# ENE4014: Programming Languages

# Lecture 3 — Basics of the OCaml Language

Woosuk Lee
2019 Spring

# Why learn ML?

Learning ML is a good way of experiencing modern language features:

- functional programming: scala, java8, haskell, python, JavaScript, etc
- value-oriented programming: scala, haskell, scheme, etc
- type inference: scala, haskell, etc
- pattern matching: scala, etc
- algebraic data types, module system, etc

# Basics of the OCaml Language

- Expressions and values
- Names and functions
- Pattern matching
- Type inference
- Tuples and lists
- Data types
- Exceptions
- Modules

# Basics of the OCaml Language

- Expressions and values
- Names and functions
- Pattern matching
- Type inference
- Tuples and lists
- Data types
- Exceptions
- Modules

$\boxed{\text{Write and run all examples in the slides by yourself!}}$

# An OCaml Program is an Expression

Statement and expressions:

# An OCaml Program is an Expression

Statement and expressions:

- A statement *does something*.
- An expression *evaluates to a value*.

# An OCaml Program is an Expression

Statement and expressions:

- A statement *does something*.
- An expression *evaluates to a value*.

Programming languages can be classified into

- statement-oriented: C, C++, Java, Python, JavaScript, etc
  - often called "imperative languages"
- expression-oriented: ML, Haskell, Scala, Lisp, etc
  - often called "functional languages"

# Arithmetic Expressions

- Arithmetic expressions evaluate to numbers: e.g., `1+2*3`, `1+5`, `7`

# Arithmetic Expressions

- Arithmetic expressions evaluate to numbers: e.g., `1+2*3`, `1+5`, `7`
- Try to evaluate expressions in the REPL:

```
# 1+2*3;;
- : int = 7
```

# Arithmetic Expressions

- Arithmetic expressions evaluate to numbers: e.g., `1+2*3`, `1+5`, `7`
- Try to evaluate expressions in the REPL:

```
# 1+2*3;;
- : int = 7
```

- Arithmetic operators on integers:

| | |
|---|---|
| $a$ `+` $b$ | addition |
| $a$ `-` $b$ | subtraction |
| $a$ `*` $b$ | multiplication |
| $a$ `/` $b$ | divide $a$ by $b$, returning the whole part |
| $a$ `mod` $b$ | divide $a$ by $b$, returning the remaining part |

# Boolean Expressions

- Boolean expressions evaluate to boolean values (i.e., true, false).

# Boolean Expressions

- Boolean expressions evaluate to boolean values (i.e., true, false).
- Try to evaluate boolean expressions:

```
# true;;
- : bool = true
# true;;
- : bool = true
# 1 > 2;;
- : bool = false
```

# Boolean Expressions

- Boolean expressions evaluate to boolean values (i.e., true, false).
- Try to evaluate boolean expressions:

```
# true;;
- : bool = true
# true;;
- : bool = true
# 1 > 2;;
- : bool = false
```

- Comparison operators produces boolean values:

| | |
|---|---|
| $a$ = $b$ | true if $a$ and $b$ are equal |
| $a$ <> $b$ | true if $a$ and $b$ are not equal |
| $a$ < $b$ | true if $a$ is less than $b$ |
| $a$ <= $b$ | true if $a$ is less than or equal to $b$ |
| $a$ > $b$ | true if $a$ is greater than $b$ |
| $a$ >= $b$ | true if $a$ is greater than or equal to $b$ |

# Boolean Operators

- Boolean expressions are combined by boolean operators:

```
# true && false;;
- : bool = false
# true || false;;
- : bool = true
# (2 > 1) && (3 > 2);;
- : bool = true
```

# ML is a Statically Typed Language

If you try to evaluate an expression that does not make sense, OCaml rejects and does not evaluate the program: e.g.,

```
# 1 + true;;
Error: This expression has type bool but an expression was
expected of type int
```

# cf) Static Types and Dynamic Types

Programming languages are classified into:

- *Statically typed languages*: type checking is done at compile-time.
  - type errors are detected before program executions
  - C, C++, Java, ML, Scala, etc
- *Dynamically typed languages*: type checking is done at run-time.
  - type errors are detected during program executions
  - Python, JavaScript, Ruby, Lisp, etc

# cf) Static Types and Dynamic Types

Programming languages are classified into:

- *Statically typed languages*: type checking is done at compile-time.
  - type errors are detected before program executions
  - C, C++, Java, ML, Scala, etc
- *Dynamically typed languages*: type checking is done at run-time.
  - type errors are detected during program executions
  - Python, JavaScript, Ruby, Lisp, etc

Statically typed languages are further classified into:

- *Type-safe languages* guarantee that compiled programs do not have type errors at run-time.
  - All type errors are detected at compile time.
  - Compiled programs do not stuck.
  - ML, Haskell, Scala
- *Unsafe languages* do not provide such a guarantee.
  - Some type errors remain at run-time.
  - C, C++

# Which one is better?

Statically typed languages:

- (+) Type errors are caught early in the development cycle.
- (+) Program execution is efficient by omitting runtime checks.
- (−) Less flexible than dynamic languages.

