

# Syntax Analysis

# Syntax Analysis is the second phase of compilation

❑ Comparison with lexical analysis:

Phase	Input	Output
<b>Lexer</b>	string of characters	string of tokens
<b>Parser</b>	string of tokens	Parse tree/AST

❑ Syntax analysis is also called **parsing**

- Because it produces a parse tree.
- AST (Abstract Syntax Tree) is a simplified parse tree.

# What is a Parse Tree?

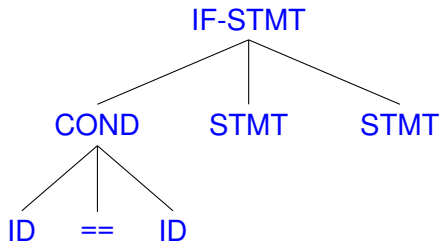
- ❑ **Parse tree:** a tree that represents grammatical structure
- ❑ Language constructs often have recursive structures

**If-stmt**  $\equiv$  **if** (EXPR) **then Stmt else Stmt fi**

**Stmt**  $\equiv$  **If-stmt** | **While-stmt** | ...

# A Parse Tree Example

- ❏ Code to be compiled:  
... if x==y then .... else ... fi
- ❏ Lexer:     ...     ...
- ❏ Parser:
  - Input: sequence of tokens  
... IF ID==ID THEN ... ELSE ... FI
  - Desired output:



# REs cannot express recursive program constructs

□ Example of a recursive construct is matching parenthesis:

# of "(" must equal # of ")"

✓  $(x+y)^*z$

✓  $((x+y)+y)^*z$

...

✓  $(\dots(((x+y)+y)+y)\dots)$

✗  $((x+y)+y)+y)^*z$

# REs cannot express recursive program constructs

□ Example of a recursive construct is matching parenthesis:

# of "(" must equal # of ")"

✓  $(x+y)^*z$

✓  $((x+y)+y)^*z$

...

✓  $(\dots(((x+y)+y)+y)\dots)$

✗  $((x+y)+y)+y)^*z$

□ Can regular expressions express this construct?

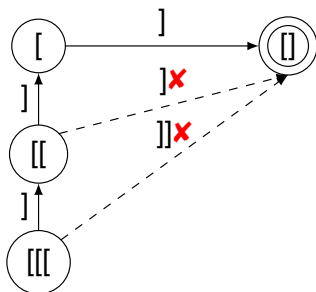
- Recall  $RL \equiv L(\text{Regular Expression}) \equiv L(\text{Finite Automata})$
- Boils down to whether an FA can accept this construct

# RE/FA is Not Powerful Enough

Describe strings with pattern  $[i]^i$  ( $i \geq 1$ )

# RE/FA is Not Powerful Enough

Describe strings with pattern  $[^i]^i$  ( $i \geq 1$ )



- “[”, “[[” are different states as only former accepts on “]”
- “[[”, “[[[” are different states as only former accepts on “]]”
- Infinite as for any  $[^i$ , there exists a  $[^{i+1}$  that is a new state
- Contradiction: no finite automaton accepts arbitrary nesting



# REs are not suitable for Syntax Analysis

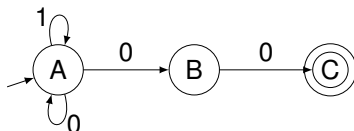
- ❑ REs cannot express recursive language constructs
- ❑ Programming languages belong to a category called CFLs
  - CFL is short for Context Free Language
  - CFLs are a strictly larger set than RLs
- ❑ To express CFLs, we need a new formalism: Grammars
- ❑ Grammars are general enough to express most languages
  - Regular Languages
  - Context Free Languages
  - Context Sensitive Languages
  - Recursively Enumerable Languages

# A Grammar defines a Language

- ❑ A grammar, along with tokens, defines a language
  - Like how English grammar defines the English language
- ❑ Grammars are defined using rigorous math just like for REs
- ❑ Recall the following definitions
  - Language: A set of strings over alphabet
  - Alphabet: A finite set of symbols
  - Empty string:  $\epsilon$
- ❑ We will also start calling strings in the language **sentences**

# An Example Grammar

- Language  $L = \{ \text{any string with "00" at the end} \}$



- Grammar  $G = \{ T, N, s, \delta \}$

where  $T = \{ 0, 1 \}$ ,  $N = \{ A, B \}$ ,  $s = A$ , and  
production rules  $\delta = \{ A \rightarrow 0A \mid 1A \mid 0B, B \rightarrow 0 \}$

- Derivation:** from grammar to language

- $A \Rightarrow 0A \Rightarrow 00B \Rightarrow 000$
- $A \Rightarrow 1A \Rightarrow 10B \Rightarrow 100$
- $A \Rightarrow 0A \Rightarrow 00A \Rightarrow 000B \Rightarrow 0000$
- $A \Rightarrow 0A \Rightarrow 01A \Rightarrow \dots$

# Grammar, formally defined

❏ A **grammar** consists of 4 components (**T**, **N**, **S**,  $\delta$ )

- T — set of **terminal** symbols
  - Leaves in the parse tree — essentially tokens
- N — set of **non-terminal** symbols
  - Internal nodes in the parse tree that expands into tokens
  - Language construct composed of one or more tokens like: statements, loops, functions, classes, ...
- S — A special non-terminal **start symbol**
  - Every string in language is derived from it
- $\delta$  — a set of **production** rules
  - “LHS  $\rightarrow$  RHS”: left-hand-side *produces* right-hand-side

# Production Rule and Derivation

## □ “LHS $\rightarrow$ RHS”

- Production rule to replace LHS with RHS
- Applied repeatedly to derive sentence from **S**

## □ $\beta \Rightarrow \alpha$ : string $\beta$ derives $\alpha$

- $\beta \Rightarrow \alpha$  — 1 step
- $\beta \Rightarrow^* \alpha$  — 0 or more steps
- $\beta \Rightarrow^+ \alpha$  — 0 or more steps

### ➤ example:

$A \Rightarrow 0A \Rightarrow 00B \Rightarrow 000$

$A \Rightarrow^* 000$

$A \Rightarrow^+ 000$

# Noam Chomsky Grammars

- Chomsky classified grammars into 4 types:

Type 0: recursive grammar

Type 1: context sensitive grammar

Type 2: context free grammar

Type 3: regular grammar

(Classification done based on form of production rules)

- The grammars produce the corresponding languages:

L(recursive grammar)  $\equiv$  recursively enumerable language

L(context sensitive grammar)  $\equiv$  context sensitive language

L(context free grammar)  $\equiv$  context free language

L(regular grammar)  $\equiv$  regular language

# Type 0: Unrestricted/Recursive Grammar

□ Type 0 grammar — unrestricted or recursive grammar

➤ Form of rules

$$\alpha \rightarrow \beta$$

where  $\alpha \in (N \cup T)^+$ ,  $\beta \in (N \cup T)^*$

➤ No restrictions on form of grammar rules

➤ Example:

$$aAB \rightarrow aCD$$

$$aAB \rightarrow aB$$

$$A \rightarrow \varepsilon$$

; erase rule is allowed

# Type 1: Context Sensitive Grammar

□ Type 1 grammar — context sensitive grammar

➤ Form of rules

$$\alpha A \beta \rightarrow \alpha \gamma \beta$$

where  $A \in N$ ,  $\alpha, \beta \in (N \cup T)^*$ ,  $\gamma \in (N \cup T)^+$

➤ Replace  $A$  by  $\gamma$  only if found in the context of  $\alpha$  and  $\beta$

➤ No erase rule

➤ Example:

$$aAB \rightarrow aCB$$



# Type 2: Context Free Grammar

□ Type 2 grammar — context free grammar

➤ Form of rules

$$A \rightarrow \gamma$$

where  $A \in N$ ,  $\gamma \in (N \cup T)^+$

➤ Can replace  $A$  by  $\gamma$  at any time — cannot specify context

# Type 2: Context Free Grammar

## □ Type 2 grammar — context free grammar

- Form of rules

$$A \rightarrow \gamma$$

where  $A \in N$ ,  $\gamma \in (N \cup T)^+$

- Can replace  $A$  by  $\gamma$  at any time — cannot specify context

## □ Are programming languages (PLs) context free ?

- Some PL constructs are context free: If-stmt, declaration
- Many are not: **def-before-use**, **matching formal/actual parameters**, etc.

# Type 3: Regular Grammar

□ Type 3 grammar — regular grammar

➤ Form of rules

$$A \rightarrow \alpha, \text{ or } A \rightarrow \alpha B$$

where  $A, B \in N, \alpha \in T$

➤ Regular grammar defines RE

➤ Can be used to define tokens for lexical analysis

➤ Example:

$$A \rightarrow 1A \mid 0$$

# Differentiate Type 2 and 3 Grammars

Language  $L1 = \{ [^i j^j \mid i, j \geq 1] \}$

➤ Regular grammar

$$\begin{aligned} S &\rightarrow [ S \mid [ T \\ T &\rightarrow ] T \mid ] \end{aligned}$$

Language  $L2 = \{ [^i j^i \mid i \geq 1] \}$

➤ Context free grammar

$$S \rightarrow [ S ] \mid [ ]$$

# Differentiate Type 1 and 2 Grammars

## Type 2 grammar (context free)

$$S \rightarrow D U$$
$$D \rightarrow \text{int } x; \quad | \quad \text{int } y;$$
$$U \rightarrow x=1; \quad | \quad y=1;$$

## Type 1 grammar (context sensitive)

$$S \rightarrow D U$$
$$D \rightarrow \text{int } x; \quad | \quad \text{int } y;$$
$$\text{int } x; U \rightarrow \text{int } x; x=1;$$
$$\text{int } y; U \rightarrow \text{int } y; y=1;$$

# Are Programming Languages Really Context Free?

## Language from type 2 grammar

- $S \Rightarrow DU \Rightarrow \text{int } x; U \Rightarrow \text{int } x; x=1;$
- $S \Rightarrow DU \Rightarrow \text{int } x; U \Rightarrow \text{int } x; y=1;$
- $S \Rightarrow DU \Rightarrow \text{int } y; U \Rightarrow \text{int } y; x=1;$
- $S \Rightarrow DU \Rightarrow \text{int } y; U \Rightarrow \text{int } y; y=1;$

