Note Tit	5180- Huffman frees 11/16/2011
	Ana - 100 - 100
	Announcements
	-txam back Monday
	/
	-HW due Sat.
	- Next HW - Checkpoint will be due right after break
	right after break
	- Final: Dec. [7 (Monday) at noon
	The transfer of the transfer o

We want to transmit information using as few bits as possible. Standard ASCII: 8 bits per character Hello: 5×8 bits total Fitended ASCII: 64 bits?

So- how can we do better? What if we don't use every character?

Use fever buts for more common characters. Penalty: Less common characters will need more bits. Problem: Variable length codes

Prefix - free codes

A ON N B EOM An unambiguous way to Send Unformation when we have characters that are not of a fixed length.

No letter's code is the prefix of another letter.

Encode: BAN

100011

Decode:

100011011 6101

BANANA message

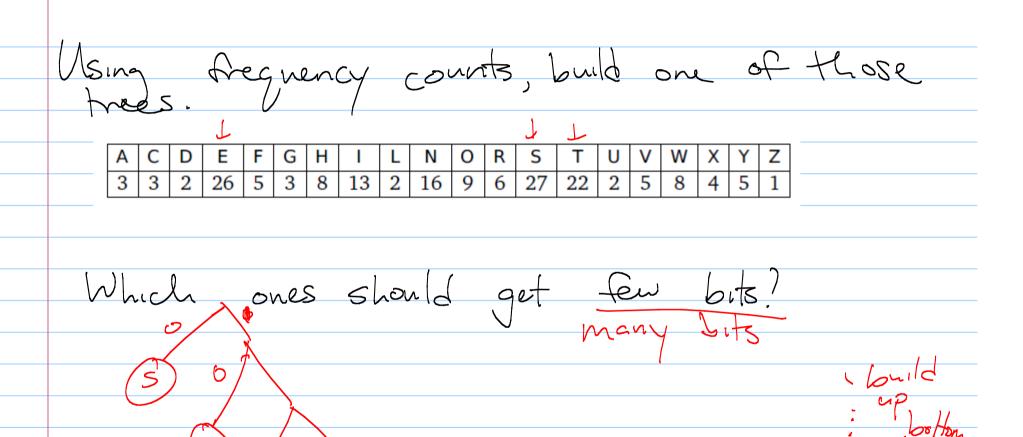
A SEOM

Fren though each letter is different length I scan a use the tree too detect letters, So how do we do this? With exact frequency counts!

This sentence contains three a's, three c's, two d's, twenty-six e's, five f's, three g's, eight h's, thirteen i's, two l's, sixteen n's, nine o's, six r's, twenty-seven s's, twenty-two t's, two u's, five v's, eight w's, four x's, five y's, and only one z.

A C P E (1-1)
3 3 2 26

pull exact letter counts



Huffman's algorithm Take the two least frequent characters.

Merge them into 1 letter, which becomes
a new "leaf". Example:

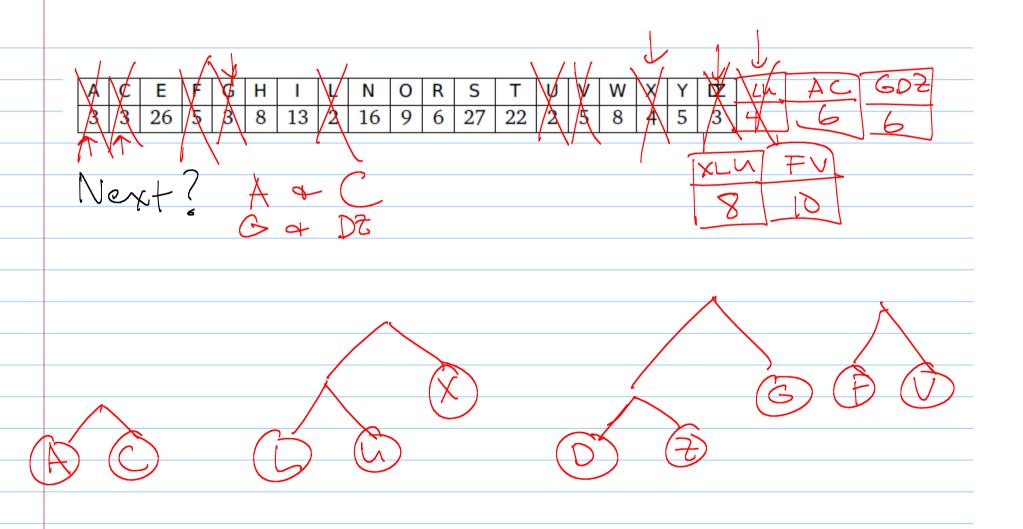
A C D E F G H I L N O R S T U V W X Y Z
3 3 2 26 5 3 8 13 2 16 9 6 27 22 2 5 8 4 5 1

