

Assignment 3 Model Solutions

1) (20 pts.) Do Exercises 40 and 50 of Section 5.1 (page 331).

Graded by : Sneha Balasubramanian

Exercise 40:

Let $P(n) \rightarrow (A_1 \cap A_2 \cap \dots \cap A_n) \cup B = (A_1 \cup B) \cap (A_2 \cup B) \cap \dots \cap (A_n \cup B)$

Basic step:

For $n=1$,

$P(1) = (A_1) \cup B = A_1 \cup B$. This is a trivial case. so we consider $n=2$ as our base case

For $n=2$,

$P(2) = (A_1 \cap A_2) \cup B = (A_1 \cup B) \cap (A_2 \cup B)$ which is true.

Assumption step:

For $n=k$,

Let us assume $P(k) \rightarrow (A_1 \cap A_2 \cap \dots \cap A_k) \cup B = (A_1 \cup B) \cap (A_2 \cup B) \cap \dots \cap (A_k \cup B)$ is true.

Inductive step:

We shall prove that $P(k+1)$ is true. (ie)

$P(k+1) \rightarrow (A_1 \cap A_2 \cap \dots \cap A_{k+1}) \cup B = (A_1 \cup B) \cap (A_2 \cup B) \cap \dots \cap (A_{k+1} \cup B)$

Let us consider $C = (A_1 \cap A_2 \cap \dots \cap A_k)$

thus, $P(k+1) \rightarrow (A_1 \cap A_2 \cap \dots \cap A_k \cap A_{k+1}) \cup B = (C \cap A_{k+1}) \cup B$

From the basic step, $(C \cap A_{k+1}) \cup B = (C \cup B) \cap (A_{k+1} \cup B)$ is true.

$(C \cup B) \cap (A_{k+1} \cup B) = ((A_1 \cap A_2 \cap \dots \cap A_k) \cup B) \cap (A_{k+1} \cup B)$

$((A_1 \cap A_2 \cap \dots \cap A_k) \cup B) \cap (A_{k+1} \cup B) = ((A_1 \cup B) \cap (A_2 \cup B) \cap \dots \cap (A_k \cup B)) \cap (A_{k+1} \cup B)$

$(A_1 \cap A_2 \cap \dots \cap A_k \cap A_{k+1}) \cup B = (A_1 \cup B) \cap (A_2 \cup B) \cap \dots \cap (A_k \cup B) \cap (A_{k+1} \cup B)$

this proves $P(k) \rightarrow P(k+1)$, making it true for any value of n .

Exercise 50:

the basic step is wrong in this problem where $n=1$. this case is fails.

$$\sum_{i=1}^n i = \frac{(n + \frac{1}{2})^2}{2}$$

for $n=1$,

$$= \sum_{i=1}^1 i = \frac{(1 + \frac{1}{2})^2}{2} = \frac{9}{8}$$

$$1 \neq \frac{9}{8}$$

thus this induction proof is wrong.

2) (20 pts.) Do Exercises 38 and 40 of Section 5.4 (page 371).

Graded by : Sneha Balasubramanian

Exercise 38:

let λ be the empty string and let the concatenation of strings w_1 and w_2 be $w_1.w_2$.

procedure concatenation(w: string, i: copies of string w)

$\lambda := ""$

if i=0 **then return** λ

else if i=1 **then return** w

else w.concatenation(w, i-1)

Exercise 40:

We have to prove that algorithm concatenation(w,i) is true for all values of i.

Basic step:

For i=0,

The algorithm will return λ meaning there are zero copies of w.

For i=1,

The algorithm will return w indicating there is one copy of w.

Assumption step:

We assume for $n=k$, the algorithm concatenation(w,k) will return w^k which gives k copies of string w.

Induction step:

For $n=k+1$,

the algorithm concatenation(w,k+1) will return w.concatenation(w,k).

We know that concatenation(w,k) will return w^k . Thus, w.concatenation(w,k) will return $w.w^k = w^{k+1}$ which gives k+1 copies of w.

