

가 장 가 까 운 1 0 0 쌍 찾 기

# 적벽대전

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CSE244 Final Project

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2017136008 권창덕

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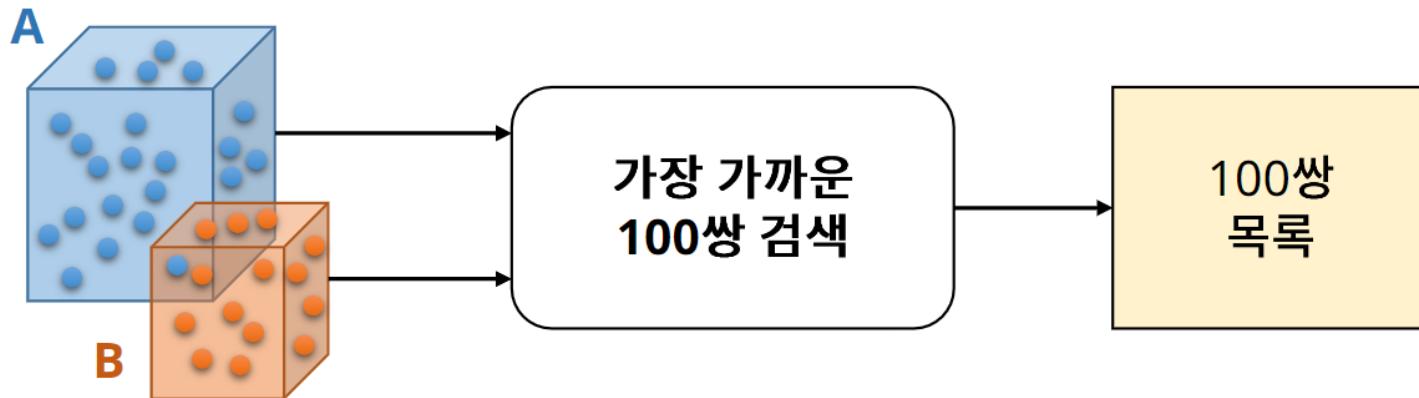
- 결과분석

01

## 문제 정의

## 01

## 문제 소개



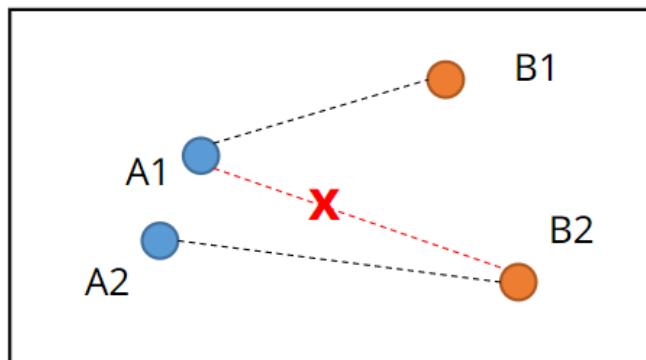
**입력** : A부대, B부대 (부대별 병사 수, x,y,z 좌표)

**출력** : 피아식별하며, 가장 가까운100쌍 목록

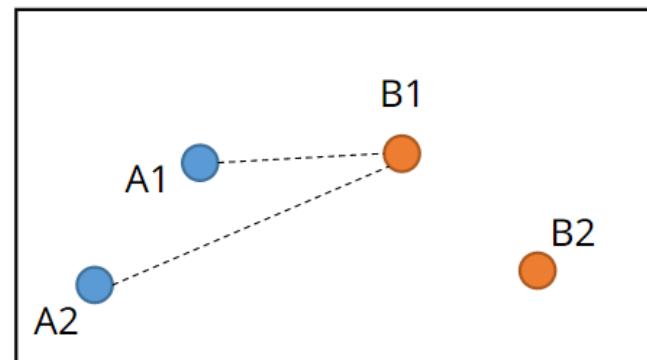
[IMG] [https://koreatechackr-my.sharepoint.com/personal/bluekds\\_koreatech\\_ac\\_kr/\\_layouts/15/onedrive.aspx?id=%2Fpersonal%2Fbluekds%5fkoreatech%5Fac%5Fkr%2FDocuments%2F%EA%B0%95%EC%9D%98%2F%EB%A9%80%ED%8B%B0%EC%BD%94%EC%96%B4%ED%94%84%EB%A1%9C%EA%B7%B8%EB%9E%98%EB%B0%8D%2F2022%5Fspring%2F%5BMP%5D%20Project02%5FCUDA%2Epdf&parent=%2Fpersonal%2Fbluekds%5fkoreatech%5Fac%5Fkr%2FDocuments%2F%EA%B0%95%EC%9D%98%2F%EB%A9%80%ED%8B%B0%EC%BD%94%EC%96%B4%ED%94%84%EB%A1%9C%EA%B7%B8%EB%9E%98%EB%B0%8D%2F2022%5Fspring&ga=1](https://koreatechackr-my.sharepoint.com/personal/bluekds_koreatech_ac_kr/_layouts/15/onedrive.aspx?id=%2Fpersonal%2Fbluekds%5fkoreatech%5Fac%5Fkr%2FDocuments%2F%EA%B0%95%EC%9D%98%2F%EB%A9%80%ED%8B%B0%EC%BD%94%EC%96%B4%ED%94%84%EB%A1%9C%EA%B7%B8%EB%9E%98%EB%B0%8D%2F2022%5Fspring%2F%5BMP%5D%20Project02%5FCUDA%2Epdf&parent=%2Fpersonal%2Fbluekds%5fkoreatech%5Fac%5Fkr%2FDocuments%2F%EA%B0%95%EC%9D%98%2F%EB%A9%80%ED%8B%B0%EC%BD%94%EC%96%B4%ED%94%84%EB%A1%9C%EA%B7%B8%EB%9E%98%EB%B0%8D%2F2022%5Fspring&ga=1)

## 01

## 문제 소개



( X )



( O )

단, A부대 기준으로 검색  
(Euclidean distance 이용)

[IMG] [https://koreatechackr-my.sharepoint.com/personal/bluekds\\_koreatech\\_ac\\_kr/\\_layouts/15/onedrive.aspx?id=%2Fpersonal%2Fbluekds%5Fkoreatech%5Fac%5Fkr%2FDocuments%2F%EA%B0%95%EC%9D%98%2F%EB%A9%80%ED%8B%B0%EC%BD%94%EC%96%B4%ED%94%84%EB%A1%9C%EA%B7%B8%EB%9E%98%EB%B0%8D%2F2022%5Fspring%2F%5BMP%5D%20Project02%5FCUDA%2Epdf&parent=%2Fpersonal%2Fbluekds%5Fkoreatech%5Fac%5Fkr%2FDocuments%2F%EA%B0%95%EC%9D%98%2F%EB%A9%80%ED%8B%B0%EC%BD%94%EC%96%B4%ED%94%84%EB%A1%9C%EA%B7%B8%EB%9E%98%EB%B0%8D%2F2022%5Fspring&ga=1](https://koreatechackr-my.sharepoint.com/personal/bluekds_koreatech_ac_kr/_layouts/15/onedrive.aspx?id=%2Fpersonal%2Fbluekds%5Fkoreatech%5Fac%5Fkr%2FDocuments%2F%EA%B0%95%EC%9D%98%2F%EB%A9%80%ED%8B%B0%EC%BD%94%EC%96%B4%ED%94%84%EB%A1%9C%EA%B7%B8%EB%9E%98%EB%B0%8D%2F2022%5Fspring%2F%5BMP%5D%20Project02%5FCUDA%2Epdf&parent=%2Fpersonal%2Fbluekds%5Fkoreatech%5Fac%5Fkr%2FDocuments%2F%EA%B0%95%EC%9D%98%2F%EB%A9%80%ED%8B%B0%EC%BD%94%EC%96%B4%ED%94%84%EB%A1%9C%EA%B7%B8%EB%9E%98%EB%B0%8D%2F2022%5Fspring&ga=1)