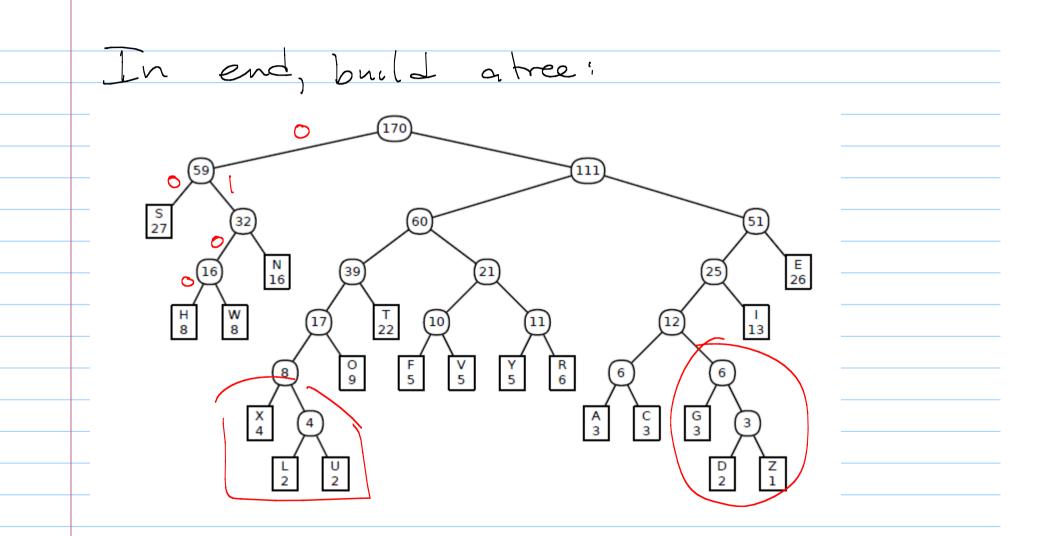
Merge D + 7: 4

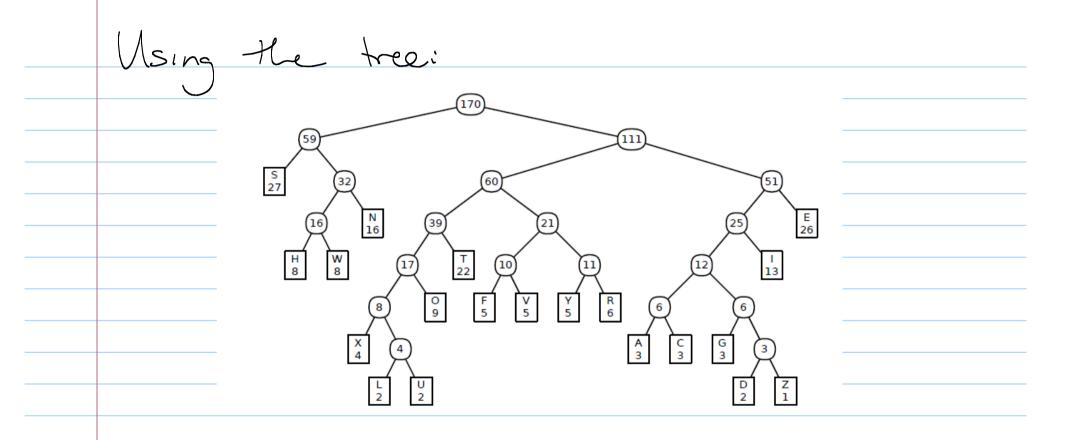
Α	С	Е	F	G	Н	- 1	L	N	0	R	S	Т	U	٧	W	X	Υ	7
3	3	26	5	3	8	13	2	16	9	6	27	22	2	5	8	4	5	3
												4					_	

