

# IBM开源技术微讲堂

## 区块链和HyperLedger系列

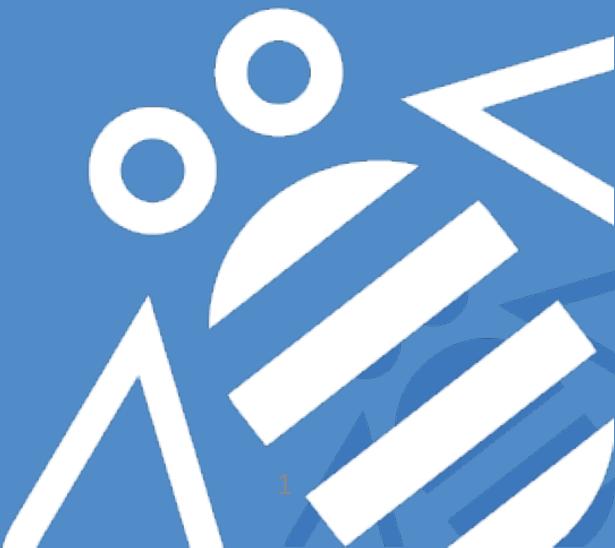
### 第六讲



扫码入群，与讲师互动

## Hyperledger中的共识机制

<http://ibm.biz/opentech-ma>



# “区块链和HyperLedger”系列公开课

- 每周四晚8点档

1. 区块链商用之道
2. HyperLedger项目与社区概览
3. HyperLedger Fabric架构解读
4. ChainCode实战
5. HyperLedger 中的共享账本
6. **HyperLedger中的共识机制**
7. HyperLedger中的隐私与安全
8. HyperLedger应用案例赏析



Q&A环节

扫码入群，与讲师互动

课程Wiki: <http://ibm.biz/opentech-ma>

往期视频: <http://list.youku.com/albumlist/show?id=49106065>



## 夏勇博士

### 个人概述

- IBM 思想领袖级顾问
- IBM大中华区GBS顾问中心负责人
- IBM大中华区GBS区块链中心负责人
- **IBM全球技术研究院院士 (IBM AoT Member)**



### 学术界

- **国际需求工程委员会理事 ( [www.ireb.org](http://www.ireb.org) )** (全球共五位理事中，唯一的华人理事 )
- **复旦大学客座教授**
- **上海交通大学 企业博士生导师**
- **会议共同主席 / 专家评审委员会成员 / 工业论坛主席 :**
  - 2017年: The 29<sup>th</sup> International Conference on Advanced Information Systems Engineering
  - 2016年: The 24<sup>th</sup> IEEE International Requirements Engineering Conference
  - 2015 和 2016 年: The Asia Pacific Requirements Engineering Symposium,
  - 2015年: 第六届CSTQB国际软件测试高峰论坛
  - 2015年: 互联网环境下的需求工程研讨会
- 多年来在国际一流期刊和会议发表有影响力论文近20篇
- **瑞士苏黎世大学经济学院计算机系博士**

### 工业界

- IBM全球技术研究院院士(IBM AoT Member, IBM GBS大中华区仅有的两名院士之一)。现担任IBM GBS顾问中心和区块链中心负责人，主要负责IBM在Fintech/区块链和认知计算方面业务和技术发展。同时，作为业务和技术总架构师和思想领袖级顾问，为摩根大通 (J.P. Morgan Chase), 新加坡星展银行(DBS)及中国工商银行和中国银行等客户提供最高管理层咨询。

之前作为IBM金融解决方案负责人，负责的团队为国际和国内一线银行提供解决方案。客户包括汇丰银行，瑞士联合银行，新澳银行，中国银行，中国建设银行，美亚保险，Suncorp等20多家金融机构。领域包括核心银行，投资银行和财富管理等等。所领导团队每年提供方案的合约总值，最高时接近上亿美元。

- 加入IBM前，在瑞士苏黎世的瑞士信贷银行总部 (Credit Suisse — <https://www.credit-suisse.com/>) “全球银行核心部”，担任总架构师。成功完成投资银行和商业银行的大型项目，包括“结构化产品业务流程”，“证券产品（股票，债券和金融衍生品）的目标系统规划”，“信贷评估系统”，“欧洲重点市场的税务系统”，“银行合规系统架构”等。具体工作包括：(1) 定义应用系统技术路线图，维护和完善公司内部的开发规范，设计流程以及系统构建标准; (2) 设计，评估和决定战略性的解决方案; (3) 企业架构和流程的设计，建模和优化; (4) 金融系统(交易)算法设计和量化分析; (5) 作为技术总负责人，负责管理和协调大型项目以及与银行最高级管理层的沟通。
- 再之前，在丹麦电讯瑞士分公司(北欧最大电信公司之一)，担任高级信息专家

### 工业证书

- 特许金融分析师 – Chartered Financial Analyst (CFA)  
三级候选人



# 议程

- Blockchain - Hyperledger
- Distributed system and its related issues
- Consensus
  - Permissioned (voting) consensus
  - BFT in details
  - Constraints related to (Vanilla) BFT
  - Consensus mechanism in Hyperledger v1.0
- Summary
- Q&A

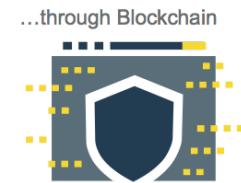


# Blockchain - Hyperledger

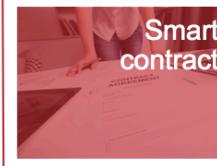
## BLOCKCHAIN EXPLAINED

### [?] What?

For the first time in over a thousand years we have an opportunity to change the mechanisms of recording value exchange...



