

Development of an AGV Algorithm for Parking Lot

Implementation of an AGV Control Algorithm and Verification via Simulation

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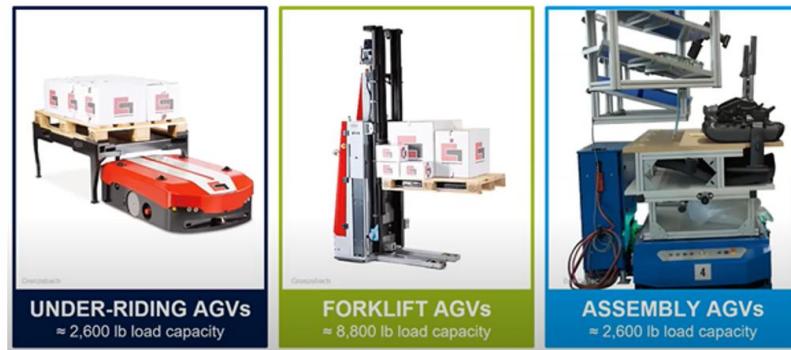
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Definition of AGV

- **What is the AGV?**

➤ Automated Guided Vehicle (무인 운반차)

- 정해진 경로를 따라 자체 동력으로만 자재를 운반하는 목적을 가진 자동화 차량



출처 : Stiles Machinery

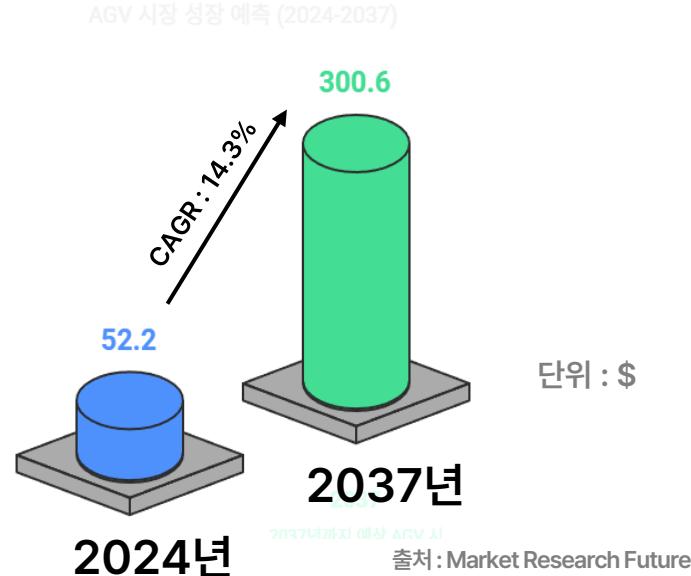
AGV Market Size

- Present Market size and its projected future growth.

➤ 2024년 AGV시장규모, 약52억 2천만 달러로 평가

➤ 2037년 **약 300억 6천만 달러** 예상

➤ 연평균 성장률 약 14.3% 예측



Utilization of AGV Control System

- Parking Robot

➤ 현대위아 vs HL만도, 주차로봇 대결 막 올랐다.



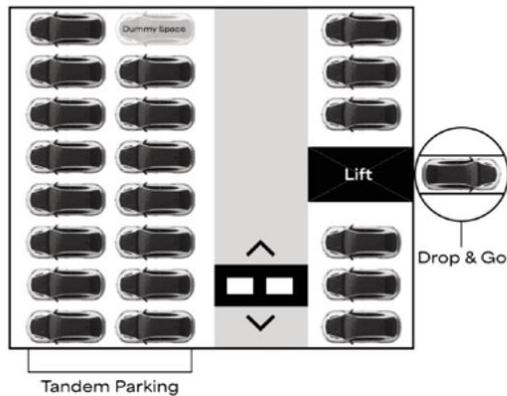
현대위아와 HL만도의 주차로봇 (사진=각사)

출처 - ZD코리아

Necessity of AGV In Parking lot

- 주차장 공간 효율성 증가

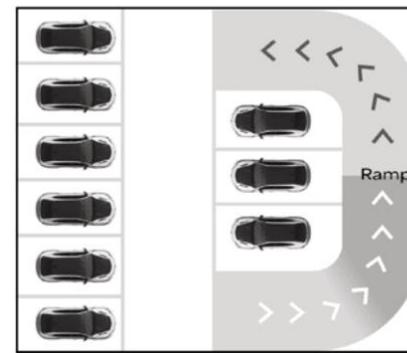
- 차량 사이 간격 최소화로 주차 밀도 향상
- 기존의 주차 도로를 주차공간으로 활용 가능



MPSSystem

25 Spaces ↑

출처 : 삼표



Conventional Parking

8 Spaces

Necessity of AGV In Parking lot

- 사고 예방

- 인적 오류로 인한 충돌 사고 제거
- 알고리즘 기반 안전 주행으로 차량 손상 방지
- 2023년 주차장 사고 건수 **1542건**, 하루 평균 **4.2건** 꼴로 발생 (출처: 경찰청)

