

STL Containers

- Sequence Containers – store sequences of values
 - vector, deque, list
- Container Adapters – specialized interfaces to general containers
 - stack, queue, priority_queue
- **Associative Containers** – use “keys” to access data rather than position (Account #, ID, SSN, ...)
 - map
 - Multimap
 - set
 - multiset

Associative Containers

- Similar to vector & list – other storage structures with operations to **access** & **modify** elements.
- Stores elements based on a **key**.
- Key can consist of **one** or **more attributes** to **uniquely identify** each element (we will assume only one attribute).
- Example: Department of Motor Vehicles (DMV) uses license-plate # to identify a vehicle.
- Main difference is that an associative container uses the **key** rather than an **index** (vector) or linear search (list) to retrieve an element.

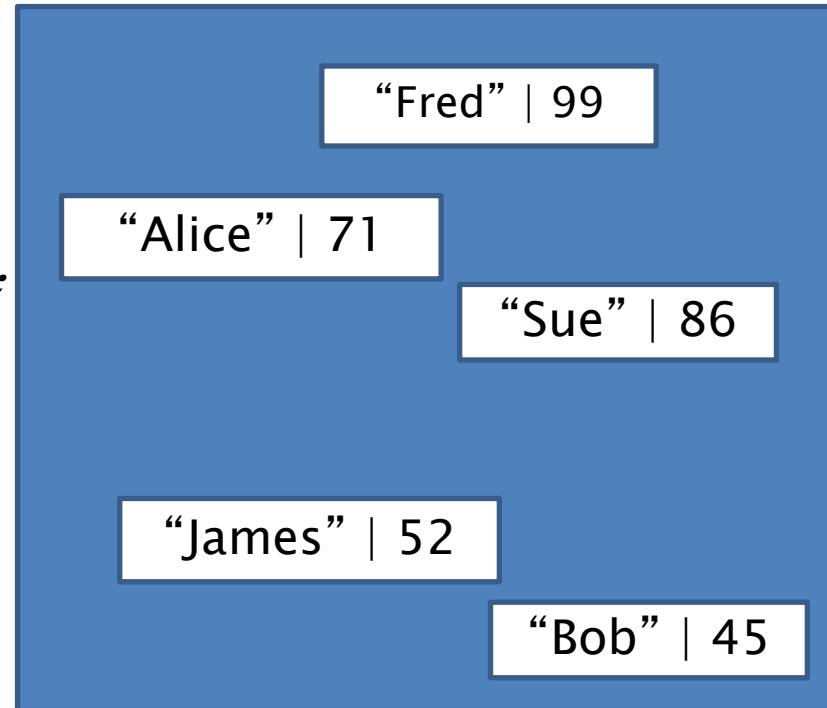
Associative Containers: unordered_map

- Stores a set of (key, value) pairs
- Each key has one value
- Implemented as a hash table

```
#include <unordered_map>  
  
//define it with  
//keys of type string  
//and values of int  
Unordered_map<string, int> um;
```

- Fast insert and delete

```
um[“Fred”] = 99;  
insert, erase
```



http://www.cplusplus.com/reference/unordered_map/unordered_map/

STL Unordered_map: Data Storage

- An STL map is implemented as a tree-structure, where each node holds a “**pair**”.
- Most important to know when retrieving data from the table
 - Some functions return the pair, not just the value
- A pair has two fields, *first* (holding the key) and *second* (holding the value)

STL Unordered_map: Data Storage

- If you have a *pair object*, you can use the following code to print the key and value:

```
cout << myPairObject.first << " " <<  
    myPairObject.second;
```

- If you have a *pointer to the pair object*, use the arrow operator instead

```
cout << myPairObject->first << " " <<  
    myPairObject->second;
```

STL Unordered_map: Data Storage

- Access element **at**
 - Returns a reference to the mapped value of the element identified with key k.
 - If k does not match the key of any element in the container, the function throws an `out_of_range` exception.
- Access element **[]**
 - If k matches the key of an element in the container, the function returns a reference to its mapped value.
 - If k does not match the key of any element in the container, the function inserts a new element with that key and returns a reference to its mapped value.

unordered_map.cpp

```
#include <iostream>
#include <unordered_map>
using namespace std;

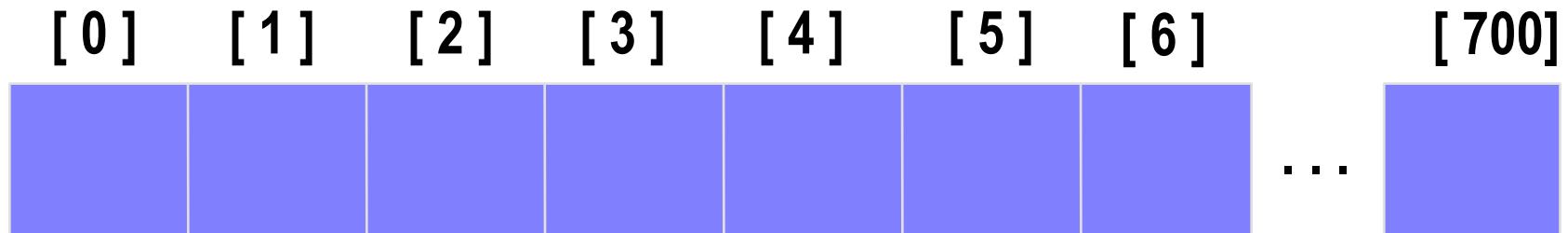
int main()
{
    // Declaring umap to be of <string, int> type
    // key will be of string type and mapped value will
    // be of double type
    unordered_map<string, int> umap;

    // inserting values by using [] operator
    umap[“Fred”] = 99;
    umap[“Sue”] = 86;
    umap[“Bob”] = 45;

    // Traversing an unordered map
    unordered_map<string, int>:: iterator itr;
    for (itr = umap.begin(); itr != umap.end(); itr++) {
        cout << itr->first << " " << itr->second << endl;
    }
}
```

What is a Hash Table ?

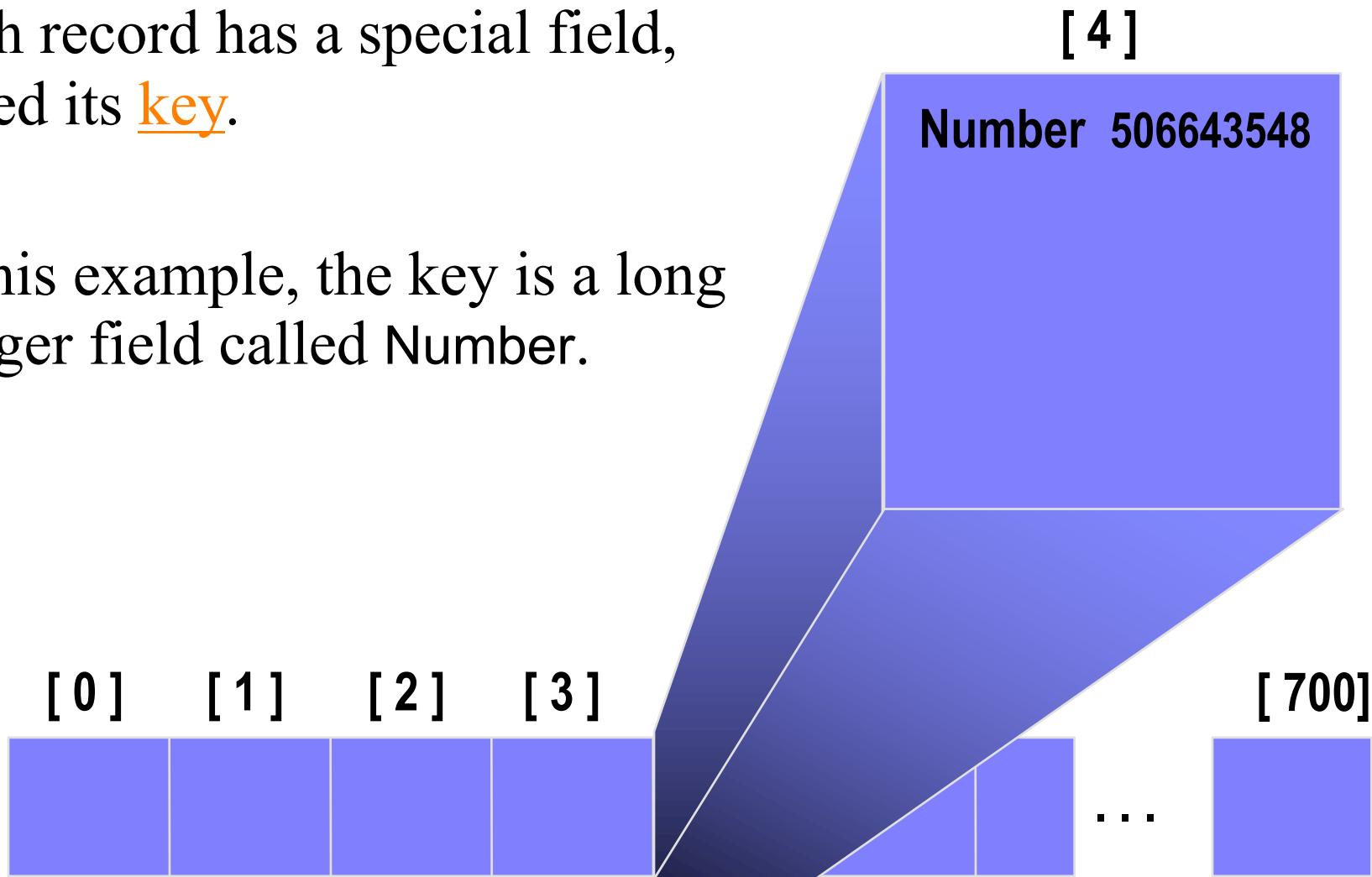
- Think about it as a Dictionary with <key,value> pairs.
- The simplest kind of hash table is an array of records.
- This example has 701 records.



