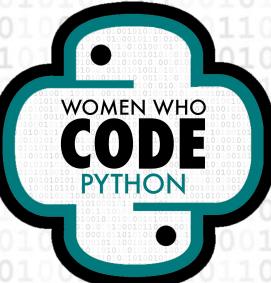


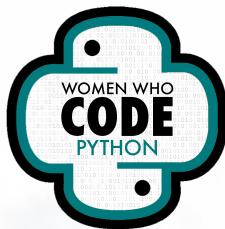
Welcome everyone!

- You can find these slides on GitHub here:
<https://github.com/WomenWhoCode/WWCodePython>
- Some housekeeping rules:
 - Everyone will be muted throughout the webinar, but there will be opportunities for participation!
 - Please share your thoughts on the chat and/or ask questions in the Q&A.
 - The entire team is here today. Please reach out to us with any technical questions!



WELCOME WOMEN WHO CODE





Women Who Code Python

**Intro to Data Structures
with Python:
Ace the Technical Interview**



Session #7: Heaps



MEET YOUR TEAM



Rishika Singh

Track Lead



Jasmeen Rajpal

Evangelist

OUR MISSION

Inspiring women to
excel in technology
careers.

WOMEN WHO
CODE



OUR VISION

A world where women are representative as technical executives, founders, VCs, board members and software engineers.

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OUR TARGET

Engineers with two or more years of experience looking for support and resources to strengthen their influence and levelup in their careers.



CODE OF CONDUCT

WWCode is an inclusive community, dedicated to providing an empowering experience for everyone who participates in or supports our community, regardless of gender, gender identity and expression, sexual orientation, ability, physical appearance, body size, race, ethnicity, age, religion, socioeconomic status, caste, creed, political affiliation, or preferred programming language(s).

Our events are intended to inspire women to excel in technology careers, and anyone who is there for this purpose is welcome. We do not tolerate harassment of members in any form. Our **Code of Conduct** applies to all WWCode events and online communities.

Read the full version and access our incident report form at womenwhocode.com/codeofconduct

230,000 Members

70 networks in 20 countries
Members in 97+ countries
10K+ events
\$1025 daily Conference tickets
\$2M Scholarships
Access to [jobs](#) + [resources](#)
Infinite connections



OUR MOVEMENT

As the world changes, we can be a connecting force that creates a sense of belonging while the world is being asked to isolate.



Upcoming Events

FRI
11
JUN

CONNECT REIMAGINE Virtual Conference Featured

► Online | 8:00 AM – 1:00 PM PDT (UTC-0700)

[Register](#)

THU
24
JUN

✨ Intro to Data Structures with Python: Ace the Technical Interview (Session #8: Graphs) ✨ Featured

► Online | Python | 5:00 PM – 6:30 PM PDT (UTC-0700)

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TUE
29
JUN

✨ Ask Me Anything with Jane Street Software Engineers ✨ Featured

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THU
01
JUL

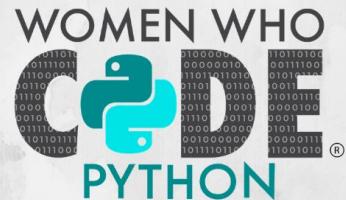
✨ Intro to Data Structures with Python: Ace the Technical Interview (Session #9: Maps & Hash Tables) ✨ Featured

► Online | Python | 5:00 PM – 6:30 PM PDT (UTC-0700)

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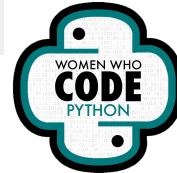
@WWCODEPYTHON

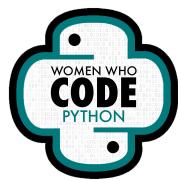
WOMENWHOCODE.COM/PYTHON



Nayeon Shin

Sophomore CS & New Media Art Major | WWCode Python Track General Volunteer





Today's Agenda



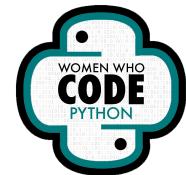
1. Queues & Trees Review
2. Priority Queue
3. Heap
 - a. Binary Heap
 - i. Structure Property
 - ii. Order Property
 - iii. Heap Operations
 - iv. Implementation
 4. Q&A
 5. Live Coding
 6. Resources

Queues & Trees Review



Queues

- On the contrary to stacks, a queue is open at both ends.
 - One end (Rear): Insertion (Enqueue)
 - The other (Front): Deletion (Dequeue)
- First-In-First-Out (FIFO)



Trees

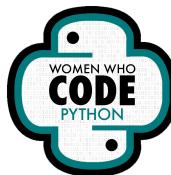
- Two definitions:

1. Nodes and edges

A tree consists of a set of nodes and a set of edges that connect pairs of nodes.

2. Recursive

A tree is either empty or consists of a root and zero or more subtrees, each of which is also a tree.

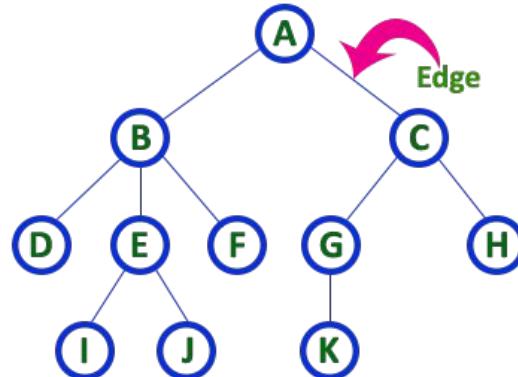


Trees

- Two definitions:

1. Nodes and edges

A tree consists of a set of nodes and a set of edges that connect pairs of nodes.



Trees

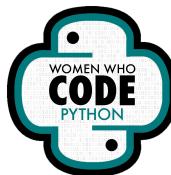
- Two definitions:

1. Nodes and edges

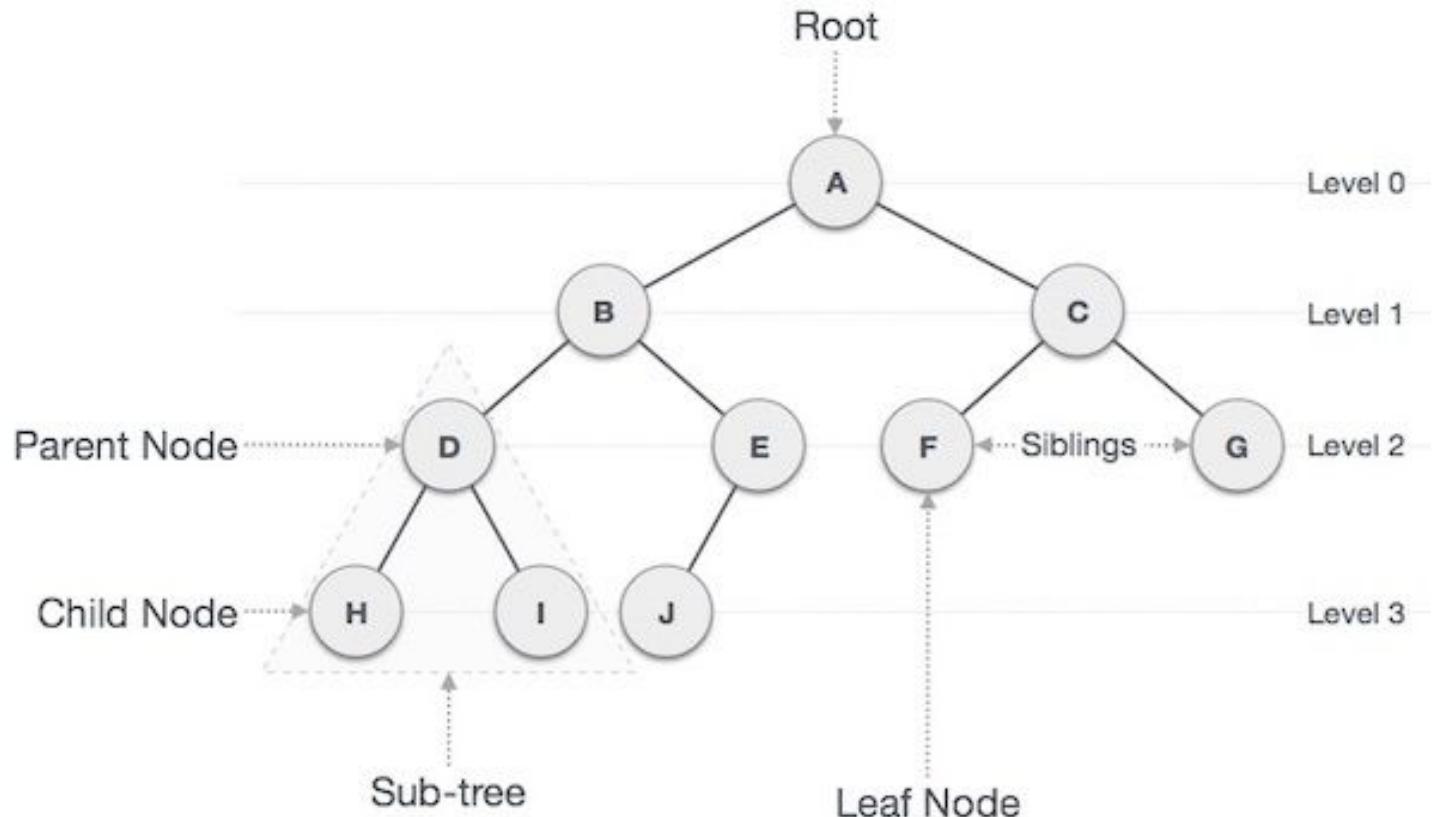
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2. Recursive

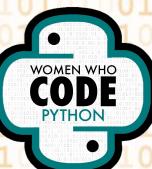
A tree is either empty or consists of a root and zero or more subtrees, each of which is also a tree.



Trees

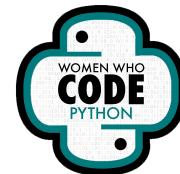
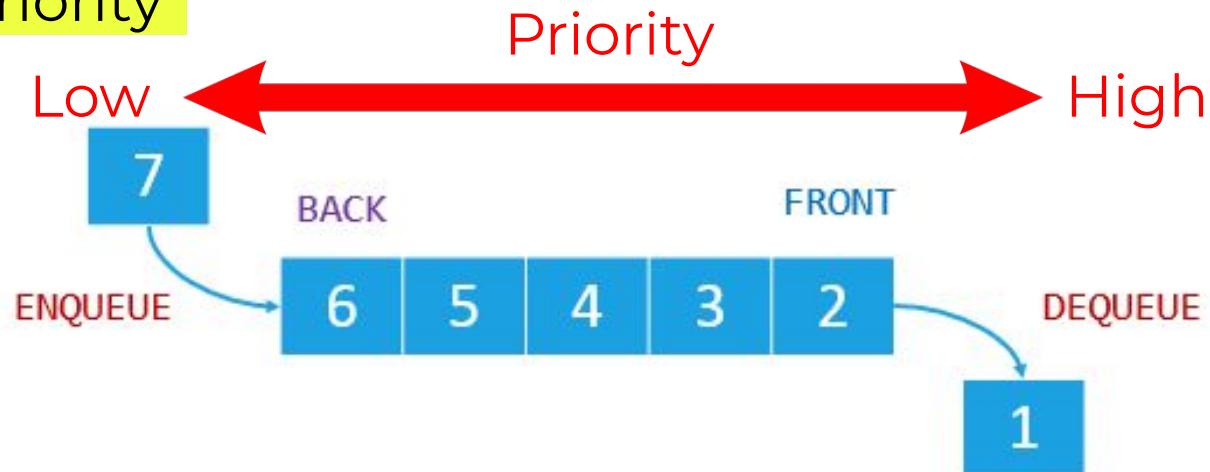


Priority Queue



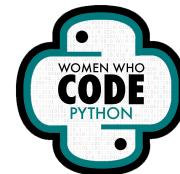
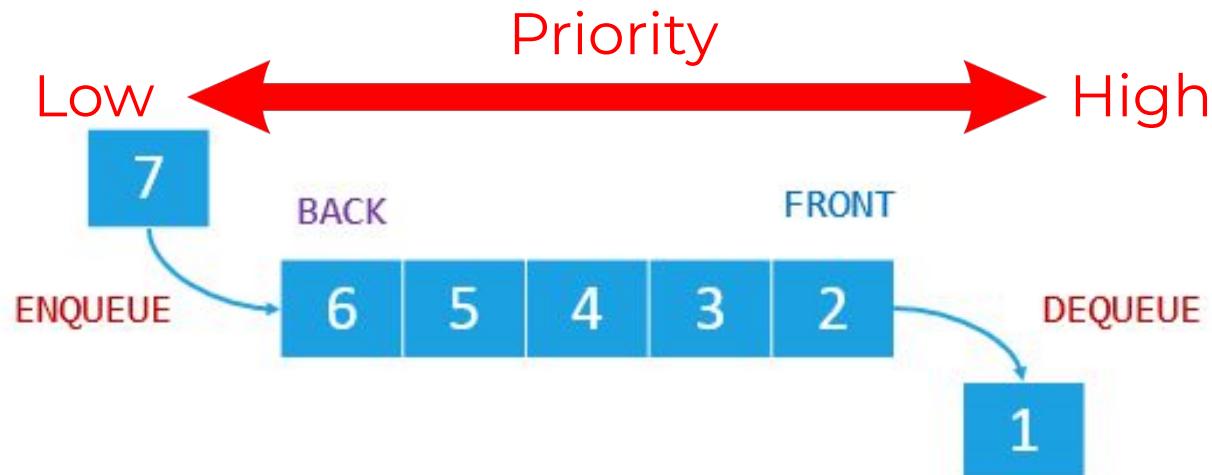
Priority Queue

