操作系统实验报告

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实验题目:

进程创建以及 Linux 进程通信

实验内容:

- 1.掌握 fork()、pipe()系统调用的形式和功能
- 2.掌握通过 fork()派生的子进程之间通过 pipe 文件的通信方式

设计原理:

用 pipe()创建一个管道文件, 然后用 fork()创建 2 个生产者进程和 2 个消费者进程, 实现它们之间通过 pipe 文件实现信息传递。

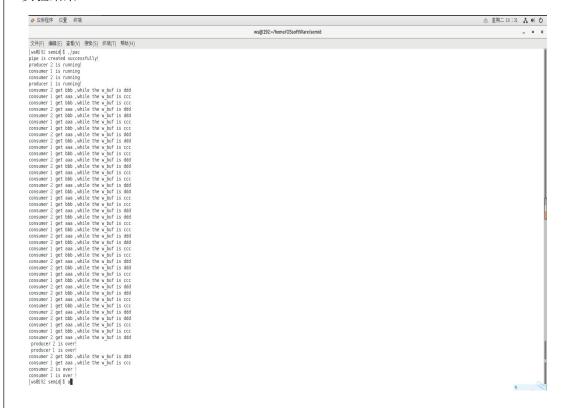
实验源码

```
#include"sys/types.h"
2. #include"sys/file.h"
3.
      #include"unistd.h"
4.
5.
   char r_buf[4];
6. char w_buf[4];
7.
     int pipe_fd[2];
8. pid_t pid1,pid2,pid3,pid4;
9.
10. int producer(int id);
11. int consumer(int id);
12.
13. int main(int argc,char **argv)
14. {
15.
         if(pipe(pipe_fd) ⟨ ∅)
16.
17.
             printf("pipe create erro");
18.
         exit(-1);
19.
20.
         else
21.
22.
             printf("pipe is created successfully!\n");
```

```
23.
             if((pid1 = fork()) == 0 )
24.
             producer(1);
25.
             if((pid2 = fork()) == 0 )
26.
                          producer(2);
27.
            if((pid3 = fork()) == 0 )
28.
                          consumer(1);
29.
             if((pid4 = fork()) == 0 )
30.
                   consumer(2);
31.
32.
       close(pipe_fd[0]);
33.
         close(pipe_fd[1]);
34.
35.
36.
      int i, pid,status;
37.
         for( i = 0 ; i < 4;i++)</pre>
38.
        pid = wait(&status);
39.
         exit(0);
40. }
41.
42. int producer(int id)
43. {
44. printf("producer %d is running! \n",id);
45.
         close(pipe_fd[0]);
46. int i = 0;
47.
         for( i = 1 ;i < 30 ;i++)</pre>
48.
49.
             sleep(3);
50.
           if(id == 1)
51.
                strcpy(w_buf,"aaa\0");
52.
           else
53.
                strcpy(w_buf,"bbb\0");
54.
            if(write(pipe_fd[1],w_buf,4) == -1)
55.
                printf("write to pipe error\n");
56.
57.
         close(pipe_fd[1]);
58.
        printf(" producer %d is over! \n",id);
59.
         exit(id);
60. <sub>}</sub>
61. int consumer( int id)
62. {
63.
64.
        close(pipe_fd[1]);
65.
         printf("consumer %d is running\n",id);
66.
         if( id == 1)
```

```
67.
              strcpy(w_buf ,"ccc\0");
68.
69.
               strcpy(w_buf ,"ddd\0");
70.
          while(1)
71.
72.
              sleep(1);
73.
              strcpy(r_buf,"eee\0");
74.
              if(read(pipe_fd[0] ,r_buf,4) == 0)
75.
76.
              printf("consumer %d get %s ,while the w_buf is %s\n",id, r_buf,w_buf);
77.
78.
          close(pipe_fd[0]);
79.
          printf("consumer %d is over !\n",id);
80.
          exit(id);
81. }
```

