Disjoint list-colorings for planar graphs

Wouter Cames van Batenburg

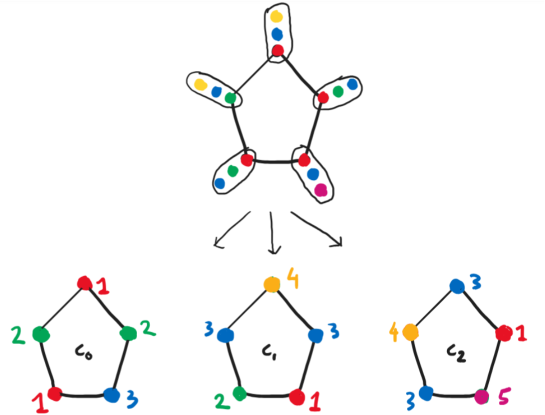
Joint work with Stijn Cambie and Xuding Zhu

Budapest, Summit280, 8 July 2024

List-coloring

Definition *(Vizing, 1976; Erdős, Rubin and Taylor, 1979)*

The **list-chromatic number** $\chi_\ell(G)$ of a graph G is the smallest integer k such that ***for every*** k**-fold list-assignment** $L : V(G) \to \binom{\mathbb{N}}{k}$, there exists an L-**coloring**, i.e. a proper vertex-coloring c s.t. $c(v) \in L(v)$ for all v.



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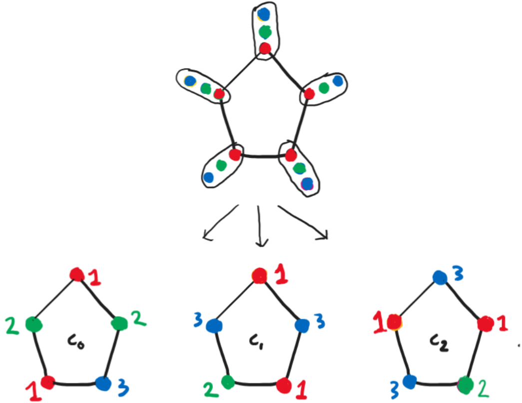

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$\chi(G)$ chromatic number of G

$\chi(G) \leq \chi_\ell(G)$ for every graph G.

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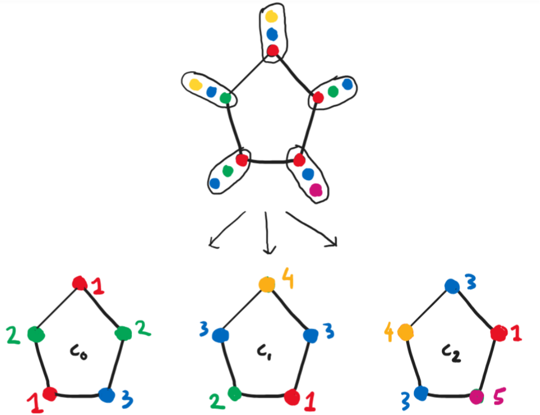

List-packing

Introduced in 2021, the **list-packing number** of a graph G has several equivalent definitions and interpretations, e.g. in terms of

- Chromatic number of certain blow-ups of G, or
- Perfect matchings of certain hypergraphs, or
- Disjoint independent transversals, or
- Disjoint list-colorings.

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This talk focuses on the list-packing number of **planar graphs**.

Plan for this talk:

Motivate list-packing (of planar graphs) from the bottom-up.

Starting from...

1. coloring
2. list-coloring
3. counting list-colorings
4. list-colorings with special requests
5. balanced probability distributions on list-colorings

...we will end up with a definition of the list-packing number, and see that it can be used to strengthen some of the literature on the above concepts.

19th and 20th century: coloring planar graphs

$\chi(G)$ the chromatic number of a graph G.

Theorem *(Appel and Haken, 1977; Grötzsch, 1959)*

For G planar:

$$\chi(G) \leq \begin{cases} 4 & \\ 3 & \text{if } G \text{ triangle-free.} \end{cases}$$

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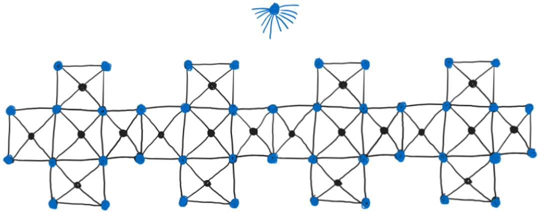

1990s: List-coloring planar graphs

For coloring...

