



Development of an AGV Algorithm for Parking Lot

Implementation of an AGV Control Algorithm and Verification via Simulation

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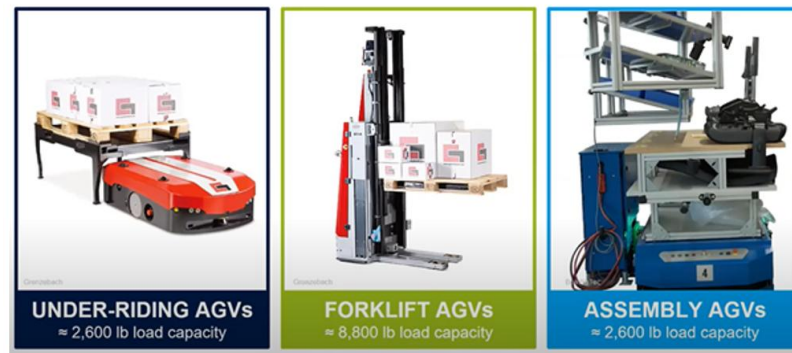
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Definition of AGV

- What is the AGV?

- Automated Guided Vehicle (무인 운반차)

- 정해진 경로를 따라 자체 동력으로만 자재를 운반하는 목적을 가진 자동화 차량



출처 : Stiles Machinery

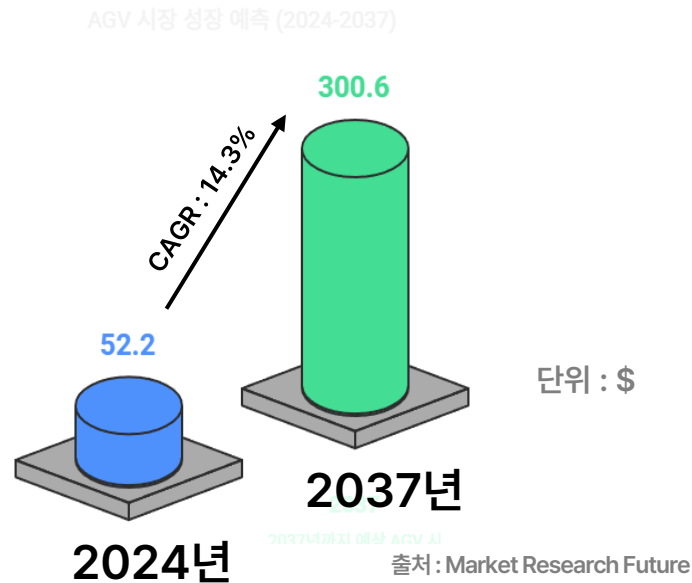
AGV Market Size

- Present Market size and its projected future growth.

- 2024년 AGV시장규모, 약52억 2천만 달러로 평가

- 2037년 약 300억 6천만 달러 예상

- 연평균 성장률 약 14.3% 예측



Utilization of AGV Control System

- Parking Robot

➤ 현대위아 vs HL만도, 주차로봇 대결 막 올랐다.



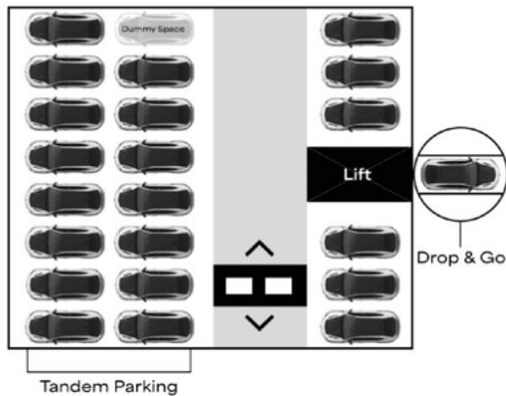
현대위아와 HL만도의 주차로봇 (사진=각사)

출처 - ZD코리아

Necessity of AGV In Parking lot

- 주차장 공간 효율성 증가

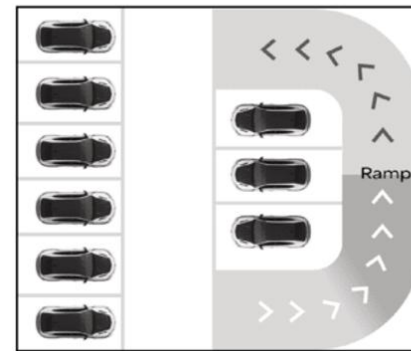
- 차량 사이 간격 최소화로 주차 밀도 향상
- 기존의 주차 도로를 주차공간으로 활용 가능



MPS System

25 Spaces ↑

출처: 삼표



Conventional Parking

8 Spaces

Necessity of AGV In Parking lot

- 사고 예방

- 인적 오류로 인한 충돌 사고 제거
- 알고리즘 기반 안전 주행으로 차량 손상 방지
- 2023년 주차장 사고 건수 **1542건**, 하루 평균 **4.2건** 꼴로 발생 (출처: 경찰청)

2023년 주차장 사고 건수

1542건



하루평균 4.2건의 주차장 사고 발생

Necessity of AGV In Parking lot

- 환경 오염 감소

- 주차장 무인화에 따른 차량 주행 배기가스 배출 '0mg'
- 주차장 주변 환경의 공기질 개선

<지하주차장의 평균 일산화탄소 농도>

$$K = \frac{0.06 \cdot X \cdot M \cdot L}{V \cdot Q}$$

단, K : 지하주차장의 평균 일산화탄소 농도 [m³/m³]

X : 자동차 출입대수 [대/h]

M : 주행시 일산화탄소 평균배출량 [m³/min·대]

L : 평균주행거리 [m/대]

V : 평균주행속도 [km/h·대]

Q : 환기량 [m³/h]

여기서, 0.06 은 환산계수이다.

