143 - Networks (pt2) mnouncements -HW2 Jue - Lab 3 up - due in I week

he I ab: IPTables

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0917



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Information - Experimentation -My DETERIab

Experiment (UP-CS448/T Begin an Experiment Begin a Risky Experiment

Experiment Options

View Activity Logfile

Swap Experiment Out

Terminate Experiment

Modify Experiment

Make Experiment Risky

Modify Traffic Shaping Modify Settings

Link Tracing/Monitoring

Event Viewer

Update All Nodes

Reboot All Nodes

Experiment List

Node Status List ImageIDs

List OSIDs

Start New Project Join Existing Project

Group:

Experiment H

ttings

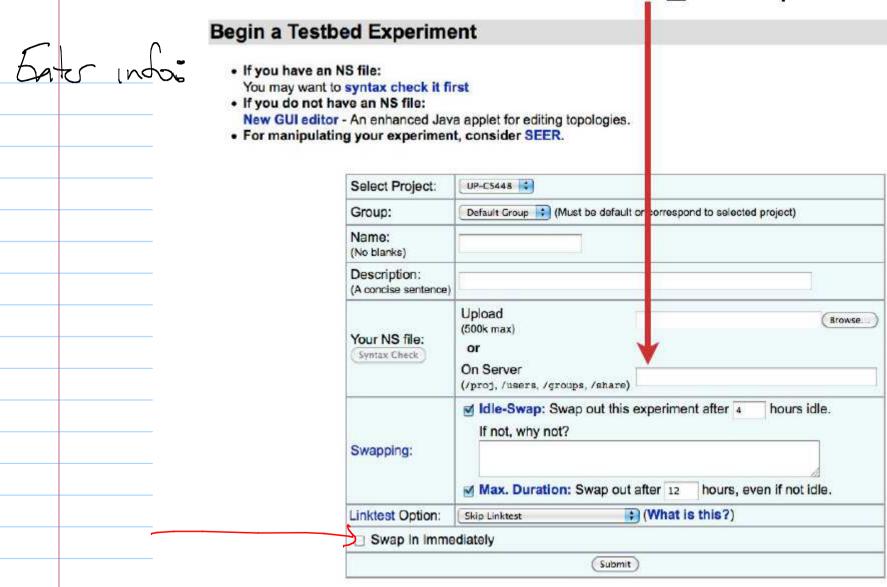
ion:

Created:

Last Swap/Mc

Idle-Swap:

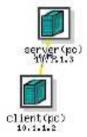
/share/education/PermissionsFirewalls_UCLA/permissions.ns



The Visualization tab will show the network architechture of this experiment.

(Will be useful in later Lbs, too.)

Settings Visualization NS File Details



Visualization NS File Install firewall on the server node. Use client node to client(pc) test the firewall.

Edit a run the frewall script
V
up448ar@server:~\$ cd /root/firewall
up448ar@server:/root/firewall\$ ls
extingui.sh extingui.sh~ firewall.sh,firewall.sh~ iptables_reset
tester.sh
make youre edits)
up448ar@server:/root/firewall\$ sudo sh firewall.sh
Starting firewall: done.
Starting Tirewall. done.

Note:

Use extinguish to reset between attempts.

up448ar@server:~\$ cd /root/firewall

up448ar@server:/root/firewall\$ ls

extingui.sh extingui.sh~ firewall.sh firewall.sh~ iptables_reset-tester.sh

up448ar@server:/root/firewall\$ sudo sh extingui.sh

up448ar@server:/root/firewall\$

TPtables; Quick intro

-L: list rules

% sudo iptables -L

Chain INPUT (policy ACCEPT) target prot opt source

destination

Chain FORWARD (policy ACCEPT) target prot opt source

destination

Chain OUTPUT (policy ACCEPT) target prot opt source

destination

-A: Append a rule to the INPUT, OUTPUT or FORWARD chain. -p: Specify a protocol. -d: Specify a destination port. -j: Specify a target: ACCEPT: Accept the packet. REJECT: Reject the packet, notify the sender. DROP: Silently ignore the packet. LOG: Log the packet.

