

# **ECE 550D**

## **Fundamentals of Computer Systems and Engineering**

### **Fall 2023**

## Combinational Logic

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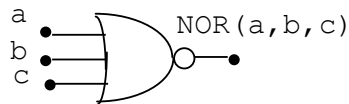
Slides are derived from work by  
Andrew Hilton, Tyler Bletsch and Rabi Yunes (Duke)

## **Last time....**

- Who can remind us what we talked about last time?
  - Electric circuit basics
    - $V_{cc} = 1$
    - Ground = 0
  - Transistors
    - PMOS
    - NMOS
  - Gates
    - Complementary PMOS + NMOS
    - Output is logical function of inputs

## Multi-input gates

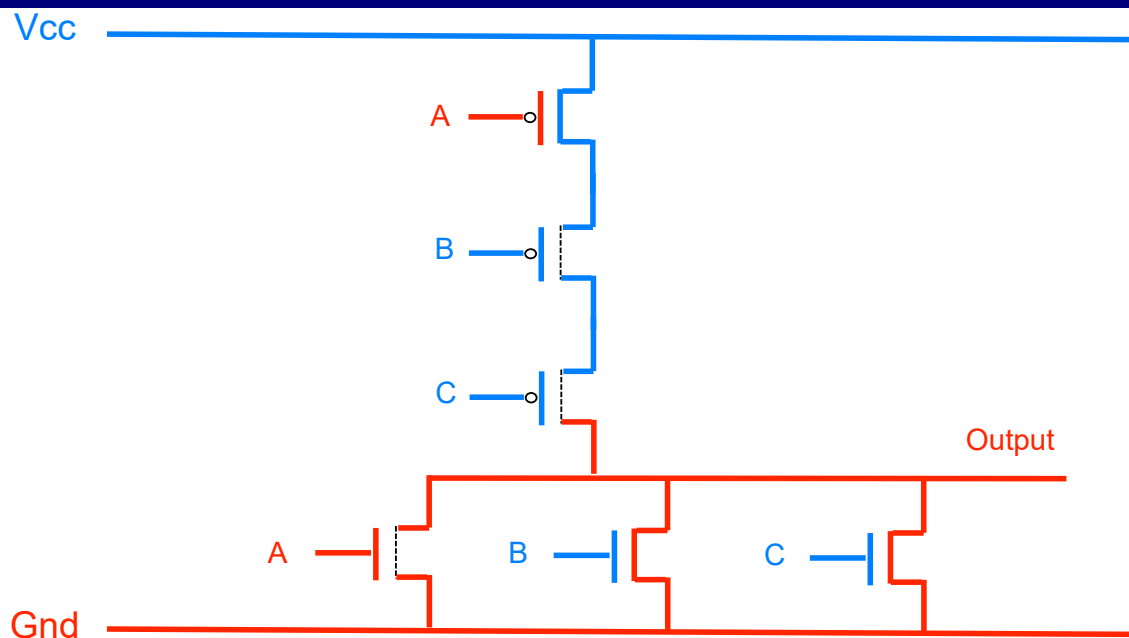
- So far gates have had 1 or 2 inputs
  - Can have more, though typically stop at 3 or 4
  - Symbols stay the same, just have more input lines



