

# ECE 443/518 – Computer Cyber Security

## Lecture 19 Proof of Work, Smart Contract

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# Outline

Proof of Work

Smart Contract

# Reading Assignment

- ▶ This lecture: Proof of Work, Smart Contract
- ▶ Next lecture: Bitcoin Security

# Outline

Proof of Work

Smart Contract

# BFT Meets Cryptocurrency

- ▶ Consensus on what the next block should be is a must for cryptocurrency.
  - ▶ Otherwise adversaries can create branches for double spending.
- ▶ Many BFT protocols are too weak to be useful here.
  - ▶ E.g. both the ones without or with digital signatures need to know how many traitors are there, but for cryptocurrency adversary can create arbitrary number of “traitor” accounts.
- ▶ The BFT propocol need to weight the participants differently.
  - ▶ So adversaries cannot simply overwhelm the protocol by creating more accounts, a.k.a. Sybil attack.

# Proof of Work (PoW)

- ▶ Who willing to generate the next block needs to perform certain amount of work before it could join the BFT protocol.
  - ▶ Typical work: for a block of hash  $s$ , find  $x$  so that  $h(s||x)$  is smaller than a threshold.
  - ▶ If  $h$  is preimage resistant, one can only find such  $x$  via brute-force.
  - ▶ Block time: the expected time for some account to find a solution  $x$ .
  - ▶ When more computational power are available, the threshold is reduced such that the block time remains unchanged.
- ▶ Who willing to participate will have an account to receive economic incentives for the work.
- ▶ Proof of Work consensus: the branch with the most of work is the correct one
  - ▶ The consensus can be reached as long as honest account owners can provide majority of work.

# Obtaining the Blockchain

- ▶ Consider an account that want to generate the next block.
  - ▶ By using a program together with its private key.
- ▶ The program connects to Internet to query the blockchain.
  - ▶ However, there could be adversaries so the program must decide if the chain it receives is valid or not.
- ▶ The genesis block: the first block of the blockchain.
  - ▶ The genesis block is assumed to be well-known, usually coded into the program directly.
  - ▶ The genesis block could contain data like cryptocurrency parameters and initial balances for certain accounts.
- ▶ The program need to validate past transactions.
  - ▶ It is necessary to accumulate balances for all accounts to decide if transactions are valid – this is possible now since our computers are actually quite powerful.
- ▶ However, recall that valid transactions along cannot prevent branches (and thus double spending).

# PoW Fork Choice

- ▶ Fork: the program may receive multiple blockchains all with the same correct genesis block.
  - ▶ They diverge somewhere back in the history.
  - ▶ Resulting from a temporary network partitioning or an attack.
- ▶ Choice: with PoW, the program should pick the blockchain with the most of the work as the correct one.
  - ▶ The work is measured as the total effort to solve the problems for all the blocks along the chain.
- ▶ 51% attack: attackers controlling more than half of the computational power could collude to cause a successful fork.
  - ▶ Suppose currently the honest accounts are at the chain  $A \rightarrow B \rightarrow \dots \rightarrow C$  from an earlier block A.
  - ▶ The attackers make a fork  $A \rightarrow B'$  and continue.
  - ▶ No matter how many blocks are there between A and C, the attackers can eventually reach  $D'$  as  $A \rightarrow B' \rightarrow \dots \rightarrow D'$ , that contains more work than the chain created by the honest accounts now as  $A \rightarrow B \rightarrow \dots \rightarrow C \rightarrow \dots \rightarrow D$



# PoW Finality

- ▶ With 51% attack, powerful attackers can revert transactions by creating successful forks.
  - ▶ It may take some time but 100% the attack will be successful.
- ▶ Can attackers with less computational power revert a block?
  - ▶ Finality: we need to define when a block is considered “final” and thus is not supposed to be changed or reverted.
  - ▶ Fake check scams are classical examples of attacks on finality for our banking system.
- ▶ Consider an attacker controlling 25% of computational power
  - ▶ Suppose the current chain is  $A \rightarrow B$  and honest accounts are working on the block  $C$ .
  - ▶ If  $B$  is considered final immediately, the attacker will attempt to make a fork  $A \rightarrow B' \rightarrow C'$  when  $A$  was ready.
  - ▶ If  $C'$  can be generated ahead of  $C$  in time, honest accounts may simply follow the chain  $A \rightarrow B' \rightarrow C'$ .
  - ▶ With 25% of computational power, this may happen with a probability of  $(\frac{25\%}{75\%})^2 = \frac{1}{9}$ .
- ▶ Practically, one should wait a few blocks to reduce the chance of having forks due to possible attacks.

# Economic Incentives

- ▶ How could PoW cryptocurrencies actually survive when finality is always probabilistic, and when powerful adversaries could have the majority of computational power?
- ▶ Economic incentives to attract honest accounts to participate in the BFT protocol.
  - ▶ Transaction fees: the account creates the next block will take all the transaction fees.
    - ▶ When there is more transactions than what the next block can hold, payers compete by paying more transaction fees.
  - ▶ Mining: the account creates the next block is allowed to award itself a predefined amount of money.
    - ▶ As a transaction with no payer.
- ▶ As a consequence, powerful adversaries have economic incentives to not cheat.
  - ▶ It is more rewarding to participate honestly than to make the cryptocurrency useless by attacking it.

