

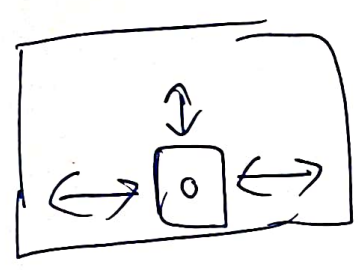
$\begin{matrix} \nearrow 3 \\ \nwarrow 1 \\ \searrow 2 \end{matrix}$

1	2	3
4	0	5

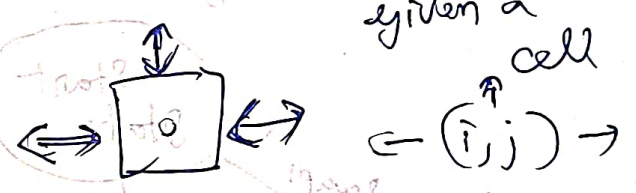
Assume that you have read the problem

1 0 3  
 4 2 5

what happened when I first saw the problem?



Possible choices of backtracking came to my mind.

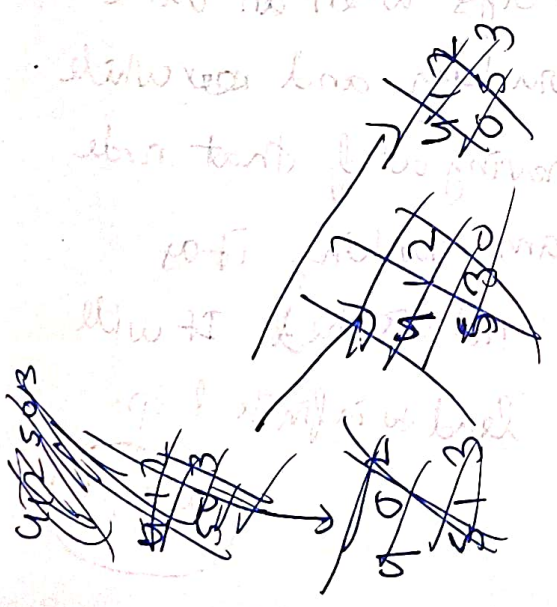


We will try to swap it with the all 4 directions of the cell.  
 We will try to see what match are we going to get on which state we are going to get next.

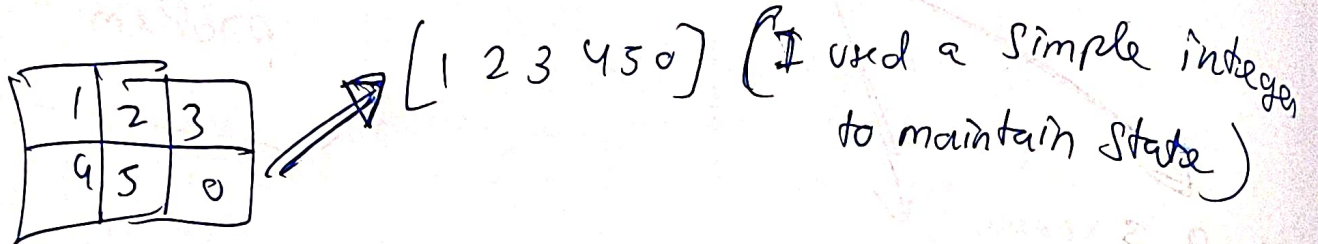
Goal is to reach

1	2	3
4	5	0

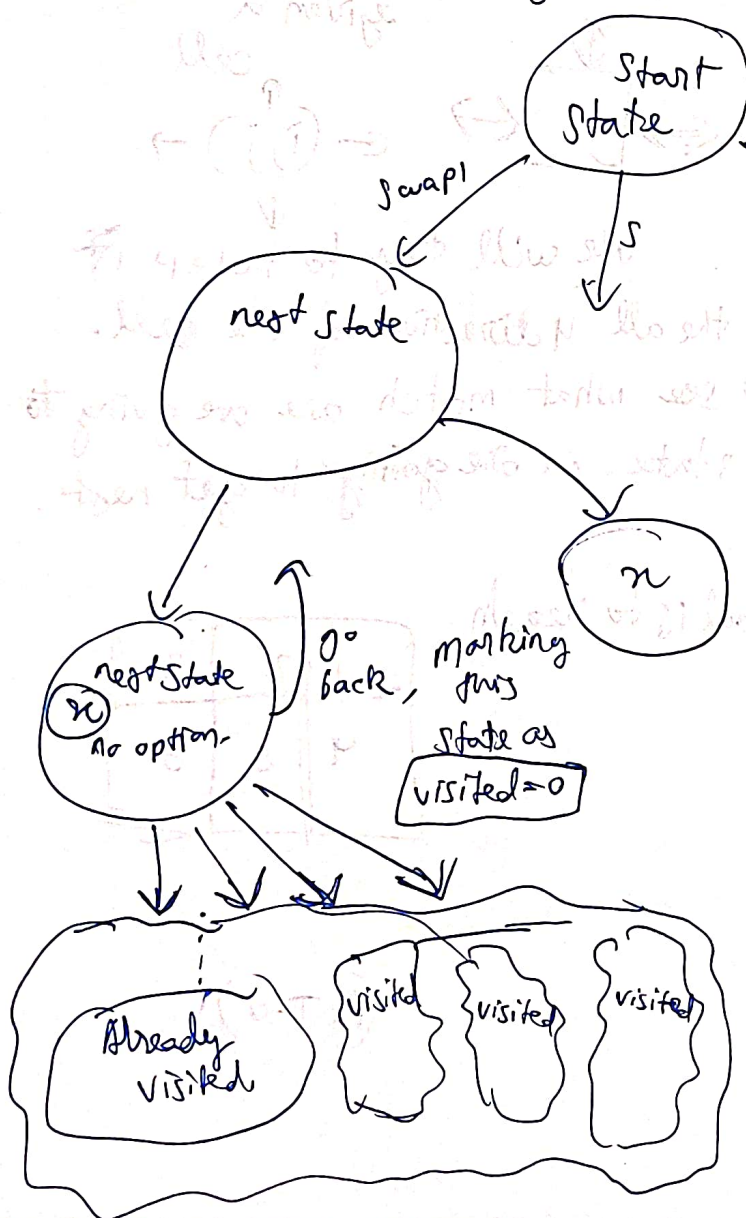
(P.T.O)