- Variation of a queue
- “Priority” queue:
 - Order of queue items is determined by their “priority”



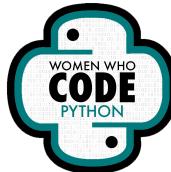
Priority Queue

- Each entry: Priority (key) + Data value
- Data remove → Element with the highest priority
 - Priority: Any kind of comparable info



Priority Queue

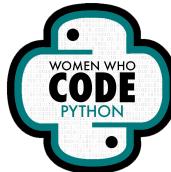
- Main methods:
 - `insert(new_item)` → Insert a new item to the queue
 - `remove_max()` → Remove & return entry with largest key
 - Return None (Null) if empty queue
 - `get_max()` → Return entry with largest key value (leave it on queue)
 - `get_size()` → Return # of queue elements
 - `is_empty()`
- Lowest value as priority → `remove_min()` & `min()`



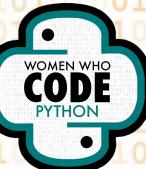
Implementation of a Priority Queue?

- Priority queue: Queue that has a priority order
 - Queue + Ordering?
 1. Implement PQ with unsorted lists
 - Expensive `max()` & `remove_max()`
 2. Implement PQ with sorted lists
 - Expensive `insert(k, v)`

Is there a better way? YES!

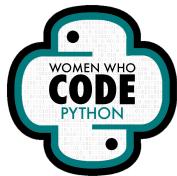


Heap



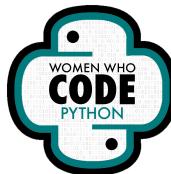
Binary Heap

- Two main properties
 1. Complete binary tree (Structure property)
 2. Ordering property



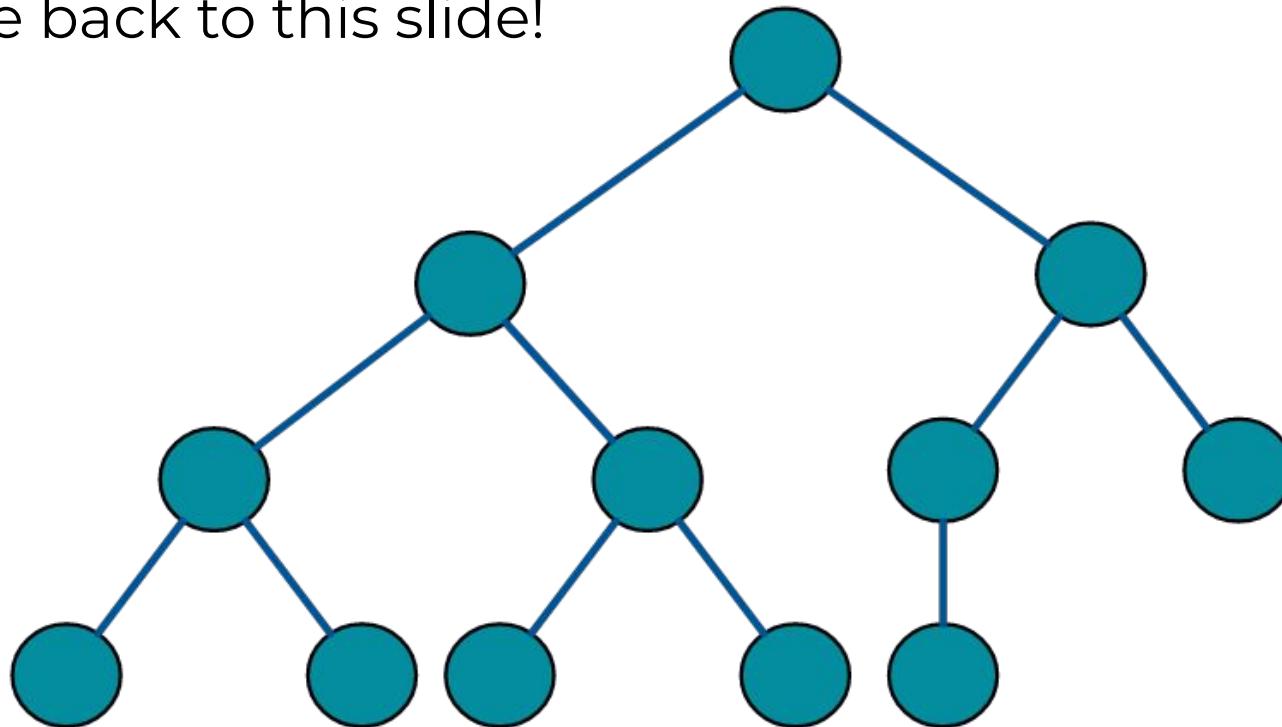
Binary Heap - Structure Property

- Two main properties
 1. Complete binary tree (Structure property)
 - For efficiency, take advantage of binary tree's logarithmic performance
 - For logarithmic performance, keep our tree balanced
 - Balanced binary tree:
Roughly the same # of nodes in the left & right



Binary Heap - Structure Property

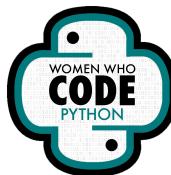
Will come back to this slide!



Binary Heap - Structure Property

- Two main properties
 1. Complete binary tree
 - Binary tree
 - Completely filled on each level except possibly lowest level
 - Lowest level: Filled from left to right
 - Any empty spots: At rightmost side of lowest level

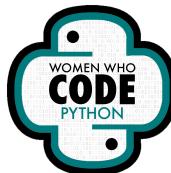
Go back to previous slide!