Multi-Agent Path Finding [MAPF]

- **Definition of MAPF**

- 여러 에이전트가 공유된 환경에서 충돌없이 이동하는 경로 집합을 찾는 문제

- **Key Components of MAPF**

- Agents
- Environment
- Start / Goal
- Time

- **Collision Types**

- 정점 충돌 [Vertex Conflict]

- 동일 시간에 동일 정점[위치] 차지



- 간선 충돌 [Edge Conflict]

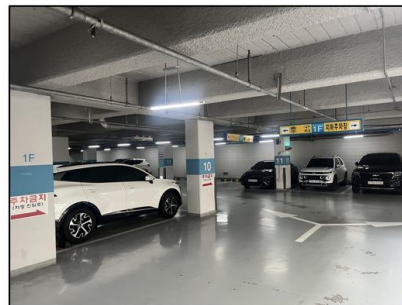
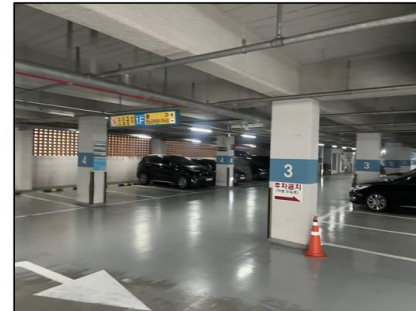
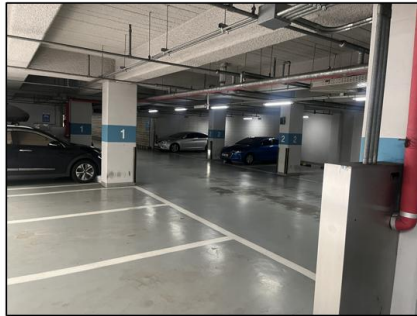
- 동일 시간 사이에 동일 간선[경로]를 교차하여 통과



Environment : 형남공학관 주차장

- 형남공학관 주차장

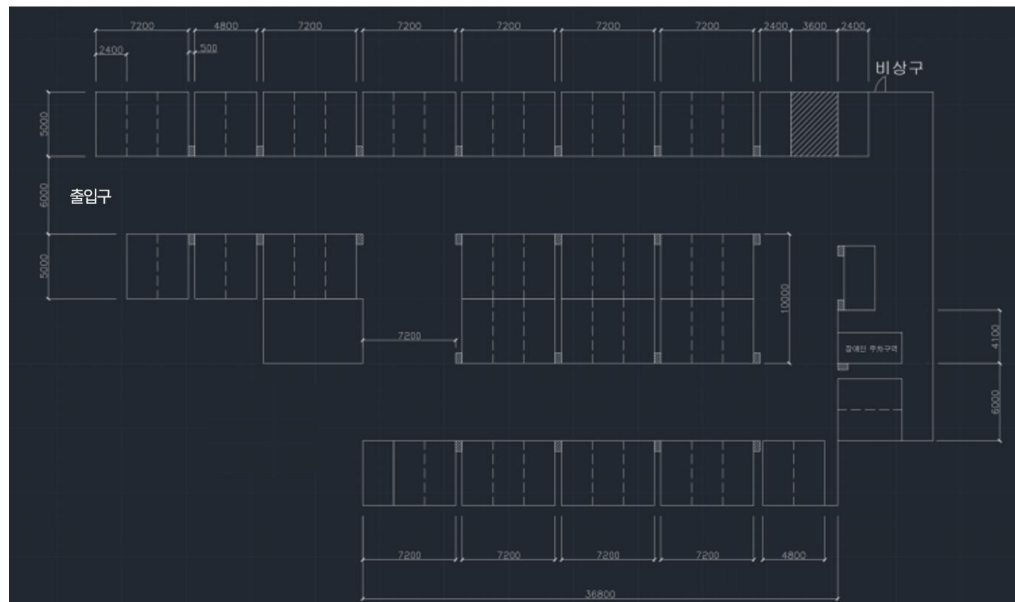
➤ 형남공학관 주차장 지하2층 전경



Environment : 형남공학관 주차장

- 형남공학관 주차장

➤ 형남공학관 주차장 지하2층 도면



[0] Priority System : Score / Stuck Boost

- **Priority Score**

- $\text{Priority}(A) = (\text{Base} \times 100) + \text{StuckBoost} - A.id$
- 출차 [RETURNING_WITH_CAR] : 3
- 충전 필요 [GOING_TO_CHARGE] : 2
- 주차 / 차량 수거 [GOING_TO PARK / GOING TO COLLECT] : 1
- 대기 / 충전 중 / 기지로 복귀 [IDLE / CHARGING / RETURNING_HOME] : 0

- **Stuck Boost Mechanism**

- $$\text{StuckBoost} = \begin{cases} 1000 & \text{if } A.stuck \geq \text{DEADLOCK_THRESHOLD} \\ A.Stuck \times 10 & \text{otherwise} \end{cases}$$

- **Tie-Breaking Mechanism**

- AGENT ID = $A.id$ [ex. $A > B > C$]

[0] Priority System : pseudo-code

- Pseudo-code[의사코드]



```
// basic priority based on agent.state
```

```
p = base(state)
```

```
//Stuck Boost
```

```
if (stuck_steps >= deadlock_threshold)
```

```
  p += stuck_boost_hard
```

```
else if (stuck_steps > 0)
```

```
  p += stuck_steps * stuck_boost_mult
```

```
// Tie-breaking mechanism
```

```
tie_break(A.id)
```

```
// plan and application
```

```
plan_step() → reservation → move → stuck update → (metrics_notify_all)
```

[1] A *algorithm : Definition

• What is the A*algorithm?

➤ A* algorithm

- 그래프 기반 최단 경로 탐색에서 가장 널리 쓰이는 알고리즘
- Shortest Path = Path with $\text{Min}[f(n)]$

➤ $g(n)$

- 출발 노드부터 현재 노드까지 실제 소요 비용

➤ $h(n)$ [Manhattan Distance]

- $h(n) = |x_n - x_{goal}| + |y_n - y_{goal}|$
- 현재 노드부터 목표 노드까지 예상 비용
- 휴리스틱 함수 활용 추정치

➤ $f(n)$

- 총 예상 비용 [Total Estimated Cost]
- $f(n) = g(n) + h(n)$



[1] A *algorithm : pseudo-code / Diagram

• pseudo-code

```
➤ function plan_astar(agents, map):  
  next_pos = get_current_positions(all_agents)  
  sorted_agents = sort_by_priority(agents)  
  
  for agent in sorted_agents:  
    if agent.is_waiting(rotation or action) then  
      continue  
  
    temp_obstacles =  
      get_higher_priority_moves(agent, next_pos) +  
      get_lower_priority_positions(agent, agents)  
    temp_pathfinder = create_pathfinder(agent.pos,  
agent.goal)  
    path = compute_path(  
      temp_pathfinder, map, temp_obstacles  
    )  
    desired_move = get_next_step(path)  
    destroy_pathfinder(temp_pathfinder)  
    next_pos[agent.id] = apply_rotation(agent, desired_move)  
  
  resolve_final_conflicts(next_pos, sorted_agents)  
  return next_pos
```



[2] D* lite algorithm : Definition

• What is the D*lite algorithm?