2023년 주차장 사고 건수

1542건



하루평균 4.2건의 주차장 사고 발생

Necessity of AGV In Parking lot

- 환경 오염 감소

- 주차장 무인화에 따른 차량 주행 배기가스 배출 '0mg'
- 주차장 주변 환경의 공기질 개선

<지하주차장의 평균 일산화탄소 농도>

$$K = \frac{0.06 \cdot X \cdot M \cdot L}{V \cdot Q}$$

단, K : 지하주차장의 평균 일산화탄소 농도 [m^3/m^3]

X : 자동차 출입대수 [대/h]

M : 주행시 일산화탄소 평균배출량 [$m^3/min\cdot 대$]

L : 평균주행거리 [m/대]

V : 평균주행속도 [$km/h\cdot 대$]

Q : 환기량 [m^3/h]

여기서, 0.06은 환산계수이다.

Multi-Agent Path Finding [MAPF]

- **Definition of MAPF**

- 여러 에이전트가 공유된 환경에서 충돌없이 이동하는 경로 집합을 찾는 문제

- **Key Components of MAPF**

- Agents
- Environment
- Start / Goal
- Time

- **Collision Types**

- 정점 충돌 [Vertex Conflict]



- 동일 시간에 동일 정점[위치] 차지

- 간선 충돌 [Edge Conflict]

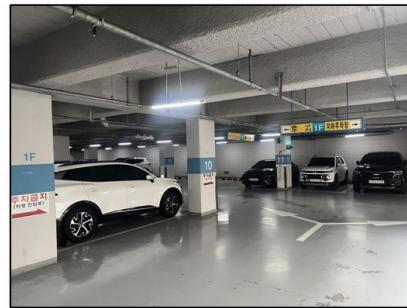
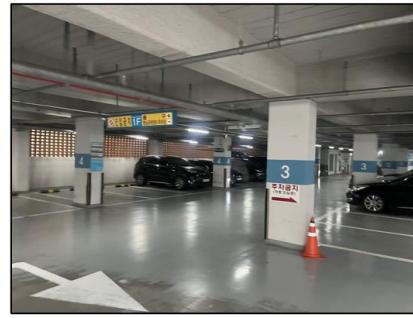
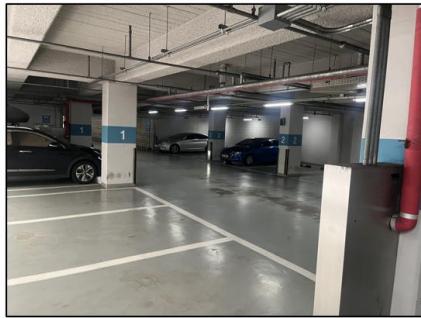
- 동일 시간 사이에 동일 간선[경로]를 교차하여 통과