2.实验结果



问题解析 1.参考代码 main 中两个 close(pipe fd[0])和 close(pipe fd[1])的作用

这两行代码的作用是关闭父进程对管道文件的读写操作,防止其影响模拟生产者、消费者的四个子进程间对管道的数据读写操作。如果 main 函数里没有这两行代码,在不同情况下会有生产者或则消费者永远等待。

学生通过修改实验指导手册的源码得到了 main 函数未写 close(pipe fd[0])和

close(pipe fd[1])时,生产者、消费者分别无限期等待的两种情况,以下分别展开分析。

一、消费者无限等待

1.修改后代码: 在源代码基础上删去 main 函数中的 close(pipe_fd[0])和 close(pipe_fd[1])两行代码即可。

2.结果展示

```
[ws@192 semid] $ ./pac
pipe is created successfully!
producer 2 is running!
consumer 1 is running
consumer 2 is running!
consumer 2 get bbb, while the w_buf is ddd
consumer 1 get aaa ,while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get bbb, while the w_buf is ccc
consumer 2 get bbb, while the w_buf is ccc
consumer 2 get bbb, while the w_buf is ccc
consumer 2 get bbb, while the w_buf is ccc
consumer 2 get aba a, while the w_buf is ccc
consumer 2 get aba a, while the w_buf is ccc
consumer 1 get aba while the w_buf is ccc
consumer 1 get aba, while the w_buf is ccc
consumer 2 get aba, while the w_buf is ccc
consumer 1 get aba, while the w_buf is ccc
consumer 2 get aba, while the w_buf is ccc
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get abb, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ddd
consumer 2 get aba, while the w_buf is ccc
consumer 2 get abb, while the w_buf is ccc
consumer 2 get abb, while the w_buf is ccc
consumer 2 get aba, while the w_buf is ccc
consumer 2 get abb, while the w_buf is ccc
consumer 2 get abb, while the w_buf is ccc
consumer 2 get abb, while the w_buf is ccc
consumer 2 get abb, while the w_buf is ccc
consumer 2 get abb, while the w_buf is ccc
consumer 2 get abb, while the w_buf is ccc
consumer 2 get abb, while the
```

图 1: consumer 无限等待图点,进程正常结束图点

3.结果分析

如果父进程没有关闭文件读写描述符,当生产者进程全部生产完物品(数据)后关闭了他自己的写文件描述符 pipe_fd[1],消费者一直在运行,消耗物品(数据),直到管道缓冲区内数据全部消耗完,此时因为父进程的写文件描述符 pipe_fd[1]未关闭,消费者进程在执行 read()操作时会被阻塞,等待父进程往管道缓冲区写数据,但是父进程没有 write()操作,因此消费者进程进入了无限期等待。

如果父进程的写文件描述符 pipe_fd[1]关闭,即此时的管道文件 Inode 的 i 写文件描述符的 count 值为 0,所有进程都关闭了写端,系统认为已经读到了数据的末尾, read()函数返回的读出字节数为 0,跳出循环,所以 consumer 不会被阻塞,无限等待。

二、生产者无限等待

1.修改后的代码

```
1. int producer(int id)
2. {
3.  printf("producer %d is running! \n",id);
4.  close(pipe_fd[0]);
5.  int i = 0;
6.  while(1)
```

```
7.
  8.
               sleep(1);
  9.
               if(id == 1)
  10.
                 strcpy(w_buf,"aaa\0");
  11.
               else
  12.
                  strcpy(w_buf,"bbb\0");
  13.
               if(write(pipe_fd[1],w_buf,4) == -1)
  14.
  15.
                      printf("write to pipe error\n");
  16.
                    break;
  17.
  18.
  19.
           close(pipe_fd[1]);
  20.
          printf(" producer %d is over! \n",id);
  21.
           exit(id);
  22. }
  23. int consumer(int id)
  24. {
  25.
  26. close(pipe fd[1]);
  27.
           printf("consumer %d is running\n",id);
  28.
         if( id == 1)
  29.
               strcpy(w_buf ,"ccc\0");
  30.
  31.
                strcpy(w_buf ,"ddd\0");
  32.
  33.
           int i ;
  34.
          for( i = 1 ;i < 10 ;i++)</pre>
  35.
  36.
           sleep(3);
  37.
               strcpy(r_buf,"eee\0");
  38.
             if(read(pipe_fd[0] ,r_buf,4) == 0)
  39.
                  break;
  40.
             printf("consumer %d get %s ,while the w_buf is %s\n",id, r_buf,w_buf);
  41.
  42.
           close(pipe_fd[0]);
  43.
           printf("consumer %d is over !\n",id);
  44.
           exit(id);
 45. }
2.结果展示
```

```
[ws@192 semid]$ ./pacf
pipe is created successfully!
producer 2 is running!
consumer 1 is running
consumer 1 is running
producer 1 is running
producer 1 is running
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 2 get aaa, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_buf is ccc
consumer 1 get bbb, while the w_
```

