I want to choose first number as the minimum itself, then it automatically becomes competitive.

K=4

I have to figure out how to write an O(n) solution for this.

2 4 3 3 5 4 9 6

1st element

in this range pick the smallest element.

2

remaining 3 elements.

2nd element

4 3 3 5 4 9 6 K=3.

in this range pick the smallest element.

remaining 2 elements.

3rd element

3 5 4 9 6 K=2

in this range stay to pick minimum.

remaining 1 element

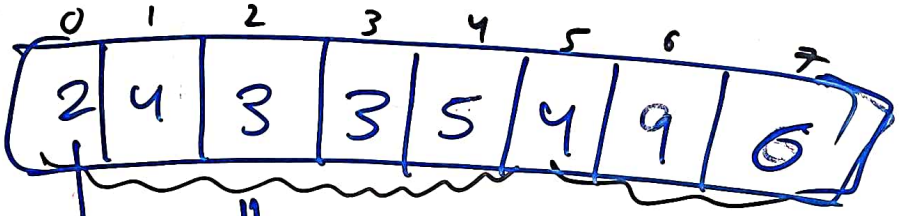
4th element

5 4 9 6

pick minimum in this

K=1

I want to choose first numbr as the minimum itself, then it automatically becomes competitive.



1st element

remaining

in this range pick the smallest element.

(2) n=8 (k=4)

[0 to n-k] { Range to pick 1st element. }

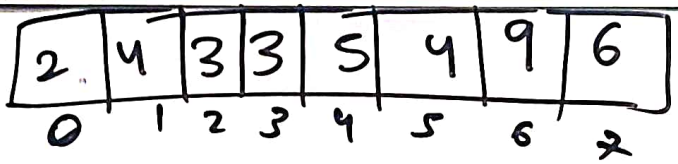
[i+1 to n-k+1] { Range to pick 2nd element }

[j+1 to n-k+2] { Range to pick 3rd element }

:

:

[p+1 to n-1] { range to pick last element }



$k = 4$.

One approach to solve it using fenwick tree. The approach is v. intuitive (intuitive approach)

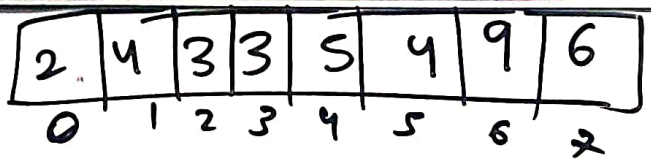
for 1st element → $\min(0, n-k) \rightarrow i$ the index which is having minimum value

for 2nd element → $\min(i, n-k+1) \rightarrow j$ the index which has minimum value in this range.

for 3rd element → $\min(j, n-k+2) \rightarrow m$

for a^{th} element → $\min(p, n-1) \rightarrow k$

In fenwick tree construct a min fenwick tree which returns the index of the smallest index in the given range.



$k=4$, $n=8$

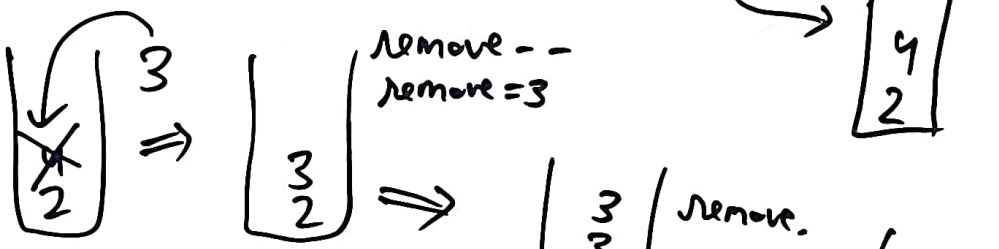
If we have to pick k elements, then we have to remove $n-k$ elements.

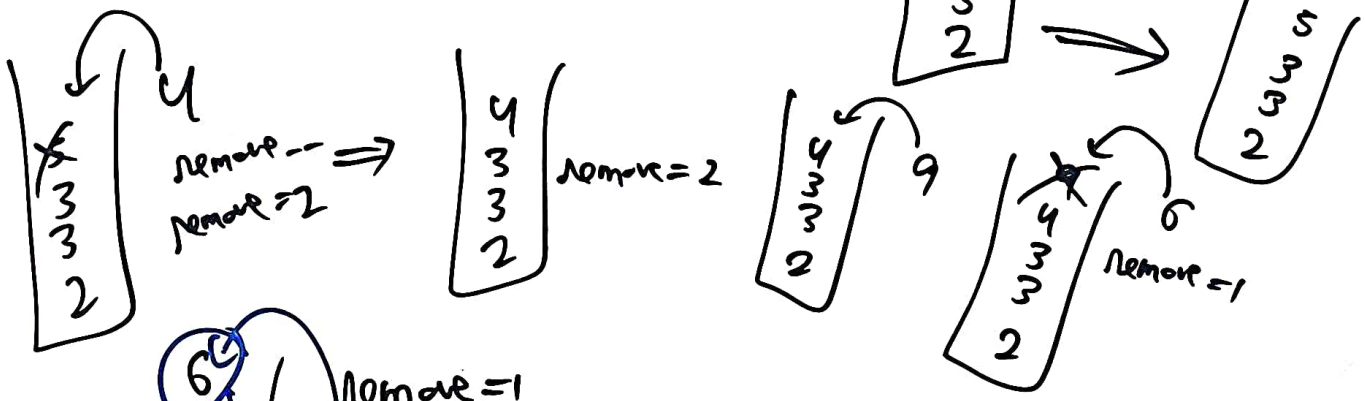
$k=4$, $n-k=4$

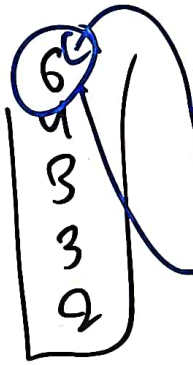
Lets use a monotonic stack to insert elements into the stack, and we will try to keep it non-decreasing.

2 4 3 3 5 4 9 6

remove = n-k = 4.





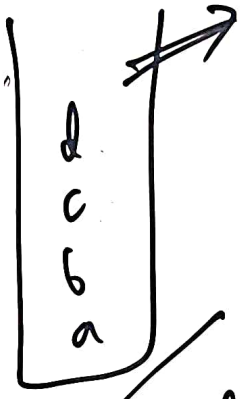


remove = 1

we still can remove 1.

Since the stack is monotonically increasing or non-decreasing we can do that

One doubt which I was having.



remove = n-k.

my doubt was that, what if I exhaust all the remove at the start of the stack itself. And at the end a big number appears in the stack like 1000.

Doubt Clearing.

But my doubt was cleared, since

a = 1 2 3 5 1000

b = 1 2 5 3 2

Seq(a) < Seq(b) this is the index where difference is happening hence we want to apply the removal in the beginning itself.