Environment : 형남공학관 주차장

- 형남공학관 주차장

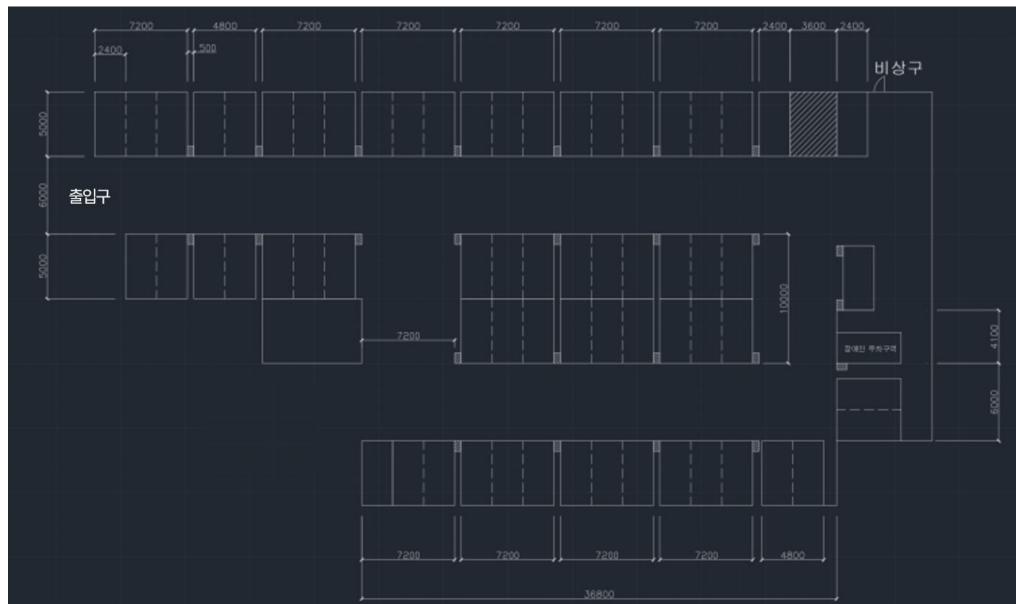
➤ 형남공학관 주차장 지하2층 전경



Environment : 형남공학관 주차장

- 형남공학관 주차장

➤ 형남공학관 주차장 지하2층 도면



[0] Priority System : Score / Stuck Boost

- Priority Score

- Priority(A) = ($Base \times 100$) + $StuckBoost - A.id$
- 출차 [RETURNING_WITH_CAR] : 3
- 충전 필요 [GOING_TO_CHARGE] : 2
- 주차 / 차량 수거 [GOING_TO PARK / GOING TO COLLECT] : 1
- 대기 / 충전 중 / 기지로 복귀 [IDLE / CHARGING / RETURNING_HOME] : 0

- Stuck Boost Mechanism

- $StuckBoost = \begin{cases} 1000 & \text{if } A.stuck \geq DEADLOCK_THRESHOLD \\ A.Stuck \times 10 & \text{otherwise} \end{cases}$

- Tie-Breaking Mechanism

- AGENT ID = $A.id$ [ex. A > B > C]

[0] Priority System : pseudo-code

- Pseudo-code[의사코드]

```
➤ // basic priority based on agent.state  
p = base(state)  
  
// Stuck Boost  
if (stuck_steps >= deadlock_threshold)  
    p += stuck_boost_hard  
else if (stuck_steps > 0)  
    p += stuck_steps * stuck_boost_mult  
  
// Tie-breaking mechanism  
tie_break(A.id)  
  
// plan and application  
plan_step() → reservation → move → stuck update → (metrics_notify_all)
```

[1] A *algorithm : Definition

• What is the A*algorithm?

➤ A* algorithm

- 그래프 기반 최단 경로 탐색에서 가장 널리 쓰이는 알고리즘
- Shortest Path = Path with Min[$f(n)$]

➤ $g(n)$

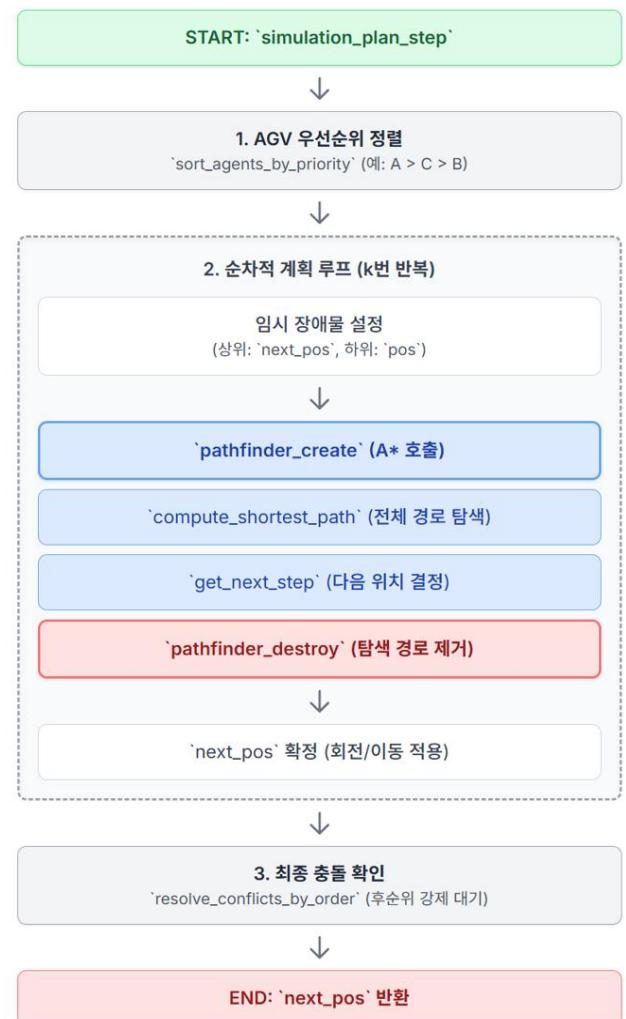
- 출발 노드부터 현재 노드까지 실제 소요 비용

➤ $h(n)$ [Manhattan Distance]

- $h(n) = |x_n - x_{goal}| + |y_n - y_{goal}|$
- 현재 노드부터 목표 노드까지 예상 비용
- 휴리스틱 함수 활용 추정치

➤ $f(n)$

- 총 예상 비용 [Total Estimated Cost]
- $f(n) = g(n) + h(n)$



[1] A *algorithm : pseudo-code / Diagram

• pseudo-code

```
> function plan_astar(agents, map):
    next_pos = get_current_positions(all_agents)
    sorted_agents = sort_by_priority(agents)

    for agent in sorted_agents:
        if agent.is_waiting(rotation or action) then
            continue

        temp_obstacles =
            get_higher_priority_moves(agent, next_pos) +
            get_lower_priority_positions(agent, agents)
        temp_pathfinder = create_pathfinder(agent.pos,
                                             agent.goal)
        path = compute_path(
            temp_pathfinder, map, temp_obstacles
        )
        desired_move = get_next_step(path)
        destroy_pathfinder(temp_pathfinder)
        next_pos[agent.id] = apply_rotation(agent, desired_move)

    resolve_final_conflicts(next_pos, sorted_agents)
    return next_pos
```



[2] D* lite algorithm : Definition

• What is the D*lite algorithm?

