# BREADTH FIRST SEARCH

**BFS** 

### Breadth-first search

- What is it good for?
- We visit every vertex exactly once
- We visit the neighbours then the neighbours of these new vertices and so on
- Running time complexity: O(V+E)
- Memory complexity is not good: we have to store lots of references
- Thats why DFS is usually preferred
- BUT it constructs a shortest path: Dijkstra algorithm does a BFS if all the edge weights are equal to 1

### Breadth-first search

#### bfs(vertex)

Queue queue vertex set visited true queue.enqueue(vertex)

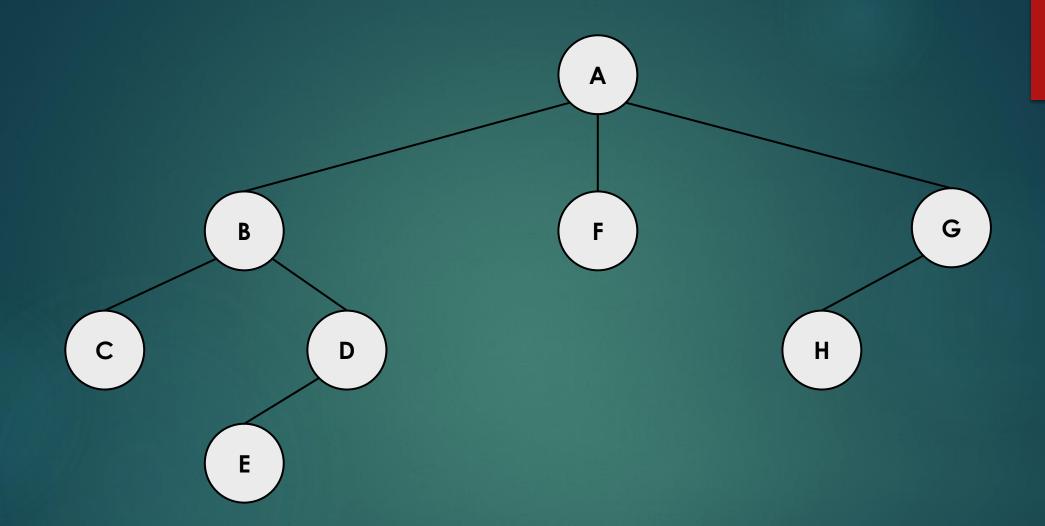
while queue not empty actual = queue.dequeue()

for v in actual neighbours
if v is not visited
v set visited true
queue.enqueue(v)

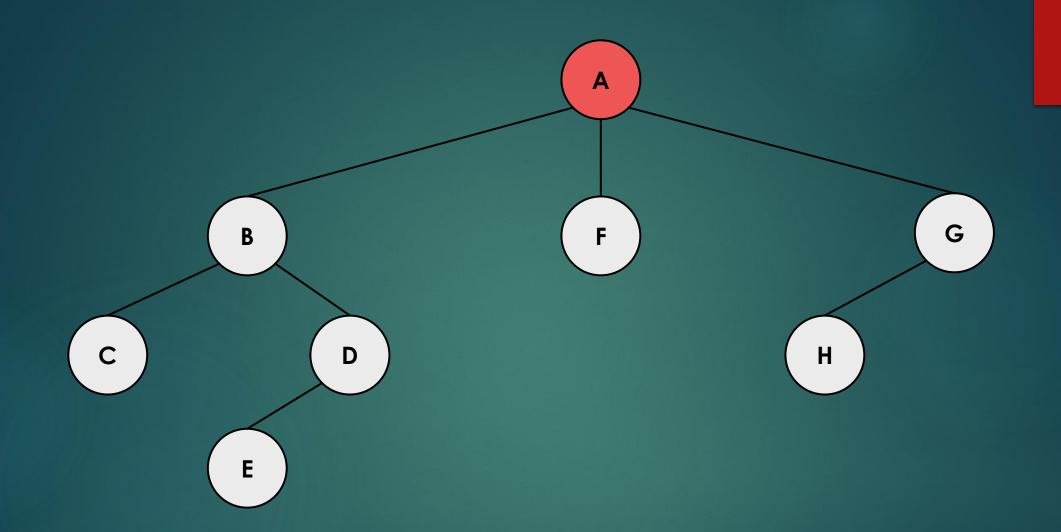
We have an empty queue at the beginning and we keep checking whether we have visited the given node or not

~ keep iterating until queue is not empty

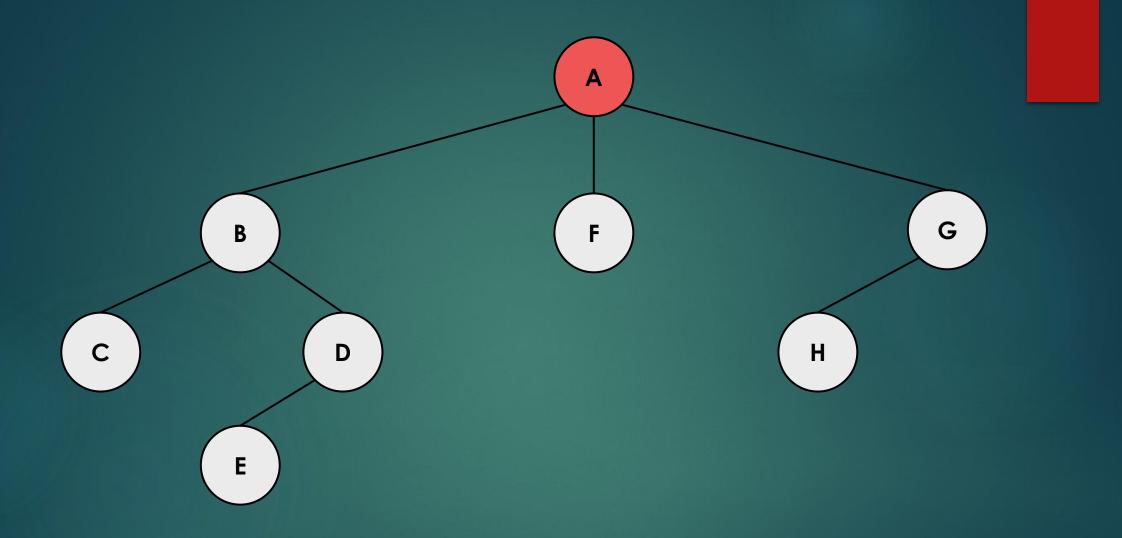
**ITERATION** 



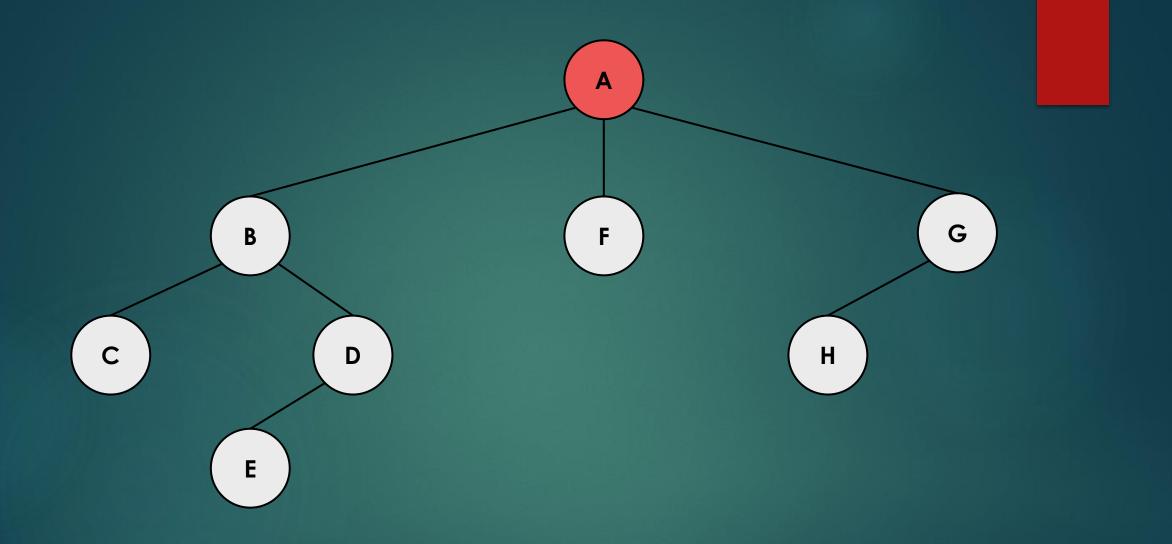
Queue:{}



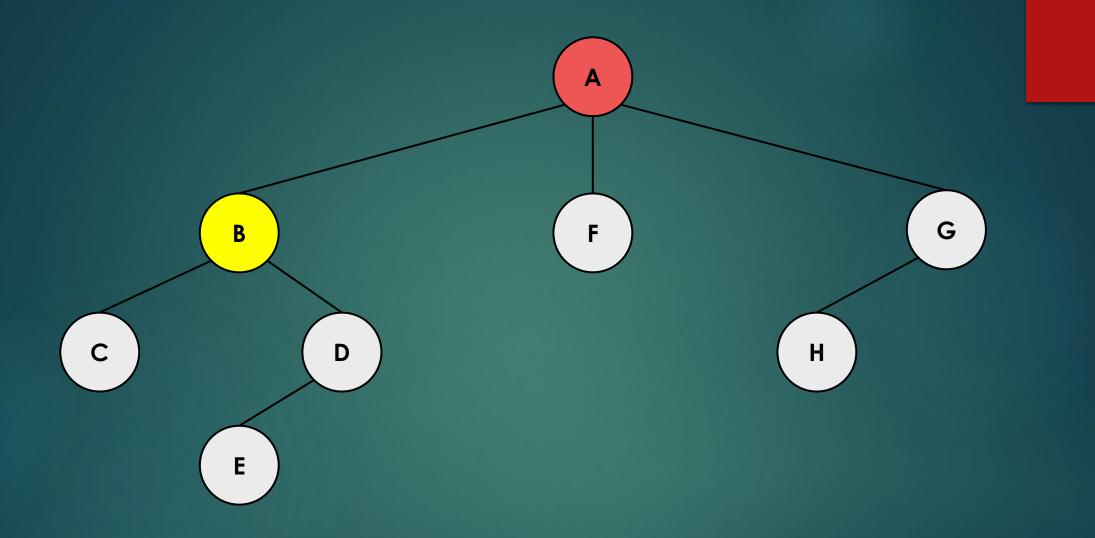
Queue: { A }



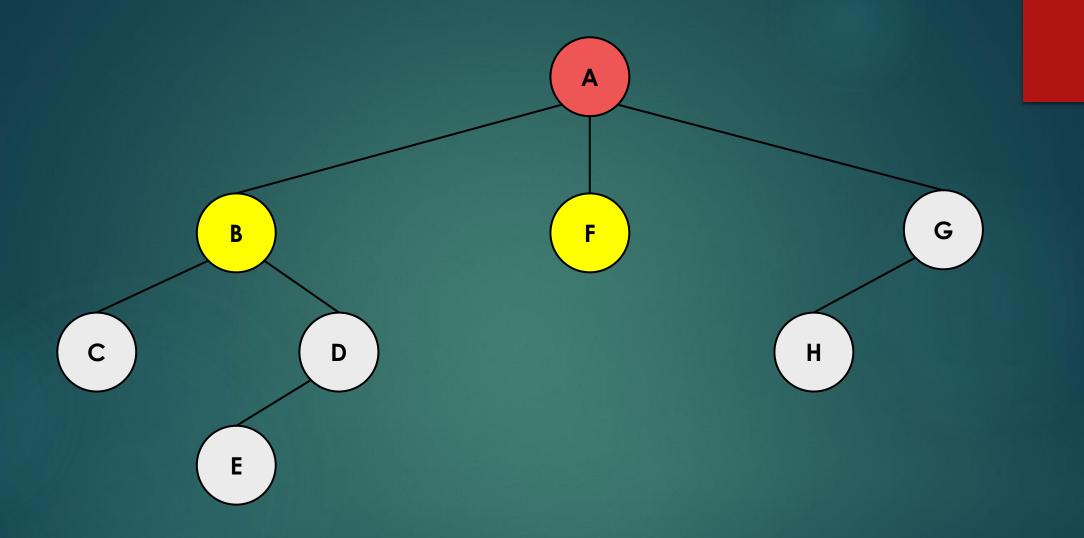
Queue: { A } dequeue node A to be able to process its