Since concatenation(w,k+1) returns the correct result if concatenation(w,k) is true, this recursive is correct for finding w^i .

3) (20 pts.) Graded by : Ramya Vulimiri

Let $P(n)$ be the statement that given n students in a class, you can find the popular student (if there is one) with at most $3(n - 1)$ questions.

Basis Step: You can show either $P(1)$ or $P(2)$. If $n=1$, then that student is the popular student, needing 0 questions.

$0 \leq 3(1 - 1) = 0 \Rightarrow P(1)$ holds. If $n=2$, say there are two students A and B. To find if there is a popular student, ask A if he follows B and ask B if he follows A - 2 questions.

$2 \leq 3(2 - 1) = 3 \Rightarrow P(2)$ holds.

Inductive Step: Assume $P(k)$ for an arbitrary $k \geq 2$, that the popular student in a class of k students can be found with $3(k-1)$ questions. Show $P(k) \rightarrow P(k+1)$, where $P(k+1) :=$ given $k+1$ students, you can find the popular student (if there is one) with at most $3(k)$ questions.

Let us pick two random students A and B from $(k+1)$ students and ask A if (s)he follows B. Either A follows B, in which case A is not a popular student or A doesn't follow B, in which case B is not a popular student. In either case, one student is eliminated as popular.

Without loss of generality, let us say A is the student who is not popular. In the remaining k people, the popular student can be found with $3(k-1)$ questions (inductive hypothesis).

Case 1: If there is no popular student among k students, then there are no popular students among $k+1$ students.

No. of questions needed = $1 + 3(k - 1) = 3k - 2 \leq 3k$.

Case 2: If there is a popular student, let's call him P, we need to determine if P is still popular among $k+1$ students. For that, we ask 2 questions, if P follows A and if A follows P.

No. of questions needed = $1 + 3(k - 1) + 2 = 3k - 3 + 3 \leq 3k$.

From case 1 and 2, $P(k+1)$ holds. Therefore, $P(n)$ holds for $n \geq 1$.

4) (15 pts.) Graded by : Ramya Vulimiri

Let $P(n)$ be the proposition that $\sum_{i=0}^{n-1} \frac{i}{2^i} = 2 - \frac{n+1}{2^{n-1}}$. We must do two things to prove that $P(n)$ is true for $n = 1, 2, 3, \dots$. Namely, we must show that $P(1)$ is true and that the conditional statement $P(k)$ implies $P(k+1)$ is true for $k = 1, 2, 3, \dots$.

BASIS STEP: $P(1)$ is true, because $\text{LHS} = \sum_{i=0}^0 \frac{i}{2^i} = \frac{0}{1} = 0$ and $\text{RHS} = 2 - \frac{1+1}{2^0} = 2 - 2 = 0$.

INDUCTIVE STEP: For the inductive hypothesis we assume that $P(k)$ holds for an arbitrary positive integer k .

That is, we assume that

$$\sum_{i=0}^{k-1} \frac{i}{2^i} = 2 - \frac{k+1}{2^{k-1}}$$

Under this assumption, it must be shown that $P(k+1)$ is true, namely, that

$$\sum_{i=0}^{(k+1)-1} \frac{i}{2^i} = 2 - \frac{(k+1)+1}{2^{(k+1)-1}}$$

$$\begin{aligned}
\sum_{i=0}^{(k+1)-1} \frac{i}{2^i} &= \sum_{i=0}^k \frac{i}{2^i} = \sum_{i=0}^{k-1} \frac{i}{2^i} + \frac{k}{2^k} \\
&= 2 - \frac{k+1}{2^{k-1}} + \frac{k}{2^k} \text{ (Applying IH)} \\
&= 2 - \frac{2(k+1) - k}{2^k} \\
&= 2 - \frac{(k+1) + 1}{2^{(k+1)-1}}
\end{aligned}$$

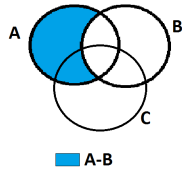
This last equation shows that $P(k+1)$ is true under the assumption that $P(k)$ is true. This completes the inductive step. We have completed the basis step and the inductive step, so by mathematical induction we know that $P(n)$ is true for all positive integers n .

