit in  $B-1$  pages, so we have:

$$B-1 \geq P(R)/B \Rightarrow \sim B^2 \geq P(R)$$

Quadratic relationship  
between *smaller*  
*relation's* size & memory!



# Hash Join Summary

- *Given enough buffer pages as on previous slide...*
- **Partitioning** requires reading + writing each page of R,S
  - $\rightarrow 2(P(R)+P(S))$  IOs
- **Matching** (with BNLJ) requires reading each page of R,S
  - $\rightarrow P(R) + P(S)$  IOs
- **Writing out results** could be as bad as  $P(R)*P(S)$ ... but probably closer to  $P(R)+P(S)$

HJ takes  $\sim 3(P(R)+P(S)) + OUT$  IOs!

### 3. The Cage Match

# Sort-Merge v. Hash Join



- ***Given enough memory***, both SMJ and HJ have performance:

$$\sim 3(P(R) + P(S)) + OUT$$



- ***"Enough" memory*** =

- SMJ:  $B^2 > \max\{P(R), P(S)\}$

- HJ:  $B^2 > \min\{P(R), P(S)\}$

Hash Join superior if relation sizes *differ greatly*. Why?

# Further Comparisons of Hash and Sort Joins

- Hash Joins are highly parallelizable.



- Sort-Merge less sensitive to data skew and result is sorted



# Summary

- Saw IO-aware join algorithms
  - Massive difference
- Memory sizes key in hash versus sort join
  - Hash Join = Little dog (depends on smaller relation)
- Skew is also a major factor