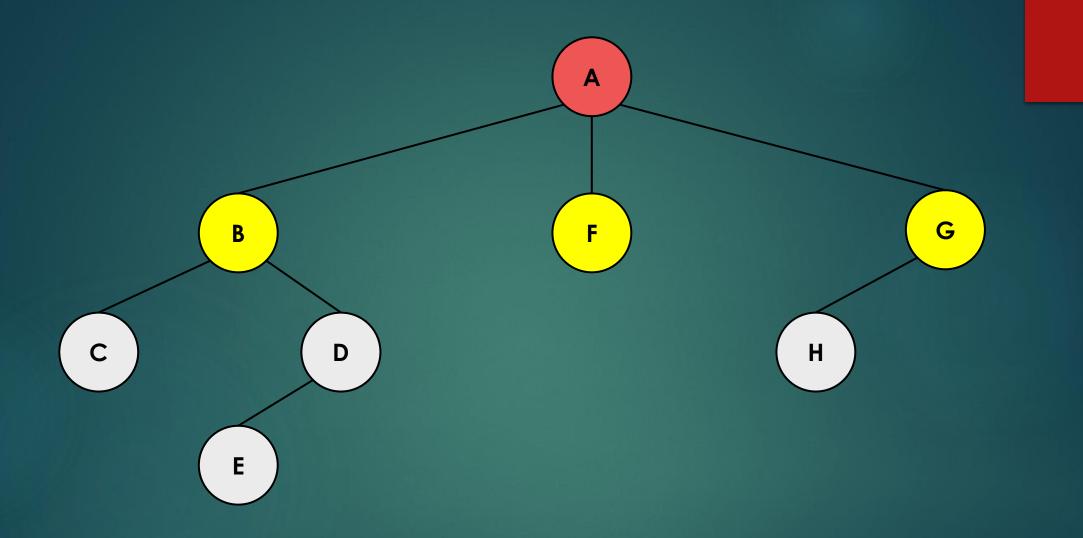
Queue: { } dequeue node A to be able to process its neighbours > visit all the children, if a node is unvisited then enqueue it !!!



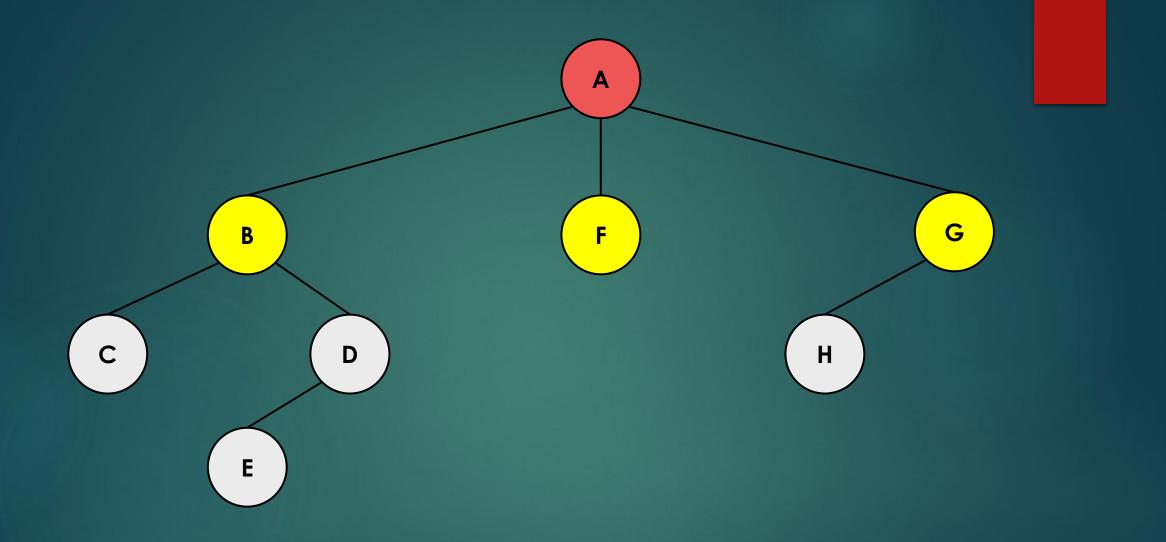
Queue: { B } node B is not visited > put it to the queue



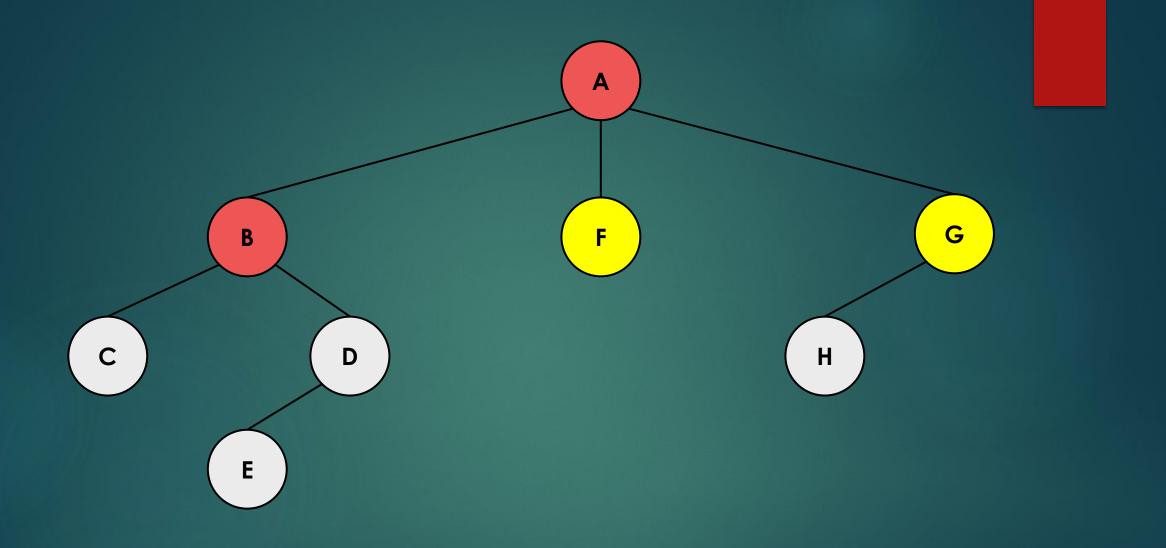
Queue: { F B } node F is not visited > put it to the queue



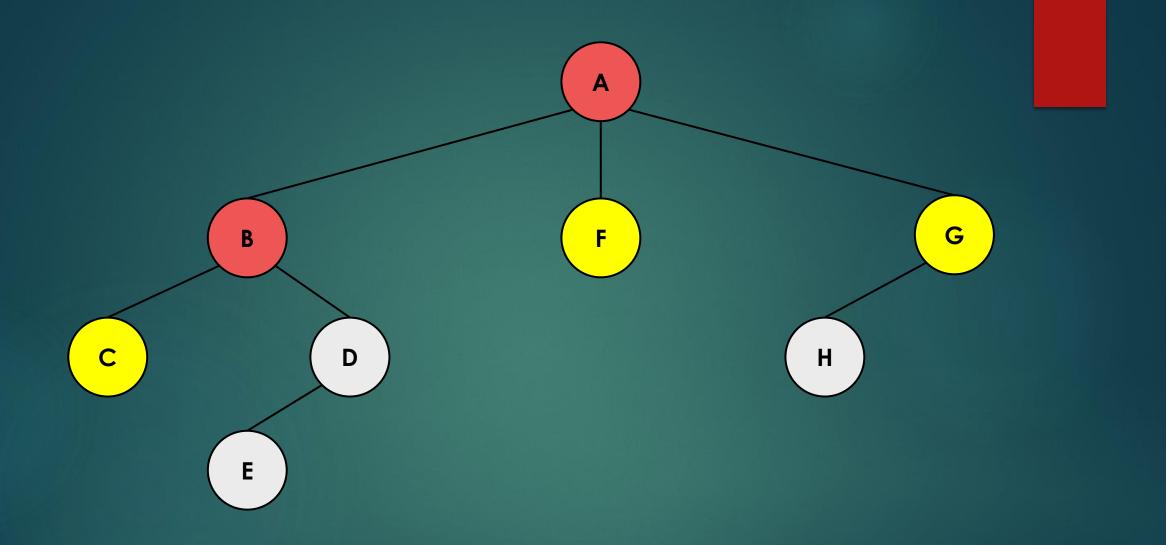
Queue: { G F B } node G is not visited > put it to the queue



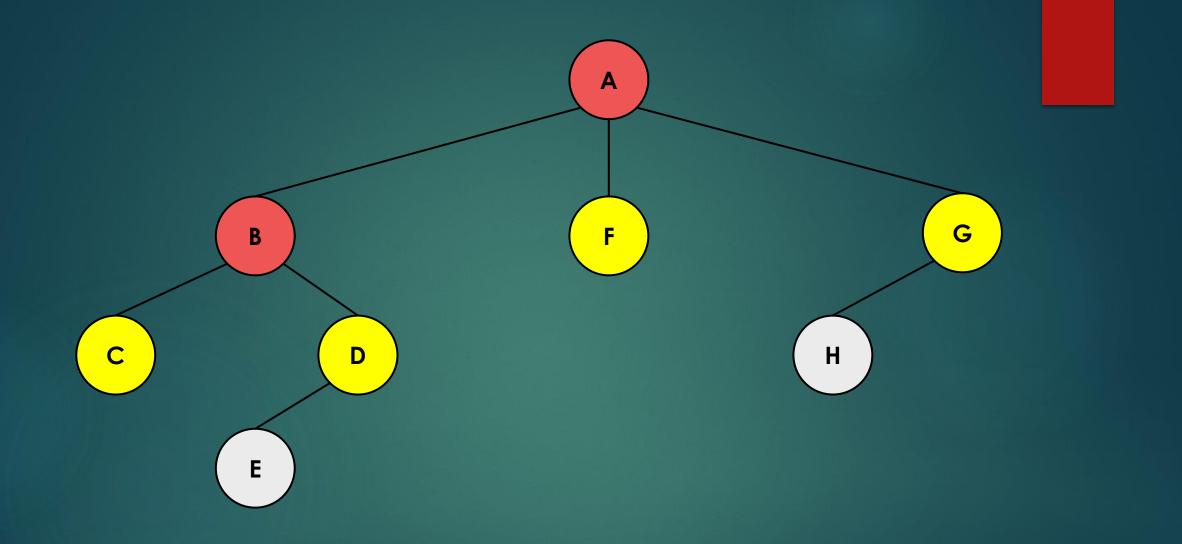
Queue: { G F B } dequeue the next node > it is B and visit its children + put them into the queue if necessary



