Problem Solving:

Sorting

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Q : Sum of pair

- Given an array A[] and a number x, check for pair in A[] with sum as x
- $A = \{1, 4, 45, 6, 10, -8\}$ and sum = 16

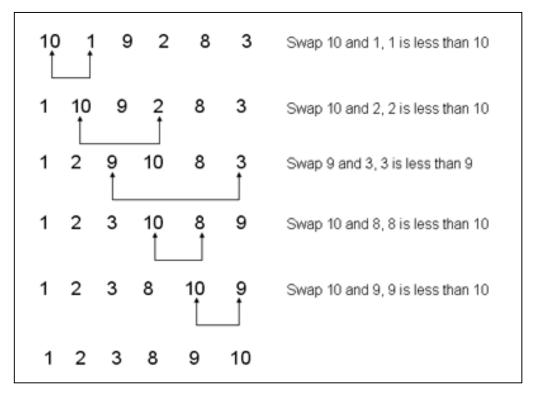
- Complexity of Exhaustive Search ?
- Better solution ?

Sorting

- The most basic technique in programming
- So many solutions for it
 - $O(n^2)$, $O(n \lg n)$, O(n)
 - depending on
 - · simplicity of mind
 - complexity of insert operation
- Algorithms
 - Selection, Insertion sort
 - Merge sort, Quick sort
 - Counting Sort
 - Bucket Sort

Selection Sort

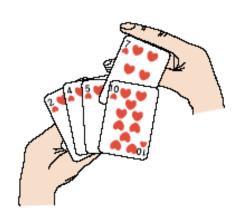
- Maintain two parts: sorted, unsorted
- Selection: in every iteration i, select the minimum (maximum) from the unsorted part, and move it to the sorted part (current index i)
- Two for-loops:
 - i = 0....n-2
 - J = i+1...n-1
 - · Keep the index of minimum
- Comparisons
 - (n-1) + (n-2) + ... + 1 = n(n-1)/2

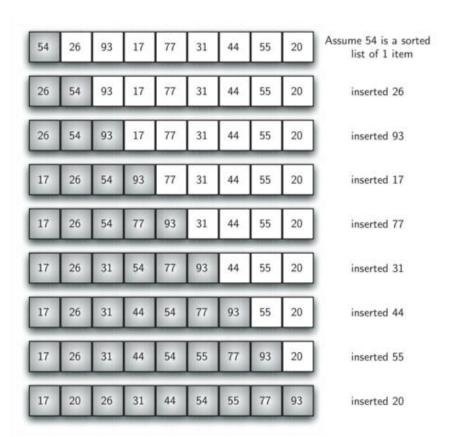


Selection Sort – the code

Insertion Sort

- Maintain two parts: sorted, unsorted
- Insertion: in every iteration i, pick i's element and insert it on the sorted part
- Two loops:
 - for-loop: i=1....n-1
 - while-loop: j=i-1, 0 and i's value < j's value



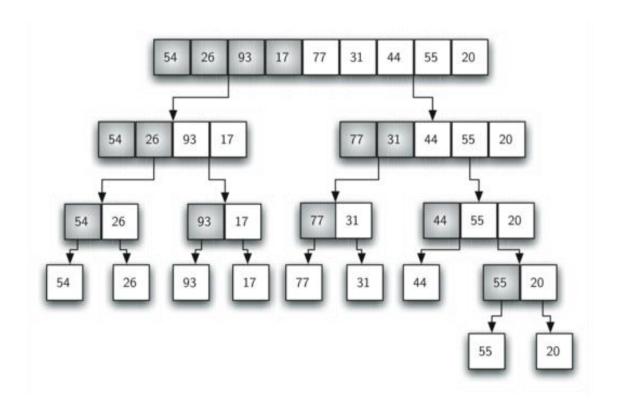


Insertion Sort – the code

```
void insertionSort(int arr[], int n) {
  int i, key, j;
  for (i = 1; i < n; i++) {
    key = arr[i];
    j = i - 1;
    while (j \ge 0 \&\& arr[j] > key) {
      arr[j + 1] = arr[j];
      \dot{j} = \dot{j} - 1;
    arr[j + 1] = key;
```