A	B	C	Ouput
0	0	0	1
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

3

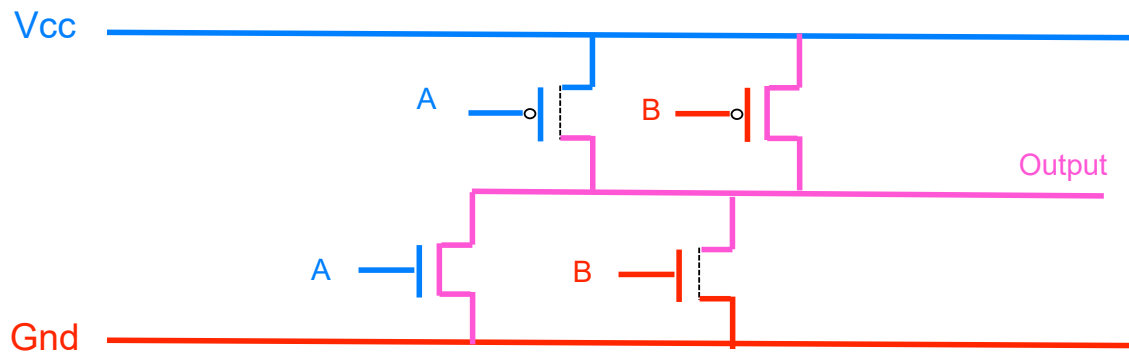
## Three input NOR Gate



- Similar to two input, more transistors
  - Slightly slower

4

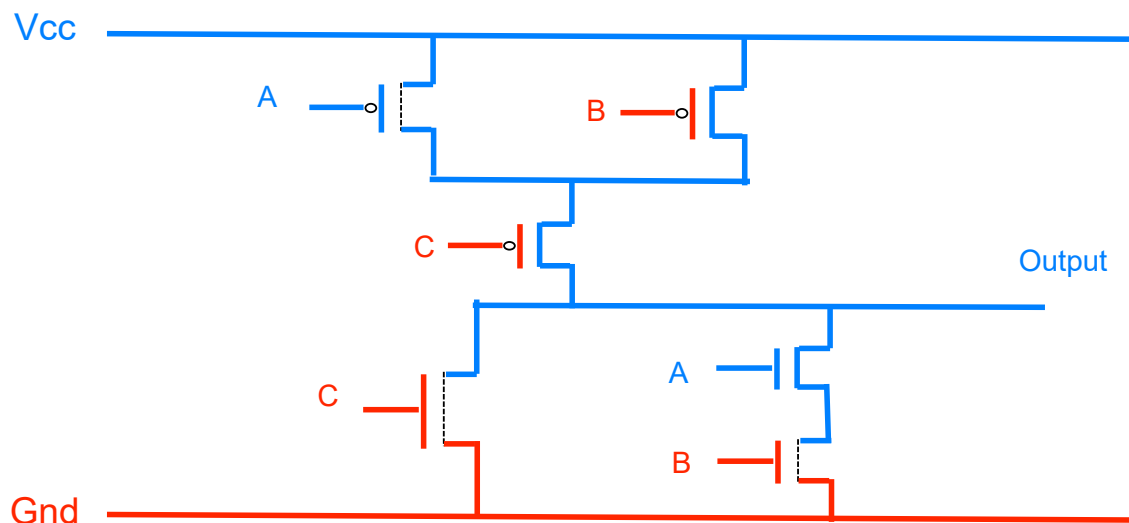
## Complementary: very important



- Complementary nature: very important
  - Without it, we have a problem
  - Here: both PMOS and NMOS in parallel...
  - $A=1, B=0$  (or  $A=0, B=1$ ) forms short-circuit
    - Chip catches fire ☹️

5

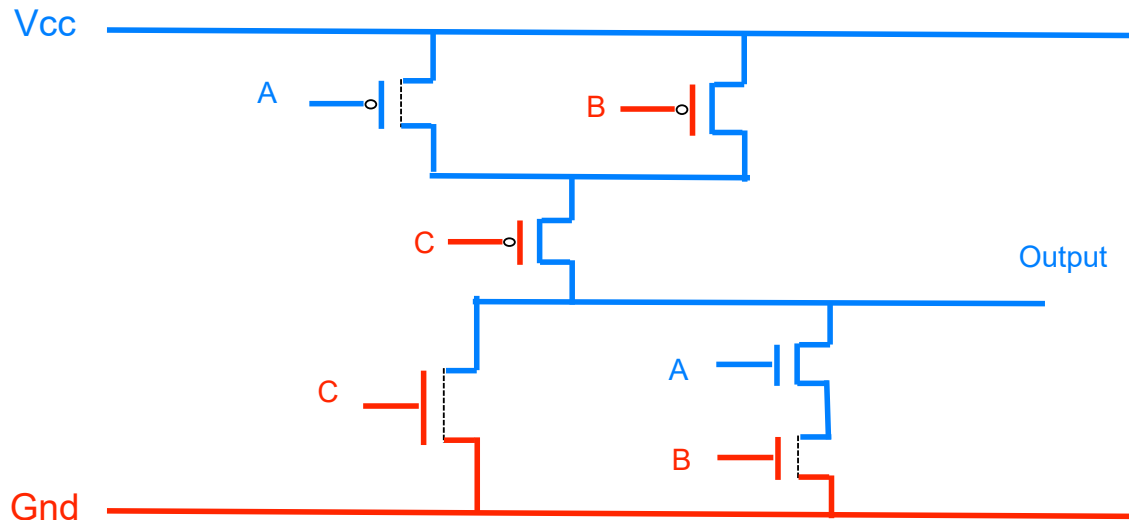
## Complementary: very important



- With more than 2 inputs, can get very complicated
  - Are the PMOS and NMOS complementary here?
  - We can go the other way: transistors  $\rightarrow$  formulas
  - Check if formulas logically equivalent