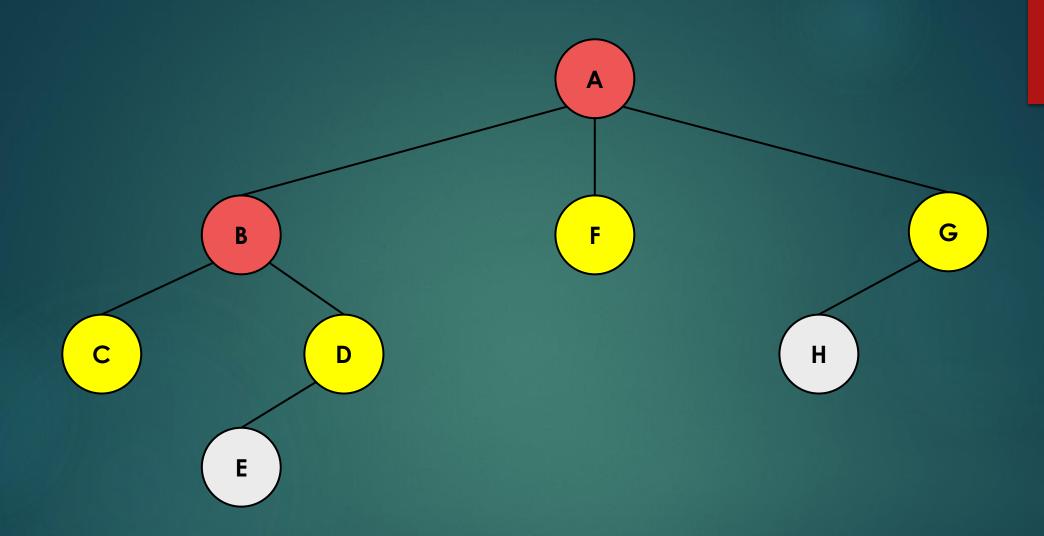
Queue: { G F } dequeue the next node > it is G and visit its children + put them into the queue if necessary



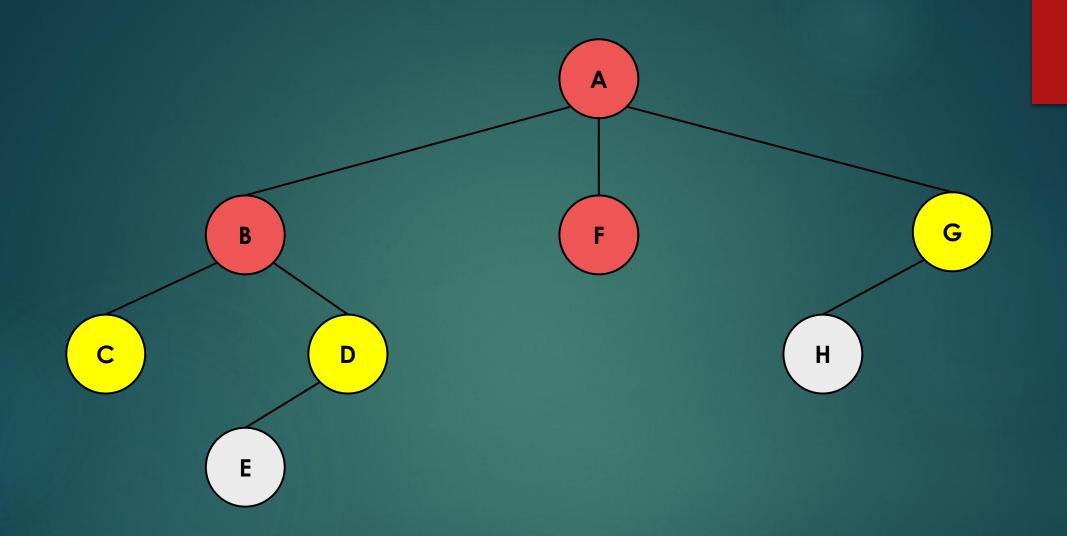
Queue: { C G F } dequeue the next node > it is G and visit its children + put them into the queue if necessary



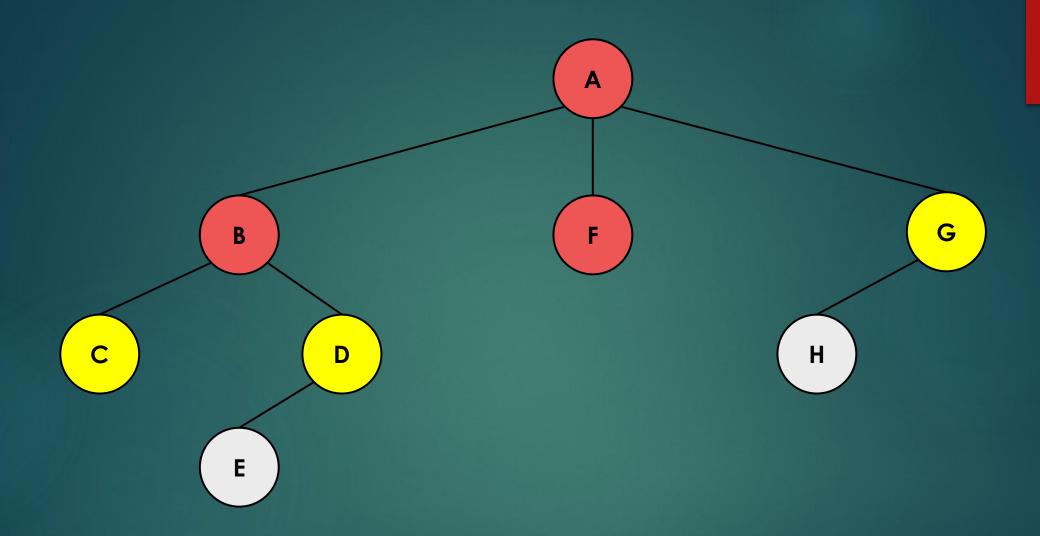
Queue: { D C G F } dequeue the next node > it is G and visit its children + put them into the queue if necessary



