**CS 354 Spring 2018**

**Lab 1: Getting Acquainted with XINU (115 pts)**

**Due: 01/30/2018 (Tue.), 11:59 PM**

**1. Objectives**

The objectives of this introductory lab are to familiarize you with the steps involved in compiling and running XINU in our lab.

**2. Readings**

1. [XINU set-up description](http://www.cs.purdue.edu/homes/cs354/xinu/xinu-setup.html)
2. Chapters 1 and 2 from the XINU textbook.

**3. Inspecting XINU's run-time environment [115 pts]**

**Problem 3.1 (30 pts)**

Follow the instructions in the [XINU set-up description](http://www.cs.purdue.edu/homes/cs354/xinu/xinu-setup.html) which compiles the XINU source code on a frontend machine (Linux PC) in the XINU Lab, grabs an unused x86 Galileo backend machine, and loads and bootstraps the compiled XINU image. Note that the frontend PC's terminal acts as a (remote) console of the dedicated backend Galileo x86 machine. If XINU bootstraps successfully, it will print a greeting message and start a simple shell called xsh. Run some commands on the shell and follow the disconnect procedure so that the backend is released. **Do not hold onto a backend: it is a shared resource.**

(a) Inside the system/ subdirectory, you will find relevant source code of XINU. The file start.S contains assembly code that is executed after XINU loads on a backend. After performing bootstrap/initialization chores (the details are not relevant at this time), a call to function nulluser() is made. Locate where in system/ the source code of nulluser() is defined. (b) Toward the end of nulluser(), it calls the function create() -- a system call as discussed in class -- to create a new running program (i.e., process) whose code is the function main() contained in main.c. All processes in XINU are created using create() except the process that runs the function nulluser() which is "self-made." That is, it is the ancestor of all other processes. Find out where in the XINU source code this ancestor of all processes is located and what name, i.e., process ID (PID), it assigns itself. To do so, track down the header file (header files are located in the include/ directory) where the PID value of the ancestor process is specified. (c) What does the ancestor process running the function nulluser() do for the rest of its existence? (d) Does nulluser() ever return after being called by the code in start.S? (e) halt() is called after nulluser() in start.S. Where is the source code of halt() located and what does it do? (f) What happens if you remove this function call from start.S? Does XINU run as before? (g) What happens if you replace the while-loop at the end of nulluser() with a call to halt()? Discuss all your findings.

**Problem 3.2 (50 pts)**

In Linux/UNIX, to create a new process that runs an executable binary, say, /bin/ls (or a.out), we first use the system call fork() to create a child process and then call the system call execve() with /bin/ls (or a.out) as an argument. In XINU we said that a newly created process is put in a "state of limbo," i.e., suspended, after it is created using create(). That is, the child process exists as an entry in the process table but it is marked as suspended (PR\_SUSP). The definition of all XINU process states is specified in process.h under include/. Being in suspended state has the consequence that XINU will not assign CPU cycles to the process until its state is changed to ready (PR\_READY) by calling resume(). Look at the code of create() and resume() and determine how this is done.

(a) What happens in Linux after a new process is created using fork()? Note, create() specifies what code to run as the child process through its first argument (a function pointer). (b) Determine who runs first: parent or child. Try to empirically gauge the answer by running test code on the frontend Linux PCs. (We will discuss scheduling of processes in modern operating systems, a complex topic, when discussing process management.) (c) As an app programmer, do you have a preference as to which process -- parent or child -- should run next in Linux? Explain your reasoning. (d) For a child process created using fork() in Linux, it has to call execve() to make the child run executable code contained in a file (e.g., /bin/ls) as part of a file system specified in the argument of execve(). In Linux, sometimes an app programmer just wants to create a new process to run an existing binary without going through the 2-step procedure of fork() and execve(). Write a wrapper function, int newProcess(const char \*filename), that internally calls fork() and execve() to make this happen where filename specifies the full pathname of an executable binary. Place the code of newProcess() in newProcess.c under system/. Annotate your code with comments to facilitate readability.

(e) How is newProcess() fundamentally different from the way XINU's create() works? (f) Ignore stack size and process priority which are specified as arguments in create(). How does Linux's clone() compare to XINU's create()? (g) How does Linux's posix\_spawn() compared to your newProcess() implementation? (h) Is there a best way to create processes? Explain your reasoning.

**Problem 3.3 (5 pts)**

Customize the welcome message printed by XINU in main() so that it prominently displays your name (last, first) and username. Preserve the new welcome message as part of your version of XINU during the rest of the semester and the labs to follow.

**Problem 3.4 (30 pts)**

Problems 3.1 and 3.2 looked at the overall picture and steps involved in starting XINU and running apps under XINU. To get a sense of how to terminate processes and XINU itself, create a process that runs the code of a function onandon() by calling create() with stack size 2048 and priority 30. onandon() should be an infinite loop that makes a XINU system call putc(CONSOLE, 'x') followed by system call sleep() with argument 2 within its loop. onandon() prints the character 'x', sleeps for 2 seconds, repeating this forever. The process running the function main(), after creating a child process that runs function onandon(), sleeps for 14 seconds, then calls XINU's kill() with the PID of the child process as argument. To verify that the child process has been terminated, check the process table data structure proctab. As in Problem 3.1, inspect the code where nulluser() resides to determine how proctab is initialized and how to access it. Put this version of main() in main1.c under system/. The same goes for onandon() in onandon.c.

Provide written answers to the above problems in Lab1Answers.pdf and place it in system/.

**Bonus problem [15 pts]**

Extend the capability of XINU's xsh (source contained in shell/) by adding a new command of your own design. Please add your new command in a new file, lab1cmd.c in the shell/ directory. Compile and test XINU with the enhanced xsh. Indicate in Lab1Answers.pdf if you have done the Bonus Problem. Provide documentation of the added feature as comments inside the source code.

*Note: The bonus problem provides an opportunity to earn extra credits that count toward the lab component of the course. It is purely optional.*

**Turn-in instructions**

Before submitting your work, make sure to double-check the [TA Notes](http://www.cs.purdue.edu/homes/ogg/cs354_spring2018.html) to ensure that additional requirements and instructions have been followed.

Electronic turn-in instructions:

        i) Go to the xinu-spring2018/compile directory and run "make clean".

        ii) Go to the directory of which your xinu-spring2018 directory is a subdirectory. (Note: please do not rename xinu-spring2018 or any of its subdirectories.)

                e.g., if /homes/bob/xinu-spring2018 is your directory structure, go to /homes/bob

        iii) Type the following command

                turnin -c cs354 -p lab1 xinu-spring2018

*Please make sure to disable all debugging output before submitting your code.*