图 2: producer 无限等待图点,进程正常结束图点

3.结果分析

如果父进程没有关闭文件读写描述符,当消费者进程消耗完指定数量的物品(数据)后关闭了他自己的读文件描述符 pipe_fd[0],此时生产者一直在运行,生产物品(数据),直到管道缓冲区内空间全部被占满,此时因为父进程的读文件描述符 pipe_fd[0]未关闭,生产者进程在执行 write()操作时会被阻塞,等待父进程从管道缓冲区读数据,但是父进程没有read()操作,因此生产者进程被阻塞,进入了无限期等待。

如果父进程的读文件描述符 pipe_fd[0]关闭,即此时的管道文件 Inode 的读文件描述符的 count 值 0,所有读进程都关闭了写端,生产者执行 write()函数时将收到内核传来的 SIFPIPE 信号,应用程序可以处理该信号,也可以忽略(默认动作则是应用程序终止)。跳出循环,所以 producer 不会被阻塞,无限等待。

问题解析 2.参考代码中

for(i=0;i<4;i++)

pid=wait(&ststus);

的作用,如果没有会是何种结果,为什么?

Wait()函数起阻塞父进程进一步执行的作用,父进程调用 wait()函数后会寻找在父进程前结束的变成僵尸进程的子进程,从而释放子进程所占用的资源。

此代码用于父进程查询等待其 fork()派生的所有子进程是否终止,在此问题中,创建两个生产者进程和两个消费者进程工作均结束后,通过 exit(id)向父进程传递该子任务结束的信息,并由 wait(&ststus)接收,由于有四个子进程所以需要循环四次。最后父进程和子进程所占的所有资源被回收。

```
[ws@192 semid|$ ./pac
pipe is created successfully!
  page 15 creates successively producer 2 is running! consumer 1 is running consumer 2 is running producer 1 is running consumer 2 get bbb ,while the w_buf is ddd fahterID:5876 consumer 1 get aga while the w_buf is consumer 2 get bbs while the w_buf is consumer 
     consumer 1 get aaa ,while the w_buf is ccc
fahterID:5876
    consumer 1 get bbb ,while the w_buf is ccc fahterID:5876
     consumer 2 get aaa ,while the w_buf is ddd
fahterID:5876
     consumer 2 get bbb ,while the w_buf is ddd fahterID:5876
       consumer 1 get aaa ,while the w_buf is ccc
fahterID:5876
     consumer 1 get bbb ,while the w_buf is ccc
fahterID:5876
     consumer 2 get aaa ,while the w_buf is ddd fahterID:5876
    consumer 2 get bbb ,while the w_buf is ddd fahterID:5876
    consumer 1 get aaa ,while the w_buf is ccc fahterID:5876
    consumer 1 get bbb ,while the w_buf is ccc fahterID:5876
anterID:5876 consumer 2 get aaa ,while the w_buf is ddd fahterID:5876 producer 2 is over! producer 1 is over! consumer 2 get bbb ,while the w_buf is ddd fahterID:5876 consumer 1 get aaa ,while the w_buf is ccc fahterID:5876 consumer 2 is over!
```

