



Cryptographic hash functions

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What is cryptographic hash function?

$$h = H(M)$$

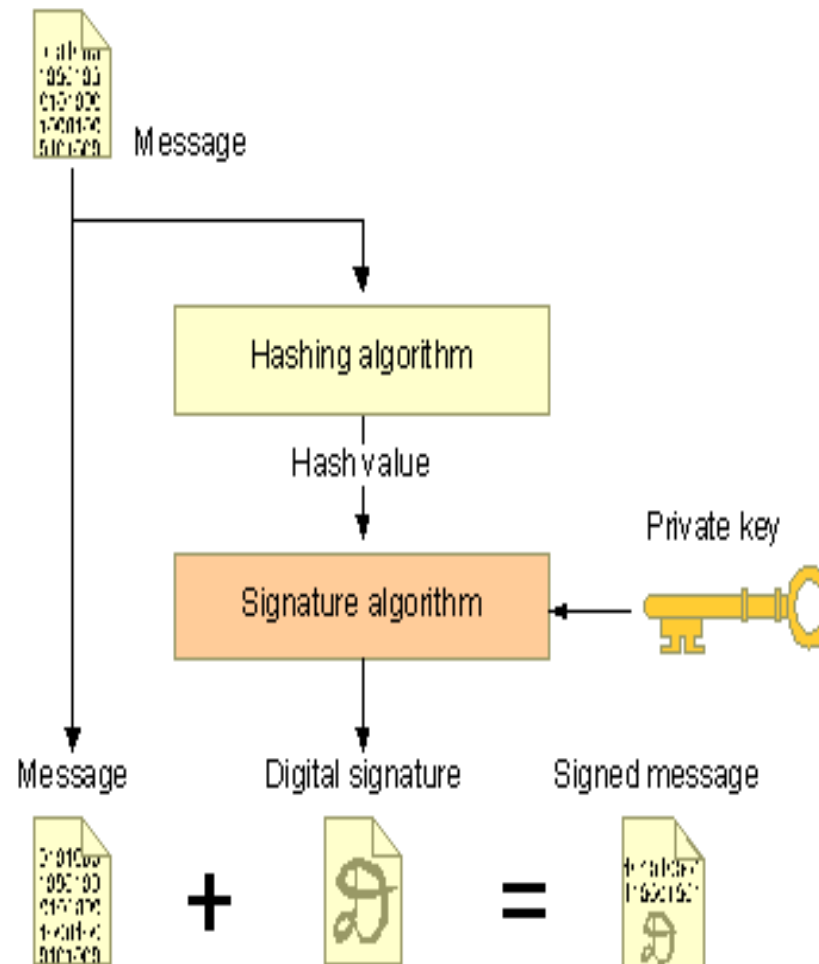
- Regular hash function (used for hash table)
 - Large input domain \rightarrow small fixed output
 - Well distributed: $P(H(x)=i) \approx 1/N$

This is not enough for cryptographic hash functions!

Hash function motivation

- Suppose Alice signs M
 - Alice sends M and $S = [M]_{\text{Alice}}$ to Bob
 - Bob verifies that $M = \{S\}_{\text{Alice}}$
- If M is big, $[M]_{\text{Alice}}$ costly to **compute & send**
- Suppose instead, Alice signs $h(M)$ where $h(M)$ is much smaller than M
 - Alice sends M and $S = [h(M)]_{\text{Alice}}$ to Bob
 - Bob verifies that $h(M) = \{S\}_{\text{Alice}}$

- First, create a message digest using a cryptographic hash
- Then, sign the message digest with your private key



Cryptographic hash functions

- A cryptographic hash function distills a message M down to a hash $H(M)$
- Additional properties
 1. **Pre-image** resistance:
 - Given $h = H(x)$, hard to find x
 2. **Weak collision** resistance:
 - Given $h = H(x)$, hard to find any x' such that $H(x') = h$
 3. **Strong collision** resistance:
 - Hard to find any pair x and y such that $H(x) = H(y)$

