

CSCI 2100

Heap Recap
Binary Trees
Search Trees



Recap

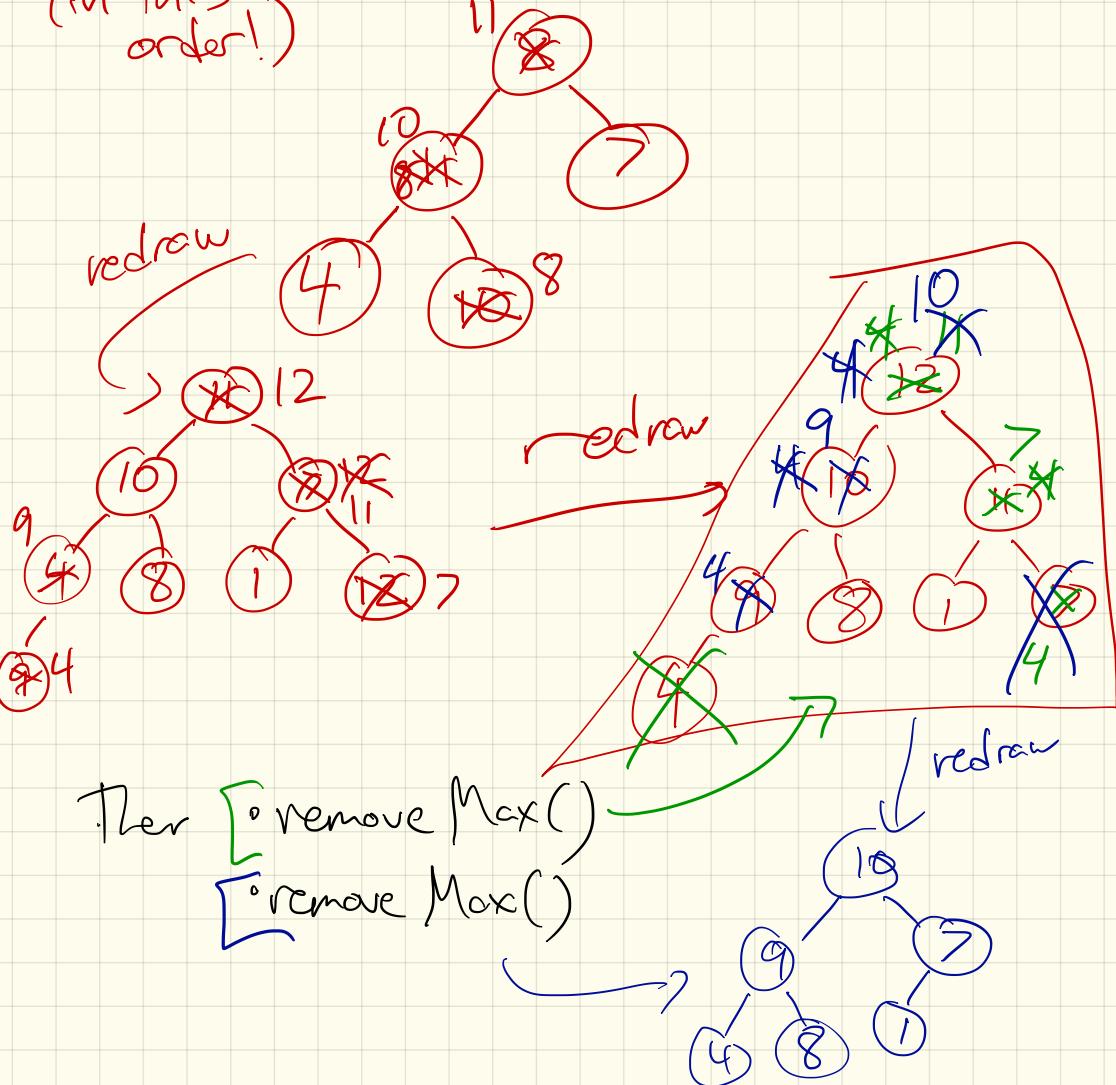
- HW due today
- Next HW - posted after class, due in 1 week
(pen + paper)
- Lab tomorrow

Last "time":

- Trees : binary tree
(height + depth)
- Priority Queues : ADT supporting:
 - insert
 - get Max
 - remove Max
- Heaps:
 - Binary tree implementation
 - value at node is
 \geq children's value
 - all operations are $O(\log n)$

Heap example :

Insert: 8, 11, 3, 4, 10, 1, 12, 9
(in this order!)