Binary Heap - Order Property

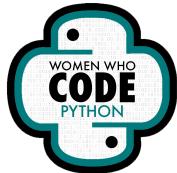
- Two main properties
 - 2. Order property
 - Min-heap or max-heap
 - Consider max-heap:

For every position p other than the root,
 $p \leq \text{key of } p\text{'s parent}$
- 1) No relative ordering between siblings
 - 2) Largest element: At the root
 - 3) Along any path from root to leaf, non-increasing order

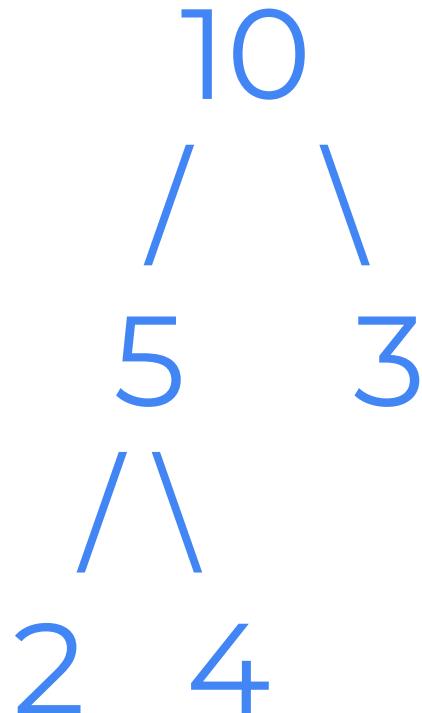


Binary Heap - Heap Operations

- Adding a new node in a max heap
 1. Add a new node at leaf
 2. while ($\text{node} > \text{parent}$) and (node is not root):
Swap node with parent

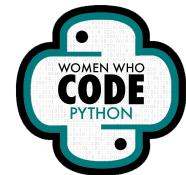


Binary Heap - Heap Operations: Add

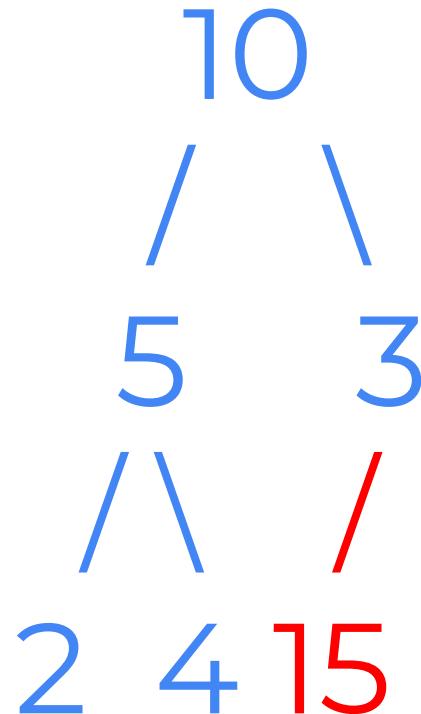


New element to be inserted: 15

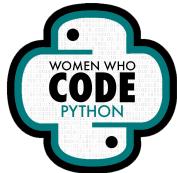
Step 1: Insert the new element at the end



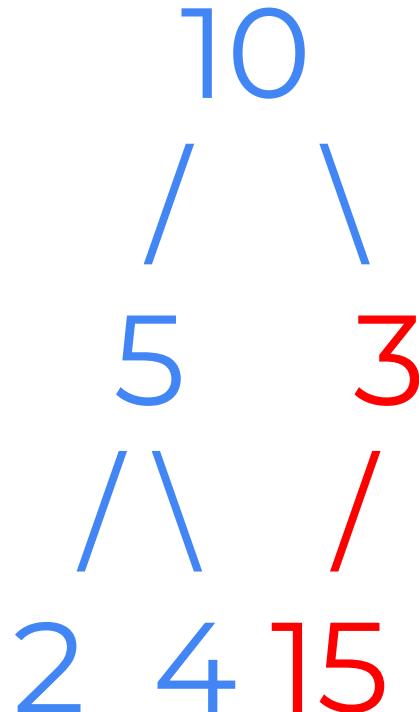
Binary Heap - Heap Operations: Add



Step 1: Insert the new element at the end

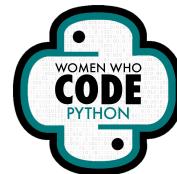


Binary Heap - Heap Operations: Add

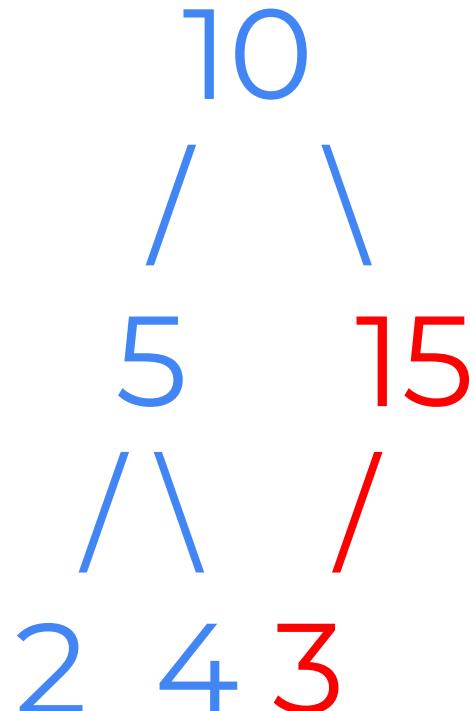


Step 2: “Heapify” the new element from the bottom.

- Heapify:
- while (`node > parent`) and
(`node` is not root):
 Swap `node` with `parent`



Binary Heap - Heap Operations: Add

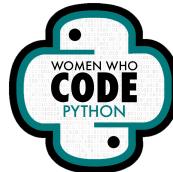


Heapify:

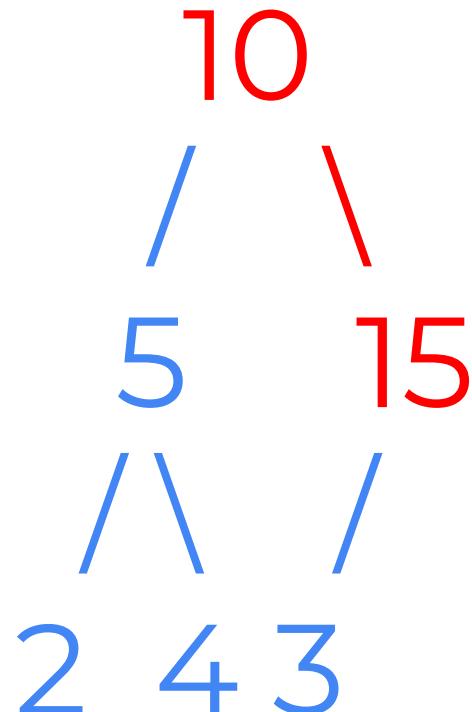
- while (`node > parent`) and
(`node` is not root):

Swap node with parent

$15 > 3 \rightarrow \text{Swap those two!}$



Binary Heap - Heap Operations: Add

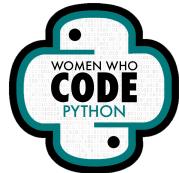


Heapify:

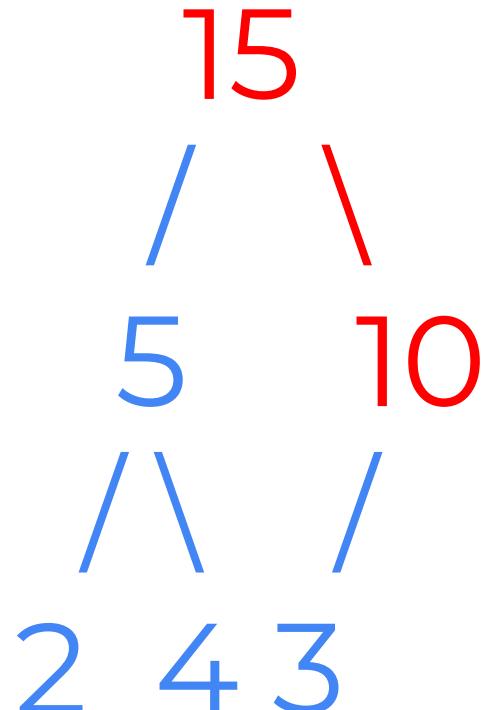
- while (`node > parent`) and
(`node` is not root):

Swap node with parent

$15 > 10 \rightarrow \text{Swap again!}$



Binary Heap - Heap Operations: Add

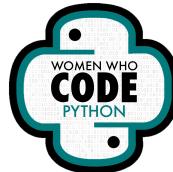


Heapify:

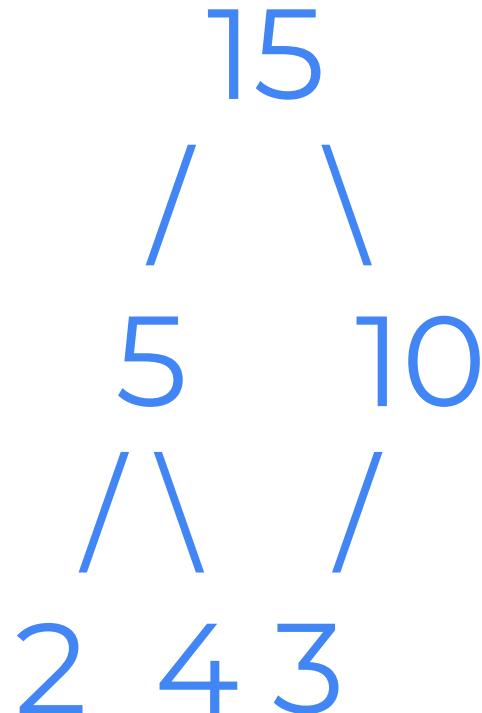
- while (`node > parent`) and
(`node` is not root):

Swap node with parent

15 is root; done.



Binary Heap - Heap Operations: Add



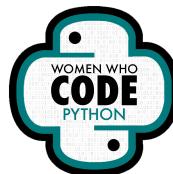
← FINAL HEAP

Heapify:

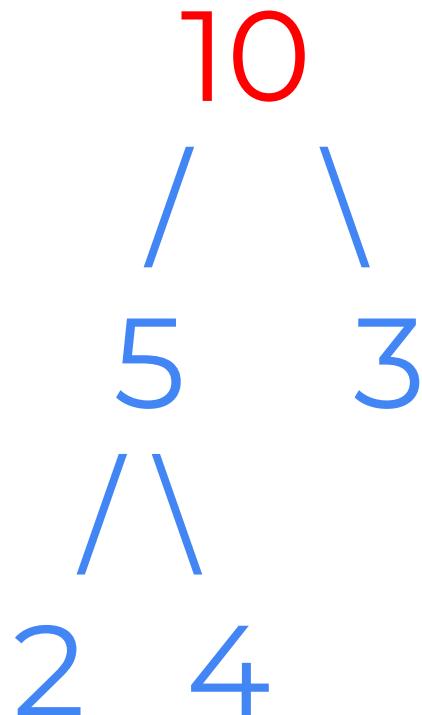
- while ($\text{node} > \text{parent}$) and
(node is not root):

Swap node with parent

Notice that the root is the largest.

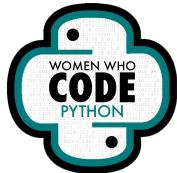


Binary Heap - Heap Operations: Remove

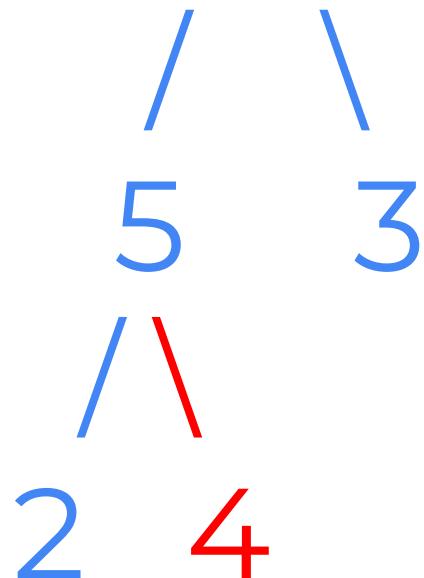


Element to be deleted: 10

Step 1: Remove the root

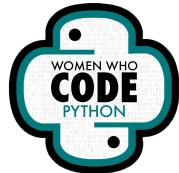


Binary Heap - Heap Operations: Remove

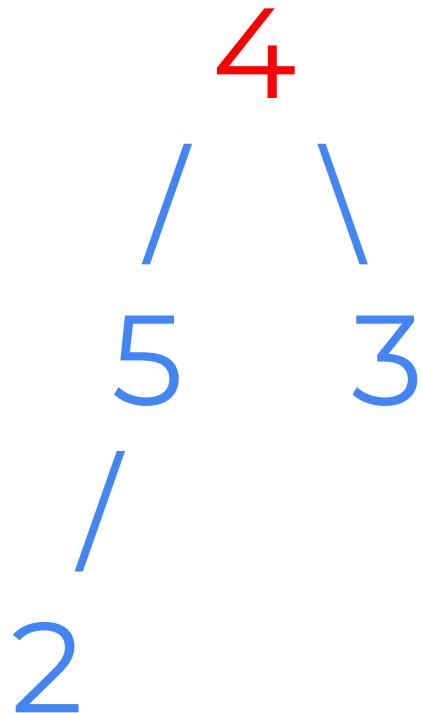


Element to be deleted: 10

Step 2: Make rightmost leaf the root

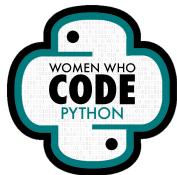


Binary Heap - Heap Operations: Remove

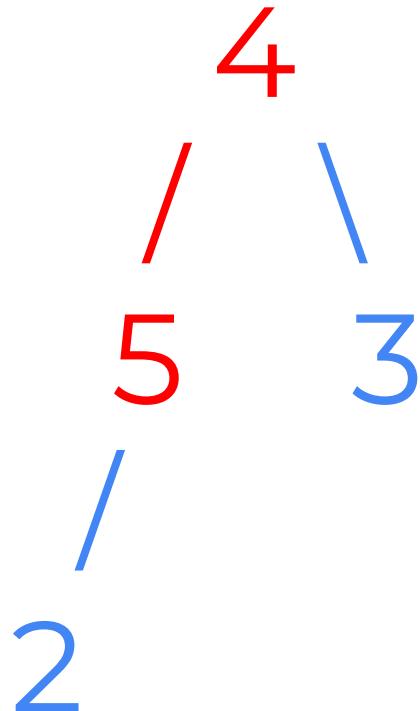


Element to be deleted: 10

Step 2: Make rightmost leaf the root



Binary Heap - Heap Operations: Remove



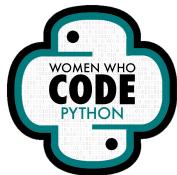
Element to be deleted: 10

Heapify:

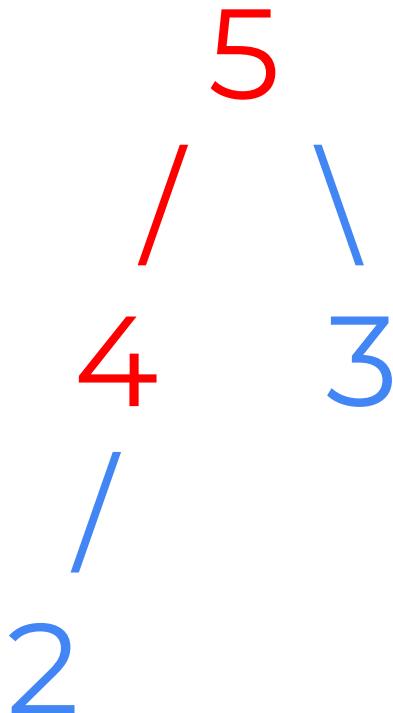
- while ($\text{node} > \text{parent}$) and
(node is not root):

Swap node with parent

Step 3: Heapify the root



Binary Heap - Heap Operations: Remove



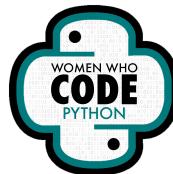
Element to be deleted: 10

Heapify:

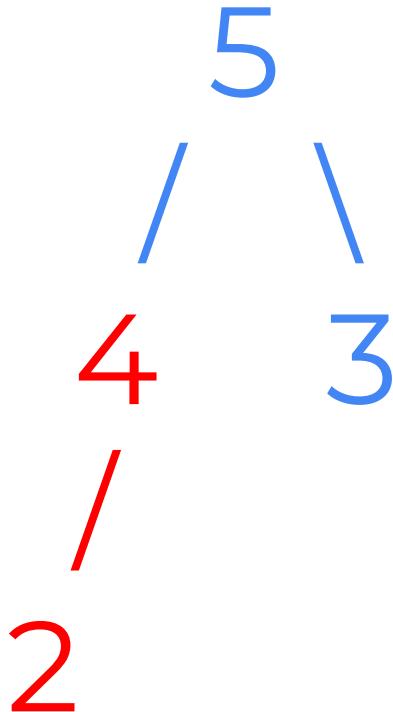
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Swap node with parent

Step 3: Heapify the root



Binary Heap - Heap Operations: Remove



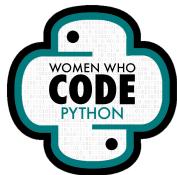
Element to be deleted: 10

Heapify:

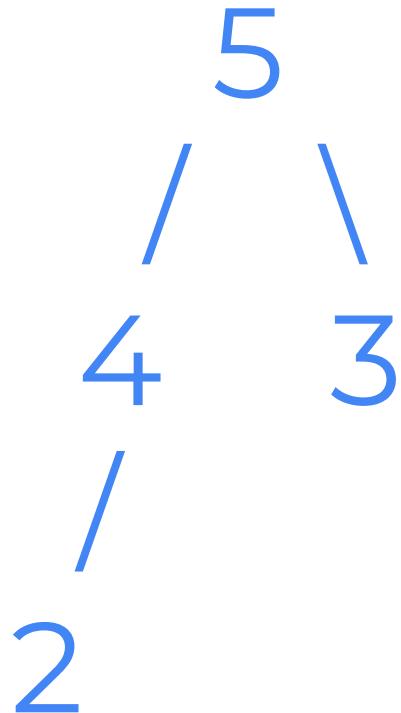
- while ($\text{node} > \text{parent}$) and
(node is not root):

Swap node with parent

Step 3: Heapify the root

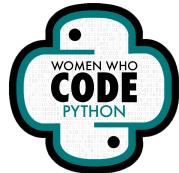


Binary Heap - Heap Operations: Remove



Element that has been deleted: 10

Done.

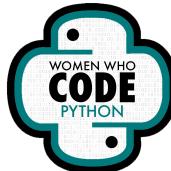


Binary Heap - Implementation

- It is yet another ADT (Abstract Data Type)!
- How should we implement it?
 1. Linked structure is possible

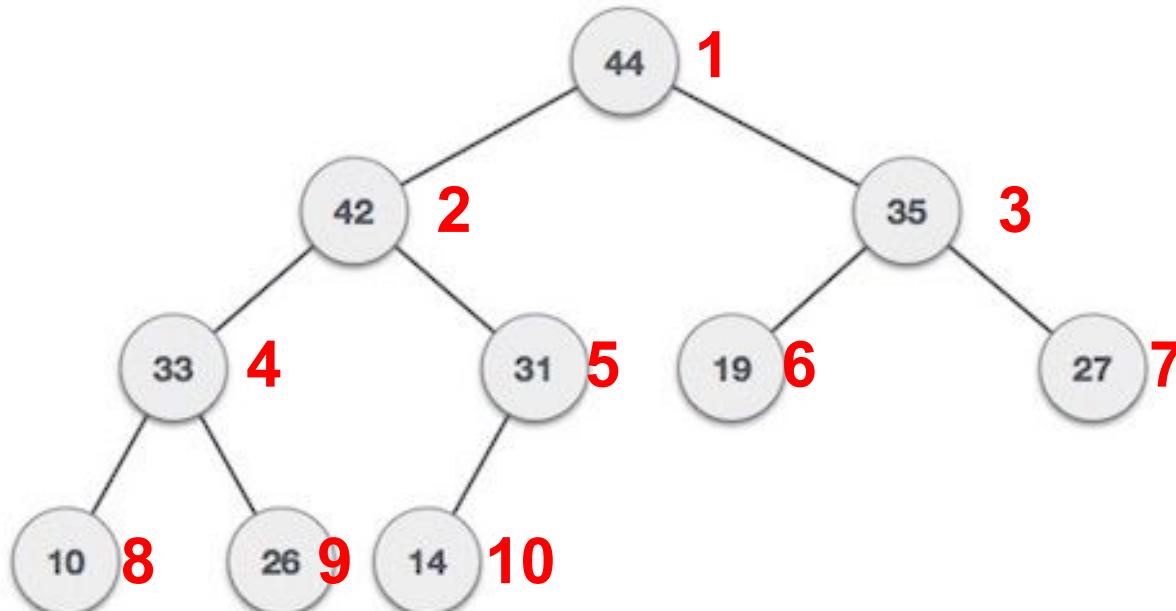
```
class HeapNode:  
    self.value ....  
    self.left ....  
    self.right ....
```