6

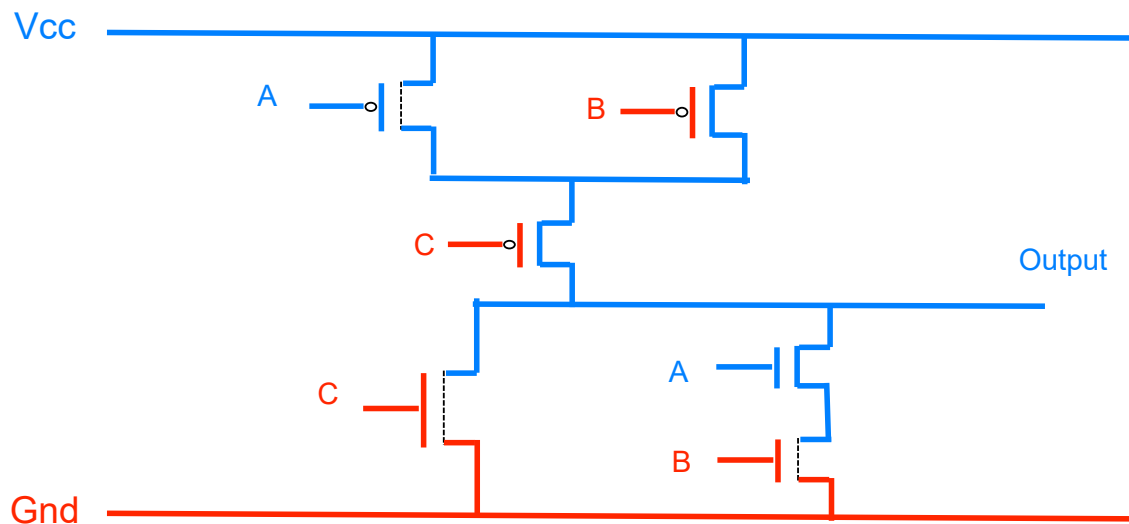
## Complementary: very important



- Everyone take a second to write down the formulas
  - PMOS: NOTs on inputs
  - NMOS: NOT around the outside

7

## Complementary: very important



- Everyone take a second to write down the formulas
  - PMOS: ((Not A) or (Not B)) and (Not C)
  - NMOS: Not (C or (A and B))
    - = (Not C) and (Not (A and B))
    - = (Not C) and ((Not A) or (Not B))

8

## What about... AND?

- Saw and did NAND, but what about AND?

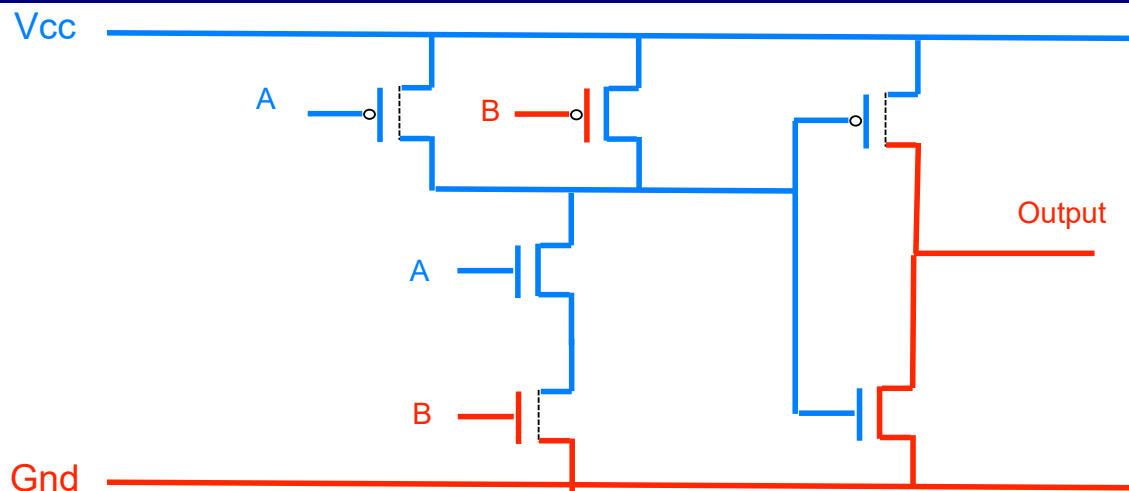
- Truth table on the right...