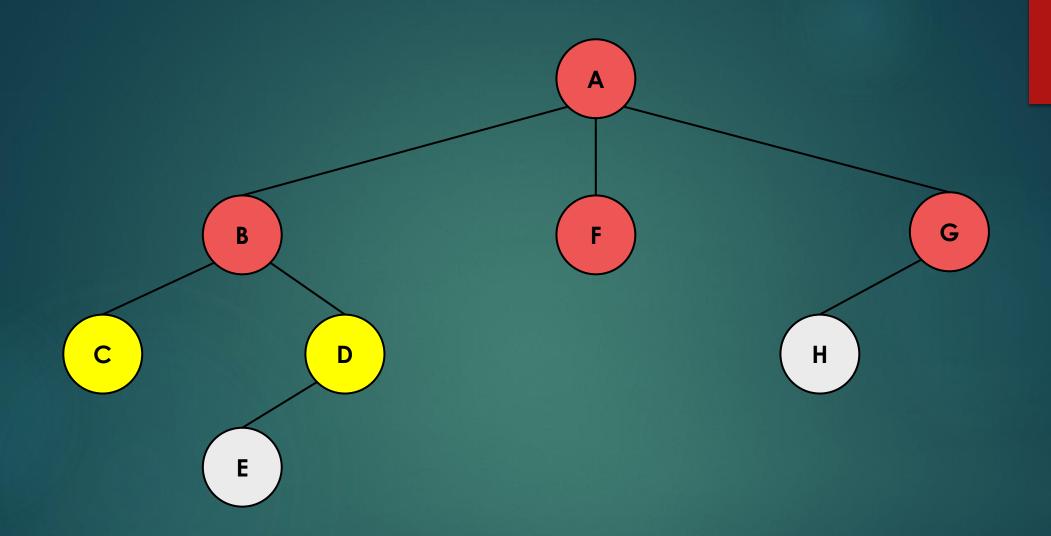
Queue: { D C G F }



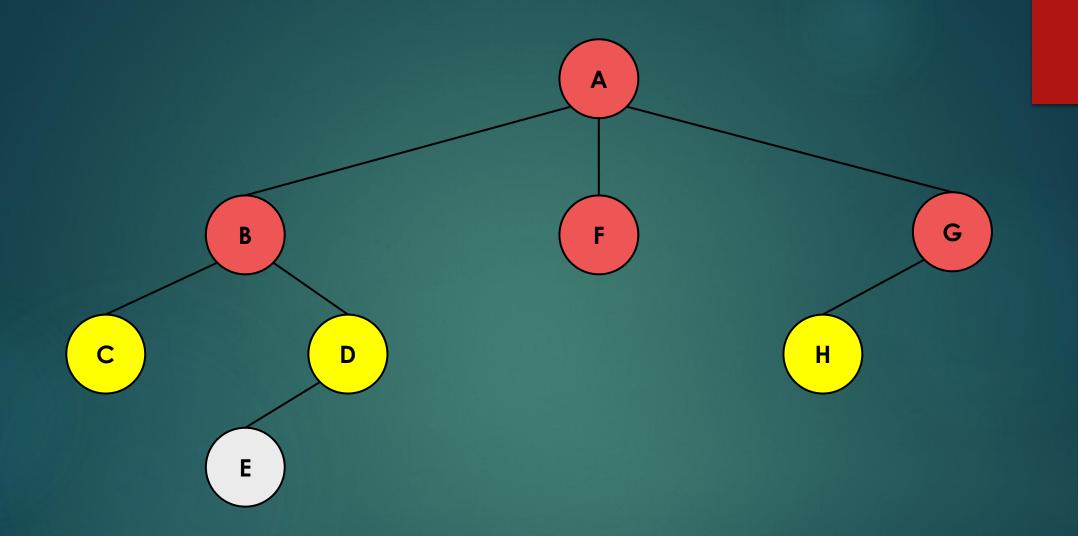
Queue: { D C G }



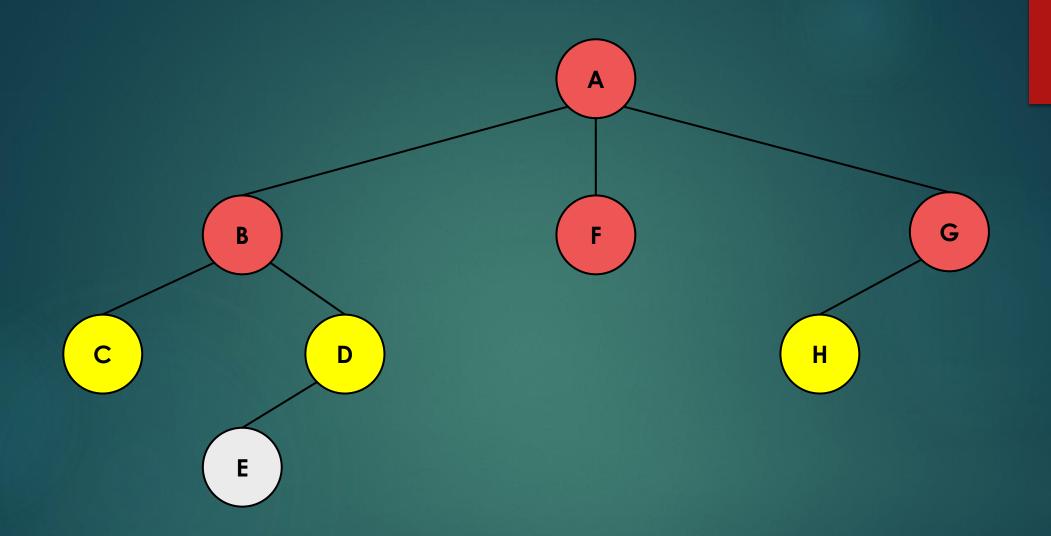
Queue: { D C G }



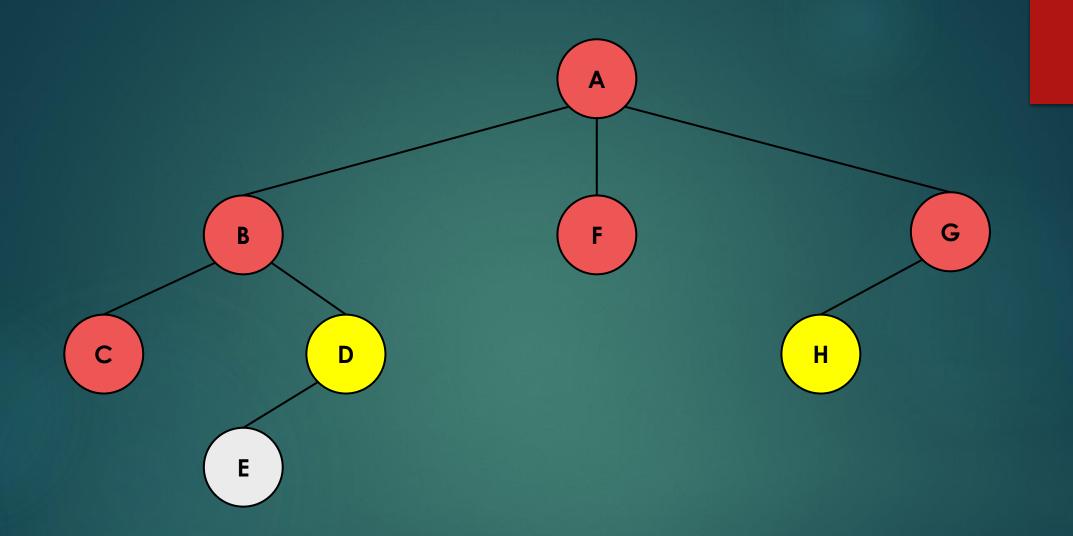
Queue: { D C }



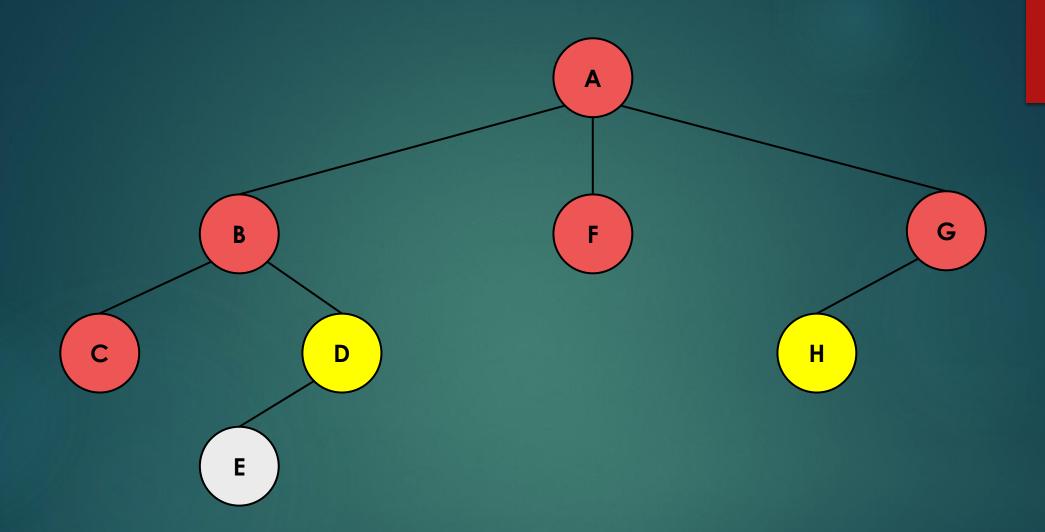
Queue: { H D C }



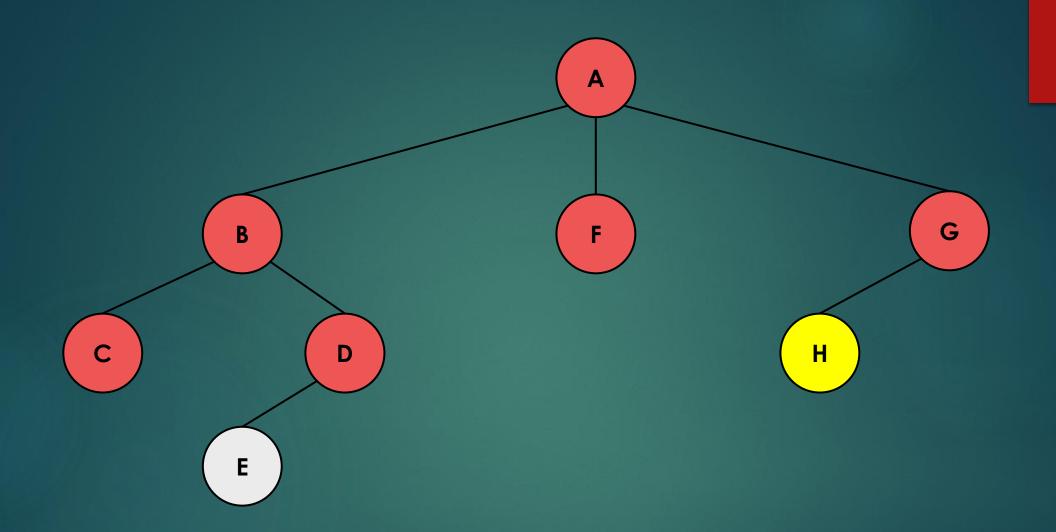
Queue: { H D C }



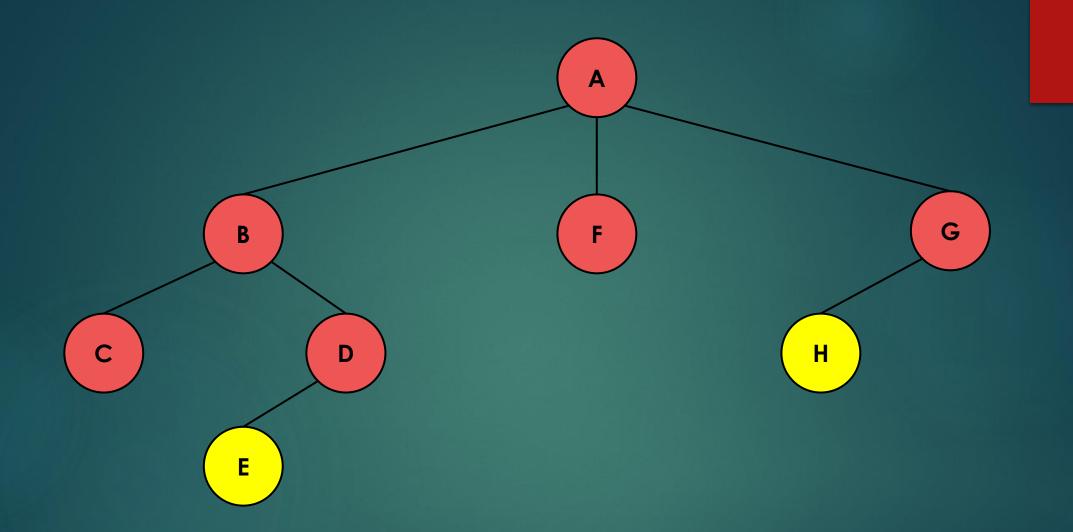
Queue: { H D }



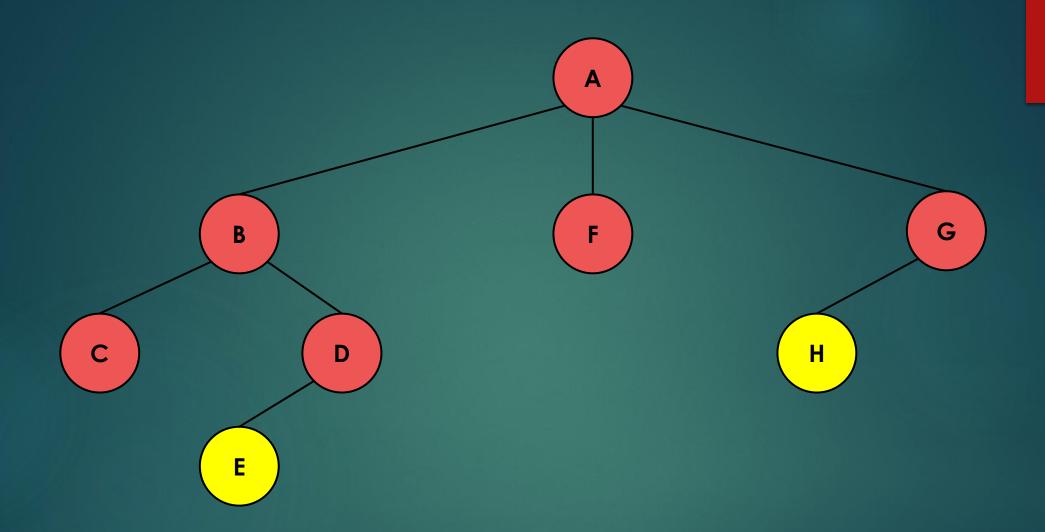