Example:

Add rule to accept in coming 55h

traffic:

iptables -A INPUT -p tcp --dport 22 -j ACCEPT

Also: restrict based on where it is from
-i: input interface.
-o: output interface.
-s: source, e.g., a machine name or an IP address.

eth4 eth8

up448ar@server:~\$ ifconfig

Link encap: Ethernet Hwaddr 00:1b:21:1e:ac:14 inet addr: 10.1.1.3 Bcast: 10.1.1.255 Mask: 255.255.255.0 inet6 addr: fe80::21b:21ff:fe1e:ac14/64 Scope:Link UP BROADCAST RUNNING MULTICAST MTU: 1500 Metric: 1 RX packets:0 errors:0 dropped:0 overruns:0 frame:0

TX packets:9 errors:0 dropped:0 overruns:0 carrier:0

collisions:0 txqueuelen:100

RX bytes: 0 (0.0 B) TX bytes: 706 (706.0 B)

Link encap: Ethernet HWaddr 00:13:72:4e:cc:6d inet addr: 192.168.1.181 Bcast: 192.168.3.255 Mask: 255.255.252.0 inet6 addr: fe80::213:72ff:fe4e:cc6d/64 Scope:Link UP BROADCAST RUNNING MULTICAST MTU: 1500 Metric: 1 RX packets:14551 errors:0 dropped:0 overruns:0 frame:0 TX packets:2131 errors:0 dropped:0 overruns:0 carrier:0 collisions:0 txqueuelen:1000

RX bytes:18471062 (18.4 MB) TX bytes:199954 (199.9 KB)

Then use this:

iptables -A INPUT -i eth4 -s server -j DROP

What does this rule do?

Drop all incomine traffic on eth4

from server

Networking?

Recall:

The Internet Protocol Suite, informally called TCP/IP, is an implementation of the OSI model.

Application					
Presentation					
Session					
Transport					
Network					
Data Link					
Physical					

OSI Model

Application

Transport

Internet

Link

TCP/IP

TCP/IP Layers

Application

Internet

Transport

Link

Determines how an application should use the network.

Handles details of data transmission at the end-points of communication.

Expresses how packets should be forwarded and how packets should be addressed (also called **Network Layer**)

TCP/IP

TCP/IP Layers

Application

Transport

Internet

Link

TCP/IP

Sometimes divided into Link Layer and Physical Layer.

Link Layer: Provides for synchronization and transfer of information. Defines how physical machines address each other.

Physical Layer: Defines electrical aspects of sending signals along a wire or wirelessly. Also addresses switch and router hardware.

We'll to cus on the **Application** addressed and Transport Arwarded over the network. Internet Link TCP/IP Headers in IPv4:
-Divided into 32-bit segments
- Headers are 5 segments long (usually), with data at the end

bit 0-3 4-7		8-13 14-15	14-15	16-18	19–31	
0	Version Header Length Differentiated Services Code Point Congest		Explicit Congestion Notification	1	Total Length	
32	Identification				Flags	Fragment Offset
64	Time to Live Protocol					Header Checksum
96	Source IP Address					
128	Destination IP Address					
160	Options (if Header Length > 5)					
160 or 192+	Data					

TPVY

10.1.3

Class A	0 net id	(7 bit) hos	st id (24 bit)		126 networks, 16 million host
Class B	10 net	id (14 bit)	host id (16	bit)	16382 networks, 65,534 hosts
Class C	110	net id (21 bit)		host id (8 bit)	2 million networks, 254 hosts
Class D	1110	multic	east (28 bit)		designed for multicasting
Class E	11110	future	use (27 bit)	1	reserved for experiments

Example: Consider the address:

10001000 11100101 11001001 0001000

Class? B

What IP? 136,229.201.8

Problem:

Thy was designed in 1981.

Classes A-C allow for under

4.3 billion address total.

(Reality - much smaller!)

Conclusion: Problem - out of space.

