

[CSED211] Introduction to Computer Software Systems

Lecture 5: Control

Prof. Jisung Park



CAOS

COMPUTER ARCHITECTURE &
OPERATING SYSTEMS LABORATORY

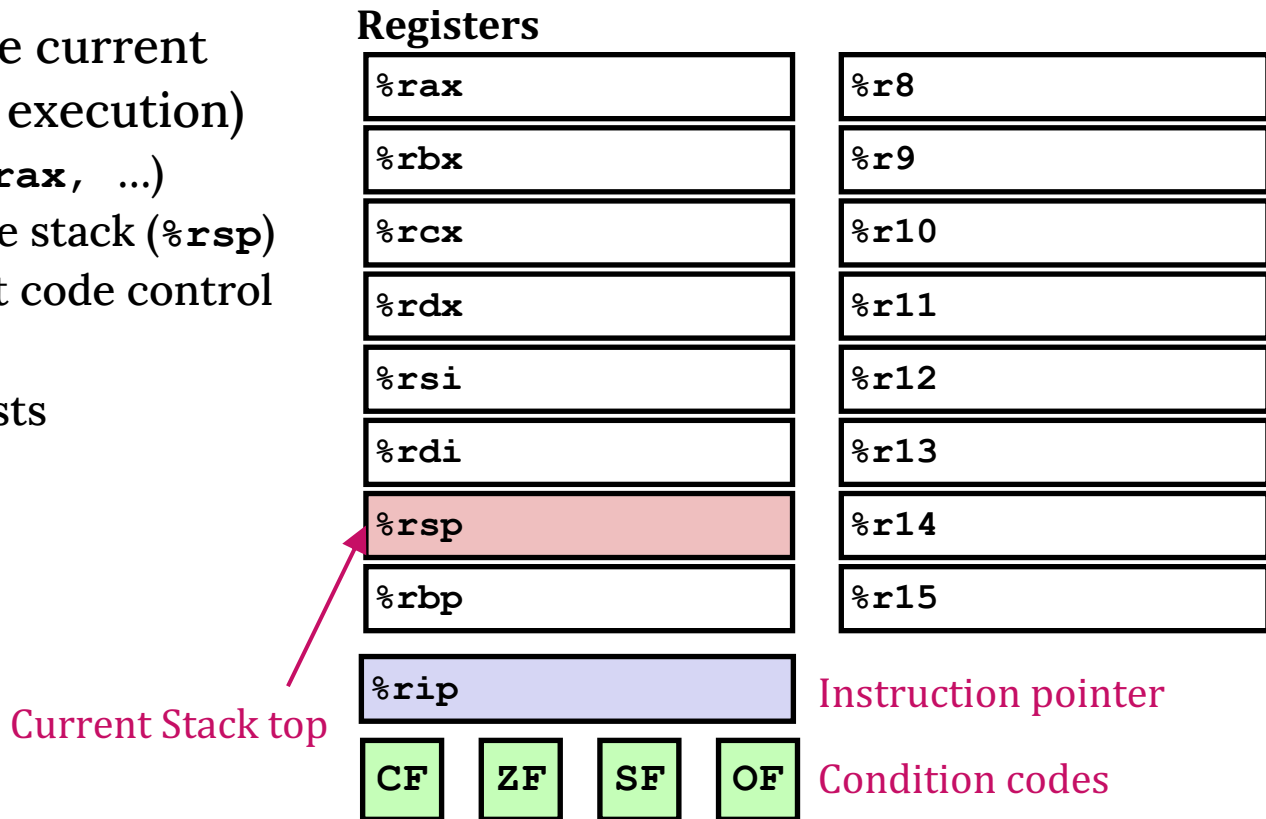
2023.09.25

Today

- Control: Condition Codes
- Conditional Branches
- Loops
- Switch Statements

Processor State (x86-64, Partial)

- Information about the current CPU status (program execution)
 - Temporary data (**%rax**, ...)
 - Location of runtime stack (**%rsp**)
 - Location of current code control point (**%rip**)
 - Status of recent tests (**CF**, **ZF**, **SF**, **OF**)



Condition Codes: Implicit Setting

- Single bit registers
 - **CF**: Carry Flag (for unsigned)
 - **SF**: Sign Flag (for signed)
 - **ZF**: Zero Flag
 - **OF**: Overflow Flag (for signed)
- Implicitly set (think of it as side effect) by arithmetic instructions