An array of records

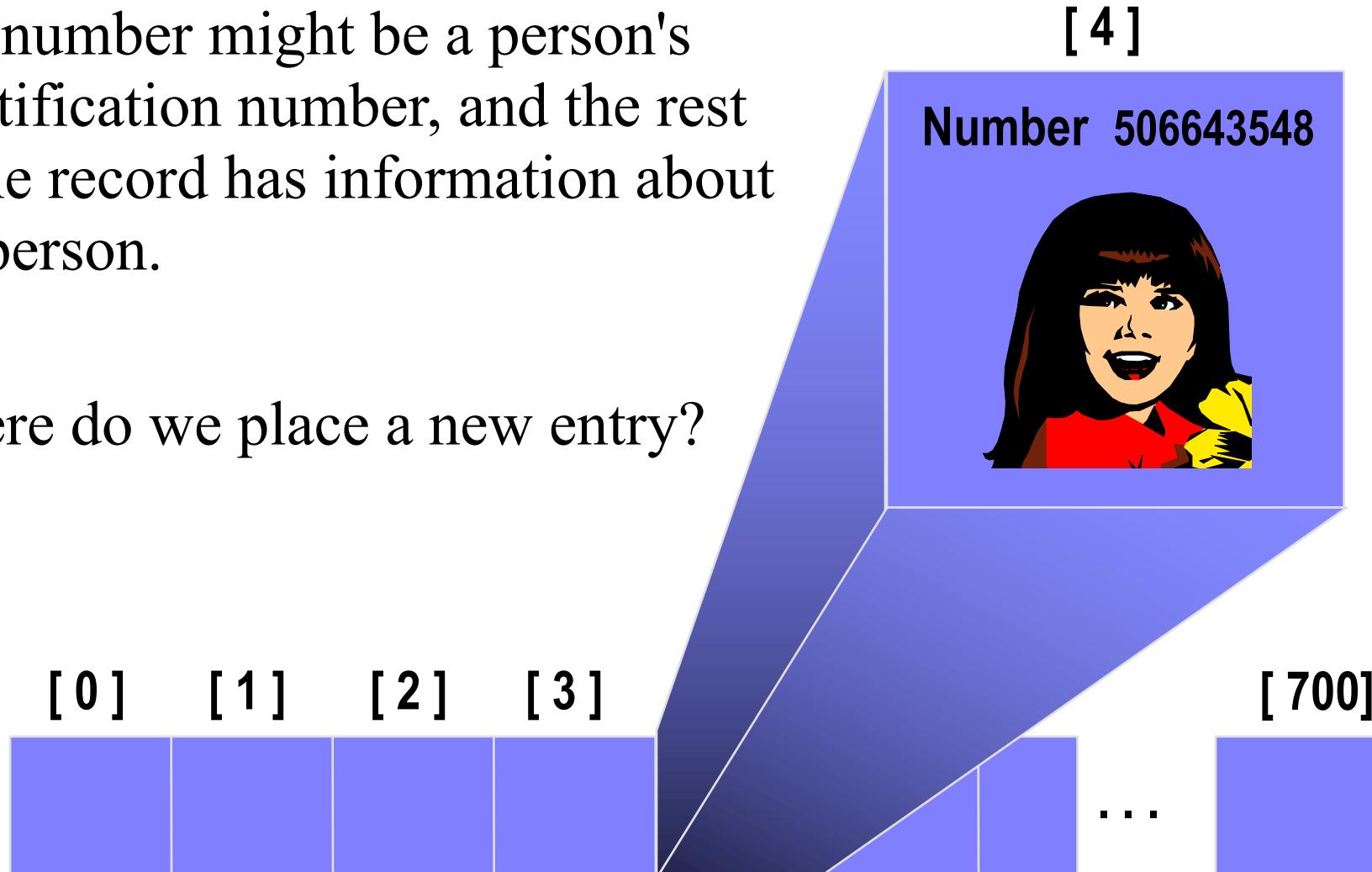
What is a Hash Table ?

- Each record has a special field, called its key.
- In this example, the key is a long integer field called Number.



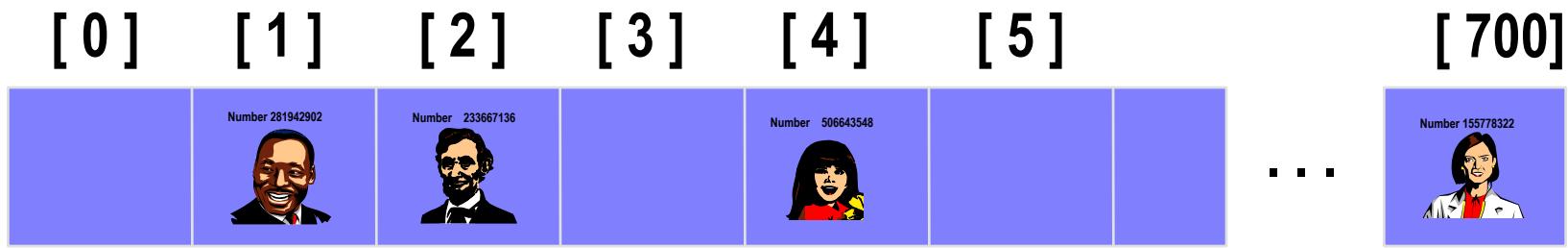
What is a Hash Table ?

- The number might be a person's identification number, and the rest of the record has information about the person.
- Where do we place a new entry?



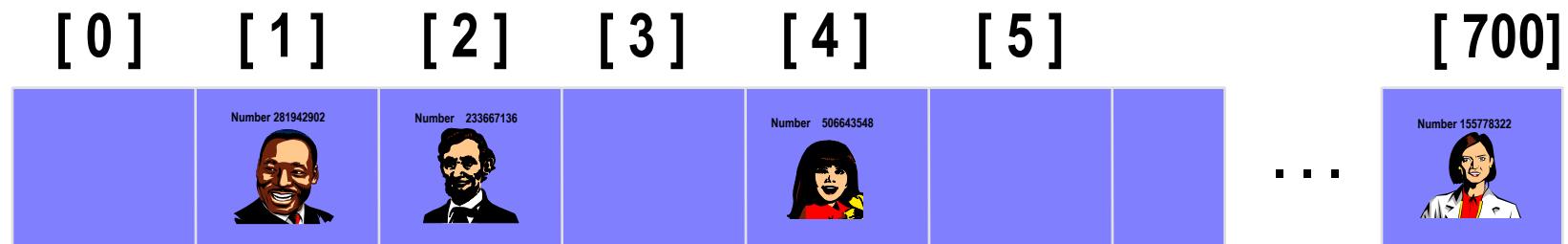
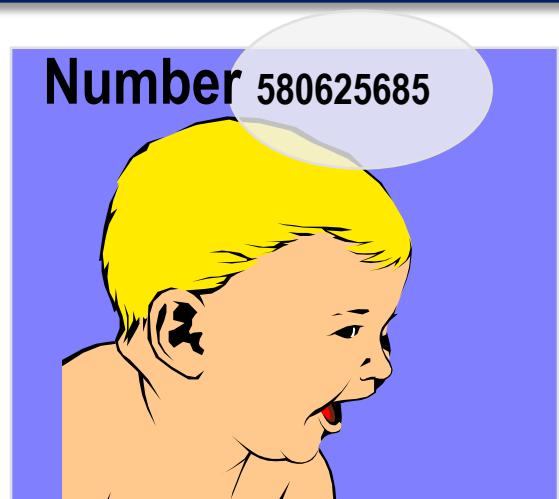
What is a Hash Table ?

