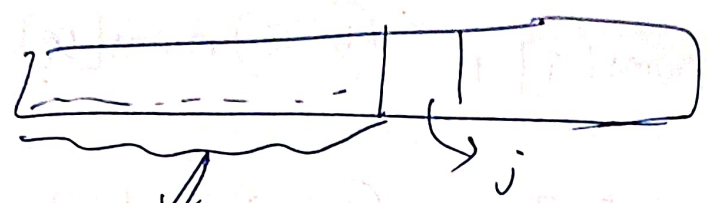


This can be done using nested for loop  $O(N^2)$  complexity

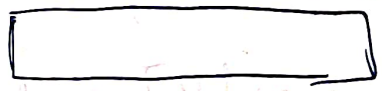


I have answers for these positions  $ans[i]$

$$ans[j] = \max \left( ans[i] \times (j-i) \times num[j] \right)$$

$\forall i \text{ from } [1 \text{ to } (j-1)]$

Is there any way to calculate this faster than  $O(N)$ ?



$[ans[i], (j-1)] \rightarrow$  we cannot maintain this as the sort order.

As  $num[j]$  is different for all indexes.

Jump  $i \rightarrow t \rightarrow k$   
 Option 1  $\rightarrow i \rightarrow j, j \rightarrow k$

$$num[j](j-i) + num[k](k-j)$$

Option 2  $\rightarrow i \rightarrow k \text{ direct}$   
 $(k-i) num[k]$



forming

$$(k-j) num[k] + (j-i) num[k]$$

P.T.O

Compare both options.

$$i \rightarrow j \rightarrow k \rightarrow (j-i) \text{nums}[j] + (k-j) \text{nums}[k]$$

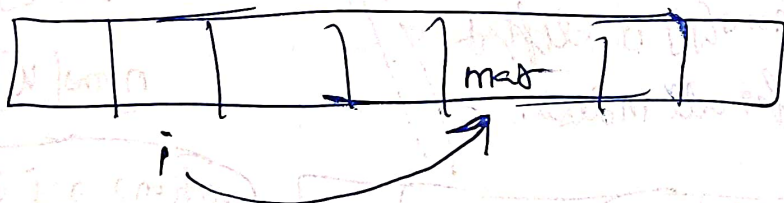
$$i \rightarrow k \rightarrow (k-i) \text{nums}[k] + (k-j) \text{nums}[k]$$

well case 1  $(\text{nums}[k] > \text{nums}[j])$

It's better to jump directly to  $k$

Case 2  $(\text{nums}[j] > \text{nums}[k])$

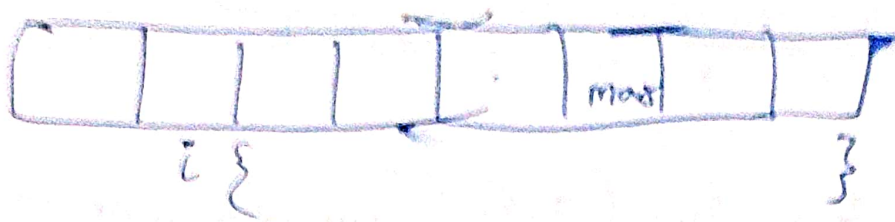
It's better to jump ~~direct~~ to  $j$  first and then  $k$ , so you pick the benefit of  $\text{nums}[j]$  over  $\text{nums}[k]$  and don't miss out on it



In brief, if you are at position  $i$ , it's better to jump to a position  $(p)$  where  $\text{nums}[p]$  is the max on the right-hand side of  $i$ .



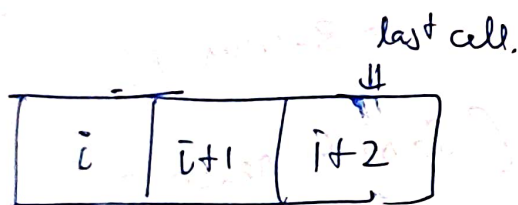
So let's consider how you can code this.



Maybe use a priority queue storing elements in the suffix,  
or a tree set  $\rightarrow$  of  $(\text{nums}[i], i)$  descending order to get  
max

~~But~~  $O(\log N)$ . So all optimisation.

Next further optimisation.



Let's say  $i+2$  is the last cell,  
and you have to jump from  $i \rightarrow i+1$   
 $\searrow$   
 $i+2$

You will jump to the cell  $i+2$  just if  $\text{nums}[i+2] > \text{nums}[i+1]$   
else jump to  $\text{nums}[i+1]$  then  $\text{nums}[i+2]$

So why don't you travel from last cell to the front,

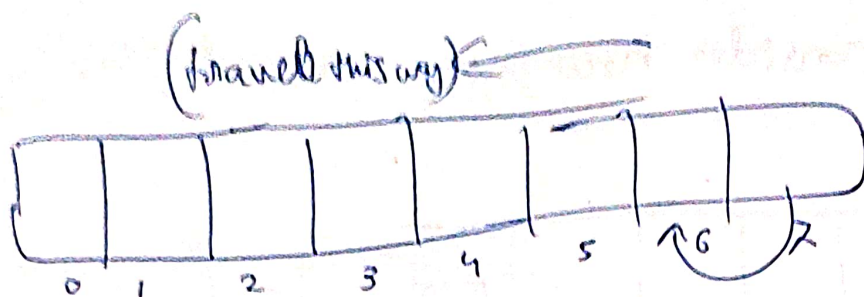
You need to keep track of max in the suffix right?

So travel from  $(\leftarrow)$  Right to left and keep for

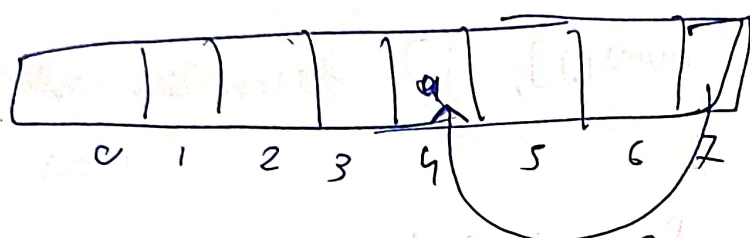
at each cell you travel ~~next~~ left ~~max~~ add it to

P.T.O.

your score



to jump from  $(7 \rightarrow 6)$  obviously  $\text{nums}[7]$  will be score



Jump from  $(7 \text{ to } 4)$

if  $\text{nums}[7]$  is maximum in that range  
the  $\text{score}(4) = (7-4) \text{nums}[7]$   
 $= 3 \cdot \text{nums}[7]$ .

So just keep a max variable  $\rightarrow$  (Score to reach 0).

~~for~~  $\text{sum} = 0$ .  $\text{max} = \text{nums}[7]$

for  $(i = \text{len} - 2; i \geq 0; i--)$  {

~~max~~  $\text{sum} += \text{max}$

$\text{max} = \text{Math.max}(\text{max}, \text{nums}[i])$

at  $(i=0)$

we will get the  $\text{max}(\text{sum})$ .