A	B	Output
0	0	0
0	1	0
1	0	0
1	1	1

- Trying to do this causes problems
    - PMOS formula: NOTs on inputs
    - NMOS formula: NOT around outside
    - ... can't seem to find a formula which works (need more NOTs):
  - AND gate is really a couple gates squished together
    - Not (Nand (A,B))
    - Nor(Not A, Not B)

9

## The AND Gate

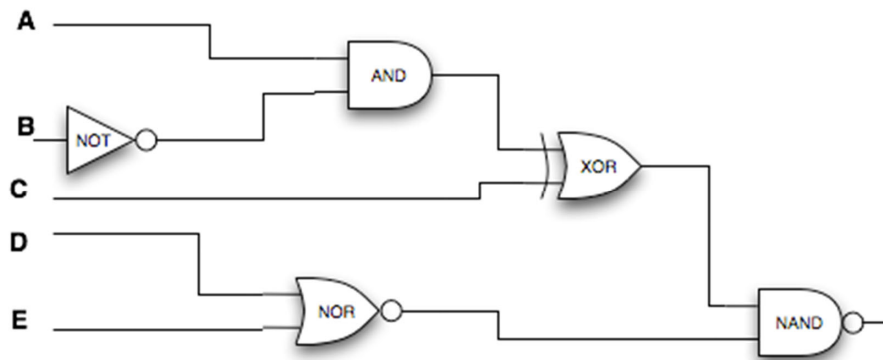


- The AND gate

- A NAND gate followed by a NOT gate
  - Also a good example of how gates connect together
    - Output of one gate goes to input of another

10

## Going forwards: Logic Gates

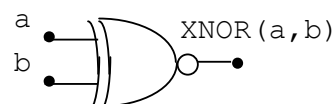
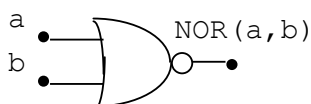
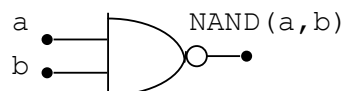
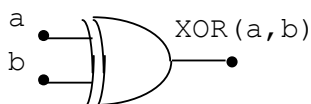
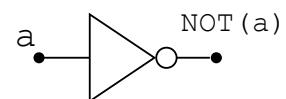
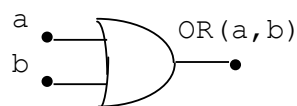
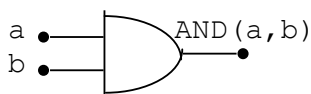


- Going forwards, will mostly design from gates
  - Abstract away transistor level implementation

11

## Boolean Gates

- Saw these gates
  - Mnemonic to remember them

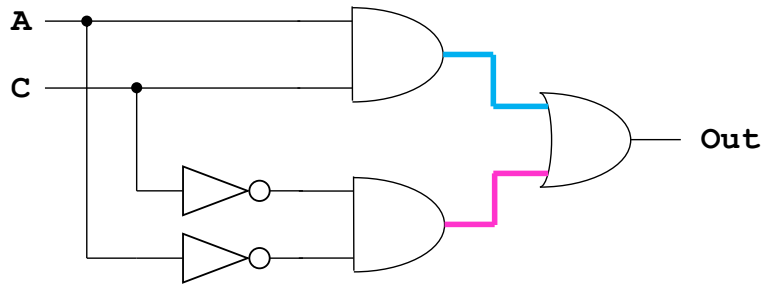


12

# Boolean Functions, Gates and Circuits

- **Circuits** are made from a network of gates.