- When a hash table is in use, some spots contain valid records, and other spots are "empty".



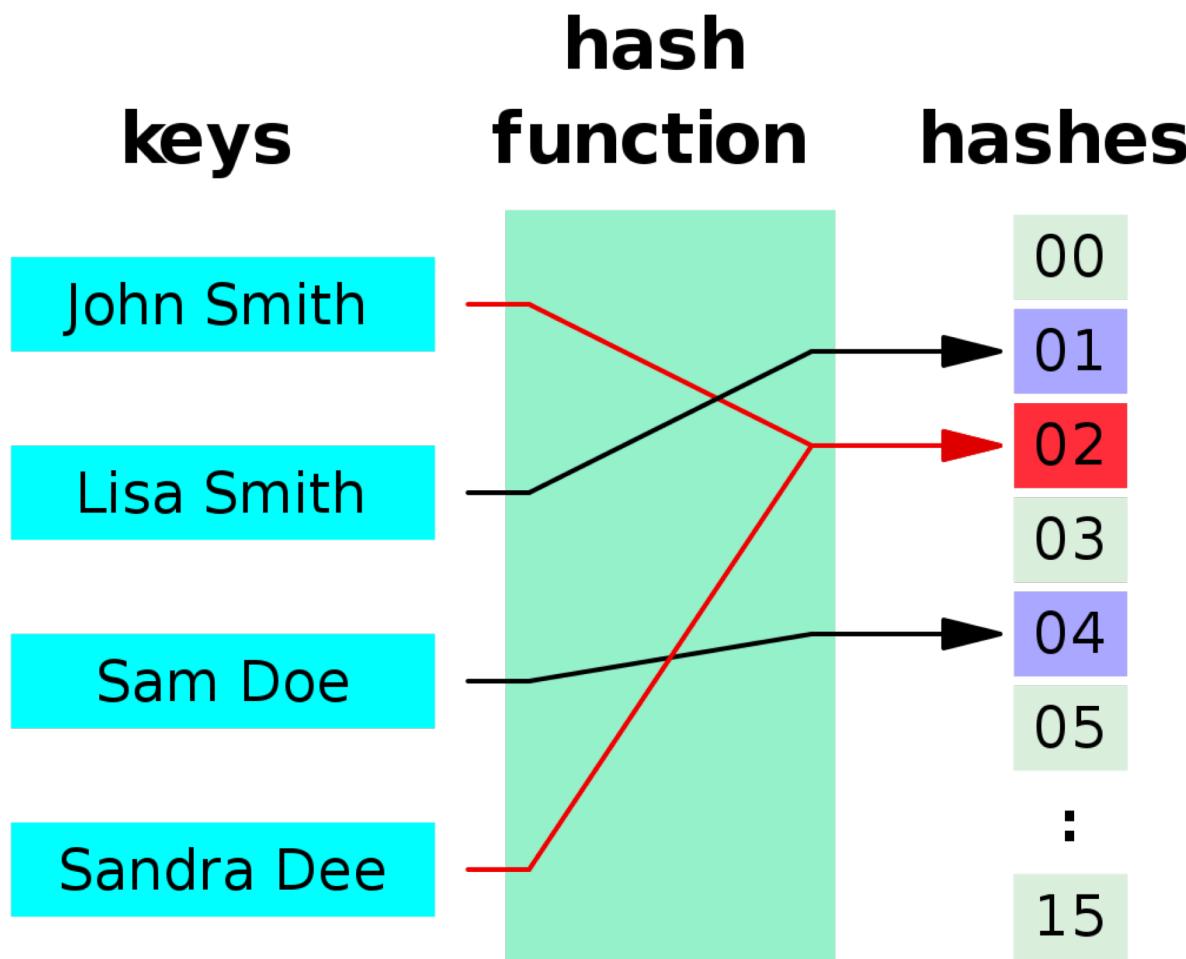
Inserting a New Record

- In order to insert a new record, the **key** must somehow be converted to an array **index**.
- The index is called the **hash value** of the key.



How can this conversion happen?

Hash Function

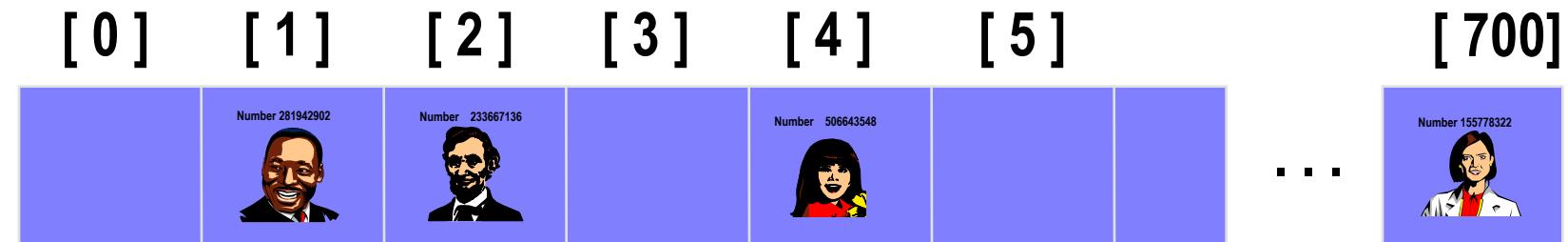
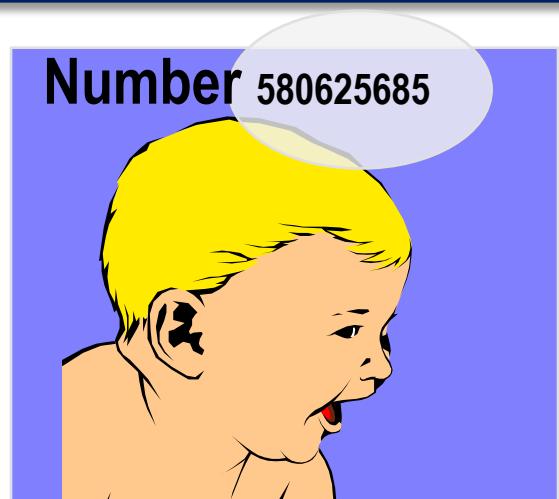


Inserting a New Record

- Typical way create a hash value:

(Number mod 701)

What is $(580625685 \text{ mod } 701)$?



Inserting a New Record

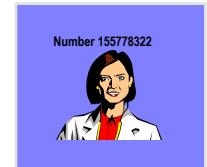
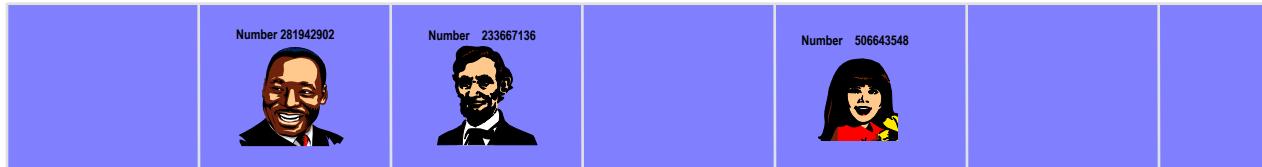
- Typical way create a hash value:

(Number mod 701)

What is $(580625685 \text{ mod } 701)$?