Dynamically typed languages:

- (−) Type errors appear at run-time, often unexpectedly.
- (+) Provide more flexible language features.
- (+) Easy and fast prototyping.

# Conversion between Different Types

- In OCaml, different types of values are distinguished:

```
# 3 + 2.0;;
Error: This expression has type float but an expression
was expected of type int
```

- Types must be explicitly converted:

```
# 3 + int_of_float 2.0;;
- : int = 5
```

- Operators for floating point numbers:

```
# 1.2 +. 2.3;;
- : float = 3.5
# 1.5 *. 2.0;;
- : float = 3.
# float_of_int 1 +. 2.2;;
- : float = 3.2
```

# Other Primitive Values

- OCaml provides six primitive values: integers, booleans, floating point numbers, characters, strings, and unit.

```
# 'c';;
- : char = 'c'
# "cose212";;
- : string = "cose212"
# ();;
- : unit = ()
```

# Conditional Expressions

$$\texttt{if } \boldsymbol{be} \texttt{ then } \boldsymbol{e_1} \texttt{ else } \boldsymbol{e_2}$$

- If $\boldsymbol{be}$ is true, the value of the conditional expression is the value of $\boldsymbol{e_1}$.
- If $\boldsymbol{be}$ is false, the value of the expression is the value of $\boldsymbol{e_2}$.
- ```
  # if 2 > 1 then 0 else 1;;
  - : int = 0
  # if 2 < 1 then 0 else 1;;
  - : int = 1
  ```
  ```

# Conditional Expressions

$$\texttt{if } \boldsymbol{be} \texttt{ then } \boldsymbol{e_1} \texttt{ else } \boldsymbol{e_2}$$

- If $\boldsymbol{be}$ is true, the value of the conditional expression is the value of $\boldsymbol{e_1}$.
- If $\boldsymbol{be}$ is false, the value of the expression is the value of $\boldsymbol{e_2}$.
- ```
  # if 2 > 1 then 0 else 1;;
  - : int = 0
  # if 2 < 1 then 0 else 1;;
  - : int = 1
  ```
- $\boldsymbol{be}$ must be a boolean expression.
- types of $\boldsymbol{e_1}$ and $\boldsymbol{e_2}$ must be equivalent.
- ```
  # if 1 then 1 else 2;;
  Error: ...
  # if true then 1 else true;;
  Error: ...
  # if true then true else false;;
  - : bool = true
  ```

# Names and Functions

- Create a global variable with the `let` keyword:

```
# let x = 3 + 4;;
val x : int = 7
```

  We say a variable `x` is *bound* to value $\mathbf{7}$.

```
# let y = x + x;;
val y : int = 14
```

# Names and Functions

- Create a global variable with the `let` keyword:

  ```
  # let x = 3 + 4;;
  val x : int = 7
  ```

  We say a variable `x` is *bound* to value $\mathbf{7}$.

  ```
  # let y = x + x;;
  val y : int = 14
  ```

- Create a local variable with `let ... in ...` construct:

$$\texttt{let } \boldsymbol{x} \texttt{ = } \boldsymbol{e_1} \texttt{ in } \boldsymbol{e_2}$$

  - $\boldsymbol{x}$ is bound to the value of $\boldsymbol{e_1}$
  - the scope of $\boldsymbol{x}$ is $\boldsymbol{e_2}$
  - the value of $\boldsymbol{e_2}$ becomes the value of the entire expression

# Examples

- ```
  # let a = 1 in a;;
  - : int = 1
  # let a = 1 in a * 2;;
  - : int = 2
  ```
- ```
  # let a = 1 in
    let b = a + a in
    let c = b + b in
      c + c;;
  - : int = 8
  ```
- ```
  # let d =
      let b = a + a in
      let c = b + b in
        c + c;;
   val d : int = 8
  ```
  ```

# Functions

- Define a function with `let`:

```
# let square x = x * x;;
 val square : int -> int = <fun>
```

# Functions

- Define a function with `let`:

```
# let square x = x * x;;
 val square : int -> int = <fun>
```

- Apply the function:

```
# square 2;;
- : int = 4
# square (2 + 5);;
- : int = 49
# square (square 2);;
- : int = 16
```