Example: `addq src, dst` \rightarrow `t = a + b`

CF set if carry/borrow out from most significant bit (unsigned overflow)

ZF set if `t == 0`

SF set if `t < 0` (as signed)

OF set if two's-complement (signed) overflow

`(a>0 && b>0 && t<0) || (a<0 && b<0 && t>=0)`

- **Not** set by `leaq` instruction

Condition Codes: Explicit Setting - Compare

- Explicit setting by compare instruction

- `cmpq src2, src1`

- `cmpq b, a` is like computing `(a-b)` without setting destination

CF set if carry/borrow out from the MSB (used for unsigned comparisons)

ZF set if `(a-b) == 0`, i.e., `a = b`

SF set if `(a-b) < 0` (as signed)

OF set if two's-complement (signed) overflow

`(a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)`

Condition Codes: Explicit Setting - Test

- Explicit setting by test instruction
 - `testq src2, src1`
 - `testq b, a` is like computing `(a&b)` without setting destination
 - Useful to have one of the operands be a mask
 - `ZF` set when `a&b == 0`
 - `SF` set when `a&b < 0`
 - `CF`, `OF` clear
- Can be used to check if the content is zero: `testq %rax, %rax`

Condition Codes: Explicit Reading - Set

- Explicit reading by **setX** instructions
 - **setX dst**: sets the least-significant byte of destination **dst** to 0 or 1 based on combinations of condition codes
 - Does **not** alter the remaining 7 bytes of **dst**

SetX	Condition	Description
sete	ZF	Equal / Zero
setne	~ZF	Not Equal / Not Zero
sets	SF	Negative
setns	~SF	Nonnegative
setg	~ (SF^OF) & ~ZF	Greater (Signed)
setge	~ (SF^OF)	Greater or Equal (Signed)
setl	(SF^OF)	Less (Signed)
setle	(SF^OF) ZF	Less or Equal (Signed)
seta	~CF & ~ZF	Above (unsigned)
setb	CF	Below (unsigned)

x86-64 Integer Registers

- Can reference the least-significant byte

%rax	%al
%rbx	%bl
%rcx	%cl
%rdx	%dl
%rsi	%sil
%rdi	%dil
%rps	%spl
%rbp	%bpl

%r8	%r8b
%r9	%r9b
%r10	%r10b
%r11	%r11b
%r12	%r12b
%r13	%r13b
%r14	%r14b
%r15	%r15b

Explicit Reading Condition Codes (Cont.)

- **setX** instructions: set a single byte based on combination of condition codes
 - One of addressable byte registers
 - Does not alter the remaining bytes
 - Typically use **movzb1** to finish job
 - 32-bit instructions also set upper 32 bits to 0

Register	Use(s)
<code>%rdi</code>	Argument x
<code>%rsi</code>	Argument y
<code>%rax</code>	Return value

```
int gt (long x, long y){  
    return x > y;  
}
```

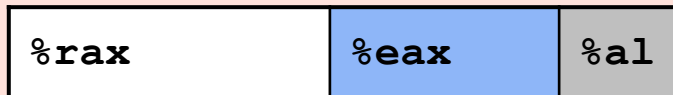
```
cmpq    %rsi, %rdi    # Compare x and y  
setg    %al           # Set when x > y  
movzb1  %al, %eax.    # Zero rest of %rax  
ret
```

Explicit Reading Condition Codes (Cont.)

- **setX** instructions: set a single byte based on combination of condition codes
 - One of addressable byte registers
 - Does not alter the remaining bytes
 - Typically use `movzbl` to finish job
 - 32-bit instructions also set upper 32 bits to 0

Beware weirdness `movzbl` (and others, i.e., `movxyz`)

`movzbl %al, %eax`



`%rax`

Return value

```
cmpq    %rsi, %rdi    # Compare x and y
setg     %al           # Set when x > y
movzbl   %al, %eax.    # Zero rest of %rax
ret
```

Today

- Control: Condition Codes
- **Conditional Branches**
- Loops
- Switch Statements