Append-only distributed system of record shared across business network



Business terms embedded in transaction database & executed with transactions

### [?] Why?



Saves time  
Transaction time from days to near instantaneous



Removes cost  
Overheads and cost intermediaries



Reduces risk  
Tampering, fraud & cyber crime

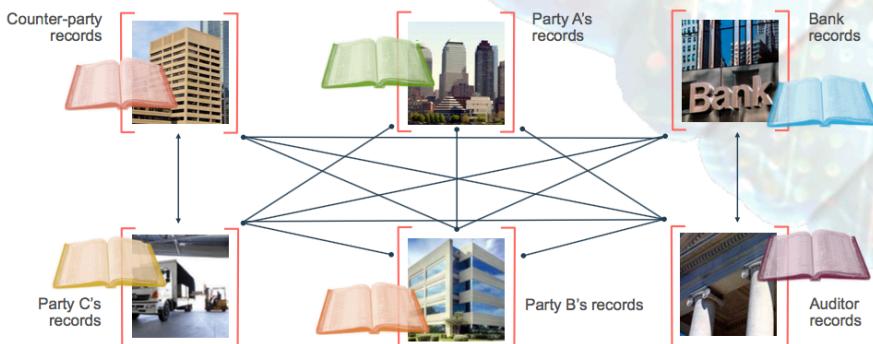


### How?

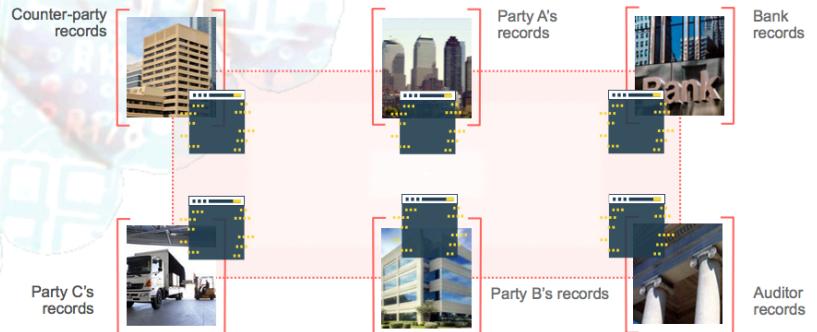
#### Engagement model



### Problem ...

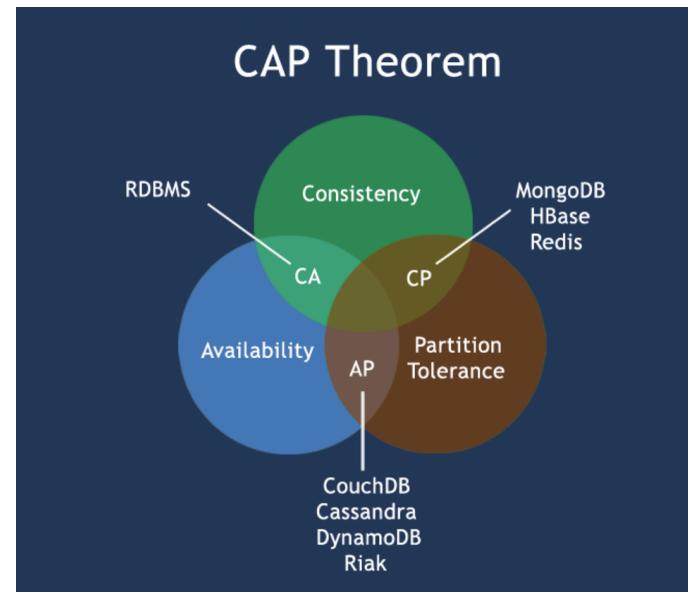


### Solution ...



# Distributed Systems

- CAPM



- Temporal Orders
  - Synchronization
  - ...
- ...

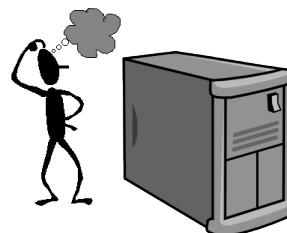


# Permissioned (voting) Consensus

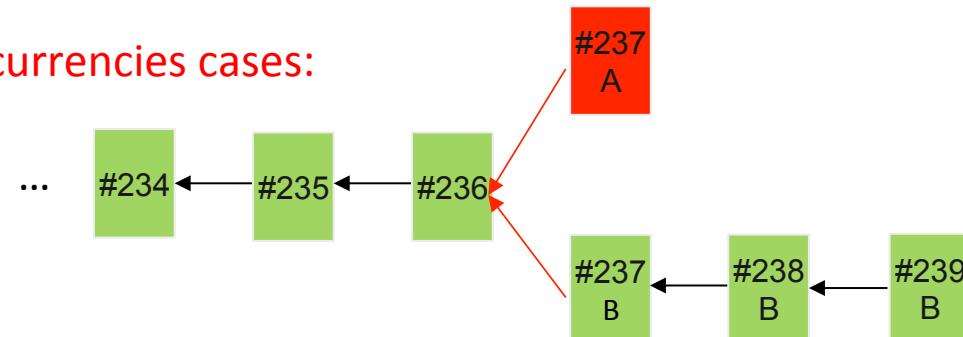
- Classical Distributed Computing protocols (since '80s)
  - Voting based
  - Consensus despite **machine** faults and (temporary) **network** partitions

