

Multicore Computing Lecture 17 - MPI



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Deadlock Pitfall (1)

```
int a[10], b[10], myrank;
MPI_Status status;
...
MPI_Comm_rank(MPI_COMM_WORLD, &myrank);
if (myrank == 0) {
    MPI_Send(a, 10, MPI_INT, 1, 1, MPI_COMM_WORLD);
    MPI_Send(b, 10, MPI_INT, 1, 2, MPI_COMM_WORLD);
}
else if (myrank == 1) {
    MPI_Recv(b, 10, MPI_INT, 0, 2, MPI_COMM_WORLD);
    MPI_Recv(a, 10, MPI_INT, 0, 1, MPI_COMM_WORLD);
}
...
```

If MPI_Send is non-buffered blocking, a deadlock occurs.

MPI standard does not specify whether the implementation of MPI_Send is buffered blocking or not.



Circular communication

Once again, a deadlock occurs if MPI_Send is non-buffered blocking.



Avoiding Deadlocks

We can break the circular wait to avoid deadlocks as follows:

```
int a[10], b[10], npes, myrank;
MPI Status status;
MPI Comm size (MPI COMM WORLD, &npes);
MPI Comm rank (MPI COMM WORLD, &myrank);
if (myrank%2 == 1) {
       MPI Send(a, 10, MPI INT, (myrank+1)%npes, 1,
               MPI COMM WORLD);
       MPI Recv(b, 10, MPI INT, (myrank-1+npes) %npes, 1,
               MPI COMM WORLD);
else {
       MPI Recv(b, 10, MPI INT, (myrank-1+npes)%npes, 1,
               MPI COMM WORLD);
       MPI Send(a, 10, MPI INT, (myrank+1)%npes, 1,
               MPI COMM WORLD);
```



Sending and Receiving Messages Simultaneously

Exchange messages

```
int MPI_Sendrecv(void *sendbuf, int sendcount,
    MPI_Datatype senddatatype, int dest, int sendtag,
    void *recvbuf, int recvcount, MPI_Datatype
    recvdatatype, int source, int recvtag,
    MPI_Comm comm, MPI_Status *status)
```

- Requires both send and receive arguments
- Avoid the circular deadlock problem
- Exchange messages using the same buffer int MPI_Sendrecv_replace(void *buf, int count, MPI_Datatype datatype, int dest, int sendtag, int source, int recvtag, MPI_Comm comm, MPI Status *status)



Non-blocking Sending and Receiving Messages

Non-blocking send and receive operations int MPI_Isend(void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm, MPI_Request *request) int MPI_Irecv(void *buf, int count, MPI_Datatype datatype, int source, int tag, MPI_Comm comm, MPI_Request *request)

- Tests whether non-blocking operation is finished int MPI_Test(MPI_Request *request, int *flag, MPI Status *status)
- Waits for the operation to complete int MPI Wait(MPI Request *request, MPI Status *status)



Avoiding Deadlocks

Using non-blocking operations remove most deadlocks

```
int a[10], b[10], myrank;
MPI Status status;
MPI Request request1, request2;
MPI Comm rank(MPI COMM WORLD, &myrank);
if (myrank == 0) {
  MPI Isend(a, 10, MPI INT, 1, 1, MPI COMM WORLD, &request1);
  MPI Isend(b, 10, MPI INT, 1, 2, MPI COMM WORLD, &request2);
else if (myrank == 1) {
  MPI Irecv(b, 10, MPI INT, 0, 2, &status, MPI COMM WORLD, &request1);
  MPI Irecv(a, 10, MPI INT, 0, 1, &status, MPI COMM WORLD, &request2);
```

 Replacing either the send or the receive operations with non-blocking counterparts fixes this deadlock.