图三: 父进程等待子进程结束,回收子进程资源

如果没有上述两行代码(即没有 wait()函数),会出现两种情况。

一、由于本程序父进程指令少,很可能会在四个子进程结束前提前结束,这四个进程就 变成了孤儿进程,此时由进程号为1的 init 进程接管,待四个进程指令执行完毕后由 init 进 程回收,这种情况是 OS 可以正常处理的情况,不会带来资源浪费的情况。

```
[ws@ 92 semid] $ ./pac
pipe is created successfully!
producer 1 is running!
producer 2 is running!
consumer 1 is running
consumer 2 is running
[ws@ 92 semid] $ consumer 2 get aaa ,while the w_buf is ddd
fabter[n.]
               1 get bbb ,while the w buf is ccc
 fahterID:
                1 get aaa ,while the w_buf is ccc
consumer 2 get bbb ,while the w_buf is ddd fahterID:1
consumer 2 get aaa ,while the w_buf is ddd fahterID:1
                1 get bbb ,while the w buf is ccc
 fahterID: 1
                1 get aaa .while the w buf is ccc
 fahterID: 1
                get bbb ,while the w buf is ddd
consumer 2 get aaa ,while the w_buf is ddd fahterID:1
consumer 1 get bbb ,while the w_buf is ccc fahterID:1
consumer 1 get aaa ,while the w_buf is ccc fahterID:1
                .
2 get bbb ,while the w_buf is ddd
TahterID:1

producer 1 is over!
producer 2 is over!
consumer 2 get aaa while the w_buf is ddd
fahterID:1
consumer 1 get bbb ,while the w_buf is ccc
fahterID:1
consumer 1 get bbb ,while the w_buf is ccc
consumer 2 is over !
```

图 4 父进程提前结束, init 接管子进程

二、如果有子进程在父进程结束前提前结束,该子进程就变成了僵尸进程:父进程由于 没有 wait()函数来等待它,如果该父进程一直运行,该子进程所占有的某些资源未被释放, 当系统中有大量僵尸进程时,会阻碍新进程的产生,这是我们应该避免的。

问题解析 3.参考代码 producer()和 consumer()中的 sleep 的作用

如上述代码所见, producer()和 consumer()每执行一次生产物品或消费物品的操作时,会调 用一次 sleep(n)操作.

- 一、四个子进程再并发运行时,无法保证先后运行顺序。四个子进程调用 sleep(n)使自己休眠,以此来保证所有 producer 和 consumer 进程已经关闭了相应的文件读写描述符,实现单方向的、先进先出的字节流传输机制。
- 二、执行 sleep(n)操作的进程(n>0)会进入等待态,让出 CPU,并在 n 秒的时间内不会再抢占 CPU(不进入就绪队列),进入等待队列;此时其他的 producer、consumer 进程可以占取 CPU 工作。这样每个进程都有占得 CPU 工作的机会,防止了在有限资源下进程饿死现象(对于本程序:即防止在 producer 产生有限资源的情况下,某一个 consumer 一口气全消耗完,而其他的 consumer 得不到资源,草草结束的现象)。

```
| wseN 22 sensid| S . /pac | pipe is created successfully! | producer 1 is unung! | producer 2 is running! | producer 2 is running! | producer 2 is running! | producer 2 is over! | consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aaa , while the w_buf is ccc consumer 1 get aab , while the w_buf is ccc consumer 1 get aab , while the w_buf is ccc consumer 1 get aab , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 1 get bbb , while the w_buf is ccc consumer 2 get bbb , while the w_buf is ccc consumer 2 get bb
```

图 5 consumer1 进程消耗完所有的产品, consumer2 无产品可消耗

三、同时实现了多进程的并发执行;减少了 producer 进程不断生成产品导致的管道缓冲空间一直处于充满的状态,从而导致 producer 进程频繁等待的现象和 consumer 进程不断消耗产品导致的管道缓冲空间一直处于空闲状态,从而导致 consumer 进程频繁等待的现象,即减少 read()操作和 write()操作被阻塞的次数,实现了生产者边生产,消费者边消费的状态,使管道文件缓冲区的空间的利用率保持在合适的范围,既不会太高也不会太低,提升了 CPU的利用率,使进程高效运行。