## Language from type 1 grammar

- $S \Rightarrow DU \Rightarrow \text{int } x; U \Rightarrow \text{int } x; x=1;$
- $S \Rightarrow DU \Rightarrow \text{int } y; U \Rightarrow \text{int } y; y=1;$

# Are Programming Languages Really Context Free?

## Language from type 2 grammar

- $S \Rightarrow DU \Rightarrow \text{int } x; U \Rightarrow \text{int } x; x=1;$
- $S \Rightarrow DU \Rightarrow \text{int } x; U \Rightarrow \text{int } x; y=1;$
- $S \Rightarrow DU \Rightarrow \text{int } y; U \Rightarrow \text{int } y; x=1;$
- $S \Rightarrow DU \Rightarrow \text{int } y; U \Rightarrow \text{int } y; y=1;$

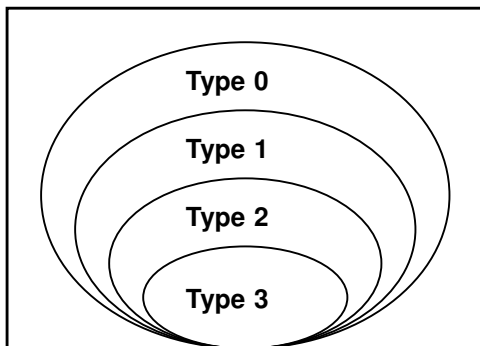
## Language from type 1 grammar

- $S \Rightarrow DU \Rightarrow \text{int } x; U \Rightarrow \text{int } x; x=1;$
- $S \Rightarrow DU \Rightarrow \text{int } y; U \Rightarrow \text{int } y; y=1;$

## PLs are context sensitive, why use CFG in parsing?

# The Chomsky Hierarchy of Grammars

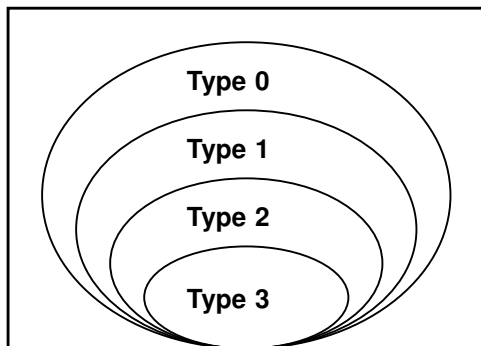
■  $RL \subset CFL \subset CSL \subset L(\text{Recursive Grammar})$





# The Chomsky Hierarchy of Grammars

■  $RL \subset CFL \subset CSL \subset L(\text{Recursive Grammar})$



■ However,  $L_y \subset L_x$  where  $L_x:[^i]^k$ —RG,  $L_y:[^i]^i$ —CFG

➤ Is it a problem?

# Context Free Grammars

# Syntax Analysis is a process of derivation

- ❑ Grammar is used to derive string or construct parser
- ❑ A **derivation** is a sequence of applications of rules
  - Starting from the **start symbol**
  - $S \Rightarrow \dots \Rightarrow \dots \Rightarrow \dots \Rightarrow (\text{sentence})$
- ❑ **Leftmost** and **Rightmost** derivations
  - At each derivation step, **leftmost** derivation always replaces the leftmost non-terminal symbol
  - **Rightmost** derivation always replaces the rightmost one

# Examples

$$E \rightarrow E * E \mid E + E \mid ( E ) \mid \text{id}$$

➤ leftmost derivation

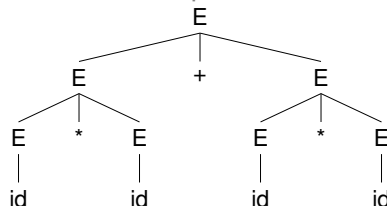
$$\begin{aligned} E &\Rightarrow E + E \Rightarrow E * E + E \Rightarrow \text{id} * E + E \Rightarrow \text{id} * \text{id} + E \Rightarrow \dots \\ &\Rightarrow \text{id} * \text{id} + \text{id} * \text{id} \end{aligned}$$

➤ rightmost derivation

$$\begin{aligned} E &\Rightarrow E + E \Rightarrow E + E * E \Rightarrow E + E * \text{id} \Rightarrow E + \text{id} * \text{id} \Rightarrow \dots \\ &\Rightarrow \text{id} * \text{id} + \text{id} * \text{id} \end{aligned}$$

# A Parse Tree represent the Derivation

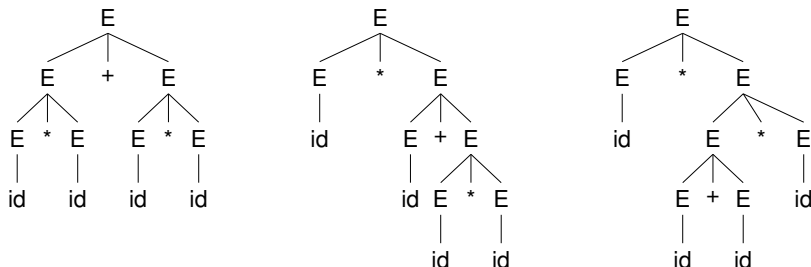
- This is the parse tree that represents both derivations:



- A parse tree
  - describes program structure (defined by the rules applied)
  - is agnostic of leftmost or rightmost derivation (as long as the same rules are applied in both)
- There are two types of nodes in a parse tree:
  - Leaves: terminals that form the sentence
  - Non-leaves: intermediate non-terminals in the derivation

# Different Rules result in different Parse Trees

Application of different rules result in different parse trees:



Note: each parse tree has a unique leftmost derivation

- First:  $E \Rightarrow E + E \Rightarrow E * E + E \Rightarrow id * E + E \Rightarrow id * id + E \Rightarrow \dots$
- Second:  $E \Rightarrow E * E \Rightarrow id * E \Rightarrow id * E + E \Rightarrow id * id + E \Rightarrow \dots$
- Third:  $E \Rightarrow E * E \Rightarrow id * E \Rightarrow id * E * E \Rightarrow id * E + E * E \Rightarrow \dots$

Same goes for rightmost derivations

# Ambiguity

- ❑ A grammar  $G$  is **ambiguous** if
  - there exist a string  $str \in L(G)$  such that
  - more than one parse tree derives  $str$ 
    - $\equiv$  there is more than leftmost derivation for  $str$
    - $\equiv$  there is more than rightmost derivation for  $str$
  
- ❑ Grammars that produce multiple parse trees is a problem
  - Each parse tree is a different interpretation of program
  
- ❑ Likely, there is an unambiguous version of the grammar
  - That accepts the same programming language
  - Programming languages are rarely inherently ambiguous

# Grammar can be rewritten to remove ambiguity

## Method I: to specify **precedence**

- build precedence into grammar, have different non-terminal for each precedence level
  - Lower precedence — relatively higher in tree (close to root)
  - Higher precedence — relatively lower in tree (far from root)
  - Same precedence — depends on associativity

$E \rightarrow E + E \mid E - E \mid E * E \mid E / E \mid E ^ E \mid ( E ) \mid id$

rewrite it to

$$E \rightarrow E + T \mid E - T \mid T$$
$$T \rightarrow T * F \mid T / F \mid F$$
$$F \rightarrow P ^ F \mid P$$
$$P \rightarrow id \mid ( E )$$



# How to Remove Ambiguity?

## ❏ Method II: to specify **associativity**

- Allow recursion only on either left or right non-terminal
  - Left associative — recursion on left non-terminal
  - Right associative — recursion on right non-terminal

## ❏ For the previous example,

$E \rightarrow E + E \dots$  ; allows both left/right associativity

rewrite it to

$E \rightarrow E + T \dots$  ; only left associativity

$F \rightarrow P \wedge F \dots$  ; only right associativity

# Ambiguity is undecidable for CFGs

- ❑ Decidable: computable using a Turing Machine
- ❑ It is **decidable** if a string is in a context free language
  - Implementing a parser is feasible for every CFL
- ❑ It is **undecidable** if a CFG is ambiguous
  - Checking ambiguity at compile time is impossible
  - Can only be checked reliably at runtime for a given string
  - In practice, tools like Yacc check for a more restricted grammar (e.g. LALR(1)) instead
    - LALR(1) is a subset of unambiguous grammars
    - Can be done easily at compile time

# The Two Outcomes of Parsing

- ❑ Outcome 1: Parser is able to derive input from grammar
  - Parser builds parse tree that represents the derivation
  
- ❑ Outcome 2: Parser is unable to derive input from grammar
  - Parser emits a syntax error with source code location

# The Two Outcomes of Parsing

- ❑ Outcome 1: Parser is able to derive input from grammar
  - Parser builds parse tree that represents the derivation
  
- ❑ Outcome 2: Parser is unable to derive input from grammar
  - Parser emits a syntax error with source code location
  
- ❑ How would you write a parser that does both well?

# Types of Parsers

## ❏ Universal parser

- Can parse any CFG e.g. Early's algorithm
- Powerful but extremely inefficient ( $O(N^3)$  where  $N$  is length of string)

## ❏ Top-down parser

- Tries to *expand* start symbol to input string
- Finds leftmost derivation
- Only works for a certain class of grammars
- Starts from root and expands into leaves
- Parser structure closely mimics grammar
  - Amenable to implementation by hand

# Types of Parsers (cont.)



## Bottom-up parser

- Tries to *reduce* the input string to the start symbol
- Finds reverse order of the rightmost derivation
- Works for wider class of grammars
- Starts at leaves and build tree in bottom-up fashion
- More amenable to generation by an automated tool

# What Output do We Want?

- ❑ The output of parsing is
  - parse tree, or
  - abstract syntax tree
  
- ❑ An **abstract syntax tree** is
  - similar to a parse tree but ignores some details
  - internal nodes may contain terminal symbols

# An Example

- Consider the grammar

$$E \rightarrow \text{int} \mid ( E ) \mid E + E$$

and an input

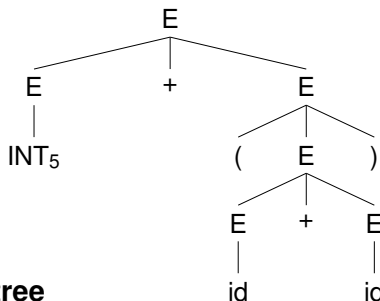
$$5 + ( 2 + 3 )$$

- After lexical analysis, we have a sequence of tokens

$$\text{INT}_5 \text{ ' + ' } ( \text{ INT}_2 \text{ ' + ' INT}_3 \text{ ' ) '}$$