Theorem *(Appel and Haken, 1977; Grötzsch, 1959)*

For G planar, the optimal bounds are:

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But for list-coloring...

Theorem *(Thomassen, 1994, 1995; Voigt, 1993, 1995; Mirzakhani, 1996)*

For G planar, the optimal bounds are:

$$\chi_\ell(G) \leq \begin{cases} 5 & \\ 4 & \text{if } G \text{ triangle-free} \\ 3 & \text{if } G \text{ girth} \geq 5. \end{cases}$$

2000s: Exponentially many L-colorings

Results on previous slide guarantee existence of at least *one* L-coloring. In fact there exist exponentially many, i.e. $\geq c^{\# V(G)}$ for some uniform $c > 1$.

Theorem *(Thomassen, 2007; Kelly and Postle, 2008)*

For G planar, a k-fold list-assignment L admits exponentially many L-colorings in each of the following cases:

$$\begin{cases} k = 5 & \\ k = 4 & \text{and } G \text{ triangle-free} \\ k = 3 & \text{and } G \text{ girth} \geq 5. \end{cases}$$

2010s and 2020s: Flexible list-colorings

Since there exist many L-colorings, can we guarantee a very nice one?

Suppose each vertex v **requests a preferred color** $R(v)$ from its list. Does there exist an L-coloring that respects a large fraction of the requests?

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Definition (Dvořák, Norin and Postle, 2019)

Graph G is ϵ-**flexible** wrt list-assignment L if for every collection of requests $(R(v) \in L(v))_{v \in V(G)}$, there exists an L-coloring c s.t.

$$c(v) = R(v)$$

for at least $\epsilon \cdot \# V(G)$ of the vertices v.

The actual definition is a bit more involved; see the paper

Example

In special case that all vertices have the **same list** $L(v) = [k]$, it easily follows that G is $\frac{1}{k}$-flexible wrt L. (Provided $k \geq \chi_\ell(G)$)

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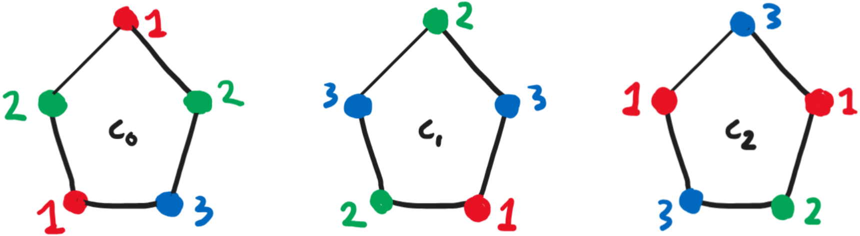
Proof sketch

Fix a k-coloring c_0, and cyclically permute it to obtain k colorings $c_0, \dots, c_{k-1}$. By pigeon hole, at least one of them satisfies $\geq \frac{1}{k} \cdot \# V(G)$ requests. □



Example with $k = 3$. Always exists coloring satisfying $\geq \lceil \frac{5}{3} \rceil = 2$ requests.

Stronger property: weighted ϵ-flexible

Definition (Dvořák, Norin and Postle, 2019)

Graph G is **weighted ϵ-flexible** wrt list-assignment L if there exists a probability distribution on L-colorings c s.t. $\forall v \in V(G), \forall x \in L(v)$:

$$\mathbb{P}(c(v) = x) \geq \epsilon.$$

Fact 1: **weighted ϵ-flexible** implies **ϵ-flexible**.

Fact 2: wrt a k-fold L, the **highest value we can hope for is $\epsilon = \frac{1}{k}$**.