Queue: { H D }



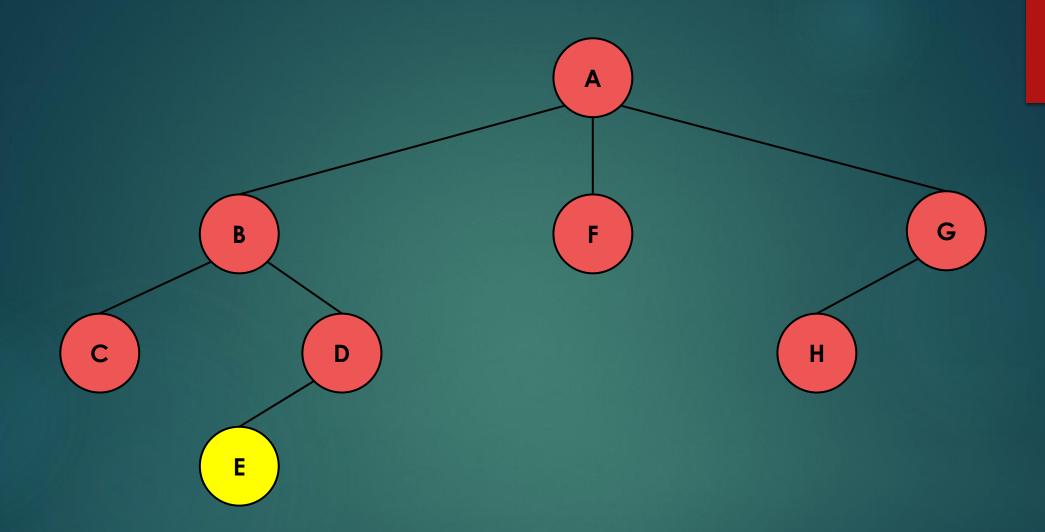
Queue: { H }



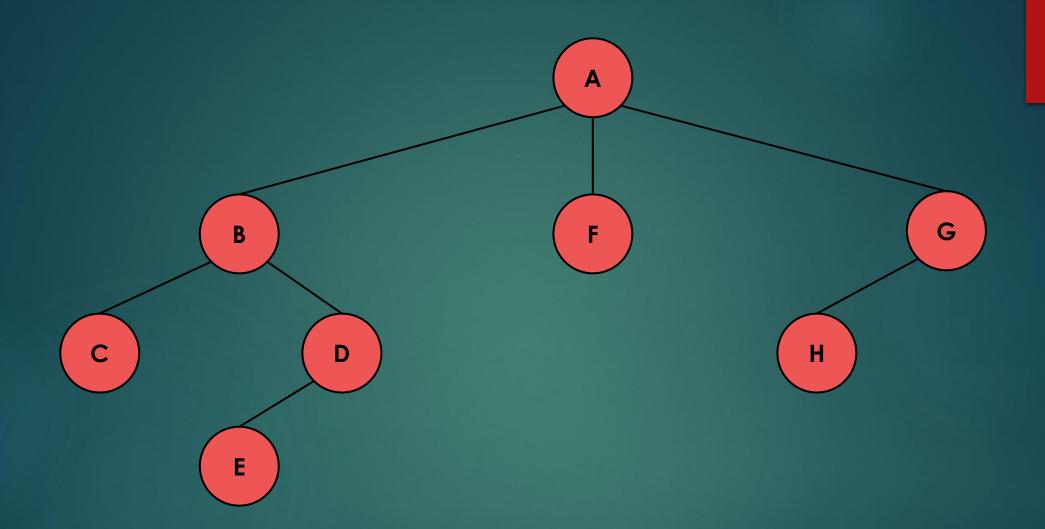
Queue: { E H }



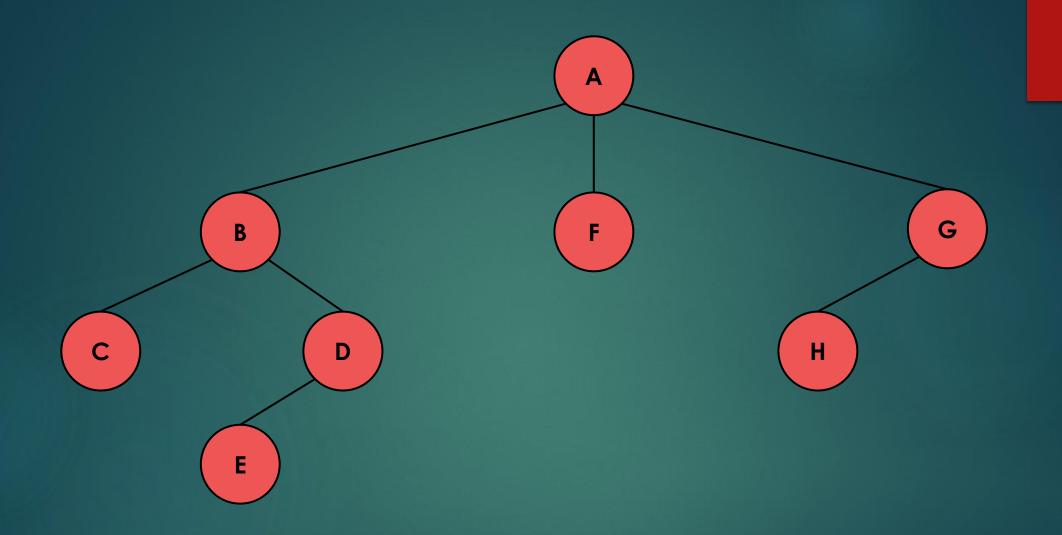
Queue: { E H }



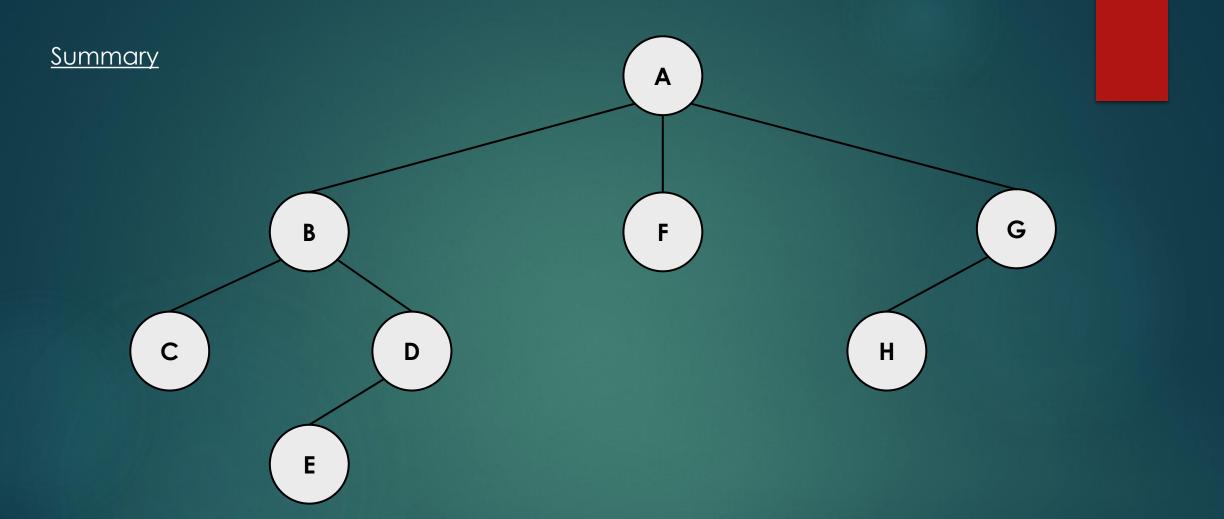
Queue: { E }

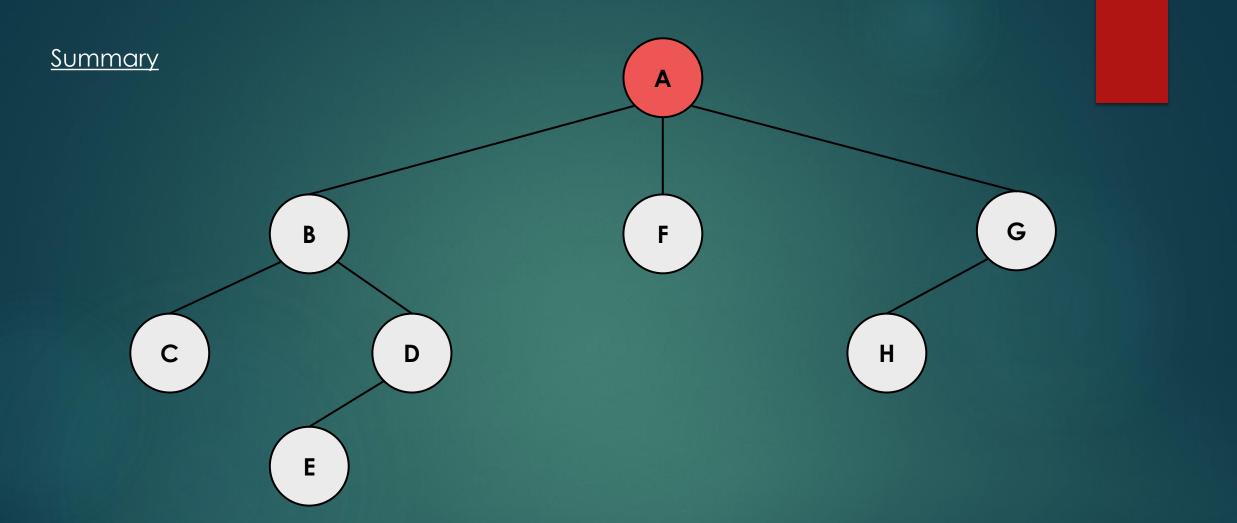


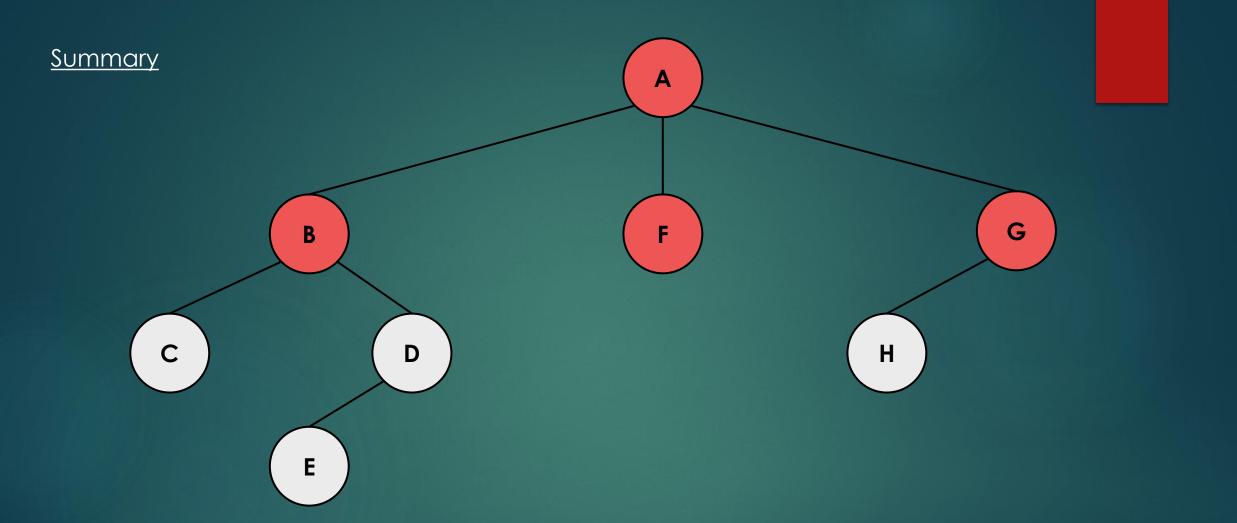
Queue: { E}

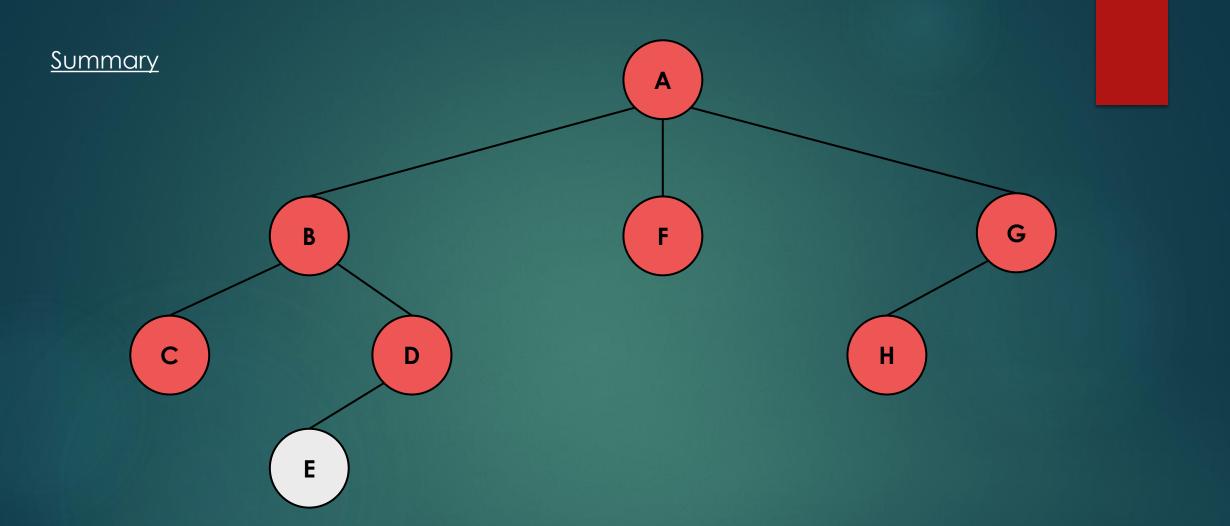


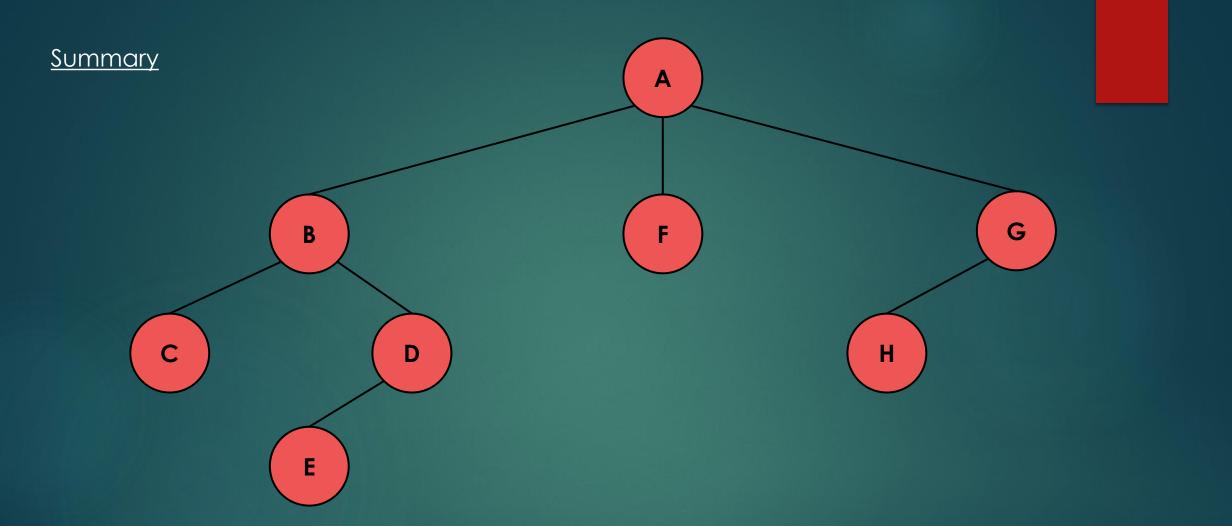
Queue: { } the queue is empty → FINISHED!!!











