COMP 3311: Database Management Systems

Lecture 21 Exercises Concurrency Control: Lock-based Protocols

Exercise 1: a) Is the schedule conflict serializab					le?		Yes	□ No					
	If yes, gi	ive th	ne equivale	nt serial sch	nedu	ıle:							
b)	Rewrite	the	schedule	according	to	strict	2PL	by	adding	lock-s(),	lock-x()	and	unlock()

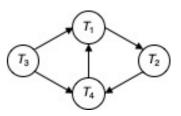
<i>T</i> ₁	<i>T</i> ₂	<i>T</i> ₃
read(X)		
	read(Y)	
	write(Y)	
		write(Z)
write(X)		
	read(X)	
	write(X)	
		read(Y)
		write(Y)
write(Z)		

Exercise 2: Which of the following statements is true about the wait-for graph (circle the correct answer)?

a) T_4 is waiting for T_3 to release a data item.

instructions.

- b) The system is in a deadlock state after removing the edge between T_2 and T_4 .
- c) The system is in a deadlock state after removing the edge between T_3 and T_4 .
- d) The system is in a deadlock state when T_1 no longer holds a data item needed by T_4 .



Name:		<i>I</i>	Student#:	Date:
_	Family/Last (PRINT)	Given/First (PRINT)		

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Exercise 3: Rewrite the following schedule according to $\underline{\mathsf{2PL}}$ by adding lock-s(), lock-x() and unlock() instructions. Is the schedule serializable? \square Yes \square No

<i>T</i> ₁	<i>T</i> ₂	<i>T</i> ₃
read(X)		
	read(X)	
		read(Y)
read(Z)		
	read(Y)	
	write(X)	
		read(X)
		write(X)
write(Z)		

Exercise 4: In which positions, A to E, can an unlock(X) instruction be inserted if the schedule is according to:

- a) strict 2PL (circle the correct answer)
 - i. {A} {B} {C} {D}
 - ii. {A} {B} {C} {D} {E}
 - iii. {A} {C} {D}
 - iv. {B} {E}
 - v. {A} {C} {D} {E}
- b) rigorous 2PL (circle the correct answer)
 - i. {A} {B} {C} {D}
 - ii. {A} {B} {C} {D} {E}
 - iii. $\{A\} \{C\} \{D\}$
 - iv. $\{B\}\{E\}$
 - v. {A} {C} {D} {E}

<i>T</i> ₁	T ₂
lock-s(X)	
read(X)	
	lock-s(X)
lock-x(Y)	
{A}	
read(Y)	
write(Y)	
	read(X)
	{C}
commit	
unlock(Y)	
{B}	
	{D}
	commit
	{E}

Ex	ercise 5: Consider the schedule shown below.
a)	Is the schedule conflict serializable? ☐ Yes ☐ No
	If yes, give the equivalent serial schedule
b)	If T_3 aborts after write(Y), which other transactions will be rolled back?
c)	If T_1 aborts after write(X), which other transactions will be rolled back?
d)	Draw the wait-for graph that results from this schedule if all locks are only exclusive-locks (lock x), no locks are released and the execution process runs to the point of lock- $x(Y)$ in \mathcal{T}_1 .
	$\overline{\tau_1}$

Wait-for Graph

e) Add lock-s(), lock-x() and unlock() instructions to the schedule below according to strict 2PL.

<i>T</i> ₁	T_2	<i>T</i> ₃	<i>T</i> ₄
read(X)			
write(X)			
	read(X)		
		read(Y)	
		write(Y)	
	write(X)		
			read(Y)
write(Y)			