## What machine faults?

- Crash faults (CFT): A machine simply stops execution and halts
  - Paxos, RAFT, Zookeeper AB,...
- Non-crash (a.k.a. Byzantine) faults (BFT)



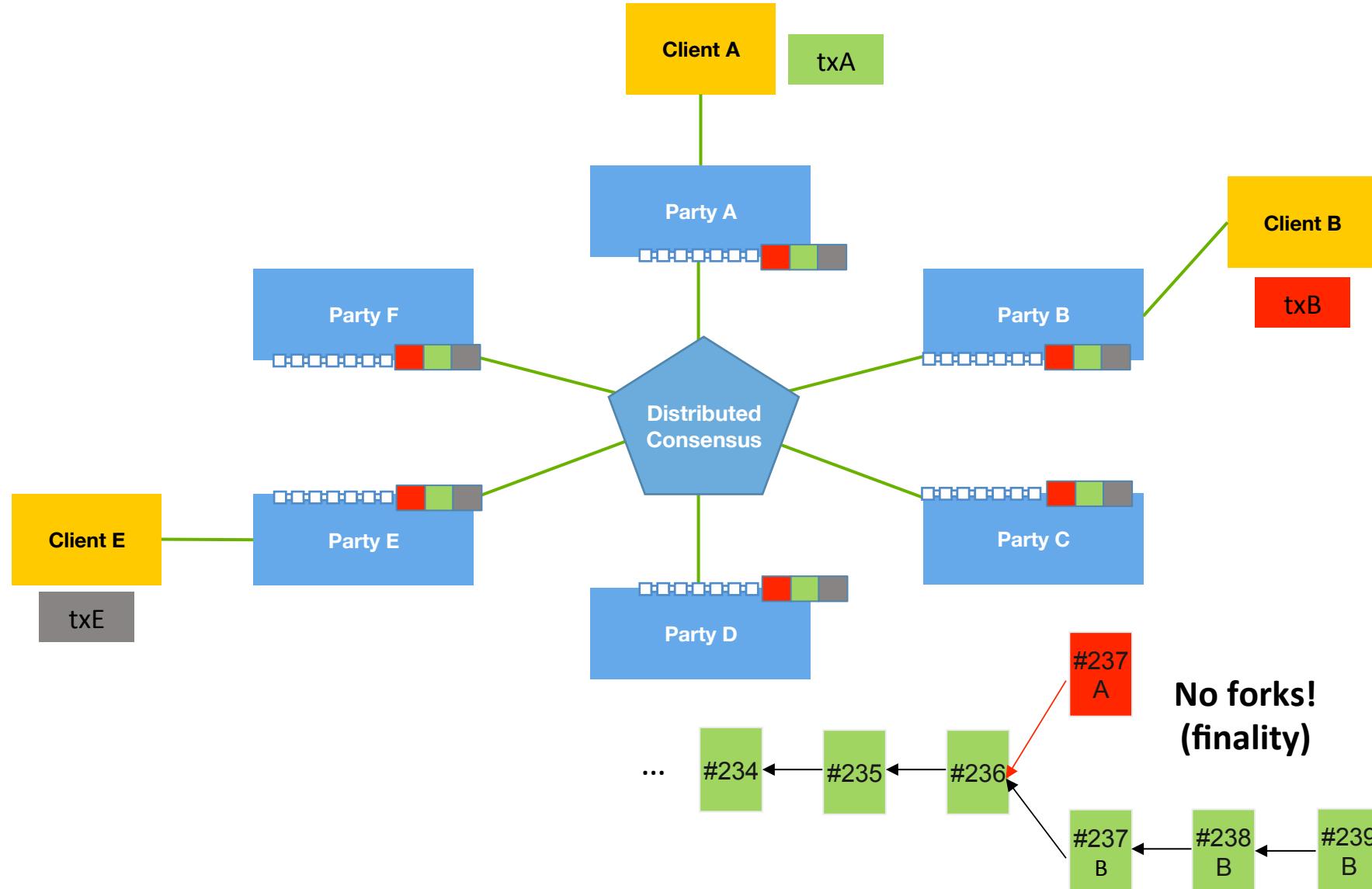
Some cryptocurrencies cases:



No forks!