➤ D* lite algorithm

- 동적 환경에서 주로 쓰이는 알고리즘
- 목표에서 역방향으로 모든 지점까지의 최단 경로 비용 계산
- Shortest Path = All Consistent path

➤ $g(s)$

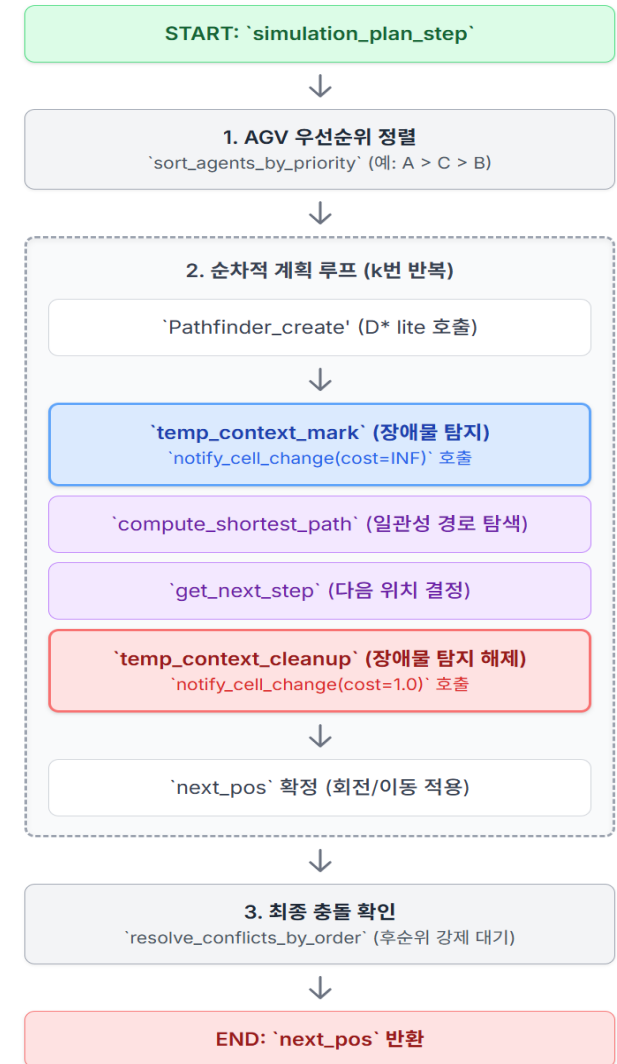
- 목표 노드부터 현재 노드까지 '알려진' 추정 소요 비용

➤ $rhs(s)$

- $g(s)$ 의 1- step 예측 값.
- $$rhs(s) = \begin{cases} 0 & \text{if } s = s_{goal} \\ \min_{s' \in Succ(s)} (c(s, s') + g(s')) & \text{otherwise} \end{cases}$$

➤ $C(s, s')$

- $S' \rightarrow S$ 비용



[2] D* lite algorithm : Definition

• What is the D*lite algorithm?

➤ Consistency

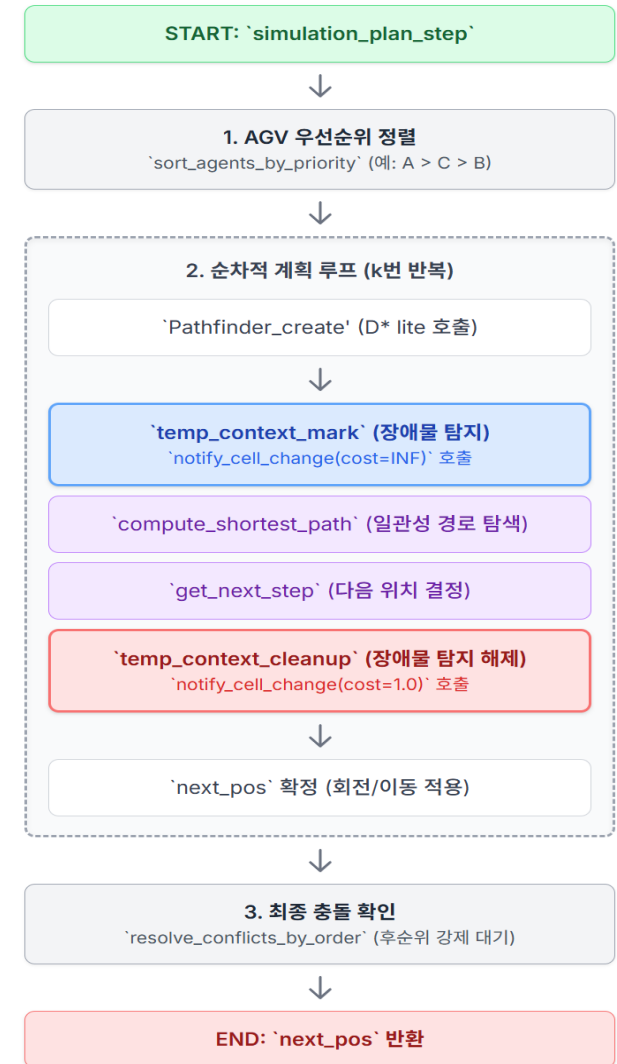
- $g(s) = rhs(s)$: Node 's' is consistent.
- $g(s) \neq rhs(s)$: Node 's' is inconsistent. -> **Update Target**

➤ Priority Key

- $Key(s) = [k_1(s); k_2(s)]$
- 우선순위 큐 사용하여 '비일관성' 노드 갱신 -> 일관성 's' 탐색
- 총 예상 경로 비용 : $k_1(s) = \min(g(s), rhs(s)) + h(s_{start}, s) + k_m$
- 현재노드 ~ 목표노드 예상 비용 : $k_2(s) = \min(g(s), rhs(s))$

➤ ComputerShortestpath()