Example: Matrix-Matrix Multiplication

A _{0,0}	A _{0,1}	A _{0,2}	A _{0,3}
A _{1,0}	A _{1,1}	A _{1,2}	A _{1,3}
A _{2,0}	A _{2,1}	A _{2,2}	A _{2,3}
A _{3,0}	A _{3,1}	A _{3,2}	A _{3,3}

B _{0,0}	B _{0,1}	B _{0,2}	B _{0,3}
B _{1,0}	B _{1,1}	B _{1,2}	B _{1,3}
B _{2,0}	B _{2,1 Å}	B _{2,2}	B _{2,3}
B _{3,0}	$\mathbf{B}_{3,1}$	$\mathbf{B}_{3,2}$	B _{3,3}

(a) Initial alignment of A

(b) Initial alignment of B



Example: Matrix-Matrix Multiplication

	A	A	A	A
Κ.	$A_{0,0}$	$A_{0,1}$	$A_{0,2}$	A _{0,3} <
	$_{A}^{\prime}$ $B_{0,0}$	$_{\mathcal{A}}^{\mathbf{B}}\mathbf{B}_{1,1}$	$_{\mathcal{A}}$ $B_{2,2}$	$_{A}^{\circ}$ B _{3,3}
≪.	A _{1,1} <	$A_{1,2}$	A _{1,3} <	A _{1,0} <
	$_{\mathcal{A}}^{\mathbf{B}}\mathbf{B}_{1,0}$	$_{\mathcal{A}}^{B}$ $\mathbf{B}_{2,1}$	$_{\mathcal{A}}$ $B_{3,2}$	$_{A}^{\circ}$ $B_{0,3}$
∀ :·	A _{2,2}	A _{2,3}	A _{2,0}	A _{2,1}
	$_{\mathcal{A}}^{\mathbf{B}}\mathbf{B}_{2,0}$	$_{\mathcal{A}}^{B}\mathbf{B}_{3,1}$	$_{\mathcal{A}}^{B}\mathrm{B}_{0,2}$	$_{\mathcal{A}}^{B}$ $\mathbf{B}_{1,3}$
≪.	A _{3,3} <	A _{3,0} <	A _{3,1} <	A _{3,2} <
	$_{1}^{\prime}$ B _{3,0}	$_{\mathcal{I}}^{B}$ $\mathbf{B}_{0,1}$	$_{\mathcal{I}}^{B}$ $\mathbf{B}_{1,2}$	$_{1}^{\prime}$ $B_{2,3}$
	\	i,	i,	1,

4	A	A	A
$<$ $A_{0,1}$	$A_{0,2}$	$A_{0,3}$	A _{0,0} <
$\mathbf{B}_{1,0}$	$\mathbf{B}_{2,1}$	$_{A}$ B _{3,2}	$_{A}^{\circ}$ B _{0,3}
$A_{1,2}$	$A_{1,3}$	$A_{1,0}$	A _{1,1} <
$\mathbf{B}_{2,0}$	$\mathbf{B}_{3,1}$	$_{A}$ $B_{0,2}$	$_{\mathcal{A}}\mathbf{B}_{1,3}$
A _{2,3}	A _{2,0} <	$A_{2,1}$	A _{2,2} <
$\mathbf{B}_{3,0}$	$_{A}^{\mathbf{B}}\mathbf{B}_{0,1}$	$_{A}$ B _{1,2}	$_{\mathcal{A}}^{B}$ B _{2,3}
A _{3,0}	$A_{3,1}$	$A_{3,2}$	A _{3,3} <
$_{\mathcal{A}}\mathbf{B}_{0,0}$	$_{\nearrow}$ $B_{1,1}$	$_{\nearrow}$ $B_{2,2}$	$_{1}^{\prime}$ $B_{3,3}$
		i ,	ţ