# Growing Permissioned Blockchains



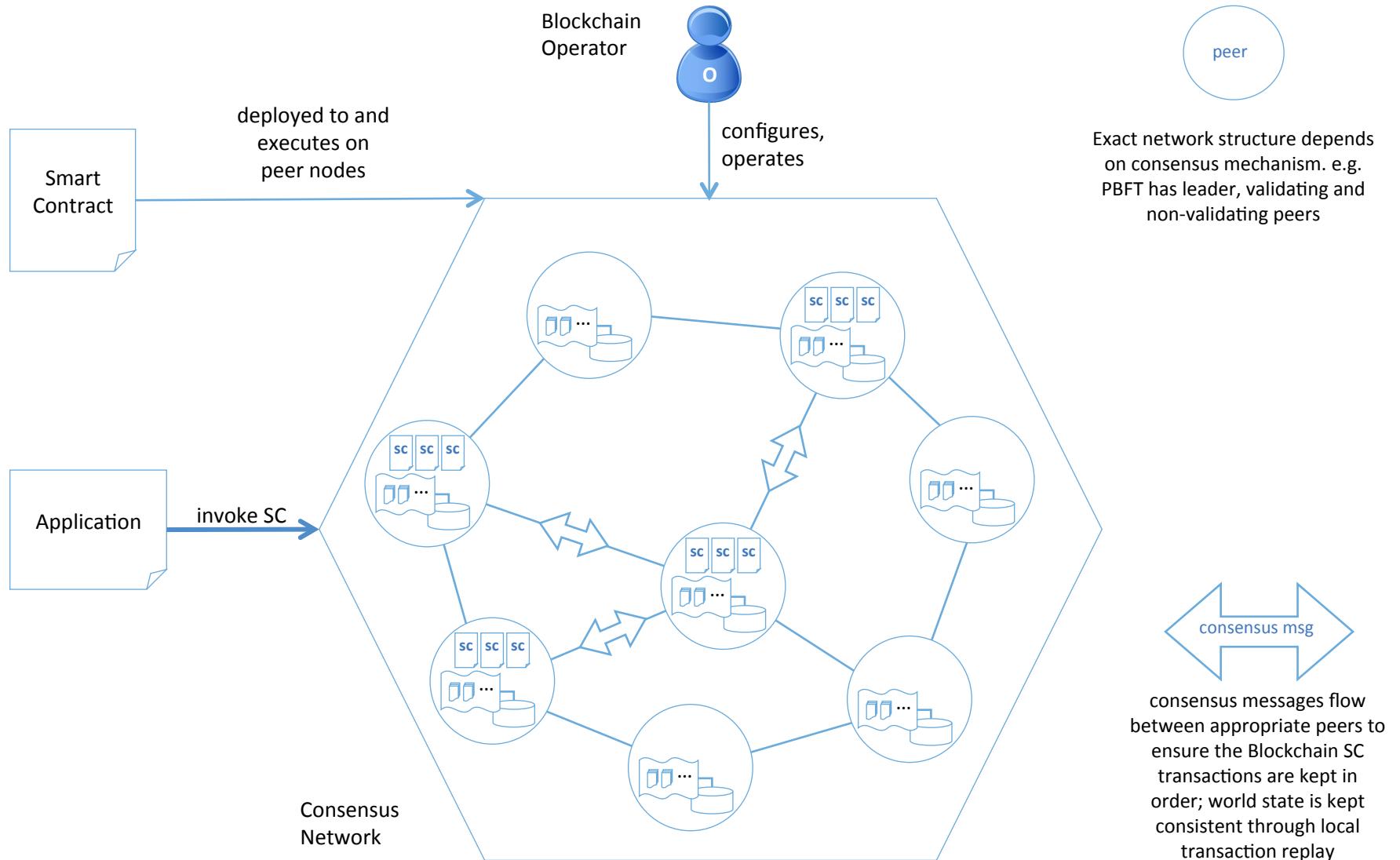
# Permissioned consensus guarantees

- (Vanilla) BFT: Up to  $n/3$  Parties can be **malicious, not follow protocol**
- CFT – Up to  $n/2$  Parties can fail by crashing
- New models – XFT [OSDI 2016]
  - Consensus with up to  $n/2$  **corrupt, Byzantine** Parties (with certain assumptions on network partitions)

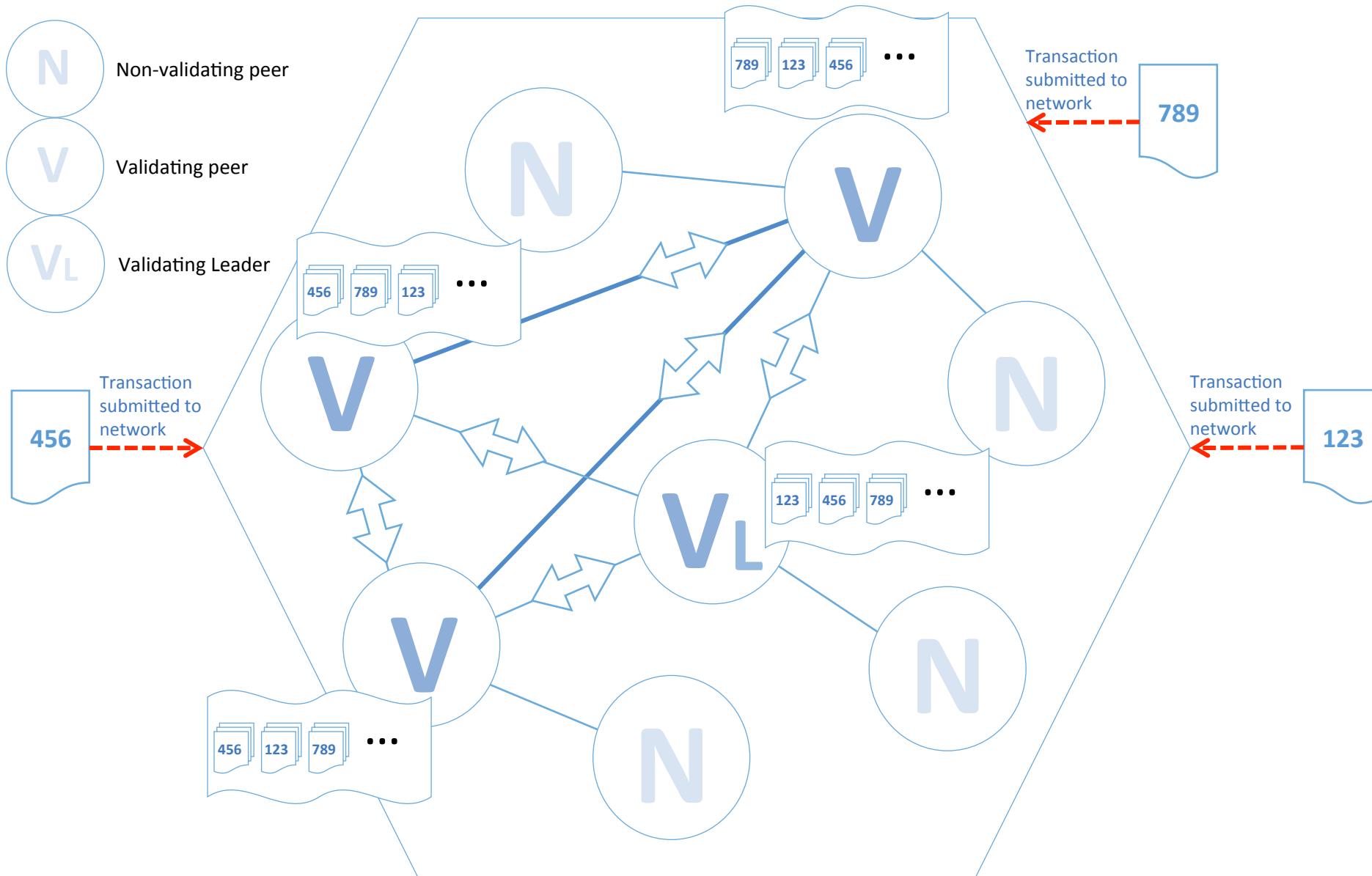
Fault model	CFT	XFT	BFT
Number of Nodes	$2f+1$	$2f+1$	$3f+1$
Tolerating Byzantine Nodes	no	yes	yes
Performance	Good	Practically as good as CFT	Worse than CFT