# Functions

- Define a function with `let`:

```
# let square x = x * x;;
 val square : int -> int = <fun>
```

- Apply the function:

```
# square 2;;
- : int = 4
# square (2 + 5);;
- : int = 49
# square (square 2);;
- : int = 16
```

- The body can be any expression:

```
# let neg x = if x < 0 then true else false;;
val neg : int -> bool = <fun>
# neg 1;;
- : bool = false
# neg (-1);;
- : bool = true
```

# Functions

- Functions with multiple arguments:

```
# let sum_of_squares x y = (square x) + (square y);;
val sum_of_squares : int -> int -> int = <fun>
# sum_of_squares 3 4;;
- : int = 25
```

# Functions

- Functions with multiple arguments:

```
# let sum_of_squares x y = (square x) + (square y);;
val sum_of_squares : int -> int -> int = <fun>
# sum_of_squares 3 4;;
- : int = 25
```

- Recursive functions are defined with `let rec` construct:

```
# let rec factorial a =
    if a = 1 then 1 else a * factorial (a - 1);;
val factorial : int -> int = <fun>
# factorial 5;;
- : int = 120
```

# Nameless Functions

- Many modern programming languages provide nameless functions, e.g., ML, Scala, Java8, JavaScript, Python, etc.
- In OCaml, a function can be defined without names:

  ```
  # fun x -> x * x;;
  - : int -> int = <fun>
  ```

  Called *nameless* or *anonymous* functions.

# Nameless Functions

- Many modern programming languages provide nameless functions, e.g., ML, Scala, Java8, JavaScript, Python, etc.
- In OCaml, a function can be defined without names:

```
# fun x -> x * x;;
- : int -> int = <fun>
```

  Called *nameless* or *anonymous* functions.
- Apply nameless function as usual:

```
# (fun x -> x * x) 2;;
- : int = 4
```

# Nameless Functions

- Many modern programming languages provide nameless functions, e.g., ML, Scala, Java8, JavaScript, Python, etc.
- In OCaml, a function can be defined without names:

  ```
  # fun x -> x * x;;
  - : int -> int = <fun>
  ```

  Called *nameless* or *anonymous* functions.
- Apply nameless function as usual:

  ```
  # (fun x -> x * x) 2;;
  - : int = 4
  ```

- A variable can be bound to functions:

  ```
  # let square = fun x -> x * x;;
  val square : int -> int = <fun>
  ```

# Nameless Functions

- Many modern programming languages provide nameless functions, e.g., ML, Scala, Java8, JavaScript, Python, etc.
- In OCaml, a function can be defined without names:

```
# fun x -> x * x;;
- : int -> int = <fun>
```

  Called *nameless* or *anonymous* functions.
- Apply nameless function as usual:

```
# (fun x -> x * x) 2;;
- : int = 4
```

- A variable can be bound to functions:

```
# let square = fun x -> x * x;;
val square : int -> int = <fun>
```

- The followings are equivalent:

```
let square = fun x -> x * x
```

```
let square x = x * x
```

# Functions are First-Class in OCaml

In programming languages, a value is *first-class*, if the value can be

- stored in a variable,
- passed as an argument of a function, and
- returned from other functions.

A language is often called *functional*, if functions are first class values, e.g., ML, Scala, Java8, JavaScript, Python, Lisp, etc.

# Functions are First-Class in OCaml

- Functions can be stored in variables:

```
# let square = fun x -> x * x;;
# square 2;;
- : int = 4
```

# Functions are First-Class in OCaml

- Functions can be stored in variables:

```
# let square = fun x -> x * x;;
# square 2;;
- : int = 4
```

- Functions can be passed to other functions:

```
# let sum_if_true test first second =
    (if test first then first else 0)
  + (if test second then second else 0);;
val sum_if_true : (int -> bool) -> int -> int -> int = <fun>
```

```
# let even x = x mod 2 = 0;;
val even : int -> bool = <fun>
# sum_if_true even 3 4;;
- : int = 4
# sum_if_true even 2 4;;
- : int = 6
```

# Functions are First-Class in OCaml

- Functions can be also returned from a procedure:

```
# let plus_a a = fun b -> a + b;;
val plus_a : int -> int -> int = <fun>
```

```
# let f = plus_a 3;;
val f : int -> int = <fun>
# f 1;;
- : int = 4
# f 2;;
- : int = 5
```

# Functions are First-Class in OCaml

- Functions can be also returned from a procedure:

```
# let plus_a a = fun b -> a + b;;
val plus_a : int -> int -> int = <fun>
```

```
# let f = plus_a 3;;
val f : int -> int = <fun>
# f 1;;
- : int = 4
# f 2;;
- : int = 5
```

Functions that manipulate functions are called *higher-order functions*.