➤ D* lite algorithm

- 동적 환경에서 주로 쓰이는 알고리즘
- 목표에서 역방향으로 모든 지점까지의 최단 경로 비용 계산
- Shortest Path = All Consistent path

➤ $g(s)$

- 목표 노드부터 현재 노드까지 '알려진' 추정 소요 비용

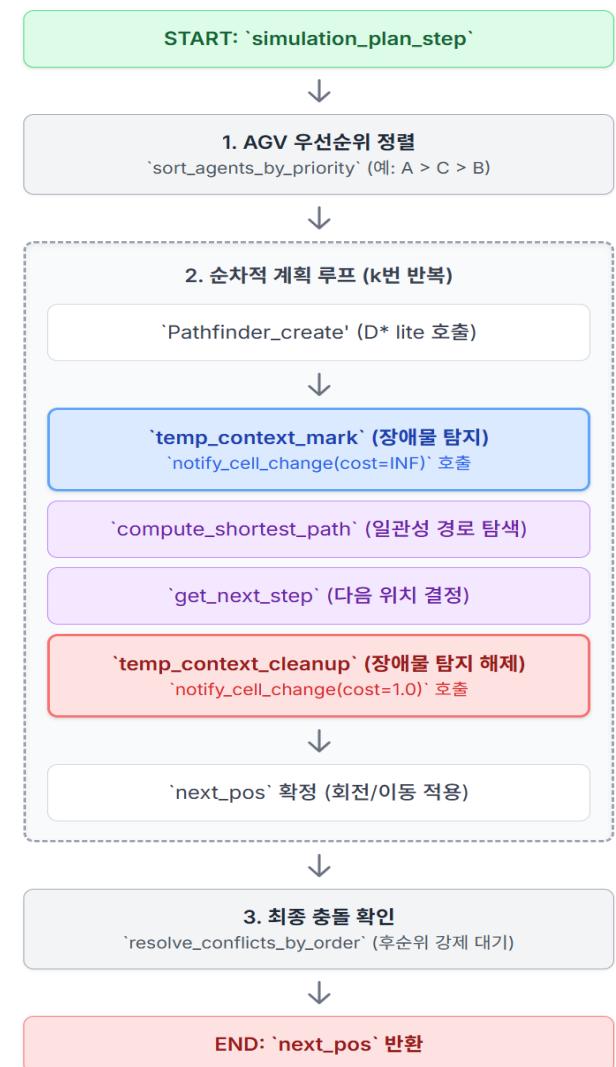
➤ $rhs(s)$

- $g(s)$ 의 1-step 예측 값.

$$- rhs(s) = \begin{cases} 0 & \text{if } s = s_{goal} \\ \min_{s' \in succ(s)} (c(s, s') + g(s')) & \text{otherwise} \end{cases}$$

➤ $c(s, s')$

- $S' \rightarrow S$ 비용



[2] D* lite algorithm : Definition

• What is the D*lite algorithm?

➤ Consistency

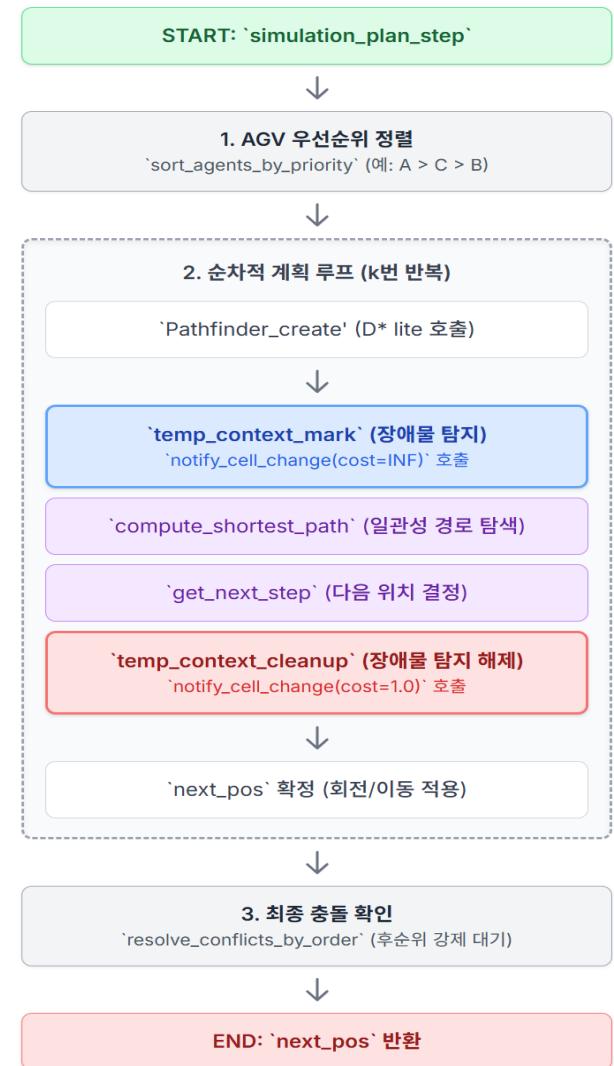
- $g(s) = rhs(s)$: Node 's' is consistent.
- $g(s) \neq rhs(s)$: Node 's' is inconsistent. -> **Update Target**