Theorem *(Dvořák, Norin and Postle, 2019; Dvořák, Masařik, Musílek, Prangrác, 2020 and 2021*; Bi and Bradshaw, 2023)

A planar G is (weighted) ϵ-flexible wrt all k-fold list-assignments L for:

$$\begin{cases} k = 7 & \\ k = 4 & \text{if } G \text{ triangle-free} \\ k = 3 & \text{if } G \text{ girth} \geq 6. \end{cases}$$

For respectively $\epsilon = 7^{-36}, \epsilon = 2^{-186}$ and $\epsilon = 2^{-30}$.

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Theorem *(Cambie, CvB., Zhu, 2023+)*

A planar G is (weighted) $\frac{1}{k}$-flexible wrt all k-fold list-assignments L for:

$$\begin{cases} k = 8 & \\ k = 5 & \text{if } G \text{ triangle-free} \\ k = 4 & \text{if } G \text{ girth} \geq 5 \text{ (=optimal)} \\ k = 3 & \text{if } G \text{ girth} \geq 6. \text{ (=optimal)} \end{cases}$$

So... what about the title of this talk? Why disjoint list-colorings?

Well, for a k-fold L...

k disjoint L-colorings

implies

weighted $\frac{1}{k}$-flexible.

Disjoint list-colorings

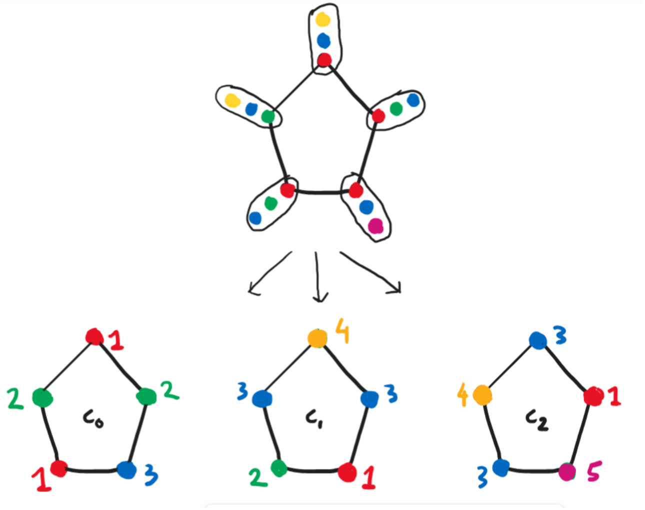
Definition (Cambie, CvB., Davies, Kang, 2021)

The **list-packing number** $\chi_\ell^\star(G)$ of a graph G is the smallest integer k such that for every k-fold list-assignment $L : V(G) \to \binom{\mathbb{N}}{k}$, there are k **disjoint** L**-colourings** $c_1, \dots, c_k$.

I.e.: for every vertex v and every color $x \in L(v)$, there is precisely one $i \in [k]$ such that $c_i(v) = x$.

Note: $\chi_\ell(G) \leq \chi_\ell^\star(G)$, since at least one L-colouring required.

Example: $\chi_\ell^\star = 3$ for cycles



Top: a 3-fold L for C_5. Bottom: three disjoint L-colorings.

From list-packing to flexibility

Observation

If $\chi_\ell^\star(G) \leq k$ then G is weighted $\frac{1}{k}$-flexible w.r.t every k-fold L-assignment.

Proof

Let $c_1, \ldots, c_k$ be k disjoint L-colorings. Among them, choose a uniformly random coloring. Then at every vertex v, every color $x \in L(v)$ has equal probability $\frac{1}{k}$ of being assigned to v. □

Thus... our results on flexibility follow from the theorem on the next slide.

Theorem (Cambie, C., Zhu 23+, and Cranston, Smith-Roberge 24+)

For every planar graph G:

$$\chi_\ell^\star(G) \leq \begin{cases} 8 & \\ 5 & \text{if triangle-free.} \\ 4 & \text{if girth at least five. } \text{(4=optimal, even for larger girth!)} \end{cases}$$

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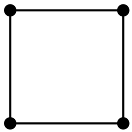
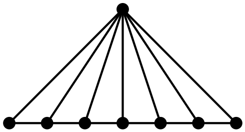
Questions

- *Does there exist a planar graph G with $\chi_\ell^\star(G) \ge 6$?*
- *Does there exist a triangle-free planar graph G with $\chi_\ell^\star(G) = 5$?*



