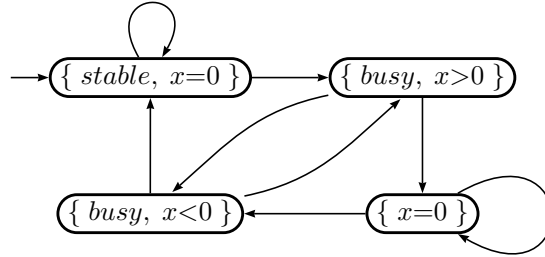


# Exercises PVL LTL

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October 31, 2018

- Below you see a Kripke structure; let's call it  $M$ . Give its explicit definition in terms of a tuple etc (see the formal definition in the slides).



- Why don't we have final states there?
  - How is the notion of 'execution' defined for a Kripke structure? And what is an 'abstract execution'?
  - Give an execution of  $M$  that satisfies the property  $\mathbf{X}(\text{busy} \mathbf{U} (x=0))$ . Does  $M$  satisfies the property?
  - So, given a property Kripke structure  $M$ , an (abstract) execution  $\Pi$ , and a property  $\phi$ , and an natural number  $i$ , what is the difference between:
    - $M \models \phi$
    - $\Pi \models \phi$
    - $\Pi, i \models \phi$
- Express the following requirements in LTL. Make the necessary assumptions if you have to; but be reasonable.
    - $P$  and  $Q$  cannot not use a resource  $r$  simultaneously.
    - Whenever  $P$  requests access to  $r$ , eventually it will get the access.
    - Whenever  $P$  requests access to  $r$ , eventually it will get the access; but only if  $P$  persists on maintaining the request.
    - $P$  cannot access  $r$  without first requesting it; and it cannot do so (make a request) without first releasing  $r$  (if it was busy using  $r$ ).
  - Construct Buchi automata representing the following LTL formulas:
    - $p \mathbf{W} q$ , where  $p, q$  are atomic propositions.
    - $\neg(x>0 \mathbf{U} x=y)$
    - $p \mathbf{U} (q \mathbf{U} r)$ , where  $p, q$  are atomic propositions.
    - $(\mathbf{X} x>0) \mathbf{U} x=y$
    - $\diamond\Box(x>0 \rightarrow x=y)$
    - $(p \mathbf{U} q) \mathbf{W} r$