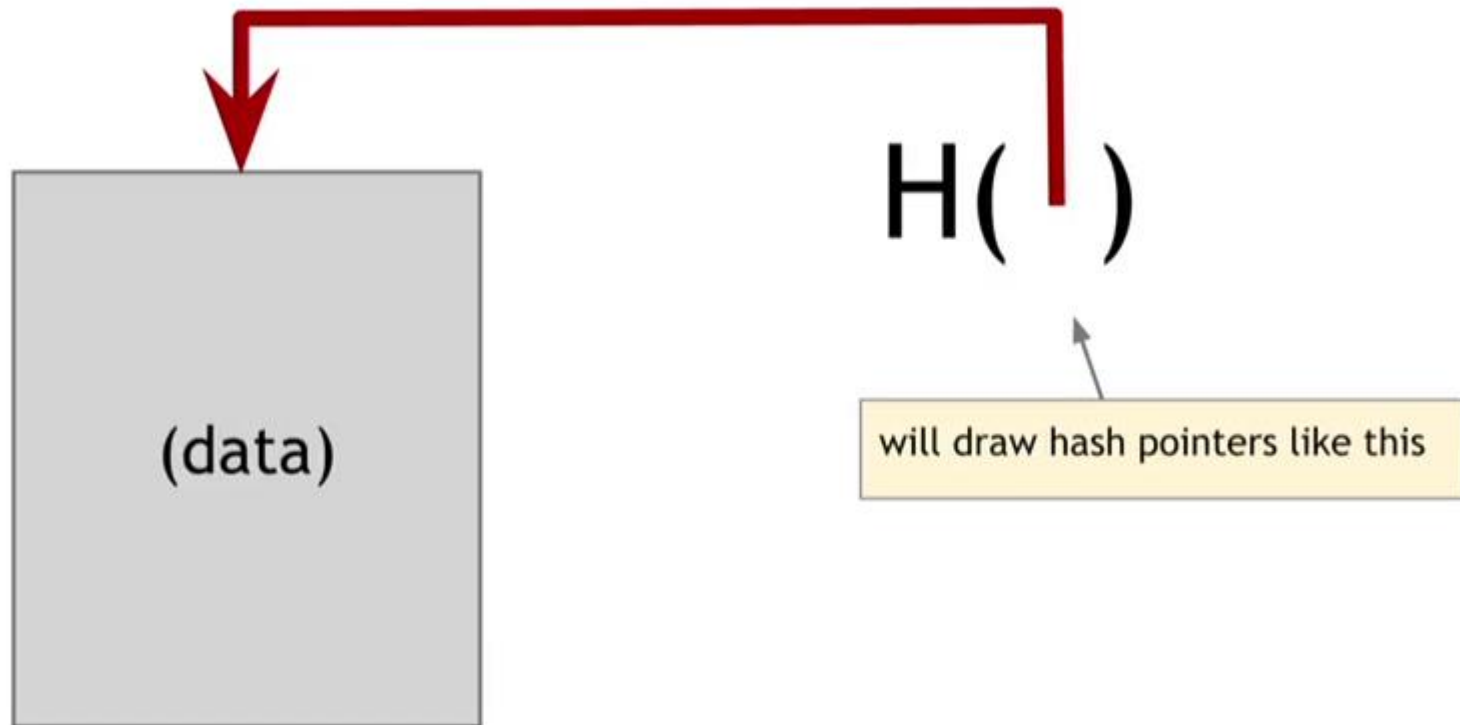
Quiz

Q. Which of these is almost a good cryptographic hash function?

- 1. Use CBC to encrypt x , take last output block
- 2. Use CTR to encrypt x , take last output block Previous blocks can be changed.
- 3. Use ECB to encrypt x , hash is XOR of all output blocks Same output if m_0 switches to m_1
- 4. Use CTR to encrypt x , hash is XOR of all output blocks No order!
Same output if m_0 switches to m_1