# BFT Consensus

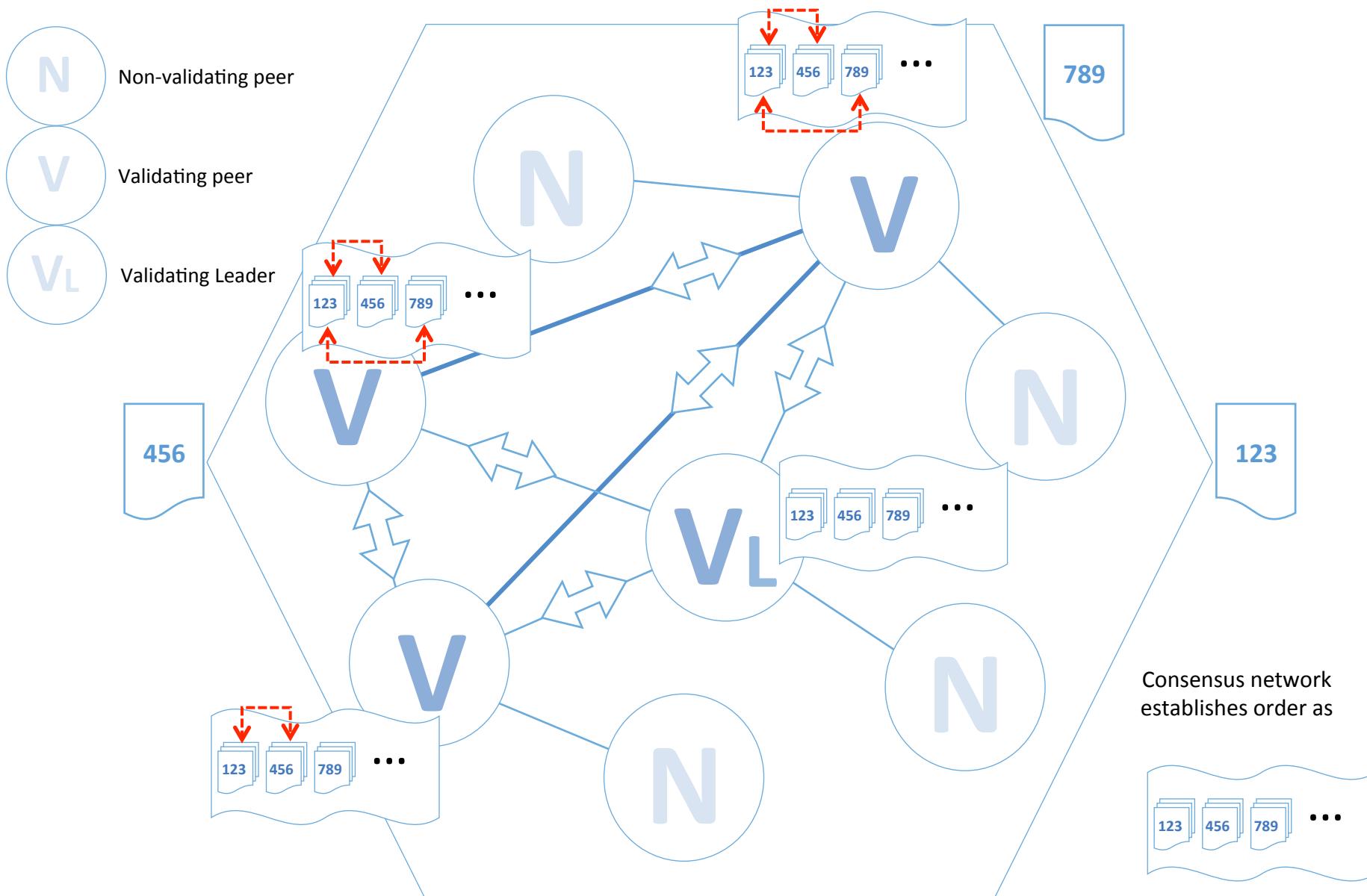
(example of PBFT [TOCS2002])