02

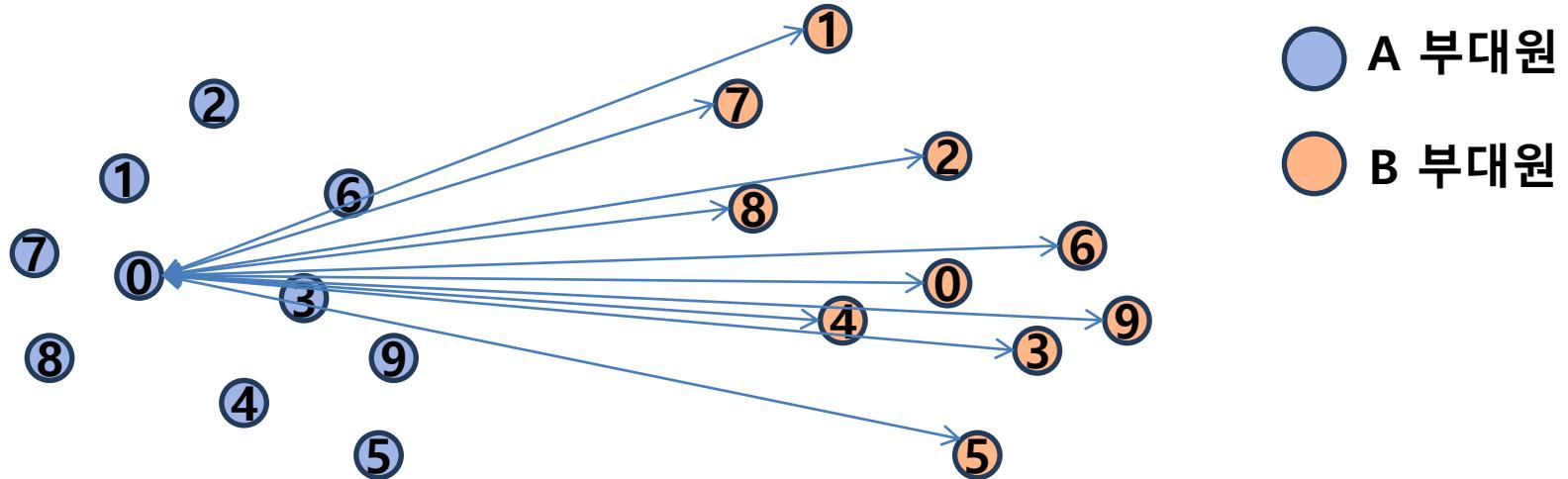
## 설계 및 구현

Parallel ver.1

**Brute-force**

02

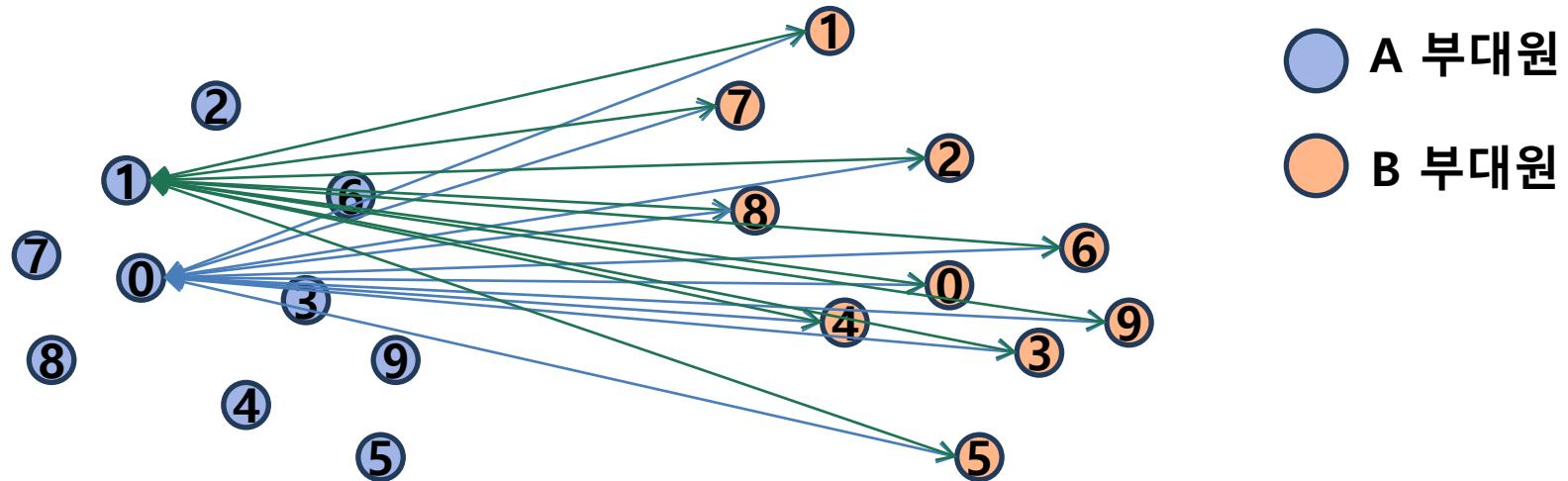
## ver1. Idea



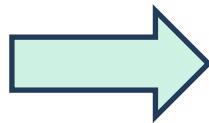
각 A부대원별, 모든 B부대원과의 거리를 계산하여, 최소거리인 B부대원을 찾음  
한 A부대원당 총 B부대원 수만큼 매칭

02

## ver1. Idea



A 수 = M  
B 수 = N

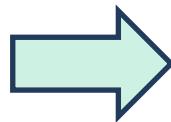


$O(M*N)$   
searching  
\*single thread

## 02

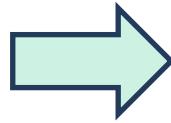
ver1. 구현

A부대원 수만큼  
GPU Thread 생성

**1D Grid & 1D Block**

Block layout : (1024, 1, 1)  
Grid layout : (ceil( $M / 1024$ ), 1, 1)

A 부대 정보  
&  
B 부대 정보

**Global Memory**

각 부대원 정보를  
[XXX..., YYY..., ZZZ...]  
형태로 저장

## 02

## ver1. 구현

Thread 0 Thread 1

...

Thread 30 Thread 31

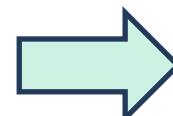
A부대원 0 A부대원 1



A부대원 30 A부대원 31



각 thread별 global index를 이용,  
한 warp이 A부대 정보를  
coalesced Mem. access



각 thread별 자신이 계산할  
A부대원 정보 획득

~~Global Memory~~

A 부대 정보

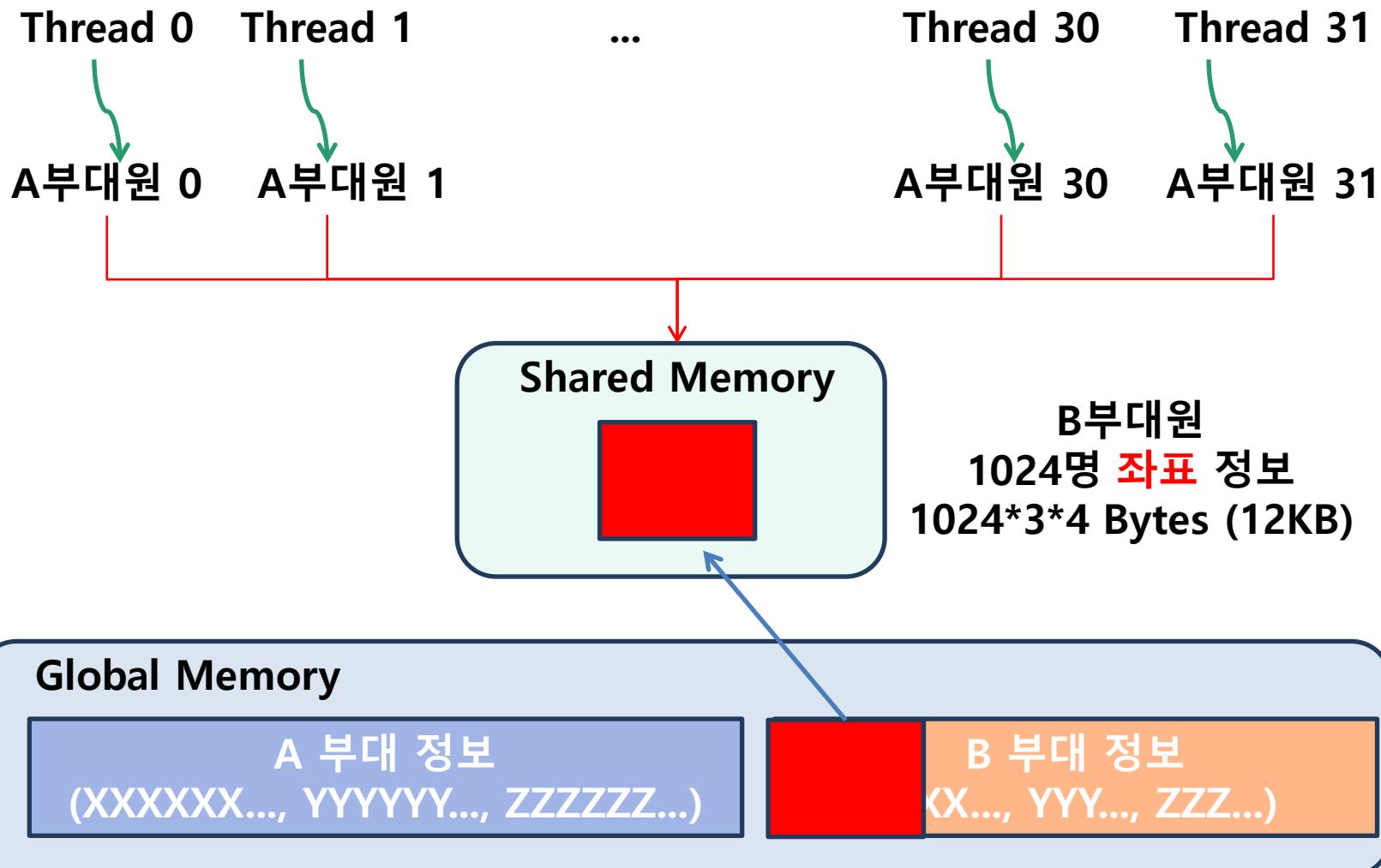
(XXXXXX..., YYYYYY..., ZZZZZZ...)