Jumping

- **jX** instructions: jump to different part of code depending on condition codes

jX	Condition	Description
jmp	1	Unconditional
je	ZF	Equal / Zero
jne	~ZF	Not Equal / Not Zero
js	SF	Negative
jns	~SF	Nonnegative
jg	~(SF^OF) & ~ZF	Greater (Signed)
jge	~(SF^OF)	Greater or Equal (Signed)
jl	(SF^OF)	Less (Signed)
jle	(SF^OF) ZF	Less or Equal (Signed)
ja	~CF & ~ZF	Above (unsigned)
jb	CF	Below (unsigned)

Conditional Branch Example: Old Style

- Generation

\$ gcc -Og -S -fno-if-conversion control.c

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
long absdiff(long x, long y){  
    long result;  
    if (x > y)  
        result = x-y;  
    else  
        result = y-x;  
    return result;  
}
```

```
absdiff:  
    cmpq    %rsi, %rdi    # x:y  
    jle     .L4  
    movq    %rdi, %rax  
    subq    %rsi, %rax  
    ret  
.L4:      # x <= y  
    movq    %rsi, %rax  
    subq    %rdi, %rax  
    ret
```

Expressing with Goto Code

- C supports goto statement
 - Jump to position designated by label

```
long absdiff(long x, long y){  
    long result;  
    if (x > y)  
        result = x-y;  
    else  
        result = y-x;  
    return result;  
}
```

```
long absdiff(long x, long y){  
    long result;  
    int ntest = x <= y;  
    if (ntest) goto Else;  
    result = x-y;  
    goto Done;  
Else:  
    result = y-x;  
Done:  
    return result;  
}
```

General Conditional Expression Translation

- C code using branch

```
val = test ? then_expr : else_expr;
```

- Goto version

```
ntest = !(test)
if (ntest) goto Else;
val = then_expr
goto Done;
Else:
    val = else_expr;
Done:
...
```

- Create separate code regions for then & else expressions
- Execute appropriate one

Conditional Moves

- Conditional move instructions
 - Instruction supports:
`if (test) dst = src`
 - Supported in post-1995 x86 processors
 - GCC tries to use them
 - But, only when known to be safe
- Why?
 - Branches are very disruptive to instruction flow through pipelines
 - Conditional moves do not require control transfer

Branch Version

```
val = test
    ? then_expr
    : else_expr;
```

Goto Version

```
result = then_expr;
eval = else_expr;
nt = !(test)
if (nt) result = eval;
return result;
```


Conditional Move Example

```
long absdiff(long x, long y){  
    long result;  
    if (x > y)  
        result = x-y;  
    else  
        result = y-x;  
    return result;  
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
absdiff:  
    cmpq    %rsi, %rdi    # x:y  
    jle     .L4  
    movq    %rdi, %rax  
    subq    %rsi, %rax  
    ret  
.L4:      # x <= y  
    movq    %rsi, %rax  
    subq    %rdi, %rax  
    ret
```

Conditional Move Example

```
long absdiff(long x, long y){
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
absdiff:
    movq    %rdi, %rax    # x
    subq    %rsi, %rax    # result = x-y
    movq    %rsi, %rdx
    subq    %rdi, %rdx    # eval = y-x
    cmpq    %rsi, %rdi    # Compare x and y
    cmovle  %rdx, %rax    # if <=, result = eval
    ret
```

Bad Cases for Conditional Move

- Expensive computations

```
val = test(x) ? hard1(x) : hard2(x);
```

- Both expressions get computed
- Only makes sense when computations are very simple

Bad Performance

- Risky computations

```
val = p ? *p : 0;
```

- Both expressions get computed
- May have undesirable effects (e.g., null-pointer exception)

Unsafe

- Computations with side effects

```
val = x > 0 ? x*=7 : x+=3;
```

- Both values get computed
- Must be side-effect free

Illegal

Exercise