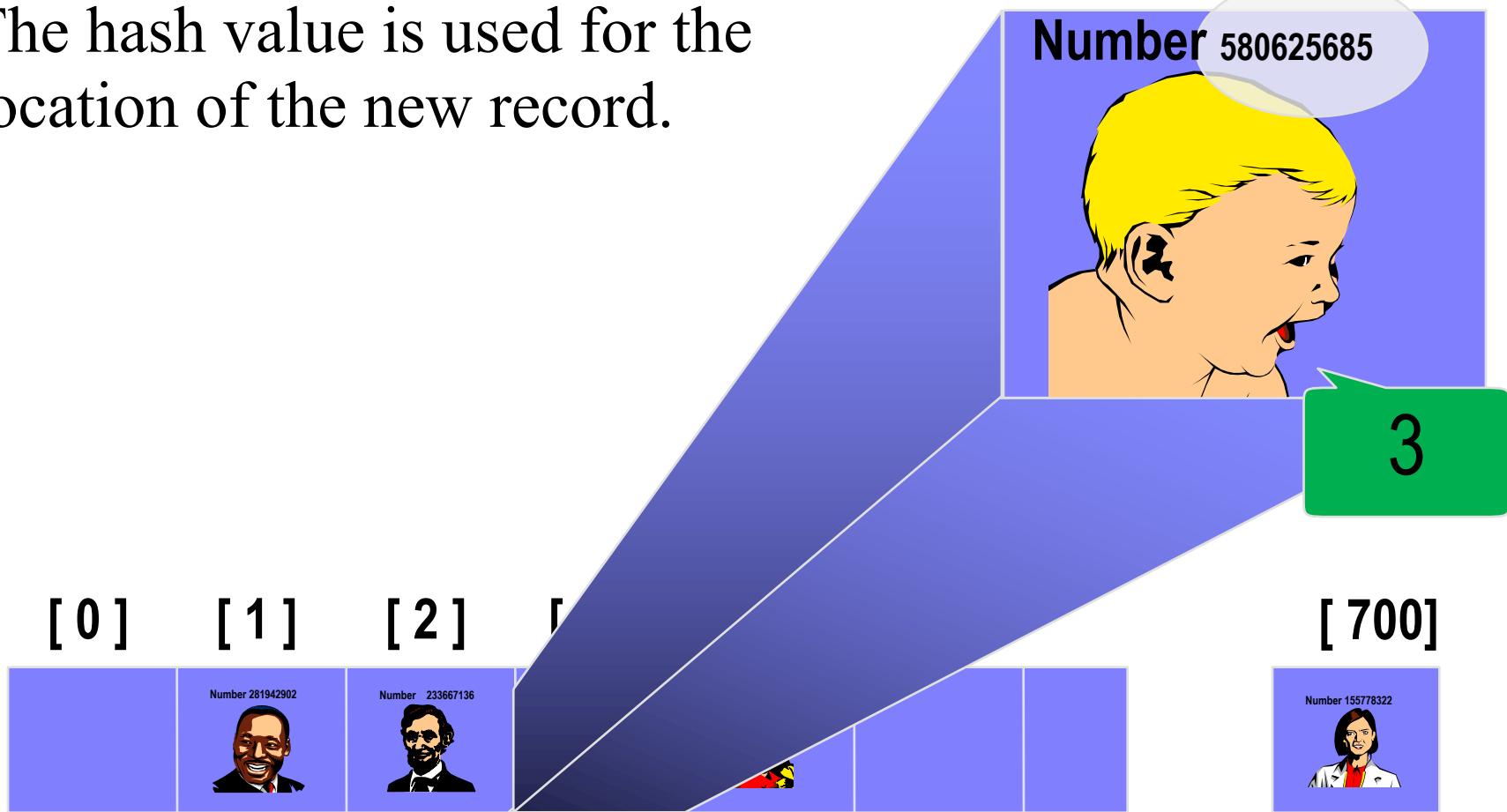


[0] [1] [2] [3] [4] [5] [700]



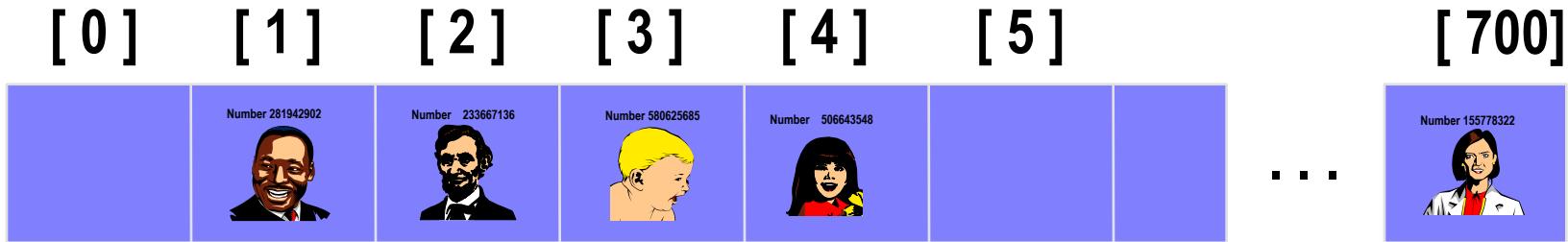
Inserting a New Record

- The hash value is used for the location of the new record.



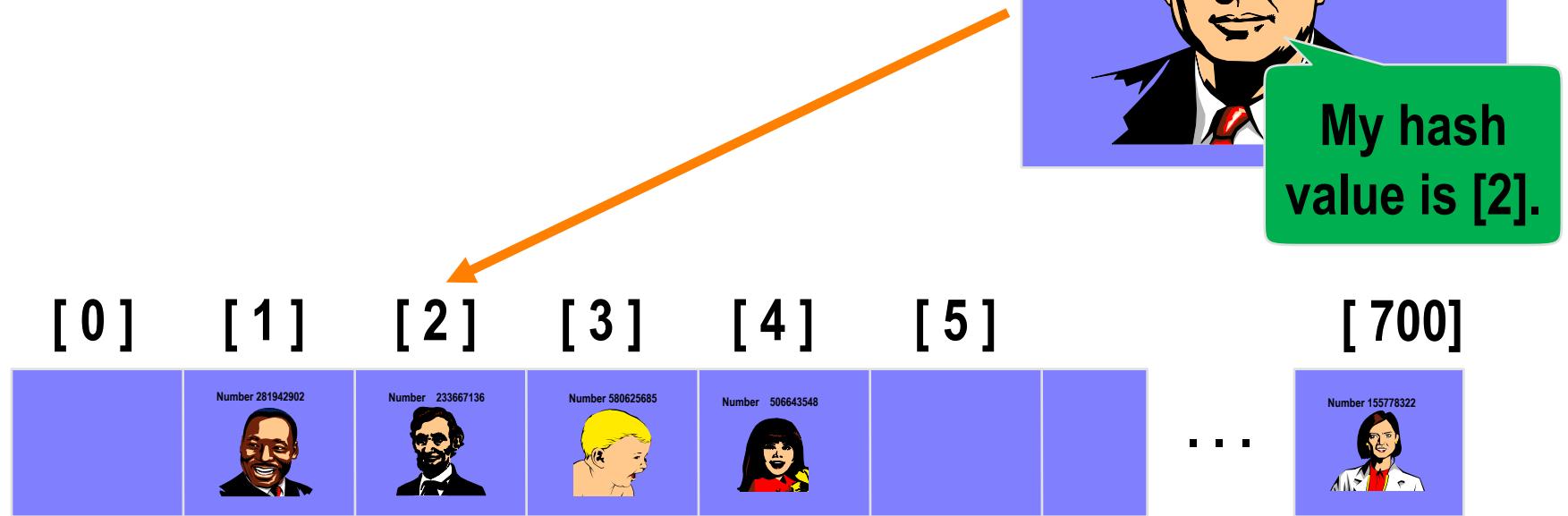
Inserting a New Record

- The hash value is used for the location of the new record.
- Do you see any problem?



Another Insertion

- Here is another new record to insert, with a hash value of 2.

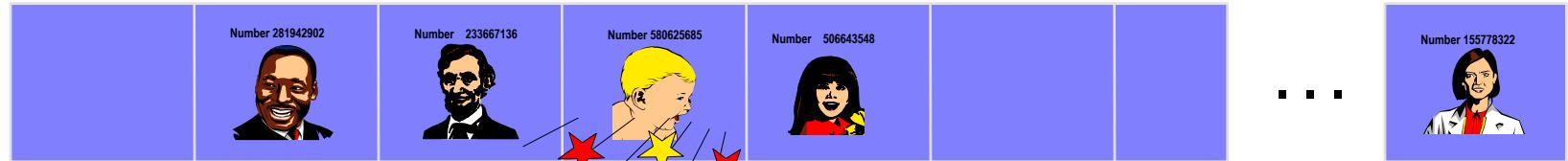


Collisions

- This is called a **collision**, because there is already another valid record at [2].

When a collision occurs,
move forward until you
find an empty spot.

