COMP 3311: Database Management Systems

Lecture 22 Exercises Concurrency Control: Timestamp-based Protocols

Exercise 1: Use the <u>single version</u>, <u>timestamp-ordering protocol</u> to complete the following non-serializable schedule assuming the timestamps 1, 2, and 3 for transactions T_1 , T_2 , and T_3 , respectively. Show where the protocol will fail. Assume initial R/W timestamp of all items is 0.

T ₁ [TS=1]	T ₂ [TS=2]	T ₃ [TS=3]
read(X)		
	read(Y)	
	write(Y)	
		write(Z)
write(X)		
	read(X)	
	write(X)	
		read(Y)
		write(Y)
write(Z)		

Exercise 2: Use the <u>multi-version</u>, <u>timestamp-ordering protocol</u> to complete the schedule of Exercise 1 assuming the timestamps 1, 2, and 3 for transactions T_1 , T_2 , and T_3 , respectively. Assume initial R/W timestamp of all items is 0.

X		
$RTS(X_0) = 0$ $WTS(X_0) = 0$		
RTS(X ₁) =	$WTS(X_1) =$	

Υ		
$RTS(Y_0) = 0 \qquad WTS(Y_0) = 0$		
RTS(Y ₁) =	$WTS(Y_1) =$	

Z			
$RTS(Z_0) = 0 \qquad WTS(Z_0) = 0$			

T ₁ [TS=1]	T ₂ [TS=2]	<i>T</i> ₃ [TS=3]
read(X)		
	read(Y)	
	write(Y)	
		write(Z)
write(X)		
	read(X)	
	write(X)	
		read(Y)
		write(Y)
write(Z)		

Name:		<u> </u>	Student#:	Date:
	Family/Last (PRINT)	Given/First (PRINT)		

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Exercise 3: The following schedule is conflict serializable.

- (a) What is the equivalent serial schedule?
- (b) Assign appropriate timestamps to the transactions T_1 , T_2 , T_3 and T_4 so that the schedule is conflict serializable according to the <u>single version</u>, <u>timestamp-based protocol</u>. Assume initial R/W timestamp of all items is 0.

T ₁ [TS=]	T ₂ [TS=]	T ₃ [TS=]	T ₄ [TS=]
read(X)				
write(X)				
		read(X)		
			read(Y)	
			write(Y)	
		write(X)		
				read(Y)
write(Y)				

Exercise 4: Use the <u>multi version, timestamp-ordering protocol</u> to complete the conflict serializable schedule of Exercise 3 assuming the timestamps 1, 2, 3, and 4 for transactions T_1 , T_2 , T_3 and T_4 , respectively. Show where the protocol will fail. Assume initial R/W timestamp of all items is 0.

X		
$RTS(X_0) = 0$	$WTS(X_0) = 0$	
RTS(X ₁) =	$WTS(X_1) =$	
RTS(X ₂) =	$WTS(X_2) =$	

Υ		
$RTS(Y_0) = 0$	$WTS(Y_0) = 0$	
RTS(Y ₁) =	$WTS(Y_1) =$	
RTS(Y ₂) =	WTS(Y ₂) =	

T ₁ [TS=1]	T ₂ [TS=2]	T ₃ [TS=3]	T ₄ [TS=4]
read(X)			
write(X)			
	read(X)		
		read(Y)	
		write(Y)	
	write(X)		
			read(Y)
write(Y)			, ,