Nonlinear Data Structure

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Office Hours:

Non Linear Data Structure

- Data structure we will consider this semister:
 - ◆ Tree(트리)
 - ◆ Binary Search Tree(이진탐색트리)
 - ◆ Graph(□래프)
- ➡ ◆ Weighted Graph(가중치그래프)
 - ◆ Sorting(정렬)
 - ◆ Balanced Search Tree(균형탐색트리)



가중치 그래프 Weighted Graph



- ◆ Minimum Spanning Tree(최소비용간선트리)
- ◆ Shortest path(최단경로)
- ◆ Topological order(토폴로지컬(위상) 순서)
- ◆ Critical Path(임계경로)

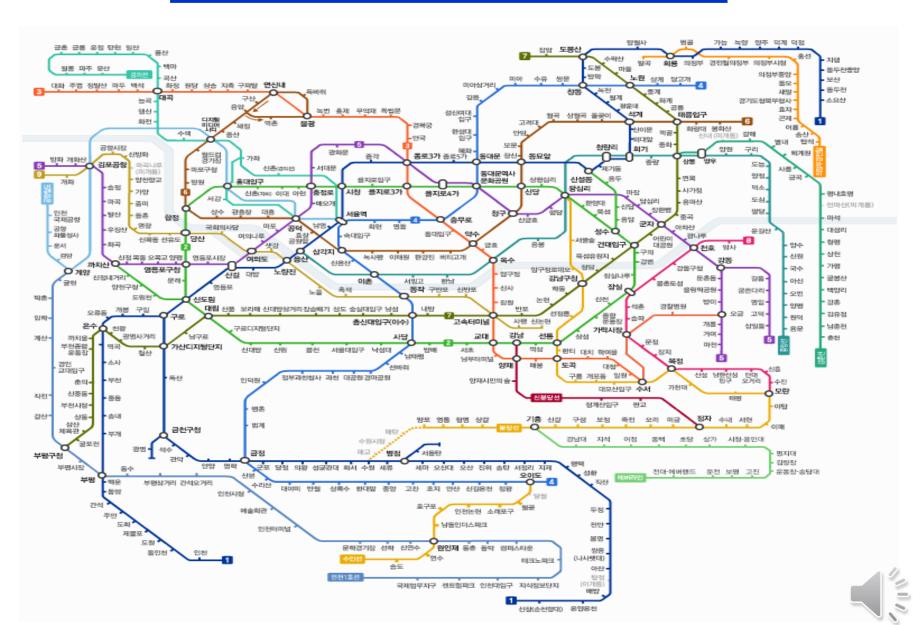


Shortest Path(최단경로)

_컴퓨터공학에서 가장 많이 사용되는 단한 가지

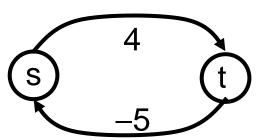
알고리즘을 고른다면 최단경로 찾는 알고리즘일 것임.

수도권 지하철 노선도



Shortest Paths Algorithms(최단경로알고리즘)

- Dijkstra's algorithm
- Bellman & Ford algorithm
- ◆ All Pairs Shortest Path Floyd-Washall algorithm
- ◆ 가중치의 합이 음이 존재하는 사이클이 있는 그래프의 경우, 어떠한 알고리즘도 최단 경로를 해결할 수 없다.





Single-Source Shortest Paths Dijkstra's algorithm

- Given: A single source vertex in a weighted, directed graph.
- Want to compute a shortest path for each possible destination.
 - Similar to BFS.
- We will assume either
 - no negative-weight edges
- Algorithm will compute a shortest-path tree.
 - Similar to BFS tree.



Dijkstra's Algorithm for Shortest Paths Algorithms(다익스트라의 최단경로알고리즘)

- ▶ Non-negative edge weight(음의 가중치를 허용하지 않음. 음의 가중치가 있는 그래프에서 잘못된 결과발생)
- Single-source shortest-paths problem: Find the shortest path from s to each vertex v. (하나의 출발점에서 다른 모든 정점까지의 최단 경로) 알고리즘은 트리모양의 출력을 제공한다.
- ◆ Like BFS: If all edge weights are equal, then use BFS, otherwise use this algorithm (너비우선탐색과 유사)
- ◆ Use Q = priority queue keyed on d[v] values(우선순위큐사용)(note: BFS uses FIFO)
- ◆ 다익스트라의 알고리즘은 greedy algorithm(탐욕적 알고리즘)



Important Aside: Greedy Algorithms

- ◆ A *greedy* algorithm always takes the best immediate or local solution while finding an answer. 탐욕 알고리즘은 특정단계에서 최적으로 보이는 답을 선택
- ◆ Greedy algorithms find optimal solutions for some optimization problems, but may find (far) less-than-optimal solutions for other optimization problems. 그러나 나중에 결과를 보면 최적의답이 아닐수있다.
 - Dijkstra's algorithm is greedy.
- ◆ When greedy algorithms work, they are usually best. 다이스트라의 알고리즘은 다행히 탐욕알고리즘이지만 최적의 해를 제공한다.



Dijkstra's Algorithm

```
Dijkstra(G)
   for each v \in V
      d[v] = \infty;
   d[s] = 0; S = \emptyset; Q = V;
   while (Q \neq \emptyset)
      u = ExtractMin(Q); //Q에서가장작은정점꺼냄
      S = S U{u};//결정된집합에 추가
      for each v \in u-Adj[]
          if (d[v] > d[u]+w(u,v))//경로 재계산
             d[v] = d[u] + w(u,v);
```



Dijkstra's Algorithm

Assumes no negative-weight edges.(음의 가중치가 없다고 가정)

Maintains a set S of vertices whose SP(shortest paths) from s has been determined.

Repeatedly selects u in V–S with minimum SP estimate (greedy choice).

Store V–S in priority queue Q.

```
Initialize(G, s);
S := \emptyset;
Q := V[G];
while Q \neq \emptyset do
u := Extract-Min(Q);
S := S \cup \{u\};
for each v \in Adi[u] do
    Relax(u, v, w)
od
od
```

Dijkstra's Algorithm

- Set all distances initially to ∞ , except the start node, which should be set to 0
- ◆ Construct a min priority
 queue(최소우선순위큐) of the nodes, with
 their distances as keys
- Repeatedly remove the minimum element, updating each of its adjacent node's distances if they are still in the queue and if the updated distance is less than the current distance

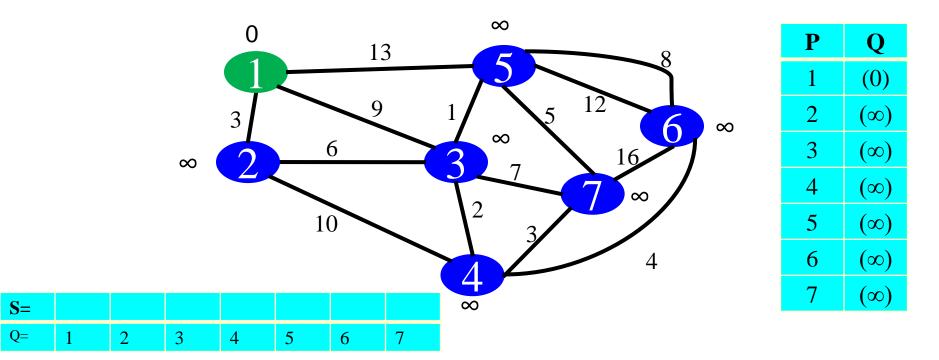


```
shortestPath(v, weight, n)
  // v는 시작점, weight는 가중치 인접 행렬, n은 정점수
  // create S[n], Dist[n]
 for (i\leftarrow 0; i< n; i\leftarrow i+1) do {
     S[i] ← false; // S를 초기화
     Dist[i] ← weight[v, i]; // Dist를 초기화
 S[v] \leftarrow true;
 Dist[v] \leftarrow 0:
 for (i←0; i<n-2; i←i+1) do { // n-2번 반복
                        // 새로운 최단 경로를 선정
     select u such that
        Dist[u] = min\{Dist[j] \mid S[j] = false and 0 \le j \le n\};
     S[u] \leftarrow true;
     for (w←0; w<n; w←w+1) do { // 확정이 안된 경로들에 대해 다시 계산
        if (S[w] = false) then {
            if (Dist[w] > (Dist[u] + weight[u, w])
                  then Dist[w] \leftarrow Dist[u] + weight[u, w];
  }}}
end shortestPath
```

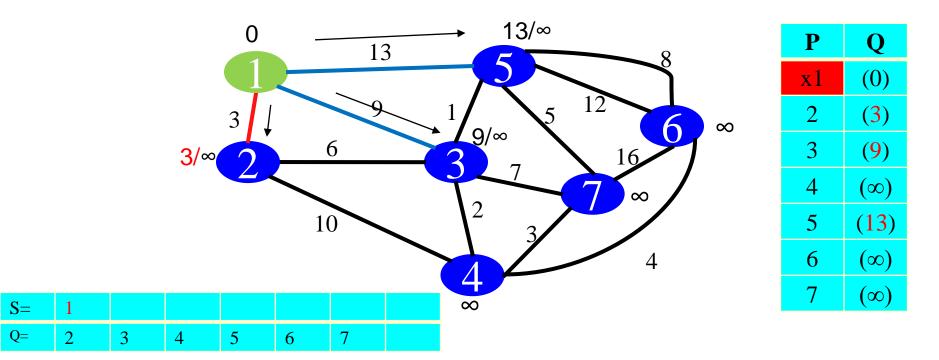
시간복잡도는 $O(E \lg V + E) V : 정점의 개수 E : 간선의 개수 = O(V^2)$