B 부대 정보

(XXX..., YYY..., ZZZ...)

02

## ver1. 구현

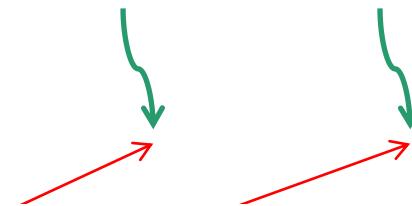
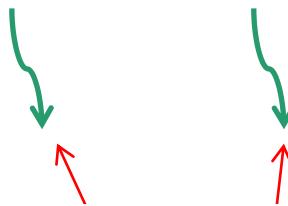


## 02

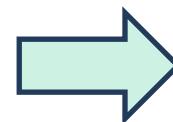
## ver1. 구현

**Thread 0   Thread 1**

...

**Thread 30****Thread 31**

블록 내 모든 Thread가  
동일 bank로 접근



**BroadCast**  
(no bank conflict)

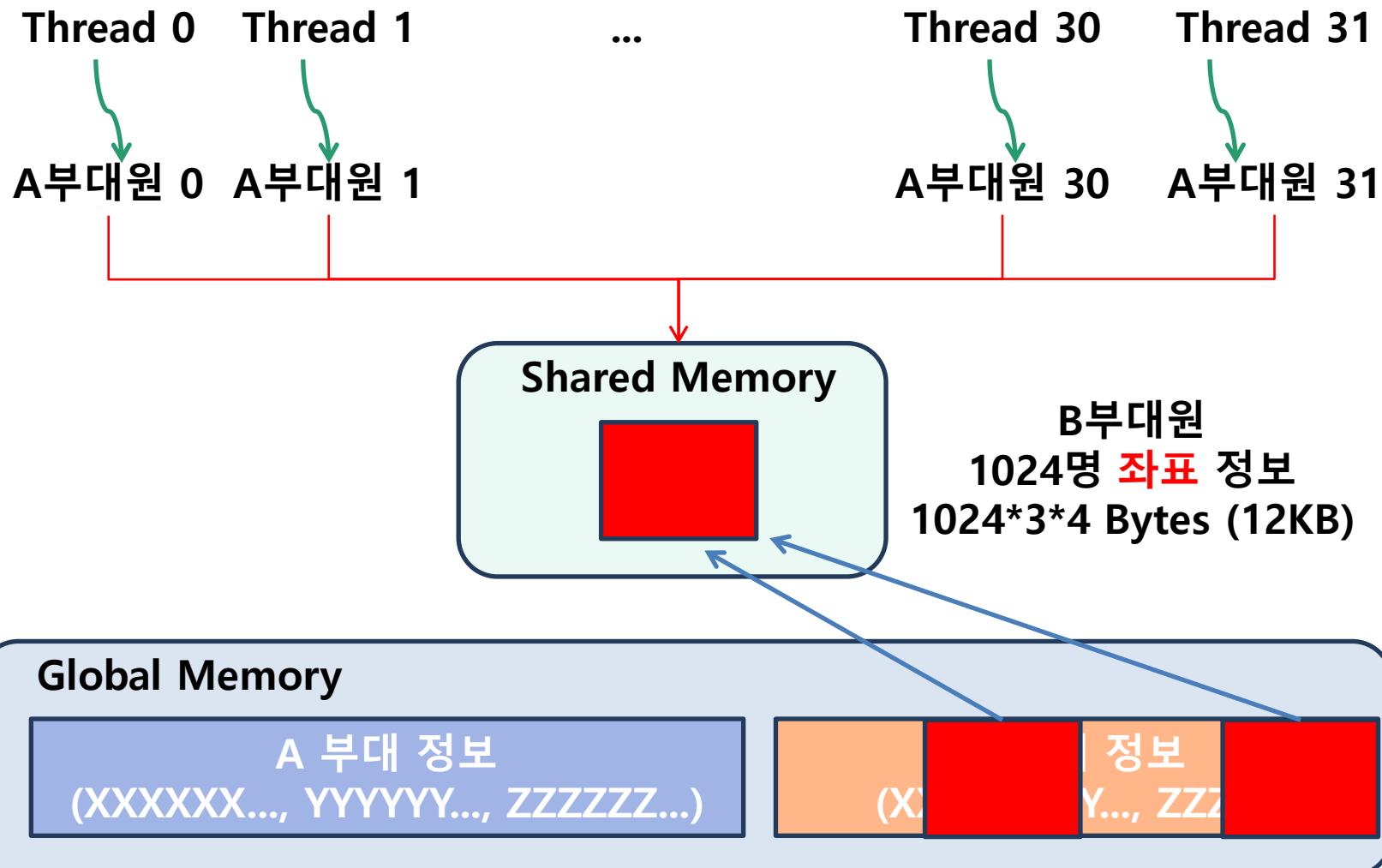
**Bank**

|   |   |   |   |     |     |     |    |    |    |
|---|---|---|---|-----|-----|-----|----|----|----|
| 0 | 1 | 2 | 3 | ... | ... | ... | 29 | 30 | 31 |
| x | y | z | x | ... |     |     | z  | x  | y  |
| z | x | y | z |     |     |     | y  | z  | x  |
| y | z | x | y |     |     |     | x  | y  | z  |

**Data****Shared Memory**

02

## ver1. 구현



Parallel ver.2

# Spatial Partitioning

02

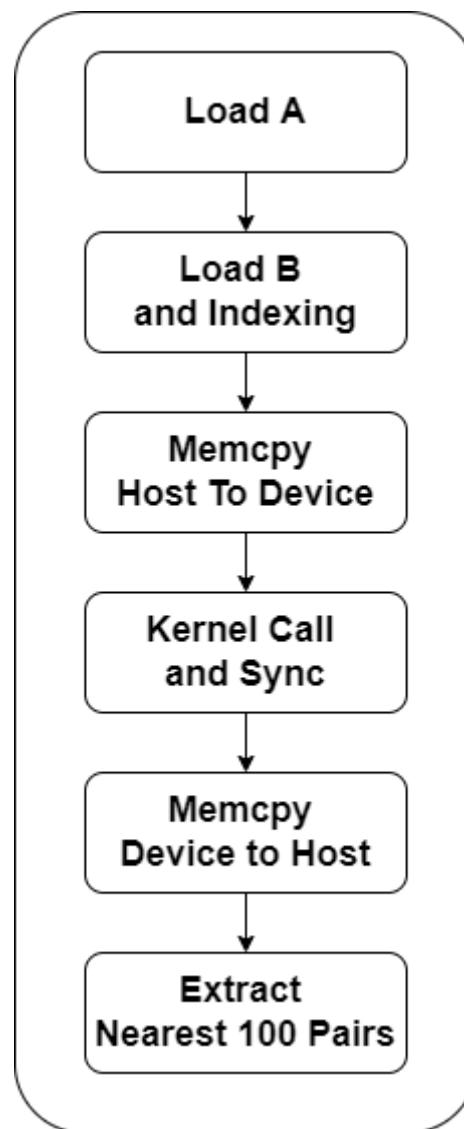
## ver2. Motivation



[IMG] <http://kkoma.net/product/%ED%98%84%EC%88%98%EB%A7%89-%EB%B0%B0%EA%B2%BD3-%EB%B0%A4%ED%95%98%EB%8A%98-%EB%B3%84-%EC%97%B0%EA%B7%B9-%EA%B3%B5%EC%97%B0-%EB%AC%B4%EB%8C%80-%ED%8F%AC%ED%86%A0%EC%A1%B4-%EB%B0%B0%EA%B2%BD/6305/>