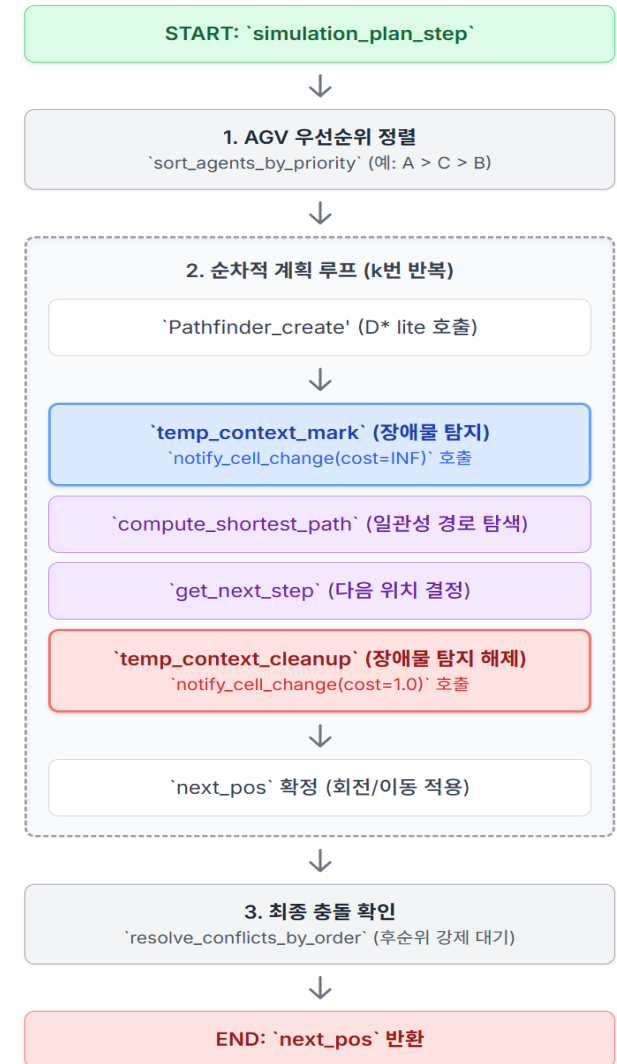
- while $\left(\min_{u \in U} key(u) < key(s_{start}) \right) \text{ or } (rhs(s_{start}) \neq g(s_{start}))$



[2] D* lite algorithm : pseudo-code / Diagram

• pseudo-code

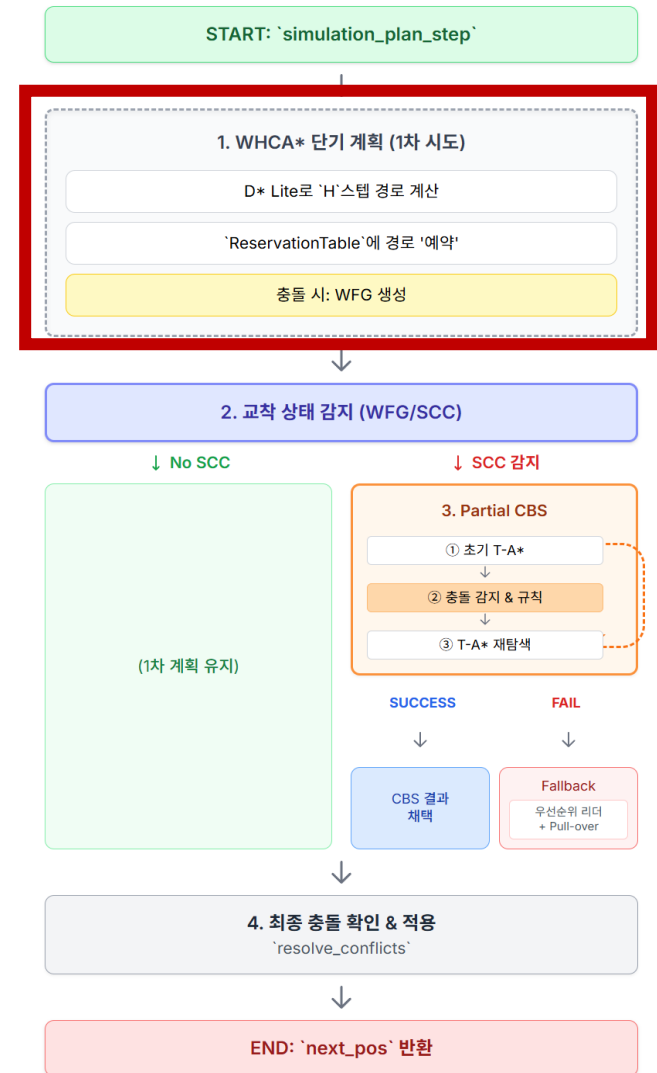
```
➤ function plan_dstar_lite(agents, map):  
  next_pos = get_current_positions(all_agents)  
  sorted_agents = sort_by_priority(agents)  
  
  for agent in sorted_agents:  
    if agent.is_waiting(rotation or action) then  
      continue  
    ensure_pathfinder_exists(agent)  
  
    temp_obstacles =  
      get_higher_priority_moves(agent, next_pos) +  
      get_lower_priority_positions(agent, agents)  
  
    notify_obstacles(agent.pathfinder, temp_obstacles)  
    path = repair_path(agent.pathfinder, map)  
    desired_move = get_next_step(path)  
    cleanup_notify_obstacles(agent.pathfinder,  
      temp_obstacles)  
    next_pos[agent.id] = apply_rotation(agent, desired_move)  
  
  resolve_final_conflicts(next_pos, sorted_agents)  
  return next_pos
```



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : WHCA*

• What is the WHCA* algorithm?

- WHCA* [Windowed Hierarchical Cooperative A*]
 - 우선순위, 예약 테이블을 활용한 MAPF 알고리즘
- Pathfinder Algorithm
 - D* lite : Provides the Ideal route
- W : Space-Time Reservation Table
 - [time][x][y]의 3차원 예약 테이블 ($0 < \text{time} < H$) *H = Horizon
- H [Hierarchical]
 - 우선순위에 따른 순차적 경로 계획
- C [Cooperative]
 - 타 에이전트 예약 노드 -> 장애물
- Limitations
 - 우선순위 기반 알고리즘 -> 'Risk of Deadlock'



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : WFG/SCC

• What is the WFG?