I started with dfs and naively was using a visited datastructure to maintain state



At the end of the dfs call  
I was again marking the visited of the  
node as false  
which was resulting in TLE.

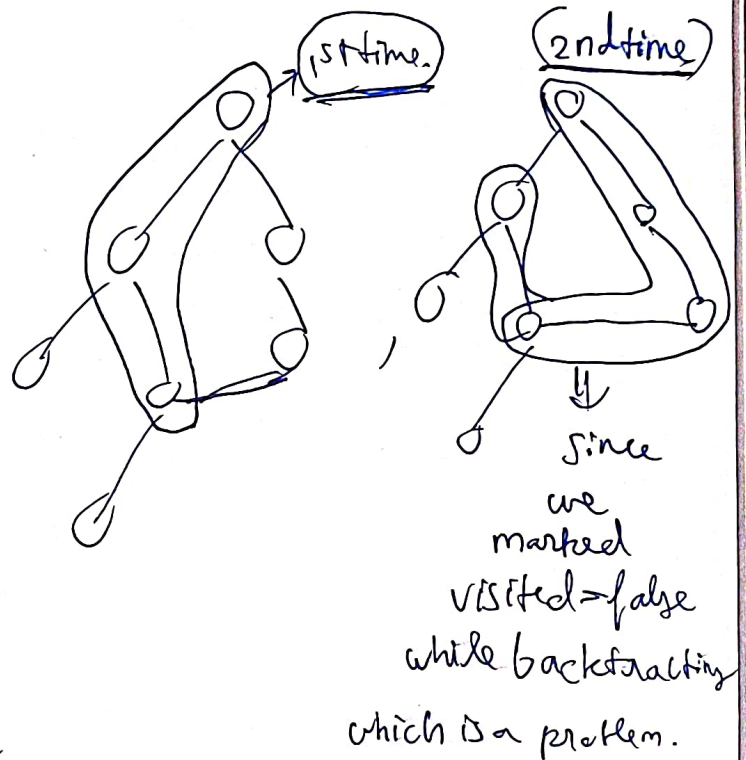
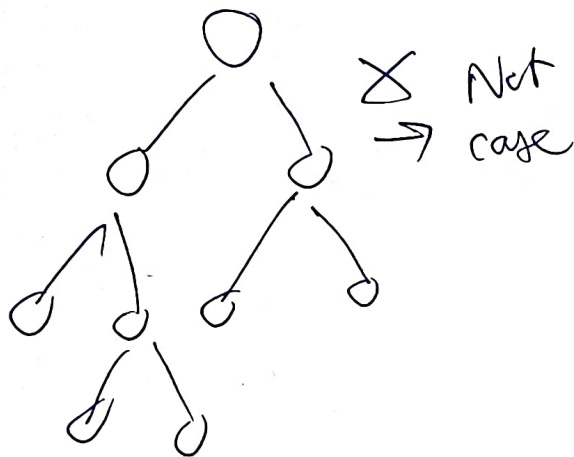


if in dfs when a back tracking and ~~over~~ while moving out of that node am marking it as not visited. It will lead to infinite loop.

Why?

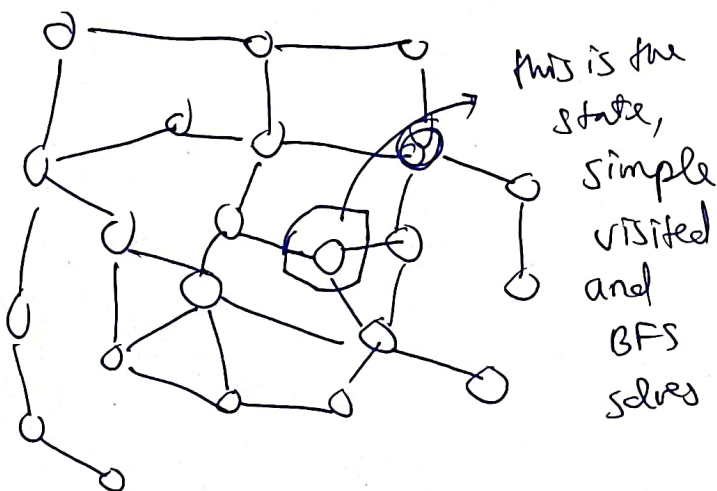


Because I made the mistake of thinking the  
 Transfers to State will represent a undirected tree.  
 But in reality, its like a undirected possible cycle  
 containing graph.



That why TLE was happening

(Soln) I saw editorial.  
 Hassle Free way, do  
 a BFS.



Another way,  
 Like in dijkstra,  
 we keep (node, distance)  
 map,  
 and we go inside that  
 node only if the distance  
 calculated presently is less  
 than the earlier one.  
 So have a (state, score)  
 map,  
 more code complexity.