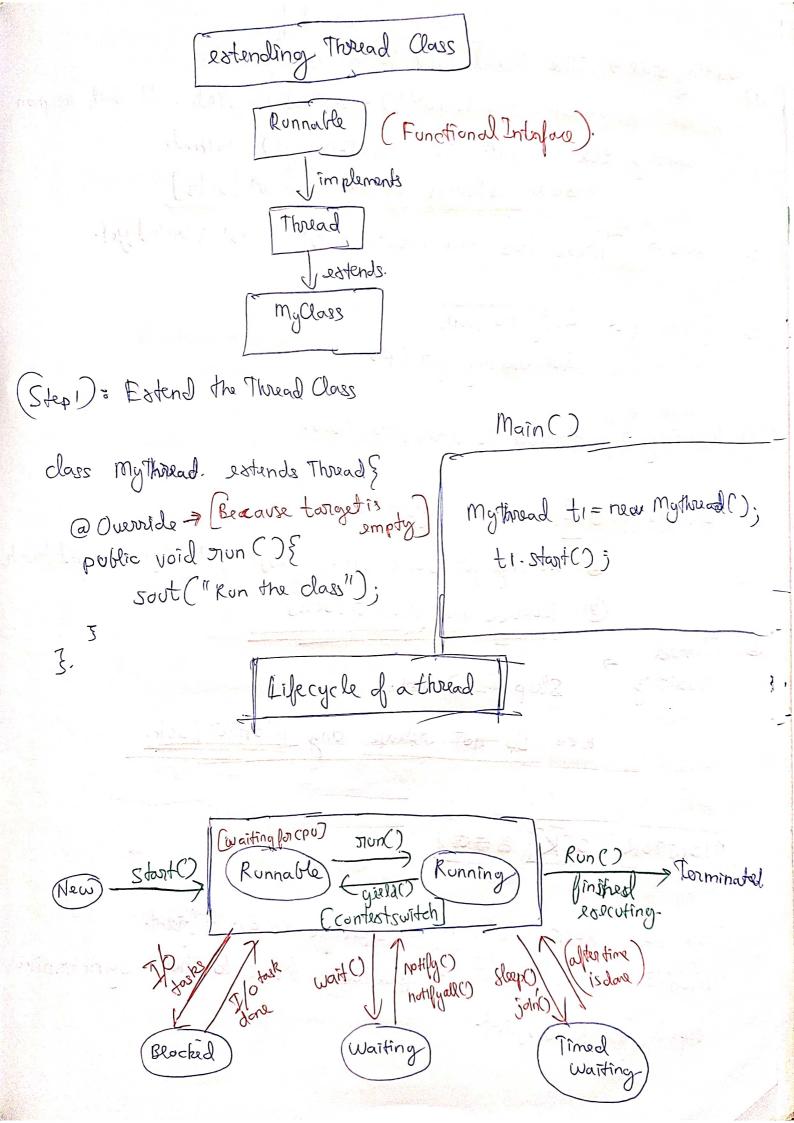


class My Rumnable thread implements Rumable } @ Oumide B() unt pion ogga Sout ("implementing kun method"); But. Placing this Object into the Inside main Methods thereads constructor give it Main () } thrad's powers My Runnable Thread Junnable Obj = new My Runnable Thread (); Thread t = new thread (runnable Obj); to start (); => thread has been starts. to Book of Internally rall + Tun () Roxprivate Runnable target Thread-class } Implementation of non method indides This constructor is Thread Class @ Overvide responsible for public void nun() } Setting the Runnable If (farget !=null) { Object. target. nun();



	0.00
States of thread	
Waiting State > The Afread will be in waiting fill.	
D when thread wait () -> method is called . It will keep	٥ŋ
waiting till we call the thread-notify() method,	
(xxx) Releases all the monitor Locks)	
2) New > Thread has been created but its not stanted yet	
3) Runnable > leady to kun,	
But waiting for (PU	
9 Running > thread is executing code.	
Blocked > 0 I/o read is going on.	
@ Waiting for Lock on an Object of Synchronised &	loc l
3 Releases all Monto Locks.	

Waiting Sleep () inms.

8 + + Do not release any Monitor Lock.

MONITOR LOCK (**

This is a lock which is specific to an object.

Monitor Lock is meant to provided lock to the synchronized block of the object.

Working of Monitor Locks. (Case I) (Thread Ti class A & Thread 15 dos A obj = rew A(); Synchronised void method 1 ()} Jobj-method1(); Thread 12. Using the Same Object Thread 2 { 3- Both one trying to ADD = 2 hour ACD) occess the synchronised block, obj. method1(); then TI will complete first and then T2. (ase 2) A obj1 = new A(); A obj2 = new ACD; Thread 2 { obj1-method1(); Since both are different > Object obj1 and Thread2} obj2. Hence thread 1 Obj2-method1 (); and thread will execute parallely together.

Code Example with Monitor Locks. public class Montdon Edample & main () { Monitor Example obj = new (public Vsynchronised void tosk1() Monitor Edample Sout ("inside tasks") sync Class Thread 1 extends Thread ? (Rock Thread-sleep (10000), @ override public void mun ()} obj-task1C); pullic void task2() { class thread 2 extends threads Sout (" before synchronised took?"); @ Override public void nun () { Synchronised (this) & Sync obj.task2c). Sout ("task2 înside synchronizad"); Bock Class thread 3 extends Thread & @d vernide poffic void non ()} public void tooks ()} 06j-task3(); South fasks"); Output. Tinside task 1 Thread 1 ti= new Thread (C); before synchronised tooks Thread 2 te= new thread 2(); Thread 3 to= new thread 3() task 3 f1-2404(); taske inside Synchronised. te-Startes; t3-Stoot ()

From the out put see, how.

task (() -> acquired the monitor lock for 10 seconds because it had the sleep ().

- We taken thread 2 tried to so execute tasks, it was able to enter the non-syne part, but had to wait for thread 1 to complete to give back the Monttor Lock to thread 2.
- thread 3 had he issues executing on it was a non-syne block. Implement Broducir, Consumer Broblem Yourself