Dijkstra's Algorithm 다익스트라의 알고리즘 예

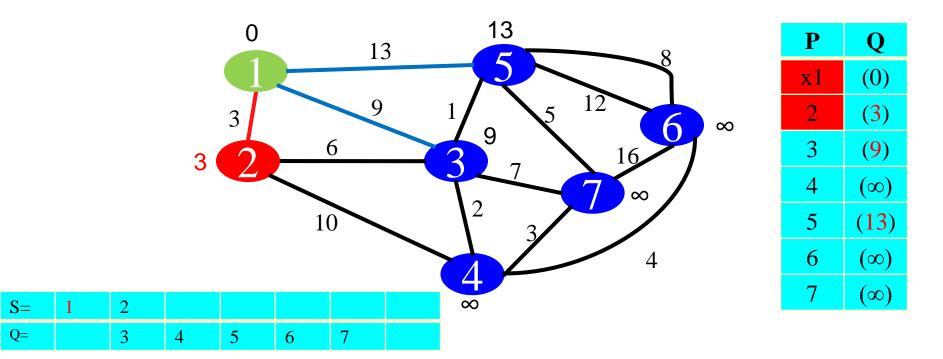




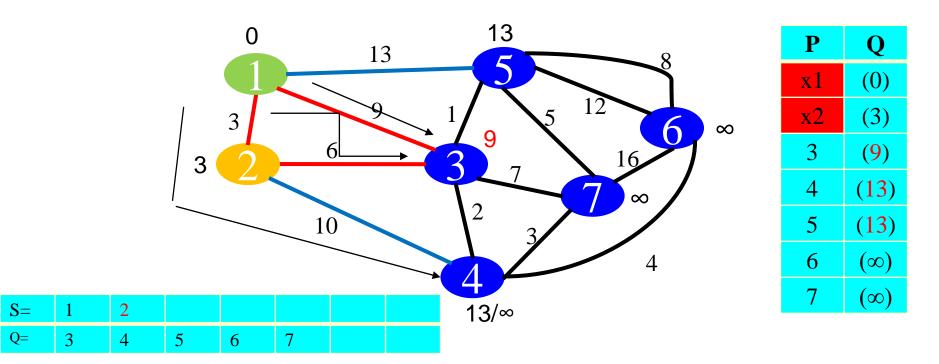




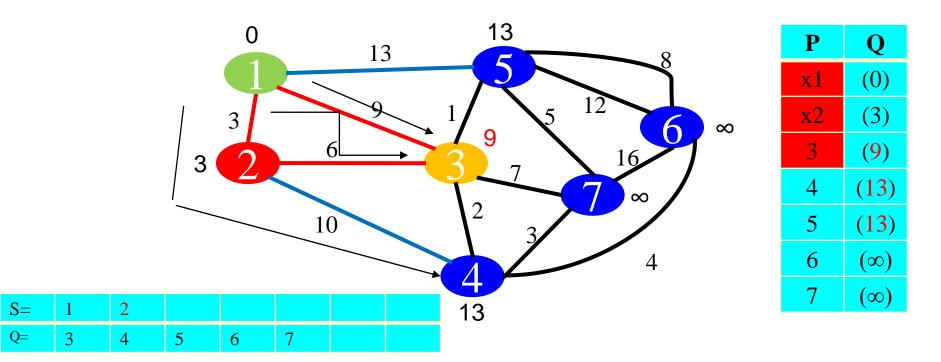




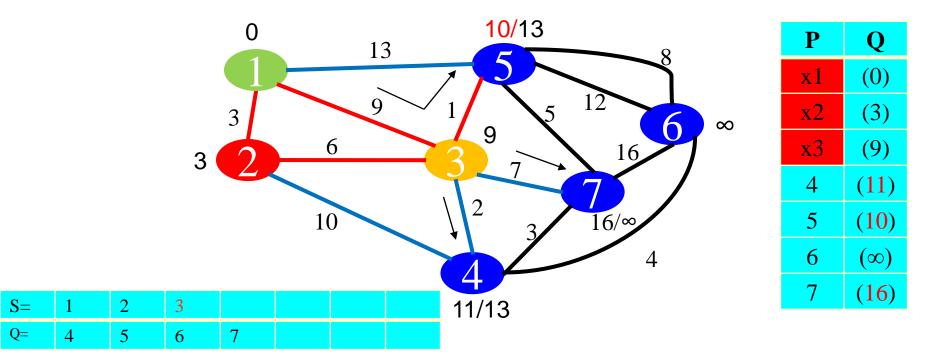




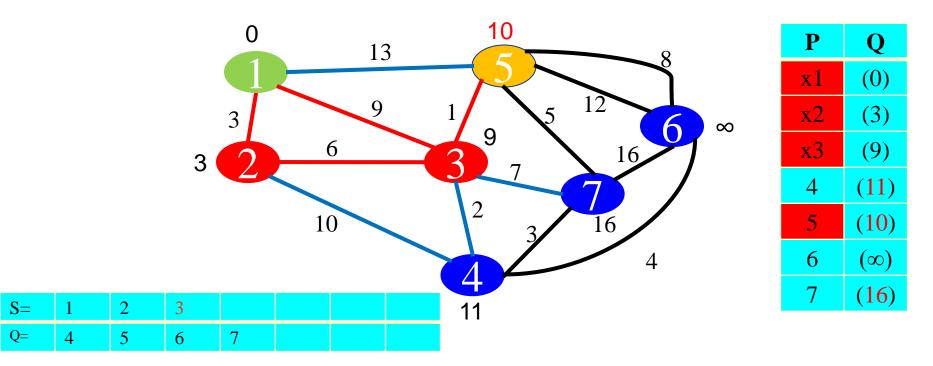




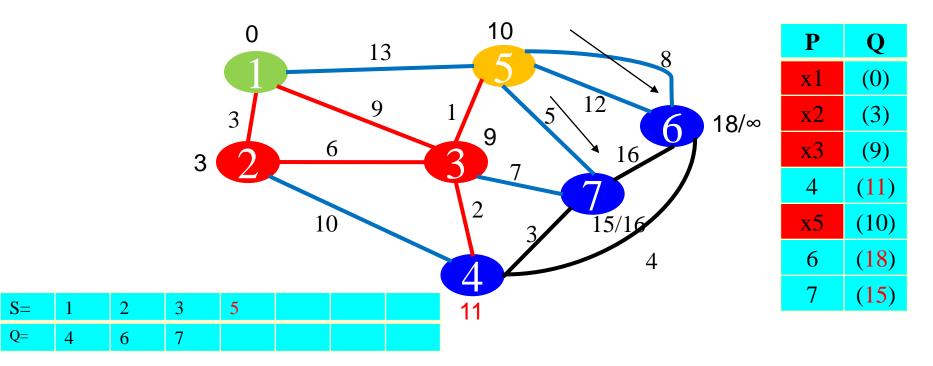




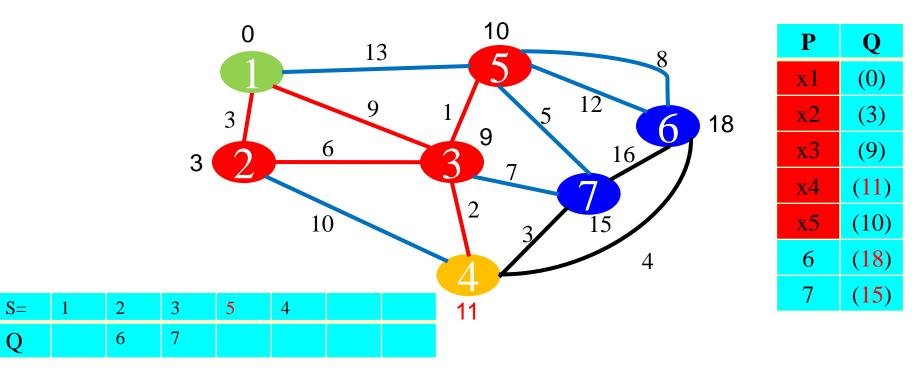




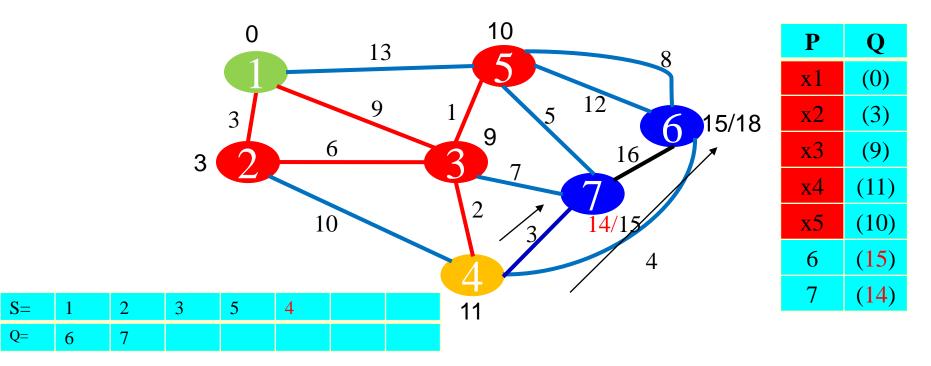




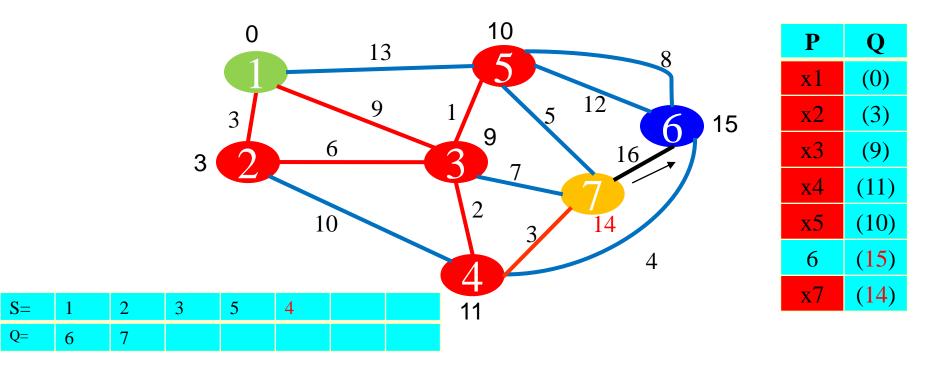




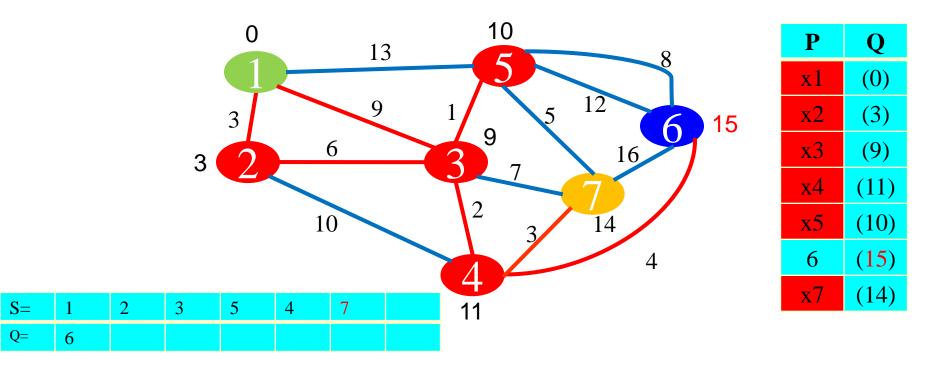




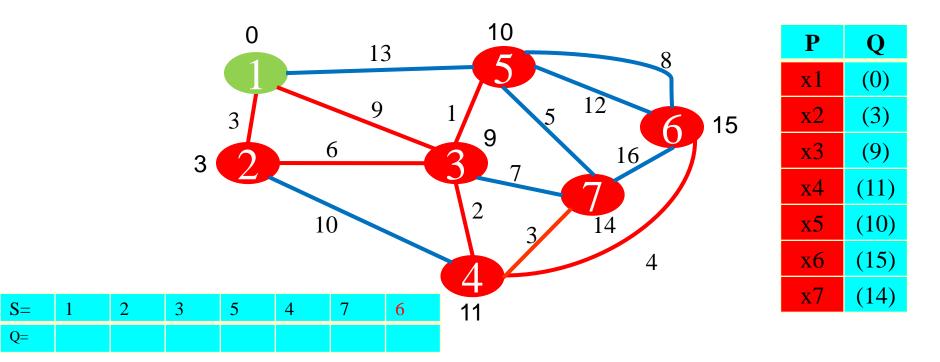




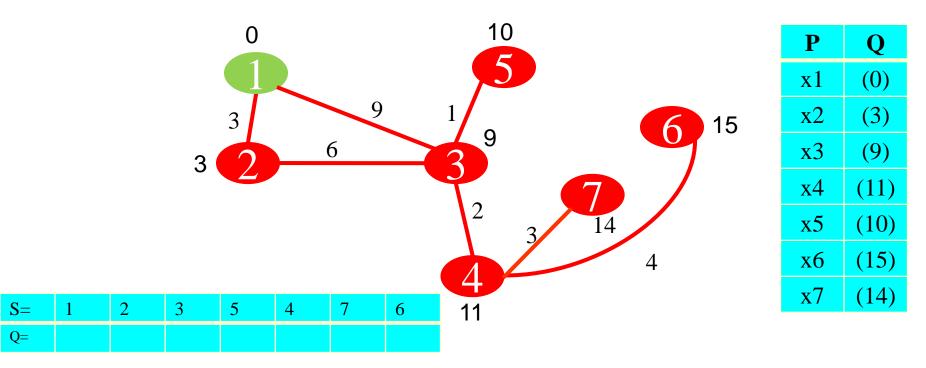








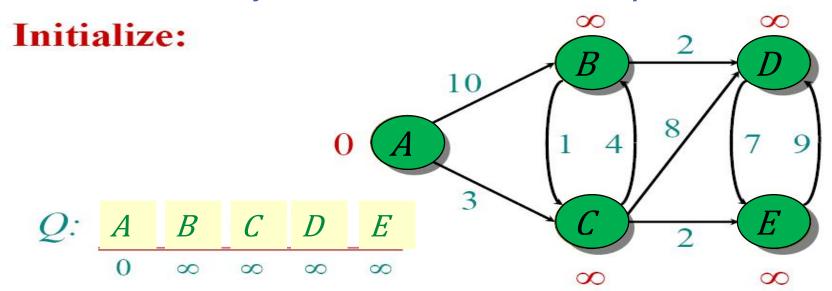




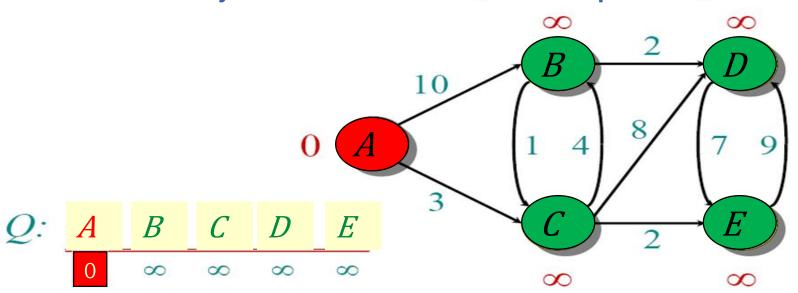
Spanning Tree (간선트리)

Dijkstra's Algorithm 다익스트라의 알고리즘 예2