# Parse Tree of the Input



## A parse tree

- Traces the operation of the parser
- Does capture the nested structure

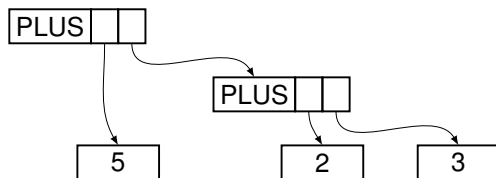


## but contains too much information

- parentheses
- single-successor nodes

# Abstract Syntax Tree

■ An **Abstract Syntax Tree (AST)** for the input



- **AST** also captures the nested structure
- **AST** abstracts from parse tree (a.k.a. concrete syntax tree)
- **AST** is more compact and contains only relevant info
- **ASTs** are used in most compilers rather than parse trees

# How are ASTs Constructed?

- ❑ Through implementation of **semantic actions**
- ❑ We already used them in project 1 to return token tuples
- ❑ To construct AST, we attach an **attribute** to each symbol  $X$ 
  - **$X.ast$**  — the constructed AST for symbol  $X$
- ❑ Extend each production rule with semantic actions, i.e.

$$X \rightarrow Y_1 Y_2 \dots Y_n \quad \{ \text{actions} \}$$

actions may define or use  $X.ast$ ,  $Y_i.ast$  ( $1 \leq i \leq n$ )

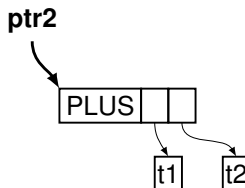
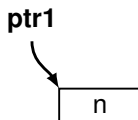
# Example

For the previous example, we have

$E \rightarrow$	<code>int</code>	<code>{ E.ast = mkleaf(int.lval) }</code>
	<code>  E1 + E2</code>	<code>{ E.ast = mkplus(E1.ast, E2.ast) }</code>
	<code>  (E1)</code>	<code>{ E.ast = E1.ast }</code>

Here, we use two pre-defined functions

- `ptr1=mkleaf(n)` — create a leaf node and assign value “n”
- `ptr2=mkplus(t1, t2)` — create a tree node and assign the root value “PLUS”, and two subtrees as t1 and t2



# AST Construction Steps

For input  $\text{INT}_5 \text{ ' + ' ( ' INT}_2 \text{ ' + ' INT}_3 \text{ ' ) '}$

Construction order given is for a top-down LL(1) parser  
(Order can change depending on parser implementation)

# AST Construction Steps

For input  $\text{INT}_5 \text{ '+' ' ( ' INT}_2 \text{ '+' INT}_3 \text{ ' )'}$

Construction order given is for a top-down LL(1) parser  
(Order can change depending on parser implementation)

$E1.ast = \text{mkleaf}(5)$



# AST Construction Steps

For input  $\text{INT}_5 \text{ '+' ' ( ' INT}_2 \text{ '+' INT}_3 \text{ ' ) '}$

Construction order given is for a top-down LL(1) parser  
(Order can change depending on parser implementation)

$E1.\text{ast} = \text{mkleaf}(5)$   $E2.\text{ast} = \text{mkleaf}(2)$



# AST Construction Steps

For input  $\text{INT}_5$  '+' '('  $\text{INT}_2$  '+'  $\text{INT}_3$  ')'

Construction order given is for a top-down LL(1) parser  
(Order can change depending on parser implementation)

$E1.\text{ast} = \text{mkleaf}(5)$   $E2.\text{ast} = \text{mkleaf}(2)$   $E3.\text{ast} = \text{mkleaf}(3)$





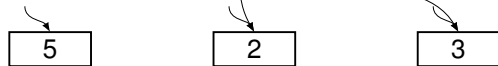
# AST Construction Steps

For input  $\text{INT}_5 \text{ '+' ' (' INT}_2 \text{ '+' INT}_3 \text{ ' )'}$

Construction order given is for a top-down LL(1) parser  
(Order can change depending on parser implementation)

$E4.\text{ast} = \text{mkplus}(E2.\text{ast}, E3.\text{ast})$

$E1.\text{ast} = \text{mkleaf}(5)$   $E2.\text{ast} = \text{mkleaf}(2)$   $E3.\text{ast} = \text{mkleaf}(3)$



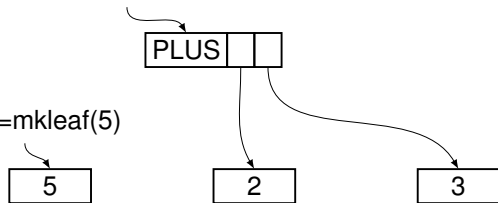
# AST Construction Steps

For input  $\text{INT}_5$  '+' '('  $\text{INT}_2$  '+'  $\text{INT}_3$  ')'

Construction order given is for a top-down LL(1) parser  
(Order can change depending on parser implementation)

$E4.\text{ast} = \text{mkplus}(E2.\text{ast}, E3.\text{ast})$

$E1.\text{ast} = \text{mkleaf}(5)$

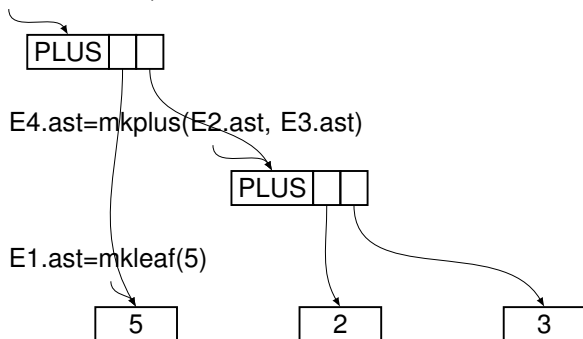


# AST Construction Steps

For input  $\text{INT}_5 \text{ '+' ' ( ' INT}_2 \text{ '+' INT}_3 \text{ ' ) '}$

Construction order given is for a top-down LL(1) parser  
(Order can change depending on parser implementation)

$E5.ast = \text{mkplus}(E1.ast, E4.ast)$

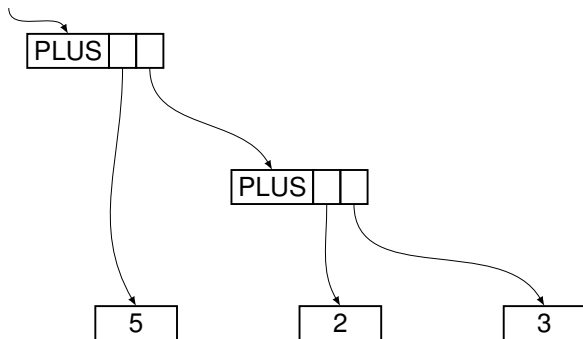


# AST Construction Steps

For input  $\text{INT}_5 \text{ '+' ' (' INT}_2 \text{ '+' INT}_3 \text{ ' )'}$

Construction order given is for a top-down LL(1) parser  
(Order can change depending on parser implementation)

$E5.ast = \text{mkplus}(E1.ast, E4.ast)$

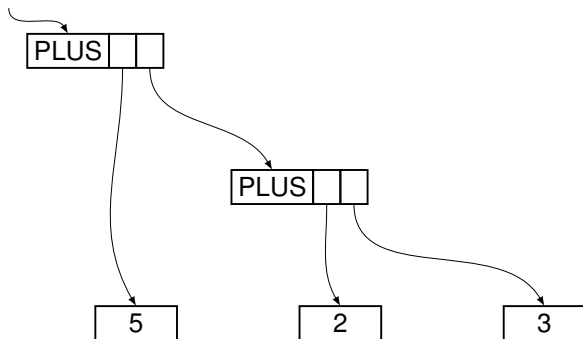


# AST Construction Steps

For input  $\text{INT}_5 \text{ '+' ' (' INT}_2 \text{ '+' INT}_3 \text{ ' )'}$

Construction order given is for a top-down LL(1) parser  
(Order can change depending on parser implementation)

$E5.ast = \text{mkplus}(E1.ast, E4.ast)$



# Summary

- ❏ Compilers specify program structure using CFG
  - Most programming languages are not context free
  - Context sensitive analysis can easily separate out to semantic analysis phase
  
- ❏ A parser uses CFG to
  - ... answer if an input  $str \in L(G)$
  - ... and build a parse tree
  - ... or build an AST instead
  - ... and pass it to the rest of compiler

# Parsing

# Parsing

- ❑ We will study two approaches
- ❑ Top-down
  - Easier to understand and implement manually
- ❑ Bottom-up
  - More powerful, can be implemented automatically



# Example

Consider a CFG grammar G

$$\begin{array}{lll} S \rightarrow A B & A \rightarrow a C & B \rightarrow b D \\ D \rightarrow d & C \rightarrow c & \end{array}$$

Actually, this language has only one sentence, i.e.

$$L(G) = \{ acbd \}$$

## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)

S

## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbD$  (2)  
 $\Rightarrow acbd$  (1)

# Example

Consider a CFG grammar G

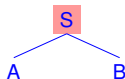
$$S \rightarrow AB \quad A \rightarrow aC \quad B \rightarrow bD$$
$$D \rightarrow d \quad C \rightarrow c$$

Actually, this language has only one sentence, i.e.

$$L(G) = \{ acbd \}$$

## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)



## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbd$  (2)  
 $\Rightarrow acbd$  (1)

# Example

Consider a CFG grammar G

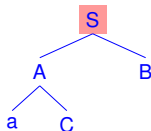
$$\begin{array}{lll} S \rightarrow AB & A \rightarrow aC & B \rightarrow bD \\ D \rightarrow d & C \rightarrow c & \end{array}$$

Actually, this language has only one sentence, i.e.

$$L(G) = \{ acbd \}$$

## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)



## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbD$  (2)  
 $\Rightarrow acbd$  (1)

# Example

Consider a CFG grammar G

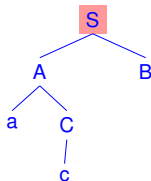
$$\begin{array}{lll} S \rightarrow AB & A \rightarrow aC & B \rightarrow bD \\ D \rightarrow d & C \rightarrow c & \end{array}$$

Actually, this language has only one sentence, i.e.

$$L(G) = \{ acbd \}$$

## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)



## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbD$  (2)  
 $\Rightarrow acbd$  (1)

# Example

Consider a CFG grammar  $G$

$$S \rightarrow AB \quad A \rightarrow aC \quad B \rightarrow bD$$
$$D \rightarrow d \quad C \rightarrow c$$

Actually, this language has only one sentence, i.e.