Hash pointer

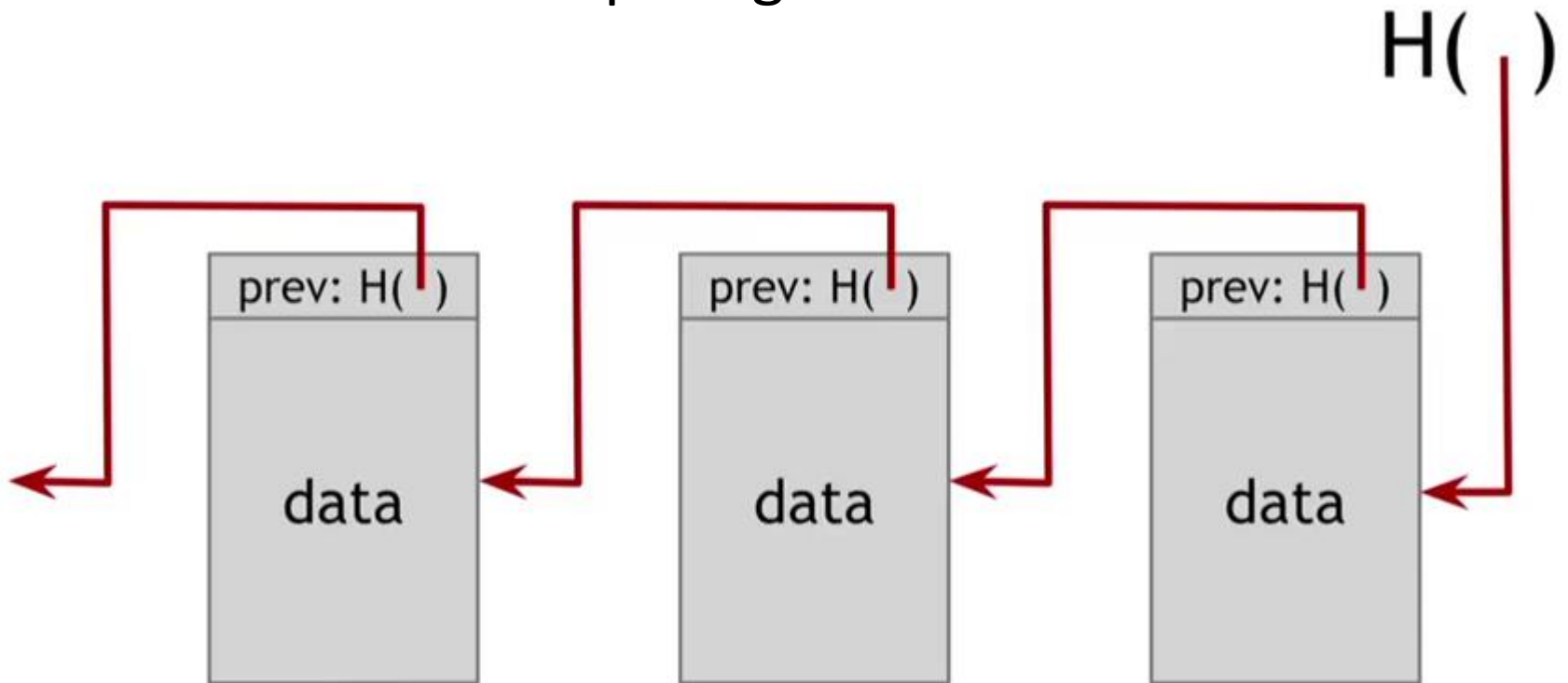
- Hash pointer is
 - pointer to where some information is stored, and
 - (cryptographic) hash of the information
- If we have a hash pointer, we can
 - ask to get the information back, and
 - verify that it hasn't changed



We can build data structures with hash pointers.