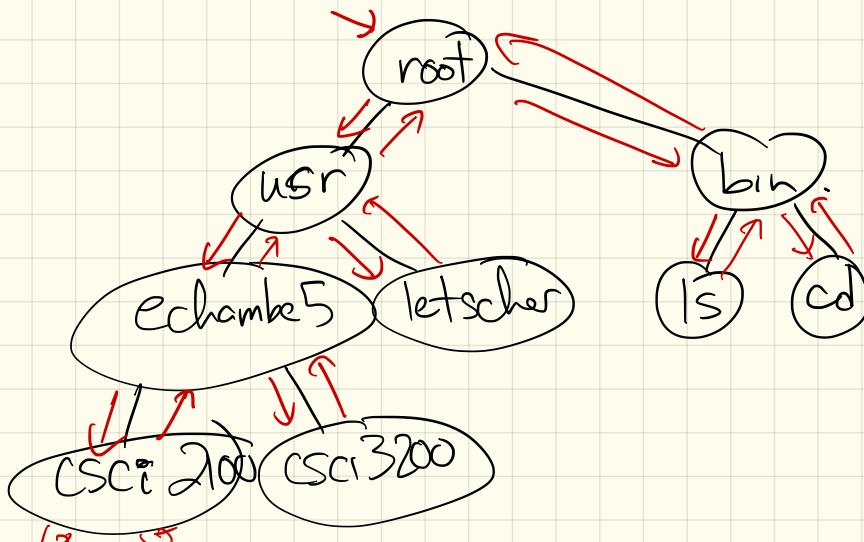


Aside: Tree traversals

3 ways to traverse a tree:
Starting at root:

Preorder (v) =

print → VISIT
Preorder ($v \rightarrow \text{left}$)
Preorder ($v \rightarrow \text{right}$)



traversal order: root, usr, ecambe5,
CSci2100, CSci3200,
letscher, bin, ls, cd,

(Start at root)

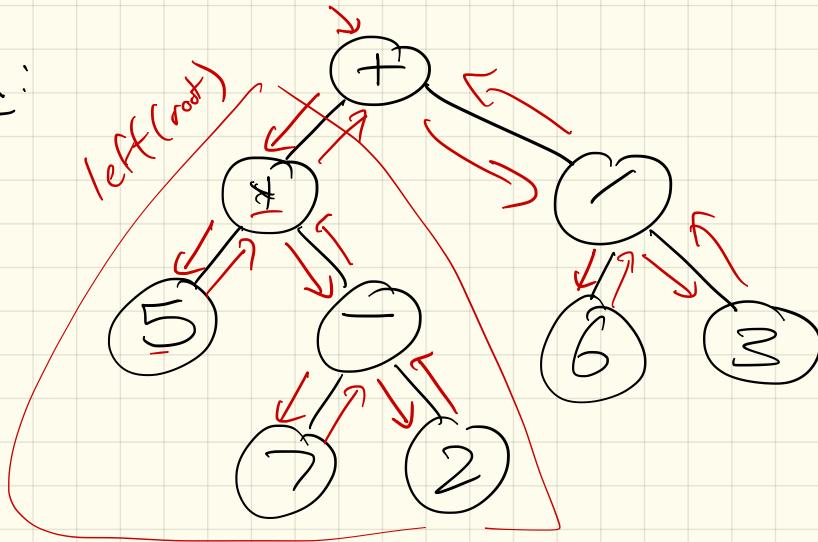
In order (v) =

Inorder ($v \rightarrow \text{left}$)

visit v

Inorder ($v \rightarrow \text{right}$)

Ex:



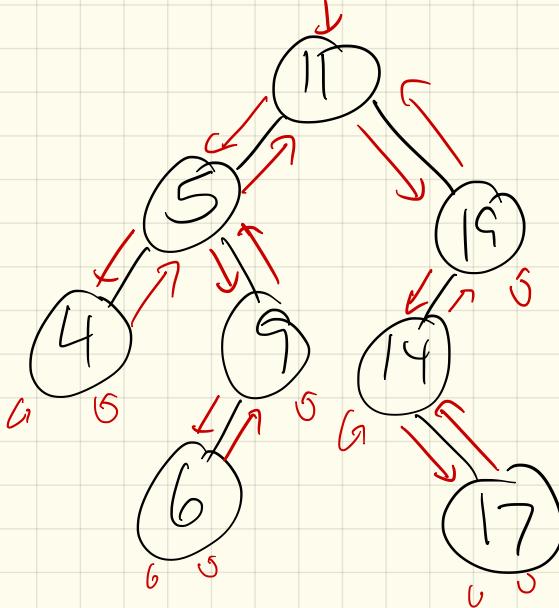
traversal: 5, *, 7, -, 2, +, 6, /, 3

"middle"

