实验进度:完成了所有必需的内容,并且可以玩打字小游戏和仙剑

1. 详细描述从测试用例中的 int \$0x80 开始一直到 HIT_GOOD_TRAP 为止的详细的系统 行为(完整描述控 制的转移过程,即相关函数的调用和关键参数传递过程),可以通过 文字或画图的方式来完成;

```
5 Disassembly of section .text:
 7 08048094 <start>:
 8 8048094: e9 00 00 00 00
                                           jmp
                                                  8048099 <main>
 9
10 08048099 <main>:
11 8048099: 55
                                           push
                                                  %ebp
12 804809a:
                89 e5
                                           mov
                                                  %esp,%ebp
13 804809c:
                                                  80480c9 < x86.get_pc_thunk.ax>
               e8 28 00 00 00
                                          call
14 80480a1:
                05 5f 1f 00 00
                                          add
                                                  $0x1f5f,%eax
15 80480a6:
                b8 04 00 00 00
                                                  $0x4,%eax
                                          mov
16 80480ab: bb 01 00 00 00 17 80480b0: b9 d0 80 04 08 18 80480b5: ba 0e 00 00 00 19 80480ba: cd 80
                                                  $0x1,%ebx
                                          mov
                                                  $0x80480d0,%ecx
                                          mov
                                                  $0xe,%edx
                                          mov
                                                  $0x80
                                          int
20 80480bc:
               b8 00 00 00 00
                                                  $0x0,%eax
                                          mov
21 80480c1: 82 b8 00 00 00 00 5d
                                                  $0x5d,0x0(%eax)
                                         cmpb
22 80480c8:
                c3
                                           ret
23
24 080480c9 <__x86.get_pc_thunk.ax>:
25 80480c9:
                8b 04 24
                                           mov
                                                   (%esp),%eax
26 80480cc:
                с3
                                           ret
```

这是 hello-inline 的测试用例的反汇编代码,可以看出在 0x80480ba 的地方进入了一个系统中断 int 0x80, 这个时候调用 int 指令进行响应。

```
4 make instr func(int b)
 5 {
       /*uint32_t imm8 = (uint32_t)instr_fetch(eip + 1, 1);
6
7
      uint32 t addr = cpu.idtr.base + imm8 * 8;
8
      GateDesc gatedesc;
9
      gatedesc.val[0] = laddr_read(addr, 4);
      //printf("des0 == %x\n", desc.val[0]);
10
11
      //printf("des1 == %x\n", desc.val[1]);
12
      gatedesc.val[1] = laddr_read(addr + 4, 4);
13
14
      uint32_t sreg = (gatedesc.selector >> 3) & 0x1fff;
15
      uint32_t vaddr = (gatedesc.offset_15_0 & 0xffff) + ((gatedesc.offset_31_16 <<</pre>
16
      uint32 t laddr = segment_translate(vaddr, sreg);
17
      uint32 t paddr = page translate(laddr);
18
19
      cpu.eip = paddr;
20
      printf("cpu.eip == %x\n", cpu.eip);
21
22
      return 0;*/
23
      printf("eip ====== %x\n",cpu.eip);
24 //
      uint8_t imm8 = (uint8_t)instr_fetch(eip + 1, 1);
25
      printf("imm8 == %x\n",imm8);
26 //
      raise sw intr(imm8);
27
28
      return 0;
29 }
```

在 int 中, 用 instr_fetch 把中断号 0x80 传给函数 raise_sw_instr 函数。

```
60 void raise_sw_intr(uint8_t intr_no) {
61    // return address is the next instruction
62    //printf("insteip:%x\n", cpu.eip);
63    cpu.eip += 2;
64    raise_intr(intr_no);
65 }
```

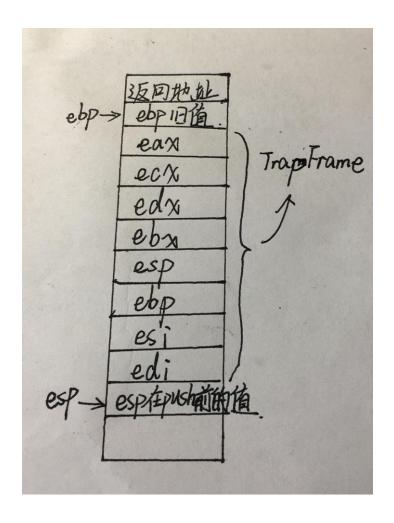
这个时候看到了 eip+=2, 因为 int 0x80 是 2 个字节的指令码, 所以这时候的 eip 就指向了在 hello-inline 用例里 int 0x80 的下一条指令。这时候调用 raise_intr。raise_intr 里面保存了 eflags cs eip 的值, 然后进入内核态, 由系统处理中断。