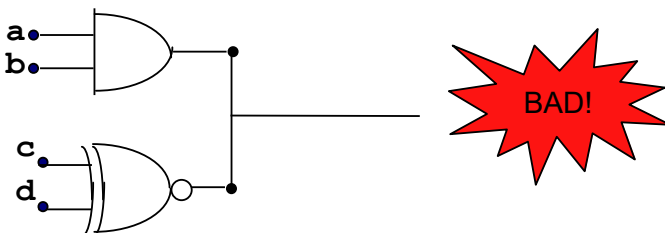
$$(!A \ \& \ !C) \mid (A \ \& \ C)$$



13

## A few more words about gates

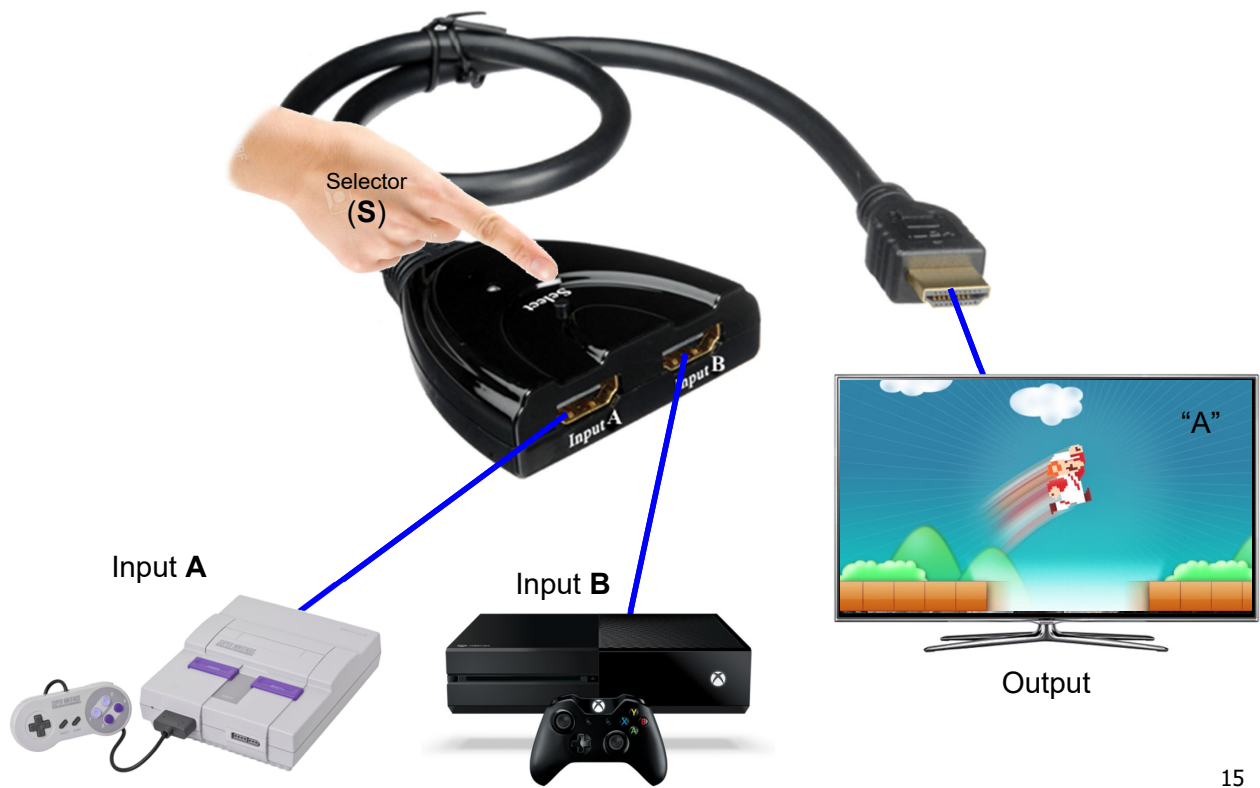
- Gates have inputs and outputs
  - If you try to hook up two outputs, get short circuit  
(Think of the transistors each gate represents)



- If you don't hook up an input, it behaves kind of randomly  
(also not good, but not set-your-chip-on-fire bad)