## **Applications**

- ▶ In artificial intelligence / machine learning it can prove to be very important: robots can discover the surrounding more easily with BFS than DFS
- It is also very important in maximum flow: Edmonds-Karp algorithm uses BFS for finding augmenting paths
- ► Cheyen's algorithm in garbage collection → it help to maintain active references on the heap memory
- It uses BFS to detect all the references on the heap
- ▶ Serialization / deserialization of a tree like structure ( for example when order does matter ) → it allows the tree to be reconstructed in an efficient manner !!!

# DEPTH FIRST SEARCH

DFS

## Depth-first search

- Depth-first search is a widely used graph traversal algorithm besides breadth-first search
- It was investigated as strategy for solving mazes by Trémaux in the 19th century
- It explores as far as possible along each branch before backtracking // BFS was a layer-by-layer algorithm
- Time complexity of traversing a graph with DFS: O(V+E)
- Memory complexity: a bit better than that of BFS !!!

#### Depth-first search

dfs(vertex)

vertex set visited true print vertex

for v in vertex neighbours
if v is not visited
dfs(v)

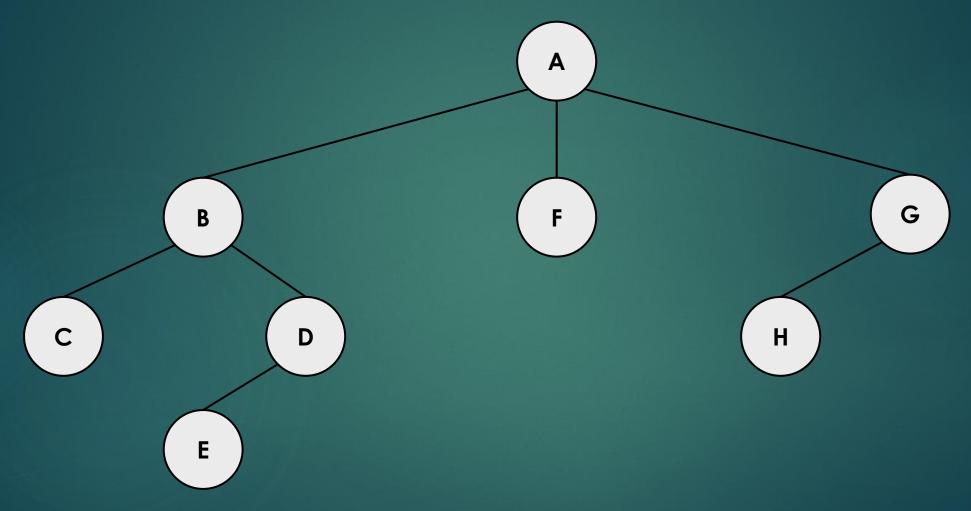
dfs(vertex)

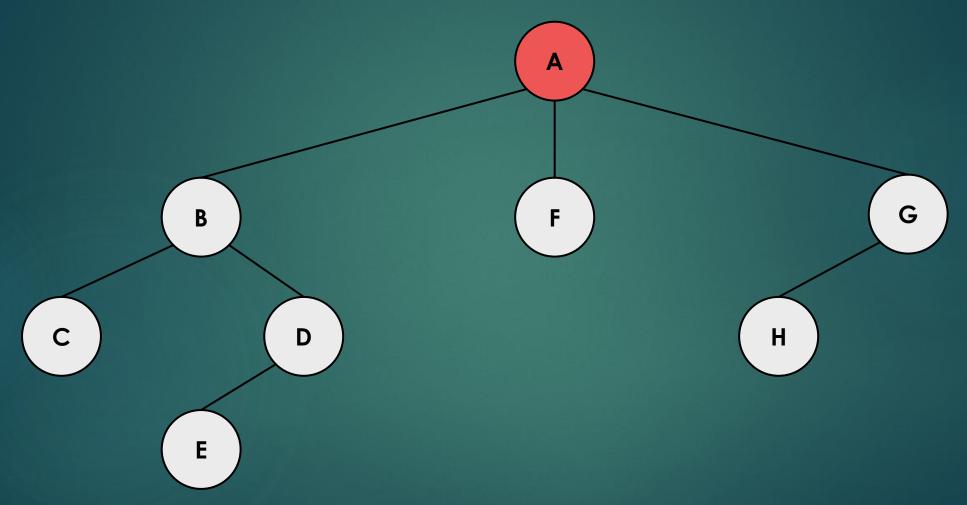
Stack stack vertex set visited true stack.push(vertex)

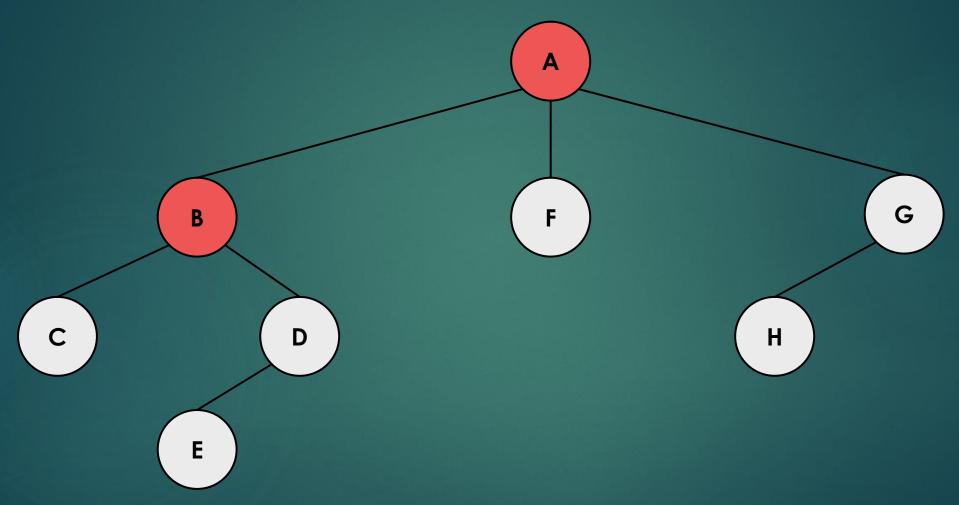
while stack not empty actual = stack.pop()

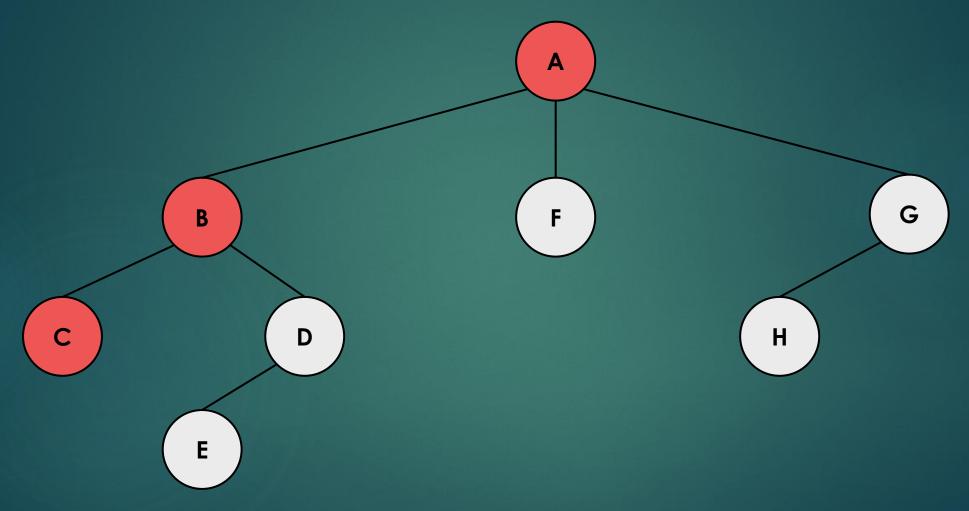
for v in actual neighbours
if v is not visited
v set visited true
stack.push(v)

