CISC 372: Parallel Computing

MPI Collectives

Stephen F. Siegel

Department of Computer and Information Sciences University of Delaware

The collective model of computation

- a collective operation is invoked by all processes in a communicator
- some processes contribute data, some receive data, some operations may be performed
- collective operations capture many commonly used parallel patterns
 - one process broadcasts data to all
 - add up values across processes and return the result to one process
 - gather data from all processes into one big array on one process
- many parallel algorithms can be expressed most easily using collectives
- collective programs are easier to understand than point-to-point
 - ▶ a higher-level, abstract, "big step" view of an algorithm
 - ▶ all procs do this, then all procs do that, then all procs do. . . .
 - vou can reason about them much like a sequential program
- ▶ MPI implementations have very optimized implementations of collectives
 - these can be guite complicated, but the implementation details are hidden from the user

MPI_Reduce

```
MPI_Reduce(sendbuf, recvbuf, count, datatype, op, root, comm)
```

```
sendbuf address of send buffer (void*)
recvbuf address of recv buffer (root only, void*)
  count number of elements in send buffer (int)
datatype data type of elements in send buffer (MPI_Datatype)
  op reduce operation (MPI_Op)
  root rank of root process (int)
  comm communicator (MPI_Comm)
```

Rank 0 sendbuf Rank 1 sendbuf Rank 2 sendbuf

<i>x</i> ₀₀	<i>x</i> ₀₁	X ₀₂
<i>x</i> ₁₀	<i>x</i> ₁₁	<i>x</i> ₁₂
<i>x</i> ₂₀	<i>x</i> ₂₁	X ₂₂

Root recybuf

 $|x_{00} + x_{10} + x_{20}| x_{01} + x_{11} + x_{21}| x_{02} + x_{12} + x_{22}$



Creating global synchronization points: MPI Barrier

```
MPI_Barrier(comm)
```

```
communicator (MPI_Comm)
comm
```

- blocks calling process until all processes in comm call MPI_Barrier
- ▶ if one process in comm calls MPI_Barrier(comm), all should
 - else deadlock ensues
- explicit barriers are rarely needed; some exceptions. . .
 - good practice: use barriers before calling MPI_Wtime
 - ensures all processes have reached that point
 - programs that use MPI_ANY_SOURCE (coming soon)
 - barriers may be necessary to control how sends are matched with received
- using barriers to control order of printing from different procs is not reliable
 - sometimes it works, sometimes it doesn't.

Broadcast: MPI Bcast

```
MPI_Bcast(buffer, count, datatype, root, comm)
```

```
address of buffer (void*)
  buffer
           number of elements in buffer (int)
   count
           data type of elements in buffer (MPI_Datatype)
datatype
           rank of root process (int)
    root
           communicator (MPI_Comm)
    comm
```

- ▶ broadcasts a message from a single process (root) to all other processes in comm
- on root, buffer acts as a send buffer; on non-root procs, buffer acts as a receive buffer
- after return, buffer will contain the same data as that on root
- is a barrier (global synchronization point) induced?
 - all procs must enter before any proc can exit? No
- see bcast.c

MPI Allreduce

```
MPI_Allreduce(sendbuf, recvbuf, count, datatype, op, comm)
             address of send buffer (void*)
   sendbuf
             address of recv buffer (void*)
   recvbuf
             number of elements in send buffer (int)
     count
             data type of elements in send buffer (MPI_Datatype)
  datatype
             reduce operation (MPI_Op)
         qo
             communicator (MPI_Comm)
       comm
```

- just like MPI_Reduce, but no root
- instead, result is returned to all processes in comm
- equivalent to MPI_Reduce followed by MPI_Bcast
- barrier? yes

MPI Scatter

```
MPI_Scatter(sendbuf, sendcount, sendtype,
             recvbuf, recvcount, recvtype, root, comm)
              address of send buffer (root only, void*)
   sendbuf
              num. elements sent to each proc (root, int)
 sendcount
             data type of send buf. elements (root, MPI_Datatype)
  sendtype
             address of receive buffer (void*)
   recybuf
             number of elements in recv buffer (int)
 recvcount
             type of data to receive (MPI_Datatype)
  recvtype
              rank of sending process (int)
      root
             communicator (MPI_Comm)
       comm
```