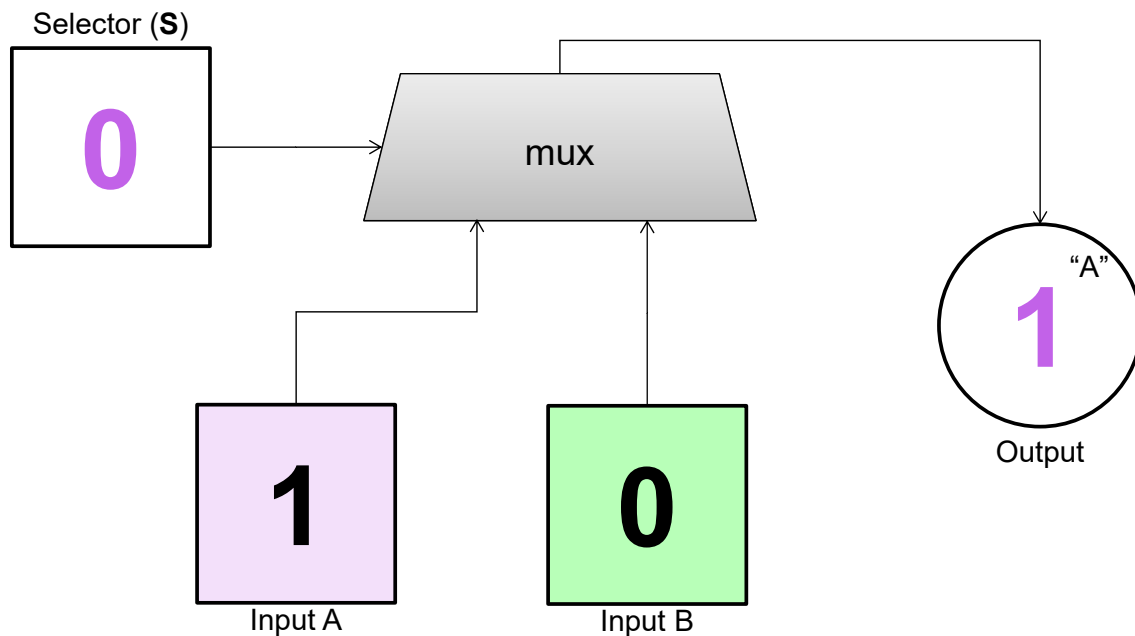
14

# Introducing the Multiplexer (“mux”)



15

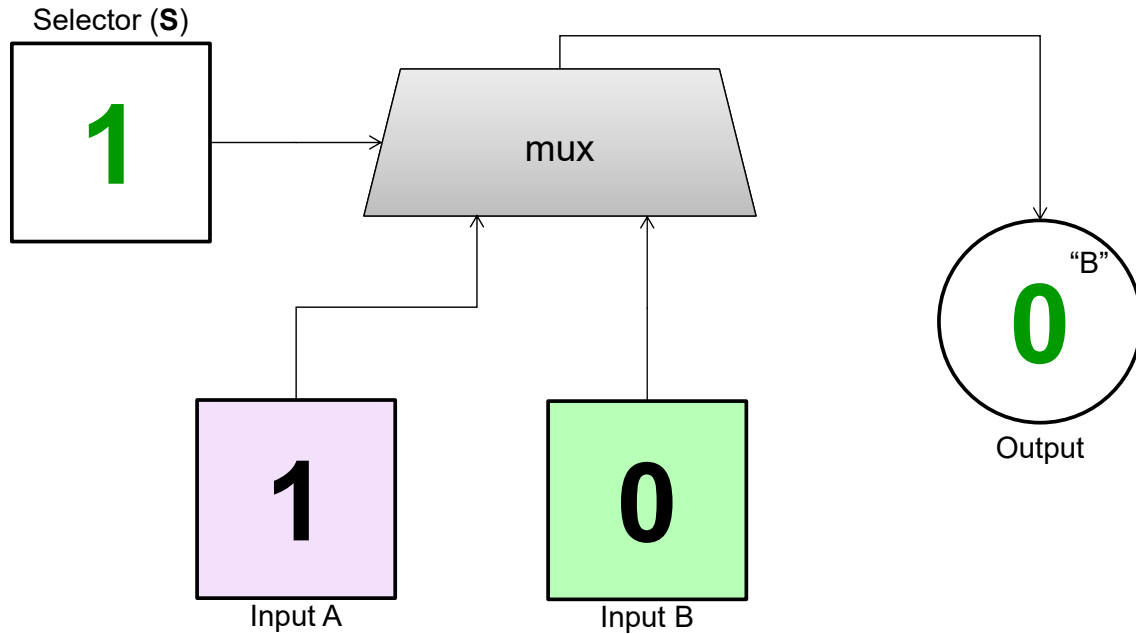
# Introducing the Multiplexer (“mux”)