```
xorq    %rax, %rax
subq    $1, %rax
cmpq    $2, %rax
setl    %al
movzbl  %al, %eax
```

%rax	SF	CF	OF	ZF

- `cmpq b, a` \rightarrow $(a-b)$ w/o setting dst
 CF set if carry/borrow out from MSB
 (used for unsigned comparisons)
 ZF set if $a == b$
 SF set if $(a-b) < 0$ (as signed)
 OF set if two's-complement (signed) overflow
- `test b, a` \rightarrow $(a \& b)$ w/o setting dest
 SF, ZF set based on result
 CF, OF cleared
- `setl` and `movzbl` do **not** modify condition codes

setX	Condition	Description
<code>sete</code>	ZF	Equal / Zero
<code>setne</code>	$\sim ZF$	Not Equal / Not Zero
<code>sets</code>	SF	Negative
<code>setns</code>	$\sim SF$	Nonnegative
<code>setg</code>	$\sim (SF \wedge OF) \ \& \ \sim ZF$	Greater (signed)
<code>setge</code>	$\sim (SF \wedge OF)$	Greater or Equal (signed)
<code>setl</code>	$SF \wedge OF$	Less (signed)
<code>setle</code>	$(SF \wedge OF) \ \ ZF$	Less or Equal (signed)
<code>seta</code>	$\sim CF \ \& \ \sim ZF$	Above (unsigned)
<code>setb</code>	CF	Below (unsigned)

Exercise

```

xorq    %rax, %rax
subq    $1, %rax
cmpq    $2, %rax
setl    %al
movzbl  %al, %eax
    
```

%rax	SF	CF	OF	ZF
0000 0000 0000 0000	0	0	0	1

- **cmpq b, a** \rightarrow (a-b) w/o setting dst
CF set if carry/borrow out from MSB
 (used for unsigned comparisons)
ZF set if **a == b**
SF set if (a-b) < 0 (as signed)
OF set if two's-complement (signed) overflow
- **test b, a** \rightarrow (a&b) w/o setting dest
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setge	~(SF^OF)	Greater or Equal (signed)
setl	SF^OF	Less (signed)
setle	(SF^OF) ZF	Less or Equal (signed)
seta	~CF & ~ZF	Above (unsigned)
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Exercise

```
xorq    %rax, %rax
subq    $1, %rax
cmpq    $2, %rax
setl    %al
movzbl  %al, %eax
```

%rax	SF	CF	OF	ZF
0000 0000 0000 0000	0	0	0	1
FFFF FFFF FFFF FFFF	1	1	0	0

- `cmpq b, a` \rightarrow $(a-b)$ w/o setting dst
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Exercise

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setl    %al
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%rax	SF	CF	OF	ZF
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FFFF FFFF FFFF FFFF	1	1	0	0
FFFF FFFF FFFF FFFF	1	0	0	0

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Exercise

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xorq    %rax, %rax
subq    $1, %rax
cmpq    $2, %rax
setl   %al
movzbl  %al, %eax
```

%rax	SF	CF	OF	ZF
0000 0000 0000 0000	0	0	0	1
FFFF FFFF FFFF FFFF	1	1	0	0
FFFF FFFF FFFF FFFF	1	0	0	0
FFFF FFFF FFFF FF 01	1	0	0	0

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Exercise

```
xorq    %rax, %rax
subq    $1, %rax
cmpq    $2, %rax
setl    %al
movzbl  %al, %eax
```

%rax	SF	CF	OF	ZF
0000 0000 0000 0000	0	0	0	1
FFFF FFFF FFFF FFFF	1	1	0	0
FFFF FFFF FFFF FFFF	1	0	0	0
FFFF FFFF FFFF FF01	1	0	0	0
0000 0000 0000 0001	1	0	0	0

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Today

- Control: Condition Codes
- Conditional Branches
- **Loops**
- Switch Statements

Do-While Loop Example

Do-While Version

```
long pcount_do(unsigned long x) {  
    long result = 0;  
    do{  
        result += x & 0x1;  
        x >>= 1;  
    } while(x);  
    return result;  
}
```

Goto Version

```
long pcount_goto(unsigned long x) {  
    long result = 0;  
    Loop:  
        result += x & 0x1;  
        x >>= 1;  
        if(x) goto Loop;  
    return result;  
}
```

- Count number of 1's in argument **x** (popcount)
- Use conditional branch to either continue looping or to exit loop

Do-While Loop Compilation

Goto Version

```
long pcount_goto(unsigned long x) {  
    long result = 0;  
Loop:  
    result += x & 0x1;  
    x >>= 1;  
    if(x) goto Loop;  
    return result;  
}
```

```
    movl    $0, %eax           # result = 0  
.L2:  
    # loop:  
    movq    %rdi, %rdx  
    andl    $1, %edx           # t = x & 0x1  
    addq    %rdx, %rax         # result += t  
    shrq    %rdi               # x >>= 1  
    jne     .L2                # if (x) goto loop  
    ret
```