- i.e., functions that take as argument functions or return functions
- greatly increase the expressiveness of the language

# Pattern Matching

- An elegant way of doing case analysis.
- E.g., using pattern-matching, the factorial function

  ```
  let rec factorial a =
    if a = 1 then 1 else a * factorial (a - 1)
  ```

  can be written as follows:

  ```
  let factorial a =
    match a with
     1 -> 1
    |_ -> a * factorial (a - 1)
  ```

# Pattern Matching

The nested if-then-else expression

```
let isabc c = if c = 'a' then true
              else if c = 'b' then true
              else if c = 'c' then true
              else false
```

can be written using pattern matching:

```
let isabc c =
  match c with
  'a' -> true
 |'b' -> true
 |'c' -> true
 | _  -> false
```

or simply,

```
 let isabc c =
  match c with
  'a' | 'b' | 'c' -> true
  | _ -> false
```

# Type Inference

In C or Java, types must be annotated:

```
public static int f(int n)
{
  int a = 2;
  return a * n;
}
```

# Type Inference

In C or Java, types must be annotated:

```
public static int f(int n)
{
  int a = 2;
  return a * n;
}
```

In OCaml, type annotations are not mandatory:

```
# let f n =
    let a = 2 in
      a * n;;
val f : int -> int = <fun>
```

# Type Inference

OCaml can infer types, no matter how complex the program is:

```
# let sum_if_true test first second =
    (if test first then first else 0)
  + (if test second then second else 0);;
val sum_if_true : (int -> bool) -> int -> int -> int = <fun>
```

# Type Inference

OCaml can infer types, no matter how complex the program is:

```
# let sum_if_true test first second =
    (if test first then first else 0)
  + (if test second then second else 0);;
val sum_if_true : (int -> bool) -> int -> int -> int = <fun>
```

OCaml compiler infers the type through the following reasoning steps:

# Type Inference

OCaml can infer types, no matter how complex the program is:

```
# let sum_if_true test first second =
    (if test first then first else 0)
  + (if test second then second else 0);;
val sum_if_true : (int -> bool) -> int -> int -> int = <fun>
```

OCaml compiler infers the type through the following reasoning steps:

1. the types of `first` and `second` must be `int`, because both branches of a conditional expression must have the same type,

# Type Inference

OCaml can infer types, no matter how complex the program is:

```
# let sum_if_true test first second =
    (if test first then first else 0)
  + (if test second then second else 0);;
val sum_if_true : (int -> bool) -> int -> int -> int = <fun>
```

OCaml compiler infers the type through the following reasoning steps:

1. the types of `first` and `second` must be `int`, because both branches of a conditional expression must have the same type,
2. the type of `test` is a function type $\alpha \to \beta$, because `test` is used as a function,

# Type Inference

OCaml can infer types, no matter how complex the program is:

```
# let sum_if_true test first second =
    (if test first then first else 0)
  + (if test second then second else 0);;
val sum_if_true : (int -> bool) -> int -> int -> int = <fun>
```

OCaml compiler infers the type through the following reasoning steps:

1. the types of `first` and `second` must be `int`, because both branches of a conditional expression must have the same type,
2. the type of `test` is a function type $\alpha \to \beta$, because `test` is used as a function,
3. $\alpha$ must be of `int`, because `test` is applied to `first`, a value of `int`,

# Type Inference

OCaml can infer types, no matter how complex the program is:

```
# let sum_if_true test first second =
    (if test first then first else 0)
  + (if test second then second else 0);;
val sum_if_true : (int -> bool) -> int -> int -> int = <fun>
```

OCaml compiler infers the type through the following reasoning steps:

1. the types of `first` and `second` must be `int`, because both branches of a conditional expression must have the same type,
2. the type of `test` is a function type $\alpha \to \beta$, because `test` is used as a function,
3. $\alpha$ must be of `int`, because `test` is applied to `first`, a value of `int`,
4. $\beta$ must be of `bool`, because conditions must be boolean expressions,

# Type Inference

OCaml can infer types, no matter how complex the program is:

```
# let sum_if_true test first second =
    (if test first then first else 0)
  + (if test second then second else 0);;
val sum_if_true : (int -> bool) -> int -> int -> int = <fun>
```

OCaml compiler infers the type through the following reasoning steps:

1. the types of `first` and `second` must be `int`, because both branches of a conditional expression must have the same type,
2. the type of `test` is a function type $\alpha \to \beta$, because `test` is used as a function,
3. $\alpha$ must be of `int`, because `test` is applied to `first`, a value of `int`,
4. $\beta$ must be of `bool`, because conditions must be boolean expressions,
5. the return value of the function has type `int`, because the two conditional expressions are of `int` and their addition gives `int`.