➤ Priority Key

- $Key(s) = [k_1(s); k_2(s)]$
- 우선순위 큐 사용하여 '비일관성' 노드 갱신 -> 일관성 's' 탐색
- 총 예상 경로 비용: $k_1(s) = \min(g(s), rhs(s)) + h(s_{start}, s) + k_m$
- 현재노드 ~ 목표노드 예상 비용: $k_2(s) = \min(g(s), rhs(s))$

➤ ComputerShortestpath()

- while $\left(\min_{u \in U} key(u) < key(s_{start}) \right) \text{ or } (rhs(s_{start}) \neq g(s_{start}))$



[2] D* lite algorithm : pseudo-code / Diagram

• pseudo-code

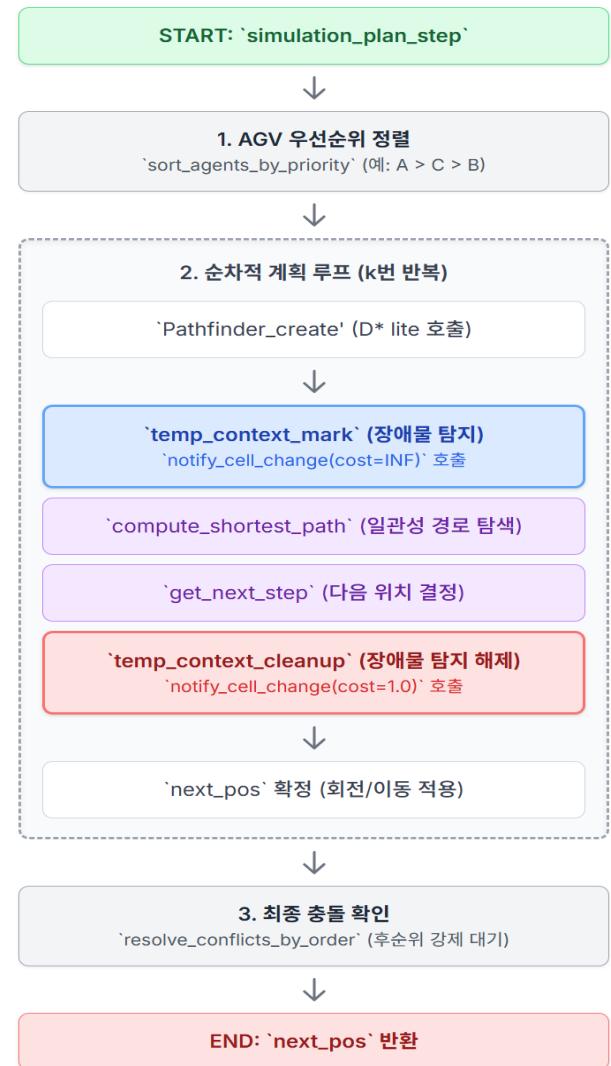
```
> function plan_dstar_lite(agents, map):
    next_pos = get_current_positions(all_agents)
    sorted_agents = sort_by_priority(agents)

    for agent in sorted_agents:
        if agent.is_waiting(rotation or action) then
            continue
        ensure_pathfinder_exists(agent)

        temp_obstacles =
            get_higher_priority_moves(agent, next_pos) +
            get_lower_priority_positions(agent, agents)

        notify_obstacles(agent.pathfinder, temp_obstacles)
        path = repair_path(agent.pathfinder, map)
        desired_move = get_next_step(path)
        cleanup_notify_obstacles(agent.pathfinder,
        temp_obstacles)
        next_pos[agent.id] = apply_rotation(agent, desired_move)

    resolve_final_conflicts(next_pos, sorted_agents)
    return next_pos
```



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : WHCA*

• What is the WHCA* algorithm?