# Proof of Stake (PoS)

- ▶ PoW consensus achieves a great success and enables a lot of honest account owners to participate.
- ▶ For a fixed block time, need to increase complexity of work.
  - ▶ So more energy is needed to generate one block.
  - ▶ hardware depreciation + energy cost + profit = mining income
- ▶ Proof of stake: accounts stake a certain amount of the cryptocurrency itself to participate in the consensus process.
  - ▶ Without the need of computing complex works (and thus consume less energy) to resist Sybil attacks.
  - ▶ Honest accounts are rewarded with transaction fees.
  - ▶ Attackers may have their staked cryptocurrency burned.

# Outline

Proof of Work

Smart Contract

# From Ledger to State Machine

- ▶ The ledger as stored in the block chain can be treated as a very simple state machine.
  - ▶ Initial state: initial account balances
  - ▶ Current state: current account balances
  - ▶ State transitions: each blockchain transaction updates account balances by addition and subtraction.
- ▶ The blockchain can support more complex state machines.
  - ▶ Allow accounts to define state variables in addition to balance.
  - ▶ Allow blockchain transactions to perform more operations on state variables than simple addition and subtraction.
- ▶ This is similar to how we build computer hardware and software to support general purpose computing need.
  - ▶ E.g. Ethereum Virtual Machine (EVM) defined by the Ethereum blockchain uses 8-bit opcode and a stack to organize its 256-bit registers, and supports high-level programming languages like Solidity.

# Smart Contract

- ▶ What are the benefits of running state machines and thus programs in a blockchain?
  - ▶ Not for efficiency: each computation needs to be executed as many times as anyone would need to validate the blockchain, using the same inputs and generating the same output.
  - ▶ Nonrepudiation: the account initiates a computation must sign the request and cannot deny so.
  - ▶ Integrity: the outcome is permanently recorded in the blockchain and cannot be reverted.
  - ▶ As long as there is no branch.
- ▶ That is what is necessary to execute a contract.
  - ▶ Smart contract: a program running inside a blockchain.

# A Smart Contract Example

```
pragma solidity 0.8.7;

contract VendingMachine {
    // Declare state variables of the contract
    address public owner;
    mapping (address => uint) public cupcakeBalances;

    // When 'VendingMachine' contract is deployed:
    // 1. set the deploying address as the owner of the contract
    // 2. set the deployed smart contract's cupcake balance to 100
    constructor() {
        owner = msg.sender;
        cupcakeBalances[address(this)] = 100;
    }
    ...
}
```

- ▶ A smart contract that you can buy cupcakes on Ethereum.
- ▶ No you don't receive an actual cupcake.
  - ▶ What you received could be treated as a ticket or token to redeem a physical cupcake somewhere.

# Smart Contract Account

```
...  
constructor() {  
    owner = msg.sender;  
    cupcakeBalances[address(this)] = 100;  
}  
...
```

- ▶ Once created, a smart contract will have its own address, as indicated by `address(this)`
- ▶ Other accounts interact with the smart contract by sending (signed) messages to the smart contract account.
- ▶ The smart contract will handle these messages in member functions.
  - ▶ `constructor` is a special one called for the first message which deploys the smart contract.



# The Message Sender

```
contract VendingMachine {  
    // Declare state variables of the contract  
    address public owner;  
    mapping (address => uint) public cupcakeBalances;  
  
    // When 'VendingMachine' contract is deployed:  
    // 1. set the deploying address as the owner of the contract  
    // 2. set the deployed smart contract's cupcake balance to 100  
    constructor() {  
        owner = msg.sender;  
        cupcakeBalances[address(this)] = 100;  
    }  
    ...  
}
```

- ▶ `msg.sender` indicates who initiates the computation.
  - ▶ The payer of cryptocurrency.
  - ▶ You pay to deploy a smart contract and to interact with it – you are consuming computational resources in the blockchain.
- ▶ The sender should in addition specify what transactions (member function) is to be performed (called).
  - ▶ E.g. one of `constructor`, `refill`, and `purchase`
  - ▶ Plus other necessary parameters.

# Transactions

```
contract VendingMachine {
    ...
    // Allow the owner to increase the smart contract's cupcake balance
    function refill(uint amount) public {
        require(msg.sender == owner, "Only the owner can refill.");
        cupcakeBalances[address(this)] += amount;
    }

    // Allow anyone to purchase cupcakes
    function purchase(uint amount) public payable {
        require(msg.value >= amount * 1 ether, "1 ETH per cupcake");
        require(cupcakeBalances[address(this)] >= amount, "Not enough in stock");
        cupcakeBalances[address(this)] -= amount;
        cupcakeBalances[msg.sender] += amount;
    }
}
```

- ▶ `msg.value` indicates money the sender pays the the contract.
  - ▶ The money is transfered from the sender address to the contract address automatically if the computation completes successfully.
- ▶ How could one withdraw money from the contract?

# Complications

- ▶ What if there is an infinite loop into a smart contract?
  - ▶ Can be exploited by adversaries to jam the blockchain.
  - ▶ In theory, we cannot detect if there is an infinite loop in a program.
  - ▶ On blockchain, we can solve the issue by limiting the number of instructions a smart contract may execute by the transaction fee the sender would like to pay.
- ▶ Since the program of a smart contract need to be deployed to the blockchain, everyone can see and analyze it.
  - ▶ Bugs in the program could be found and exploited by adversaries.

# Summary

- ▶ Both proof of work (PoW) and proof of stake (PoS) work as the consensus mechanism for cryptocurrencies.
- ▶ Smart contracts are programs running inside a blockchain, reacting to blockchain events.