$\chi_\ell = 2 < 3 = \chi_\ell^\star$



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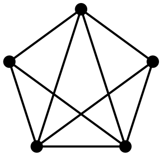


$\chi_\ell = 4 = \chi_\ell^\star$

Going mad

Definition maximum average degree

mad(G):= max$\{$average degree of $H \mid H$ subgraph of $G\}$

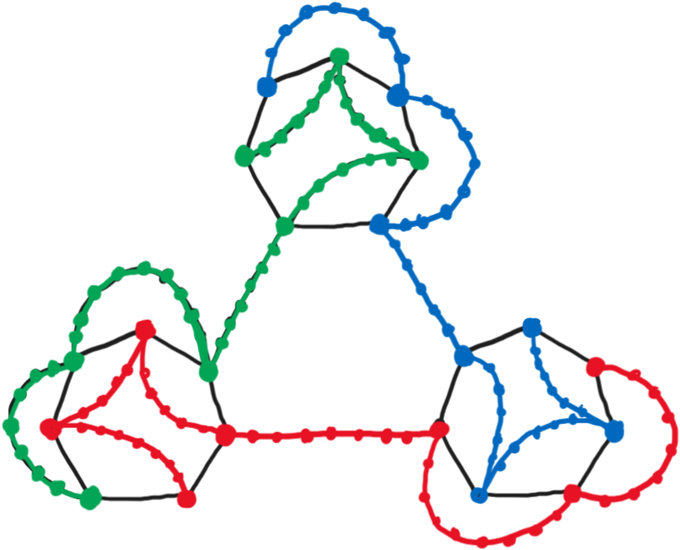
Planar graphs satisfy mad$(G) < 6$.
In this way, our upper bounds for planar graphs follow from

Theorem (Cambie, CvB, Zhu 23+)

For every graph G,

$$\chi_\ell^\star(G) \le \begin{cases} 8 & \text{if } \mathrm{mad}(G) < 6 \\ 5 & \text{if } \mathrm{mad}(G) < 4 \\ 4 & \text{if } \mathrm{mad}(G) < 10/3 \end{cases}$$

Planar graph with arbitrarily large girth, yet $\chi_\ell^* > 3$



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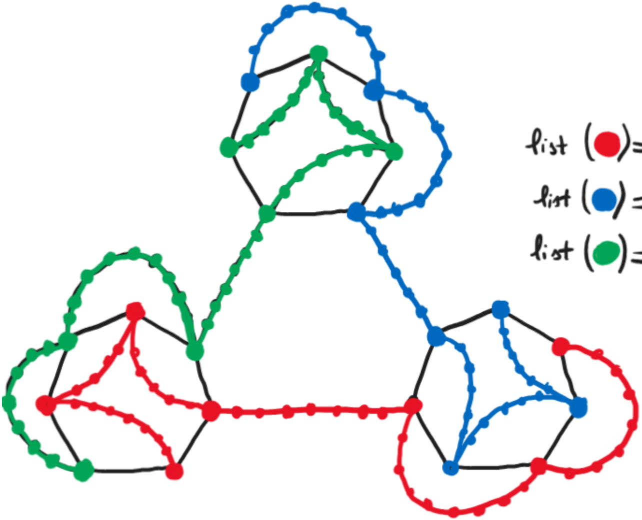

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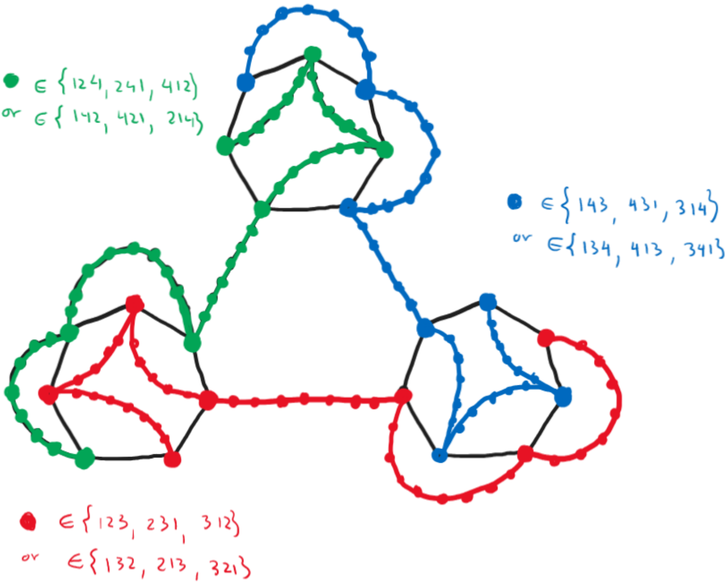