(c) A and B after initial alignment

(d) Submatrix locations after first shift



Example: Matrix-Matrix Multiplication

	A	A	A	1	
≪:	A _{0,2} <	A _{0,3} <	A _{0,0} <	A _{0,1} <	
	$_{\mathcal{A}}^{B}\mathrm{B}_{2,0}$	$_{\mathcal{A}}\mathbf{B}_{3,1}$	$_{\mathcal{A}}^{B}$ $\mathbf{B}_{0,2}$	$_{\mathcal{A}}\mathbf{B}_{1,3}$	
≪:	A _{1,3} <	$A_{1,0}$	A _{1,1} <	$A_{1,2}$	
	$_{\mathcal{A}}^{B}\mathrm{B}_{3,0}$	$_{A}^{\wedge}$ $B_{0,1}$	$_{\mathcal{A}}^{B}\mathrm{B}_{1,2}$	$_{A}$ B _{2,3}	
≪:	A _{2,0} <	$A_{2,1} < 1$	A _{2,2} <	A _{2,3} <-	***
	$_{\mathcal{A}}\mathbf{B}_{0,0}$	$_{A}^{\uparrow}$ B _{1,1}	$_{\mathcal{A}}^{B}$ B _{2,2}	$_{A}$ B _{3,3}	
≪:-	A _{3,1} <	A _{3,2} <	A _{3,3} <	A _{3,0} <	
	$_{1}$ $B_{1,0}$	$_{1}^{\prime}$ $B_{2,1}$	$_{1}^{\prime}$ B _{3,2}	$_{1}B_{0,3}$	
	1	:	1	1	

A _{0,3}	${ m A}_{0,0} \ { m B}_{0,1}$	A _{0,1}	A _{0,2}
B _{3,0}		B _{1,2}	B _{2,3}
A _{1,0}	$A_{1,1} B_{1,1}$	A _{1,2}	A _{1,3}
B _{0,0}		B _{2,2}	B _{3,3}
A _{2,1}	A _{2,2}	A _{2,3}	A _{2,0}
B _{1,0}	B _{2,1}	B _{3,2}	B _{0,3}
A _{3,2} B _{2,0}	A _{3,3} B _{3,1}	A _{3,0} B _{0,2}	A _{3,1} B _{1,3}

(e) Submatrix locations after second shift (f) Submatrix locations after third shift



Using blocking communications

```
/* Get into the main computation loop */
for (i=0; i<dims[0]; i++) {
  MatrixMultiply(nlocal, a, b, c); /*c=c+a*b*/
  /* Shift matrix a left by one */
  MPI Sendrecv replace(a, nlocal*nlocal, MPI DOUBLE,
      leftrank, 1, rightrank, 1, comm 2d, &status);
  /* Shift matrix b up by one */
  MPI Sendrecv replace(b, nlocal*nlocal, MPI DOUBLE,
      uprank, 1, downrank, 1, comm 2d, &status);
```

Code snippet of Cannon's matrix-matrix multiplication



Using non-blocking communications

```
/* Get into the main computation loop */
for (i=0; i<dims[0]; i++) {
  MPI Isend(a buffers[i%2], nlocal*nlocal, MPI DOUBLE,
      leftrank, 1, comm 2d, &reqs[0]);
  MPI Isend(b buffers[i%2], nlocal*nlocal, MPI DOUBLE,
      uprank, 1, comm 2d, &reqs[1]);
  MPI Irecv(a buffers[(i+1)%2], nlocal*nlocal, MPI DOUBLE,
      rightrank, 1, comm 2d, &reqs[2]);
  MPI Irecv(b buffers[(i+1)%2], nlocal*nlocal, MPI DOUBLE,
      downrank, 1, comm 2d, &reqs[3]);
  /* c = c + a*b */
  MatrixMultiply(nlocal, a buffers[i%2], b buffers[i%2], c);
  for (j=0; j<4; j++)
    MPI Wait(&reqs[j], &status);
```