➤ WFG [Wait-For Graph]

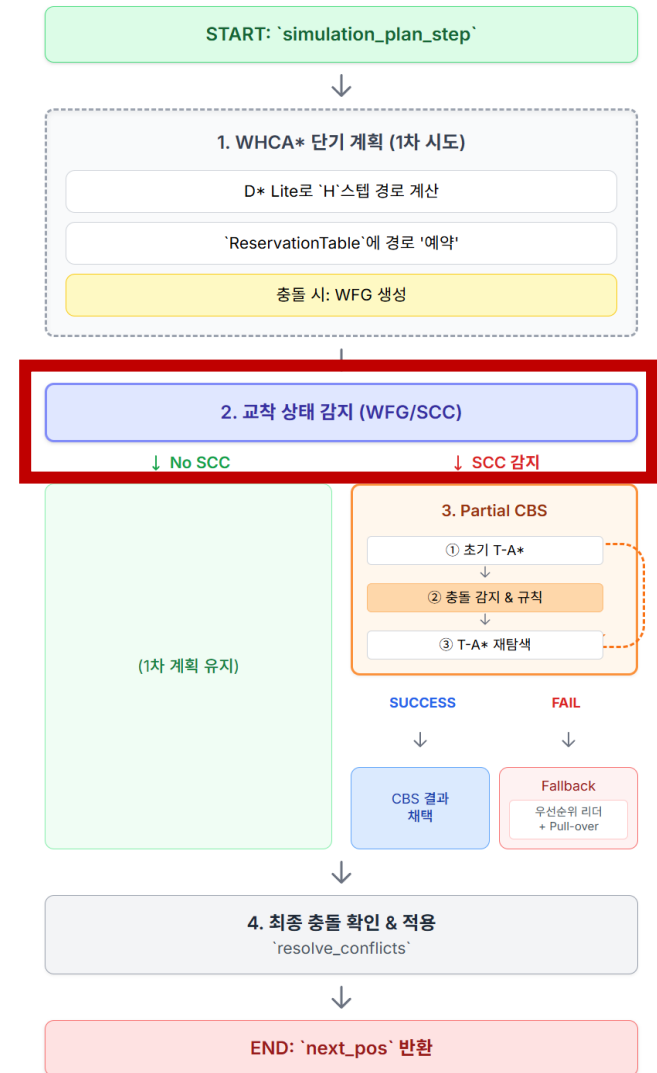
- 현재 시스템 대기 상태 표현하는 자료구조.

➤ Role

- 교착 상태 탐지위한 자료구조

➤ Components

- $G = (V, E)$
- V [Vertex / Node]: 정점의 집합 (운용 AGV: $\{A_1, A_2, \dots\}$)
- E [Edge]: 간선의 집합 (대기 관계: $\{(A_i, A_j) \mid A_i \text{ waits for } A_j\}$)



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : WFG/SCC

• What is the SCC?

➤ SCC [Strongly Connected Components]

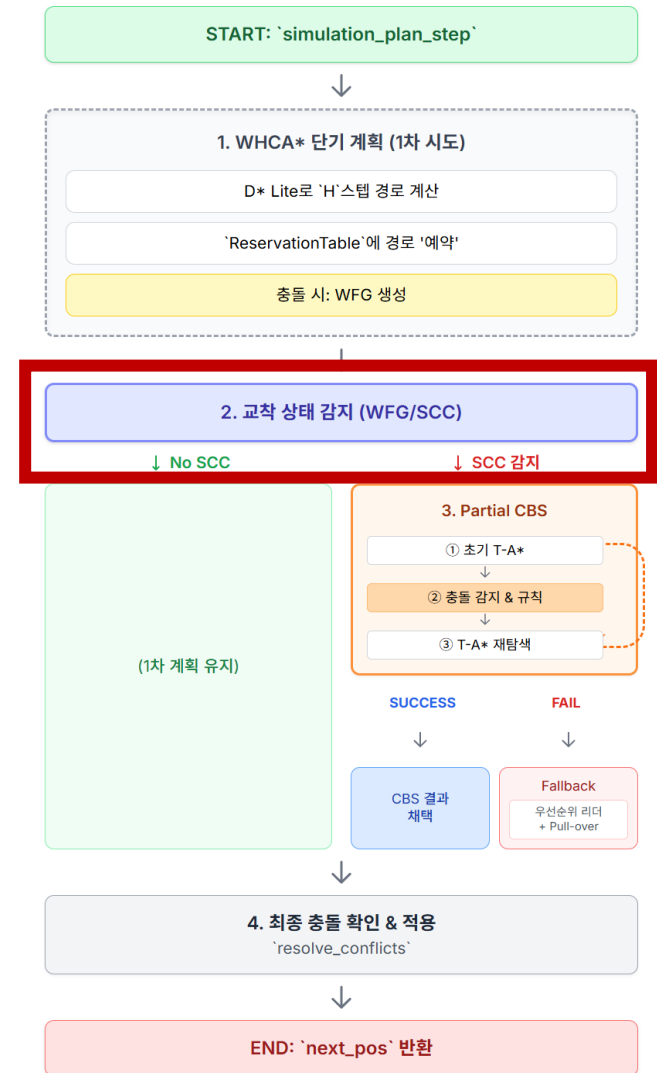
- 그래프 내에서 서로에게 도달 가능한 정점들의 집합
- 크기 2이상 SCC 존재 [ex. $A_i \leftrightarrow A_j$], -> 교착 상태 의미

➤ Role

- 자료구조에서 교착상태 감지

➤ DeadLock Detetction

- Deadlock $\Leftrightarrow \exists \text{ SCC } C \subseteq V \text{ s.t. } |C| \geq 2$



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : CBS

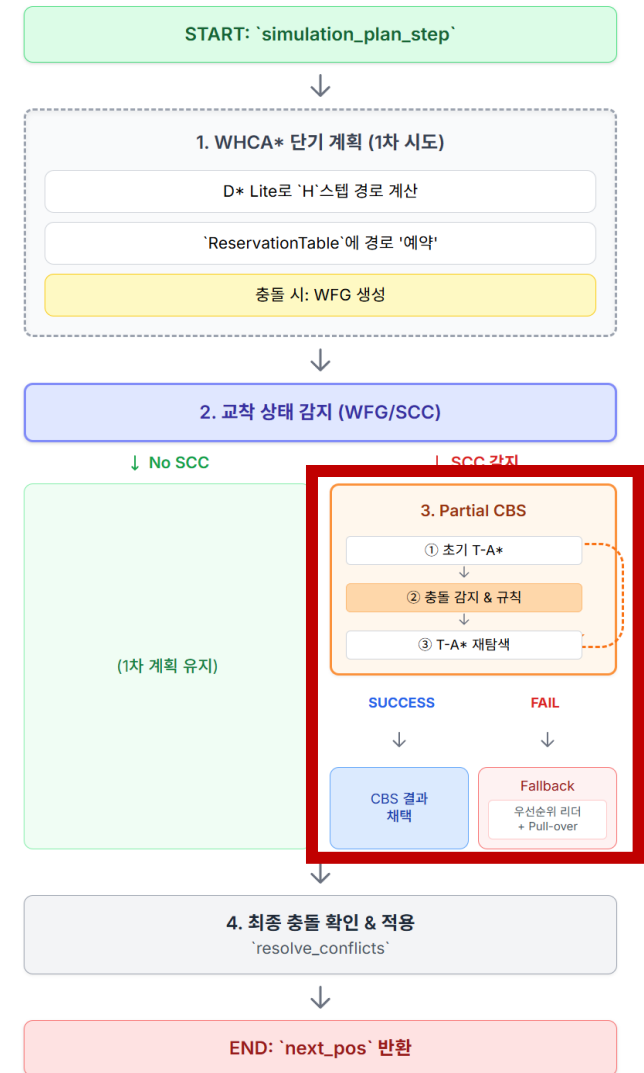
• What is the CBS algorithm?