Postorder(v) =

post order($v \rightarrow \text{left}$)
post order($v \rightarrow \text{right}$)
visit v

Ex:

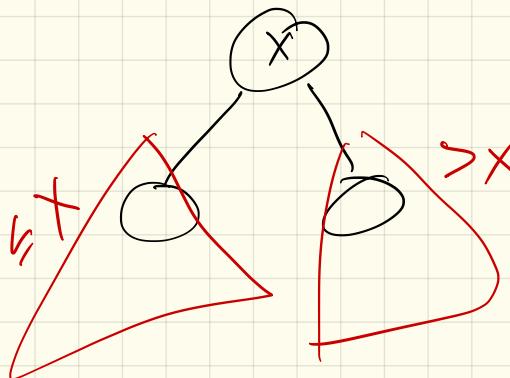


traversal : 4, 6, 9, 5, 17, 14, 19,
 11

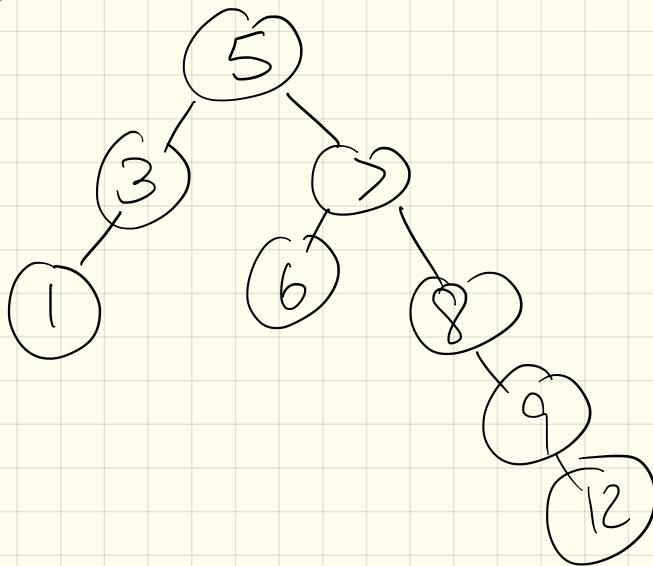
Next: Binary Search Trees

A binary tree where we maintain the following:

- The value at any node is \geq its left child and $<$ its right child.