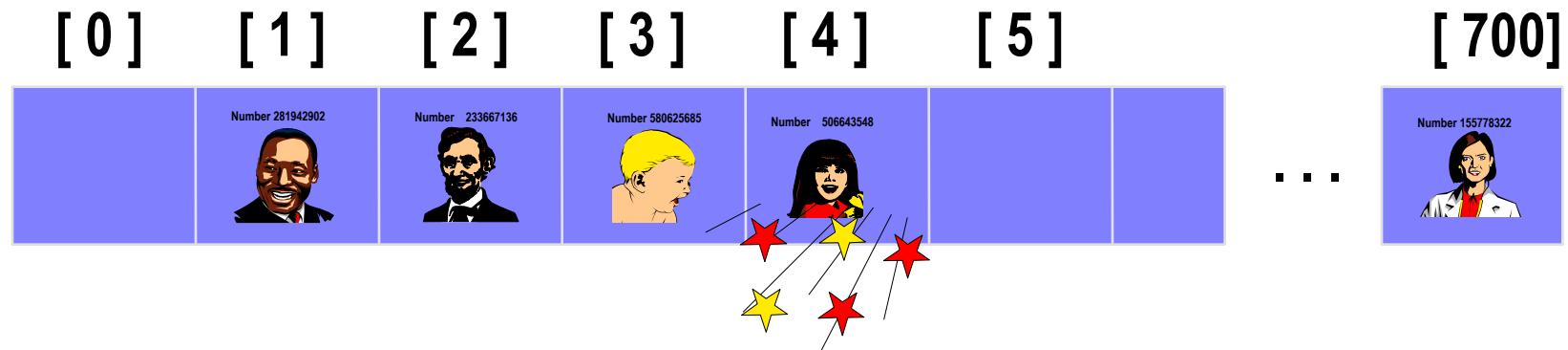
[0] [1] [2] [3] [4] [5] ... [700]



Collisions

- This is called a **collision**, because there is already another valid record at [2].

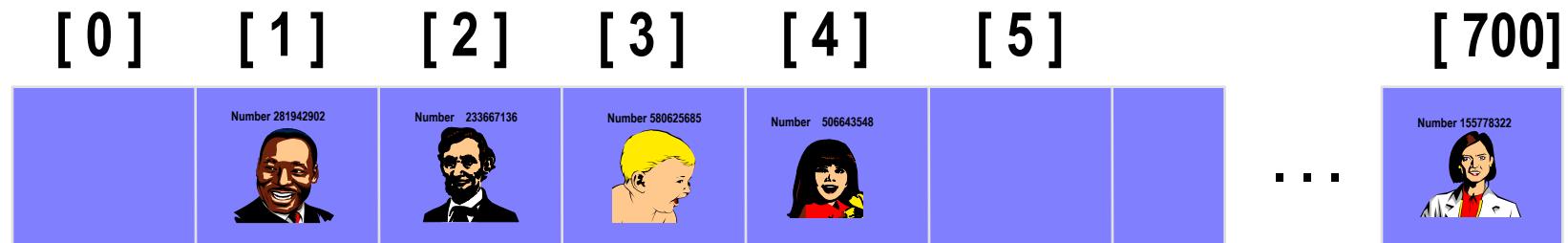
When a collision occurs,
move forward until you
find an empty spot.



Collisions

- This is called a **collision**, because there is already another valid record at [2].

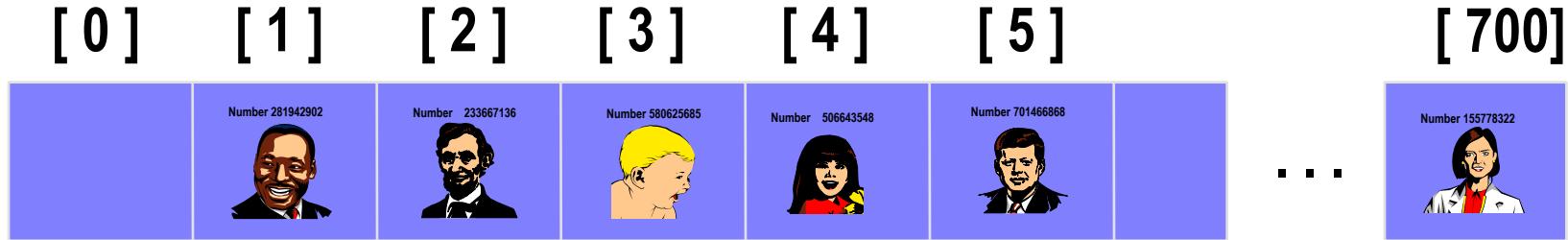
When a collision occurs,
move forward until you
find an empty spot.



Collisions

- This is called a **collision**, because there is already another valid record at [2].

The new record goes in the empty spot.

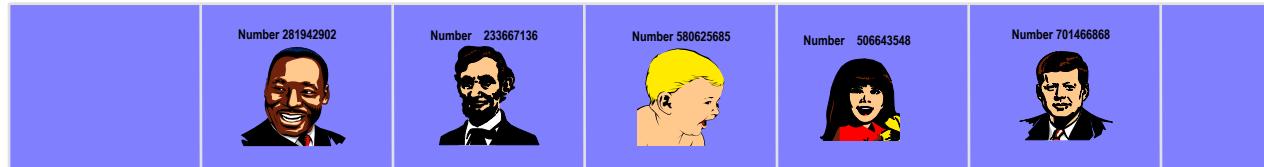


Searching for a Key

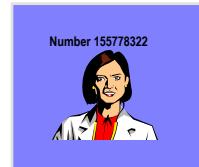
- The data that's attached to a key can be found fairly quickly.



[0] [1] [2] [3] [4] [5] [700]

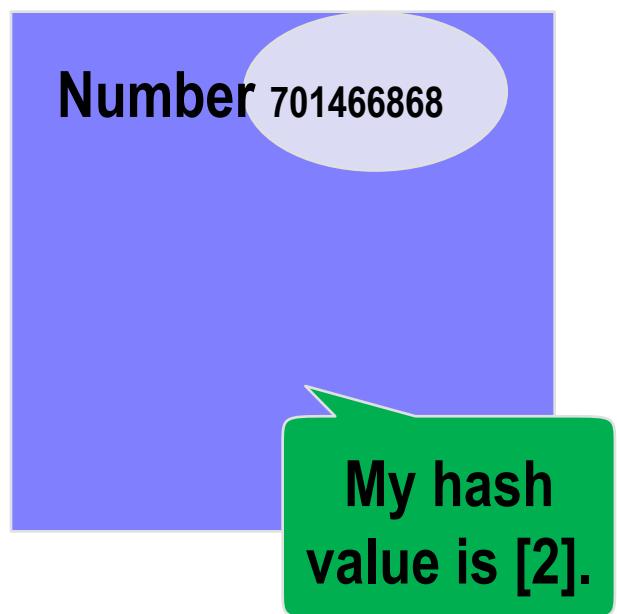


...



Searching for a Key

- Calculate the hash value.
- Check that location of the array for the key.



Not me.

[0]

[1]

[2]

[3]

[4]

[5]

[700]



...



Searching for a Key

- Keep moving forward until you find the key, or you reach an empty spot.



My hash value is [2].

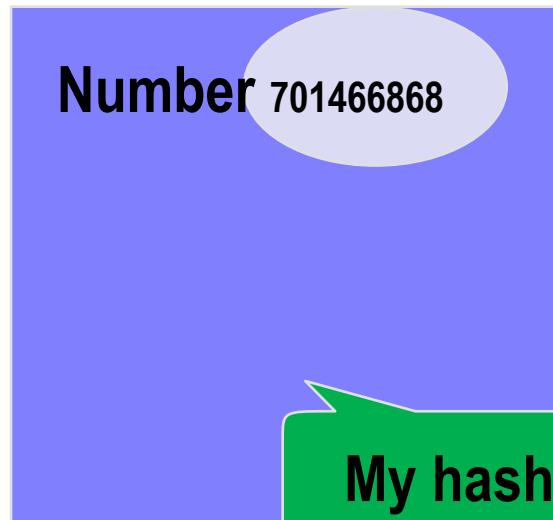
Not me.

[0] [1] [2] [3] [4] [5] ... [700]



Searching for a Key

- Keep moving forward until you find the key, or you reach an empty spot.



My hash value is [2].

Not me.

[0]

[1]

[2]

[3]

[4]

[5]

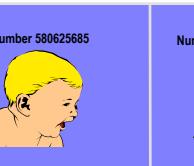
[700]



Number 281942902



Number 233667136



Number 580625685



Number 506643548



Number 701466868

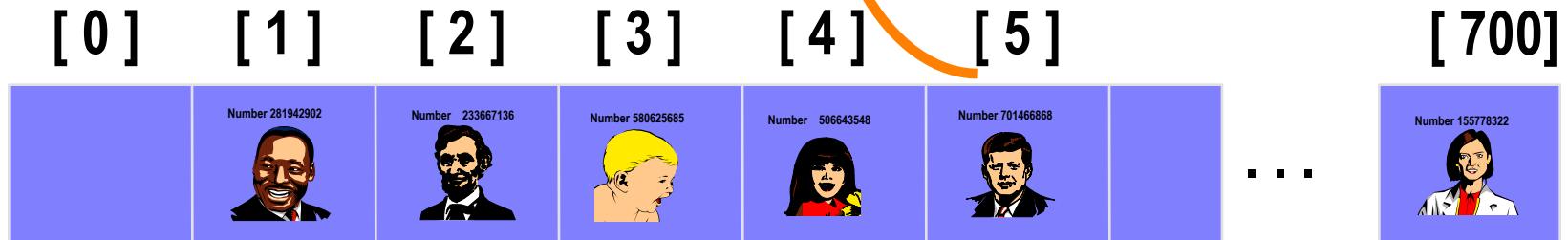
...