02

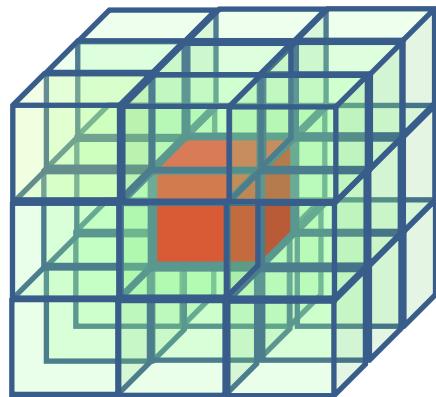
## 전체 흐름



\*\* kernel 외 전부 serial

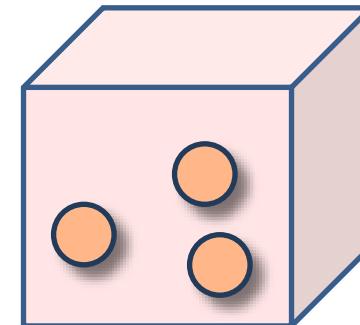
02

## ver2. Idea



큐브 구조체 배열

< 전체 공간 >



큐브 구조체

- B의 갯수

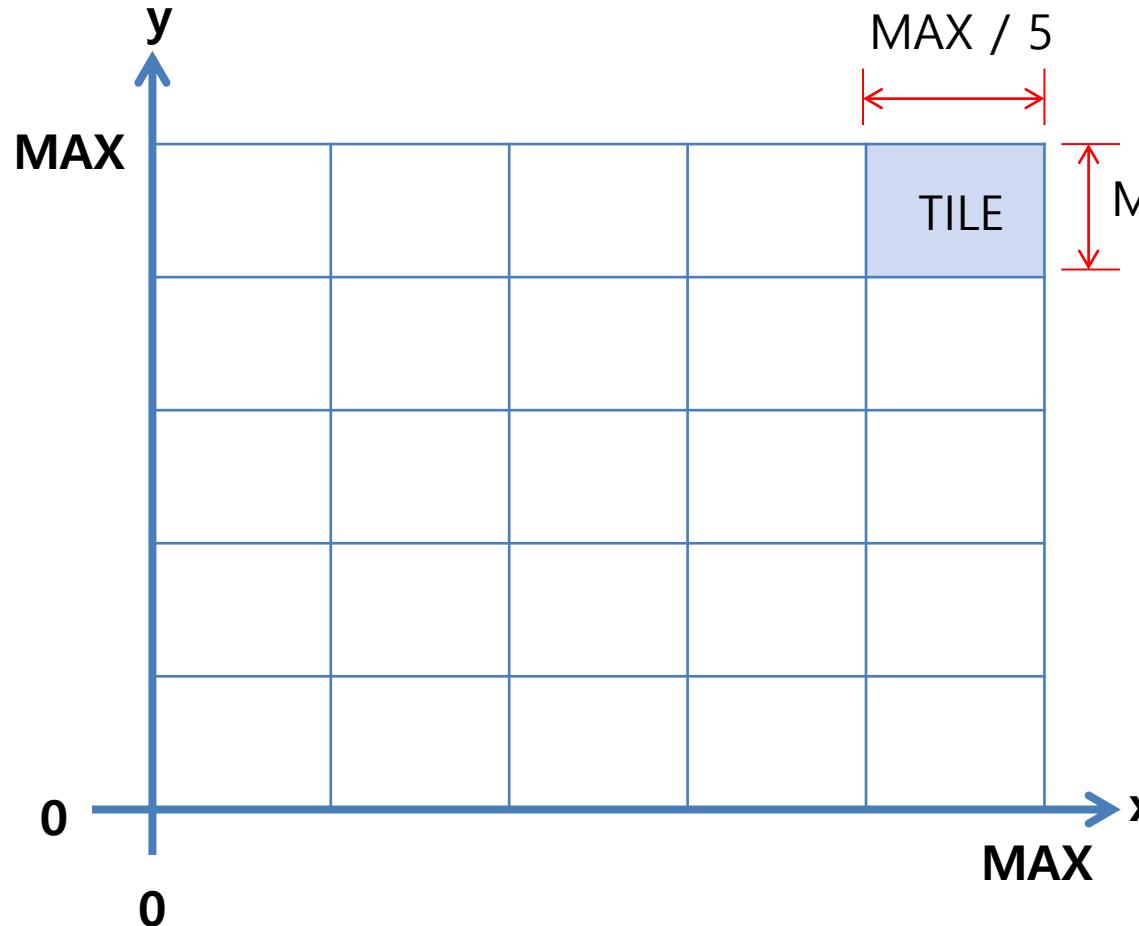
- B 부대원(ID, 좌표) 배열

< 분할된 공간 >

각 B의  $(x,y,z)$  좌표를 이용해  
큐브 구조체 배열(큐브 배열)에서 저장 위치(큐브)를 결정

02

## ver2. Idea



A 부대원

B 부대원

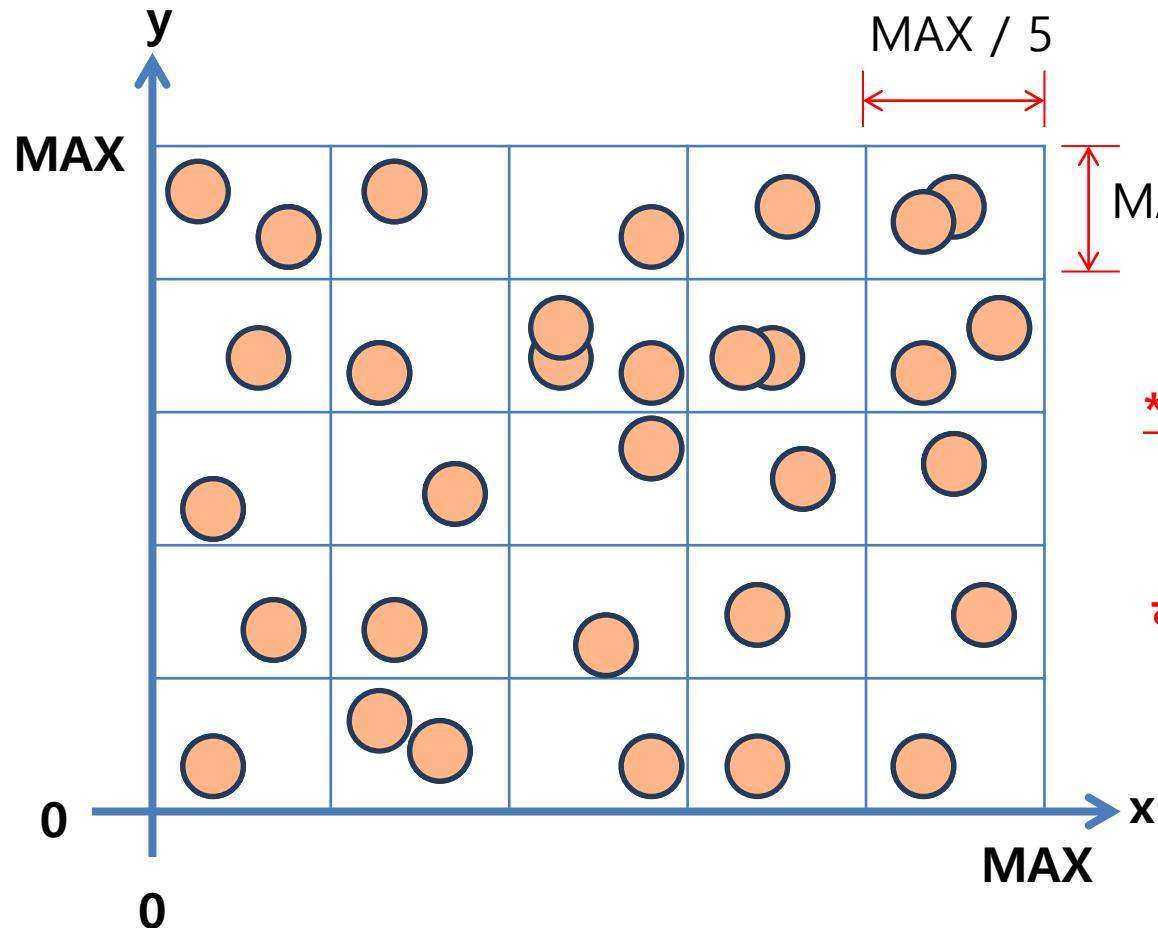
만약, B의 수 = 25,

$$\sqrt{25} = 5,$$

각 축을 5등분  
→ 전체 평면이  
25등분됨

02

## ver2. Idea



● A 부대원

● B 부대원

$MAX / 5$

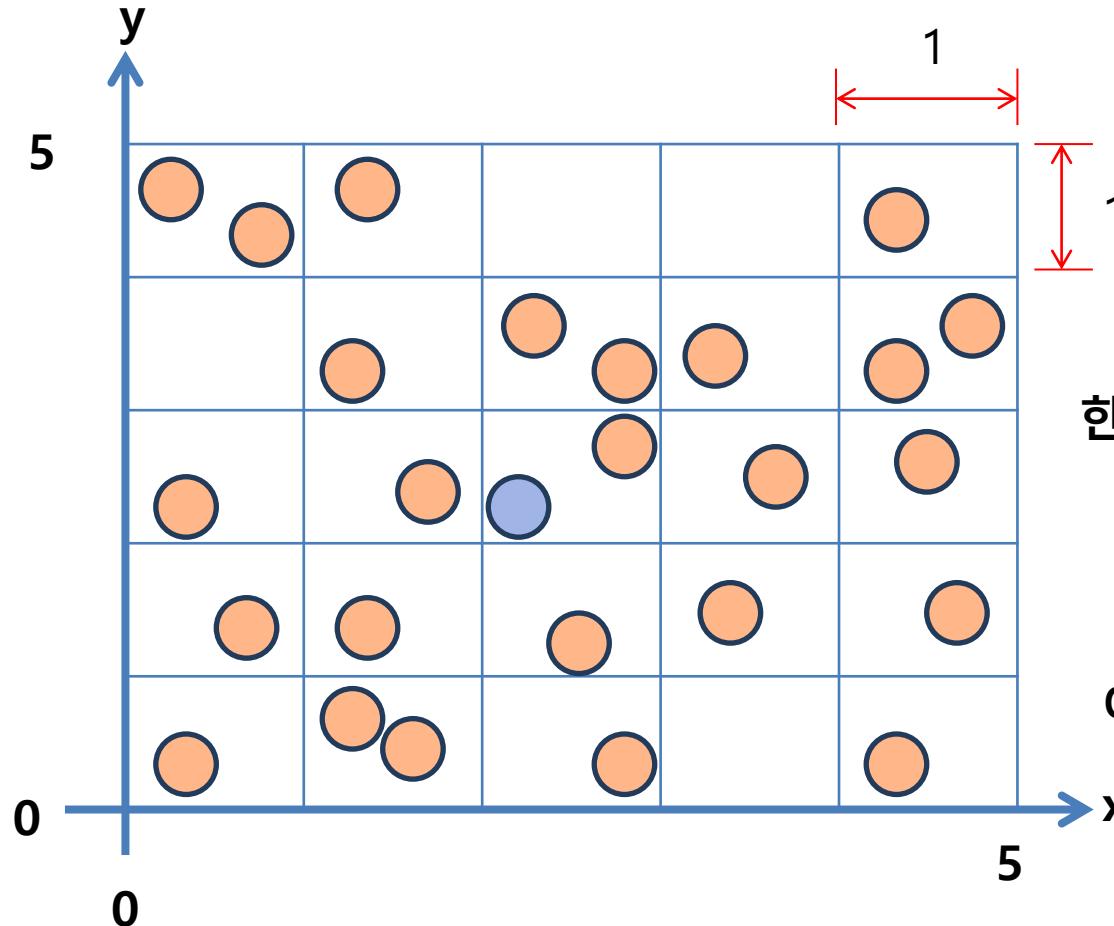