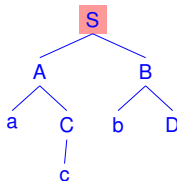
$$L(G) = \{acbd\}$$

## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)

## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbd$  (2)  
 $\Rightarrow acbd$  (1)



# Example

Consider a CFG grammar G

$$\begin{array}{lll} S \rightarrow AB & A \rightarrow aC & B \rightarrow bD \\ D \rightarrow d & C \rightarrow c & \end{array}$$

Actually, this language has only one sentence, i.e.

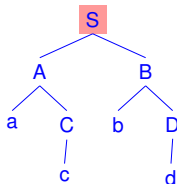
$$L(G) = \{ acbd \}$$

## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)

## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbd$  (2)  
 $\Rightarrow acbd$  (1)



# Example

Consider a CFG grammar G

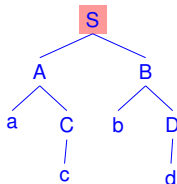
$$\begin{array}{lll} S \rightarrow AB & A \rightarrow aC & B \rightarrow bD \\ D \rightarrow d & C \rightarrow c & \end{array}$$

Actually, this language has only one sentence, i.e.

$$L(G) = \{ acbd \}$$

## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)



## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbd$  (2)  
 $\Rightarrow acbd$  (1)

a c b d

# Example

Consider a CFG grammar G

$$S \rightarrow AB \quad A \rightarrow aC \quad B \rightarrow bD$$

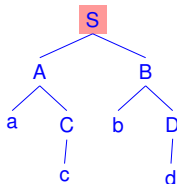
$$D \rightarrow d \quad C \rightarrow c$$

Actually, this language has only one sentence, i.e.

$$L(G) = \{acbd\}$$

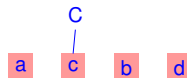
## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)



## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbd$  (2)  
 $\Rightarrow acbd$  (1)





# Example

Consider a CFG grammar G

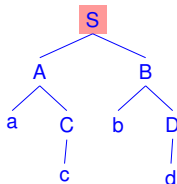
$S \rightarrow AB$      $A \rightarrow aC$      $B \rightarrow bD$   
 $D \rightarrow d$      $C \rightarrow c$

Actually, this language has only one sentence, i.e.

$L(G) = \{acbd\}$

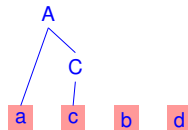
## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)



## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbd$  (2)  
 $\Rightarrow acbd$  (1)



# Example

Consider a CFG grammar G

$S \rightarrow AB$       $A \rightarrow aC$       $B \rightarrow bD$

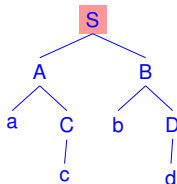
$D \rightarrow d$       $C \rightarrow c$

Actually, this language has only one sentence, i.e.

$L(G) = \{acbd\}$

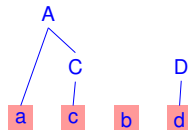
## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)



## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbd$  (2)  
 $\Rightarrow acbd$  (1)



# Example

Consider a CFG grammar G

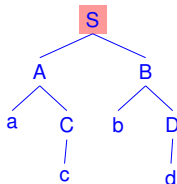
$$S \rightarrow AB \quad A \rightarrow aC \quad B \rightarrow bD$$
$$D \rightarrow d \quad C \rightarrow c$$

Actually, this language has only one sentence, i.e.

$$L(G) = \{acbd\}$$

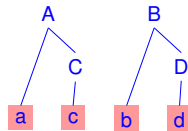
## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)



## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbd$  (2)  
 $\Rightarrow acbd$  (1)



# Example

Consider a CFG grammar  $G$

$S \rightarrow AB$        $A \rightarrow aC$        $B \rightarrow bD$

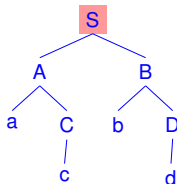
$D \rightarrow d$        $C \rightarrow c$

Actually, this language has only one sentence, i.e.

$L(G) = \{acbd\}$

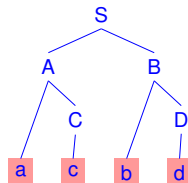
## Leftmost Derivation:

$S \Rightarrow AB$  (1)  
 $\Rightarrow aCB$  (2)  
 $\Rightarrow acB$  (3)  
 $\Rightarrow acbD$  (4)  
 $\Rightarrow acbd$  (5)



## Rightmost Derivation:

$S \Rightarrow AB$  (5)  
 $\Rightarrow AbD$  (4)  
 $\Rightarrow Abd$  (3)  
 $\Rightarrow aCbd$  (2)  
 $\Rightarrow acbd$  (1)



# Top Down Parsers

# Backtracking or Predictive?

- ❑ How does a parser choose between production rules?
  - Given  $A \rightarrow \alpha|\beta$ , expand  $A$  to  $\alpha$  or  $\beta$ ?

- ❑ Backtracking parser: exhaustively tries all rules
  - When input mismatch, backtrack to alternative rule
  - Con Non-linear time due to exhaustive search
  - Con Complex to roll back semantic actions on backtrack
  - Pro Can parse most CFGs (except left-recursion)

- ❑ Predictive parser: predict correct rule using lookahead
  - Looks ahead  $k$  input symbols to make prediction
  - Con Can parse only a subset of CFGs (dependent on  $k$ )
  - Pro Linear time as only correct derivations are done
  - Pro Simple structure as there is no need to backtrack

- ❑ Parsers can be backtracking or predictive (or both).

# Recursive Descent or Table Driven?

- ❑ How is the parser implementation done?
  - Hand-coded parsers are typically recursive descent
  - Auto-generated parsers are table driven
  
- ❑ Recursive descent parser: each non-terminal is a function
  - Function is in charge of expanding non-terminal
  - Descends parse tree via recursive calls to non-terminals
  - Hand-written but easier to customize and control
  - Typically uses backtracking rather than prediction
  
- ❑ Table driven parser: uses a table of predictions
  - Similar to lexer, uses a table to decide on next production
  - Table indexed by non-terminal and  $k$  lookahead symbols
  - Similar to lexer, table can be generated from grammar
  - Always predictive but can use backtracking if needed

# Backtracking Example

$$E \rightarrow T + E \mid T$$
$$T \rightarrow \text{int} * T \mid \text{int} \mid ( E )$$

input string:     `int * int`

start symbol:     `E`

initial parse tree is   `E`



# Backtracking Example

$$E \rightarrow T + E \mid T$$
$$T \rightarrow \text{int} * T \mid \text{int} \mid ( E )$$

input string:     `int * int`

start symbol:     `E`

initial parse tree is   `E`

 Assume: when there are alternative rules, try right rule first

# Parsing Sequence using Backtracking

E

# Parsing Sequence using Backtracking

$E \Rightarrow T$

– pick right most rule  $E \rightarrow T$

# Parsing Sequence using Backtracking

$E \Rightarrow T \Rightarrow (E)$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$

# Parsing Sequence using Backtracking

$E \Rightarrow T \Rightarrow (E)$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”

# Parsing Sequence using Backtracking

$E \Rightarrow T \Rightarrow (E)$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- failure, backtrack one level

# Parsing Sequence using Backtracking

$E \Rightarrow T \Rightarrow \cancel{(E)}$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- failure, backtrack one level

# Parsing Sequence using Backtracking

$E \Rightarrow T \Rightarrow \cancel{(E)}$

$\Rightarrow \text{int}$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int}$
- “int” matches input “int”



# Parsing Sequence using Backtracking

$E \Rightarrow T \Rightarrow (E)$

$\Rightarrow \text{int}$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- failure, backtrack one level
- pick  $T \rightarrow \text{int}$
- “int” matches input “int”
- however, we expect more tokens
- failure, backtrack one level

# Parsing Sequence using Backtracking

$E \Rightarrow T \Rightarrow (E)$

$\Rightarrow \text{int}$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- failure, backtrack one level
- pick  $T \rightarrow \text{int}$
- “int” matches input “int”
- however, we expect more tokens
- failure, backtrack one level

# Parsing Sequence using Backtracking

$E \Rightarrow T \Rightarrow (E)$

$\Rightarrow \text{int}$

$\Rightarrow \text{int} * T$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- failure, backtrack one level
- pick  $T \rightarrow \text{int}$
- “int” matches input “int”
- however, we expect more tokens
- failure, backtrack one level
- pick  $T \rightarrow \text{int} * T$

# Parsing Sequence using Backtracking

$$E \Rightarrow T \Rightarrow \cancel{(E)}$$

$$\Rightarrow \cancel{\text{int}}$$

$$\Rightarrow \text{int} * T \Rightarrow \text{int} * (E)$$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int}$
- “int” matches input “int”
- however, we expect more tokens
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int} * T$
- pick  $T \rightarrow \text{int} * (E)$

# Parsing Sequence using Backtracking

$$E \Rightarrow T \Rightarrow (E)$$

$$\Rightarrow \text{int}$$

$$\Rightarrow \text{int} * T \Rightarrow \text{int} * (E)$$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int}$
- “int” matches input “int”
- however, we expect more tokens
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int} * T$
- pick  $T \rightarrow \text{int} * (E)$
- “(” does not match input “int”
- **failure, backtrack one level**

# Parsing Sequence using Backtracking

$$E \Rightarrow T \Rightarrow (E)$$

$$\Rightarrow \text{int}$$

$$\Rightarrow \text{int} * T \Rightarrow \text{int} * (E)$$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int}$
- “int” matches input “int”
- however, we expect more tokens
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int} * T$
- pick  $T \rightarrow \text{int} * (E)$
- “(” does not match input “int”
- **failure, backtrack one level**

# Parsing Sequence using Backtracking

$$E \Rightarrow T \Rightarrow \cancel{(E)}$$

$$\Rightarrow \cancel{\text{int}}$$

$$\Rightarrow \text{int} * T \Rightarrow \cancel{\text{int} * (E)}$$

$$\Rightarrow \text{int} * \text{int}$$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int}$
- “int” matches input “int”
- however, we expect more tokens
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int} * T$
- pick  $T \rightarrow \text{int} * (E)$
- “(” does not match input “int”
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int}$

# Parsing Sequence using Backtracking

$$E \Rightarrow T \Rightarrow \cancel{(E)}$$

$$\Rightarrow \cancel{\text{int}}$$

$$\Rightarrow \text{int} * T \Rightarrow \cancel{\text{int} * (E)}$$

$$\Rightarrow \text{int} * \text{int}$$

- pick right most rule  $E \rightarrow T$
- pick right most rule  $T \rightarrow (E)$
- “(” does not match “int”
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int}$
- “int” matches input “int”
- however, we expect more tokens
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int} * T$
- pick  $T \rightarrow \text{int} * (E)$
- “(” does not match input “int”
- **failure, backtrack one level**
- pick  $T \rightarrow \text{int}$
- **match, accept**



# Recursive Descent Parser with Backtracking

# Recursive Descent Parsing Implementation

- ❑ When expanding a non-terminal, try all productions until
  - A production is found that generates a portion of the input, or
  - No production is found that generates a portion of the input, in which case backtrack to previous non-terminal
  
- ❑ Create a function for each non-terminal
  1. For RHS of each production rule,
    - a. For a terminal, match with input symbol and consume
    - b. For a non-terminal, call function for that non-terminal
    - c. If match succeeds for entire RHS, return success
    - d. If match fails, regurgitate input and try next RHS
  2. If match succeeds for any rule, apply that rule to LHS
  
- ❑ If entire input string matched with start symbol, success!