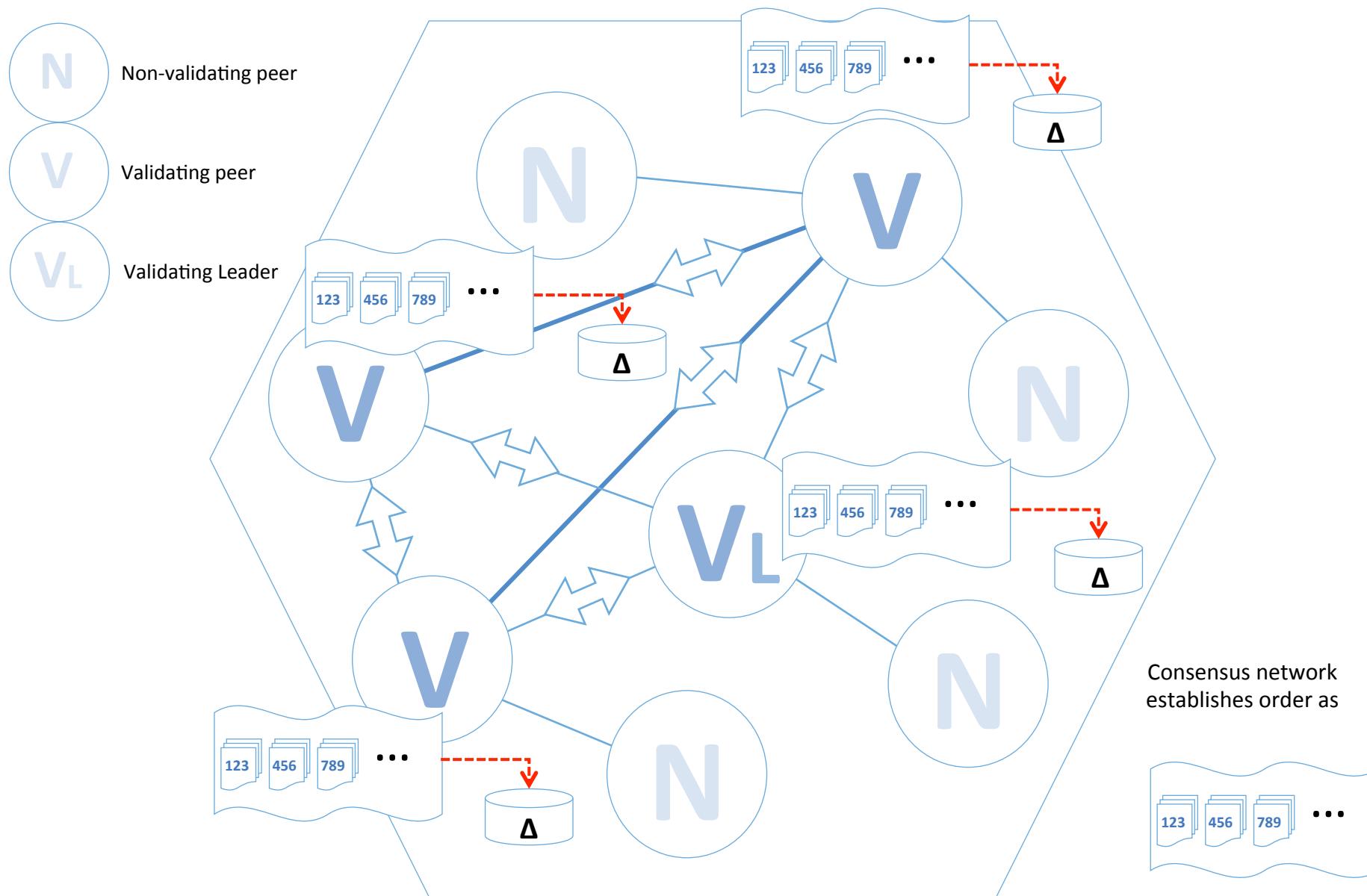
# How a PBFT Network Works (1/4) – Submission