$MAX / 5$

Army Generator →  
\* uniform int distribution  
\* RANGE MAX

이상적인 상황에서,  
한 Tile 당 존재할 B의 수  
의 기대값은 1개

02

## ver2. Idea



A 부대원

B 부대원

만약, MAX=5이고,

한 A의 좌표가 (2.1, 2.1)이면,

$$X' = 2.1 / 1$$

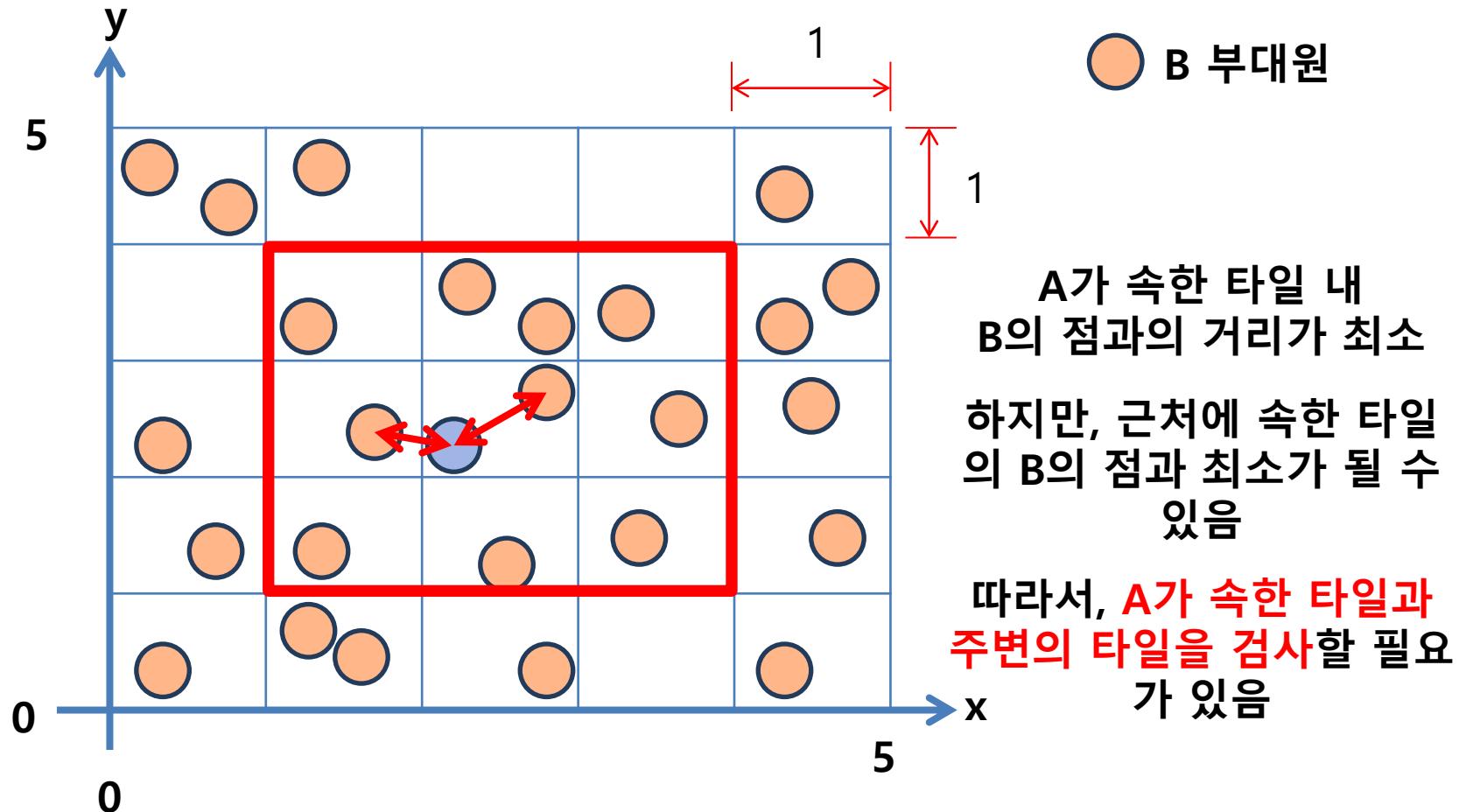
$$Y' = 2.1 / 1$$

\* 정수 나눗셈

이 A가 속한 타일의 인덱스  
 $(X', Y') = (2, 2)$

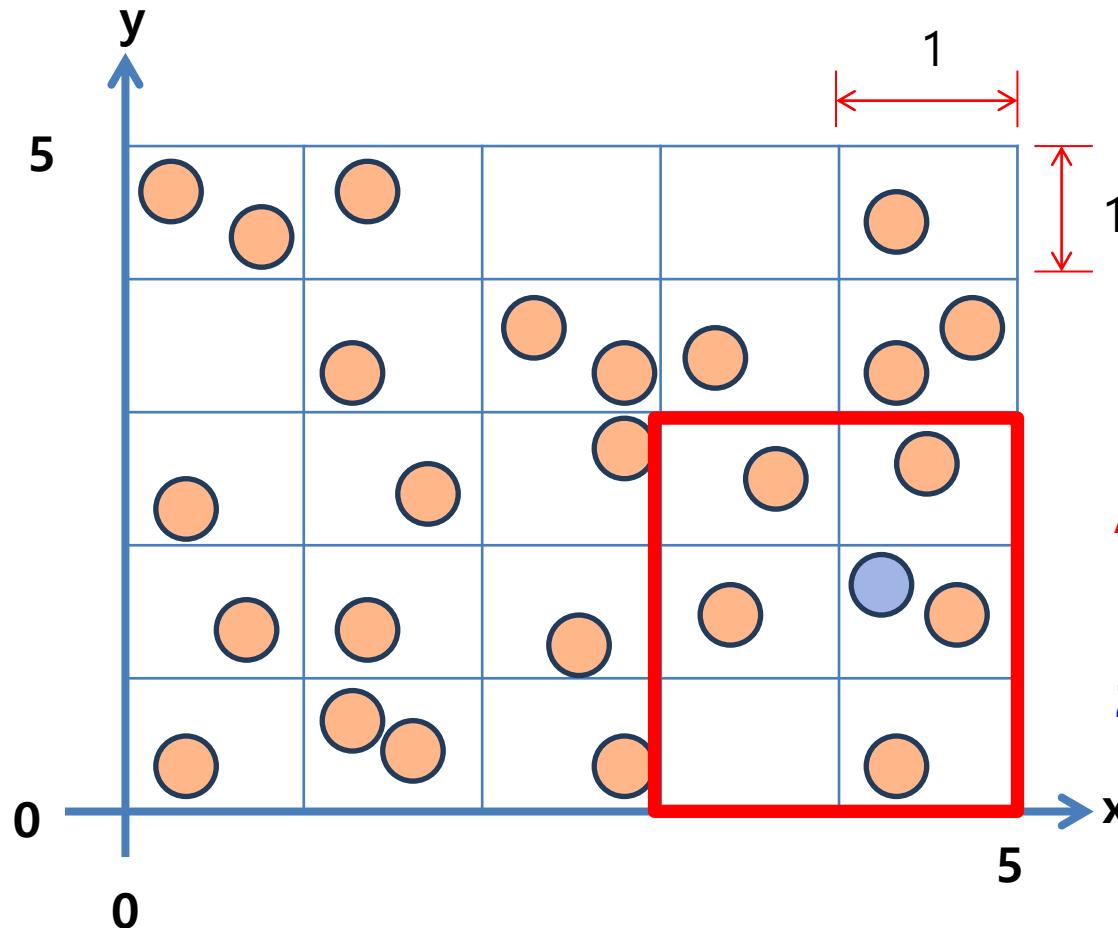
02

## ver2. Idea



02

## ver2. Idea



다른 A부대원 역시  
A가 속한 타일과 주변의  
타일을 검사

25개의 B 중 5개만 검색

## 02

## ver2. Idea

B의 수 =  $N$

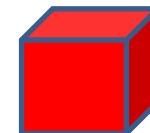