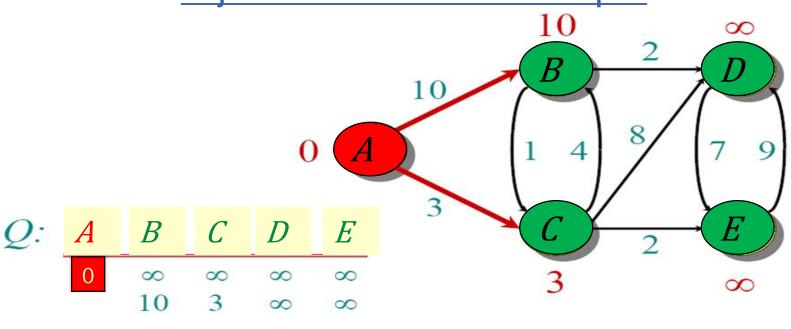




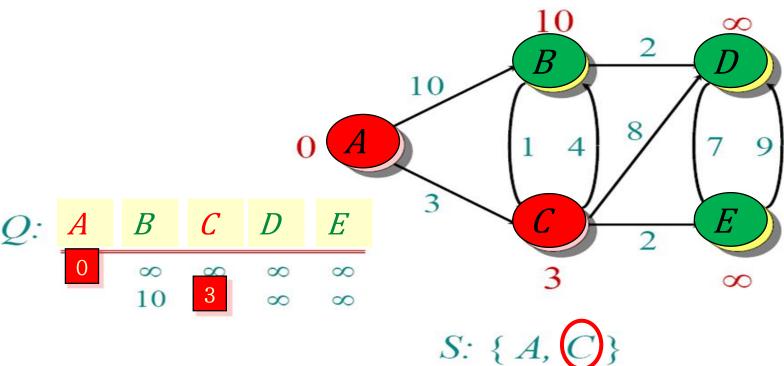






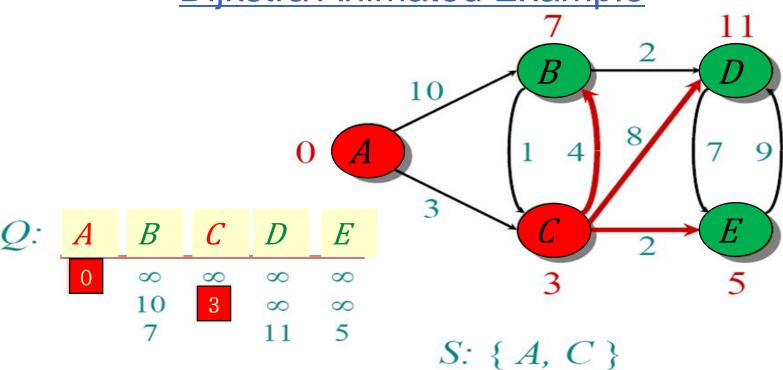






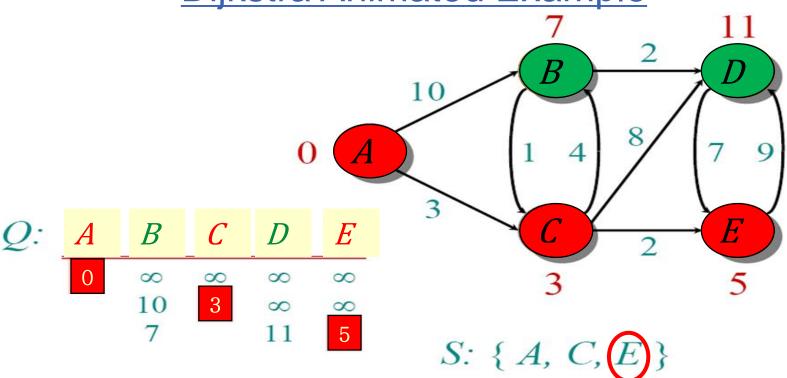




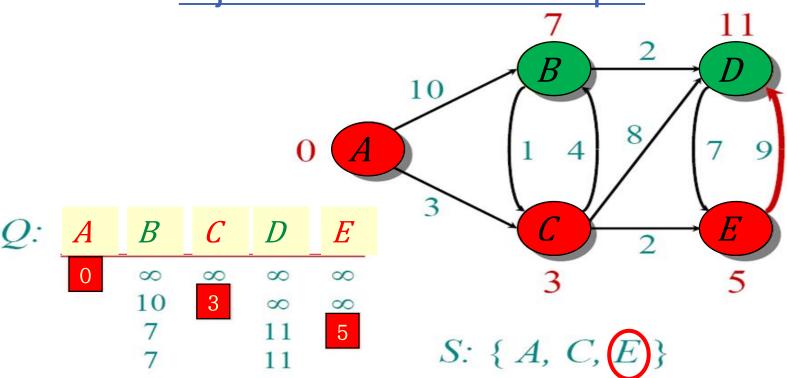


$$S: \{A, C\}$$

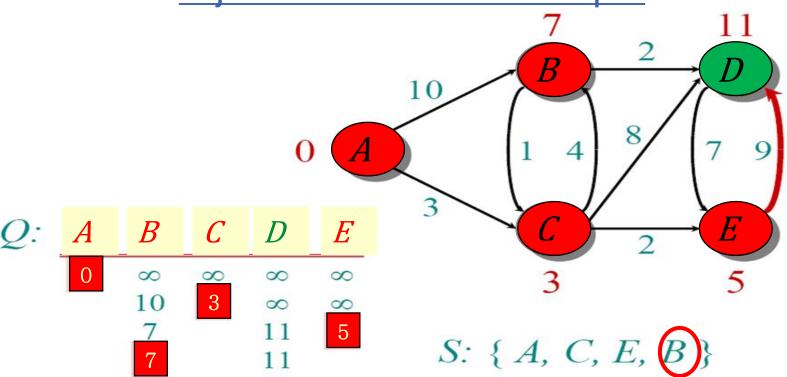




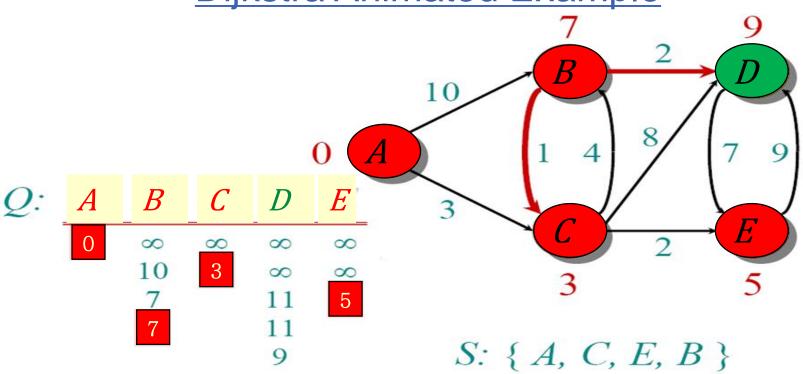




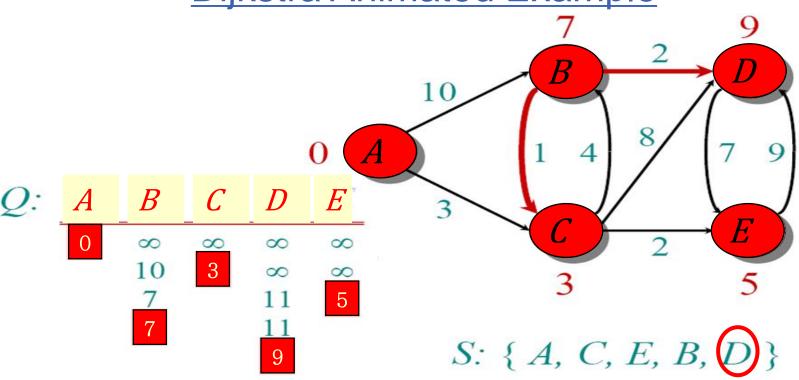








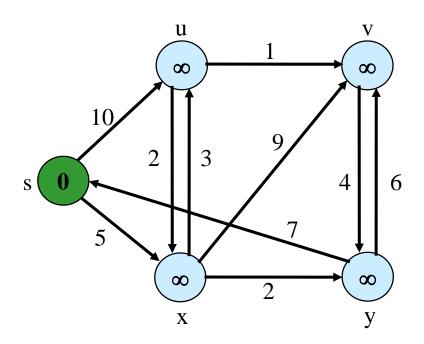






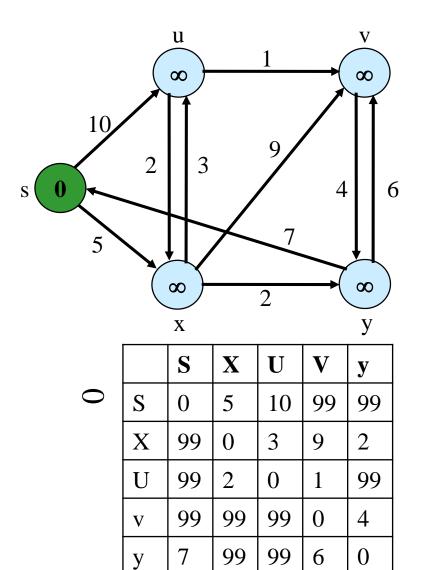
Dijkstra's Algorithm 다익스트라의 알고리즘 예3





 $S: \{S\}$

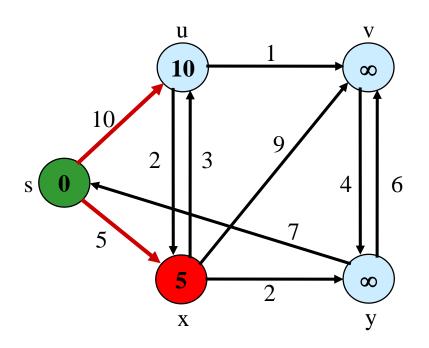


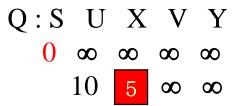


 $Q: S \quad U \quad X \quad V \quad Y$ $0 \quad \infty \quad \infty \quad \infty \quad \infty$