2. Another way?

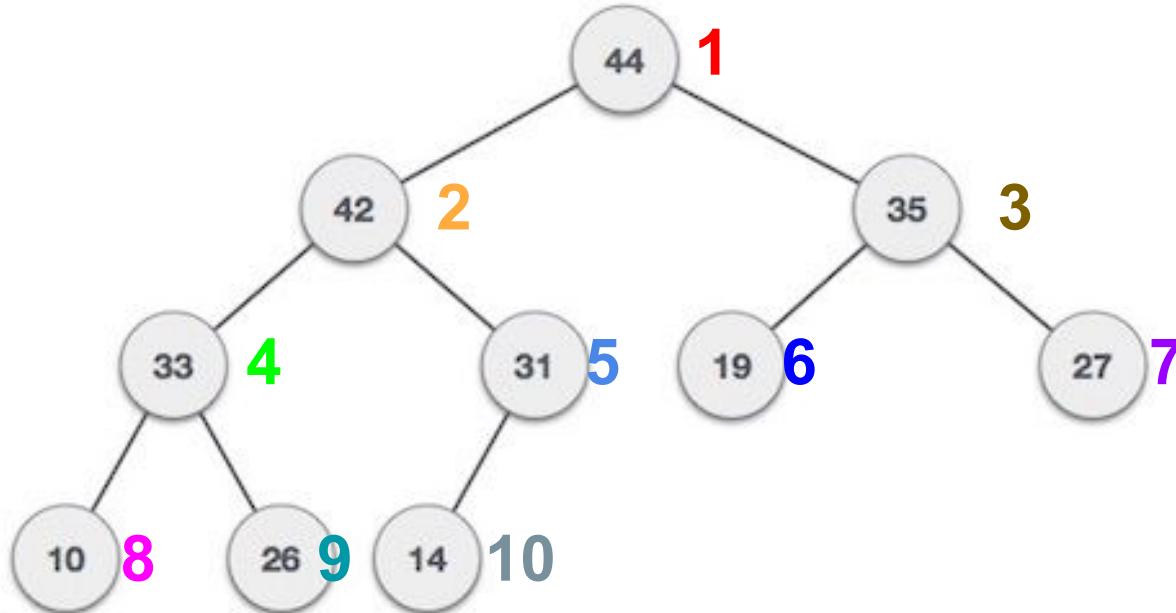


Binary Heap - Implementation

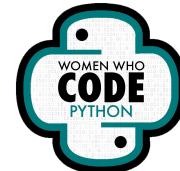
- Use **lists!**
 - Order nodes top to bottom, left to right



Binary Heap - Implementation

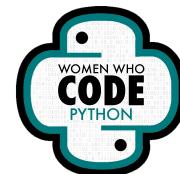
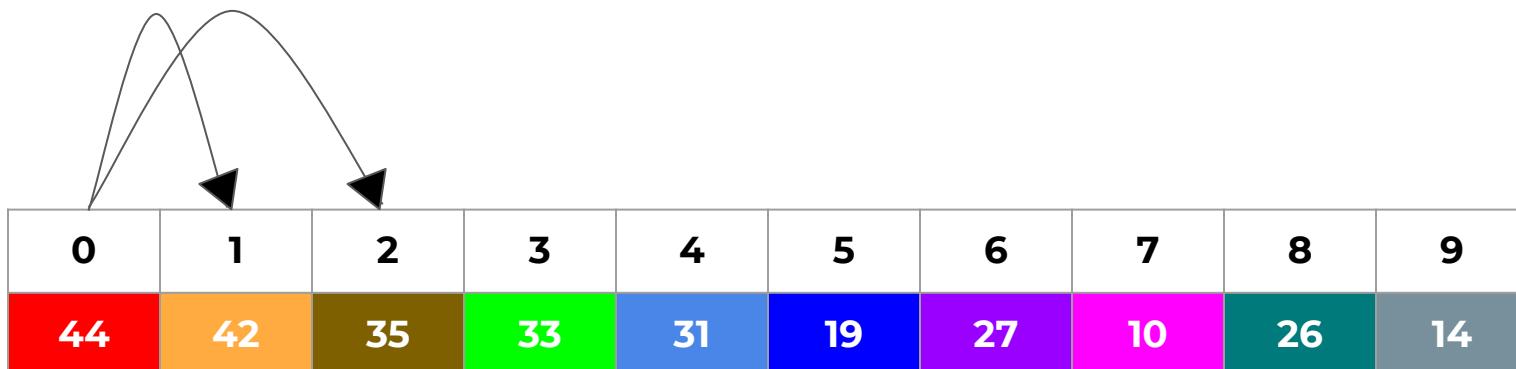


0	1	2	3	4	5	6	7	8	9
44	42	35	33	31	19	27	10	26	14



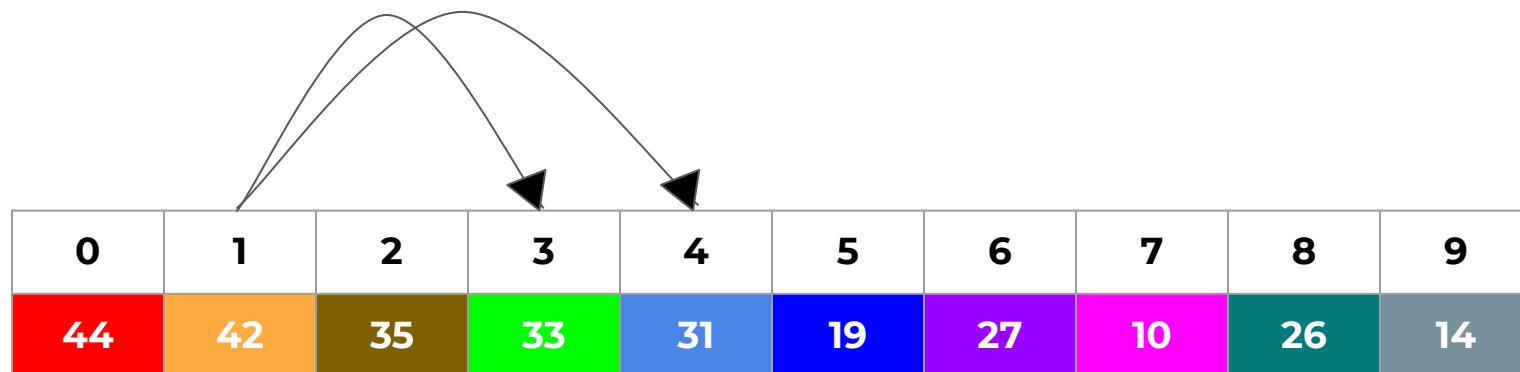
Binary Heap - Implementation

- For any node **i**
 - Left child is at $(2 * i) + 1$
 - Right child is at $(2 * i) + 2$
 - Parent is at **math.ceil(i / 2) - 1; (i - 1) // 2**