# A Hand-coded Recursive Descent Parser

Sample implementation of parser for previous grammar:

$$E \rightarrow T + E \mid T$$

$$T \rightarrow \text{int} * T \mid \text{int} \mid ( E )$$

```
char fetchNext() {
    // Fetch one character
}
void regurgitate() {
    // Unfetch one character
}

bool expr() {
rule1:
    if(!term()) goto rule2;
    if (fetchNext()!=AddNum) {
        regurgitate();
        goto rule2;
    }
    if(!expr()) goto rule2;
    return true;
rule2:
    if (!term()) return false;
    return true;
}
```

```
bool term() {
rule1:
    if (fetchNext()!=IntNum) {
        regurgitate();
        goto rule2;
    }
    if (fetchNext()!=StarNum) {
        regurgitate();
        goto rule2;
    }
    if(!term()) goto rule2;
    return true;
rule2:
    if (fetchNext()!=IntNum) {
        regurgitate();
        goto rule3;
    }
    return true;
rule3:
    ...
}
```

# Recursive Descent has a Left Recursion Problem

❑ Recursive descent doesn't work if grammar is left recursive

❑ Why is left recursion a problem?

- For left recursive grammar

$$A \rightarrow A b \mid c$$

- We may repeatedly choose to apply  $A b$

$$A \Rightarrow A b \Rightarrow A b b \dots$$

- Sentence can grow indefinitely w/o consuming input
- How do you know when to stop recursion and choose  $c$ ?

# Recursive Descent has a Left Recursion Problem

- ❑ Recursive descent doesn't work if grammar is left recursive
- ❑ Why is left recursion a problem?
  - For left recursive grammar
$$A \rightarrow A b \mid c$$
  - We may repeatedly choose to apply  $A b$ 
$$A \Rightarrow A b \Rightarrow A b b \dots$$
  - Sentence can grow indefinitely w/o consuming input
  - How do you know when to stop recursion and choose  $c$ ?
- ❑ Rewrite the grammar so that it is right recursive
  - Which expresses the same language

# Removing Left Recursion

- All immediate left recursion can be eliminated this way:

$$A \rightarrow A x \mid y$$

change to

$$A \rightarrow y A'$$

$$A' \rightarrow x A' \mid \epsilon$$

- Not all left recursion is immediate  
(Recursion may involve multiple non-terminals)

$$A \rightarrow BC \mid D$$

$$B \rightarrow AE \mid F$$

... see Section 4.3 for *elimination of general left recursion*

... (not required for this course)

# Table Driven Parser using Predictions

# Predictive Parsers can avoid Backtracking

- ❑ Predict correct production rule based on  $k$  lookahead
  - Backtracking can be avoided if grammar limited to  $LL(k)$
  
- ❑  $LL(k)$  Parser
  - L — left to right scan
  - L — leftmost derivation
  - $k$  —  $k$  symbols of lookahead
  - A predictive parser that uses  $k$  lookahead tokens
  
- ❑  $LL(k)$  Grammar
  - A grammar parse-able by  $LL(k)$  parser with no backtracking
  
- ❑  $LL(k)$  Language
  - A language that can be expressed as a  $LL(k)$  grammar
  - $LL(k)$  languages are a restricted subset of CFLs
  - But many languages are  $LL(k)$ . In fact, many are  $LL(1)$ !



# Left factoring can make grammars LL(1)

## □ An LL(1) grammar

- First terminal of every alternative production is unique

$$\begin{aligned} A &\rightarrow a B D \mid b B B \\ B &\rightarrow c \mid b c e \\ D &\rightarrow d \end{aligned}$$

## □ What if no LL(1)? Left factor to make it LL(1)!

- What if production rules for  $A$  was changed to below?

$$A \rightarrow a B D \mid a B B$$

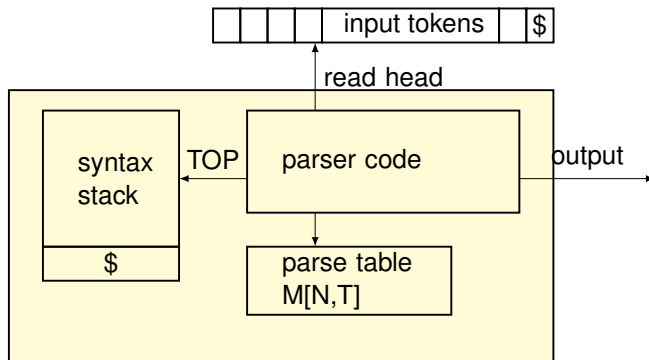
- Left factor  $a B$  to enable prediction:

$$\begin{aligned} A &\rightarrow a B A' \mid a B A' \\ A' &\rightarrow D \mid B \end{aligned}$$

## □ In general, if you see $A \rightarrow \alpha\beta \mid \alpha\gamma$ , change to:

$$\begin{aligned} A &\rightarrow \alpha A' \\ A' &\rightarrow \beta \mid \gamma \end{aligned}$$

# A Table Driven Pushdown Automaton



Syntax stack — hold right hand side (RHS) of grammar rules

Parse table  $M[A,b]$  — an entry containing rule “ $A \rightarrow \dots$ ” or error

Parser code — next action based on **(current token, stack top)**

Table can be automatically generated from grammar (just like lexers)

# A Sample Parse Table

	int	*	+	(	)	\$
E	$E \rightarrow TX$			$E \rightarrow TX$		
X			$X \rightarrow +E$		$X \rightarrow \epsilon$	$X \rightarrow \epsilon$
T	$T \rightarrow \text{int } Y$			$T \rightarrow (E)$		
Y		$Y \rightarrow *T$	$Y \rightarrow \epsilon$		$Y \rightarrow \epsilon$	$Y \rightarrow \epsilon$

☐ Predicts rule based on (current non-terminal, lookahead)

- **First column** lists all non-terminals
- **First row** lists all possible terminals and \$
- A table entry contains one production (one prediction)

☐ What if an entry has more than one production?

- Means that this grammar is not LL(1)
- A parser can handle this situation by either:
  - Throwing an error to grammar writer to fix the problem
  - Resorting to backtracking to try out both productions

# Pseudocode for Table-Driven Parser

**X** — symbol at the top of the syntax stack

**a** — current input symbol

## Parsing based on **(X,a)**

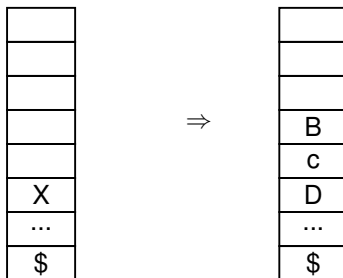
- If  $X == a == \$$ , then
  - parser halts with “success”
- If  $X == a \neq \$$ , then
  - pop X from stack **and** advance input head
- If  $X \neq a$ , then
  - Case (a): if  $X \in T$ , then
    - parser halts with “failed”, input rejected
  - Case (b): if  $X \in N$ ,  $M[X,a] = “X \rightarrow RHS”$ 
    - pop X **and** push RHS to stack in reverse order


# Push RHS in Reverse Order

**X** — symbol at the top of the syntax stack

**a** — current input symbol

if  $M[X,a] = "X \rightarrow B \ c \ D"$



 Why? Because that is the order of leftmost derivation.

# Applying LL(1) Parsing to a Grammar

❏ Given our old grammar

$$E \rightarrow T + E \mid T$$
$$T \rightarrow \text{int} * T \mid \text{int} \mid ( E )$$

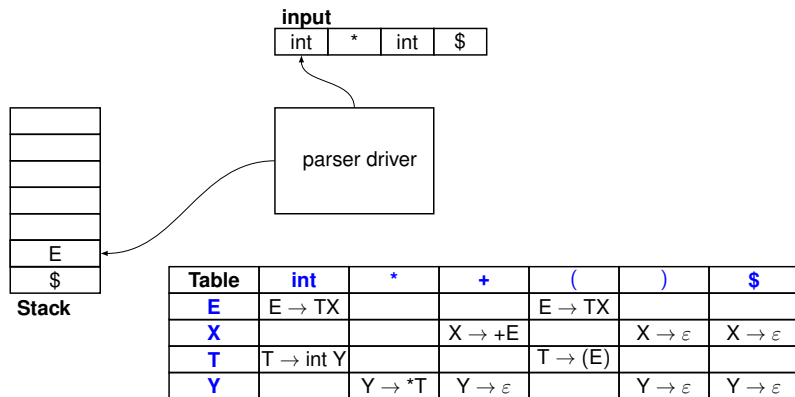
➤ Requires left factoring of  $T$  and  $\text{int}$

❏ After rewriting grammar, we have

$$E \rightarrow T X$$
$$X \rightarrow + E \mid \epsilon$$
$$T \rightarrow \text{int} Y \mid ( E )$$
$$Y \rightarrow * T \mid \epsilon$$