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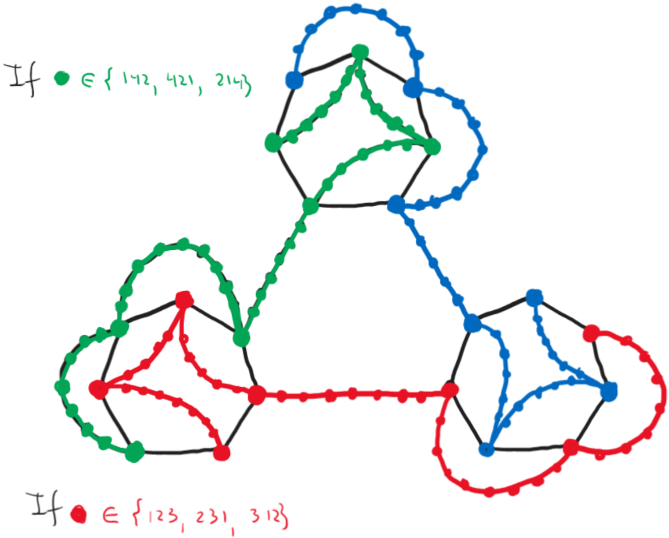

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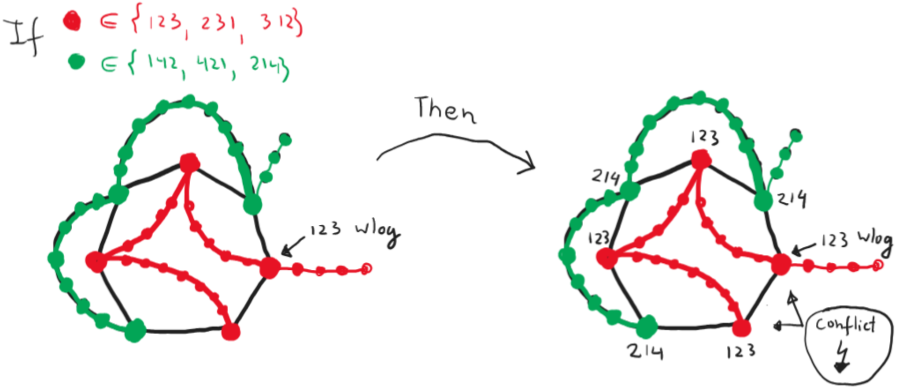

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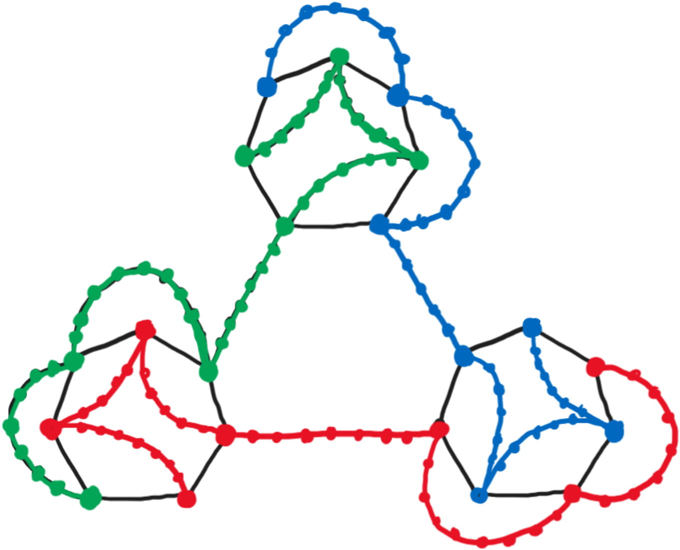

Planar girth ≥ 6 graphs

Construction on previous slide shows that $\chi_\ell^\star \leq 3$ does NOT hold for every planar large girth graph. Despite that, we still have...