然后进入 vecsys 这个调用,调用响应的处理函数 asm_do_irg。

```
135 c00300ca <asm do irq>:
136 c00300ca:
                60
                                          pusha
137 c00300cb:
                54
                                          push
                                                 %esp
138 c00300cc:
                e8 fc 1b 00 00
                                         call
                                                 c0031ccd <irg handle>
139 c00300d1:
                83 c4 04
                                         add
                                                 $0x4,%esp
140 c00300d4:
                61
                                         popa
141 c00300d5:
                83 c4 08
                                         add
                                                 $0x8,%esp
142 c00300d8:
                cf
                                         iret
143
```

pusha 保存之前寄存器的值,然后调用 irq_handle 进行具体的处理,最后 popa 还原之前的寄存器的值,最后 iret 将 eip cs eflags 的现场还原出来,这时候将控制权转移进用户态,eip 还原出来的值就是在 hello-inline 用例里 int 0x80 的下一条指令,这时用户进程继续执行,最后 HIT_GOOD_TRAP。

2、在描述过程中,回答 kernel/src/irq/do_irq.S 中的 push %esp 起什么作用,画出在 call irq_handle 之前,系统栈的内容和 esp 的位置,指出 TrapFrame 对应系统栈的哪一段内容。



3、详细描述 NEMU 和 Kernel 响应时钟中断的过程和先前的系统调用过程不同之处在哪里?相同的地方 又在哪里?可以通过文字或画图的方式来完成

相同的地方就是系统在发生中断后,都会保存现场,然后将控制权转移到内核态去处理,处理完毕后恢复现场,将控制权转移给用户态继续执行进程。

不同的地方就是,先前的系统调用是通过 int 指令来实现的,而时钟中断是每次执行完一条指令后调用 do_intr()函数查看并处理中断事件,属于系统实时监测中断的产生再做相应的处理。

4. 注册监听键盘事件是怎么完成的?

```
48
49
       // loop to get keyboard events
       SDL_Event e;
50
       uint64_t jiffy = 0;
51
52 #endif
53
      // initializing the display finish
54
       initialized = true;
       //For tracking if we want to quit
55
       while(initialized) {
57 #ifdef HAS DEVICE TIMER
58
           timer intr();
59 #endif
60 #ifdef HAS_DEVICE_VGA
          // update the screen
61
           if(jiffy % (SDL_HZ / VGA_HZ) == 0)
62
               update_screen();
63
64 #endif
65 #if defined(HAS_DEVICE_VGA) || defined(HAS_DEVICE_KEYBOARD)
66  //Read all the events that occured
           while(SDL PollEvent(&e)){
67
68 #if defined(HAS DEVICE VGA) || defined(HAS DEVICE KEYBOARD)
69
               //If user closes the window
70
               if (e.type == SDL QUIT){
                   nemu state = NEMU STOP;
71
72
73 #endif
74 #ifdef HAS DEVICE KEYBOARD
75
               else if (e.type == SDL KEYDOWN){
                    keyboard_down(e.key.keysym.sym);
76
77
               } else if (e.type == SDL KEYUP) {
                   keyboard up(e.key.keysym.sym);
78
79
80 #endif
81
           jiffy++;
82
83 #endif
           SDL_Delay(1000 / SDL_HZ); // 100Hz
84
85
86 #ifdef HAS DEVICE VGA
      close_vga();
87
```

通过一个循环,每 0.01 秒系统都会发生一个信号给 nemu,每当信号来的时候,系统就会检测有没有事情发生,然后对于键盘事件,通过键盘的捕获 keyboard_up 和 keyboard_down,来监听键盘事件的产生。