# Type Annotation

Explicit type annotations are possible:

```
# let sum_if_true (test : int -> bool) (x : int) (y : int) : int =
    (if test x then x else 0) + (if test y then y else 0);;
val sum_if_true : (int -> bool) -> int -> int -> int = <fun>
```

# Type Annotation

Explicit type annotations are possible:

```
# let sum_if_true (test : int -> bool) (x : int) (y : int) : int =
    (if test x then x else 0) + (if test y then y else 0);;
val sum_if_true : (int -> bool) -> int -> int -> int = <fun>
```

If the annotation is wrong, OCaml finds the error and report it:

```
# let sum_if_true (test : int -> int) (x : int) (y : int) : int =
    (if test x then x else 0) + (if test y then y else 0);;
Error: The expression (test x) has type int but an expression
was expected of type bool
```

# Polymorphic Types

- What is the type of the program?

  ```
  let id x = x
  ```

# Polymorphic Types

- What is the type of the program?

  ```
  let id x = x
  ```

  See how OCaml infers its type:

  ```
  # let id x = x;;
  val id : 'a -> 'a = <fun>
  ```

# Polymorphic Types

- What is the type of the program?

```
let id x = x
```

See how OCaml infers its type:

```
# let id x = x;;
val id : 'a -> 'a = <fun>
```

The function works for values of any type:

```
# id 1;;
- : int = 1
# id "abc";;
- : string = "abc"
# id true;;
- : bool = true
```

# Polymorphic Types

- What is the type of the program?

```
let id x = x
```

See how OCaml infers its type:

```
# let id x = x;;
val id : 'a -> 'a = <fun>
```

The function works for values of any type:

```
# id 1;;
- : int = 1
# id "abc";;
- : string = "abc"
# id true;;
- : bool = true
```

- Such a function is called *polymorphic* and `'a` is a *type variable*.

# Polymorphic Types

Quiz) What is the type of the function?

```
let first_if_true test x y =
  if test x then x else y
```

# Tuples

- An ordered collection of values, each of which can be a different types, e.g.,

```
# let x = (1, "one");;
val x : int * string = (1, "one")
```

# Tuples

- An ordered collection of values, each of which can be a different types, e.g.,

```
# let x = (1, "one");;
val x : int * string = (1, "one")
# let y = (2, "two", true);;
val y : int * string * bool = (2, "two", true)
```

# Tuples

- An ordered collection of values, each of which can be a different types, e.g.,

```
# let x = (1, "one");;
val x : int * string = (1, "one")
# let y = (2, "two", true);;
val y : int * string * bool = (2, "two", true)
```

- Extract each component using pattern-matching:

```
# let fst p = match p with (x,_) -> x;;
val fst : 'a * 'b -> 'a = <fun>
# let snd p = match p with (_,x) -> x;;
val snd : 'a * 'b -> 'b = <fun>
```

# Tuples

- An ordered collection of values, each of which can be a different types, e.g.,

```
# let x = (1, "one");;
val x : int * string = (1, "one")
# let y = (2, "two", true);;
val y : int * string * bool = (2, "two", true)
```

- Extract each component using pattern-matching:

```
# let fst p = match p with (x,_) -> x;;
val fst : 'a * 'b -> 'a = <fun>
# let snd p = match p with (_,x) -> x;;
val snd : 'a * 'b -> 'b = <fun>
```

or equivalently,

```
# let fst (x,_) = x;;
val fst : 'a * 'b -> 'a = <fun>
# let snd (_,x) = x;;
val snd : 'a * 'b -> 'b = <fun>
```

# Tuples

- Patterns can be used in `let`:

```
# let p = (1, true);;
val p : int * bool = (1, true)
# let (x,y) = p;;
val x : int = 1
val y : bool = true
```

# Lists

- A finite sequence of elements, each of which has the same type, e.g.,

  ```
  [1; 2; 3]
  ```

  is a list of integers:

# Lists

- A finite sequence of elements, each of which has the same type, e.g.,

  `[1; 2; 3]`

  is a list of integers:

```
# [1; 2; 3];;
- : int list = [1; 2; 3]
```

# Lists

- A finite sequence of elements, each of which has the same type, e.g.,

  `[1; 2; 3]`

  is a list of integers:

  ```
  # [1; 2; 3];;
  - : int list = [1; 2; 3]
  ```

  Note that
  - all elements must have the same type, e.g., `[1; true; 2]` is not a list,
  - the elements are ordered, e.g., `[1; 2; 3]` $\neq$ `[2; 3; 1]`, and
  - the first element is called *head*, the rest *tail*.