- WHCA* [Windowed Hierarchical Cooperative A*]
 - 우선순위, 예약 테이블을 활용한 MAPF 알고리즘
- Pathfinder Algorithm
 - D* lite : Provides the Ideal route
- W : Space-Time Reservation Table
 - [time][x][y]의 3차원 예약 테이블 ($0 < \text{time} < H$) * $H = \text{Horizon}$
- H [Hierachical]
 - 우선순위에 따른 순차적 경로 계획
- C [Cooperative]
 - 타 에이전트 예약 노드 -> 장애물
- Limitations
 - 우선순위 기반 알고리즘 -> 'Risk of Deadlock'



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : WFG/SCC

- What is the WFG?

➤ WFG [Wait-For Graph]

- 현재 시스템 대기 상태 표현하는 자료구조.

➤ Role

- 교착 상태 탐지위한 자료구조

➤ Components

- $G = (V, E)$
- V [Vertex / Node]: 정점의 집합 (운용 AGV: $\{A_1, A_2, \dots\}$)
- E [Edge] : 간선의 집합 (대기 관계: $\{(A_i, A_j) \mid A_i \text{ waits for } A_j\}$)



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : WFG/SCC

- What is the SCC?

➤ SCC [Strongly Connected Components]

- 그래프 내에서 서로에게 도달 가능한 정점들의 집합
- 크기 2이상 SCC 존재 [ex. $A_i \leftrightarrow A_j$], \rightarrow 교착 상태 의미

➤ Role

- 자료구조에서 교착상태 감지

➤ DeadLock Detection

- Deadlock $\Leftrightarrow \exists$ SCC $C \subseteq V$ s.t. $|C| \geq 2$



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : CBS

- What is the CBS algorithm?

➤ CBS [Conflict – Based Search]

- 지역적 충돌 해소 위한 MAPF 알고리즘
- 소수 에이전트 간의 충돌을 효율적으로 해결함.
- 시스템의 계산 효율성을 향상.

➤ 저수준 [Low – Level] : Time A* -> 개별 경로 탐색기

- [time][x][y] 구조에서 A* 탐색 수행
- State (상태): $s = (x, y, t)$
- 평가 함수 (f-value): $f(s) = g(s) + h(s)$
- $g(s)$: $g(x, y, t) = t$ ($(x_0, y_0, 0) \rightarrow (x, y, t)$: 비용 t)
- $h(s)$: $h(x, y, t) = |x - x_{goal}| + |y - y_{goal}|$



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : CBS

- What is the CBS algorithm?

- 고수준 [High – Level] : 충돌 해결 위한 제약 조건 결정
- Cost(N) 가장 낮은 노드 N -> 우선순위 큐에서 선택

$$Cost(N) = \sum_{i \in Agents} PathCost(P_i)$$

- N: 충돌 트리의 노드 [제약 조건의 집합]
 - P_i : 노드 N의 제약 조건 & 저수준 탐색 기반 에이전트 i의 경로
 - $PathCost(P_i)$: 경로 비용
- $$\alpha \cdot \{moves\} + \beta \cdot \{waits\} + \gamma \cdot h_{\{residual\}}$$

➤ 제약 조건 (Constraint) 예시:

- High Level에서 (*Agent A*, (5, 5), $t = 3$) 제약이 내려오면,
- Low Level은 (5,5,3) 상태를 고려하지 않는다.



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : DeadLock resolve

- 1. 교착 감지

➤ WFG & SCC 활용하여 교착 상태 식별

- 2. Partical CBS

➤ SCC 식별 그룹에 Partial CBS 적용

- 3. 우선순위 기반 해결

➤ CBS 실패할 시, 우선순위 기반 해결

➤ A*, D* lite와 동일한 우선순위 기준

- 4. Pull – over

➤ Max Stuck value AGV -> Leader AGV

➤ 우선순위 실패할 시, Leader 제외 AGV 한 칸씩 후진.



[3] D* Lite + WHCA* + WFG/SCC + Partial CBS : pseudo-code

```
function plan_WHCAstar(agents, map):
    // 1. WHCA* (First attempt)
    rt = create_reservation_table()
    wf_edges = []
    sorted_agents = sort_by_priority(agents)

    for agent in sorted_agents:
        // (Plan H_steps using D* Lite)
        path = dstar_find_path(agent.pf, h_steps, rt)

        // (Build WFG - Wait-For Graph)
        if has_conflict(path) with other_agent then
            add_wait_edge(wf_edges, agent, other_agent)
            provisional_plan[agent.id] = agent.pos // wait
        else
            reserve_path(rt, path)
            provisional_plan[agent.id] = path[1] // move forward

    // 2. wfg/scc (Deadlock detection)
    sccmask = detect_scc(wf_edges)

    // 3. partial cbs (Deadlock resolution branch)
    if sccmask > 0 then
        // (Deadlock occurred)
        deadlocked_group = get_agents_from_mask(sccmask)
        // (Attempt negotiation)
        negotiated_plans, ok
            = run_partial_cbs(deadlocked_group, rt)

        if ok == true then
            // (Negotiation success: Overwrite with CBS results)
            overwrite_next_pos(provisional_plan, negotiated_plans)
        else
            // (Negotiation failure: Only leader moves)
            fallback_leader_moves(provisional_plan, deadlocked_group)
            // No deadlock. Keep the first-attempt plan
        else : (provisional_plan)

    // 4. Final
    resolve_final_conflicts(provisional_plan, sorted_agents)

    return provisional_plan
```