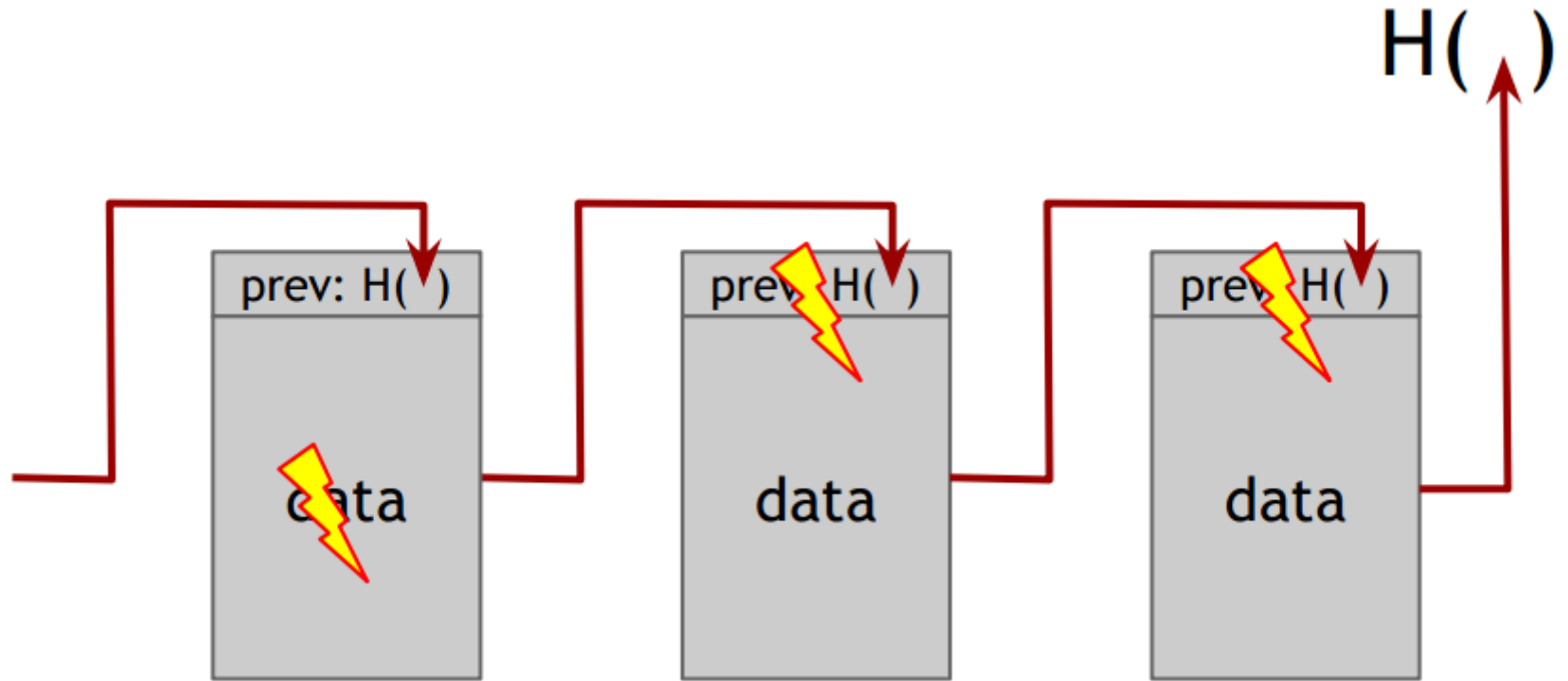
Hash chain

We can detect tampering.