Code snippet of Cannon's matrix-matrix multiplication

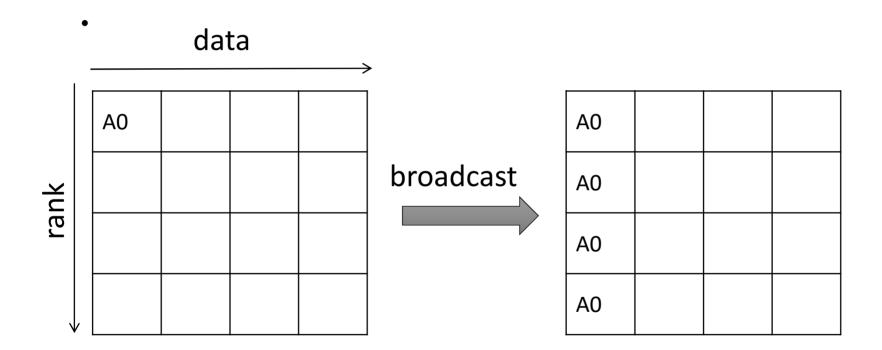


Collective Communication

- MPI provides an extensive set of functions of collective communication operations
- Operations are defined over a group corresponding to the communicator
- All processors in a communicator must call these operations
- Simple collective communication: barrier int MPI_Barrier(MPI_Comm comm)
 - Waits until all processes arrive

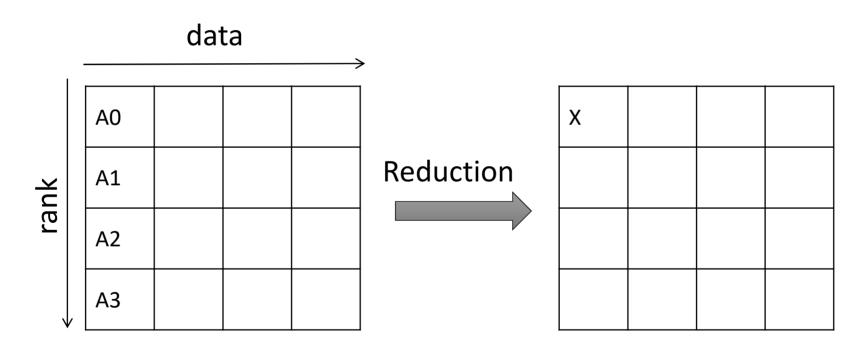


One-to-All Broadcast





All-to-One Reduction



X = A0 op A1 op A2 op A3



Predefined Reduction Operations

Operation	Meaning	Datatypes
MPI_MAX	Maximum	C integers and floating point
MPI_MIN	Minimum	C integers and floating point
MPI_SUM	Sum	C integers and floating point
MPI_PROD	Product	C integers and floating point
MPI_LAND	Logical AND	C integers
MPI_BAND	Bit-wise AND	C integers and byte
MPI_LOR	Logical OR	C integers
MPI_BOR	Bit-wise OR	C integers and byte
MPI_LXOR	Logical XOR	C integers
MPI_BXOR	Bit-wise XOR	C integers and byte
MPI_MAXLOC	max value-location	Data-pairs
MPI_MINLOC	min value-location	Data-pairs



Collective Communications (1)

- All processes must call the same collective function.
 - Ex. MPI_Recv() in P0 + MPI_Reduce() in P1 (X)
- All processes must send to the same target process
 - Ex.

Process 0

Process 1



Collective Communications (2)

- Recv buffer argument is only used in the destination process
 - But, other processes should provide the argument, even if it is NULL
 - Ex.

Process 0

Process 1

```
MPI_Reduce(in_buf,
NULL,
1,
MPI_CHAR,
MPI_SUM,
0,
MPI_COMM_WORLD);
```



Collective vs. Point-to-Point Comm.

- Point-to-point communication
 - MPI_Send/Recv are matched on the basis of tags and ranks
- Collective communication
 - Do NOT use tags
 - They're matched solely on the basis of receiver's rank and order



MPI_MAXLOC and MPI_MINLOC

- MPI_MAXLOC
 - Combines pairs of values (v_i, l_i)
 - Returns the pair (v, l)
 - -v is the maximum among all v_i 's
 - -l is the corresponding l_i
 - (if there are more than one, it is the smallest among all these l_i 's).
- MPI_MINLOC does the same, except for minimum value of v_i .

```
    Value
    15
    17
    11
    12
    17
    11

    Process
    0
    1
    2
    3
    4
    5
```

```
MinLoc(Value, Process) = (11, 2)
MaxLoc(Value, Process) = (17, 1)
```



MPI_MAXLOC and MPI_MINLOC

MPI datatypes for data-pairs used with the MPI_MAXLOC and MPI_MINLOC reduction operations.