각 축을  $\sqrt[3]{N}$  등분

→ 전체 공간  $N$ 등분  
( $=N$ 개의 큐브)

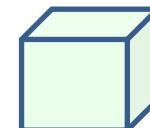
큐브 모서리 길이  
 $= \text{RANGE\_MAX} / \sqrt[3]{N}$

부대원 좌표  $(x,y,z)$  각각  
큐브 모서리 길이로 정수 나누셈

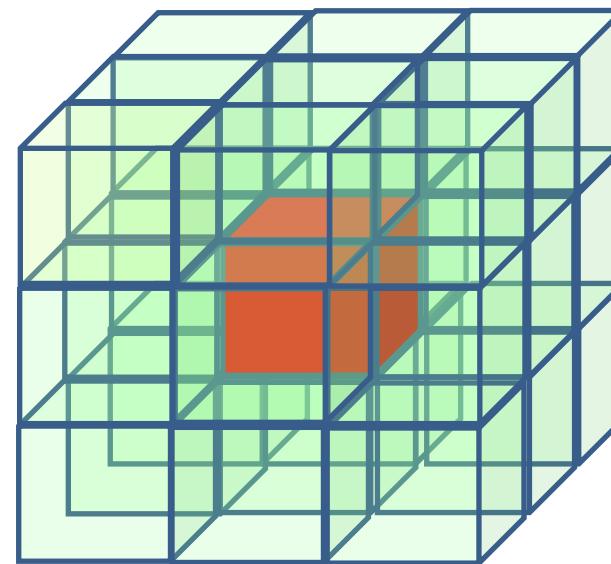
몫 → 해당 부대원이 존재하는  
큐브 위치를 알 수 있음



A가 속한 큐브



주변 큐브



인접한 큐브를 포함  
총 27개( $3 \times 3 \times 3$ )의 큐브  
조사

02

## ver2. Idea

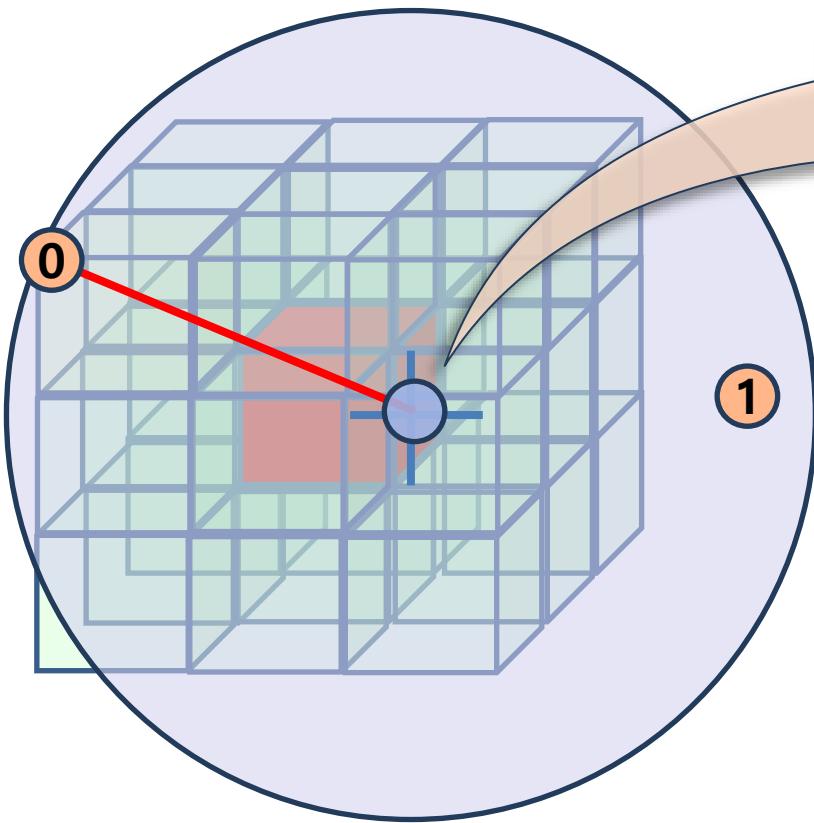
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총 27개( $3 \times 3 \times 3$ ) 큐브 내  
B가 하나라도 존재하지 않으면,