# Lists

- A finite sequence of elements, each of which has the same type, e.g.,

  `[1; 2; 3]`

  is a list of integers:

  ```
  # [1; 2; 3];;
  - : int list = [1; 2; 3]
  ```

  Note that

  - all elements must have the same type, e.g., `[1; true; 2]` is not a list,
  - the elements are ordered, e.g., `[1; 2; 3]` $\neq$ `[2; 3; 1]`, and
  - the first element is called *head*, the rest *tail*.

- `[]`: the empty list, i.e., nil. What are head and tail of `[]`?

# Lists

- A finite sequence of elements, each of which has the same type, e.g.,

  `[1; 2; 3]`

  is a list of integers:

  ```
  # [1; 2; 3];;
  - : int list = [1; 2; 3]
  ```

  Note that

  - all elements must have the same type, e.g., `[1; true; 2]` is not a list,
  - the elements are ordered, e.g., `[1; 2; 3]` $\neq$ `[2; 3; 1]`, and
  - the first element is called *head*, the rest *tail*.
- `[]`: the empty list, i.e., nil. What are head and tail of `[]`?
- `[5]`: a list with a single element. What are head and tail of `[5]`?

# List Examples

- ```
  # [1;2;3;4;5];;
  - : int list = [1; 2; 3; 4; 5]
  ```
  ```

# List Examples

- ```
  # [1;2;3;4;5];;
  - : int list = [1; 2; 3; 4; 5]
  ```
- ```
  # ["OCaml"; "Java"; "C"];;
  - : string list = ["OCaml"; "Java"; "C"]
  ```

# List Examples

- `# [1;2;3;4;5];;`
  `- : int list = [1; 2; 3; 4; 5]`
- `# ["OCaml"; "Java"; "C"];;`
  `- : string list = ["OCaml"; "Java"; "C"]`
- `# [(1,"one"); (2,"two"); (3,"three")];;`
  `- : (int * string) list = [(1, "one"); (2, "two"); (3, "three")`

# List Examples

- ```
  # [1;2;3;4;5];;
  - : int list = [1; 2; 3; 4; 5]
  ```
- ```
  # ["OCaml"; "Java"; "C"];;
  - : string list = ["OCaml"; "Java"; "C"]
  ```
- ```
  # [(1,"one"); (2,"two"); (3,"three")];;
  - : (int * string) list = [(1, "one"); (2, "two"); (3, "three")
  ```
- ```
  # [[1;2;3];[2;3;4];[4;5;6]];;
  - : int list list = [[1; 2; 3]; [2; 3; 4]; [4; 5; 6]]
  ```

# List Examples

- `# [1;2;3;4;5];;`
  `- : int list = [1; 2; 3; 4; 5]`
- `# ["OCaml"; "Java"; "C"];;`
  `- : string list = ["OCaml"; "Java"; "C"]`
- `# [(1,"one"); (2,"two"); (3,"three")];;`
  `- : (int * string) list = [(1, "one"); (2, "two"); (3, "three")]`
- `# [[1;2;3];[2;3;4];[4;5;6]];;`
  `- : int list list = [[1; 2; 3]; [2; 3; 4]; [4; 5; 6]]`
- `# [1;"OCaml";3] ;;`
  `Error: This expression has type string but an expression was expected of type int`

# List Operators

- `::` (cons): add a single element to the front of a list, e.g.,

```
# 1::[2;3];;
- : int list = [1; 2; 3]
# 1::2::3::[];;
- : int list = [1; 2; 3]
```

  (`[1; 2; 3]` is a shorthand for `1::2::3::[]`)

# List Operators

- `::` (cons): add a single element to the front of a list, e.g.,

  ```
  # 1::[2;3];;
  - : int list = [1; 2; 3]
  # 1::2::3::[];;
  - : int list = [1; 2; 3]
  ```

  (`[1; 2; 3]` is a shorthand for `1::2::3::[]`)
- `@` (append): combine two lists, e.g.,

  ```
  # [1; 2] @ [3; 4; 5];;
  - : int list = [1; 2; 3; 4; 5]
  ```

# Patterns for Lists

Pattern matching is useful for manipulating lists.

- A function to check if a list is empty:

```
# let isnil l =
    match l with
     [] -> true
    |_  -> false;;
val isnil : 'a list -> bool = <fun>
# isnil [1];;
- : bool = false
# isnil [];;
- : bool = true
```

# Patterns for Lists

- A function that computes the length of lists:

```
# let rec length l =
    match l with
     [] -> 0
    |h::t -> 1 + length t;;
val length : 'a list -> int = <fun>
# length [1;2;3];;
- : int = 3
```

# Patterns for Lists

- A function that computes the length of lists:

```
# let rec length l =
    match l with
     [] -> 0
    |h::t -> 1 + length t;;
val length : 'a list -> int = <fun>
# length [1;2;3];;
- : int = 3
```

We can replace pattern `h` by `_`:

```
let rec length l =
  match l with
   [] -> 0
  |_::t -> 1 + length t;;
```

# Data Types

- If data elements are finite, just enumerate them, e.g., “days”:

```
# type days = Mon | Tue | Wed | Thu | Fri | Sat | Sun;;
type days = Mon | Tue | Wed | Thu | Fri | Sat | Sun
```

# Data Types

- If data elements are finite, just enumerate them, e.g., “days”:

```
# type days = Mon | Tue | Wed | Thu | Fri | Sat | Sun;;
type days = Mon | Tue | Wed | Thu | Fri | Sat | Sun
```

  Construct values of the type:

```
# Mon;;
- : days = Mon
# Tue;;
- : days = Tue
```

# Data Types

- If data elements are finite, just enumerate them, e.g., "days":

```
# type days = Mon | Tue | Wed | Thu | Fri | Sat | Sun;;
type days = Mon | Tue | Wed | Thu | Fri | Sat | Sun
```

Construct values of the type:

```
# Mon;;
- : days = Mon
# Tue;;
- : days = Tue
```

A function that manipulates the defined data:

```
# let nextday d =
  match d with
  | Mon -> Tue | Tue -> Wed | Wed -> Thu | Thu -> Fri
  | Fri -> Sa  | Sat -> Sun | Sun -> Mon ;;
val nextday : days -> days = <fun>
# nextday Mon;;
- : days = Tue
```

# Data Types

- Constructors can be associated with values, e.g.,

```
# type shape = Rect of int * int | Circle of int;;
type shape = Rect of int * int | Circle of int
```

# Data Types

- Constructors can be associated with values, e.g.,

```
# type shape = Rect of int * int | Circle of int;;
type shape = Rect of int * int | Circle of int
```

Construct values of the type:

```
# Rect (2,3);;
- : shape = Rect (2, 3)
# Circle 5;;
- : shape = Circle 5
```

# Data Types

- Constructors can be associated with values, e.g.,

```
# type shape = Rect of int * int | Circle of int;;
type shape = Rect of int * int | Circle of int
```

Construct values of the type:

```
# Rect (2,3);;
- : shape = Rect (2, 3)
# Circle 5;;
- : shape = Circle 5
```

A function that manipulates the data:

```
# let area s =
    match s with
      Rect (w,h) -> w * h
    | Circle r -> r * r * 3;;
val area : shape -> int = <fun>
# area (Rect (2,3));;
- : int = 6
# area (Circle 5);;
- : int = 75
```

# Data Types

- Inductive data types, e.g.,

```
# type mylist = Nil | List of int * mylist;;
type mylist = Nil | List of int * mylist
```

# Data Types

- Inductive data types, e.g.,

  ```
  # type mylist = Nil | List of int * mylist;;
  type mylist = Nil | List of int * mylist
  ```

  Construct values of the type:

  ```
  # Nil;;
  - : mylist = Nil
  # List (1, Nil);;
  - : mylist = List (1, Nil)
  # List (1, List (2, Nil));;
  - : mylist = List (1, List (2, Nil))
  ```

# Data Types

- Inductive data types, e.g.,

```
# type mylist = Nil | List of int * mylist;;
type mylist = Nil | List of int * mylist
```

Construct values of the type:

```
# Nil;;
- : mylist = Nil
# List (1, Nil);;
- : mylist = List (1, Nil)
# List (1, List (2, Nil));;
- : mylist = List (1, List (2, Nil))
```

A function that manipulates the data:

```
# let rec mylength l =
    match l with
     Nil -> 0
    |List (_,l') -> 1 + mylength l';;
val mylength : mylist -> int = <fun>
# mylength (List (1, List (2, Nil)));;
- : int = 2
```

# Data Types

- Another inductive data type: binary tree

```
# type btree = Empty | Node of int * btree * btree
```

# Data Types

- Another inductive data type: binary tree

  ```
  # type btree = Empty | Node of int * btree * btree
  ```

  Construct values of the type:

  ```
  # Empty;;
  - : btree = Empty
  # let t1 = Node (1, Empty, Empty);;
  val t1 : btree = Node (1, Empty, Empty)
  # let t2 = Node (1, t1, Node (3, Empty,Empty));;
  val t2 : btree = Node (1, Node (1, E, E), Node (3, E, E))
  ```

# Data Types

- Another inductive data type: binary tree