Number 155778322

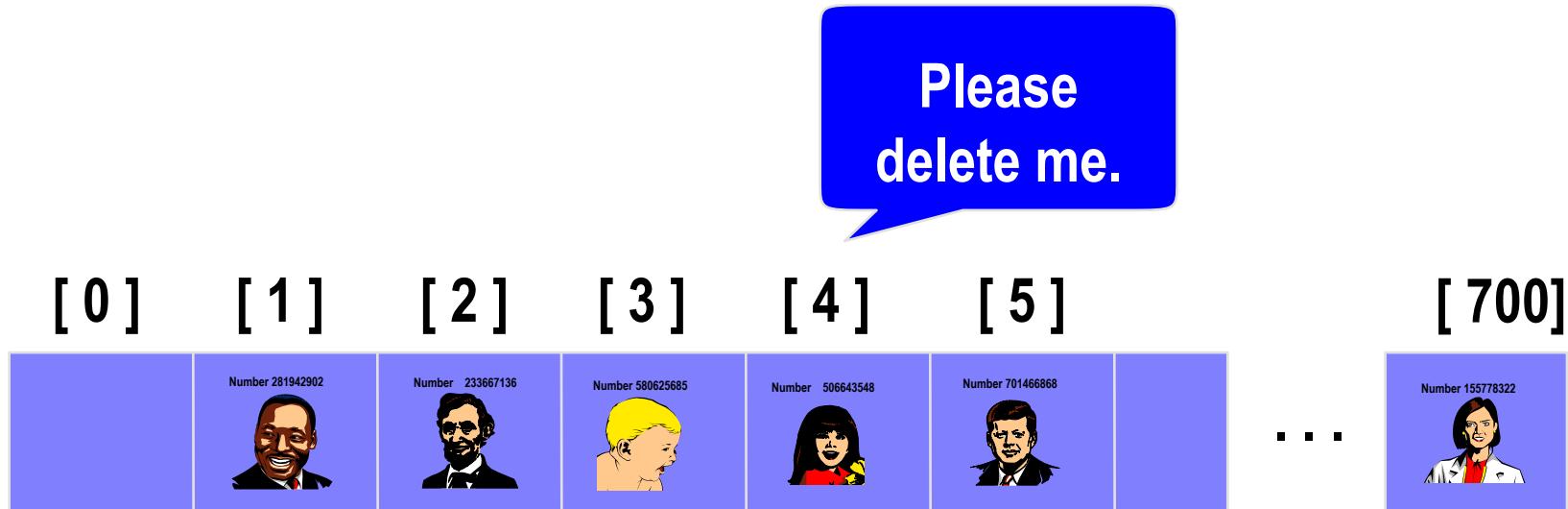
Searching for a Key

- When the item is found, the information can be copied to the necessary location.
- What if you reach an empty spot or the end of array?



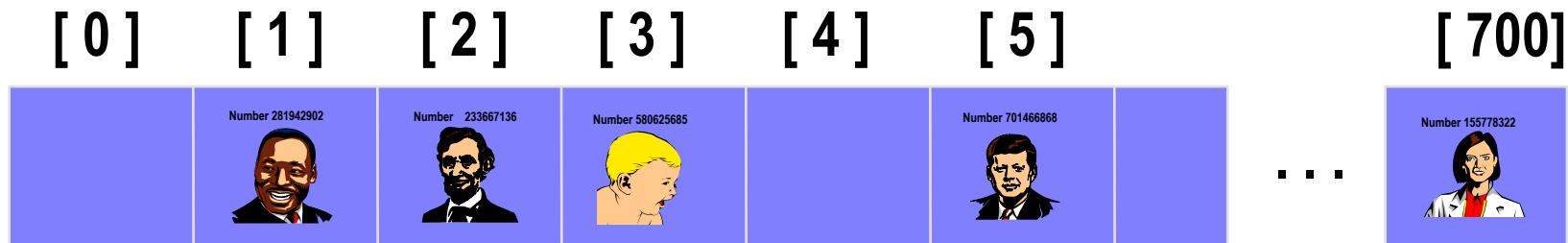
Deleting a Record

- Records may also be deleted from a hash table.



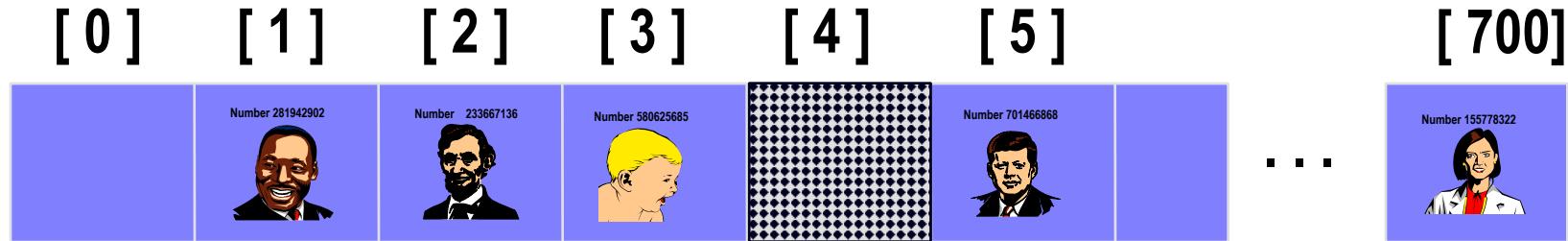
Deleting a Record

- Records may also be deleted from a hash table.
- But the location must not be left as an ordinary "**empty spot**" since that could interfere with searches.



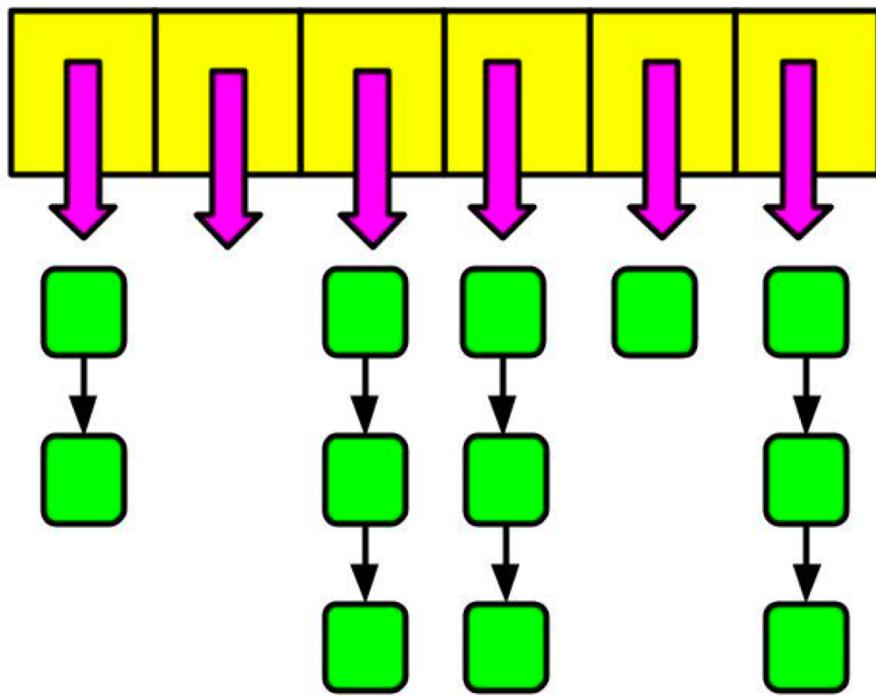
Deleting a Record

- Records may also be deleted from a hash table.
- But the location must not be left as an ordinary "**empty spot**" since that could interfere with searches.
- The location must be marked in some special way so that a search can tell that the spot used to have something in it.