Theorem (Cambie, CvB, Zhu, 2023+)

Every planar girth ≥ 6 graph is weighted $\frac{1}{3}$-flexible wrt every 3-fold L.

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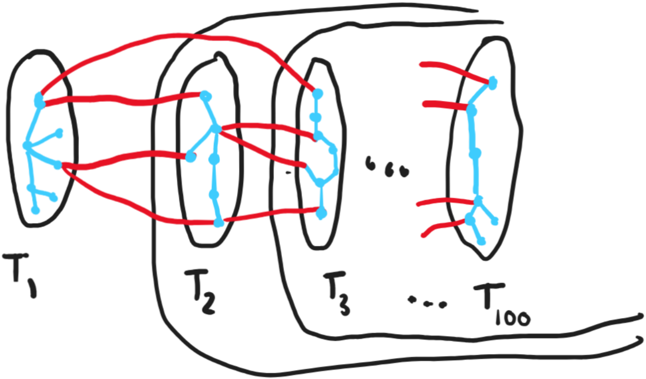
Theorem (Cambie, CvB, Zhu, 2023+)

Every planar girth ≥ 6 graph is weighted $\frac{1}{3}$-flexible wrt every 3-fold L.

Proof sketch

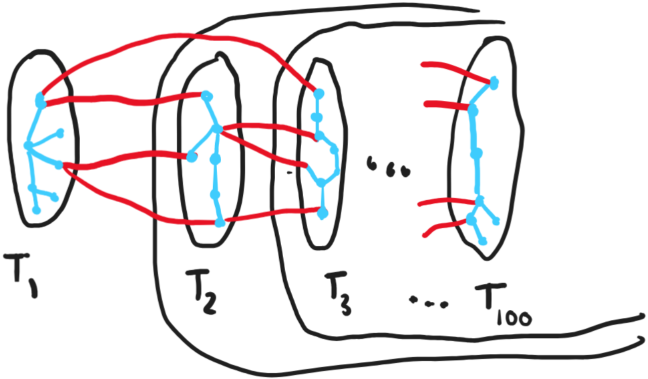
- By a Technical Lemma, it suffices to prove that G contains an induced subtree T of which every vertex has at most one neighbour in $G - V(T)$.
- Euler's formula and a subtle global discharging argument yield T. $\square$

Technical lemma; a tree-layering



Suppose G has a layering into **induced subtrees** $T_1, T_2, \ldots$, such that each vertex in T_i has *at most one neighbour* in the layers $T_1, \ldots, T_{i-1}$ to its left. Then the whole graph is **weighted $\frac{1}{3}$-flexible**.

Technical lemma; a tree-layering



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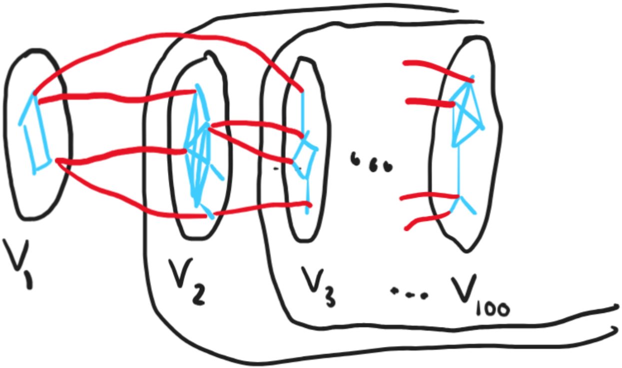
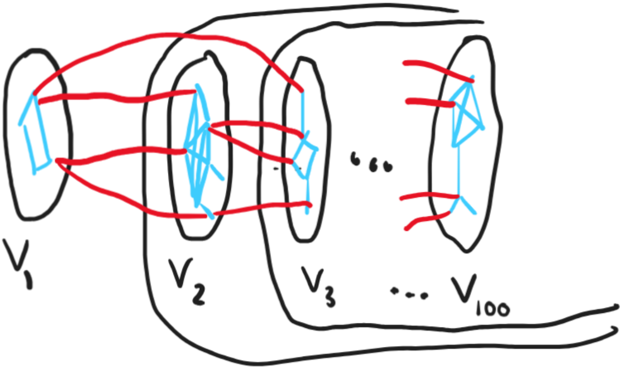
This works because trees are weighted $\frac{1}{2}$-flexible.