  ```
  # type btree = Empty | Node of int * btree * btree
  ```

  Construct values of the type:

  ```
  # Empty;;
  - : btree = Empty
  # let t1 = Node (1, Empty, Empty);;
  val t1 : btree = Node (1, Empty, Empty)
  # let t2 = Node (1, t1, Node (3, Empty,Empty));;
  val t2 : btree = Node (1, Node (1, E, E), Node (3, E, E))
  ```

  Write function mem that checks if an integer element is in a tree:

  ```
  # let rec mem: int -> btree -> bool =
    fun n t -> ...
  val mem : int -> btree -> bool = <fun>
  # mem 1 t1;;
  - : bool = true
  # mem 4 t2;;
  - : bool = false
  ```

# Exceptions

- An exception means a run-time error: e.g.,

```
# let div a b = a / b;;
val div : int -> int -> int = <fun>
# div 10 5;;
- : int = 2
# div 10 0;;
Exception: Division_by_zero.
```

# Exceptions

- An exception means a run-time error: e.g.,

```
# let div a b = a / b;;
val div : int -> int -> int = <fun>
# div 10 5;;
- : int = 2
# div 10 0;;
Exception: Division_by_zero.
```

- The exception can be handled with `try ... with` constructs.

```
# let div a b =
    try
      a / b
    with Division_by_zero -> 0;;
val div : int -> int -> int = <fun>
# div 10 5;;
- : int = 2
# div 10 0;;
- : int = 0
```

# Exceptions

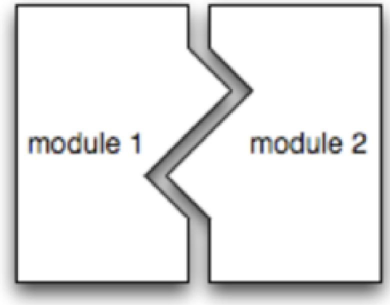
- User-defined exceptions: e.g.,

```
# exception Problem;;
exception Problem
# let div a b =
    if b = 0 then raise Problem
    else a / b;;
val div : int -> int -> int = <fun>
# div 10 5;;
- : int = 2
# div 10 0;;
Exception: Problem.
# try
    div 10 0
  with Problem -> 0;;
- : int = 0
```

# Module System

ML supports modular programming where code comprises independent modules

- developed separately
- understand behavior of module in isolation
- reason locally, not globally



# Module System

Features that support modularity

- *Namespaces*: enables the name "foo" in one namespace to have a distinct meaning from "foo" in another namespace
- *Abstraction*: hides some information while revealing other information
- *Code reuse*: enables code from one module to be used as part of another module without having to copy that code

Example: Java features for modularity

- *Namespaces*: class, package
- *Abstraction*: interface, access modifiers (private, protected)
- *Code reuse*: inheritance

# Module System

An elegant module system:

- *Structure* is a collection of types, exceptions, values, and functions, i.e., implementation details.
- *Signature* is the *interface* of the structure.
- *Functor* enables code reuse (not covered in this class)

# Example

The interface of a queue data structure:

- empty: the empty queue
- isempty: the boolean-valued test of whether q is empty
- enq(q,x): the queue obtained by inserting x on the end of q
- deq(q): the queue obtained by removing the front element of q (also returns the front element)
- print(q): show the contents of q
- E: the exception raised by deq if the queue is empty

# Example

The signature of the queue data structure:

```
module type IntQueue =
sig
  type t
  exception E
  val empty : t
  val is_empty : t -> bool
  val enq : t -> int -> t
  val deq : t -> int * t
  val print : t -> unit
end
```

# Example

An implementation:

```
module IntQueue : IntQueue =
struct
  type t = int list
  exception E
  let empty = []
  let enq q x = q @ [x]
  let is_empty q = q = []
  let deq q = match q with [] -> raise E | h::t -> (h, t)
  let rec print q =
    match q with
     [] -> print_string "\n"
    |h::t -> print_int h; print_string " "; print t
end
```

# Example

The module can be used as follows:

```
let q0 = IntQueue.empty
let q1 = IntQueue.enq q0 1
let q2 = IntQueue.enq q1 2
let (_,q3) = IntQueue.deq q2
let _  = IntQueue.print q1
let _  = IntQueue.print q2
let _  = IntQueue.print q3
```

The program prints:

```
1
1 2
2
```

# Example

The OCaml module system ensures the abstraction layer of the program:

```
let q4 = q1 @ [2]
```

produces a compile error:

```
Error: This expression has type IntQueue.t
       but an expression was expected of type 'a list
```