Merge La U

D+Z, together have Geg. (+2







 $\frac{1001}{\mathsf{T}} \frac{0100}{\mathsf{H}} \frac{1101}{\mathsf{I}} \frac{00}{\mathsf{S}} \frac{00}{\mathsf{S}} \frac{111}{\mathsf{E}} \frac{011}{\mathsf{N}} \frac{1001}{\mathsf{T}} \frac{111}{\mathsf{E}} \frac{011}{\mathsf{N}} \frac{110001}{\mathsf{C}} \frac{111}{\mathsf{E}} \frac{110001}{\mathsf{C}} \frac{10001}{\mathsf{O}} \frac{011}{\mathsf{N}} \frac{1001}{\mathsf{T}} \frac{110000}{\mathsf{A}} \frac{1101}{\mathsf{I}} \dots$

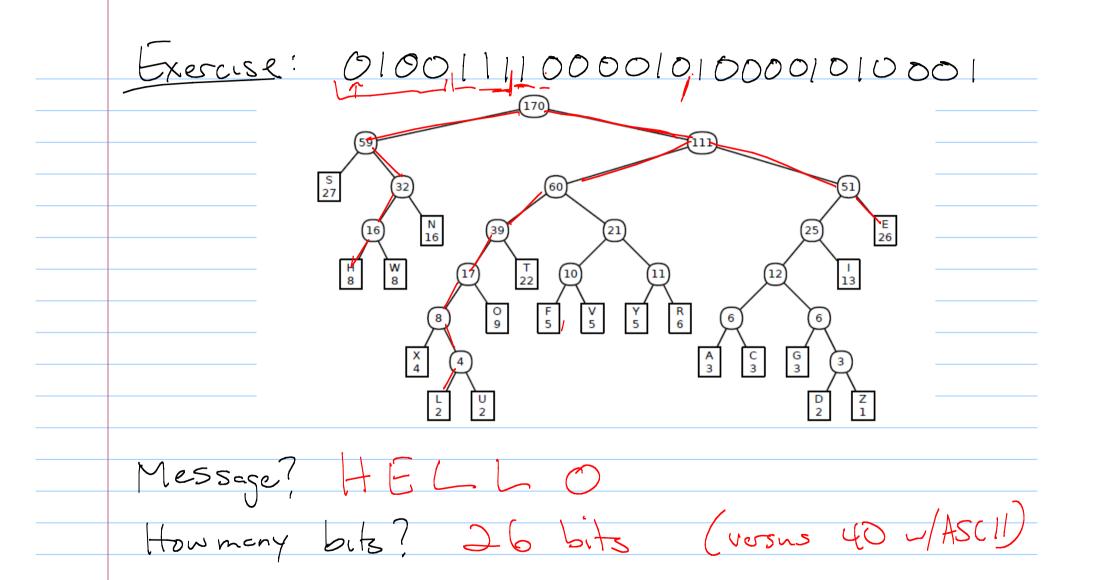
How many bits?

char.	Α	С	D	Е	F	G	Н	Т	L	N	0	R	S	Т	U	٧	W	Χ	Υ	Z
freq.	3	3	2	26	5	3	8	13	2	16	9	6	27	22	2	5	8	4	5	1
depth	6	6	7	3	5	6	4	4	7	3	4	4	2	4	7	5	4	6	5	7
total	18	18	14	78	25	18	32	52	14	48	36	24	54	88	14	25	32	24	25	7

total = 646 bits

How many bits would ASCII use to send these 170 letters?

170×8 (bigger)



Thm: Huffman codes are optimal, in the sense that they use the fewest # of bits possible.

(60 take 314 to see the proof, or read supplemental notes on the schedule page.)

This is a greedy algorithm.

Next program: Decode Given an input, which describes a free and a set of bits which are a message; 1) Create the tree) Use it to de code the message