Technical lemma; more general



A layering of the vertices of G. If each layer V_i is **weighted $\frac{1}{k-1}$-flexible**, and each vertex in V_i has *at most one neighbour* in the layers $V_1, \ldots, V_{i-1}$ to its left, then the whole graph is **weighted $\frac{1}{k}$-flexible**.

Technical lemma; more general



More applications?

Any graph class with a nice layered structure;

Cartesian products, graphs with bounded treedepth, . . .

Summary

List-packing lies at the base of a sequence of implications.

$$
\begin{aligned}
\chi_\ell^\star(G) \le k \;\Leftrightarrow\;& k \text{ disjoint } L\text{-colorings}_{\text{ wrt every } k\text{-fold } L}\\
\Rightarrow\;& \text{weighted } \frac{1}{k}\text{-flexible}_{\text{ wrt every } k\text{-fold } L}\\
\Rightarrow\;& \frac{1}{k}\text{-flexible}_{\text{ wrt every } k\text{-fold } L}\\
\Rightarrow\;& L\text{-coloring}_{\text{ wrt every } k\text{-fold } L} && \Leftrightarrow \chi_\ell(G) \le k\\
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&\qquad\qquad\qquad\qquad\qquad\qquad\quad \Rightarrow \chi(G) \le k.
\end{aligned}
$$

- Through bounding $\chi^\star_\ell$, we improved results on weighted ϵ-flexibility, reaching the optimal value $\epsilon = \frac{1}{k}$.
- In some cases we directly proved weighted $\frac{1}{k}$-flexibility, via a graph layering argument.

Many open problems

What is the best-possible upper bound on $\chi_\ell^\star(G)$ if G is ...

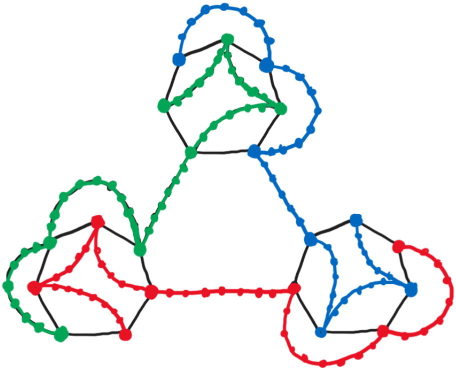
- Planar. $5, 6, 7$ or 8?
- Planar triangle-free. 4 or 5?
- Planar bipartite. $3, 4$ or 5?
- Bounded treewidth?
- K_t-minor-free?

Main open problem on listpacking:

$$\chi_\ell^\star(G) \leq C \cdot \chi_\ell(G) \quad ?$$

Thank you!

Slides are at woutercvb.github.io



Cambie, S., Cames van Batenburg, W. and Zhu, X., Disjoint list-colorings for planar graphs, arxiv:2312.17233

Technical lemma- exact statement of special case

The following holds for *every* graph G:

Key technical lemma

Let $k \geq 2$. If there is an induced subgraph T of G s.t.

1. Every vertex of T has at most one neighbour in $G - V(T)$;
2. T is weighted $\frac{1}{k-1}$-flexible;
3. $G - V(T)$ is weighted $\frac{1}{k}$-flexible;

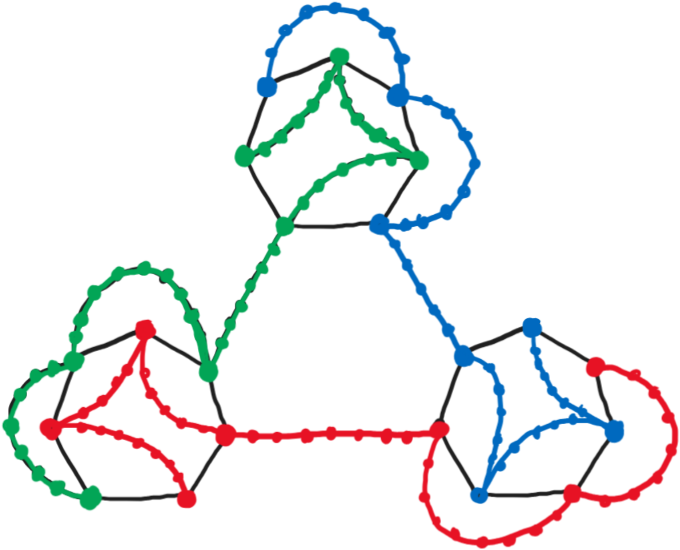
Then G is weighted $\frac{1}{k}$-flexible.