99:infinite

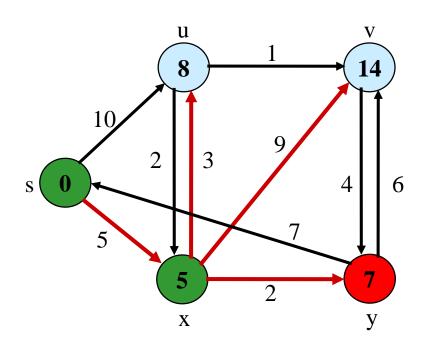


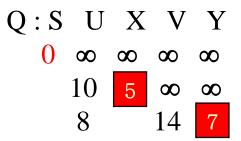




$$S: \{S, X\}$$

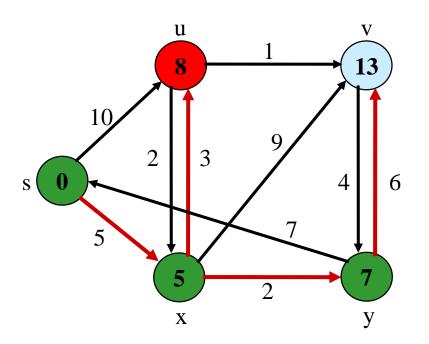






$$S: \{ S, X, \frac{Y}{Y} \}$$

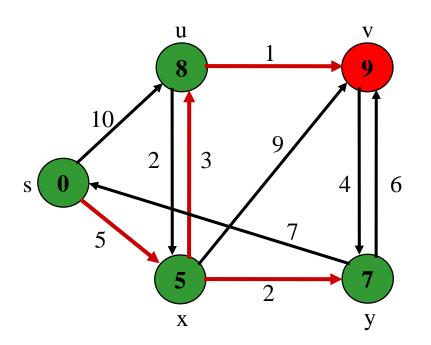


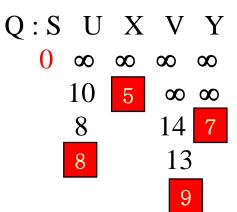


$$Q: S \ U \ X \ V \ Y \\ 0 \ \infty \ \infty \ \infty \ \infty \\ 10 \ 5 \ \infty \ \infty \\ 8 \ 14 \ 7 \\ 8 \ 13$$