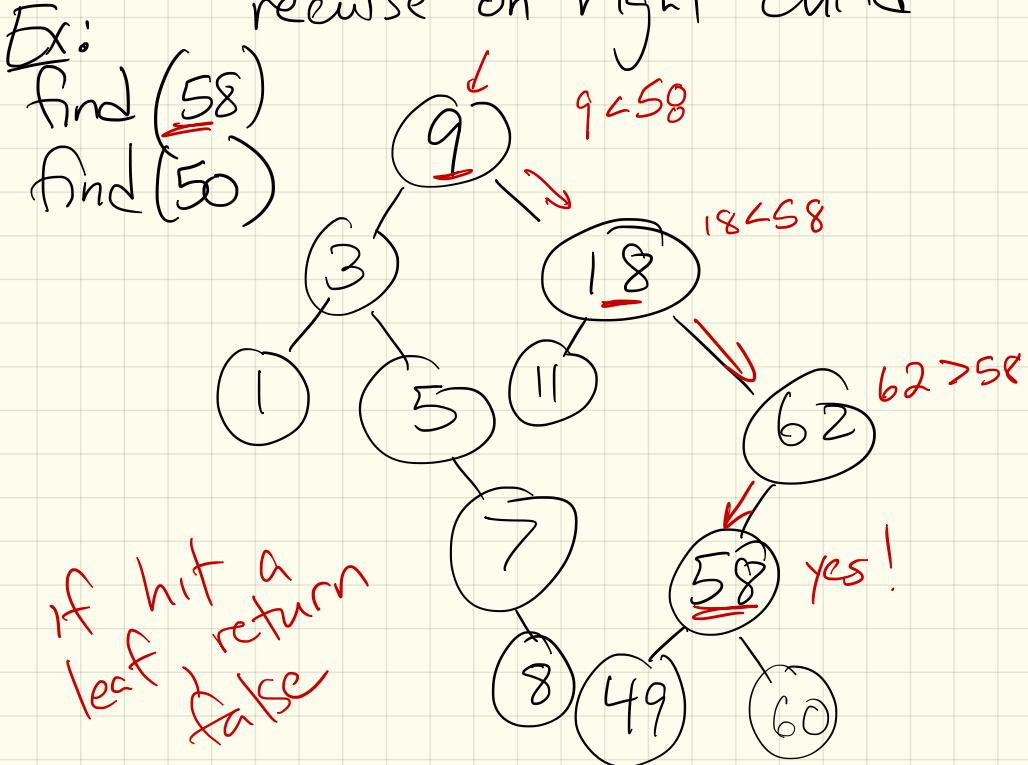


Ex:



Finding in a BST:

- Check if root = target value
return true
- If root $>$ target
 - If left != NULL
else recurse on left child
- If root $<$ target
recurse on right child

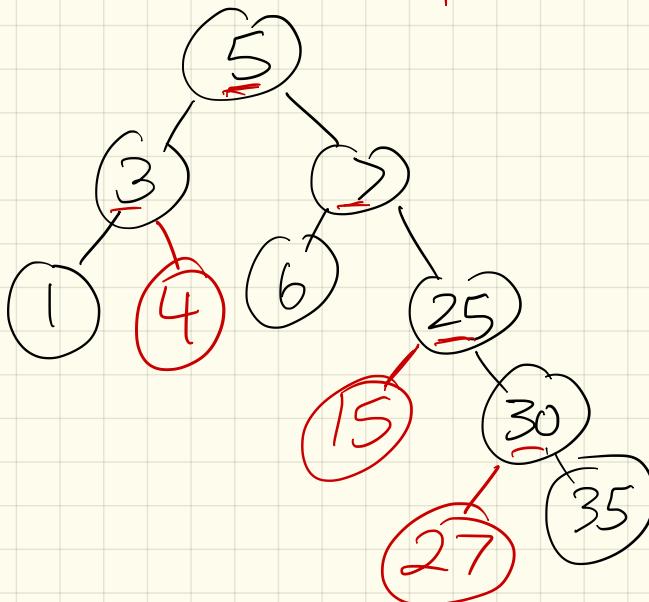


Inserting:

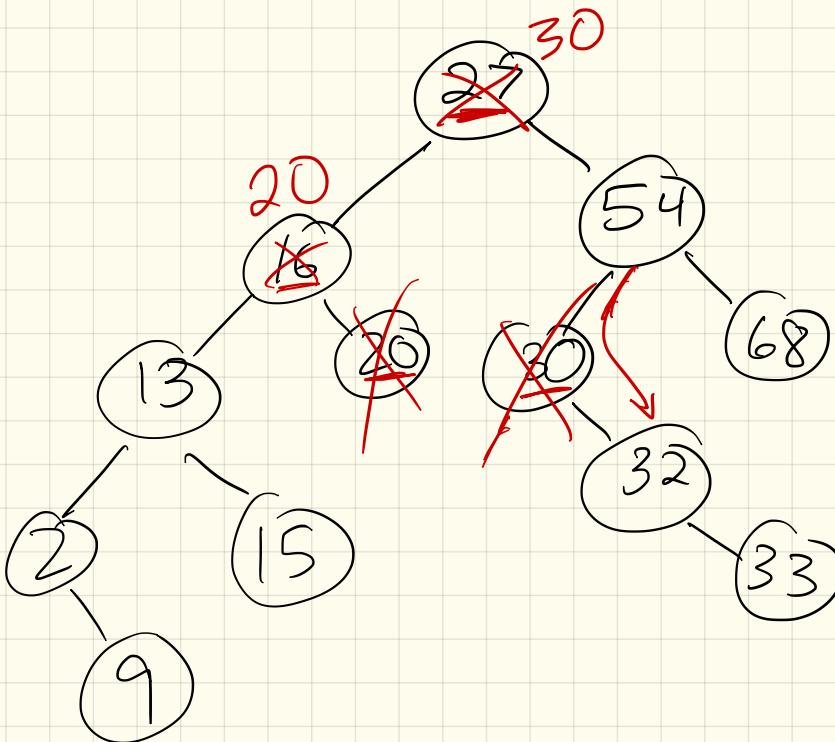
Given a BST, insert is done by finding the (unique) leaf location where the value fits.