➤ CBS [Conflict – Based Search]

- 지역적 충돌 해소 위한 MAPF 알고리즘
- 소수 에이전트 간의 충돌을 효율적으로 해결함.
- 시스템의 계산 효율성을 향상.

➤ 저수준 [Low – Level] : Time A* -> 개별 경로 탐색기

- [time][x][y] 구조에서 A* 탐색 수행
- State (상태): $s = (x, y, t)$
- 평가 함수 (f-value): $f(s) = g(s) + h(s)$
- $g(s)$: $g(x, y, t) = t \ ((x_0, y_0, 0) \rightarrow (x, y, t) : \text{비용 } t)$
- $h(s)$: $h(x, y, t) = |x - x_{goal}| + |y - y_{goal}|$



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : CBS

• What is the CBS algorithm?

➤ 고수준 [High-Level] : 충돌 해결 위한 제약 조건 결정

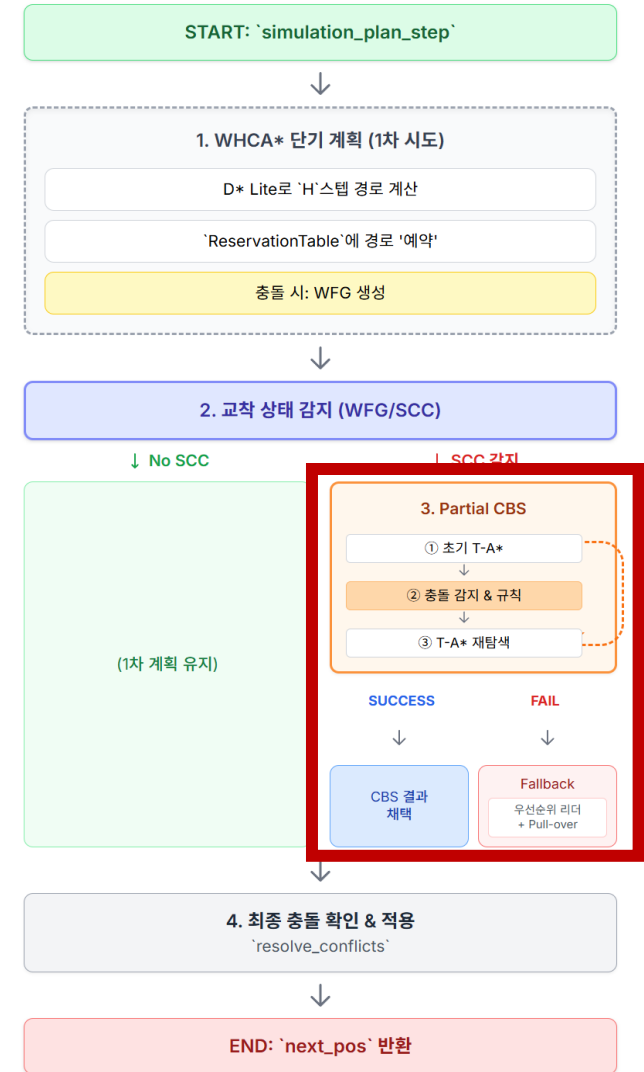
- Cost(N) 가장 낮은 노드 N -> 우선순위 큐에서 선택

$$Cost(N) = \sum_{i \in Agents} PathCost(P_i)$$

- N: 충돌 트리의 노드 [제약 조건의 집합]
- P_i : 노드 N의 제약 조건 & 저수준 탐색 기반 에이전트 i의 경로
- $PathCost(P_i)$: 경로 비용
 $\alpha \cdot \text{"moves"} + \beta \cdot \text{"waits"} + \gamma \cdot h_{\text{residual}}$

➤ 제약 조건 (Constraint) 예시:

- High Level에서 (Agent A, (5, 5), $t = 3$) 제약이 내려오면,
- Low Level은 (5,5,3) 상태를 고려하지 않는다.