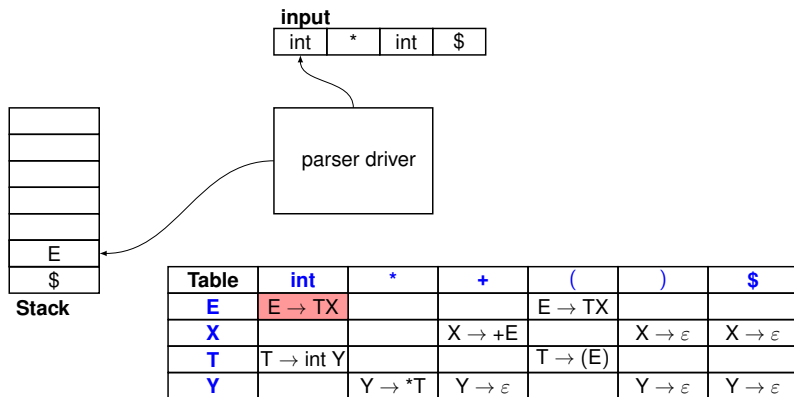
# Using the Parse Table

■ To recognize “int \* int”



# Using the Parse Table

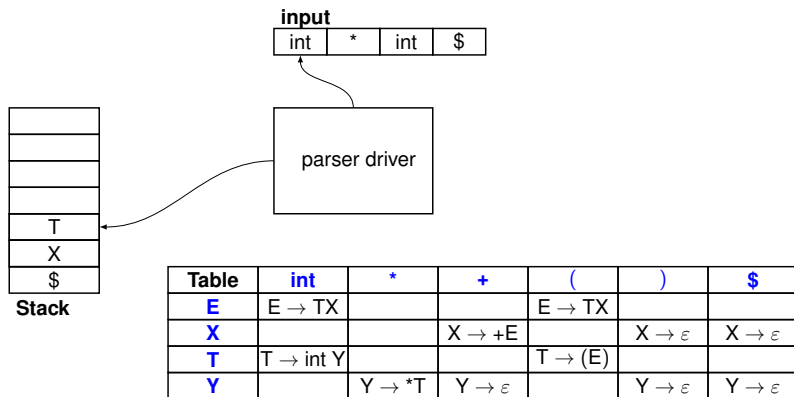
■ To recognize “int \* int”





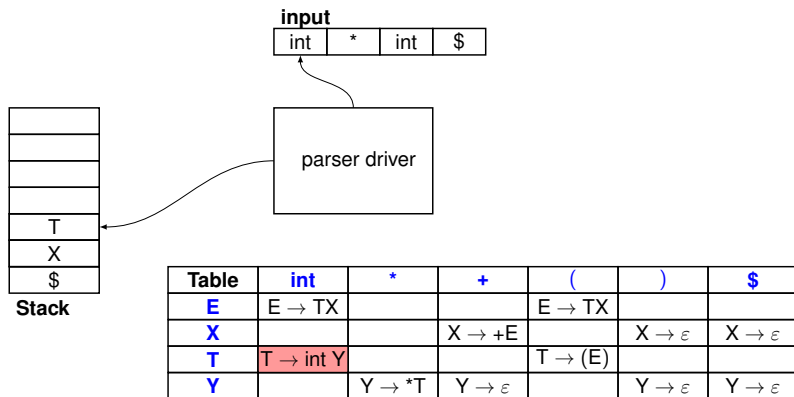
# Using the Parse Table

■ To recognize “int \* int”



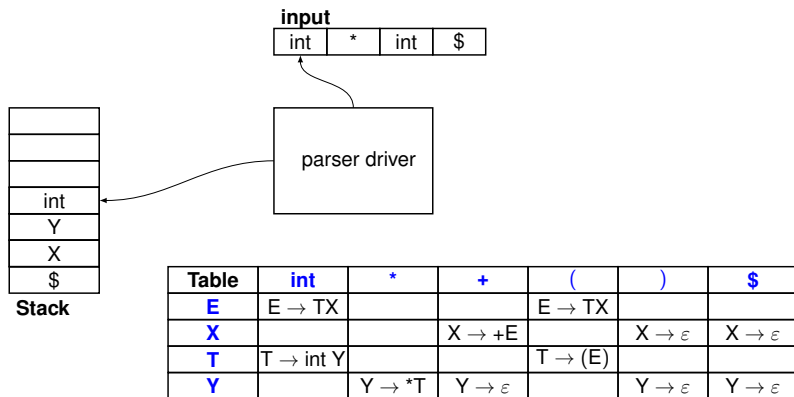
# Using the Parse Table

■ To recognize “int \* int”



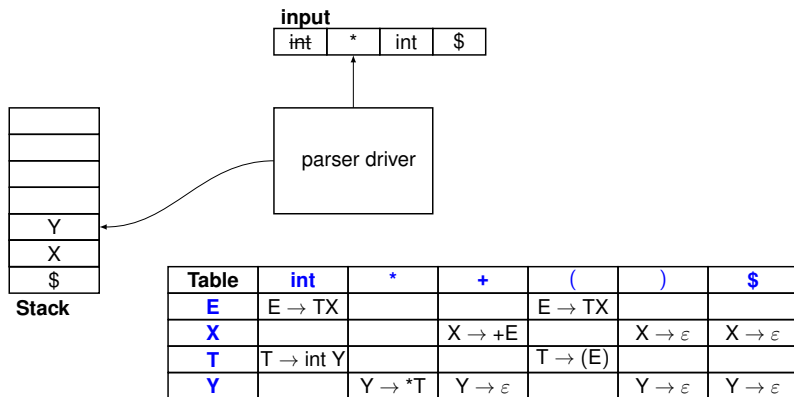
# Using the Parse Table

□ To recognize “int \* int”



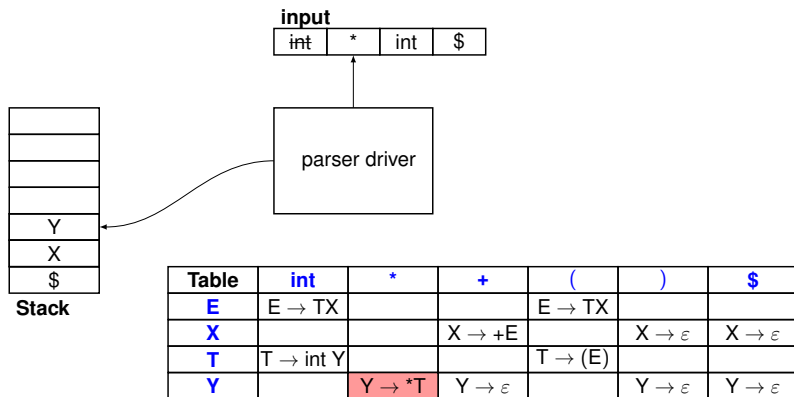
# Using the Parse Table

■ To recognize “int \* int”



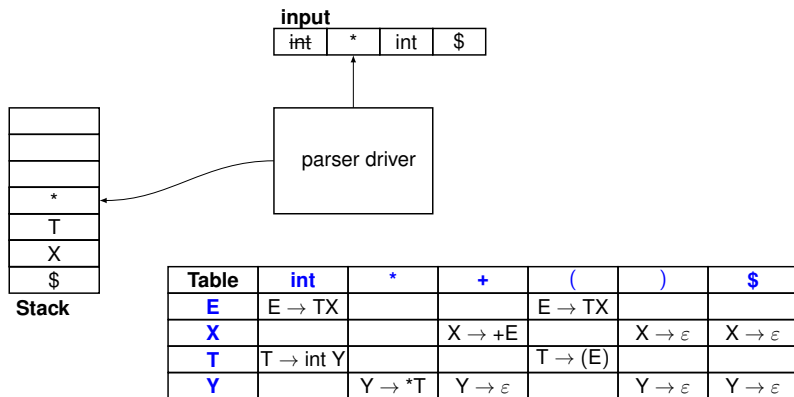
# Using the Parse Table

■ To recognize “int \* int”



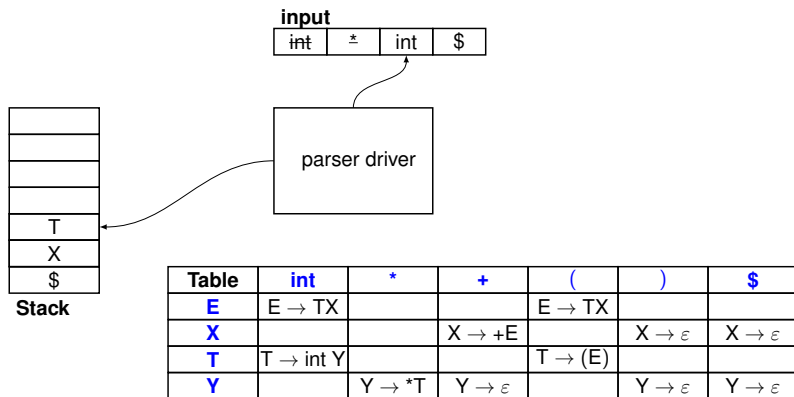
# Using the Parse Table

□ To recognize “int \* int”



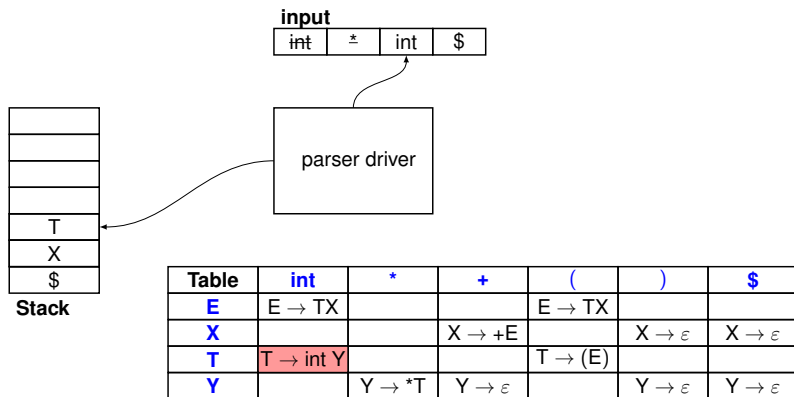
# Using the Parse Table

■ To recognize “int \* int”