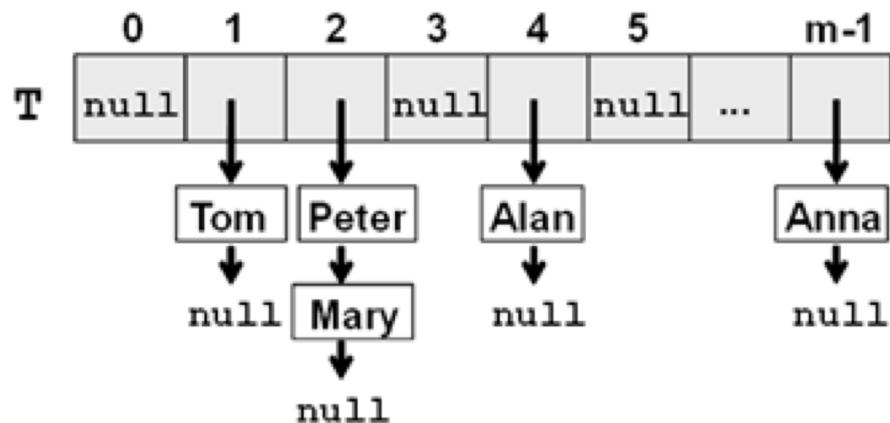


List based Hash Table

Can we do better?



$h("Alan") = 4$
 $h("Peter") = 2$
 $h("Tom") = 1$
 $h("Mary") = 2$
 $h("Anna") = m-1$



Hash Map

- Implement Hash Map

- Hash function just returns the remainder when the key is divided by the hash table size.
- Hash entry (node) has key and value structure.
- In addition, the class contains *search(key)* function to access mapped value by key, *insert(key,value)* function to put key-value pair in table and *remove(key)* function to remove hash node by key.
- For collision resolution, separate chaining strategy could be used.

Map Methods

Delegate operations to a list-based map at each cell:

Algorithm $\text{search}(k)$:

Output: The value of the key k or null

```
bucket = h(k)  
return  $A[\text{bucket}].\text{search}(k)$  {delegate the search to the list-based map at  $A[h(k)]$ }
```

Algorithm $\text{insert}(k, v)$:

```
bucket = h(k)  
 $t = A[\text{bucket}].\text{insert}(k, v)$   
 $n = n + 1$   
return  $t$ 
```

{delegate the put to the list-based map at $A[h(k)]$ }

Algorithm $\text{remove}(k)$:

```
bucket = h(k)  
 $A[\text{bucket}].\text{remove}(k)$   
 $n = n - 1$ 
```

{delegate the remove to the list-based map at $A[h(k)]$ }

Implement Simple Hash Map: testHashMap.cpp

```
1 #include<iostream>
2
3 using namespace std;
4
5 const int TABLE_SIZE = 128;
6
7 // HashEntry Class Declaration
8 class HashEntry
9 {
10 public:
11     int key;
12     int value;
13     HashEntry(int key, int value)
14     {
15         this->key = key;
16         this->value = value;
17     }
18 };
19
20 //HashMap Class Declaration
21 class HashMap
22 {
23 private:
24     HashEntry **table;
25 public:
26     HashMap()
27     {
28         table = new HashEntry * [TABLE_SIZE];
29         for (int i = 0; i < TABLE_SIZE; i++)
30         {
31             table[i] = NULL;
32         }
33     }
34
35     // Hash Function
36     int HashFunc(int key)
37     {
38         return key % TABLE_SIZE;
39     }
40 }
```

```
41     void Insert(int key, int value)
42     {
43         int hash = HashFunc(key);
44         while (table[hash] != NULL && table[hash]->key != key)
45         {
46             hash = HashFunc(hash + 1);
47         }
48         if (table[hash] == NULL)
49             delete table[hash];
50         table[hash] = new HashEntry(key, value);
51     }
52
53     // Search Element at a key
54     int Search(int key)
55     {
56         int hash = HashFunc(key);
57         while (table[hash] != NULL && table[hash]->key != key)
58         {
59             hash = HashFunc(hash + 1);
60         }
61         if (table[hash] == NULL)
62             return -1;
63         else
64             return table[hash]->value;
65     }
66
67     // Remove Element at a key
68     void Remove(int key)
69     {
70         int hash = HashFunc(key);
71         while (table[hash] != NULL)
72         {
73             if (table[hash]->key == key)
74                 break;
75             hash = HashFunc(hash + 1);
76         }
77         if (table[hash] == NULL)
78         {
79             cout<<"No Element found at key "<<key<<endl;
80             return;
81         }
82         else
83         {
84             delete table[hash];
85         }
86         cout<<"Element Deleted"<<endl;
87     }
88 }
```

```

89     ~HashMap()
90     {
91         for (int i = 0; i < TABLE_SIZE; i++)
92         {
93             if (table[i] != NULL)
94                 delete table[i];
95             delete[] table;
96         }
97     }
98 };
99
100 int main()
101 {
102     HashMap hash;
103     int key, value;
104     int choice;
105     while (1)
106     {
107         cout<<"\n-----" << endl;
108         cout<<"Operations on Hash Table" << endl;
109         cout<<"\n-----" << endl;
110         cout<<"1.Insert element into the table" << endl;
111         cout<<"2.Search element from the key" << endl;
112         cout<<"3.Delete element at a key" << endl;
113         cout<<"4.Exit" << endl;
114         cout<<"Enter your choice: ";
115         cin>>choice;
116
117         switch(choice)
118         {
119             case 1:
120                 cout<<"Enter element to be inserted: ";
121                 cin>>value;
122                 cout<<"Enter key at which element to be inserted: ";
123                 cin>>key;
124                 hash.Insert(key, value);
125                 break;
126             case 2:
127                 cout<<"Enter key of the element to be searched: ";
128                 cin>>key;
129                 if (hash.Search(key) == -1)
130                 {
131                     cout<<"No element found at key " << key << endl;
132                     continue;
133                 }
134             else
135             {
136                 cout<<"Element at key " << key << " : ";
137                 cout<<hash.Search(key) << endl;
138             }
139             break;
140         case 3:
141             cout<<"Enter key of the element to be deleted: ";
142             cin>>key;
143             hash.Remove(key);
144             break;
145         case 4:
146             exit(1);
147         default:
148             cout<<"\nEnter correct option\n";
149         }
150     }
151     return 0;
152 }

```

Double Hashing

Double hashing is a collision resolving technique in open addressed hash tables. Double hashing uses the idea of applying a second hash function to key when a collision occurs.

Double hashing can be done using:

$$(\text{hash1(key)} + i * \text{hash2(key)}) \% \text{TABLE_SIZE}$$

First hash function is typically: $\text{hash1(key)} = \text{key \% TABLE_SIZE}$

Second hash function can be: $\text{hash2(key)} = \text{PRIME} - (\text{key \% PRIME})$ where PRIME is a prime smaller than the TABLE_SIZE.

Double Hashing

insert(76)
 $76 \% 7 = 6$

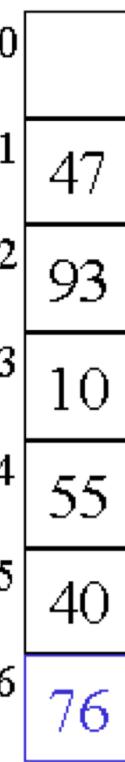
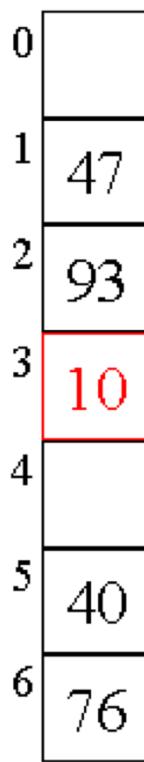
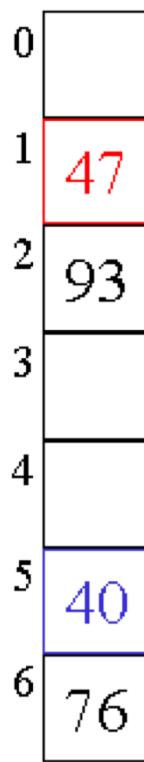
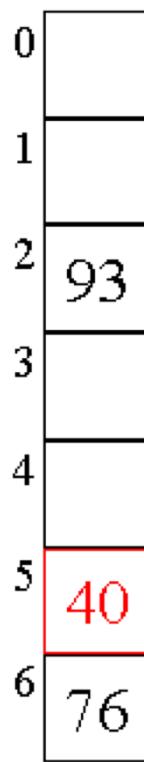
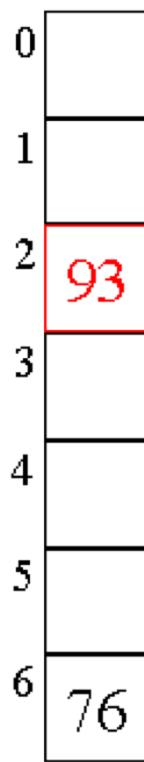
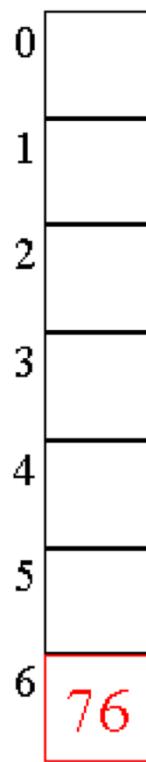
insert(93)
 $93 \% 7 = 2$