5. 从键盘按下一个键到控制台输出对应的字符,系统的执行过程是什么?如果涉及与之前报告重复的内容,简单引用之前的内容即可。

```
x #Tucrnae __deAtce\TQSDa_btc.u
4 static wint8 t scan code buf;
6 // called by the nemu sdl thread on detecting a key down event
7 void keyboard_down(uint32_t sym) {
      // put the scan code into the buffer
      scan code buf = sym2scancode[sym >> 8][sym & 0xff];
10
      // issue an iterrupt
11
      i8259 raise intr(KEYBOARD IRQ);
      // maybe the kernel will be interested and come to read on the data port
12
13 }
14
15 // called by the nemu sdl thread on detecting a key up event
16 void keyboard up(uint32 t sym) {
17
      // put the scan code into the buffer
      scan code buf = sym2scancode[sym >> 8][sym & 0xff] | 0x80;
19
      // issue an iterrupt
20
      i8259 raise intr(KEYBOARD IRQ);
      // maybe the kernel will be interested and come to read on the data port
21
22 }
23
24 // called when the kernel issues an 'in' instruction on the keyboard's data port
25 make pio handler(handler keyboard data) {
      if(!is write) {
26
27
          // only read allowed, and we do not consider race condition here
28
          write io port(port, len, scan code buf);
29
      }
30 }
```

利用前面的键盘事件监听功能去捕获键盘的扫描码, 然后调用 i8259_raise_intr()来触发中断, 然后去查表来调用相应的处理函数。

```
1 #----|---entry-----|-errorcode-|----id-----|---handler---|
2 .globl vec0; vec0: pushl $0; pushl
                                              $0; jmp asm do irq
3 .globl vec1;
                 vec1:
                          pushl $0;
                                    pushl
                                              $1; jmp asm_do_irq
4 .globl vec2;
                 vec2: pushl $0;
                                    pushl
                                              $2; jmp asm_do_irq
5 .globl vec3;
                 vec3: pushl $0;
                                    pushl
                                              $3; jmp asm_do_irq
6 .globl vec4;
                 vec4: pushl $0;
                                    pushl
                                             $4; jmp asm do irq
7 .globl vec5;
                 vec5: pushl $0;
                                    pushl
                                             $5; jmp asm do irq
8 .globl vec6;
                 vec6:
                          pushl $0;
                                             $6; jmp asm do irq
                                    pushl
                 vec7:
9 .globl vec7;
                          pushl $0;
                                              $7; jmp asm do irq
                                    pushl
10 .globl vec8;
                                              $8; jmp asm do irq
                 vec8:
                                     pushl
11 .globl vec9;
                          pushl $0;
                 vec9:
                                     pushl
                                             $9; jmp asm do irq
12 .globl vec10;
                 vec10:
                                     pushl
                                             $10; jmp asm do irq
13 .globl vec11;
                 vecl1:
                                     pushl
                                             $11; jmp asm do irq
14 .globl vec12;
                  vec12:
                                     pushl
                                             $12; jmp asm do irq
15 .globl vec13;
                                             $13; jmp asm do irq
                 vec13:
                                     pushl
16 .globl vec14;
                  vec14:
                                     pushl
                                             $14; jmp asm do irq
17
18 .globl vecsys; vecsys:
                          pushl $0; pushl $0x80; jmp asm do irq
19
20 .globl irq0;
                   irq0:
                          pushl $0;
                                    pushl $1000; jmp asm_do_irq
                          pushl $0;
                                    pushl $1001; jmp asm_do_irq
21 .globl irq1;
                  irq1:
22 .globl irq14;
                                    pushl $1014; jmp asm do irq
                 irq14:
                          pushl $0;
23 .globl irq_empty;
24
              irq_empty:
                          pushl $0; pushl $-1; jmp asm_do_irq
25
26 .globl asm_do_irq
27 .extern irq handle
```

然后调用 irq_handle 进行处理

```
32 void irq_handle(TrapFrame *tf) {
33
       int irq = tf->irq;
34
35
       if (irq < 0) {
36
           panic("Unhandled exception!");
37
       } else if (irq == 0x80) [
38
           do_syscall(tf);
39
       | else if (irq < 1000) {</pre>
40
           panic("Unexpected exception #%d at eip = %x", irq, tf->eip);
41
       } else if (irq >= 1000) {
           int irq_id = irq - 1000;
//Log("irq = %d", irq);
42
43
44
           //Log("irq_id = %d\n", irq_id);
45
           assert(irq id < NR HARD INTR);
46
47
           //if(irq_id == 0) panic("You have hit a timer interrupt, remove this
48
49
           struct IRQ t *f = handles[irq id];
50
           while (f != NULL) { /* call handlers one by one */
51
               f->routine();
52
               f = f->next;
53
54
           }
      }
55
56 }
```

这时就会调用显示的相关函数,将字符输出到屏幕上,最后再回到用户程序继续执行。