- similar to broadcast: the root sends, everyone else receives
- but: root sends a different block of sendbuf to each proc
 - rank 0 gets the first sendcount elements, rank 1 gets the next sendcount elements . . .
- number of elements in sendbuf is nprocs*sendcount; see scatter.c

MPI Scattery

```
MPI_Scatterv(sendbuf, sendcounts, displs, sendtype,
              recvbuf, recvcount, recvtype, root, comm)
              address of send buffer (root only, void*)
    sendbuf
              num. elements sent to each proc (root only, int[nprocs])
 sendcounts
              displacements for each proc (root only, int[nprocs])
     displs
              data type of send buf. elements (root only, MPI_Datatype)
   sendtype
    recvbuf
              address of receive buffer (void*)
              number of elements in recv buffer (int)
  recvcount
              type of data to receive (MPI_Datatype)
   recvtype
              rank of sending process (int)
        root
              communicator (MPI_Comm)
        comm
```

- generalizes MPI_Scatter: the amount of data sent to each process can vary
- \triangleright sendcounts[i] = number of elements to send to proc i
- displs[i] = offset of send buffer for proc i relative to sendbuf

MPI_Gather

```
MPI_Gather(sendbuf, sendcount, sendtype,
            recvbuf, recvcount, recvtype, root, comm)
              address of send buffer (void*)
   sendbuf
              num. elements to send (int)
 sendcount
             data type of send buf. elements (MPI_Datatype)
  sendtype
              address of receive buffer (root only, void*)
   recybuf
              number of elements to recv from each proc (root, int)
 recvcount
             type of data to receive (root, MPI_Datatype)
  recvtype
              rank of receiving process (int)
      root
              communicator (MPI_Comm)
       comm
inverse of MPI_Scatter: everyone sends to root, root receives from everyone
```

- inverse of MP1_Scatter. everyone sends to root, root receives from ever
- root receives into a different block of recybuf for each proc
 - rank 0's message goes into first recvcount elements
 - rank 1's message goes into next recvcount elements ...
- number of elements in recvbuf is nprocs*recvcount

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MPI_Allgather MPI_Allgather(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtvpe, comm) address of send buffer (void*) sendbuf num. elements to send (int) sendcount data type of send buf. elements (MPI_Datatype) sendtype address of receive buffer (void*) recybuf number of elements to recv from each proc (int) recvcount type of data to receive (MPI_Datatype) recvtype communicator (MPI_Comm) comm

like MPI_Gather done once for each proc

MPI Gatherv MPI_Gatherv(sendbuf, sendcount, sendtype, recvbuf, recvcounts, displs, recvtype, root, comm) address of send buffer (void*) sendbuf num. elements to send (int) sendcount data type of send buf. elements (MPI_Datatype) sendtype address of receive buffer (root, void*) recybuf number of elements to recv from each proc (root, int[nprocs]) recvcounts displacements of receive buffers (root, int[nprocs]) displs type of data to receive (root, MPI_Datatype) recvtype rank of receiving process (int) root communicator (MPI_Comm) comm generalizes MPI_Gather the amount sent by each process can vary

- see gatherv.c
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MPI Alltoall

```
MPI_Alltoall(sendbuf, sendcount, sendtype,
              recvbuf, recvcount, recvtype, comm)
             address of send buffer (void*)
   sendbuf
             num. elements to send (int)
 sendcount
             data type of send buf. elements (MPI_Datatype)
  sendtype
             address of receive buffer (void*)
   recybuf
             number of elements to receive on each proc (int)
 recvcount
             type of element to receive (MPI_Datatype)
  recvtype
             communicator (MPI_Comm)
       comm
every process sends distinct buffers to all others
```

- amount of data sent to each process is the same
- symmetric (no root)

MPI Alltoally MPI_Alltoallv(sendbuf, sendcounts, sdispls, sendtype, recvbuf, recvcounts, rdispls, recvtype, comm) address of send buffer (void*) sendbuf num. elements to send to others (int[nprocs]) sendcounts displacements of send buffers (int[nprocs]) sdispls data type of send buf. elements (MPI_Datatype) sendtype recybuf address of receive buffer (void*) number of elements to receive from others (int[nprocs]) recvcounts displacements of receive buffers (int[nprocs]) rdispls type of element to receive (MPI_Datatype) recutype communicator (MPI Comm) comm generalizes MPI_Alltoall

- ► see emily.c
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