5) (15 points) Graded By : Rashmi Soni

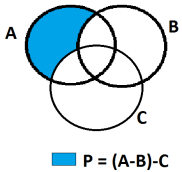
Given, the set $P = (A-B)-C$ and $Q = (A-C)-(B-C)$

1) Drawing the venn diagram

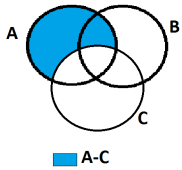
For A-B :



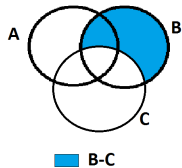
So, the venn diagram for $P = (A-B)-C$:



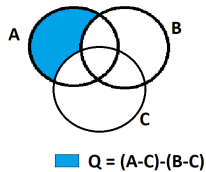
Now, drawing the venn diagram for A-C :



Now, drawing the venn diagram for B-C :



Thus, the venn diagram for $Q = (A-C)-(B-C)$ is :



We can see that the venn diagrams of sets P and Q are equal. Thus, P=Q.

2)Constructing membership table

A	B	C	A - B	P = (A - B) - C	A - C	B - C	Q = (A - C) - (B - C)
1	1	1	0	0	0	0	0
1	1	0	0	0	1	1	0
1	0	1	1	0	0	0	0
1	0	0	1	1	1	0	1
0	1	1	0	0	0	0	0
0	1	0	0	0	0	1	0
0	0	1	0	0	0	0	0
0	0	0	0	0	0	0	0

So, we can conclude from the membership table that P = Q.

3)Using set identities with set builder notation

We can write,

$$A-B = \{x|x \in A \wedge x \notin B\}$$

$$\text{So, } (A-B)-C = \{x|x \in (A - B) \wedge x \notin C\}$$

$$= \{x|(x \in A \wedge x \notin B) \wedge x \notin C\}$$

Now,

$$A-C = \{x|x \in A \wedge x \notin C\}$$

$$B-C = \{x|x \in B \wedge x \notin C\}$$

Thus, we can write

$$(A-C)-(B-C) = \{x|x \in (A - C) \wedge x \notin (B - C)\}$$

$$= \{x|(x \in A \wedge x \notin C) \wedge \neg(x \in B \wedge x \notin C)\}$$

$$= \{x|(x \in A \wedge x \notin C) \wedge (x \notin B \vee x \in C)\}$$

$$= \{x|((x \in A \wedge x \notin C) \wedge x \notin B) \vee ((x \in A \wedge x \notin C) \wedge x \in C)\}$$

$$= \{x|((x \in A \wedge x \notin C) \wedge x \notin B) \vee (x \in A \wedge (x \notin C \wedge x \in C))\}$$

$$= \{x|((x \in A \wedge x \notin C) \wedge x \notin B) \vee (x \in A \wedge \emptyset)\}$$

$$= \{x|((x \in A \wedge x \notin C) \wedge x \notin B) \vee (\emptyset)\}$$

$$= \{x|((x \in A \wedge x \notin B) \wedge x \notin C)\}$$

Thus,

$$(A-B)-C = (A-C)-(B-C)$$

6) (10 points) Graded By : Rashmi Soni

If A,B, and C are sets, whether the symmetric difference is associative: $A \oplus (B \oplus C) = (A \oplus B) \oplus C$

Constructing the membership table:

A	B	C	$B \oplus C$	$A \oplus (B \oplus C)$	$A \oplus B$	$(A \oplus B) \oplus C$
1	1	1	0	1	0	1
1	1	0	1	0	0	0
1	0	1	1	0	1	0
1	0	0	0	1	1	1
0	1	1	0	0	1	0
0	1	0	1	1	1	1
0	0	1	1	1	0	1
0	0	0	0	0	0	0

From the membership table we see that the column values for $(A \oplus (B \oplus C))$ and $((A \oplus B) \oplus C)$ are equal, so $A \oplus (B \oplus C) = (A \oplus B) \oplus C$.