MPI Datatype	C Datatype
MPI_2INT	pair of ints
MPI_SHORT_INT	short and int
MPI_LONG_INT	long and int
MPI_LONG_DOUBLE_INT	long double and int
MPI_FLOAT_INT	float and int
MPI_DOUBLE_INT	double and int

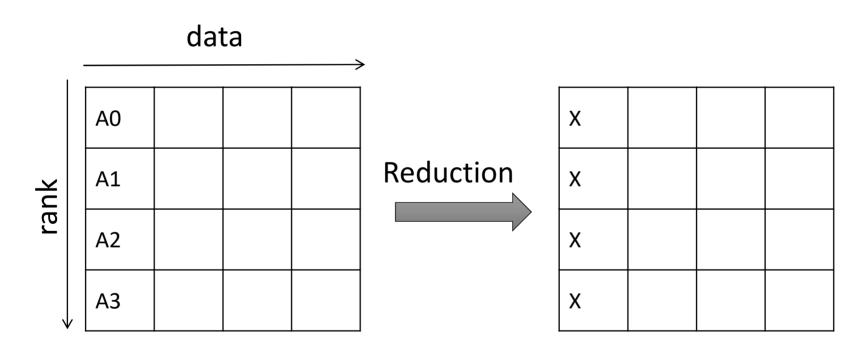


MPI_MAXLOC and MPI_MINLOC

```
double ain[30], aout[30];
int ind[30];
struct {
    double val;
    int rank;
} in[30], out[30];
for (i=0; i<30; ++i) {
    in[i].val = ain[i];
    in[i].rank = myrank;
MPI_Reduce( in, out, 30, MPI_DOUBLE_INT, MPI_MAXLOC, root, comm );
if (myrank == root) {
    for (i=0; i<30; ++i) {
        aout[i] = out[i].val;
        ind[i] = out[i].rank;
```

All-to-All Reduction

int MPI_Allreduce(void *sendbuf, void *recvbuf,
 int count, MPI_Datatype datatype, MPI_Op op,
 MPI_Comm comm)



X = A0 op A1 op A2 op A3



Prefix Operation (Inclusive)

X0		
X1		
X2		
Х3		

X0 = A0,

X1 = A0 op A1,

X2 = A0 op A1 op A2,

X3 = A0 op A1 op A2 op A3



Prefix Operation (Exclusive)

Scan	(op)