Simulation Map

- 형남공학관 주차장

➤ AGV: 1~4

➤ Parking : 61

- MAP2

➤ AGV: 4

➤ Parking : 390

- ## • MAP3

➤ AGV: 8

➤ Parking : 900

Simulation Result

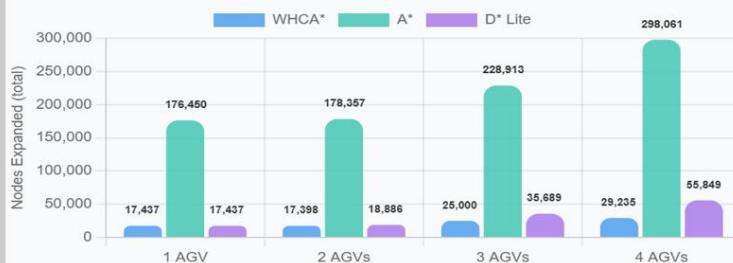


Simulation Result

MAP

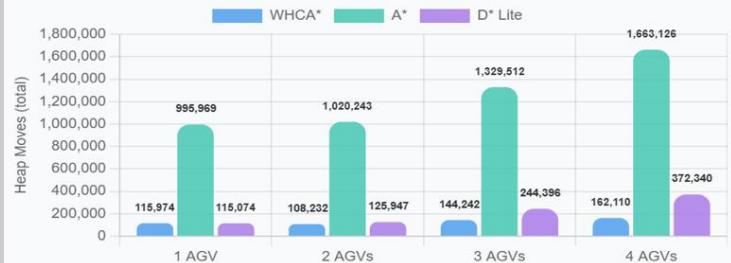
형남공학관 주차장
Parking : 61

Nodes Expanded

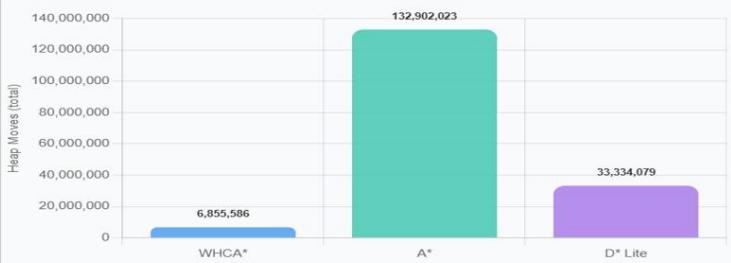
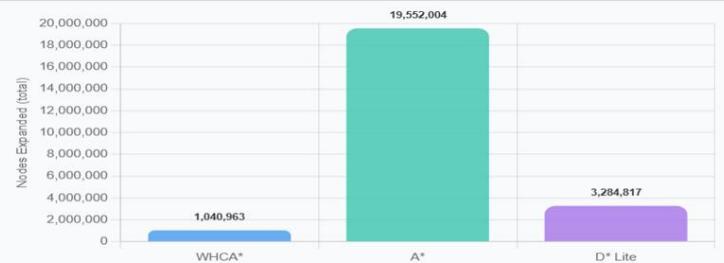
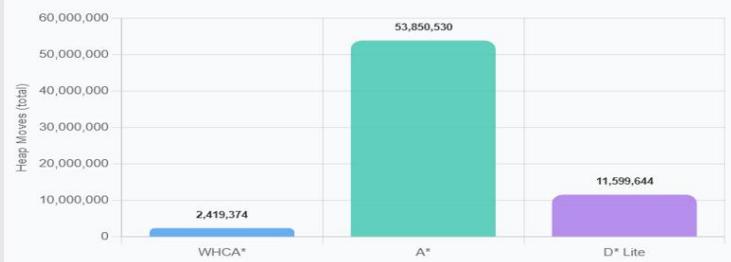
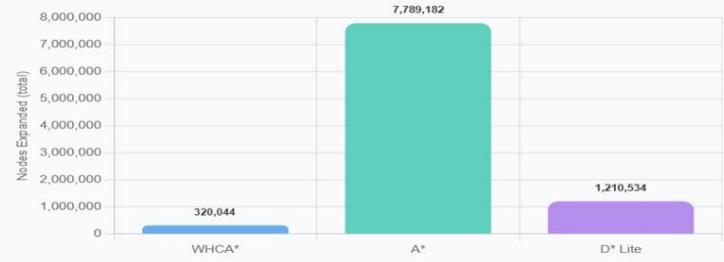