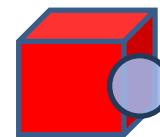
가장 가까운 100쌍 내 속하지 않을 것으로 추정

02

## ver2. Idea



검색 영역 내 B가 하나 존재하고,  
가장 멀리 떨어진 경우



- A 부대원
- B 부대원

0번은 검색 영역에 속하지만,  
1번은 검색 영역에 속하지 않음  
하지만, 1번 점과의 거리가 더 짧음

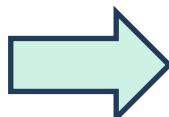
100 << A의 수

\* A: 100만, B: 40만 기준  
큐브 모서리 길이 = 약 14,170  
GT 100쌍 중 가장 큰 거리 = 395.30

## 02

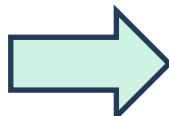
## ver2. Idea

A부대원 수만큼  
GPU Thread 생성



**1D Grid & 1D Block**  
Block layout : (1024, 1, 1)  
Grid layout : (ceil(M / 1024), 1, 1)

A 부대 정보  
&  
**B가 할당된 큐브 배열**



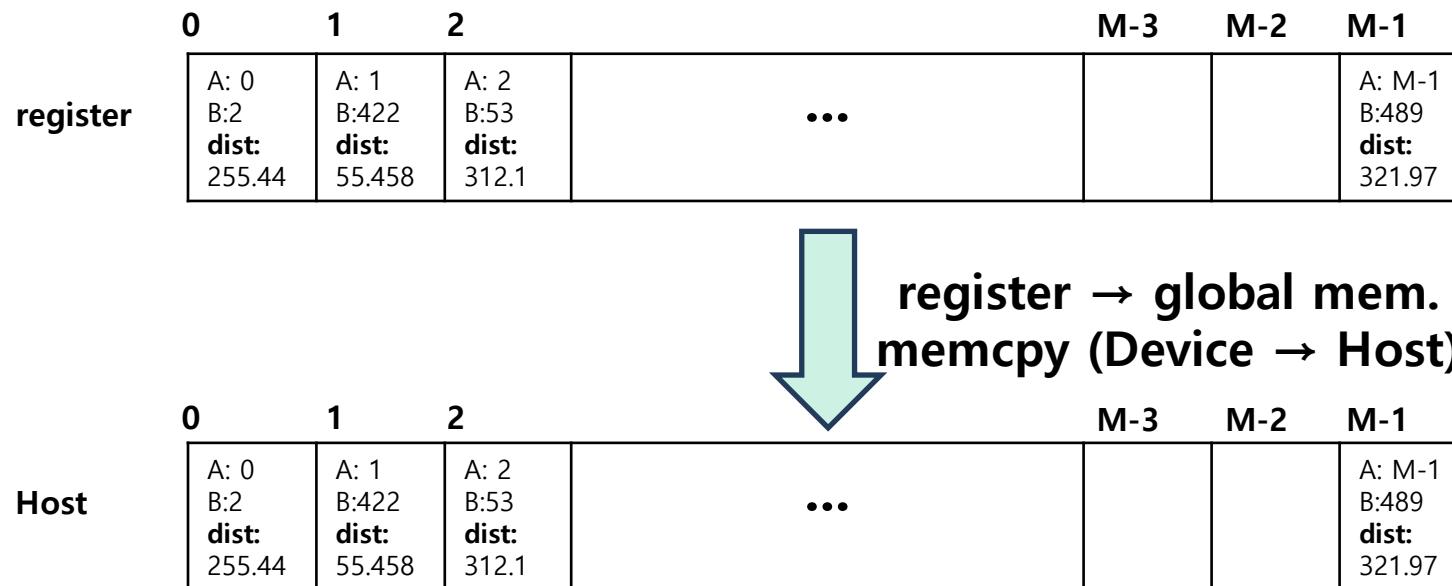
Global Memory

각 Thread별  
총 27개(3x3x3)의 큐브 search

가장 가까운 B ID, 거리를  
자신의 register에 기록

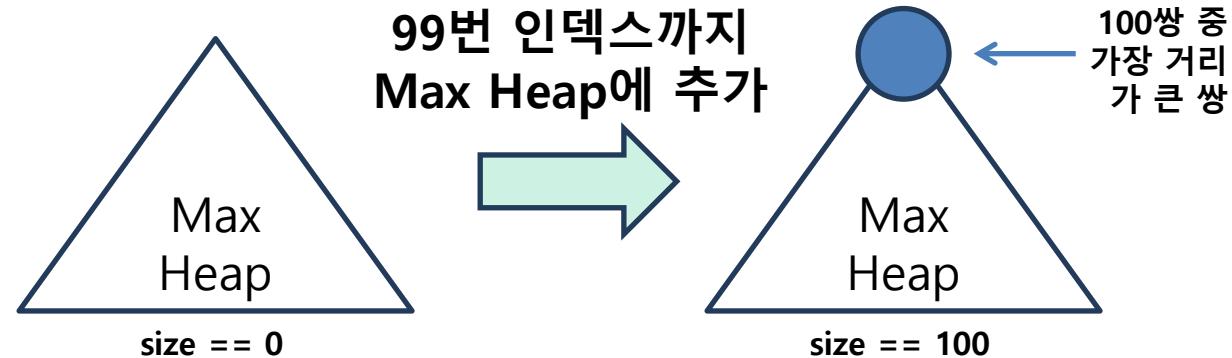
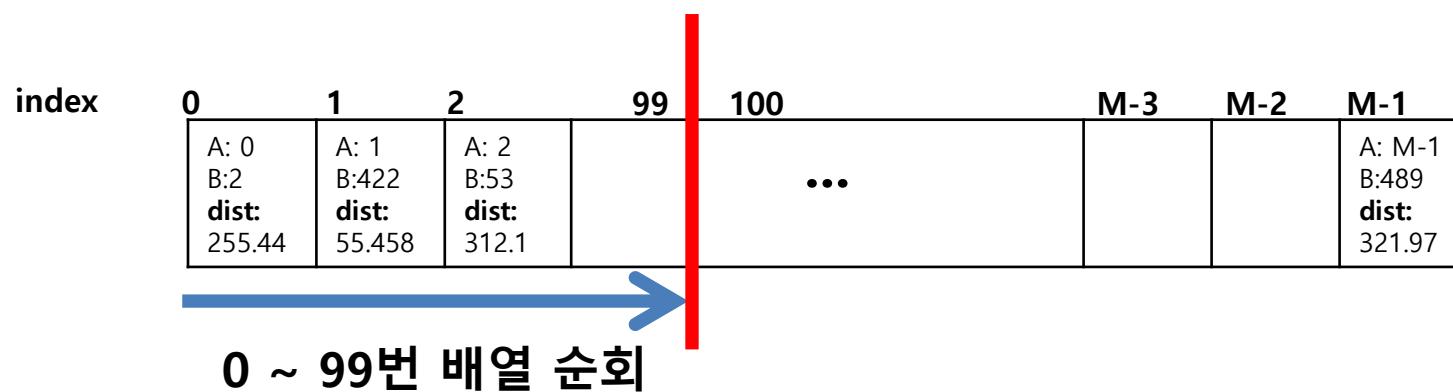
02