X0		
X1		
X2		
Х3		

$$X0 = 0$$

$$X1 = A0$$

$$X2 = A0 \text{ op } A1$$

$$X3 = A0 \text{ op } A1 \text{ op } A2$$



Scatter and Gather

Gather data at one process

Scatter data from source to all processes

data

ı						
		A0	A1	A2	А3	
rank						
٢						
	,					

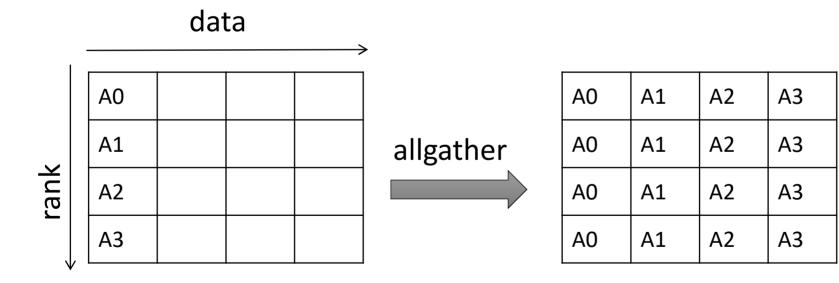
scatter
gather
/

A0		
A1		
A2		
A3		



Allgather

Gather and scatter them to all processes

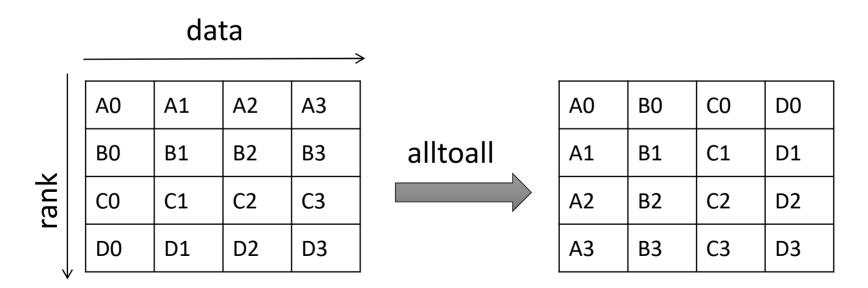




All-to-All

The all-to-all personalized communication

Analogous to a matrix transpose





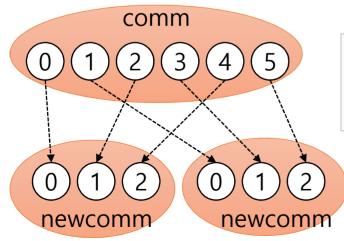
Communicators

- All MPI communication is based on a communicator which contains a context and a group
 - Contexts define a safe communication space for messagepassing – viewed as system-managed tags
 - Contexts allow different libraries to co-exist
 - Group is just a set of processes
 - Processes are always referred to by the unique rank in a group
- Pre-defined communicators
 - MPI_COMM_WORLD
 - MPI_COMM_NULL // initial value, cannot be used as comm
 - MPI_COMM_SELF // contains only the local calling process



Communicator Manipulation

- Duplicate communicator
 - MPI Comm dup(comm, newcomm)
 - Create a new context with similar structure
- Partition the group into disjoint subgroups
 - MPI_Comm_split(comm, color, key, newcomm)
 - Each sub-comm. contains the processes with the same color
 - The rank in the sub-communicator is defined by the key



```
color = (rank % 2 == 0)? 0 : 1;
key = rank / 2;
MPI_Comm_split(comm, color, key, &newcomm);
```



Communicator Manipulation – con't

- Obtain an existing group and free a group
 - MPI_Comm_group(comm, group) create a group having processes in the specified communicator
 - MPI_Group_free(group) free a group
- New group can be created by specifying members
 - MPI_Group_incl(), MPI_Group_excl()
 - MPI_Group_range_incl(), MPI_Group_range_excl()
 - MPI_Group_union(), MPI_Group_intersect()
 - MPI_Group_compare(), MPI_Group_translate_ranks()
- Subdivide a communicator
 - MPI Comm create(comm, group, newcomm)



Creating Cartesian Topologies

Creates cartesian topologies

- Creates a new communicator with dims dimensions.
 - ndims = number of dimensions
 - dims = vector of length of each dimension
 - periods = vector indicates which dims are periodic (wraparound link)
 - reorder = flag ranking may be reordered

P0	P1
P2	Р3

P0, P1, P2, P3

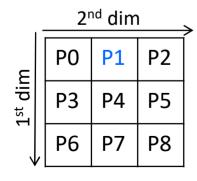


Using Cartesian Topologies

- Sending and receiving messages still require ranks (or process IDs)
- Convert ranks to cartesian coordinates and vice-versa

```
int MPI_Cart_coord(MPI_Comm comm_cart, int rank, int maxdims,
  int *coords)
int MPI_Cart_rank(MPI_Comm comm_cart, int *coords, int *rank)
```

The most common operation on cartesian topologies is a shift int MPI_Cart_shift(MPI_Comm comm_cart, int dir, int s_step, int *rank source, int *rank dest)



Examples:

P1 calls MPI_Cart_coord()
$$\rightarrow$$
 (0, 1)
MPI_Cart_rank({0, 1}) \rightarrow 1 (P1)
MPI_Cart_shift(0 (direction), 1 (step)) \rightarrow src:P0, dst:P2
(assuming all dims are periodic)