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : DeadLock resolve

• 1. 교착 감지

- WFG & SCC 활용하여 교착 상태 식별

• 2. Partical CBS

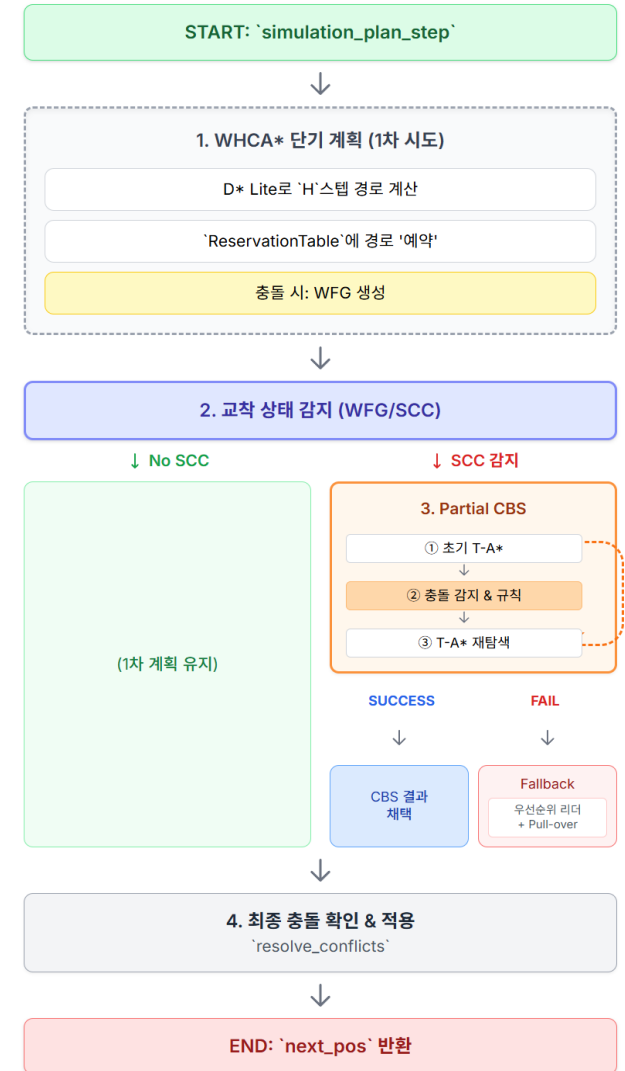
- SCC 식별 그룹에 Partial CBS 적용

• 3. 우선순위 기반 해결

- CBS 실패할 시, 우선순위 기반 해결
- A*, D* lite와 동일한 우선순위 기준

• 4. Pull – over

- Max Stuck value AGV -> Leader AGV
- 우선순위 실패할 시, Leader 제외 AGV 한 칸씩 후진.



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : pseudo-code

function plan_WHCAstar(agents, map):

// 1. WHCA (First attempt)*

rt = create_reservation_table()

wf_edges = []

sorted_agents = sort_by_priority(agents)

for agent in sorted_agents:

// (Plan H_steps using D Lite)*

path = dstar_find_path(agent.pf, h_steps, rt)

// (Build WFG - Wait-For Graph)

if has_conflict(path) with other_agent then

add_wait_edge(wf_edges, agent, other_agent)

provisional_plan[agent.id] = agent.pos *// wait*

else

reserve_path(rt, path)

provisional_plan[agent.id] = path[1] *// move forward*

// 2. wfg/scc (Deadlock detection)

sccmask = detect_scc(wf_edges)

// 3. partial cbs (Deadlock resolution branch)

if sccmask > 0 then

// (Deadlock occurred)

deadlocked_group = get_agents_from_mask(sccmask)

// (Attempt negotiation)

negotiated_plans, ok

= run_partial_cbs(deadlocked_group, rt)

if ok == true then

// (Negotiation success: Overwrite with CBS results)

overwrite_next_pos(provisional_plan, negotiated_plans)

else

// (Negotiation failure: Only leader moves)

fallback_leader_moves(provisional_plan, deadlocked_group)

// No deadlock. Keep the first-attempt plan

else : (provisional_plan)

// 4. Final

resolve_final_conflicts(provisional_plan, sorted_agents)