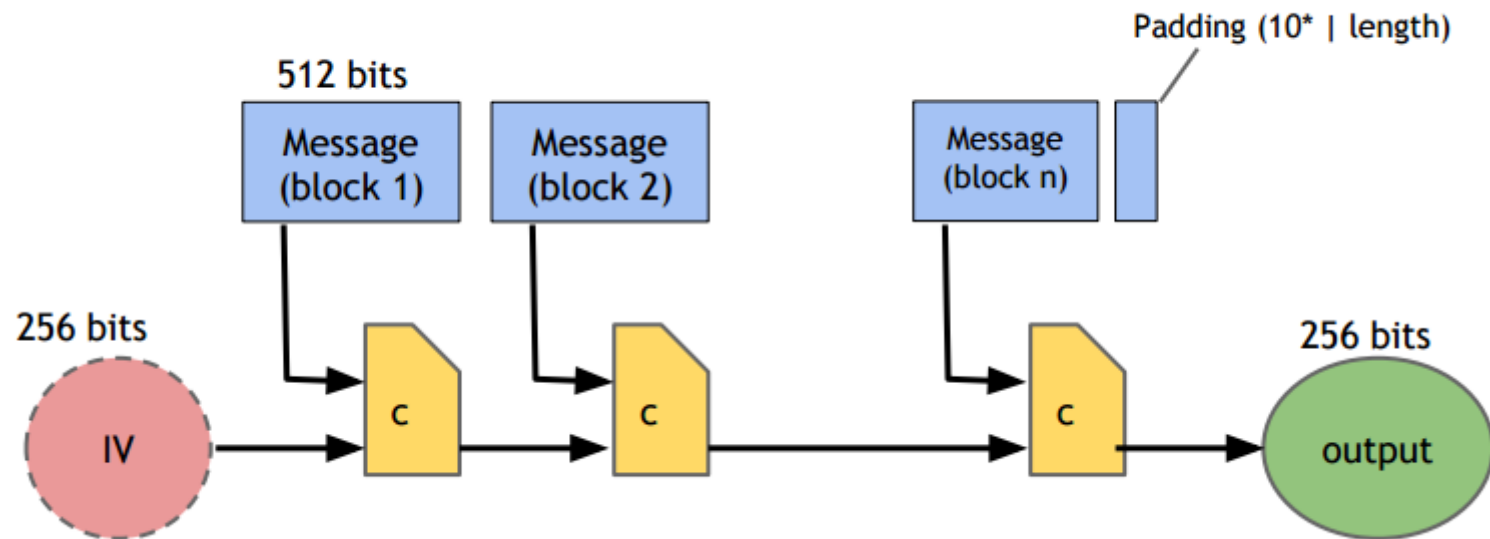
Use case: tamper-evident log

Propagation of modifications



Modifications to any data will propagate forever.