# Using the Parse Table

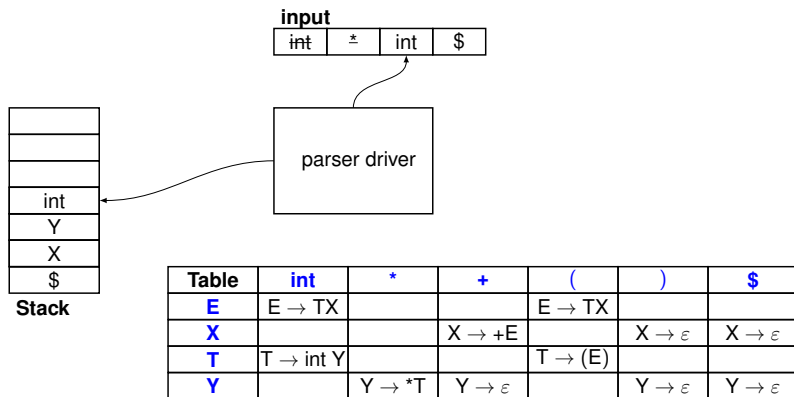
■ To recognize “int \* int”





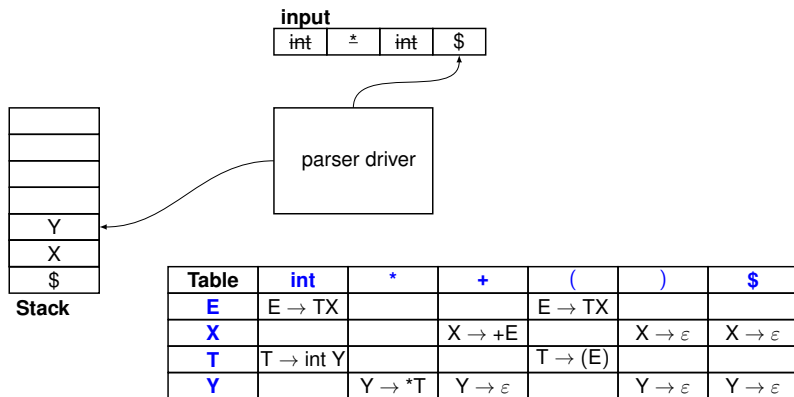
# Using the Parse Table

□ To recognize “int \* int”



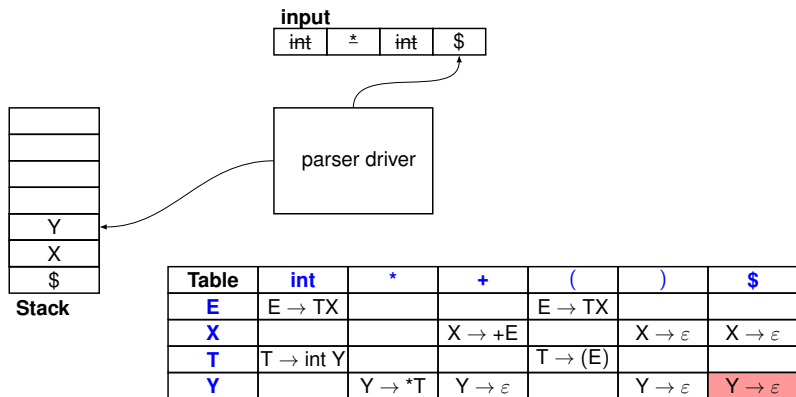
# Using the Parse Table

■ To recognize “int \* int”



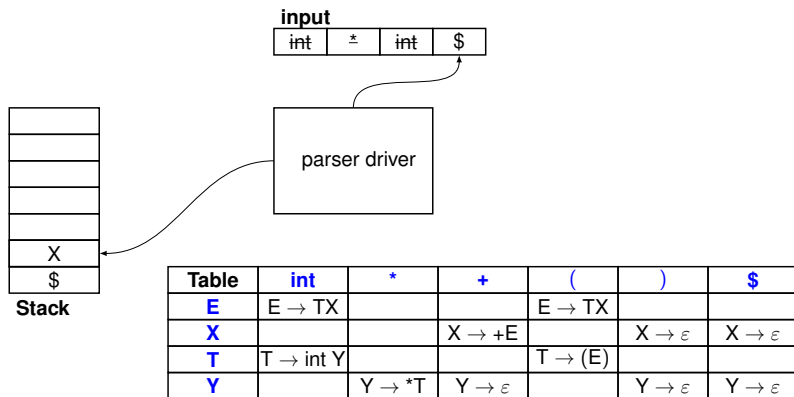
# Using the Parse Table

■ To recognize “int \* int”



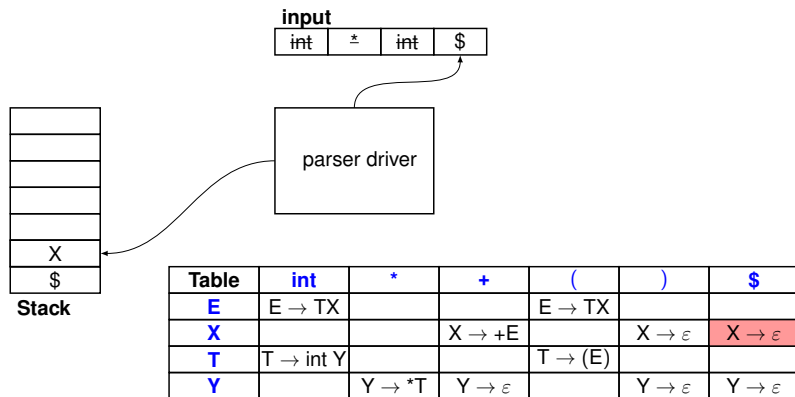
# Using the Parse Table

■ To recognize “int \* int”



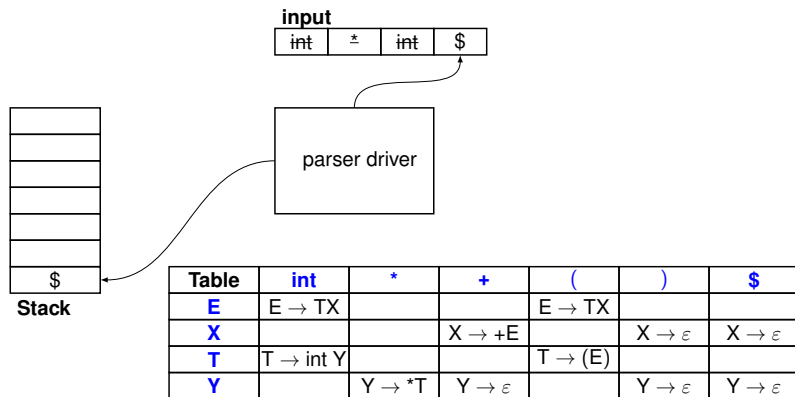
# Using the Parse Table

■ To recognize “int \* int”



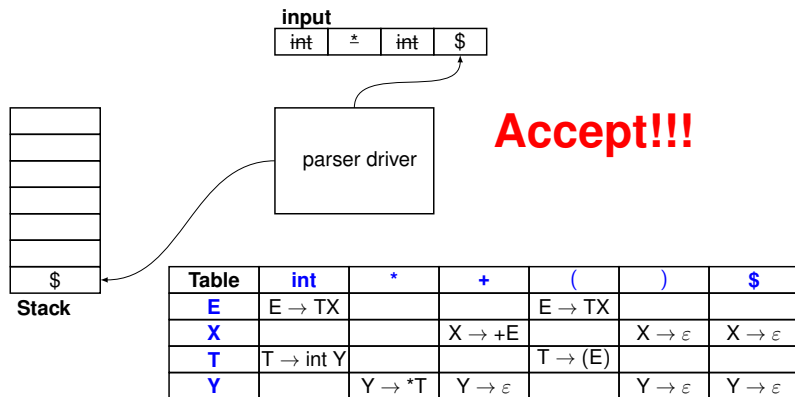
# Using the Parse Table

■ To recognize “int \* int”



# Using the Parse Table

■ To recognize “int \* int”



# Recognition Sequence

It is possible to write in a action list

Stack	Input	Action
E \$	int * int \$	$E \rightarrow TX$
T X \$	int * int \$	$T \rightarrow \text{int } Y$
int Y X \$	int * int \$	terminal
Y X \$	* int \$	$Y \rightarrow * T$
* T X \$	* int \$	terminal
T X \$	int \$	$T \rightarrow \text{int } Y$
int Y X \$	int \$	terminal
Y X \$	\$	$Y \rightarrow \epsilon$
X \$	\$	$X \rightarrow \epsilon$
\$	\$	halt and accept



# First step in building Parse Table: First and Follow Sets

□  $\text{First}(\alpha) = \{t \mid \alpha \Rightarrow *t\beta\}$

➤ Set of terminals that can start a string derived from  $\alpha$ .

□  $\text{Follow}(\alpha) = \{t \mid S \Rightarrow *\alpha t\beta\}$

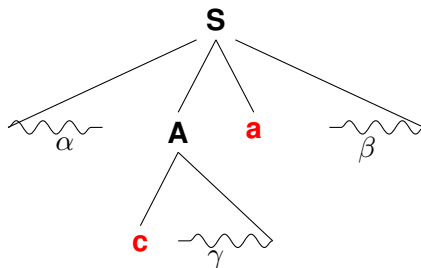
➤ Set of terminals that can follow  $\alpha$  in some derivation.

□ Given rule  $A \rightarrow \alpha$ ,

➤ Choose  $A \rightarrow \alpha$  for all terminals in  $\text{First}(\alpha)$

➤ Choose  $A \rightarrow \alpha$  for all terminals in  $\text{Follow}(A)$ ,  
if and only if  $\alpha \Rightarrow *\epsilon$

# Intuitive Meaning of **First** and **Follow**



$c \in \text{First}(A)$

$a \in \text{Follow}(A)$

 Why is the Follow Set important?

# Calculating $\text{First}(\alpha)$

- Given  $A \rightarrow \alpha$ , let's calculate  $\text{First}(\alpha)$ .
  - $\alpha$  is string  $Y_1 Y_2 Y_3 \dots Y_m$  of terminals and non-terminals.
  - For all  $Y_i$ , if  $Y_i$  is a terminal  $t$ , then  $\text{First}(Y_i) = t$
  - For all non-terminal  $Y_i$ , recursively calculate  $\text{First}(Y_i)$   
(Using below algorithm, replacing  $\alpha$  with  $Y_i$ )
  - Either way, we can assume we know  $\text{First}(Y_i)$  for all  $i$
  
- Apply following rules until no terminal or  $\varepsilon$  can be added
  - 1). Add  $(\text{First}(Y_1) - \varepsilon)$  to  $\text{First}(\alpha)$ .
  - 2). If  $\text{First}(Y_1), \dots, \text{First}(Y_{k-1})$  all contain  $\varepsilon$ ,  
then add  $(\sum_{1 \leq i \leq k} \text{First}(Y_i) - \varepsilon)$  to  $\text{First}(\alpha)$ .
  - 3). If  $\text{First}(Y_1), \dots, \text{First}(Y_m)$  all contain  $\varepsilon$ ,  
then add  $\varepsilon$  to  $\text{First}(\alpha)$ .