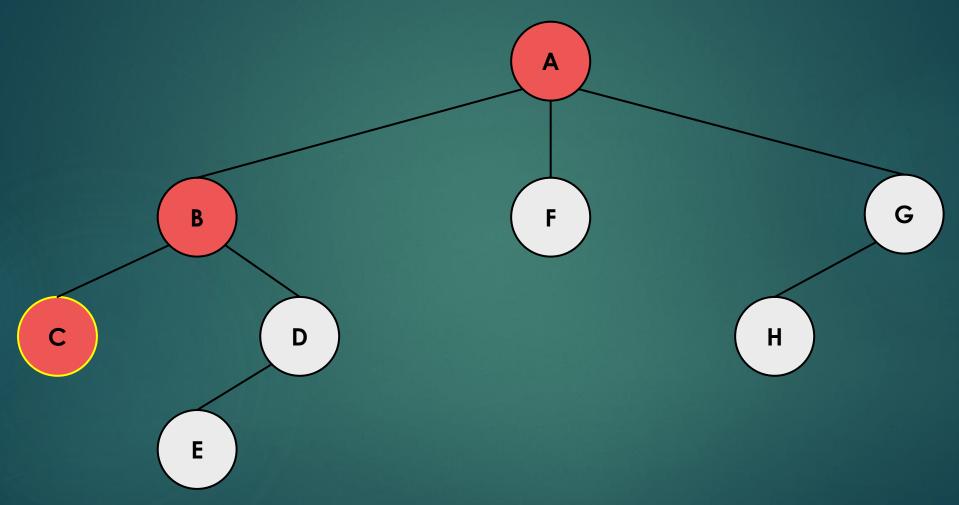
RECURSION ITERATION

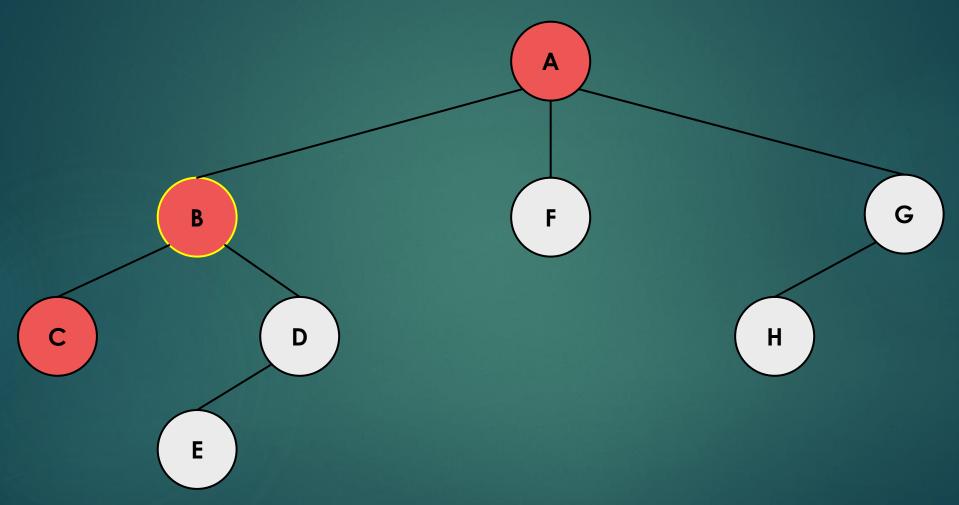


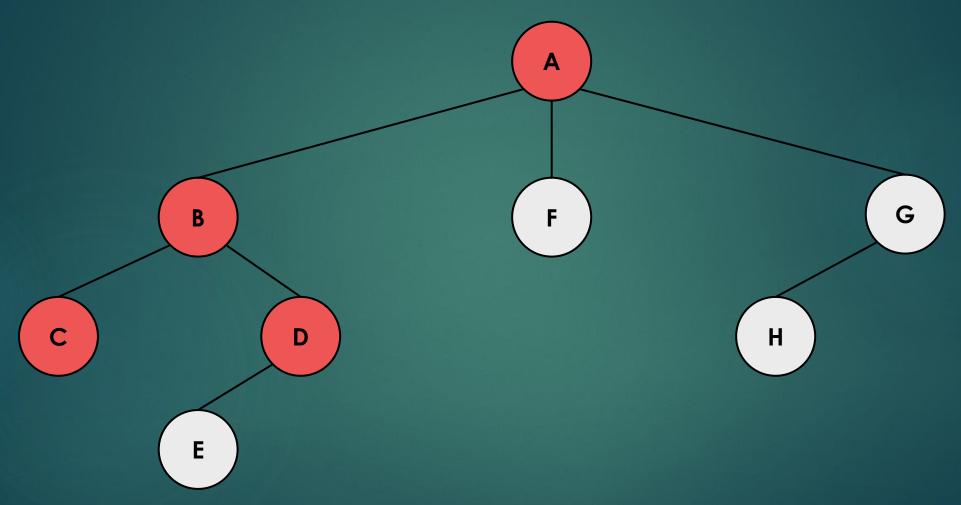


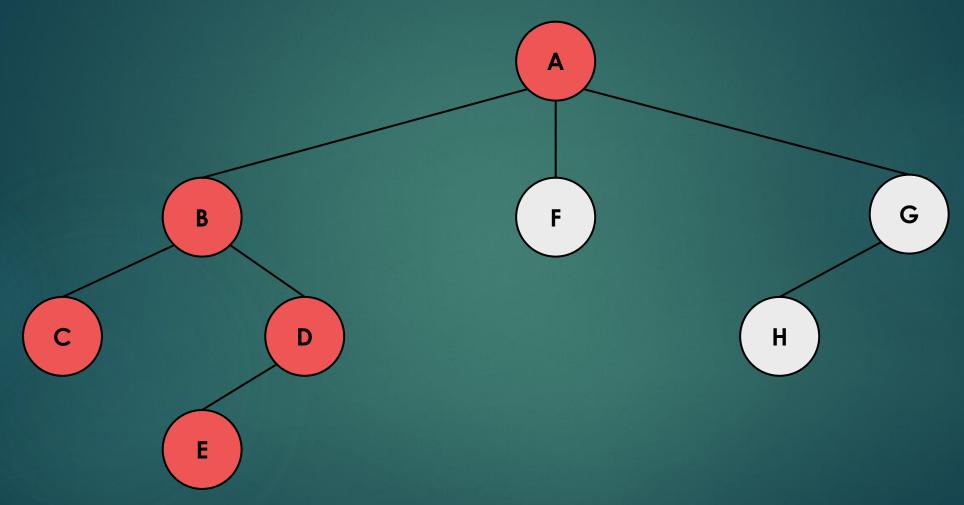


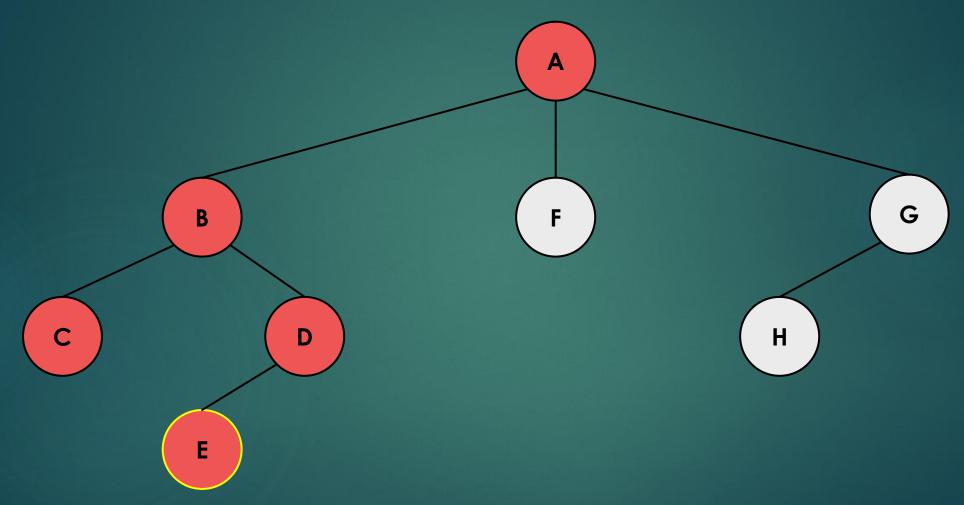


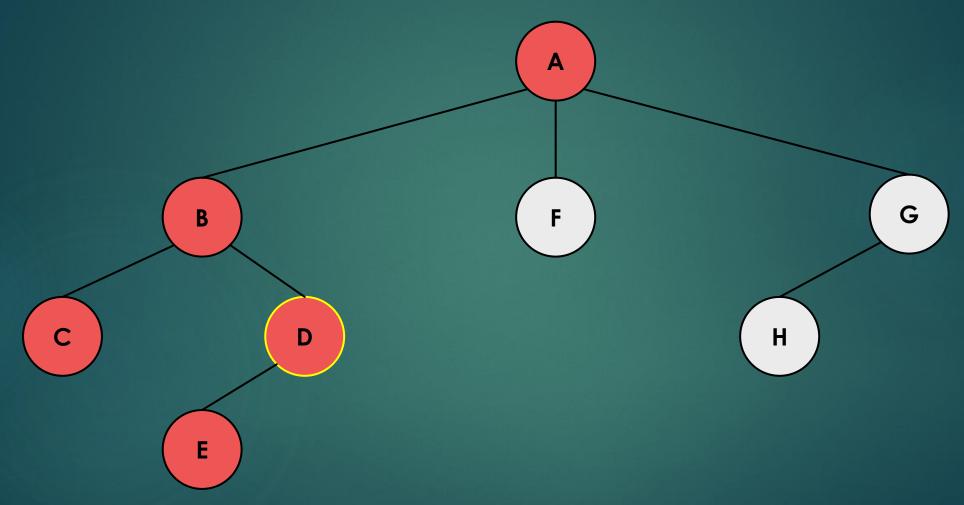


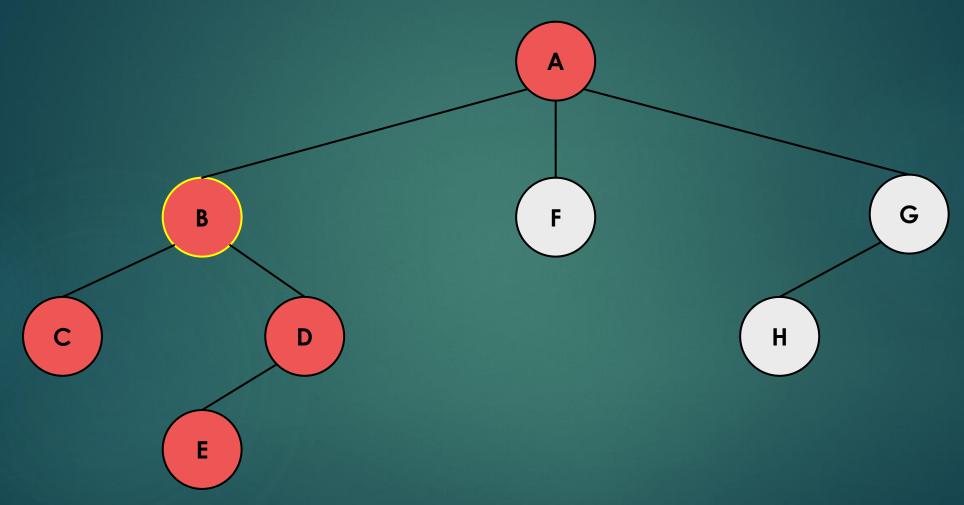


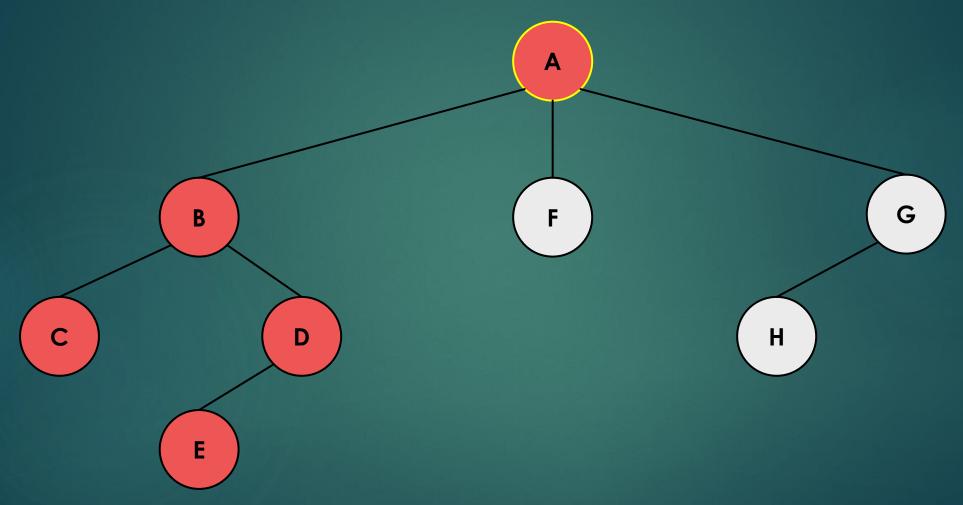


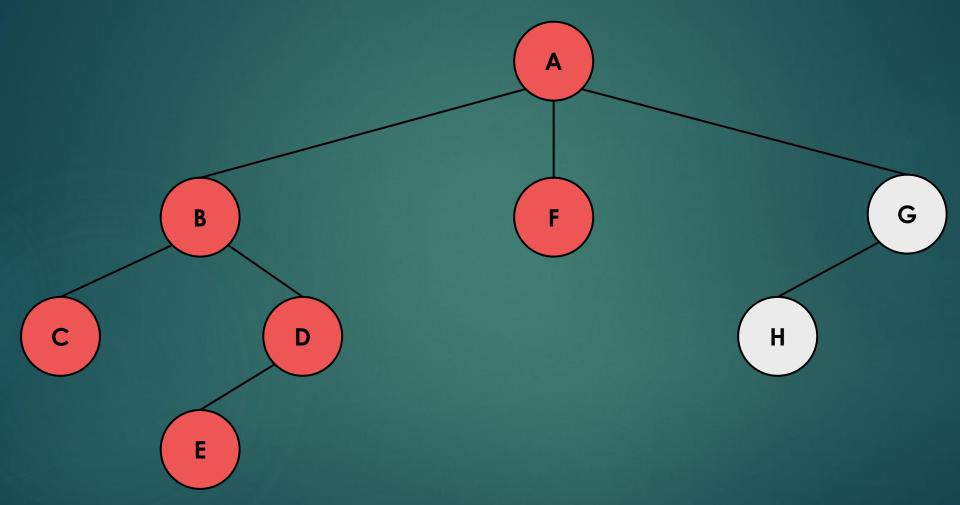


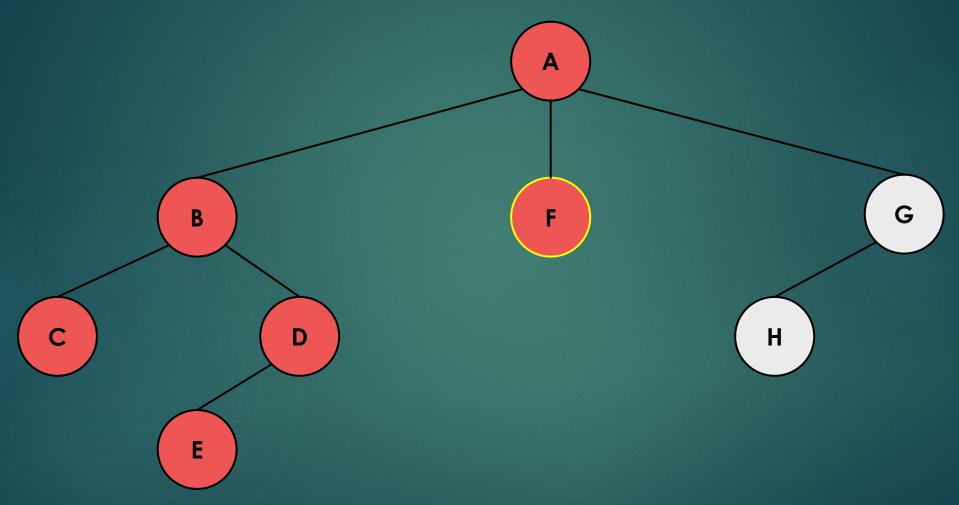


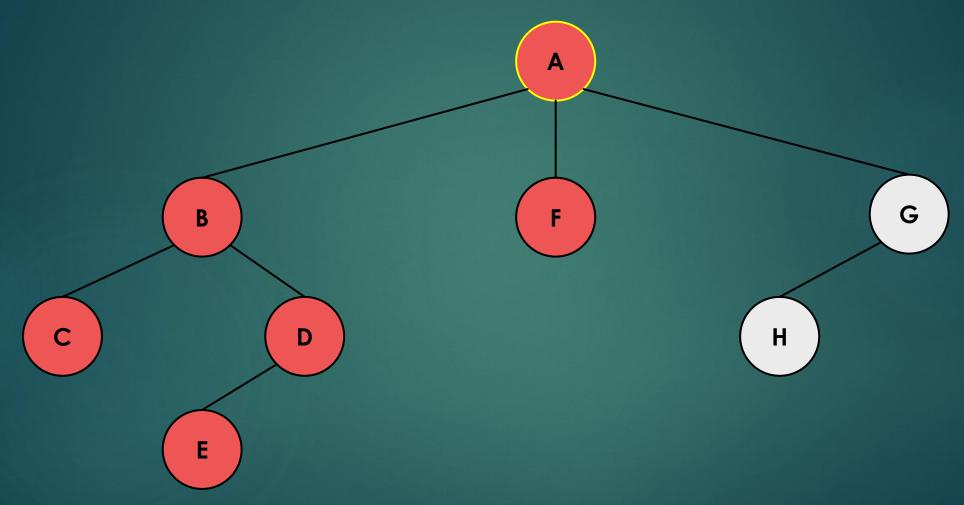


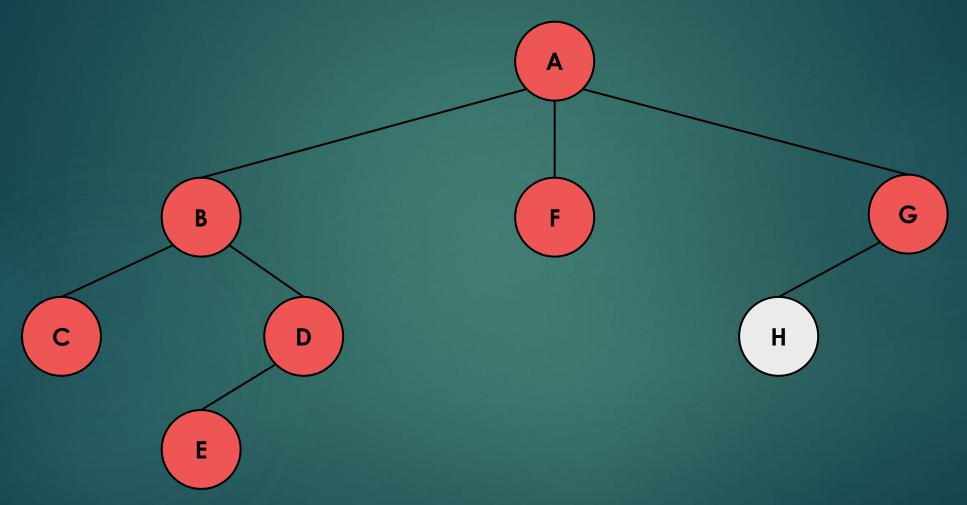