$$S: \{ S, X, Y, U \}$$

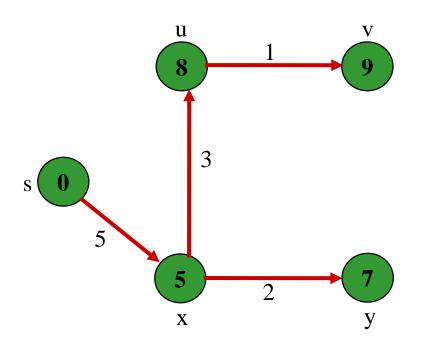






$$S: \{ S, X, Y, U, V \}$$



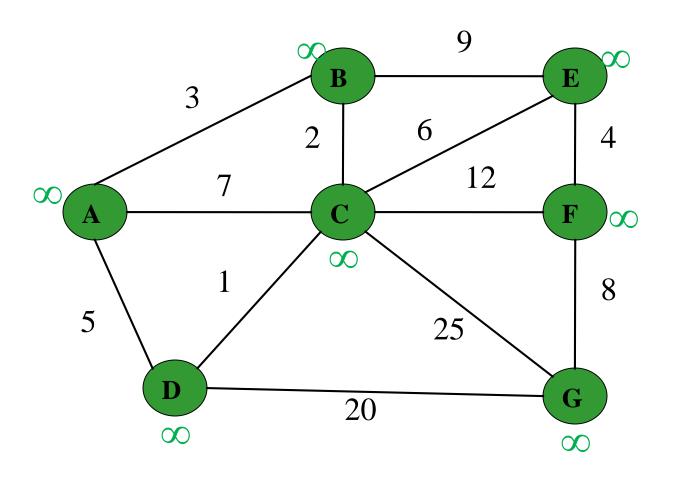


 $S: \{ S, X, Y, U, V \}$

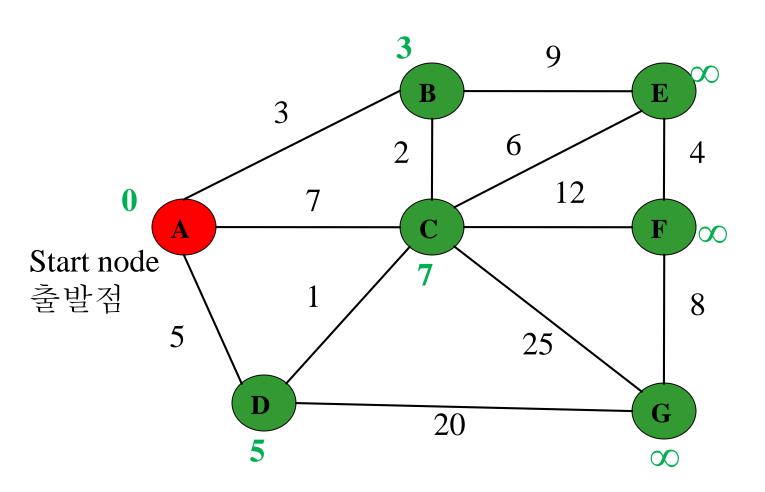


Dijkstra's Algorithm 다익스트라의 알고리즘 예4

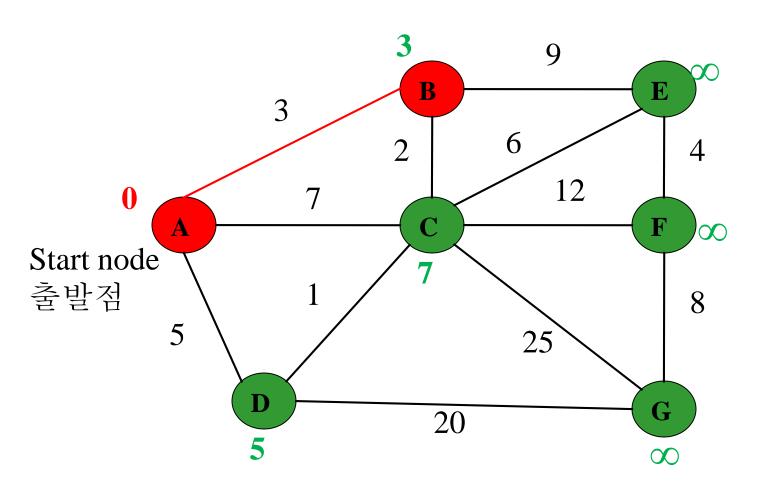




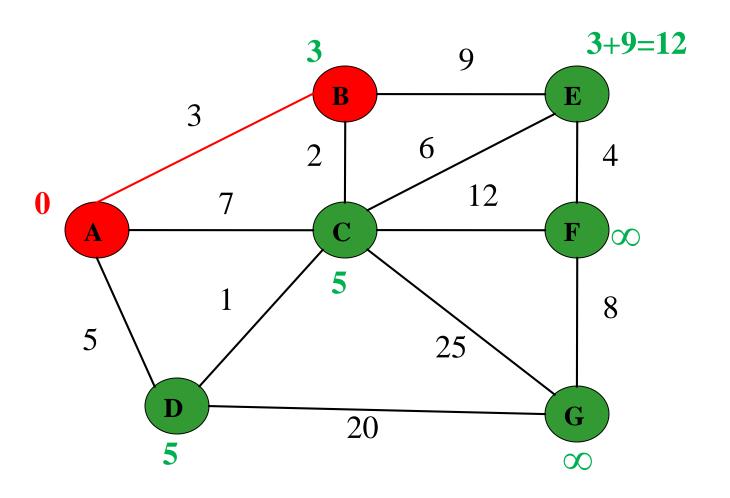




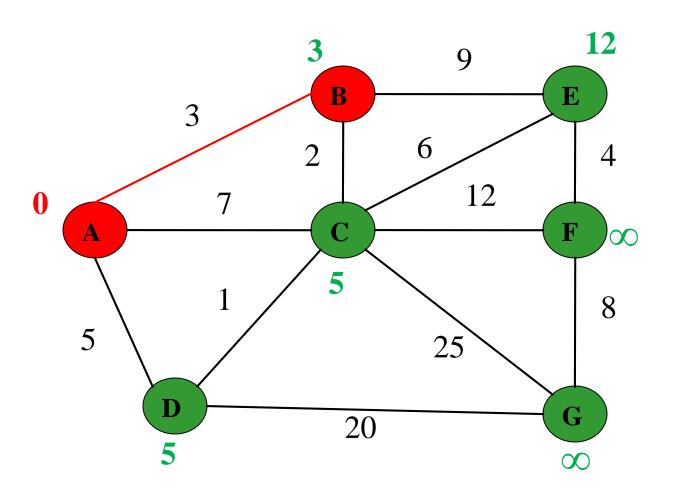




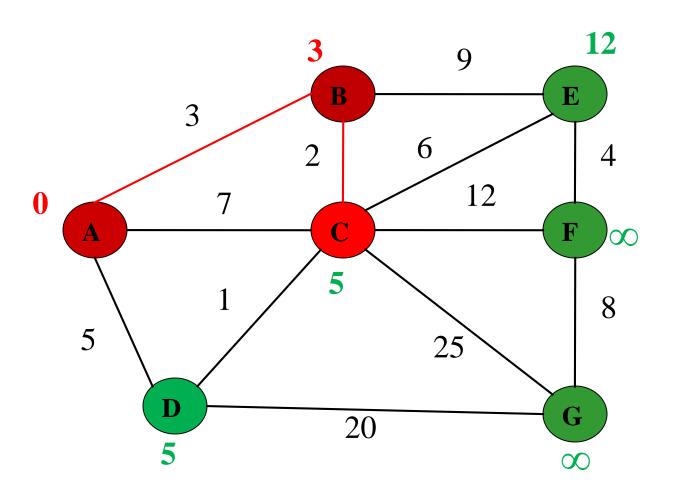




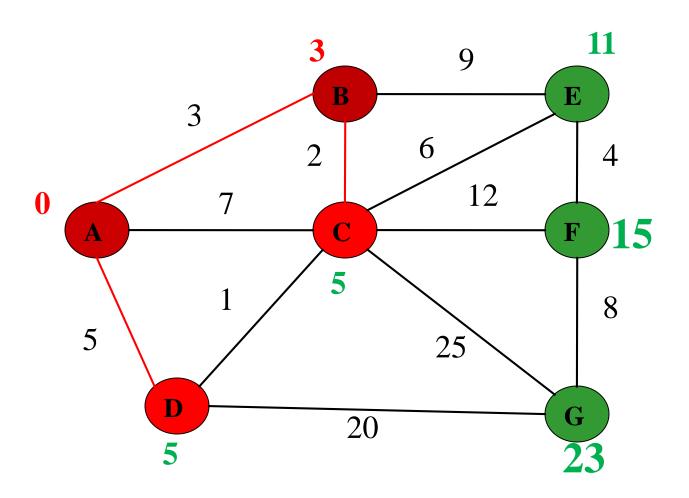




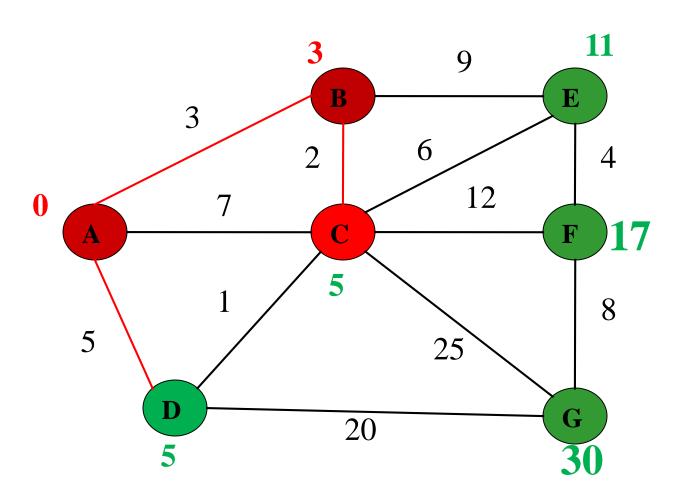




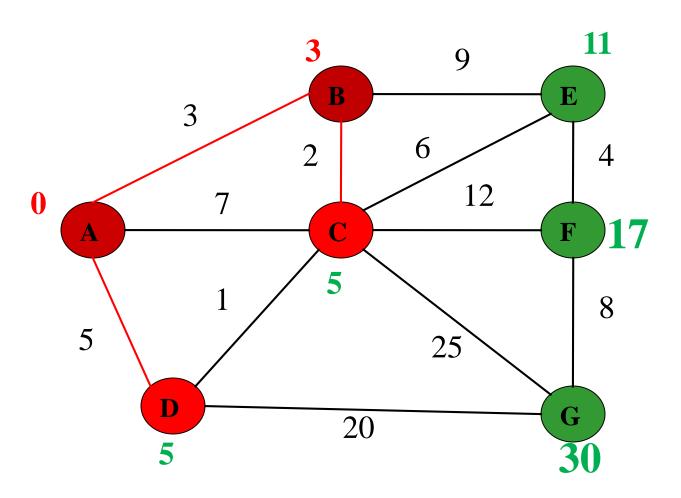




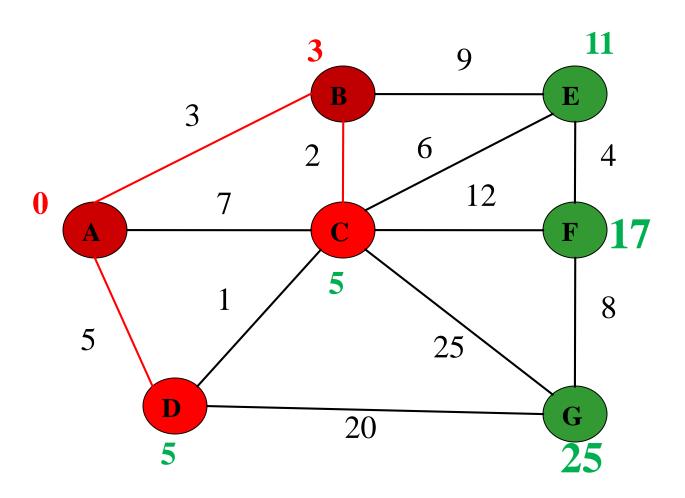




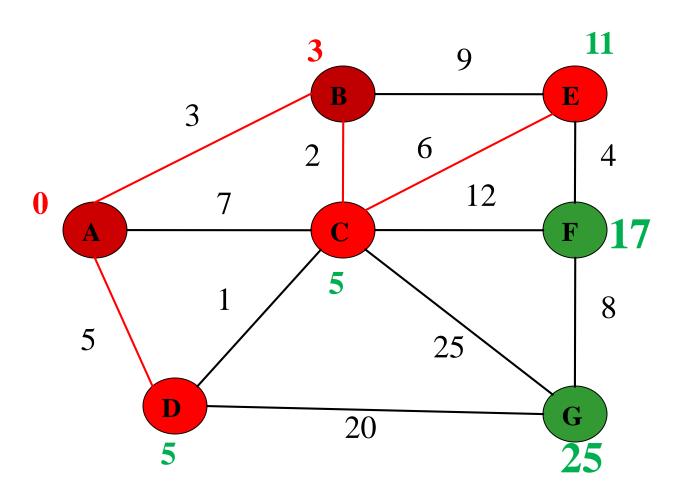




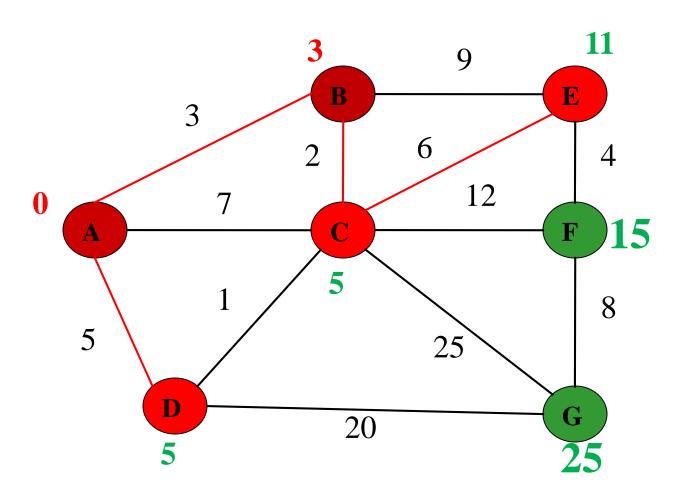




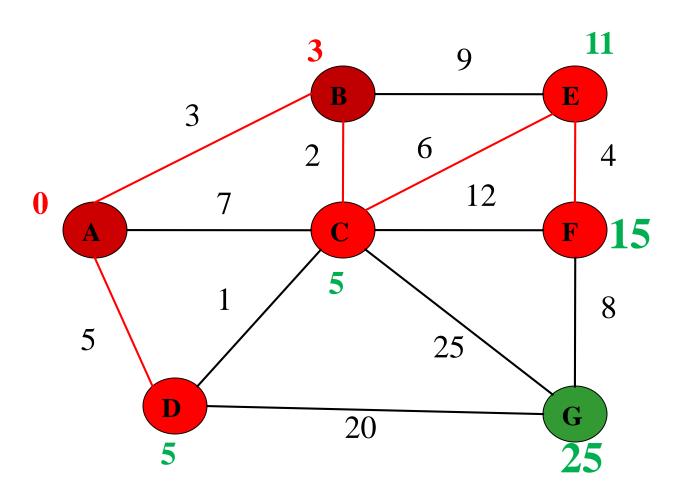




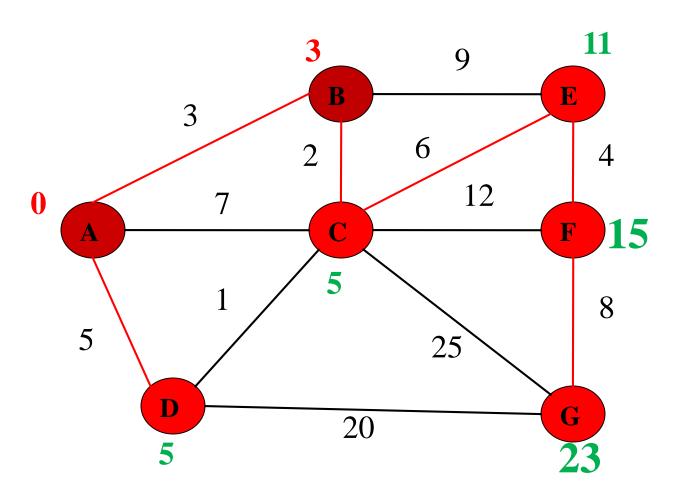




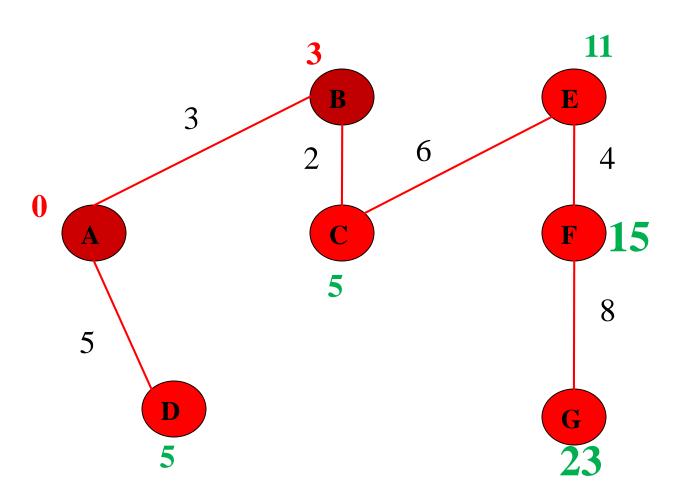












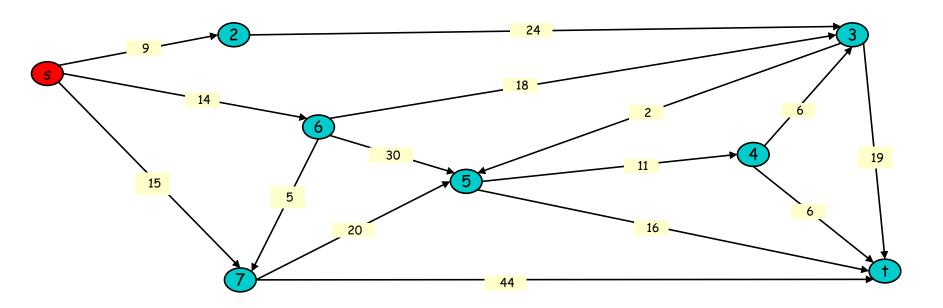
Spanning Tree (간선트리)