16

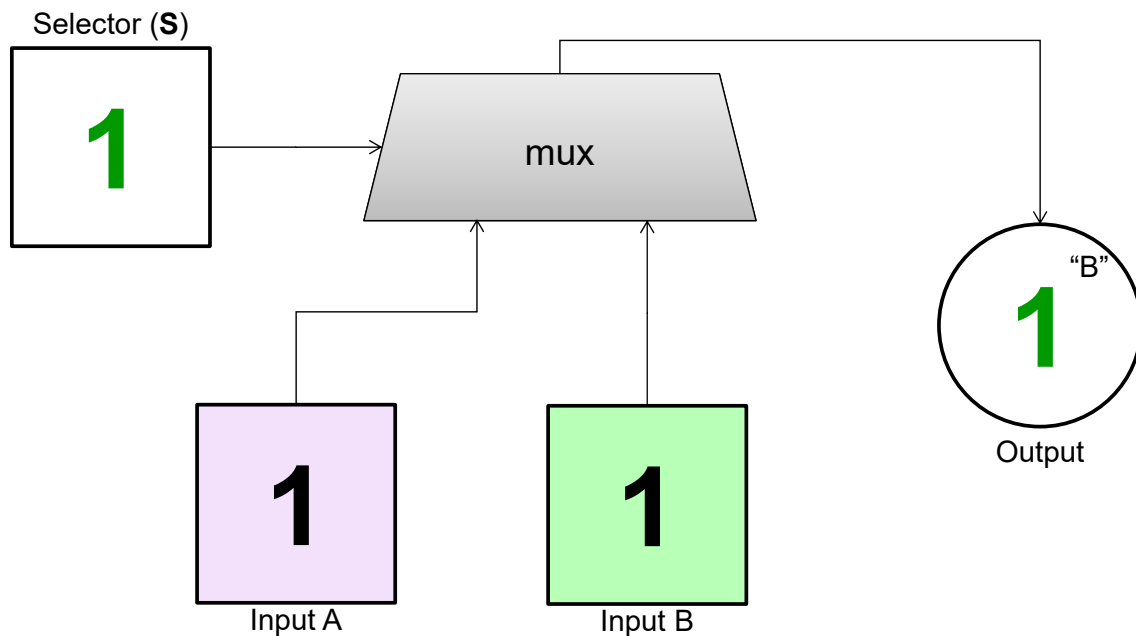


## Introducing the Multiplexer (“mux”)



17

## Introducing the Multiplexer (“mux”)



18

## Let's Make a Useful Circuit

- Pick between 2 inputs (called 2-to-1 MUX)
  - Short for multiplexor

- What might we do first?

- Make a truth table?
  - S is selector:
    - S=0, pick A
    - S=1, pick B

A	B	S	Output
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

- Next: how to get formula?

### Sum-of-products

19

## Sum of Products

- Find the rows where the output is 1
- Write a formula that exactly specifies each row
- OR these all together
  - Possible ways to get 1.

	A	B	S	Output
	0	0	0	0
	0	0	1	0
	0	1	0	0
$\neg A \ \& \ B \ \& \ S$ →	<b>0</b>	<b>1</b>	<b>1</b>	<b>1</b>
$A \ \& \ \neg B \ \& \ \neg S$ →	<b>1</b>	<b>0</b>	<b>0</b>	<b>1</b>
	1	0	1	0
$A \ \& \ B \ \& \ \neg S$ →	<b>1</b>	<b>1</b>	<b>0</b>	<b>1</b>
$A \ \& \ B \ \& \ S$ →	<b>1</b>	<b>1</b>	<b>1</b>	<b>1</b>

20

## Let's Make a Useful Circuit

- **Sum-of-products:**

$(\neg A \ \& \ B \ \& \ S) \mid$   
 $(A \ \& \ \neg B \ \& \ \neg S) \mid$   
 $(A \ \& \ B \ \& \ \neg S) \mid$   
 $(A \ \& \ B \ \& \ S)$

- This is long, though.  
Need to **simplify**.

A	B	S	Output
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

21

## Simplifying The Formula

- Simplifying this formula:

$(\neg A \ \& \ B \ \& \ S) \mid$

$(A \ \& \ \neg B \ \& \ \neg S) \mid$

$(A \ \& \ B \ \& \ \neg S) \mid$

$(A \ \& \ B \ \& \ S)$

← B doesn't matter

22

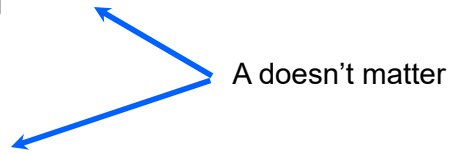
## Simplifying The Formula

- Simplifying this formula:

$(\neg A \ \& \ B \ \& \ S) \mid$

$(A \ \& \ \neg S) \mid$

$(A \ \& \ B \ \& \ S)$



23

## Simplifying The Formula

- Simplifying this formula:

$(A \ \& \ \neg S) \mid$

$(B \ \& \ S)$

24

# Let's Make a Useful Circuit

- Simplified formula:

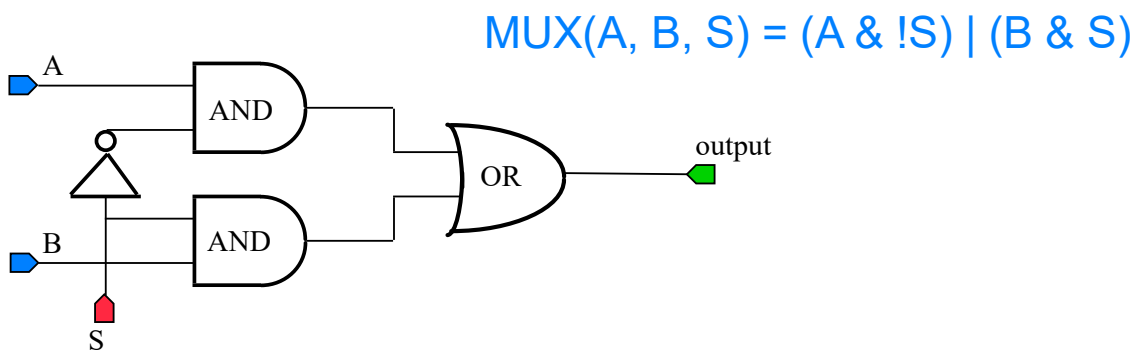
$$(A \ \& \ !S) \ | \ (B \ \& \ S)$$

A	B	S	Output
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	1

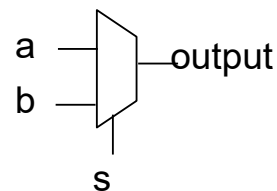
25

## Circuit Example: 2x1 MUX

Draw it in gates:

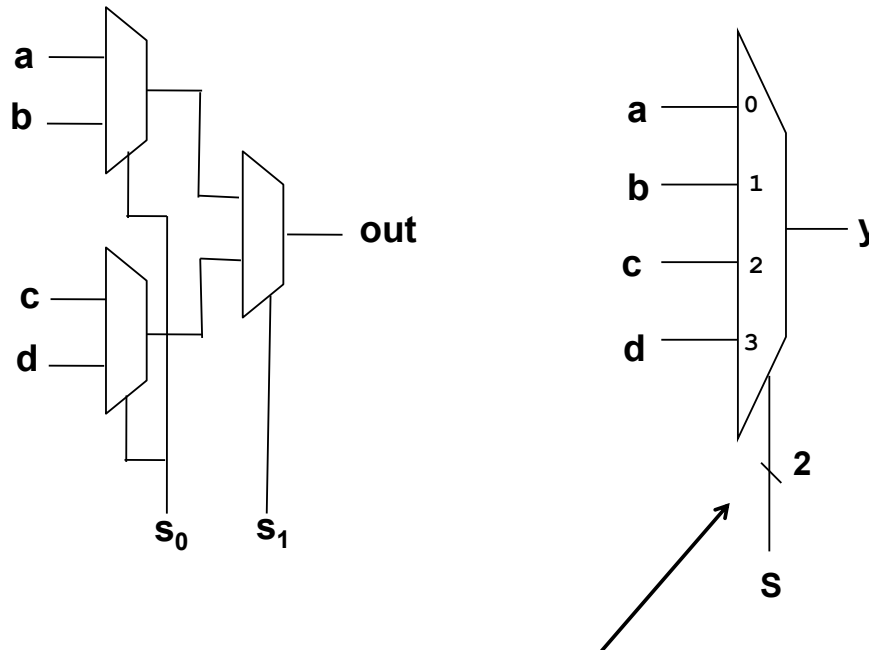


So common, we give it  
its own symbol:



26

## Example 4x1 MUX



The / 2 on the wire means "2 bits"

27

## Boolean Function Simplification

- Boolean expressions can be simplified by using the following rules (bitwise logical):

- $A \& A = A$
- $A \& 0 = 0$
- $A \& 1 = A$
- $A \& !A = 0$

$$A \mid A = A$$

$$A \mid 0 = A$$

$$A \mid 1 = 1$$

$$A \mid !A = 1$$

- $!!A = A$

- $\&$  and  $\mid$  are both commutative and associative
- $\&$  and  $\mid$  can be distributed:  $A \& (B \mid C) = (A \& B) \mid (A \& C)$
- $\&$  and  $\mid$  can be subsumed:  $A \mid (A \& B) = A$

- Can typically just let synthesis tools do this dirty work, but good to know

28

# DeMorgan's Laws

- De Morgan's laws

$$\neg (A \ \& \ B) = (\neg A) \ | \ (\neg B)$$

$$\neg (A \ | \ B) = (\neg A) \ \& \ (\neg B)$$

29

## Wrap Up

- Combinatorial Logic
  - Putting gates together
  - Sum-of-products
  - Simplification
  - Muxes

30