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(a) Answer: $(n+1)/2n$

$$1 - 1/k^2 = (k^2 - 1)/k^2 = (k-1)(k+1)/k^2 = (1 \cdot 3/2^2)(2 \cdot 4/3^2)(3 \cdot 5/4^2) \dots (n-1)(n+1) = (n+1)/2n$$

(b) Answer: 4

$$3^1 \bmod 7 = 3, 3^2 \bmod 7 = 2, 3^3 \bmod 7 = 6, \text{ and then } 3^6 \bmod 7 = 6^2 \bmod 7 = (-1)^2 \bmod 7 = 1; 3^{1000} = (3^6)^{166} (3^4),$$

$$3^{1000} \bmod 7 = (-1^3)(-1^2)^{166} = -3 \bmod 7 = 4$$

(c) Answer: 1

$$\frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \frac{1}{16} + \dots = 1$$

(d) Answer: 2

$$(\lg 81 / \lg 7) / (\lg 7 / \lg 9) = (\lg 81 / \lg 9) = \log_9 81 = 2$$

(e) Answer: $4n$

$$\log_2 2^{(2 \cdot 2n)} = \log_2 2^{4n} = 4n$$

(f) Answer: 1

$$\log_{17} (221/13) = 1$$

4

Answer: $a_n = (n+1)!$

$$A_1 = 2 = 2!, A_2 = 6 = 3!, A_3 = 24 = 4! \dots a_n = (n+1)!$$

Proof:

Basic case: for $n=0$, $A_0 = (0+1)! = 1!$

Induced case: As $a_n = (n+1)!$ We need to prove $a_{(n+1)} = (n+2)!$ By the definition formula, we have $a_{(n+1)} = a_n + (n+1)!(n+1)$ Thus $a_n = (n+1)!$

$A_n + (n+1)!(n+1) = (n+1)! + (n+1)!(n+1) = (n+1)!(1+n+1) = (n+1)!(n+2) = (n+2)!$, so $a_{(n+1)} = (n+2)!$, provided $a_n = (n+1)!$

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(a) Answer: $f(n) = \theta(g(n))$

Proof: As $f(n) = n$ and $g(n) = 2n+1$ are both linear, there exists a c such $c \cdot g(n)$ could be both the upper bound and the lower bound of $f(n)$ for larger n . Thus $f(n) = \theta(g(n))$.

(b) Answer: $f(n) = \omega(g(n))$

Proof: $g(n) = \sqrt{n} = n^{1/2}$. Because $g(n)$ has lower dimension than $f(n)$, for large n , $c \cdot g(n)$ could only be the lower bound of $f(n)$. Thus $f(n) = \omega(g(n))$.

(c) Answer: $f(n) = O(g(n))$

(d) Answer: $f(n) = O(g(n))$

6 (a) Answer: $T(n) = 1 + 5(\log_2 n)$

As $T(2) = T(1) + 5 = 6$, $T(4) = T(2) + 5 = 11$, $T(8) = T(4) + 5 = 16$. $T(16) = T(8) + 5 = 21$, so $T(n) = 1 + 5(\log_2 n)$

(b)

Answer: $T(n) = \text{sum of } 1/n \text{ (n from 1 to n)}$

$T(1) = T(0) + 1 = 1$, $T(2) = T(1) + 1/2$, $T(3) = T(2) + 1/3$ Thus as the answer showed.

H_n (harmonic series up to n)

(c) Answer: Basic case: For $n=1$, $T(n)=1$ and $T(1) = 1 + 5(\log_2 1) = 1$ as showed.

Induced case: As $T(n) = 1 + 5(\log_2 n)$, we need to prove $T(2n) = 1 + 5(\log_2 (2n))$.

$T(2n) = 1 + 5(\log_2 (2n)) = 1 + (5(\log_2 2) + 5(\log_2 n)) = 1 + 5 + 5(\log_2 n) = 6 + 5(\log_2 n)$

By the definition, $T(2n) = T(n) + 5 = 1 + 5(\log_2 n) + 5 = 6 + 5(\log_2 n)$

Thus we get $T(2n) = 1 + 5(\log_2 n)$

Q.E.D

7 (a) formula: $T(n) = T(n/2) + c_2$

Base case $T(1) = c_1$

Solution: $c(\log_2 n) + c_1$

(b) formula: $T(n) = 2T(n/2) + c_2n + c_3$

Base case: $T(1) = c_1$

Solution: $n \lg n$

8 (a) Answer: process: $(2, 12) - (4, 6) - (16, 3) - 16(256, 1) - 256(256^2, 0) - 1$

So the result is $256 * 16 * 1$.

(b) Try to find the n th power of x . For instance, if $x=2$, $n=12$, it will find out 2^{12} .

(c) $T(0) = d$

$T(n) = T([n/2]) + c$

(d) $T(n) = c * (\log_2 n) + d$