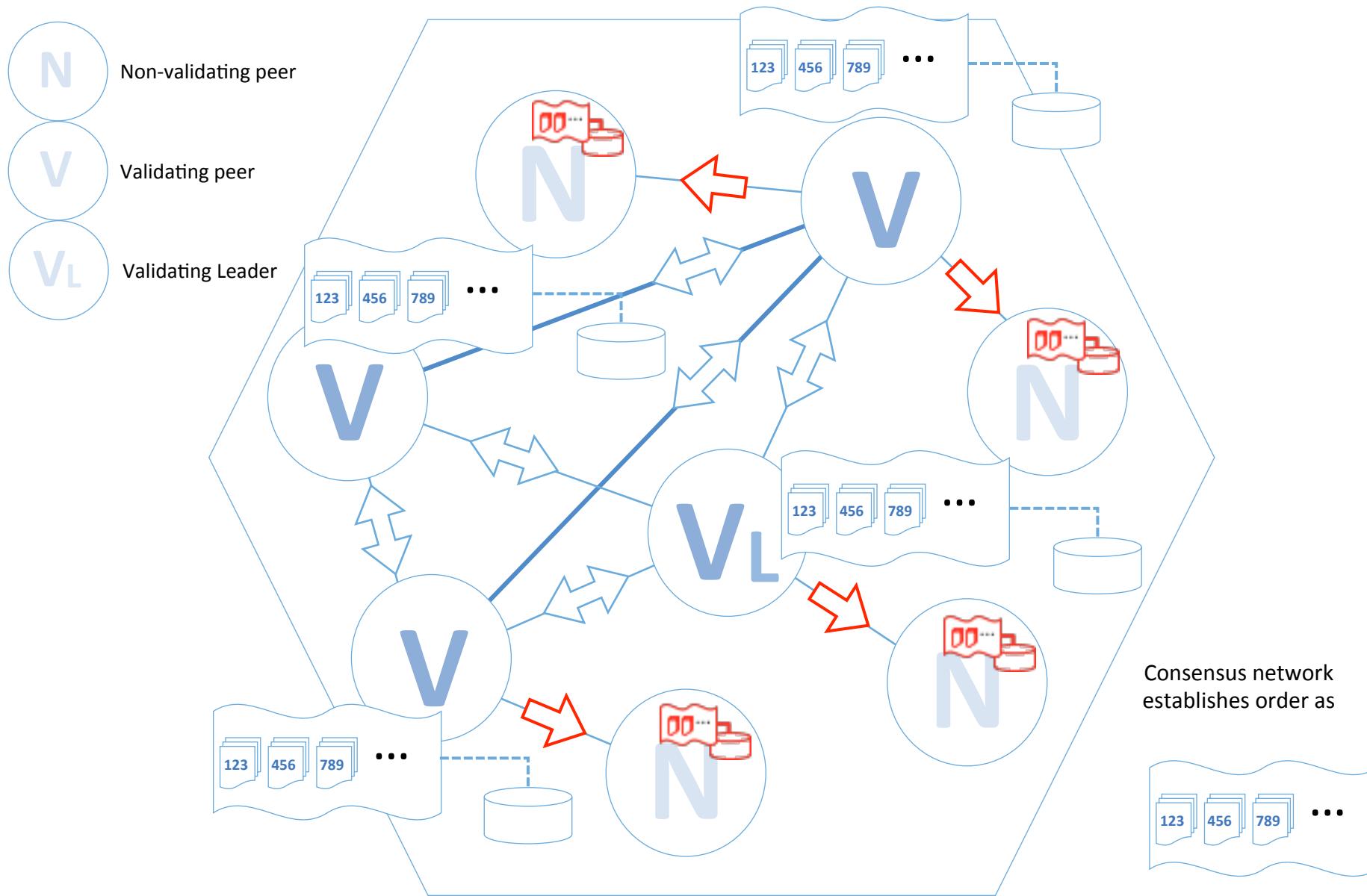
# How a PBFT Network Works (2/4) – Ordering