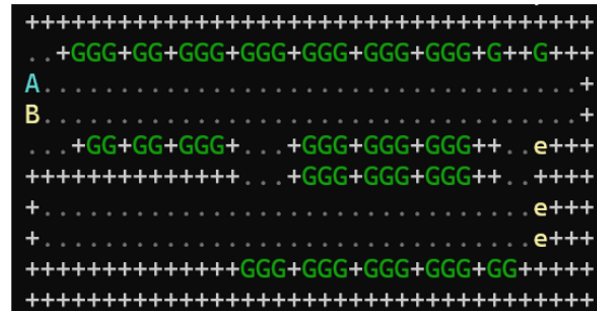
return provisional_plan

Simulation Map

- **형남공학관 주차장**

- AGV : 1~4

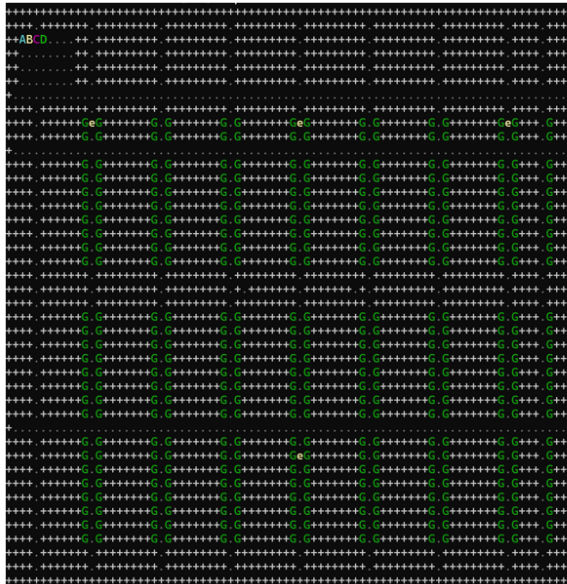
- Parking : 61



- **MAP2**

- AGV : 4

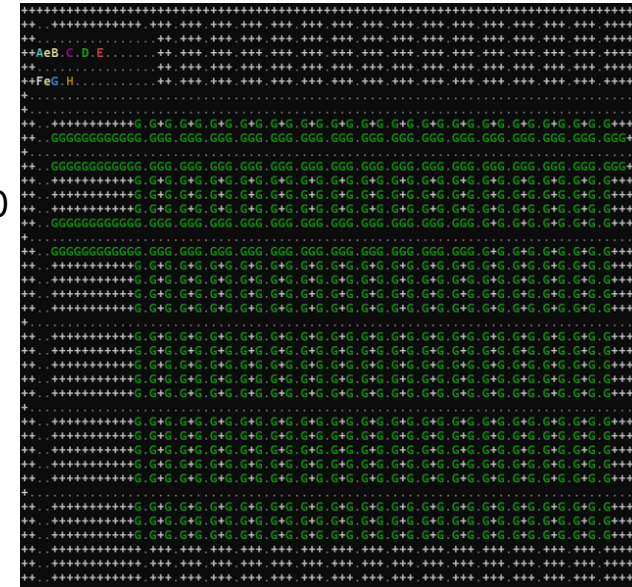
- Parking : 390



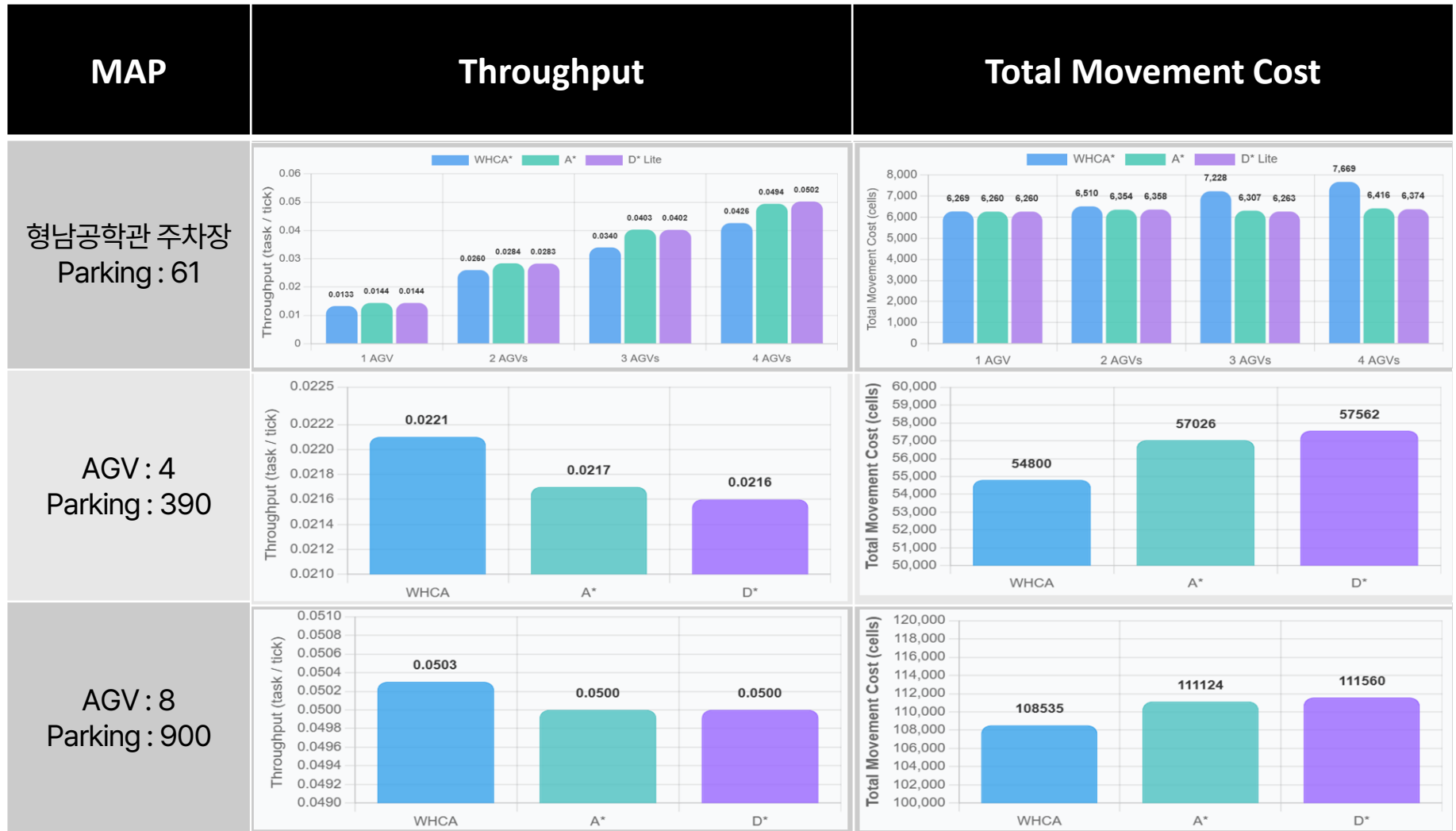
- **MAP3**

- AGV : 8

- Parking : 900



Simulation Result



Simulation Result

MAP	Nodes Expanded	Heap Moves																																								
형남공학관 주차장 Parking : 61	<table><tr><th>AGVs</th><th>WHCA*</th><th>A*</th><th>D* Lite</th></tr><tr><td>1 AGV</td><td>17,437</td><td>176,450</td><td>17,437</td></tr><tr><td>2 AGVs</td><td>17,398</td><td>178,357</td><td>18,886</td></tr><tr><td>3 AGVs</td><td>25,000</td><td>228,913</td><td>35,689</td></tr><tr><td>4 AGVs</td><td>29,235</td><td>298,061</td><td>55,849</td></tr></table>	AGVs	WHCA*	A*	D* Lite	1 AGV	17,437	176,450	17,437	2 AGVs	17,398	178,357	18,886	3 AGVs	25,000	228,913	35,689	4 AGVs	29,235	298,061	55,849	<table><tr><th>AGVs</th><th>WHCA*</th><th>A*</th><th>D* Lite</th></tr><tr><td>1 AGV</td><td>115,974</td><td>995,969</td><td>115,074</td></tr><tr><td>2 AGVs</td><td>108,232</td><td>1,020,243</td><td>125,947</td></tr><tr><td>3 AGVs</td><td>144,242</td><td>1,329,512</td><td>244,396</td></tr><tr><td>4 AGVs</td><td>162,110</td><td>1,663,126</td><td>372,340</td></tr></table>	AGVs	WHCA*	A*	D* Lite	1 AGV	115,974	995,969	115,074	2 AGVs	108,232	1,020,243	125,947	3 AGVs	144,242	1,329,512	244,396	4 AGVs	162,110	1,663,126	372,340
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AGV : 4 Parking : 390	<table><tr><th>Algorithm</th><th>Nodes Expanded (total)</th></tr><tr><td>WHCA*</td><td>320,044</td></tr><tr><td>A*</td><td>7,789,182</td></tr><tr><td>D* Lite</td><td>1,210,534</td></tr></table>	Algorithm	Nodes Expanded (total)	WHCA*	320,044	A*	7,789,182	D* Lite	1,210,534	<table><tr><th>Algorithm</th><th>Heap Moves (total)</th></tr><tr><td>WHCA*</td><td>2,419,374</td></tr><tr><td>A*</td><td>53,850,530</td></tr><tr><td>D* Lite</td><td>11,599,644</td></tr></table>	Algorithm	Heap Moves (total)	WHCA*	2,419,374	A*	53,850,530	D* Lite	11,599,644																								
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Simulation Result



Conclusion

- **Which algorithm has the highest throughput in a low-density environment?**
 - A* algorithm And D* lite Algorithm
- **Which algorithm has the highest throughput in a High-density environment?**
 - WHCA* + WFG/SCC + CBS Algorithm
- **Which algorithm requires the least computation under the same conditions?**
 - 1st : WHCA* + WFG/SCC + CBS Algorithm
 - 2nd : D* Lite Algorithm
 - 3rd : A* Algorithm
- **"The optimal algorithm varies depending on the context."**

[illegible][illegible]

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