Mad question

Question

Let $k \in \mathbb{N}$. Does $\mathrm{mad}(G) < k$ imply $\chi_\ell^\star(G) \leq k + 1$?

Planar graph with arbitrarily large girth, yet $\chi_\ell^* > 3$



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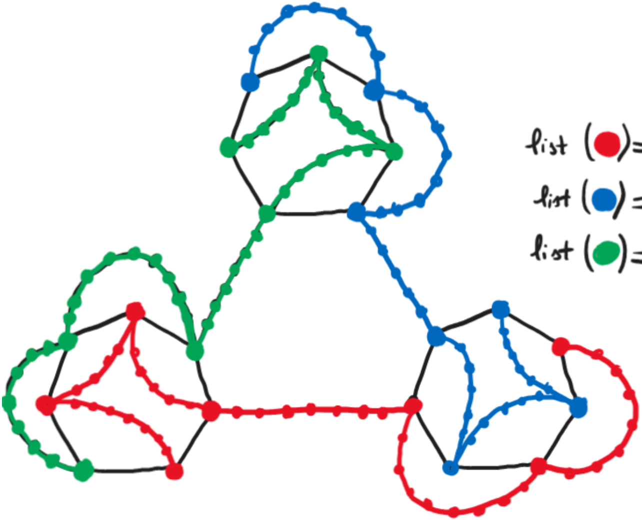

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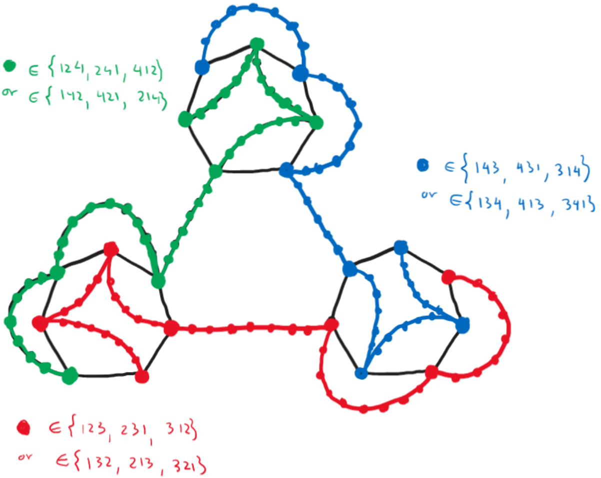

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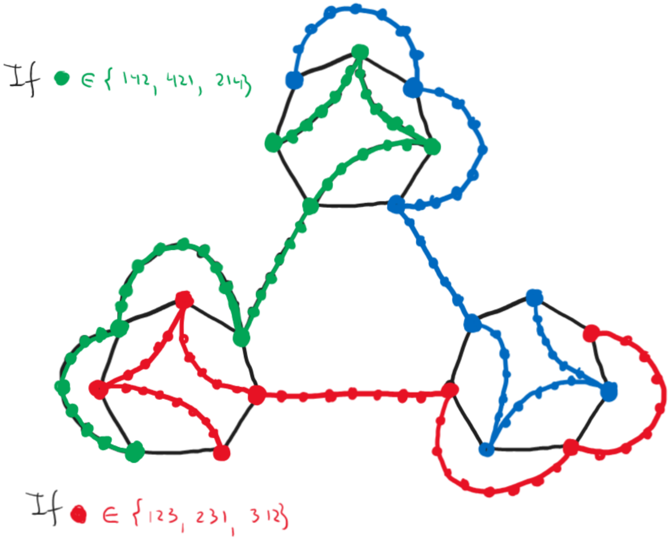

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