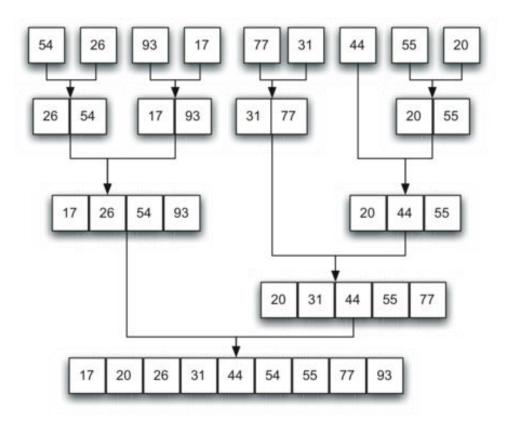
Merge sort – Divide & Conquer

• Divide



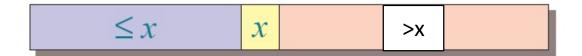
Merge sort – Divide & Conquer

Conquer (merge)



Quick Sort

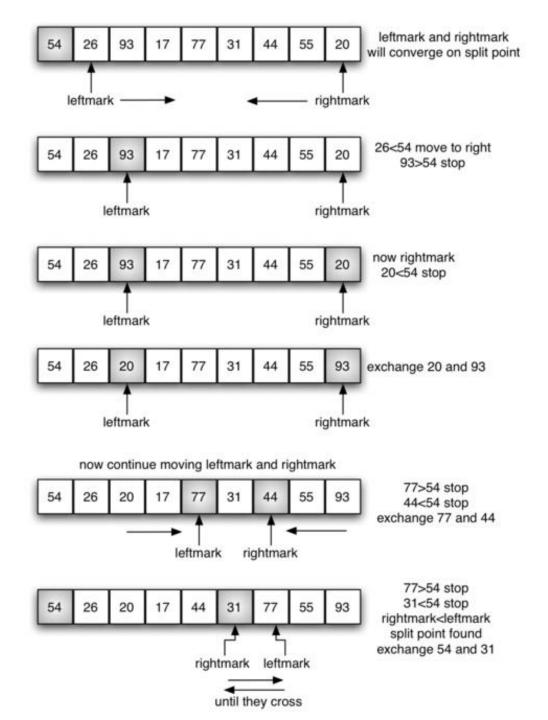
- Divide-and-Conquer
 - Divide the array into two parts
 - the value of x is called pivot



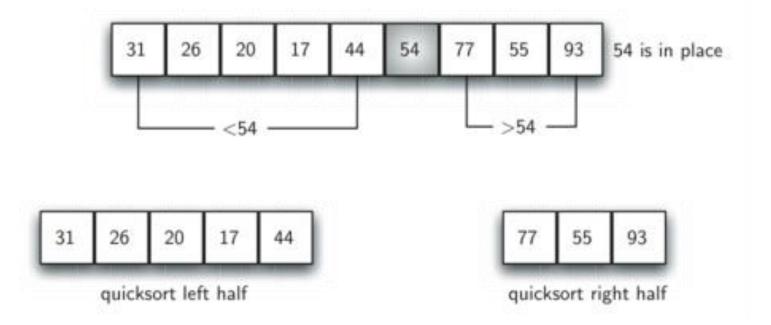
- Conquer
 - do the same to each divided subarray
 - Combine

partition

• pivot: 54



partition



```
void quicksort(int a[], int l, int h)
{
    int p;
    if((h-1)>0) {
        p = partition(a, l, h);
        quicksort(a, l, p-1);
        quicksort(a, p+1, h);
    }
}
int partit
```

```
int partition(int a[], int first, int last) {
    int pivot, left, right;
    pivot = first;
    left = first;
    right = last;
    while (left < right) {</pre>
       while(a[left] <= a[pivot] && left < last)</pre>
         left++;
       while(a[right] > a[pivot] && right > first)
         right--;
       if(left < right) {</pre>
         Swap(&a[left], &a[right]);
    Swap(&a[pivot], &a[right]);
    return right;
}
```

Library

```
#include <stdlib.h>
void qsort(void *base, size_t nel, size_t width,
      int (*compare) (const void *, const void *));

    array: base

    number of elements: nel

    size of each element: width bytes

    you should provide compare function

int mycompare(int *i, int *j)
{
        if (*i > *j) return (1);
        elseif (*i < *j) return (-1);
        else return (0);
qsort((char *) a, cnt, sizeof(int), mycompare);
```

Vito's Family

The input consists of several test cases. The first line contains the number of test cases.

For each test case you will be given the integer number of relatives r (0 < r < 500) and the street numbers (also integers) $s_1, s_2, \ldots, s_i, \ldots, s_r$ where they live (0 < s_i < 30,000). Note that several relatives might live at the same street number.

For each test case, your program must write the minimal sum of distances from the optimal Vito's house to each one of his relatives. The distance between two street numbers s_i and s_j is $d_{ij} = |s_i - s_j|$.

```
Sample Input
```

2

2 2 4

3 2 4 6

Sample Output

2

4