Ex: Insert (4)
Insert (27)
Insert (15)

find, but when
hit a null, insert
the new node



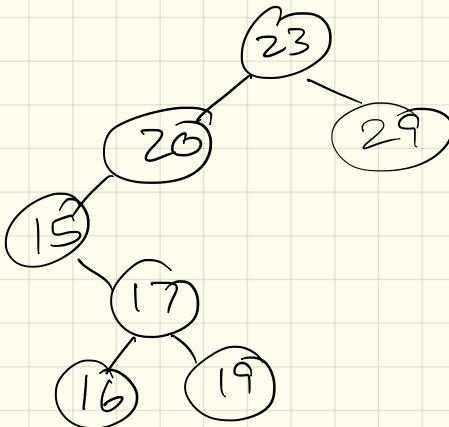
Delete: More Complex



remove(16) : delete 20's node
copy 20 into 16's node
remove(27) copy 30 up to root

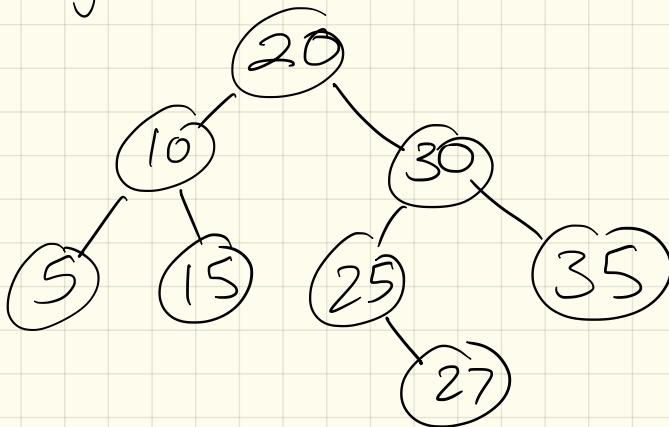
Cases:

- Target could be a leaf
- Target could have only 1 child: remove(15)



Cases (cont) :

- Target could have 2 children



remove (20) :

Remove(x):

Find (x)

If x is a leaf
delete x 's node

else if x has one child

can remove x 's node
and "promote" its
only child to its spot

Else:

Find smallest value $> x$

(Note: this is next value
in an inorder traversal!)

Copy that value to x 's spot
& remove that node,
promoting its right child
in its spot (if it has one)
Note: it will have no
left child

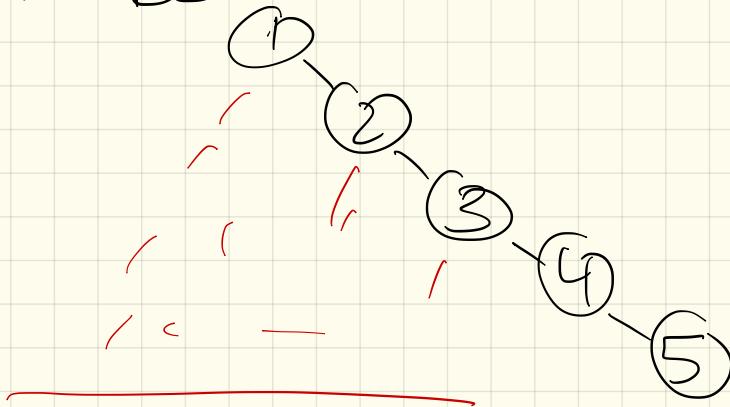
Code:

- Pointer based.

Reason: we need to move around entire subtrees.

Also, tree is not complete!

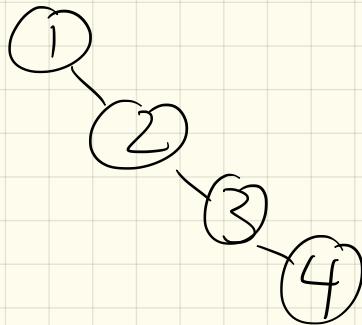
Valid BST:



not space efficient

Note: BSTs are not unique!

Consider



Can you make another BST
with these values?