

Problems Solvable using this template

- [3. Longest Substring Without Repeating Characters](#)
- [159. Longest Substring with At Most Two Distinct Characters \(Medium\)](#)
- [340. Longest Substring with At Most K Distinct Characters](#)
- [424. Longest Repeating Character Replacement](#)
- [487. Max Consecutive Ones II](#)
- [713. Subarray Product Less Than K](#)
- [1004. Max Consecutive Ones III](#)
- [1208. Get Equal Substrings Within Budget \(Medium\)](#)
- [1493. Longest Subarray of 1's After Deleting One Element](#)
- [1695. Maximum Erasure Value](#)
- [1838. Frequency of the Most Frequent Element](#)
- [2009. Minimum Number of Operations to Make Array Continuous](#)
- [2024. Maximize the Confusion of an Exam](#)

The following problems are also solvable using the shrinkable template with the ["At Most to Equal" trick](#)

- [930. Binary Subarrays With Sum \(Medium\)](#)
- [992. Subarrays with K Different Integers](#)
- [1248. Count Number of Nice Subarrays \(Medium\)](#)
- [2062. Count Vowel Substrings of a String \(Easy\)](#)

Template 1: Sliding Window (Shrinkable)

The best template I've found so far:

```
int i = 0, j = 0, ans = 0;
for (; j < N; ++j)
{
    // CODE: use A[j] to update state which might make the window invalid
    for (; invalid(); ++i) { // when invalid, keep shrinking the left edge until it's valid again
        // CODE: update state using A[i]
    }
    ans = max(ans, j - i + 1); // the window [i, j] is the maximum window we've found thus far
}
return ans;
```

Template 2: Sliding Window (Non-shrinkable)

```
int i = 0, j = 0;
for (; j < N; ++j) {
    // CODE: use A[j] to update state which might make the window invalid
    if (invalid()) { // Increment the left edge ONLY when the window is invalid. In this way, the window GROWS when it's valid, and SHIFTS
                     // when it's invalid
        // CODE: update state using A[i]
        ++i;
    }
    // after '++j' in the for loop, this window '[i, j]' of length 'j - i' MIGHT be valid.
}
return j - i; // There must be a maximum window of size 'j - i'.
```

Essentially, we GROW the window when it's valid, and SHIFT the window when it's invalid.

Note that there is only a SINGLE for loop now!

Apply these templates to other problems

[1493. Longest Subarray of 1's After Deleting One Element \(Medium\)](#)

Sliding Window (Shrinkable)

1. What's state? cnt as the number of 0s in the window.
2. What's invalid? cnt > 1 is invalid.

```
// OJ: https://leetcode.com/problems/longest-subarray-of-1s-after-deleting-one-element/
// Author: github.com/lz1124631x
// Time: O(N)
// Space: O(1)
class Solution {
public:
    int longestSubarray(vector<int>& A) {
        int i = 0, j = 0, N = A.size(), cnt = 0, ans = 0;
        for (; j < N; ++j) {
            cnt += A[j] == 0;
            while (cnt > 1) cnt -= A[i++] == 0;
            ans = max(ans, j - i); // note that the window is of size 'j - i + 1'. We use 'j - i' here because we need to delete a number.
        }
        return ans;
    }
};
```

Sliding Window (Non-shrinkable)

```
// OJ: https://leetcode.com/problems/longest-subarray-of-1s-after-deleting-one-element/
// Author: github.com/lzl124631x
// Time: O(N)
// Space: O(1)
class Solution {
public:
    int longestSubarray(vector<int>& A) {
        int i = 0, j = 0, N = A.size(), cnt = 0;
        for (; j < N; ++j) {
            cnt += A[j] == 0;
            if (cnt > 1) cnt -= A[i++] == 0;
        }
        return j - i - 1;
    }
};
```

3. Longest Substring Without Repeating Characters (Medium)

Sliding Window (Shrinkable)

1. state: cnt[ch] is the number of occurrence of character ch in window.
2. invalid: cnt[s[j]] > 1 is invalid.

```
// OJ: https://leetcode.com/problems/longest-substring-without-repeating-characters/
// Author: github.com/lzl124631x
// Time: O(N)
// Space: O(1)
class Solution {
public:
    int lengthOfLongestSubstring(string s) {
        int i = 0, j = 0, N = s.size(), ans = 0, cnt[128] = {};
        for (; j < N; ++j) {
            cnt[s[j]]++;
            while (cnt[s[j]] > 1) cnt[s[i++]]--;
            ans = max(ans, j - i + 1);
        }
        return ans;
    }
};
```

Sliding Window (Non-shrinkable)

Note that since the non-shrinkable window might include multiple duplicates, we need to add a variable to our state.

1. state: dup is the number of different kinds of characters that has duplicate in the window. For example, if window contains aabbc, then dup = 2 because a and b has duplicates.
2. invalid: dup > 0 is invalid

```
// OJ: https://leetcode.com/problems/longest-substring-without-repeating-characters/  
// Author: github.com/lzl124631x  
// Time: O(N)  
// Space: O(1)  
class Solution {  
public:  
    int lengthOfLongestSubstring(string s) {  
        int i = 0, j = 0, N = s.size(), cnt[128] = {}, dup = 0;  
        for (; j < N; ++j) {  
            dup += ++cnt[s[j]] == 2;  
            if (dup) dup -= --cnt[s[i++]] == 1;  
        }  
        return j - i;  
    }  
};
```

713. Subarray Product Less Than K (Medium)

Sliding Window (Shrinkable)

- state: prod is the product of the numbers in window
- invalid: prod >= k is invalid.

Note that since we want to make sure the window [i, j] is valid at the end of the for loop, we need $i \leq j$ check for the inner for loop. $i == j + 1$ means this window is empty.

Each maximum window [i, j] can generate $j - i + 1$ valid subarrays, so we need to add $j - i + 1$ to the answer.

```

// OJ: https://leetcode.com/problems/subarray-product-less-than-k/
// Author: github.com/lzl124631x
// Time: O(N)
// Space: O(1)
class Solution {
public:
    int numSubarrayProductLessThanK(vector<int>& A, int k) {
        if (k == 0) return 0;
        long i = 0, j = 0, N = A.size(), prod = 1, ans = 0;
        for (; j < N; ++j) {
            prod *= A[j];
            while (i <= j && prod >= k) prod /= A[i++];
            ans += j - i + 1;
        }
        return ans;
    }
};

```

The non-shrinkable template is not applicable here since we need to the length of each maximum window ending at each position