

Lab 3 — Discussion

Baker's algorithm is centred around the `checkSafe()` method in our `BankImpl` class.

This safety algorithm first initialises its vectors, `available`, `maximum`, `allocation` and `need`. The space complexity from this is $O(n*m)$, where n is the number of processes and m is the number of resources.

This initialisation in the `checkSafe()` method also uses for-loops to set the values for each of the vectors, `temp_avail`, `temp_need`, `temp_allocation`, `work` and `finish`, taking a time complexity of $O(n*m)$.

The next step in this algorithm checks for an unfinished process that could be completed with the currently available resources (a.k.a `Finish[i] = false && Need[i] <= Work`). This checking process can finish with $\Omega(n)$ if all processes do not satisfy the `Need[i] <= Work` condition. However, if every process is unfinished and all the needed resources can be satisfied by the available resources, the time complexity for this step will be $O(m*m*n)$.

In the final step of `checkSafe()`, it is checked if every process is completed, taking a time complexity of $O(n)$ and $\Omega(1)$.

However, there is a case where the very first resource requested by the customer is greater than the `Need`, allowing the entire `checkSafe()` to be bypassed. This allows the best case time complexity to be reduced to $\Omega(1)$.

In conclusion, the space complexity of Banker's algorithm is $O(n*m)$, and the time complexity is $O(m*m*n)$ and $\Omega(1)$.