## ver2. Idea



## 02

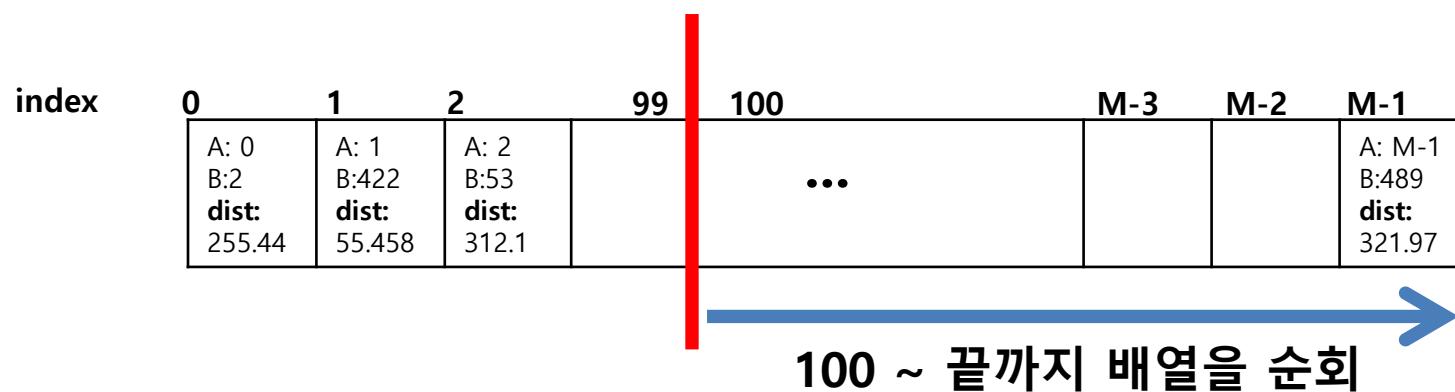
## ver2. Idea



거리가 가까운 100쌍을  
유지하기 위한  
자료구조

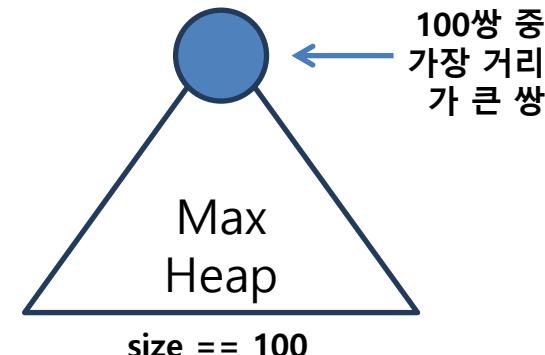
## 02

## ver2. Idea



현재 순회 중인 원소의 거리가  
Max Heap의 root 노드의  
거리보다 작을 경우,

root 노드를 pop하고,  
해당 원소를 push

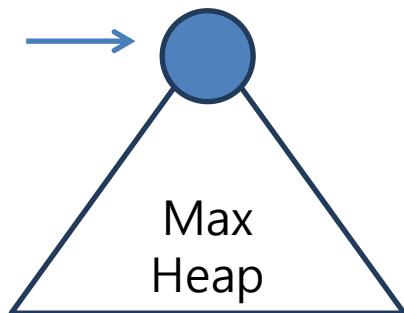


최종 Max Heap엔  
거리가 가까운 100쌍만 남음

02

## ver2. Idea

Max Heap에 저장된 쌍 중  
가장 거리가 큰 쌍



Max Heap을 pop하면서,  
최종 결과 배열의 역순으로 저장

| index  | 0                              | 1                                 | 2                              | ... | 97 | 98 | 99                               |
|--------|--------------------------------|-----------------------------------|--------------------------------|-----|----|----|----------------------------------|
| RESULT | A: 573<br>B:2<br>dist:<br>2.54 | A: 521<br>B:422<br>dist:<br>5.458 | A: 12<br>B:53<br>dist:<br>12.1 | ... |    |    | A: 1<br>B:489<br>dist:<br>321.97 |

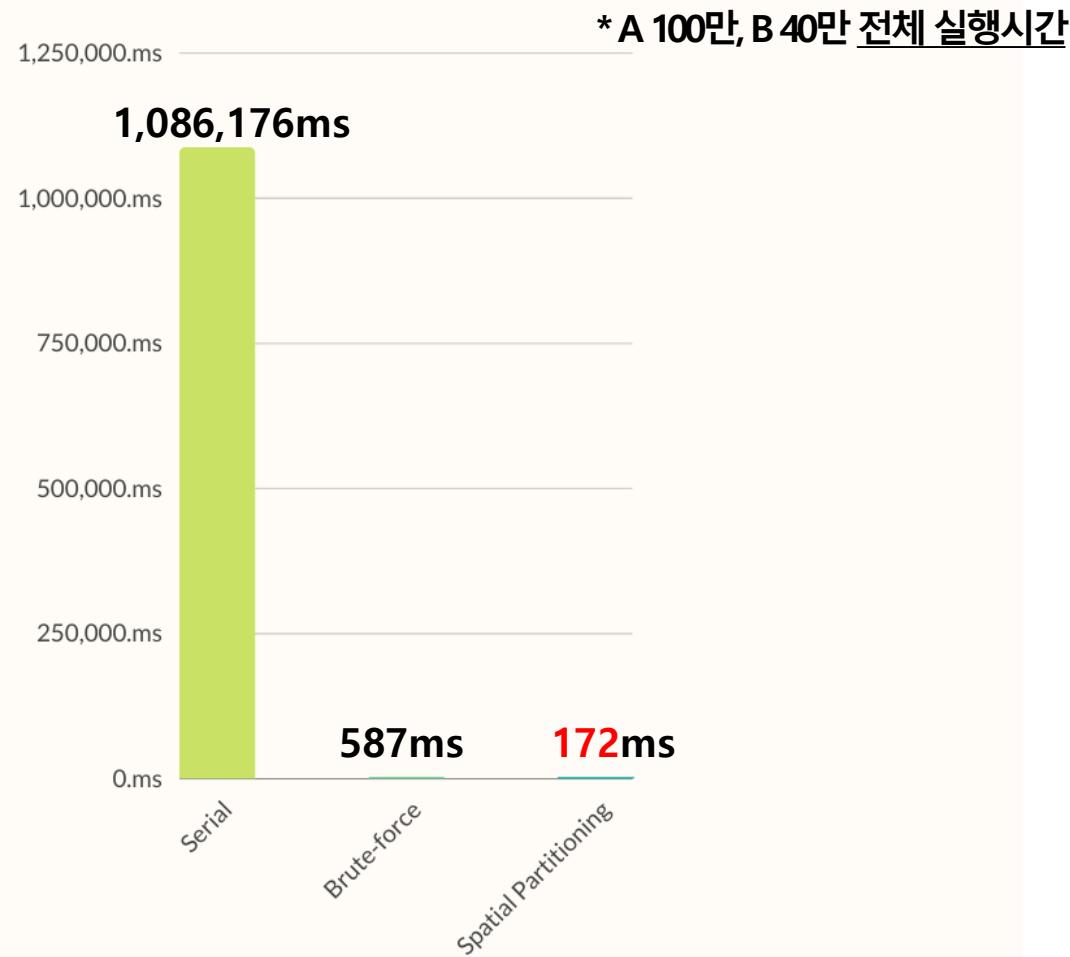
03

## 결과 및 분석

03

## 결과 분석

실습 환경 : Ryzen 7 5800 X / RTX 2080 Ti

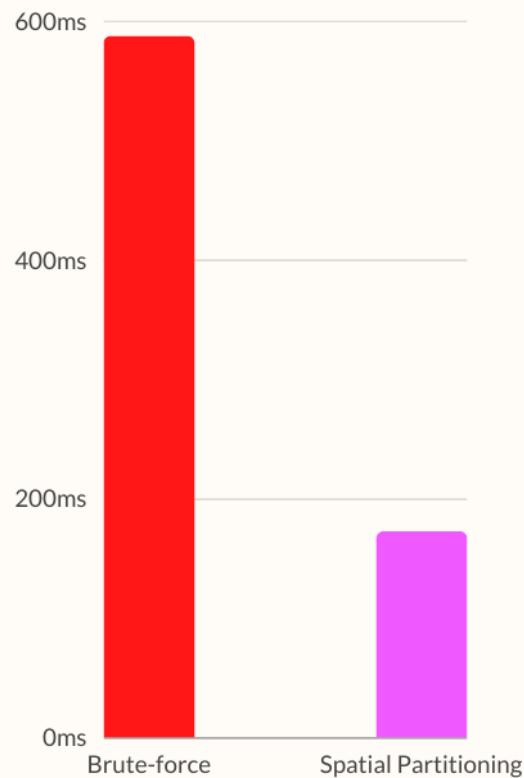


## 03

## 결과 분석

실습 환경 : Ryzen 7 5800 X / RTX 2080 Ti

\* A 100만, B 40만 전체 실행시간



Parallel ver.1 (Brute-force)  
약 587ms

Serial 대비 성능향상 약 1,850배  
정확도 : 100%

Parallel ver.2 (Spatial Partitioning)  
약 172ms

Serial 대비 성능향상 약 6,314배  
정확도 : 100%

03

## 결과 분석

실습 환경 : Ryzen 7 5800 X / RTX 2080 Ti

\* A 100만, B 40만 전체 실행시간

\* 단위 : ms

|                   | 1회        | 2회        | 3회        | 4회        | 5회        | 평균              |
|-------------------|-----------|-----------|-----------|-----------|-----------|-----------------|
| Parallel<br>ver.1 | 587.64890 | 594.12080 | 564.26280 | 585.85450 | 594.34940 | <b>585.2473</b> |
| Parallel<br>ver.2 | 172.96660 | 177.01230 | 173.69550 | 197.92030 | 205.06940 | <b>185.3328</b> |

평균 성능 기준 **약 3.16 배** 성능 향상

**THANK  
YOU**



## Shared memory bank conflicts

- Shared memory is as fast as registers if there are no bank conflicts
- Use the bank checker macro in the SDK to check for conflicts
  - warp\_serialize signal can usually be used to check for conflicts
- The fast case:
  - If all threads of a half-warp access different banks, there is no bank conflict
  - If all threads of a half-warp read the identical address, there is no bank conflict (broadcast)
- The slow case:
  - Bank Conflict: multiple threads in the same half-warp access the same bank
  - Must serialize the accesses
  - Cost = max # of simultaneous accesses to a single bank