Register	Use(s)
%rdi	Argument x
%rax	result

General Do-While Loop Translation

Do-While Version

```
do  
    body_expr  
while(test);
```



Goto Version

```
Loop:  
    body_expr  
    if(test) goto Loop;
```

```
body_expr:  
    {  
        statement1;  
        statement2;  
        ...  
        statementn;  
    }
```

General While Translation#1

- **Jump-to-middle** translation
- Used with -Og

While Version

```
while (test)  
    body_expr
```



Goto Version

```
goto Test;  
Loop:  
    body_expr  
Test:  
    if (test) goto Loop;
```

While Loop Example#1

While Version

```
long pcount_while(unsigned long x){
    long result = 0;
    while(x){
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

Goto (Jump-to-Middle) Version

```
long pcount_goto_jtm(unsigned long x){
    long result = 0;
    goto Test;
Loop:
    result += x & 0x1;
    x >>= 1;
Test:
    if(x) goto Loop;
    return result;
}
```

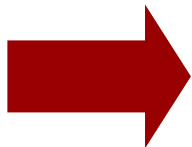
- Compare to do-while version of the function
- Initial `goto` starts loop at **Test**

General While Translation #2

- Do-while conversion
- Used with -O1

While Version

```
while (test)
    body_expr
```



Do-While Version

```
if (!test)
    goto Done;
do
    body_expr
while (test)
Done:
```



Goto Version

```
if (!test)
    goto Done;
Loop:
    body_expr
    if (test) goto Loop;
Done:
```


While Loop Example#1

While Version

```
long pcount_while(unsigned long x){
    long result = 0;
    while(x){
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

Goto Version (via Do-While Conversion)

```
long pcount_goto_dw(unsigned long x){
    long result = 0;
    if(!x) goto Done;
Loop:
    result += x & 0x1;
    x >>= 1;
    if(x) goto Loop;
Done:
    return result;
}
```

- Compare to do-while version of the function
- Initial `goto` starts loop at **Test**

While Loop Example#1

Goto (Jump-to-Middle) Version

```
long pcount_goto_jtm(unsigned long x) {  
    long result = 0;  
    goto Test;  
  
Loop:  
    result += x & 0x1;  
    x >>= 1;  
  
Test:  
    if(x) goto Loop;  
    return result;  
}
```

Goto Version (via Do-While Conversion)

```
long pcount_goto_dw(unsigned long x) {  
    long result = 0;  
    if(!x) goto Done;  
  
Loop:  
    result += x & 0x1;  
    x >>= 1;  
    if(x) goto Loop;  
  
Done:  
    return result;  
}
```

- Compare to do-while version of the function
- Initial `goto` starts loop at **Test**

General For Loop Translation

For Loop Version

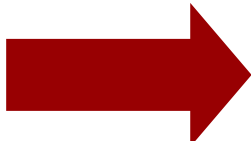
```
for(init; test; update)
    body_expr
```

```
#define WSIZE 8*sizeof(long)
long pcount_for(unsigned long x){
    size_t i;
    long result = 0;
    for(i = 0; i < WSIZE; i++){
        unsigned bit = (x >> i) & 0x1;
        result += bit;
    }
    return result;
}
```

General For Loop Translation: While Conversion

For Loop Version

```
for(init; test; update)
    body_expr
```



While Version

```
init;
while(test){
    body_expr
    update;
}
```

```
#define WSIZE 8*sizeof(long)
long pcount_for(unsigned long x){
    size_t i;
    long result = 0;
    for(i = 0; i < WSIZE; i++){
        unsigned bit = (x >> i) & 0x1;
        result += bit;
    }
    return result;
}
```

```
long pcount_for_while(unsigned long x){
    size_t i;
    long result = 0;
    i = 0;
    while (i < WSIZE){
        unsigned bit = (x >> i) & 0x1;
        result += bit;
        i++;
    }
    ...
}
```

For Loop Do-While Conversion

For Loop Version

```
#define WSIZE 8*sizeof(long)
long pcount_for(unsigned long x){
    size_t i;
    long result = 0;
    for(i = 0; i < WSIZE; i++){
        unsigned bit = (x >> i) & 0x1;
        result += bit;
    }
    return result;
}
```