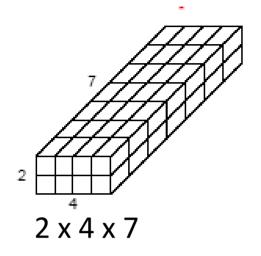
Splitting Cartesian Topologies

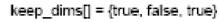
Partition a Cartesian topology to form lower-dimensional grids:

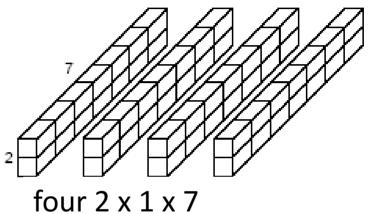
- keep_dims[i] determines whether to split or not the ith dimension
- The coordinate of a process in a sub-topology
 - Derived from its coordinate in the original topology
 - Disregarding the coordinates that correspond to the dimensions that were not retained
 - Example: $(2, 3) \rightarrow (2)$ if dims = {true, false}

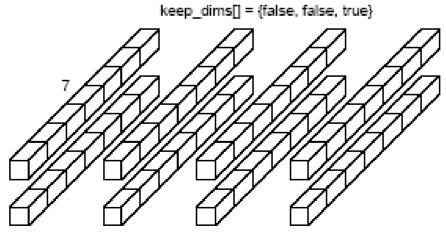


Splitting Cartesian Topologies









eight 1 x 1 x 7



Splitting Cartesian Topologies: Example

```
int nrow, mcol, i, lastrow, p, root;
int Iam, id2D, colID, ndim;
int coords1D[2], coords2D[2], dims[2], aij[1], alocal[3];
int belongs[2], periods[2], reorder;
MPI Comm comm2D, commcol;
/* Starts MPI processes ... */
MPI Init(&argc, &argv);
                                     /* starts MPI */
MPI Comm rank (MPI COMM WORLD, &Iam); /* get current process id */
MPI Comm size(MPI COMM WORLD, &p); /* get number of processes */
nrow = 3; mcol = 2; ndim = 2;
root = 0; periods[0] = 1; periods[1] = 0; reorder = 1;
/* create cartesian topology for processes */
                       /* number of rows */
dims[0] = nrow;
dims[1] = mcol;
                       /* number of columns */
MPI Cart create (MPI COMM WORLD, ndim, dims, periods, reorder, &comm2D);
MPI Comm rank(comm2D, &id2D);
MPI Cart coords(comm2D, id2D, ndim, coords2D);
/* Create 1D column subgrids */
                       /* this dimension belongs to subgrid */
belongs[0] = 1;
belongs[1] = 0;
MPI Cart sub(comm2D, belongs, &commcol);
MPI Comm rank(commcol, &colID);
MPI Cart coords(commcol, colID, 1, coords1D);
```

Limitations of MPI Data Types

- Only primitive data types can be exchanged through MPI_Send/Recv
- Many programs use more complex data structures
 - Ex. struct in C

MPI datatype	C datatype	
MPI_CHAR	signed char	
MPI_SHORT	signed short int	
MPI_INT	signed int	
MPI_LONG	signed long int	
MPI_LONG_LONG	signed long long int	
MPI_UNSIGNED_CHAR	unsigned char	
MPI_UNSIGNED_SHORT	unsigned short int	
MPI_UNSIGNED	unsigned int	
MPI_UNSIGNED_LONG	unsigned long int	
MPI_FLOAT	float	
MPI_DOUBLE	double	
MPI_LONG_DOUBLE	long double	
MPI_BYTE	11 75 11	
MPI_PACKED		

Basic data types in MPI



MPI Derived Data Types

- To make more complex data types to be exchanged through MPI communication methods
 - MPI should know the size of a data structure
 - MPI should know the members within the data structure
 - Location
 - Size of each member

```
struct a {
    MPI_DOUBLE x[2];
    MPI_DOUBLE y[2];
    MPI_LONG value[2];
};
```