Dijkstra's Algorithm 다익스트라의 알고리즘 예5



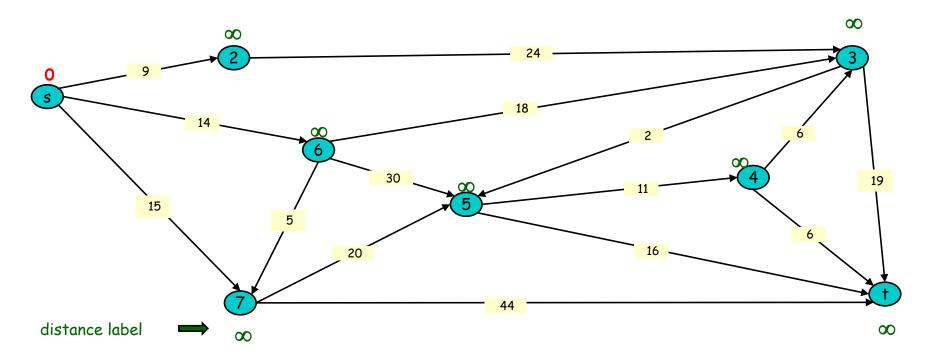
- Find shortest path from s to t.
- ◆ BFS(너비우선탐색)





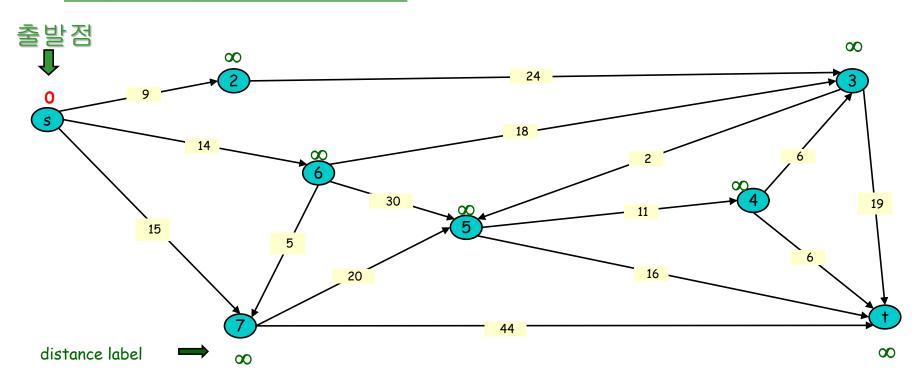
S	2	3	4	5	6	7	t
0	œ	∞	œ	œ	œ	œ	∞

S = { } 해답집합 PQ = { s, 2, 3, 4, 5, 6, 7, † } 우선순위큐



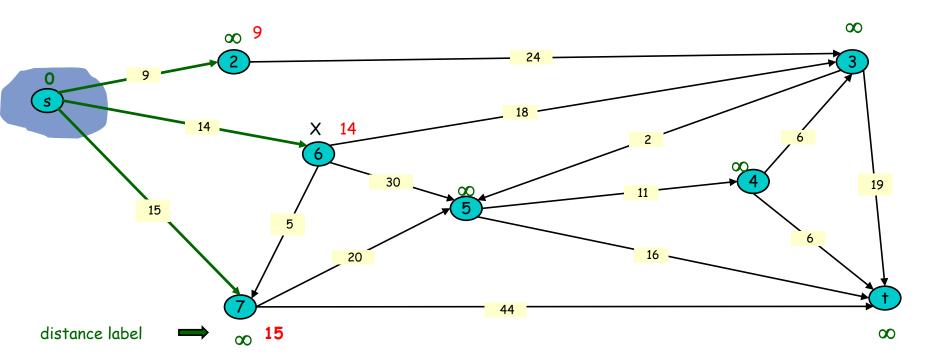


S	2	3	4	5	6	7	t
0	œ	œ	œ	œ	œ	œ	∞



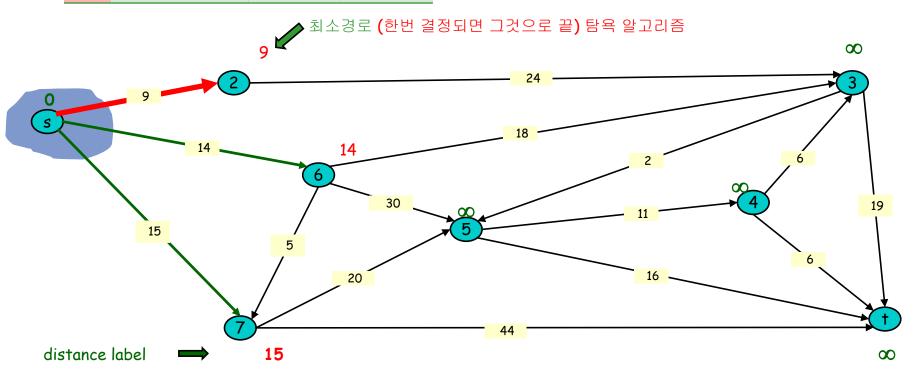


S	2	3	4	5	6	7	t
0	9	œ	œ	œ	14	15	œ



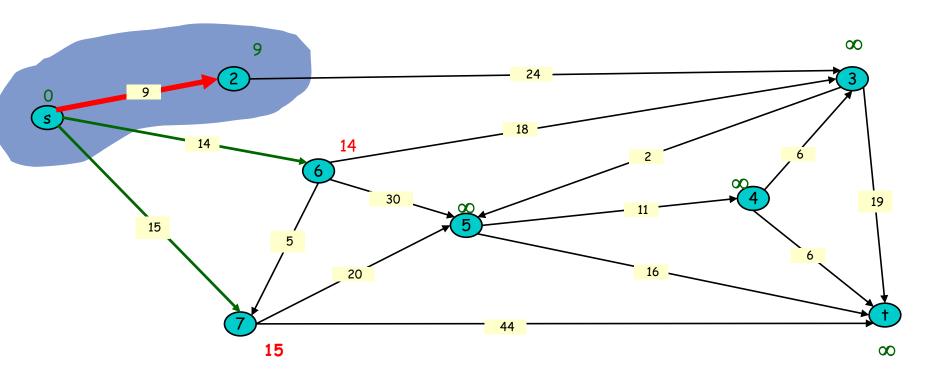


S	2	3	4	5	6	7	t
0	9	œ	∞	œ	14	15	∞



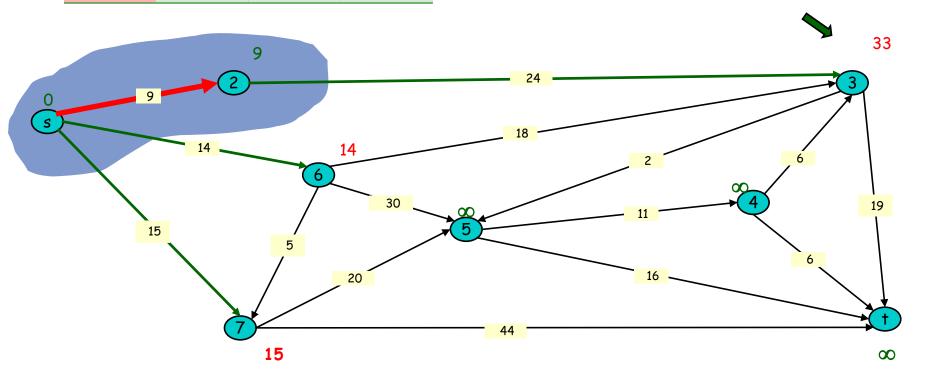


S	2	3	4	5	6	7	t
0	9	∞	∞	œ	14	15	∞



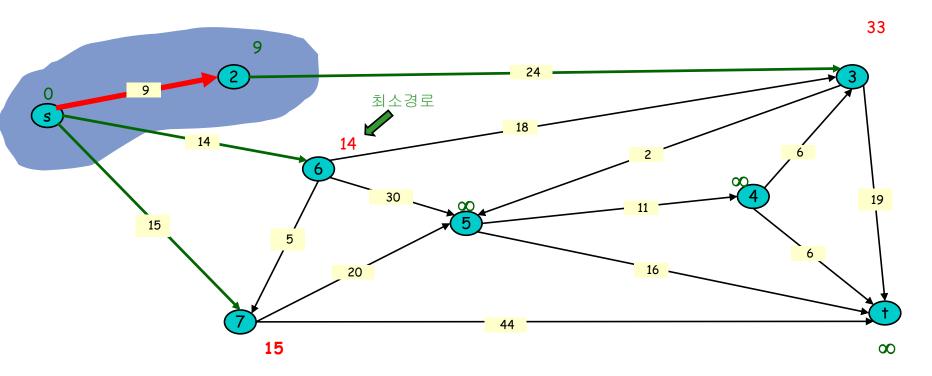


S	2	3	4	5	6	7	t
0	9	33	œ	œ	14	15	00



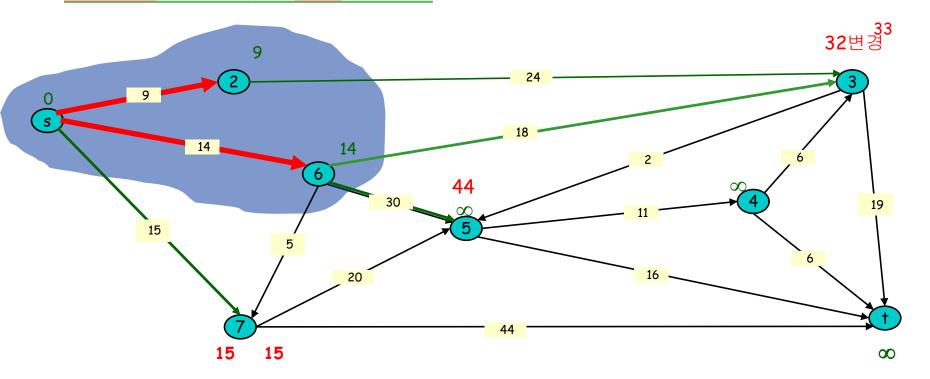


S	2	3	4	5	6	7	t
0	9	33	∞	œ	14	15	∞



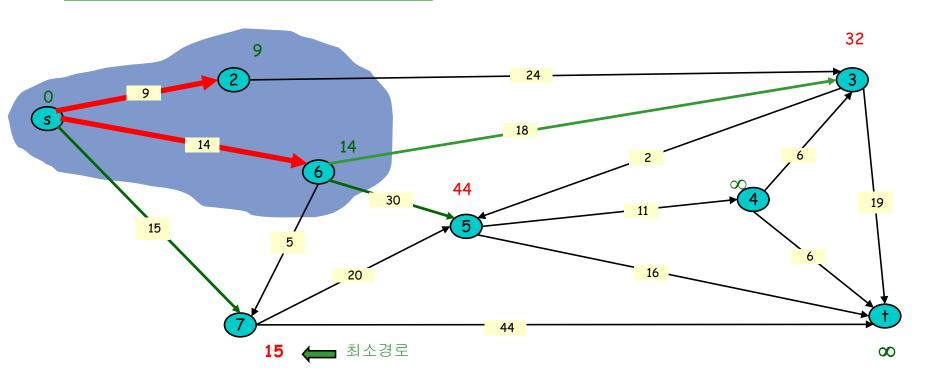


S	2	3	4	5	6	7	t
0	9	32	∞	44	14	15	00





