

## Intuition

if we are able to find  $k$  pairs with maximum difference  $d$  and after that among all the  $d$  minimum will be answer.

## Approach

1. Sort the given nums array in ascending order.
2. Initialize  $lo$  to 0 and  $hi$  to the maximum difference between the largest and smallest elements in the sorted array.
3. Perform a binary search loop while  $lo$  is less than  $hi$ .
4. Calculate the midpoint  $mid$  between  $lo$  and  $hi$ .
5. Count the number of pairs with a difference less than or equal to  $mid$  while ensuring that no index appears more than once among the pairs.
6. If the count is greater than or equal to  $p$ , update  $hi$  to  $mid$ .
7. Otherwise, update  $lo$  to  $mid + 1$ .
8. Return the value of  $lo$  as the minimum maximum difference that allows forming at least  $p$  valid pairs.

## Complexity

- Time complexity:  $O(n \cdot \log(d))$  where  $d$  is maximum difference.
- Space complexity:  $O(1)$