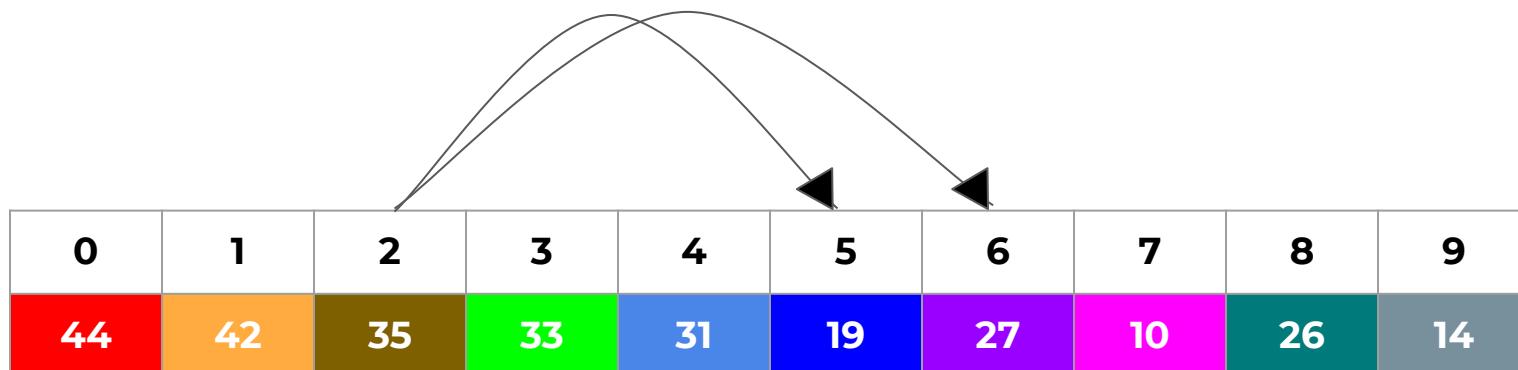
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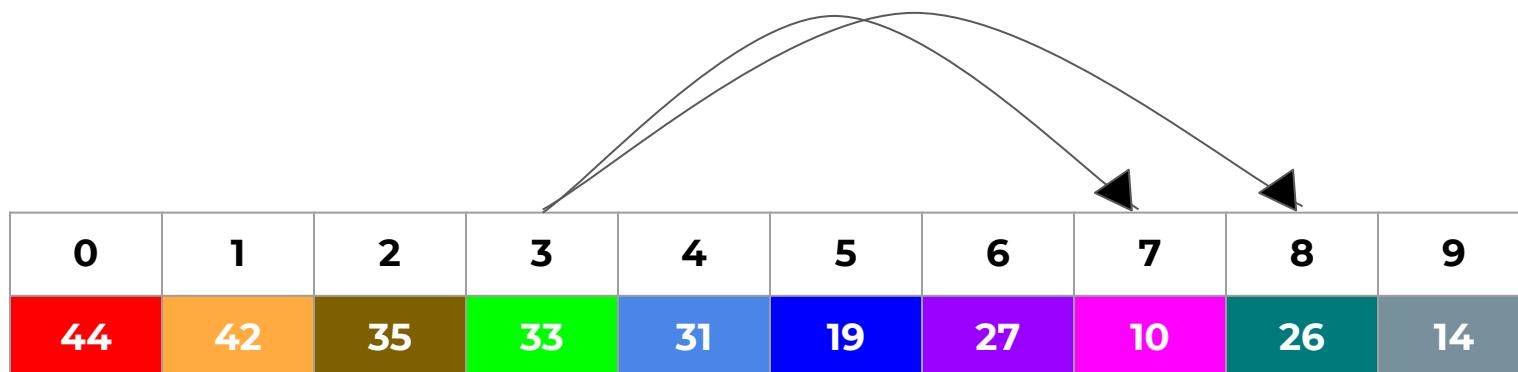
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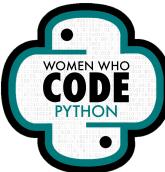
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- For any node **i**
 - Left child is at $(2 * i) + 1$
 - Right child is at $(2 * i) + 2$
 - Parent is at **math.ceil(i / 2) - 1; (i - 1) // 2**

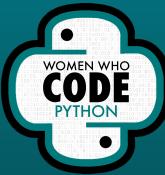


Binary Heap

- Other heaps the other way
 - Max-heap (max at the root):
For every position p other than the root,
 $p \leq$ key of p 's parent
 - Min-heap (min at the root):
For every position p other than the root,
 $p \geq$ key of p 's parent

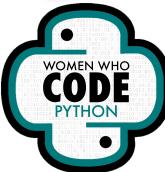


Q&A Time!



Time for Live Coding!

- <https://drive.google.com/file/d/1DMjHnILhb4b-KoqiGx7OaudAwRdLX5WY/view?usp=sharing>

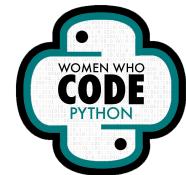
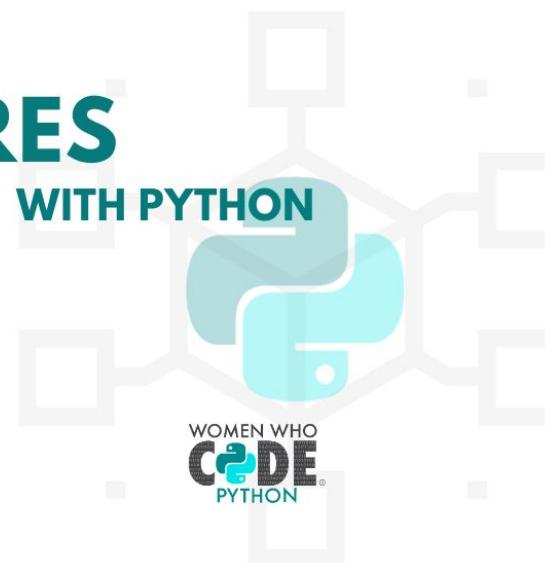


Next Session!

INTRO TO
**DATA
STRUCTURES**

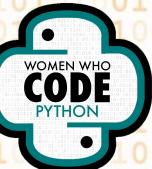
ACE THE
TECHNICAL
INTERVIEW

THU. JUNE 17TH
@ 8:00PM EDT



Questions?

Join our Slack channel:
#intro-data-structures-stdy-grp



Thank You!