insert(40)
 $40 \% 7 = 5$

insert(47)
 $47 \% 7 = 5$
 $5 - (47 \% 5) = 3$

insert(10)
 $10 \% 7 = 3$

insert(55)
 $55 \% 7 = 6$
 $5 - (55 \% 5) = 5$



probes: 1

1

1

2

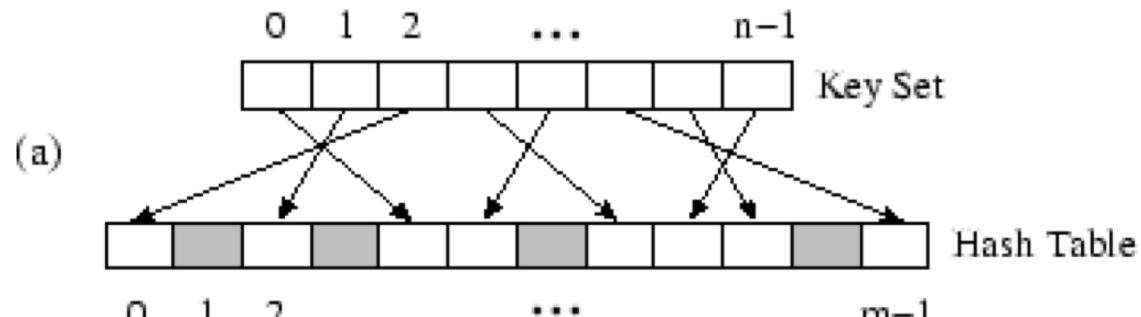
1

2

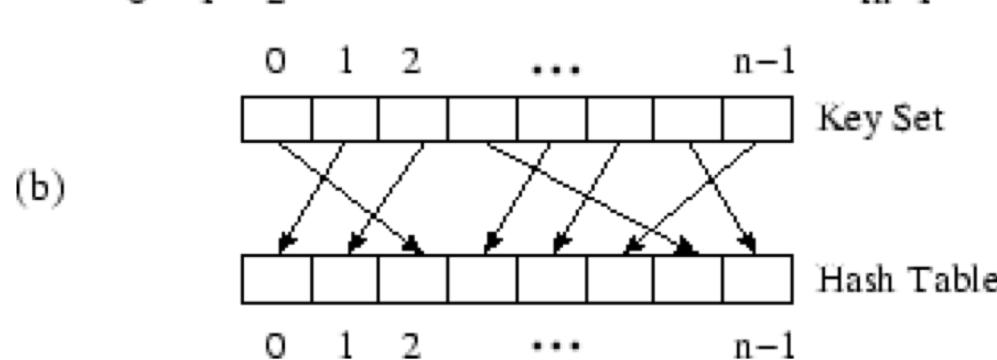
Good Hash Function

- A hash function is **perfect** if it maps items to buckets with no collisions.
- The number of items and all keys should be known to design a perfect hashing that results in constant complexity for insert, search, and remove operations.

Perfect hashing function



Minimal perfect hashing function



<https://emn178.github.io/online-tools/sha256.html>

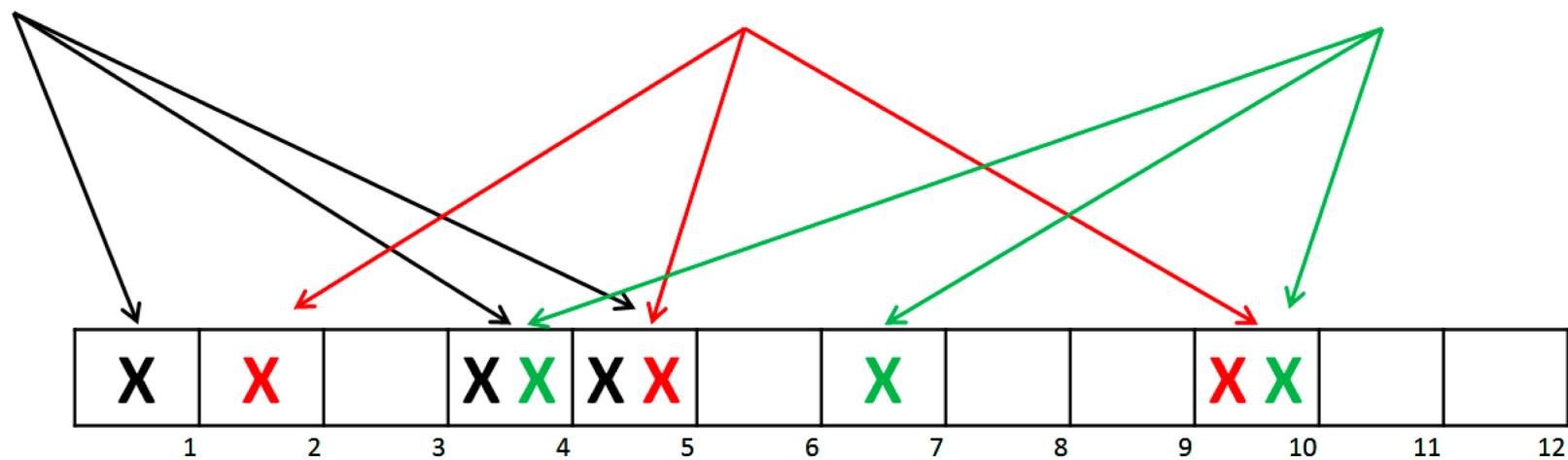
Bloom filter

A space-efficient data structure designed to tell you, rapidly and memory-efficiently, whether an element is present in a set or not with high probability.

$h_1(\text{"oracle"}) = 1$
 $h_2(\text{"oracle"}) = 4$
 $h_3(\text{"oracle"}) = 5$

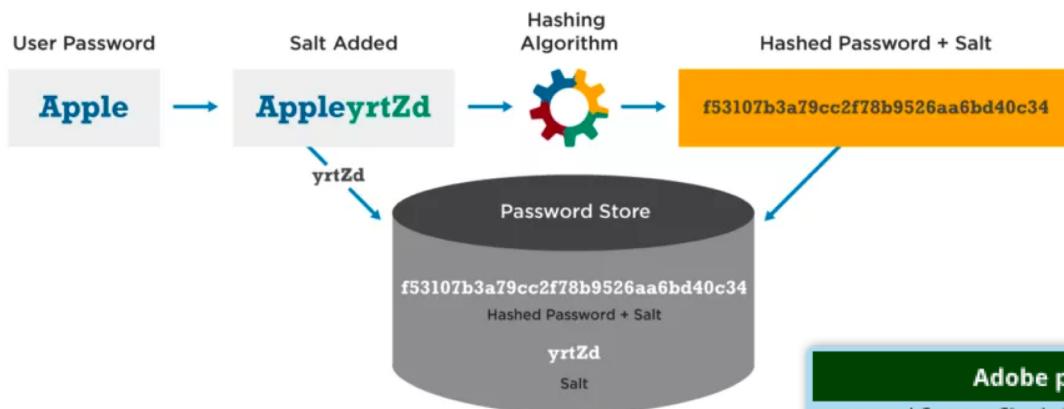
$h_1(\text{"database"}) = 2$
 $h_2(\text{"database"}) = 5$
 $h_3(\text{"database"}) = 10$

$h_1(\text{"filter"}) = 4$
 $h_2(\text{"filter"}) = 7$
 $h_3(\text{"filter"}) = 10$



Storing user passwords

Password Hash Salting



Hashing is used extensively in computer security:

- Authentication
- Signature
- Storing password
- Blockchain

Adobe password data	Password hint
110edf2294fb8bf4	-> numbers 123456
110edf2294fb8bf4	-> ==123456
110edf2294fb8bf4	-> c'est "123456"
8fda7e1f0b56593f e2a311ba09ab4707	-> numbers
8fda7e1f0b56593f e2a311ba09ab4707	-> 1-8
8fda7e1f0b56593f e2a311ba09ab4707	-> 8digit
2fca9b003de39778 e2a311ba09ab4707	-> the password is password
2fca9b003de39778 e2a311ba09ab4707	-> password
2fca9b003de39778 e2a311ba09ab4707	-> rhymes with assword
e5d8efed9088db0b	-> q w e r t y
e5d8efed9088db0b	-> ytrewq tagurpidi
e5d8efed9088db0b	-> 6 long qwert
ecba98cca55eabc2	-> sixxone
ecba98cca55eabc2	-> 1*6
ecba98cca55eabc2	-> sixones

Adobe's password database format made many users' passwords easy to recover