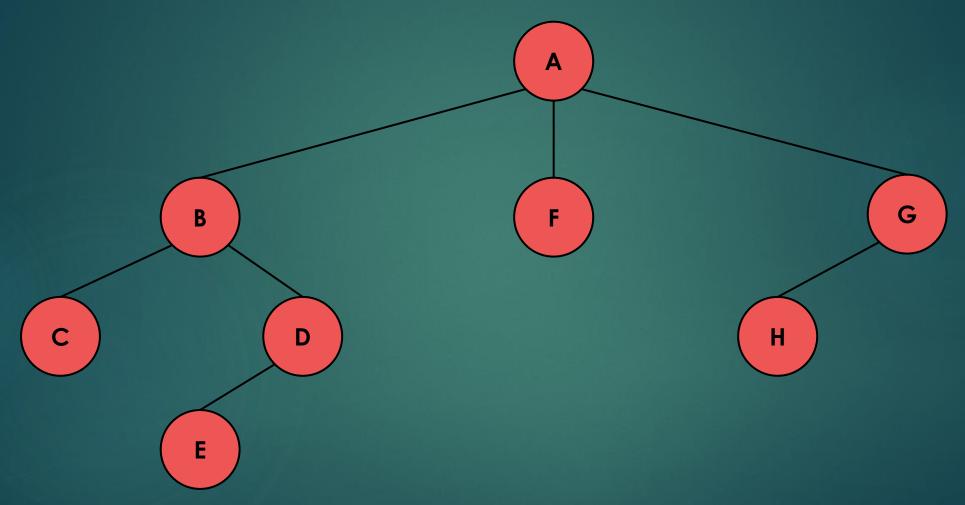


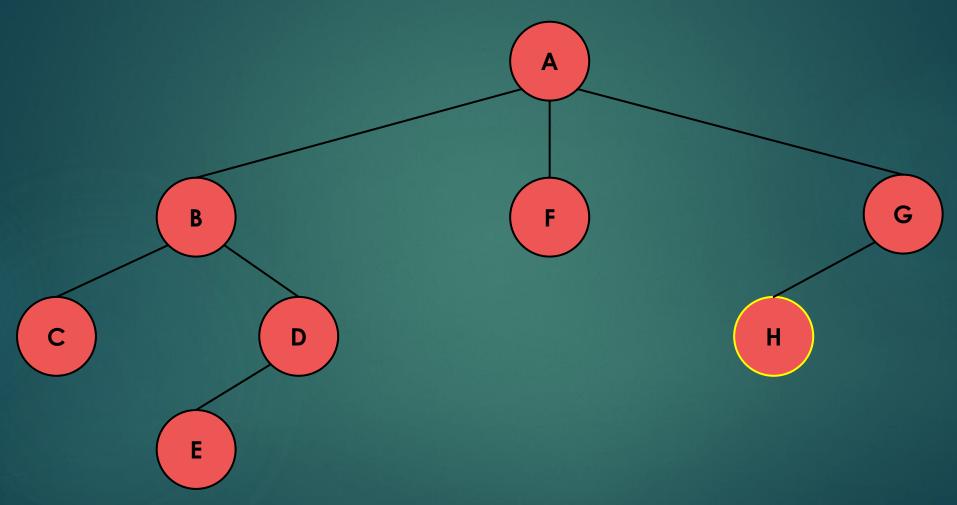


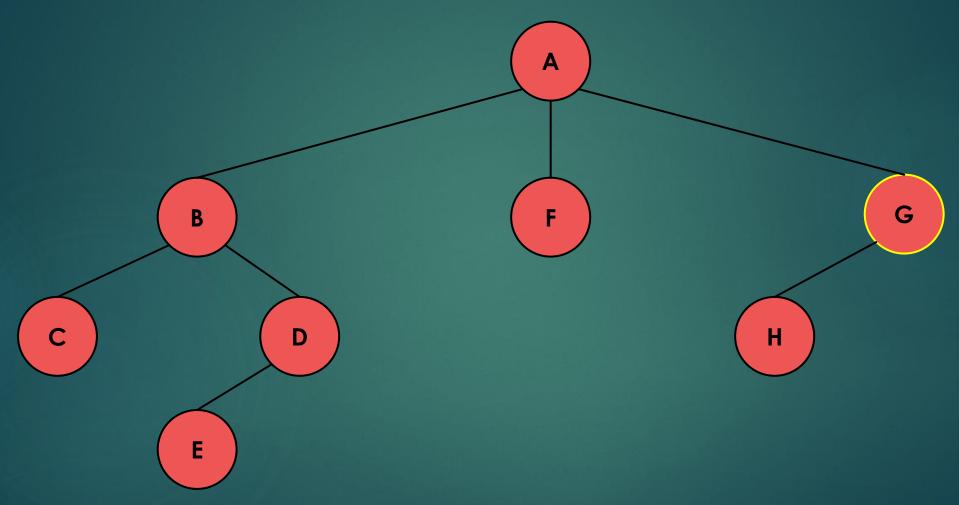


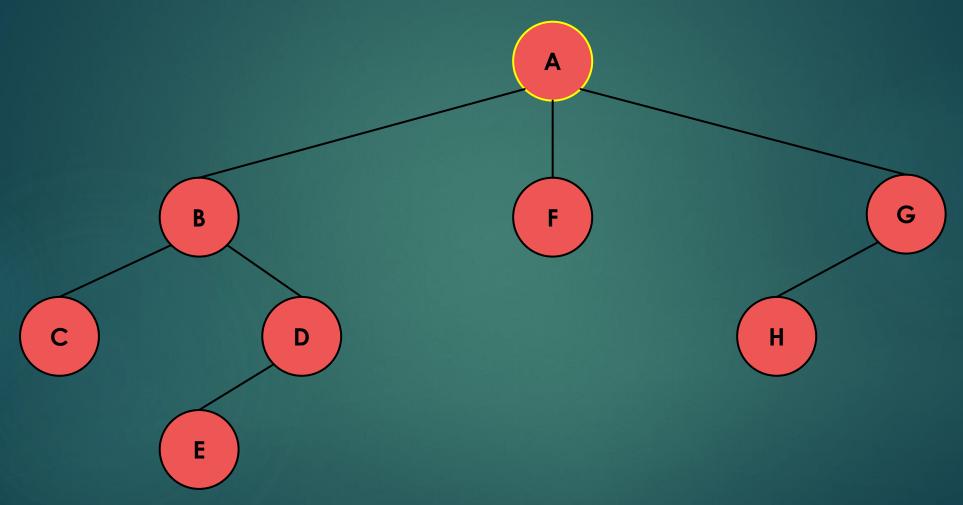


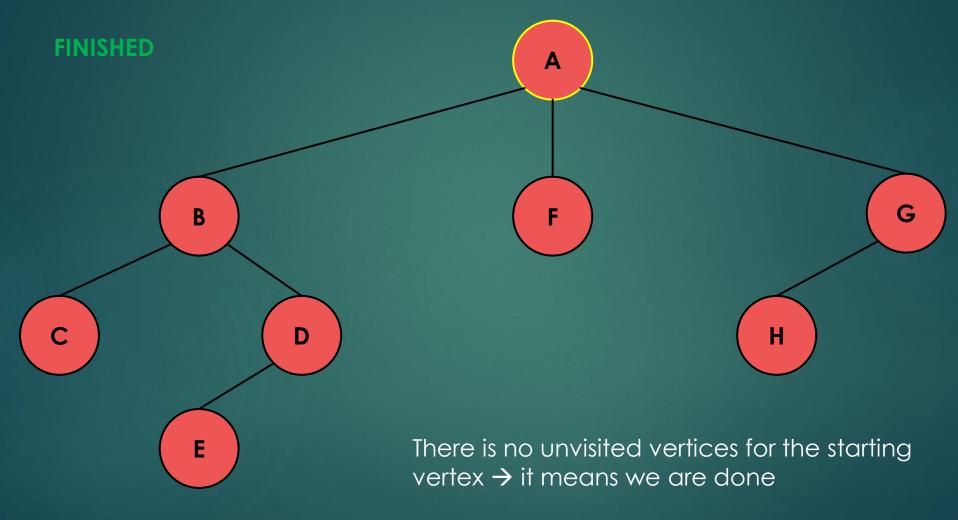












#### <u>Applications</u>

- Topological ordering
- Kosaraju algorithm for finding strongly connected components in a graph which can be proved to be very important in recommendation systems (youtube)
- Detecting cycles (checking whether a graph is a DAG or not)
- Generating mazes OR finding way out of a maze