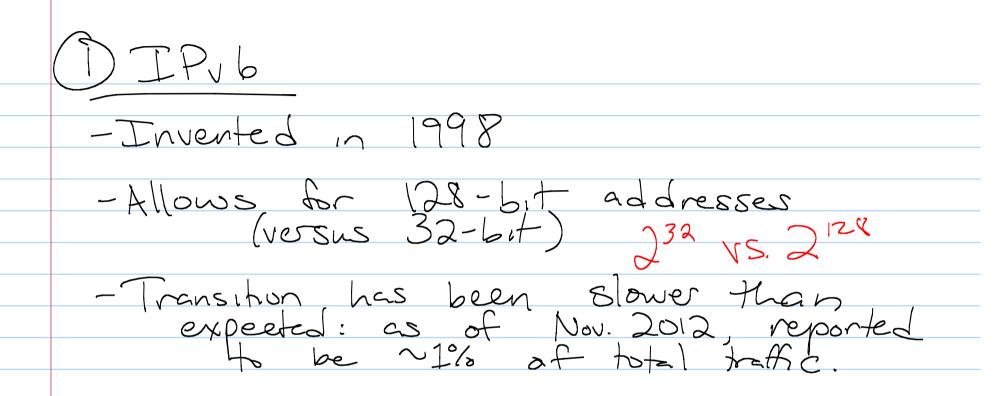
Solutions

(D) IPV6

(2) NAT

3 Subnething

None is a perfect are but all have been used to offset issues.



Thub details:

- Packet headers are twice as long.

- However, processing is actually simpler + fasted at routers.

- Privacy extensions exist to "hide" ideality: OS generates random host didentifier.

2) Network Address Translation A router stands between a private network a outside would Every internal IP address maps to a single TP/port which is all ontoide sees. (Combines well with Frewall Functionality.) Pros of NAT: - effectively a frenc! - only need I IP

Cons: - Single point of failure - significant overhead

(3) 5	Subnetting		
Ther	e is a ju	ump betwee	n class B and
Class B	10 net id (14 bit)	host id (16 bit)	16382 networks, 65,534 hosts
Class C	110 net id (21 bit)	host id (8 bit)	2 million networks, 254 hosts
Mo	any larger Subdivide	networks of	actually

Example: 1 2 15 16 24 25 31 Subnet ID Net ID Host ID 0 class B # hosts in each: 27

Subnets cont;

Every computer gets a subnet mask,

eg 2550255.255.128

[[[[[[]]]]]]

As well as an IP: 128.96.34.15

lake the bitwise AND to get the subnet versus host id: 1 2 15 16 24 25 31 Subnet ID Net ID Host ID

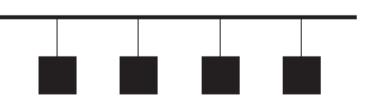
Local Area Networks

A LAN is a "small interconnection infrastructure that typically uses a shared transmission medium".

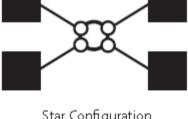
From Computer and Communication Networks by N. Mir

Note: A single LAN may actually be huge. "Local" is relative, but generally these all connect to same youter or switch.

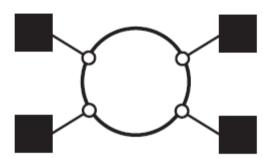
LAN Topologies



Bus Configuration

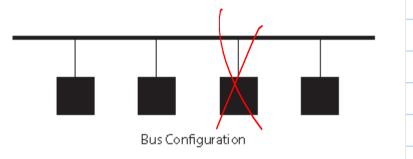


Star Configuration



Ring Configuration

Bus Configuration

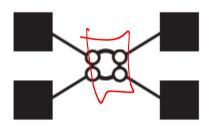


Transmissions are propagated on the bus in both directions.

All users (locally)
will receive all
packets.

Sender ques packet to hepeater. Repeater forwards until reaches destination. If reaches original Source, not Upproverded any further. Ring Configuration

Star configuration



Star Configuration

Center of "Ster" is a mult-port hub or switch.

Frames are sent to center, which either broadcasts or sends to targeted destination.

None of these have much security! are vulnerable to attack and to eaves drapping. To address this, we'll due down a level to the Link (More rest time!)

Application

Transport

Internet

Link

TCP/IP