

Solution

The candidates are $\{pass, d, A\}$, $\{pass, e, A\}$, $\{pass, d, B\}$, $\{pass, e, B\}$. Let's just try them in order.

$\{pass, d, A\}$. Let's see if this is rational for Player 1 by comparing $pass$ to $play$. We get expected payoffs of

$$\begin{aligned} pass : \quad & \frac{1}{5}[2] + \frac{1}{5}[2] + \frac{3}{5}[2] = 2, \\ play : \quad & \frac{1}{5}[4] + \frac{1}{5}[4] + \frac{3}{5}[1] = 2.2. \end{aligned}$$

Oh no! This can't be a WSE; sequential rationality fails for Player 1.

$\{pass, e, A\}$. Let's see if this is rational for Player 1 by comparing $pass$ to $play$. We get expected payoffs of

$$\begin{aligned} pass : \quad & \frac{1}{5}[2] + \frac{1}{5}[2] + \frac{3}{5}[2] = 2, \\ play : \quad & \frac{1}{5}[4] + \frac{1}{5}[2] + \frac{3}{5}[0] = 1.2. \end{aligned}$$

Okay, so at least Player 1's strategy is sequentially rational. But what about the other two meatballs?

Because Player 1 passes always, information set $\{u, v\}$ is never reached. Therefore it faces no Bayesian restrictions – the beliefs can be anything. Accordingly, let's specify

$$\mu = \begin{pmatrix} u & v \\ p & 1-p \end{pmatrix}.$$

We can set the beliefs ourselves in such a way that will make A preferred to B , if such beliefs are feasible (i.e. are legit probabilities). Compare the expected values of A and B , given by

$$\begin{aligned} A : \quad & p[2] + (1-p)[4] = 4 - 2p, \\ B : \quad & p[0] + (1-p)[5] = 5 - 5p. \end{aligned}$$

Then A is rational when $4 - 2p \geq 5 - 5p$, that is, when $p \geq 1/3$.

Similarly, because Player 1 always passes, information set $\{x, y\}$ is never reached, so let's

write

$$\mu = \begin{pmatrix} x & y \\ q & 1 - q \end{pmatrix}.$$

We need e to be rational. Compare the expected values of d and e , given by

$$d : q[10] + (1 - q)[2] = 8q + 2,$$

$$e : q[3] + (1 - q)[4] = 4 - q.$$

Then e is rational when $4 - q \geq 8q + 2$, that is, when $q \leq 2/9$.

So we have found a WSE; all players are acting rationally, given what the other players are doing and what their beliefs are.

$$\mu = \left(\begin{array}{ccc|cc} r & s & t & x & y & u & v \\ \frac{1}{5} & \frac{1}{5} & \frac{3}{5} & q & 1 - q & p & 1 - p \end{array} \right).$$

such that $p \geq 1/3$ and $q \leq 2/9$. So there are an infinite number of beliefs that constitute a WSE. A particularly elegant solution is to just have $p = 1$ and $q = 0$.

{pass, d, B}. Same routine. Let's see if this is rational for Player 1 by comparing *pass* to *play*. We get expected payoffs of

$$pass : \frac{1}{5}[2] + \frac{1}{5}[2] + \frac{3}{5}[2] = 2,$$

$$play : \frac{1}{5}[0] + \frac{1}{5}[5] + \frac{3}{5}[1] = 1.6.$$

Okay, so at least Player 1's strategy is sequentially rational. But what about the other meatballs?

Because Player 1 passes always, information set $\{u, v\}$ is never reached. Therefore it faces no Bayesian restrictions – the beliefs can be anything. Accordingly, let's specify

$$\mu = \begin{pmatrix} u & v \\ p & 1 - p \end{pmatrix}.$$

We can set the beliefs ourselves in such a way that will make B preferred to A , if such beliefs

are feasible (i.e. are legit probabilities). Compare the expected values of A and B , given by

$$A : p[2] + (1 - p)[4] = 4 - 2p,$$

$$B : p[0] + (1 - p)[5] = 5 - 5p.$$

Then B is rational when $4 - 2p \leq 5 - 5p$, that is, when $p \leq 1/3$.

Similarly, because Player 1 always passes, information set $\{x, y\}$ is never reached, so let's write

$$\mu = \begin{pmatrix} x & y \\ q & 1 - q \end{pmatrix}.$$

We need d to be rational. Compare the expected values of d and e , given by

$$d : q[0] + (1 - q)[2] = 2 - 2q,$$

$$e : q[3] + (1 - q)[4] = 4 - q.$$

Then d is rational when $2 - 2q \geq 4 - q$, that is, when $q \leq -2$. So, never. (Notice that when B is played, e strictly dominates d ; had we noticed this initially, we wouldn't have had to bother looking at the other two players.) This cannot be sequentially rational for Player 2.

$\{pass, e, B\}$. Same routine. Let's see if this is rational for Player 1 by comparing *pass* to *play*. We get expected payoffs of

$$pass : \frac{1}{5}[2] + \frac{1}{5}[2] + \frac{3}{5}[2] = 2,$$

$$play : \frac{1}{5}[0] + \frac{1}{5}[2] + \frac{3}{5}[0] = 0.4.$$

Okay, so at least Player 1's strategy is sequentially rational. But what about the other meatballs?

Because Player 1 passes always, information set $\{u, v\}$ is never reached. Therefore it faces no Bayesian restrictions – the beliefs can be anything. Accordingly, let's specify

$$\mu = \begin{pmatrix} u & v \\ p & 1 - p \end{pmatrix}.$$

We can set the beliefs ourselves in such a way that will make B preferred to A , if such beliefs

are feasible (i.e. are legit probabilities). Compare the expected values of A and B , given by

$$A : p[2] + (1 - p)[4] = 4 - 2p,$$

$$B : p[0] + (1 - p)[5] = 5 - 5p.$$

Then B is rational when $4 - 2p \leq 5 - 5p$, that is, when $p \leq 1/3$.

Similarly, because Player 1 always passes, information set $\{x, y\}$ is never reached, so let's write

$$\mu = \begin{pmatrix} x & y \\ q & 1 - q \end{pmatrix}.$$

We need e to be rational. Well, we just found above that if B is played, then e strictly dominates d . Since B is played in this strategy profile, we can conclude that e is rational for any beliefs Player 2 might have. Although if you want to get specific about it again, compare the expected values of d and e , given by

$$d : q[0] + (1 - q)[2] = 2 - 2q,$$

$$e : q[3] + (1 - q)[4] = 4 - q.$$

So e is rational when $4 - q \geq 2 - 2q$, that is, when $q \geq -2$. Since q is a probability and can't be negative, this is always satisfied.

So we have found another WSE; all players are acting rationally, given what the other players are doing and what their beliefs are.

$$\sigma = \left(\begin{array}{cc|cc|cc} \text{pass} & \text{play} & d & e & A & B \\ 1 & 0 & 0 & 1 & 0 & 1 \end{array} \right), \quad \mu = \left(\begin{array}{ccc|cc|cc} r & s & t & x & y & u & v \\ \frac{1}{5} & \frac{1}{5} & \frac{3}{5} & q & 1 - q & p & 1 - p \end{array} \right),$$

such that $q \in [0, 1]$ and $p \leq 1/3$. Having $p = 0$ and $q = 0$ or $q = 1$ are particularly elegant cases.