# How a PBFT Network Works (3/4) – Execution

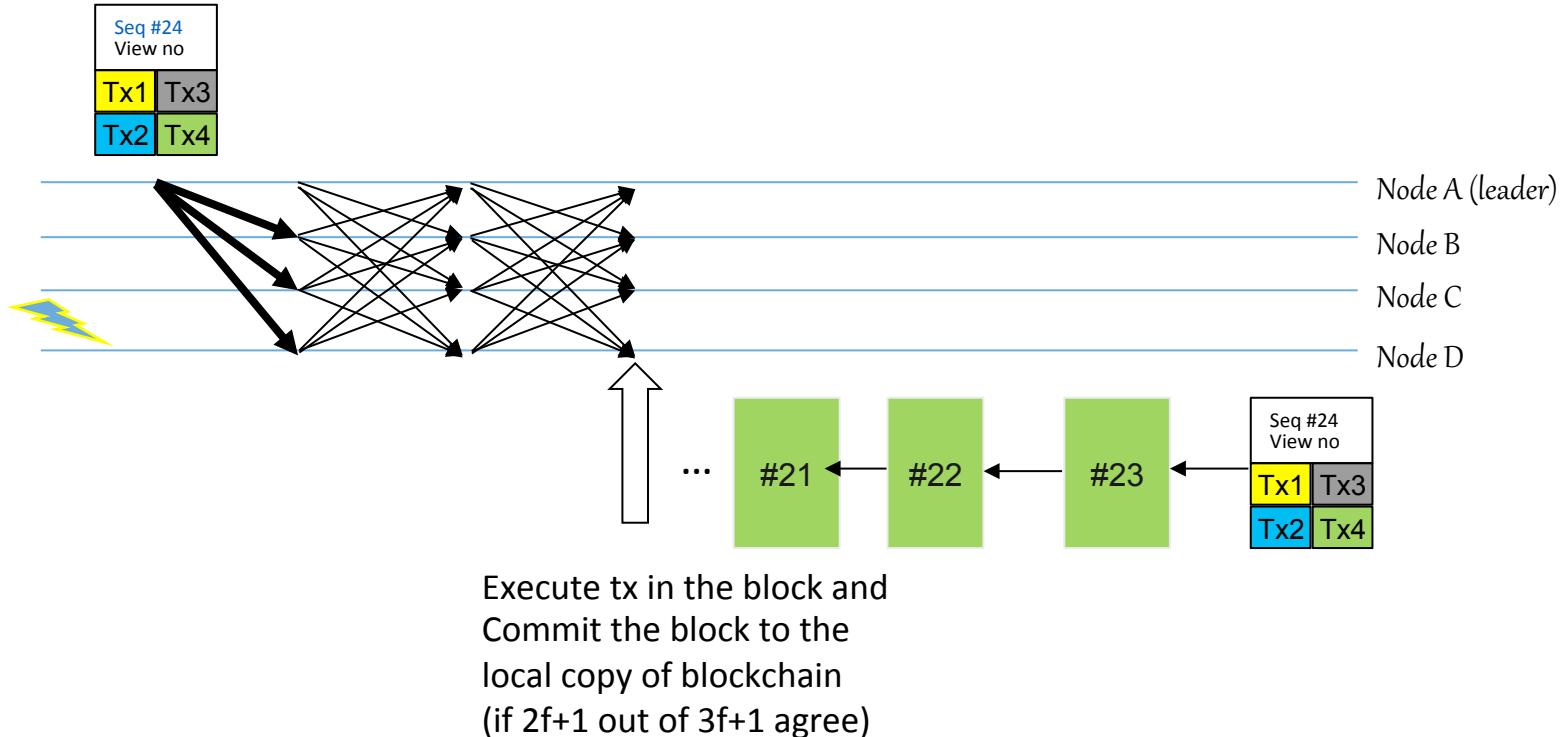


# How a PBFT Network Works (4/4) – Distribution



# BFT Consensus

(example of PBFT [TOCS2002])



Many other things burden the implementation (it is not simple as it might look)

- Non-deterministic value
- Leader election
- State transfer (new, slow Party)
- Reconfiguration

# Nodes and roles

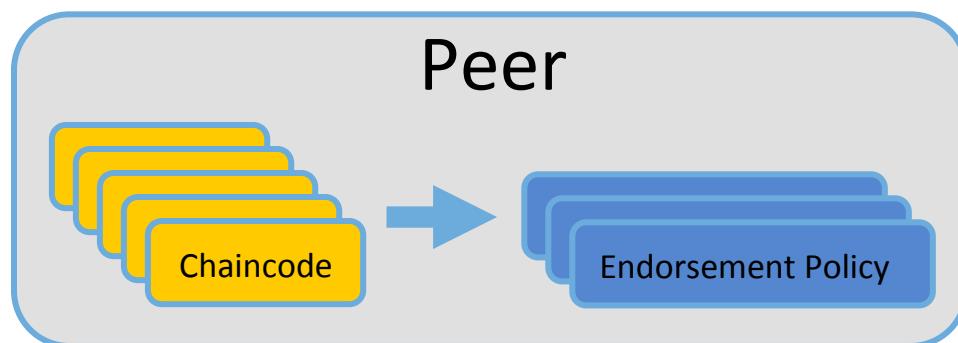
	<b>Peer:</b> Commits transactions, maintains ledger and state
	<b>Endorsing peer:</b> Specialised peer that receives a transaction proposal for endorsement, responds granting or denying endorsement
	<b>Ordering peer:</b> Approves the inclusion of transaction blocks into the ledger and communicates with peer and endorsing peer nodes



# Endorsement Policies

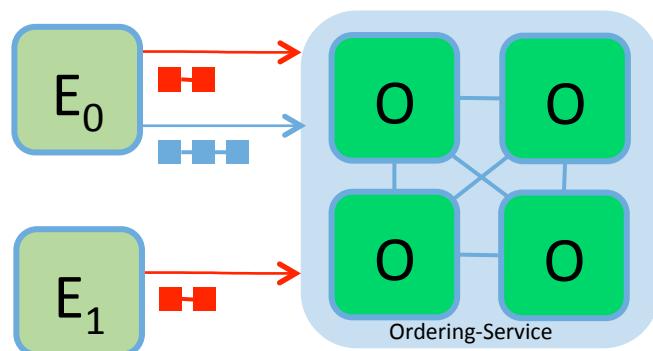
An endorsement policy describes the conditions by which a transaction can be endorsed. A transaction can only be considered valid if it has been endorsed according to the policy.

- Peers maintain a set of endorsement policies
- An endorsement policy is specified on deployment of chaincode



# Channels

Nodes send/receive messages to the ordering-service via channels.

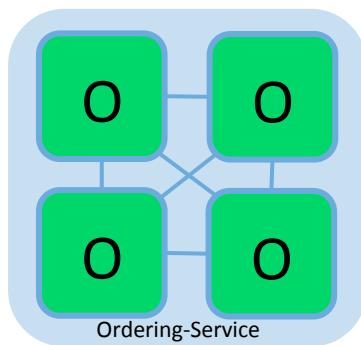


- Enables chaincode privacy
  - Chaincode deployed to certain nodes
- Messages partitioned into separate channels
  - Transactions stored depending on node and channel
- Nodes can connect to one or more channels



# Ordering Services

The ordering service packages transactions into blocks to be delivered to peers. Communication with the service is via channels.



Different configuration options for the ordering service include:

- **SOLO**
  - Single node for development
- **Kafka / Zookeeper**
  - 1:n nodes providing Crash Fault Tolerance
  - Odd number of nodes recommended
- **SBFT**
  - 1:n nodes providing Byzantine Fault Tolerance



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