

Theorem 0.1. (*Ray-Chaudhuri-Wilson*) Give a subset L of s nonnegative integers, Let F be a family of subsets of $[n]$ such that the intersection of any two members of F has a cardinality contained in L . Then

$$(0.2) \quad |F| \leq \sum_{i=0}^s \binom{n}{i}.$$

Proposition 0.3. Let $k, i \in \mathbb{N}$ and $1 \leq i \leq k/2$, there exists k graph property on v vertices that has sensitivity $O(v^{k-i})$ and block sensitivity $\Omega(v^k)$.

Proof. Define f be the graph property that there exists a $K_{k+1}^{(k)}$ inside the graph such that $|K_{k+1}^{(k)} \cap E| \leq i - 1$ for any edge E lies not entirely inside the clique. We claim that this is the desired k -graph property.

First we calculate the block sensitivity. Consider the empty graph, there are $\binom{v}{k+1}$ such cliques and for every clique chosen we eliminate any other cliques with more than $k - 1$ points in common. In this way we get at least

$$(0.4) \quad \frac{\binom{v}{k+1}}{\binom{k+1}{k}(v-k)} = c * v^k$$

many cliques and any two of them have no more than $k - 1$ points in common, which guarantees that they are disjoint sensitive blocks. This shows that block sensitivity of f is $\Omega(v^k)$.

Then we want to show that the sensitivity of f is $O(v^{k-i})$. We calculate the sensitivity by looking at $s^0(f)$ and $s^1(f)$ separately

When $f = 1$, there is a desired $K_{k+1}^{(k)}$ clique inside the graph. To change the value of f , we need to either remove an edge from $K_{k+1}^{(k)}$ or add an edge with more than $i - 1$ common vertices with the clique. We have

$$(0.5) \quad s^1(f) \leq \binom{k+1}{2} + \binom{k+1}{i} \binom{v-i}{k-i} \leq C * v^{k-i},$$

for some constant C .

When $f = 0$, there doesn't exist such $K_{k+1}^{(k)}$ in the graph. We call a $k + 1$ -tuple $\{v_1, \dots, v_{k+1}\}$ sensitive if adding or removing an edge from the graph will make $\{v_1, \dots, v_{k+1}\}$ the vertices of a desired clique. Any sensitive edge is associated with a sensitive tuple by this definition. Also, we can show that there is precisely 1 sensitive edge associated with each sensitive tuple since if there are two sensitive edges associated with a $k + 1$ -tuple, we can't construct a desired clique by just removing or adding just one edge. Thus we have a injection from the set of sensitive edges into the set of sensitive tuples, and let \mathcal{F} denote the set of all sensitive tuples, we have

$$(0.6) \quad s^0(f) \leq |\mathcal{F}|.$$

Given two sensitive tuples, if they have more than $i - 1$ vertices in common, without loss of generality, assume that $\{v_1, \dots, v_i\}$ are vertices in common. If both of them can form a desired clique by adding an edge to the graph, there are no less than 1 edge through these i points and adding any edge will not remove these edges, which contradicts the fact that these two tuples can form a desired clique by adding an edge. If there exists one tuple can form a desired clique by removing an edge, there exists no less than 2 edges through $\{v_1, \dots, v_i\}$ in one of the tuples, and thus another tuple can't form a desired clique by either removing or adding exactly

1 edge, which contradicts the fact that the tuple is sensitive. Then given any two sensitive tuples, they have at most $k - i - 1$ common vertices.

Then \mathcal{F} is a family of subsets of $[v]$ and any of the two subsets in F have at most $i - 1$ intersections. By using **Theorem 0.1** with $L = \{0, 1, \dots, i - 1\}$ and $n = v$, we have

$$(0.7) \quad s^0(f) \leq |F| \leq \sum_{j=1}^i \binom{v}{j} \leq C' * v^i,$$

for some constant C' .

From above $s^0(f) = O(v^{k-i})$ and $s^1(f) = O(v^i)$ we conclude that $s(f) = O(v^{k-i})$ since $i \leq k/2$. Hence, f is a k -graph property with sensitivity v^{k-i} and block sensitivity v^k . \square

Corollary 0.8. *When k is even, there exists k graph property on v vertices that has sensitivity $O(v^{k/2})$ and block sensitivity $\Omega(v^k)$. When k is odd, there exists k graph property on v vertices that has sensitivity $O(v^{(k+1)/2})$ and block sensitivity $\Omega(v^k)$.*

Proof. When k is even, choose $i = k/2$ and when k is odd, choose $i = (k - 1)/2$ in above proposition. \square