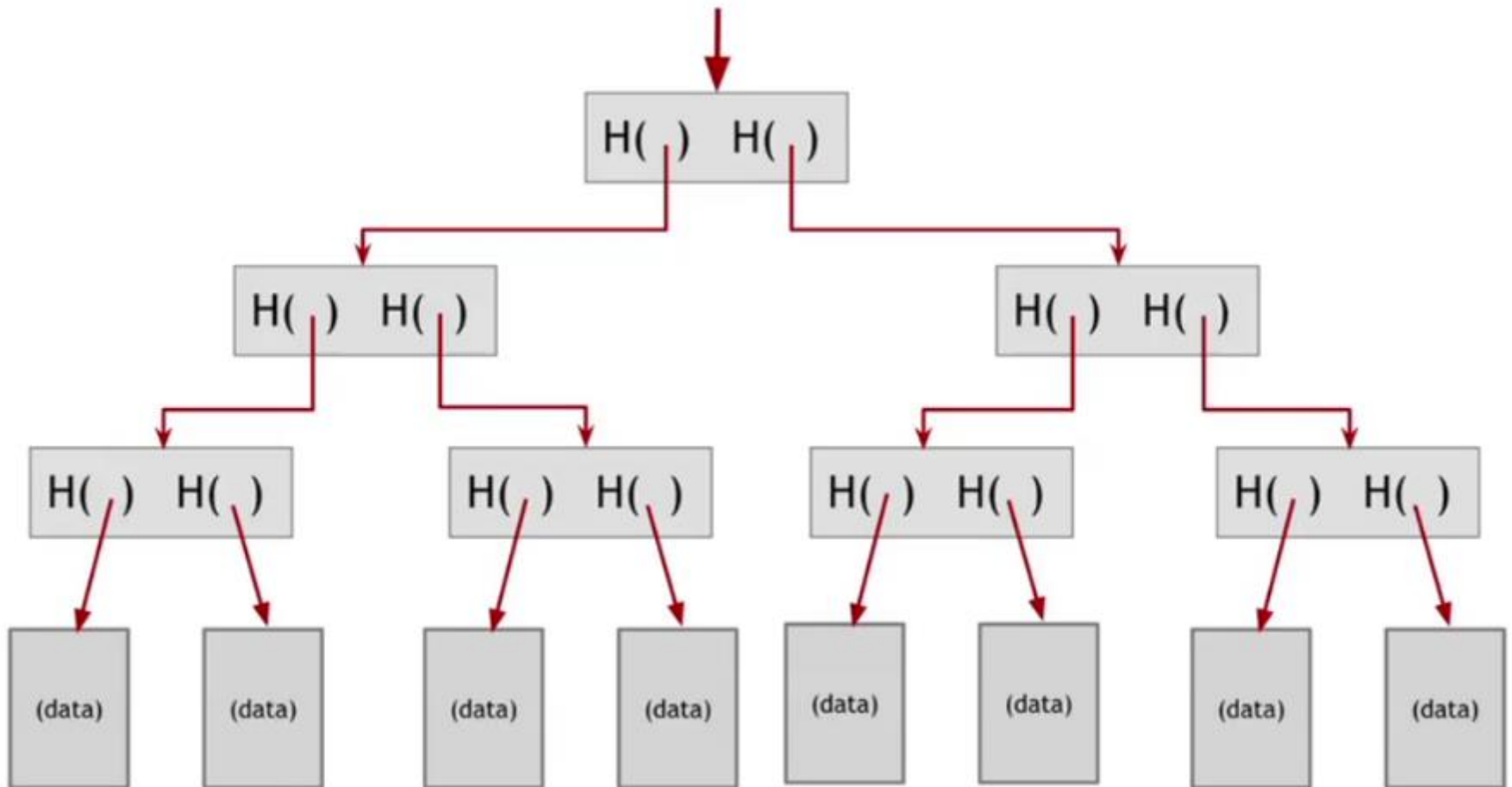
Merkle-Damgård construction (e.g., with SHA-256)



Theorem: If c is *collision-free*, then the hash is *collision-resistant*.

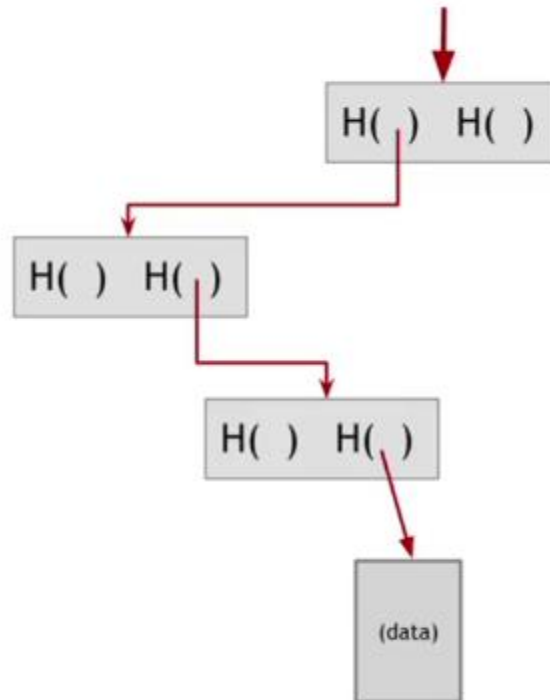
Binary tree with hash pointers

We call this “Merkle tree”.



Advantages of Merkle tree

- Tree holds many items but just need to remember the root hash.
- We can verify membership in $O(\log n)$ time/space.



Questions?