AGV : 4
Parking : 390

Heap Moves



AGV : 8
Parking : 900



Simulation Result

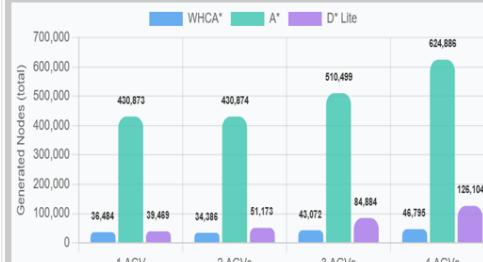
MAP

Generated Noes

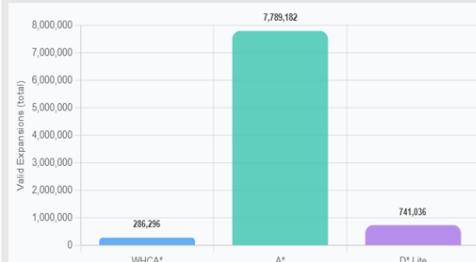
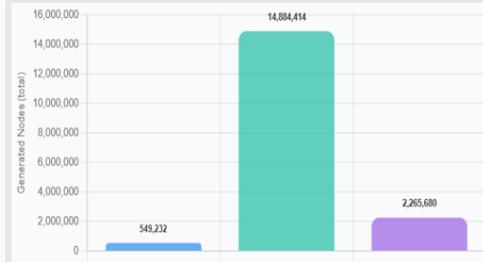
Valid Expansions

Valid Expansion Ratio

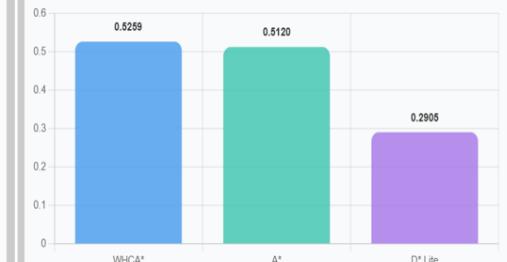
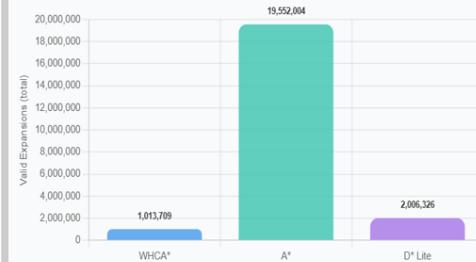
형남공학관 주차장
Parking : 61



AGV : 4
Parking : 390



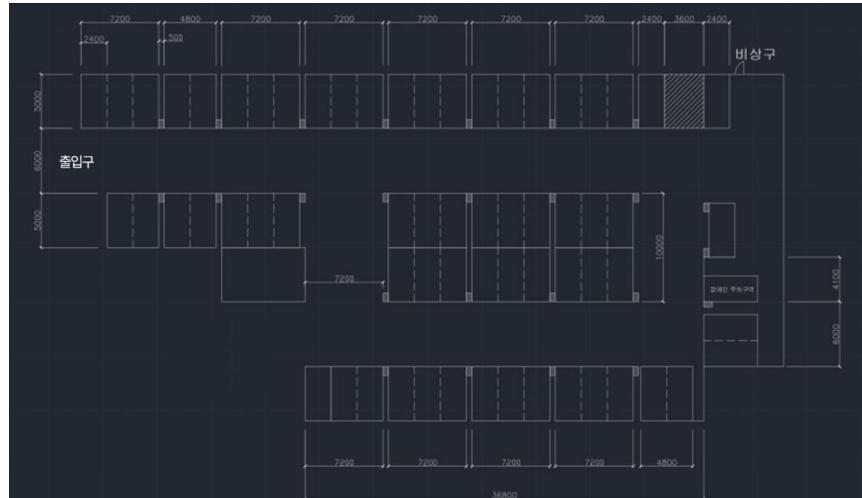
AGV : 8
Parking : 900



Conclusion

- **Which algorithm has the highest throughput in a low-density environment?**
 - A* algorithm And D* lite Algorithm
- **Which algorithm has the highest throughput in a High-density environment?**
 - WHCA* + WFG/SCC + CBS Algorithm
- **Which algorithm requires the least computation under the same conditions?**
 - 1st : WHCA* + WFG/SCC + CBS Algorithm
 - 2nd : D* Itie Algorithm
 - 3rd : A* Algorithm
- **"The optimal algorithm varies depending on the context."**

시뮬레이션 영상 [형남 공학관 주차장 + (WHCA*+WFG/SCC +CBS)]



시뮬레이션 영상 [Map2 / Map3]

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