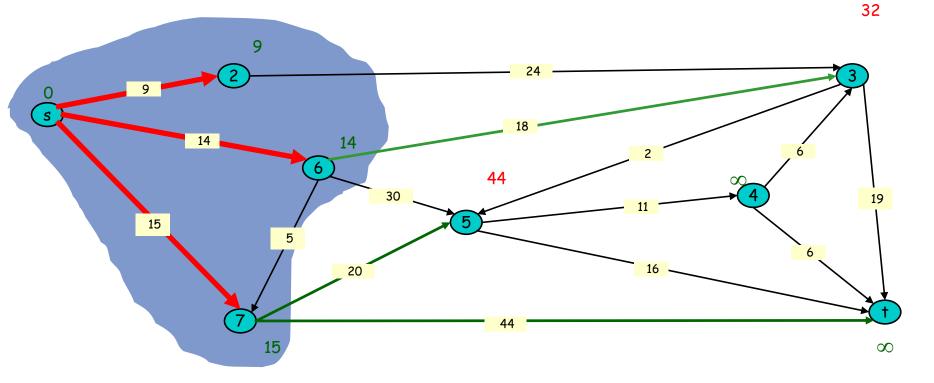
S	2	3	4	5	6	7	t
0	9	32	œ	44	14	15	∞



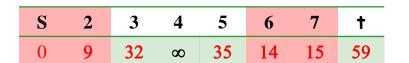


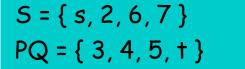
S	2	3	4	5	6	7	t
0	9	32	∞	35	14	15	59

S = { s, 2, 6, 7 } PQ = { 3, 4, 5, † }

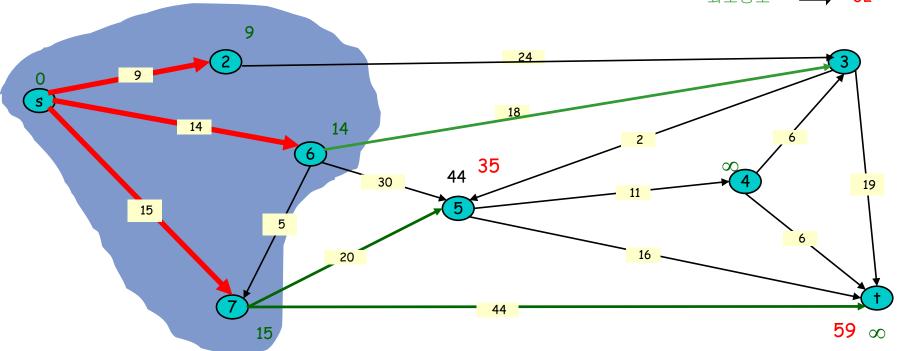




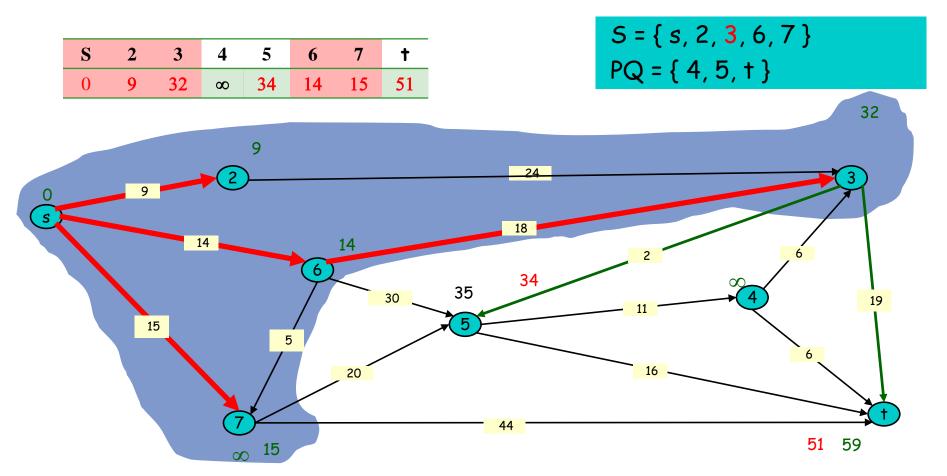




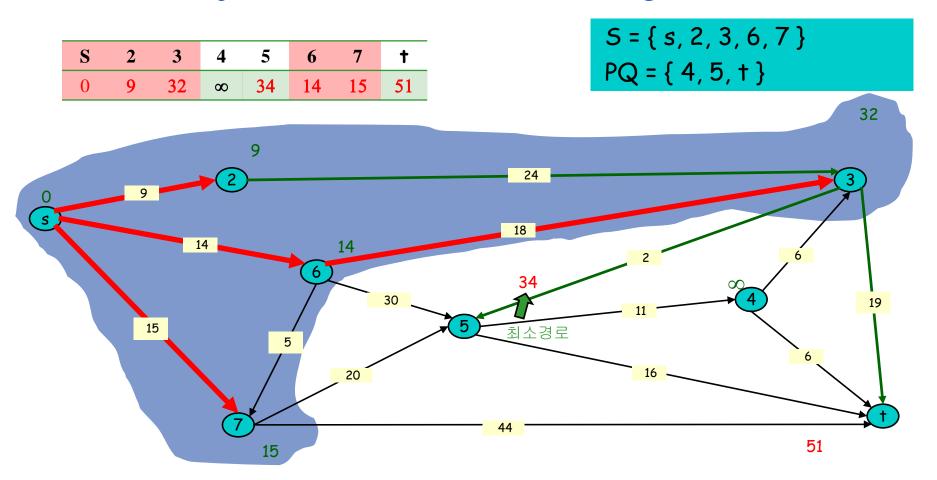
최소경로 🗪 37







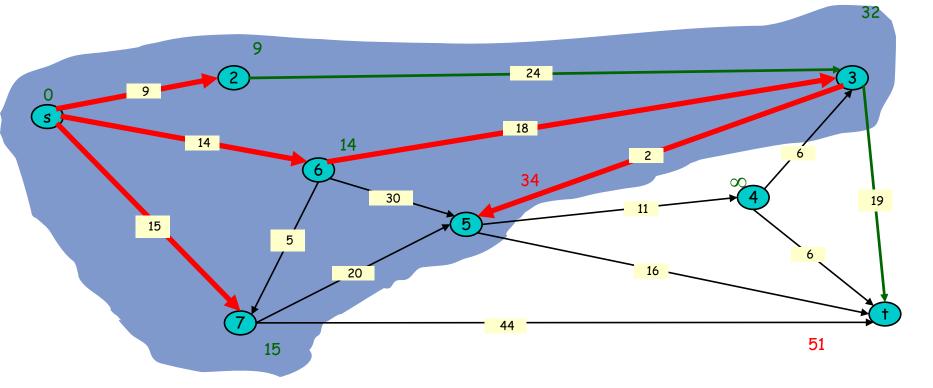






S	2	3	4	5	6	7	t
0	9	32	œ	34	14	15	51

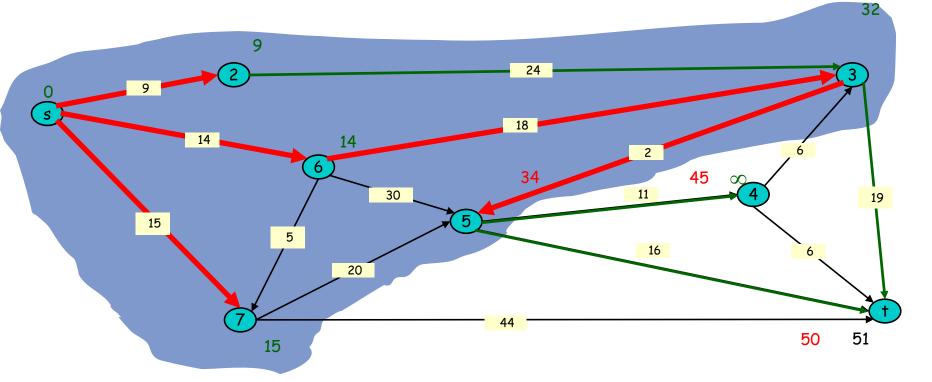
 $S = \{s, 2, 3, 5, 6, 7\}$ PQ = $\{4, +\}$





S	2	3	4	5	6	7	†
0	9	32	œ	34	14	15	51

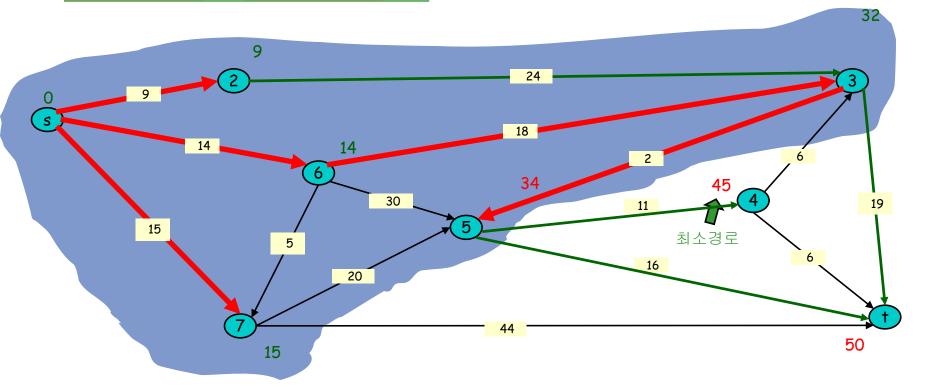
 $S = \{s, 2, 3, 5, 6, 7\}$ PQ = $\{4, +\}$





S	2	3	4	5	6	7	t
0	9	32	œ	34	14	15	51

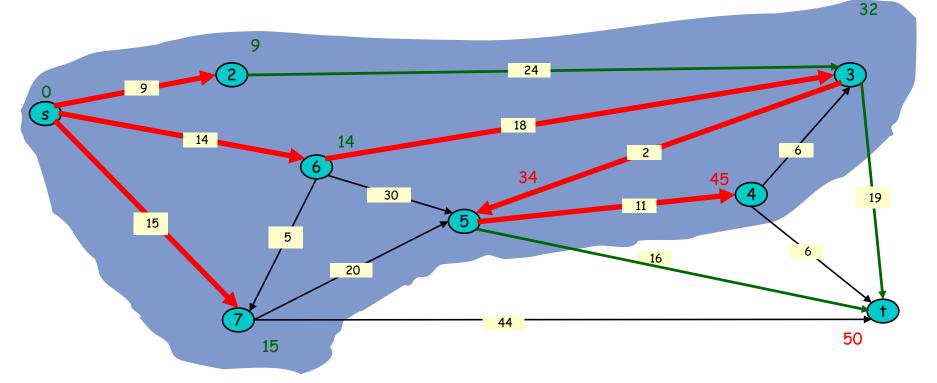
 $S = \{s, 2, 3, 5, 6, 7\}$ PQ = $\{4, +\}$