Member	Offset in bytes
Х	0
У	16
value	32



MPI_Type create_struct

Builds a derived datatype consisting of individual elements

- array_of_block_lengths
 - Each member can be either a variable or an array
 - Ex. {2, 2, 2};
- array_of_displacements
 - Offsets of each member from start address
 - Ex. {0, 16, 32}
- array_of_types
 - Types of each member
 - Ex. {MPI_DOUBLE, MPI_DOUBLE, MPI_LONG}

```
struct a {
    MPI_DOUBLE x[2];
    MPI_DOUBLE y[2];
    MPI_LONG value[2];
};
```



MPI_Type create_struct

```
int main(int argc, char *argv[])
                                                                                  struct Partstruct
  struct Partstruct particle[1000];
                                                                                     char c;
  int i, j, myrank;
                                                                                     double d[6]:
  MPI Status status;
                                                                                     char b[7];
  MPI Datatype Particletype;
                                                                                  };
  MPI Datatype type[3] = { MPI CHAR, MPI DOUBLE, MPI CHAR };
  int blocklen[3] = \{1, 6, 7\};
  MPI Aint disp[3];
  MPI Init(&argc, &argv);
  disp[0] = &particle[0].c - &particle[0];
  disp[1] = &particle[0].d - &particle[0];
  disp[2] = &particle[0].b - &particle[0];
  MPI Type create struct(3, blocklen, disp, type, &Particletype);
  MPI Type commit(&Particletype);
  MPI Comm rank(MPI COMM WORLD, &myrank);
  if (myrank == 0)
    MPI Send(particle, 1000, Particletype, 1, 123, MPI COMM WORLD);
  else if (myrank == 1)
    MPI Recv(particle, 1000, Particletype, 0, 123, MPI COMM WORLD, &status);
  MPI Finalize();
```

- To know the address of the memory location referenced by location_p
- The address is stored in an integer variable of type MPI_Aint

```
int MPI_Get_address(
    void* location_p /* in */,
    MPI_Aint* address_p /* out */);
```

```
struct a {
    MPI_DOUBLE x[2];
    MPI_DOUBLE y[2];
    MPI_LONG value[2];
};
```

```
struct a a;
MPI_Get_address(&a.x, &x_addr);
MPI_Get_address(&a.y, &y_addr);
MPI_Get_address(&a.value, &value_addr);
array_of_displacements[0] = x_addr - &a;
array_of_displacements[1] = y_addr - &a;
array_of_dispalcements[2] = value_addr - &a;
```

Other methods

- MPI_Type_commit
 - To let MPI know the new data type
 - After calling this function, the new data type can be used in MPI communication methods

```
int MPI_Type_commit(MPI_Datatype* new_mpi_t_p /* in/out */);
```

- MPI_Type_free
 - When the new data type is no longer used, this function frees any additional storages used for the new data type

```
int MPI_Type_free(MPI_Datatype* old_mpi_t_p /* in/out */);
```



MPI_Pack/Unpack

- An alternative method to send/receive a complex data structure
- Pack multiple data types into a single buffer
- One pair of MPI_Send & MPI_Recv
- Sender and receiver have to know which data types are packed in the single buffer

```
buffer = malloc()
MPI_Pack(A)
MPI_Recv(buffer, MPI_PACKED)
MPI_Pack(B)
MPI_Unpack(A)
MPI_Unpack(B)
MPI_Unpack(B)
MPI_Unpack(B)
MPI_Unpack(C)

A B C
Pack
Sender
Buffer
Buffer
Receiver
```



Concluding Remarks

- MPI or the Message-Passing Interface
 - An interface of parallel programming in distributed memory system
 - Supports C, C++, and Fortran
 - Many MPI implementations
 - Ex, OpenMPI, MPICH2, Intel MPI
- SPMD program
- Message passing
 - Communicator
 - Point-to-point communication
 - Collective communication
 - Safe use of communication is important
 - Ex. MPI_Sendrecv()

