Dueling bandits

immediate

September 26, 2014

We define the set $T_p = \{2^{p-1}, \dots, 2^p - 1\} = \{s \in \mathbb{N} : \lfloor \log_2 s \rfloor = p - 1\}.$

Algorithm 1: Improved Doubler

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initialization x_1 fixed in X, \mathcal{L} = \{x_1\}, \hat{f_0} = 0; t \leftarrow 1; p \leftarrow 1; while true do

for j = 1 to 2^{p-1} do

choose x_t uniformly from \mathcal{L}; y_t \leftarrow \text{advance}(S); play (x_t; y_t), observe choice b_t; feedback \left(S; b_t + \hat{f}_{p-1}\right); t \leftarrow t + 1;

\mathcal{L} the multi-set of arms played as y_t in the last for-loop; \hat{f_p} \leftarrow \hat{f_p} + \sum_{s \in T_p} b_s/2^{p-1} - 1/2; p \leftarrow p + 1;
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Observe that if $t \in T_p$

$$\mathbb{E}\left[b_t\left|\left\{y_s, s \in T_{p-1}\right\}, y_t\right.\right] \ = \ \sum_{s \in T_{p-1}} \frac{\mu\left(y_t\right) - \mu\left(y_s\right) + 1}{2^{p-1}} = \frac{\mu\left(y_t\right) + 1}{2} - \sum_{s \in T_{p-1}} \frac{\mu\left(y_s\right)}{2^{p-1}} \; ,$$

and that

$$\mathbb{E}\left[\sum_{s \in T_{p-1}} b_s/2^{p-2} - 1/2 \left| \bigcup_{r=p-2}^{p-1} \{y_s, s \in T_r\} \right] = \sum_{s \in T_{p-1}} \frac{\mu(y_s)}{2^{p-1}} - \sum_{s \in T_{p-2}} \frac{\mu(y_s)}{2^{p-2}} \right].$$

Let us denote $f_t = b_t + \hat{f}_{p-1} = b_t + \sum_{r=1}^{\lfloor \log_2 t \rfloor} \sum_{s \in T_r} b_s / 2^{r-1} - \lfloor \log_2 t \rfloor / 2$ the feedback that we introduce in S. Using the recurrence defining \hat{f}_p we obtain

$$\mathbb{E}\left[f_{t} \middle| x_{1}, \bigcup_{r=1}^{p-1} \{y_{s}, s \in T_{r}\}, y_{t}\right] = \frac{\mu(y_{t}) - \mu(x_{1}) + 1}{2}.$$

Since the above right term is $\sigma(x_1, y_t)$ - measurable we conclude that

$$\mathbb{E}\left[f_t|x_1,y_t\right] = \frac{\mu(y_t) - \mu(x_1) + 1}{2}.$$

Let $y_{t_1} = \dots = y_{t_k}$ and let $f = \sum_{j=1}^k f_{t_j}/k$. Observe that $f = \sum_{s=1}^{t_k} a_s b_s/t_k - \sum_{j=1}^k \log_2 t_j/2k$ with $a_s \in [0, A_s]$. We will later specify this bound. Since the $a_s b_s$ are independent and $\mathbb{P}[a_s b_s \in [0, A_s]] = 1$ we can apply the Hoeffding's inequality:

$$\mathbb{P}\left[f - \mathbb{E}\left[f\right] \ge \epsilon \left| x_1, y_{t_1} = \dots = y_{t_k} \right| \le \exp\left(-\frac{2t_k^2 \epsilon^2}{\sum_{s=1}^{t_k} A_s^2}\right)$$

Assume the convention $t_0=1$ and set $S_j=\{s\in\mathbb{N}:2^{\lfloor\log_2t_{j-1}\rfloor}\leq s<2^{\lfloor\log_2t_{j}\rfloor}\}$ for $j=1,\ldots,k$ and $S_{k+1}=\{s\in\mathbb{N}:2^{\lfloor\log_2t_k\rfloor}\leq s< t_k\}$. For each $1\leq j\leq k+1$, $s\in S_j$, $A_s=t_k((k-j+1)/2^{\lfloor\log_2s\rfloor}+1)/k$ if $s=t_i$ for some $1\leq i< j$ and $A_s=t_k(k-j+1)/(2^{\lfloor\log_2s\rfloor}k)$ otherwise. We say that $t_j\in S_{t(j)}$. The function t is obviously non decreasing, $t(j)\geq j+1$ and t(k)=k+1.

$$\begin{split} \frac{k^2}{t_k^2} \sum_{s=1}^{t_k} A_s^2 &= \sum_{j=1}^{k+1} \sum_{s \in S_j} \frac{(k-j+1)^2}{2^{2\lfloor \log_2 s \rfloor}} + \sum_{j=1}^k \frac{k-t(j)+1}{2^{\lfloor \log_2 t_j \rfloor -1}} + k = \\ & \sum_{j=1}^k (k-j+1)^2 \left(\frac{1}{2^{\lfloor \log_2 t_{j-1} \rfloor +1}} - \frac{1}{2^{\lfloor \log_2 t_j \rfloor +1}} \right) + \sum_{j=1}^k \frac{4(k-t(j))+4}{2^{\lfloor \log_2 t_j \rfloor +1}} + k \\ &= \frac{k^2}{2} - \sum_{j=1}^k \frac{2(k-j)+1}{2^{\lfloor \log_2 t_j \rfloor +1}} + \sum_{j=1}^k \frac{4(k-t(j))+4}{2^{\lfloor \log_2 t_j \rfloor +1}} + k \; . \end{split}$$

This implies that

$$k^2 \geq \frac{k^2}{2} + \sum_{j=1}^k \frac{2(k-j)-1}{2^{\lfloor \log_2 t_j \rfloor + 1}} + k \geq \frac{k^2}{t_k^2} \sum_{s=1}^{t_k} A_s^2 \geq \frac{k^2}{2} - \sum_{j=1}^k \frac{2(k-j)+1}{2^{\lfloor \log_2 t_j \rfloor + 1}} \geq \frac{k^2}{2} \left(1 - \frac{1}{t_k}\right) \; .$$

To obtain the convergence of Improved Doubler using UCB at the same rate as UCB we need that $\sum_{s=1}^{t_k} A_s^2 = O(t_k^2/k)$, but we just showed that it is not possible.

$$\sum_{s=1}^{t} t^{-\beta/s} \tag{0.1}$$