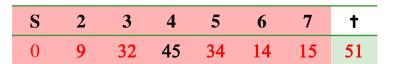


S	2	3	4	5	6	7	t
0	9	32	45	34	14	15	51

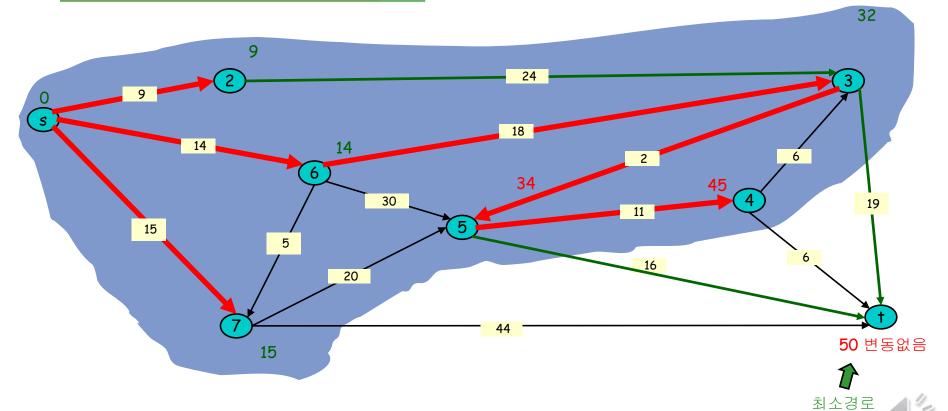
 $S = \{ s, 2, 3, 4, 5, 6, 7 \}$ PQ = $\{ t \}$





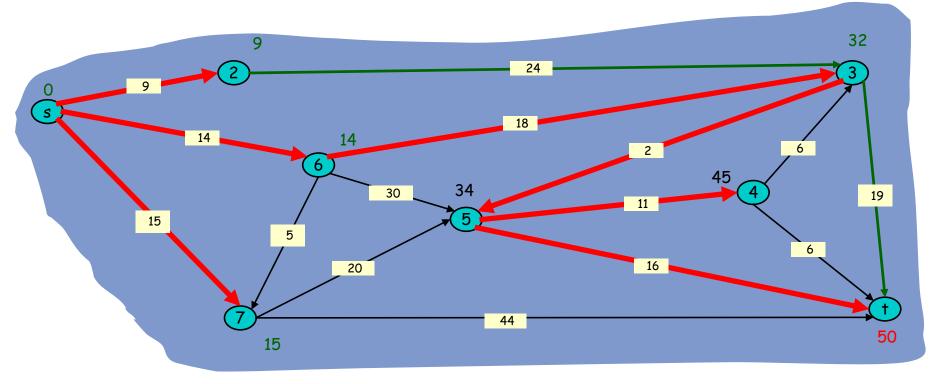


 $S = \{ s, 2, 3, 4, 5, 6, 7 \}$ PQ = $\{ t \}$



S	2	3	4	5	6	7	t
0	9	32	45	34	14	15	51

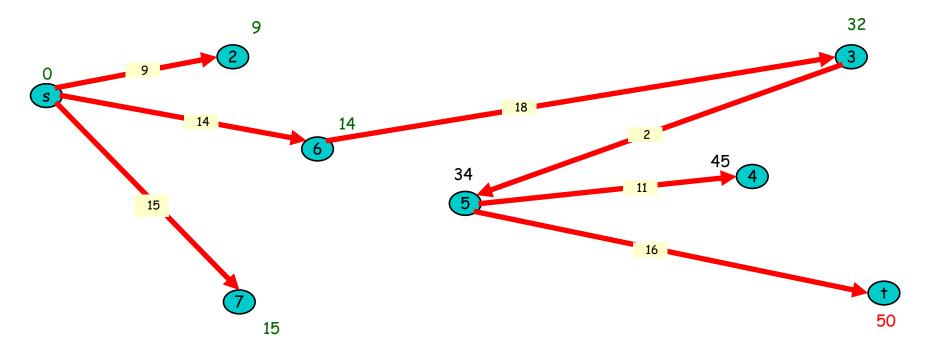
 $S = \{ s, 2, 3, 4, 5, 6, 7, t \}$ PQ = $\{ \}$





S	2	3	4	5	6	7	†
0	9	32	45	34	14	15	51

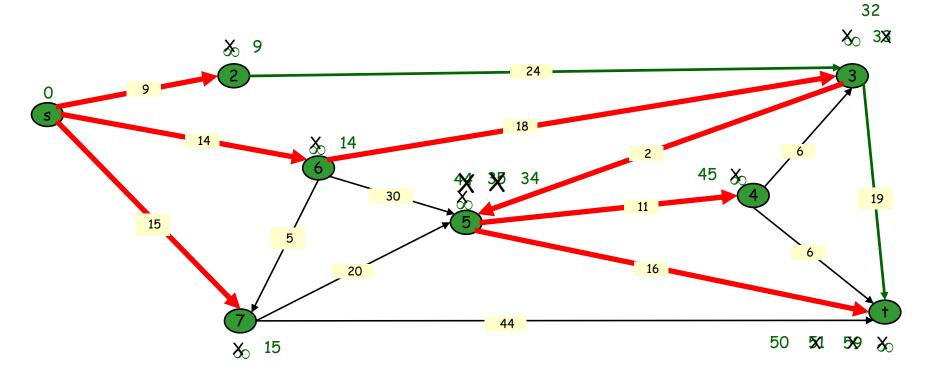
S = { s, 2, 3, 4, 5, 6, 7,† } PQ = {}



Spanning Tree (간선트리)



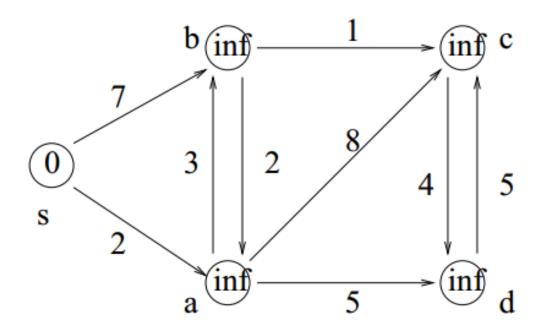
S	2	3	4	5	6	7	t
0	9	32	45	34	14	15	50





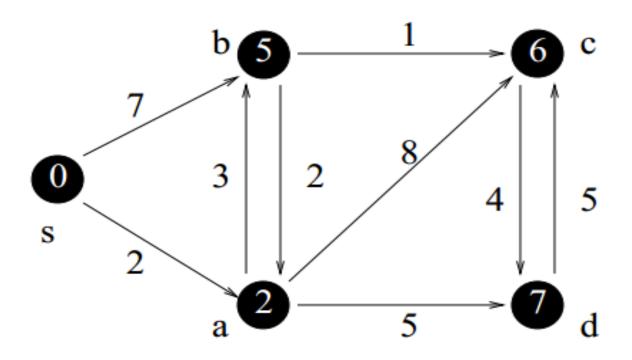
EXERCISE 1

Find the shortest path from 0 to every vertices on this graph





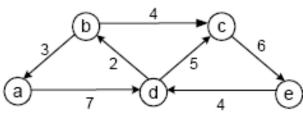
Solution 1



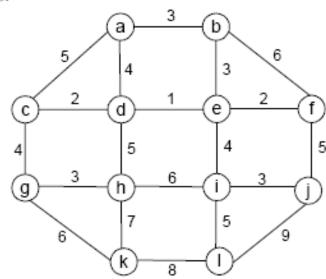


2. Solve the following instances of the single-source shortest-paths problem with vertex a as the source:

a.

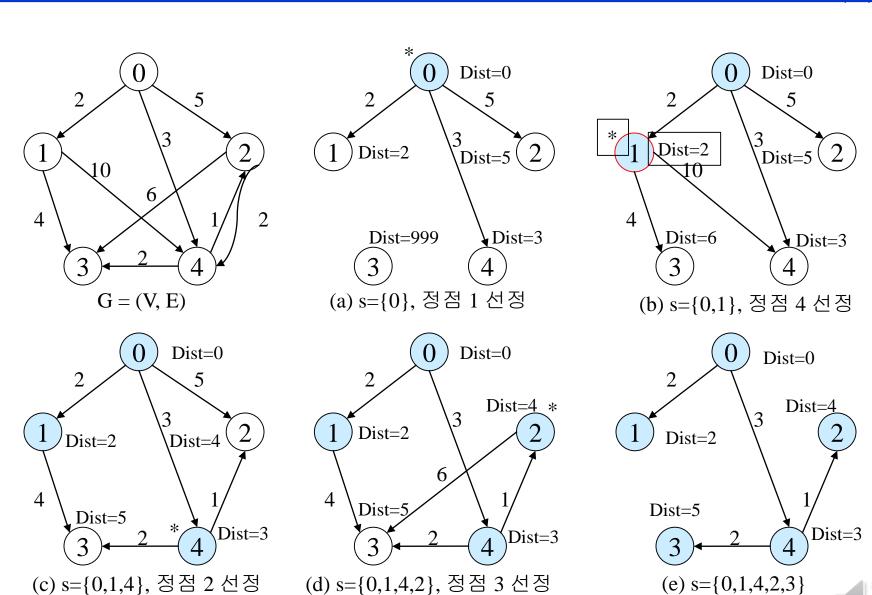


b.



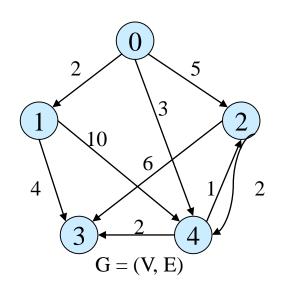


하나의 정점에서 다른 모든 정점까지의 최단경로(3)



최단 경로 계산의 예

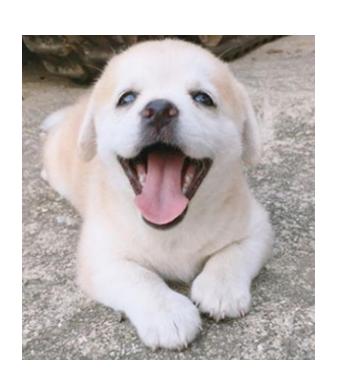
하나의 정점에서 다른 모든 정점까지의 최단경로(6)

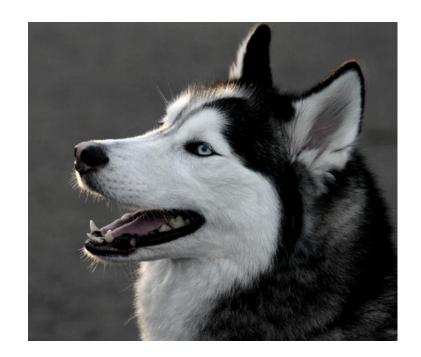


for 루프	선택된 정점	S=true인 정점			Dist[i]	
101 —	정점		[0]	[1]	[2]	[3]	[4]
초기화		[0]	0	2	5	999	3
1	[1]	[0],[1]	0	2	5	6	3
2	[4]	[0],[1],[4]	0	2	4	5	3
3	[2]	[0],[1],[4],[2]	0	2	4	5	3

그래프 G에 대한 shortestPath 수행 내용









감사합니다.