Goto Version (via Do-While Conversion)

```
long pcount_for_goto_dw(unsigned long x) {
    size_t i;
    long result = 0;
    i = 0; // init
    if(!(i < WSIZE)) goto Done; // !test
Loop:
    unsigned bit = (x >> i) & 0x1; // body
    result += bit; // body
    i++; // update
    if(i < WSIZE) goto Loop; // test
Done:
    return result;
}
```

For Loop Do-While Conversion

For Loop Version

```
#define WSIZE 8*sizeof(long)

long pcount_for(unsigned long x){
    size_t i;
    long result = 0;
    for(i = 0; i < WSIZE; i++){
        unsigned bit = (x >> i) & 0x1;
        result += bit;
    }
    return result;
}
```

Goto Version (via Do-While Conversion)

```
long pcount_for_goto_dw(unsigned long x) {
    size_t i;
    long result = 0;
    i = 0; // init
if(!(i < WSIZE)) goto Done; // !test
Loop:
    unsigned bit = (x >> i) & 0x1; // body
    result += bit; // body
    i++; // update
    if(i < WSIZE) goto Loop; // test
Done:
    return result;
}
```

- Initial test can be optimized away

Today

- Control: Condition Codes
- Conditional Branches
- Loops
- Switch Statements

Switch Statement Example

- Multiple case labels
 - e.g., cases 5 & 6
- Fall through cases
 - e.g., case 2
- Missing cases
 - e.g., case 4

```
long switch_eg(long x, long y, long z){  
    long w = 1;  
    switch(x) {  
        case 1:  
            w = y*z; break;  
        case 2:  
            w = y/z; // Fall through  
        case 3:  
            w += z; break;  
        case 5:  
        case 6:  
            w -= z; break;  
        default:  
            w = 2;  
    }  
    return w;  
}
```


Jump Table Structure

Switch Form

```
switch(x) {  
  case val_0:  
    cblock0  
  case val_1:  
    cblock1  
    ...  
  case val_n-1:  
    cblock(n-1)  
}
```

Translation (Extended C)

```
goto *JTab[x];
```

Jump Table

JTab:

Targ0
Targ1
Targ2
...
Targ(n-1)

Jump Targets

Targ0:

Code Block 0

Targ1:

Code Block 1

Targ2:

Code Block 2

...

Targ(n-1):

Code Block (n-1)

Switch Statement Example

```
long switch_eg(long x, long y, long z){  
    long w = 1;  
    switch(x){  
        ...  
    }  
    return w;  
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Setup

```
switch_eg:  
    movq    %rdx, %rcx  
    cmpq    $6, %rdi      # x:6  
    ja      .L8            # default  
    jmp     *.L4(, %rdi, 8) # goto *JTab[x]
```

Note that `w` is not initialized here

Switch Statement Example

```
long switch_eg(long x, long y, long z){  
    long w = 1;  
    switch(x){  
        ...  
    }  
    return w;  
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Setup

```
switch_eg:  
    movq    %rdx, %rcx  
    cmpq    $6, %rdi      # x:6  
    ja      .L8            # default  
    jmp     (*).L4(, %rdi, 8) # goto *JTab[x]
```

Jump table

```
.section      .rodata # read only  
    .align 8  
.L4:  
    .quad    .L8 # x = 0  
    .quad    .L3 # x = 1  
    .quad    .L5 # x = 2  
    .quad    .L9 # x = 3  
    .quad    .L8 # x = 4  
    .quad    .L7 # x = 5  
    .quad    .L7 # x = 6
```

Indirect Jump

Assembly Setup Explanation

- Table structure
 - Each target requires 8 bytes
 - Base address at `.L4`
- Jumping
 - **Direct:** `jmp .L8`
 - Jump target is denoted by label `.L8`
 - **Indirect:** `jmp *.L4(,%rdi,8)`
 - Must scale by factor of 8:
an address (pointer) is 8 bytes
 - Fetch target from the effective address
`.L4 + x*8`, only for $0 \leq x \leq 6$