# Calculating Follow(A)

□  $\text{Follow}(\alpha) = \{t \mid S \Rightarrow * \alpha t \beta\}$

Intuition: if  $X \rightarrow A B$ , then  $\text{First}(B) \subseteq \text{Follow}(A)$

little trickier because B may be  $\varepsilon$  i.e.  $B \Rightarrow * \varepsilon$

□ Apply following rules until no terminal or  $\varepsilon$  can be added

- 1).  $\$ \in \text{Follow}(S)$ , where S is the start symbol.  
e.g.  $\text{Follow}(E) = \{\$ \dots\}$ .
- 2). Look at the occurrence of a non-terminal on the right hand side of a production which is followed by something  
If  $\dots \rightarrow \alpha A \beta$ , then  $\text{First}(\beta) - \{\varepsilon\} \subseteq \text{Follow}(A)$
- 3). Look at N on the RHS that is not followed by anything,  
if  $(X \rightarrow \alpha A)$  or  $(X \rightarrow \alpha A \beta \text{ and } \varepsilon \in \text{First}(\beta))$ ,  
then  $\text{Follow}(X) \subseteq \text{Follow}(A)$

# Calculating First and Follow Sets for the example

$$\begin{array}{lcl} E & \rightarrow & T X \\ X & \rightarrow & + E \mid \varepsilon \\ T & \rightarrow & \text{int } Y \mid ( E ) \\ Y & \rightarrow & * T \mid \varepsilon \end{array}$$

□ Start by calculating the First Sets for all RHSs

- First( $T X$ )
- First( $+ E$ ), First( $\varepsilon$ )
- First( $\text{int } Y$ ), First( $( E )$ )
- First( $* T$ ), First( $\varepsilon$ )

□ If any of the above First Sets contains  $\varepsilon$ ,  
calculate the Follow Set for corresponding non-terminal

# Calculating First and Follow Sets for the example

$$\begin{aligned}
 E &\rightarrow T X \\
 X &\rightarrow + E \mid \varepsilon \\
 T &\rightarrow \text{int } Y \mid ( E ) \\
 Y &\rightarrow * T \mid \varepsilon
 \end{aligned}$$

Symbol	First
(	(
)	)
+	+
*	*
int	int
E	(, int
X	+, $\varepsilon$
T	(, int
Y	*, $\varepsilon$

RHS	First
T X	(, int
+ E	+
$\varepsilon$	$\varepsilon$
int Y	int
( E )	(
* T	*
$\varepsilon$	$\varepsilon$

# Calculating First and Follow Sets for the example

$$\begin{aligned}
 E &\rightarrow TX \\
 X &\rightarrow +E \mid \epsilon \\
 T &\rightarrow \text{int } Y \mid (E) \\
 Y &\rightarrow *T \mid \epsilon
 \end{aligned}$$

Symbol	First
(	(
)	)
+	+
*	*
int	int
E	(, int
X	+, $\epsilon$
T	(, int
Y	*, $\epsilon$

RHS	First
$TX$	(, int
$+E$	+
$\epsilon$	$\epsilon$
$\text{int } Y$	int
$(E)$	(
$*T$	*
$\epsilon$	$\epsilon$

Non-terminal	Follow
X	\$,)
Y	\$,),+
E	\$,)
T	\$,),+

# Construction of LL(1) Parse Table

- To construct the parse table, we check each  $A \rightarrow \alpha$
- For each terminal  $a \in \text{First}(\alpha)$ , then add  $A \rightarrow \alpha$  to  $M[A, a]$ .
  - If  $\varepsilon \in \text{First}(\alpha)$ , then  
for each terminal  $b \in \text{Follow}(A)$ , add  $A \rightarrow \alpha$  to  $M[A, b]$ .
  - If  $\varepsilon \in \text{First}(\alpha)$  and  $\$ \in \text{Follow}(A)$ , then add  $A \rightarrow \alpha$  to  $M[A, \$]$ .



# Example

$E \rightarrow TX$   
 $X \rightarrow +E \mid \epsilon$   
 $T \rightarrow \text{int } Y \mid (E)$   
 $Y \rightarrow *T \mid \epsilon$

RHS	First
$TX$	(,int
$+E$	+
$\epsilon$	$\epsilon$
$\text{int } Y$	int
$(E)$	(
$*T$	*
$\epsilon$	$\epsilon$

Non-terminal	Follow
$X$	\$,)
$Y$	\$,),+

Table	int	*	+	(	)	\$
<b>E</b>	$E \rightarrow TX$			$E \rightarrow TX$		
<b>X</b>			$X \rightarrow +E$		$X \rightarrow \epsilon$	$X \rightarrow \epsilon$
<b>T</b>	$T \rightarrow \text{int } Y$			$T \rightarrow (E)$		
<b>Y</b>		$Y \rightarrow *T$	$Y \rightarrow \epsilon$		$Y \rightarrow \epsilon$	$Y \rightarrow \epsilon$

# Determine if Grammar G is LL(1)

## Observation

If a grammar is LL(1), then each of its LL(1) table entry contains at most one rule. Otherwise, it is not LL(1).

## Two methods to determine if a grammar is LL(1) or not

(1) Construct LL(1) table, and check if there is a multi-rule entry  
or

(2) Checking each rule as if the table is getting constructed.

G is LL1(1) **iff** for a rule  $A \rightarrow \alpha | \beta$

➤  $\text{First}(\alpha) \cap \text{First}(\beta) = \phi$

➤ at most one of  $\alpha$  and  $\beta$  can derive  $\varepsilon$

➤ If  $\beta$  derives  $\varepsilon$ , then  $\text{First}(\alpha) \cap \text{Follow}(A) = \phi$

# Left-recursion disqualifies grammar for LL(1)

- ❑ Recall recursive descent had trouble with left-recursion.
- ❑ Table-driven parsers have a similar problem.
- ❑ Left-recursion is of the form  $A \rightarrow Ab|a$  or  $A \rightarrow Ab|\varepsilon$ 
  - For  $A \rightarrow Ab|a$ ,  $\text{First}(Ab) \cap \text{First}(a) = \{a\}$
  - For  $A \rightarrow Ab|\varepsilon$ ,  $\text{First}(Ab) \cap \text{Follow}(A) = \{b\}$
  - Either way, an ambiguity in prediction
- ❑ Even if prediction can be made with more lookahead,
  - Sentence can grow indefinitely w/o consuming input
  - We may repeatedly choose to apply  $A \rightarrow Ab$ :  
 $A \Rightarrow A b \Rightarrow A b b \dots$
  - Same stack explosion problem as with recursive descent

# Dealing with Non-LL(1) Grammars

(1) Likely still an LL(1) language. Massage to LL(1) grammar:

- Apply left-factoring
- Apply left-recursion removal

(2) If (1) fails, the possibilities are...

- Grammar just needs a little more lookahead  
(May need LL(k) parser where  $k > 1$  or backtracking)
- Grammar is ambiguous (multiple parse trees)

□ How do we deal with ambiguous grammars then?

- Note: left-factoring and left-recursion removal don't help
- Expressing precedence and associativity in grammar helps

# Ambiguous not just non-LL(1)

- Some grammars are not LL(1) even after left-factoring and left-recursion removal

$S \rightarrow \text{if } C \text{ then } S \mid \text{if } C \text{ then } S \text{ else } S \mid a \text{ (other statements)}$   
 $C \rightarrow b$

change to

$S \rightarrow \text{if } C \text{ then } S X \mid a$   
 $X \rightarrow \text{else } S \mid \epsilon$   
 $C \rightarrow b$

**problem sentence: “if b then if b then a else a”**

$\text{First}(X) = \{\text{else}, \epsilon\}$

From  $S \rightarrow \text{if } C \text{ then } S X$ ,  $\text{Follow}(S) \subseteq \text{Follow}(X)$

$\text{Follow}(X) = \{\text{else}, \$\}$

For  $X \rightarrow \text{else } S \mid \epsilon$ ,  $\text{First}(\text{else } S) \cap \text{Follow}(X) = \{\text{else}\}$

- Such grammars are potentially ambiguous

# Removing Ambiguity

- ❑ We want to express precedence of if-then-else over if-then.
- ❑ How would you rewrite grammar to express precedence?

$S \rightarrow \text{if } C \text{ then } S \mid S_2$

$S_2 \rightarrow \text{if } C \text{ then } S_2 \text{ else } S \mid a$

$C \rightarrow b$

- ❑ Now grammar is unambiguous but it is not LL(k) for any k
  - Intuitively, must lookahead until 'else' to choose rule for 'S'
  - That lookahead may be an arbitrary number of tokens
- ❑ Changing the grammar to be perfectly unambiguous
  - Can be very taxing for programmers to specify correctly
  - May still result in grammar not suitable for LL(1) parsing
- ❑ More practical to encode precedence rules into parser
  - E.g. Always choose  $X \rightarrow \text{else } S$  over  $X \rightarrow \varepsilon$  on 'else' token

# LL(1) Time and Space Complexity

- ❑ LL(1) parsers operate in linear time and space relative to the length of input.
- ❑ Time: each token is processed constant number of times
  - Why?
- ❑ Space: stack space required is at max the length of input
  - If  $X \rightarrow \varepsilon$  rules removed (easily done by substitution)
  - Why?
- ❑ How about LL(k)?
  - Same time complexity as the same argument applies
  - Space complexity is  $O(T^k)$ , where  $T$  is number of terminals (if constructing the parse table naively)

# ANTLR: A modern LL(k) parser

- ❑ A free open source top-down LL(k) parser (antlr.org)
  - Used in Apache Groovy, Jython, MySQL Workbench, ...
  
- ❑ Reduces table space by expressing lookahead as DFA
  - A DFA decides on which rule for each non-terminal
  - DFA can express arbitrarily long lookahead compactly
  
- ❑ If you are interested, refer to this paper:  
LL(\*): The Foundation of the ANTLR Parser Generator  
<https://www.antlr.org/papers/LL-star-PLDI11.pdf>