Jump table

```
.section      .rodata # read only
    .align 8
.L4:
    .quad     .L8 # x = 0
    .quad     .L3 # x = 1
    .quad     .L5 # x = 2
    .quad     .L9 # x = 3
    .quad     .L8 # x = 4
    .quad     .L7 # x = 5
    .quad     .L7 # x = 6
```

Jump Table: with the Example C Code

```
long switch_eg(long x, long y, long z){
    long w = 1;
    switch(x) {
    case 1:      // .L3
        w = y*z; break;
    case 2:      // .L5
        w = y/z; // Fall through
    case 3:      // .L9
        w += z; break;
    case 5:      // .L7
    case 6:      // .L7
        w -= z; break;
    default:    // .L8
        w = 2;
    }
    return w;
}
```

Jump table

```
.section      .rodata # read only
    .align 8
.L4:
    .quad     .L8 # x = 0
    .quad     .L3 # x = 1
    .quad     .L5 # x = 2
    .quad     .L9 # x = 3
    .quad     .L8 # x = 4
    .quad     .L7 # x = 5
    .quad     .L7 # x = 6
```

Code Blocks: $x==1$

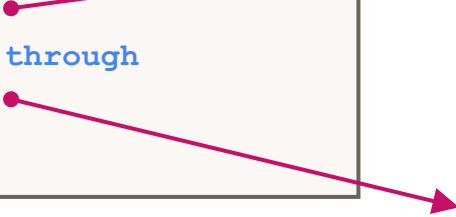
```
switch(x) {  
case 1:    // .L3  
    w = y*z; break;
```

```
.L3:  
movq    %rsi, %rax # ret = y  
imulq   %rdx, %rax # ret *= z  
ret
```


Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value

Handling Fall-Through

```
switch(x) {  
  ...  
  case 2:      // .L5  
    w = y/z;   // Fall through  
  case 3:      // .L9  
    w += z; break;
```



```
case 2:  
  w = y/z;  
  goto Merge;
```



```
case 3:  
  w = 1;
```

```
Merge:  
  w += z;
```

Code Blocks: $x == 2$ and $x == 3$

```
switch(x) {
...
case 2:      // .L5
    w = y/z; // Fall through
case 3:      // .L9
    w += z; break;
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	Return value
%rcx	Argument z

```
.L5:                                # Case 2
    movq    %rsi, %rax
    cqto    # extend %rax
    idivq   %rcx    # y/z
    jmp     .L6     # goto merge
.L9:                                # Case 3
    movl    $1, %eax # w = 1
.L6:                                # Merge:
    addq    %rcx, %rax # w += z
    ret
```

Instruction	Effect	Description
imulq S	$R[\%rdx]:R[\%rax] \leftarrow S \times R[\%rax]$	Signed full multiply
mulq S	$R[\%rdx]:R[\%rax] \leftarrow S \times R[\%rax]$	Unsigned full multiply
cltq	$R[\%rax] \leftarrow \text{SignExtend}(R[\%eax])$	Convert %eax to quad word
cqto	$R[\%rdx]:R[\%rax] \leftarrow \text{SignExtend}(R[\%rax])$	Convert to oct word
idivq S	$R[\%rdx] \leftarrow R[\%rdx]:R[\%rax] \bmod S;$ $R[\%rax] \leftarrow R[\%rdx]:R[\%rax] \div S$	Signed divide
divq S	$R[\%rdx] \leftarrow R[\%rdx]:R[\%rax] \bmod S;$ $R[\%rax] \leftarrow R[\%rdx]:R[\%rax] \div S$	Unsigned divide

Code Blocks: `x == 5`, `x == 6` and default

```
switch(x) {  
...  
case 5:      // .L7  
case 6:      // .L7  
    w -= z; break;  
default:     // .L8  
    w = 2;  
}
```

```
.L7:                # Case 5,6  
    movl    $1, %eax    # ret = w = 1  
    subq    %rdx, %rax  # ret -= z  
    ret  
.L8:                # Default:  
    movl    $2, %eax    # ret = 2  
    ret
```

Register	Use(s)
%rdi	Argument <code>x</code>
%rsi	Argument <code>y</code>
%rdx	Argument <code>z</code>
%rax	Return value
%rcx	Argument <code>z</code>

Finding Jump Table in Binary

```
00000000004005e0 <switch_eg>:
4005e0:    48 89 d1                mov     %rdx,%rcx
4005e3:    48 83 ff 06            cmp     $0x6,%rdi
4005e7:    77 2b                  ja      400614 <switch_eg+0x34>
4005e9:    ff 24 fd f0 07 40 00   jmpq    *0x4007f0(,%rdi,8)
4005f0:    48 89 f0                mov     %rsi,%rax
4005f3:    48 0f af c2            imul    %rdx,%rax
4005f7:    c3                     retq
4005f8:    48 89 f0                mov     %rsi,%rax
4005fb:    48 99                  cqto
4005fd:    48 f7 f9                idiv    %rcx
400600:    eb 05                  jmp     400607 <switch_eg+0x27>
400602:    b8 01 00 00 00        mov     $0x1,%eax
400607:    48 01 c8                add     %rcx,%rax
40060a:    c3                     retq
40060b:    b8 01 00 00 00        mov     $0x1,%eax
400610:    48 29 d0                sub     %rdx,%rax
400613:    c3                     retq
400614:    b8 02 00 00 00        mov     $0x2,%eax
400619:    c3                     retq
```

Finding Jump Table in Binary

00000000004005e0 <switch_eg>:

```

4005e0:  48 89 d1      mov    %rdx,%rcx
4005e3:  48 83 ff 06   cmp    $0x6,%rdi
4005e7:  77 2b        ja     400614 <switch_eg+0x34>
4005e9:  ff 24 fd f0 07 40 00 jmpq   *0x4007f0(,%rdi,8)

```

x==1

4005f0: 48 89 f0 % gdb switch
(gdb) x /8xg 0x4007f0

x==2

4005f8: 48 89 f0 0x4007f0: 0x0000000000400614 0x00000000004005f0
 4005fb: 48 99 0x400800: 0x00000000004005f8 0x0000000000400602
 4005fd: 48 f7 f9 0x400810: 0x0000000000400614 0x000000000040060b
 400600: eb 05 0x400820: 0x000000000040060b 0x2c646c25203d207

x==3

400602: b8 01 00 00 00 jmp 400607 <switch_eg+0x27>
 400607: 48 01 c8 mov \$0x1,%eax
 40060a: c3 add %rcx,%rax
 40060b: b8 01 00 00 00 retq
 400610: 48 29 d0 mov \$0x1,%eax
 400613: c3 sub %rdx,%rax
 400614: b8 02 00 00 00 retq
 400619: c3 retq

x==5, x==6

default

Finding Jump Table in Binary

	00000000004005e0 <switch_eg>:	
	4005e0: 48 89 d1	mov %rdx,%rcx
	4005e3: 48 83 ff 06	cmp \$0x6,%rdi
	4005e7: 77 2b	ja 400614 <switch_eg+0x34>
	4005e9: ff 24 fd f0 07 40 00	jmpq *0x4007f0(,%rdi,8)
x==1	4005f0: 48 89 f0	mov %rsi,%rax
	4005f3: 48 0f af c2	imul %rdx,%rax
	4005f7: c3	retq
x==2	4005f8: 48 89 f0	mov %rsi,%rax
	4005fb: 48 99	cqto
	4005fd: 48 f7 f9	idiv %rcx
	400600: eb 05	jmp 400607 <switch_eg+0x27>
x==3	400602: b8 01 00 00 00	mov \$0x1,%eax
	400607: 48 01 c8	add %rcx,%rax
	40060a: c3	retq
x==5, x==6	40060b: b8 01 00 00 00	mov \$0x1,%eax
	400610: 48 29 d0	sub %rdx,%rax
	400613: c3	retq
default	400614: b8 02 00 00 00	mov \$0x2,%eax
	400619: c3	retq

Summary

- C control
 - If-then-else
 - Do-while, while, for
 - Switch
- Assembler control
 - Conditional jump
 - Conditional move
 - Indirect jump (via jump tables)
 - Compiler generates code sequence to implement more complex control
- Standard techniques
 - Loops converted to do-while or jump-to-middle form
 - Large switch statements use jump tables
 - Sparse switch statements may use **decision trees** (if-elseif-elseif-else)

[CSED211] Introduction to Computer Software Systems

Lecture 5: Control

Prof. Jisung Park



CAOS

COMPUTER ARCHITECTURE &
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2023.09.25