



Birthday attacks

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Birthday attacks

- This is the best “generic” collision attack on a hash function $H: \{0,1\}^* \rightarrow \{0,1\}^n$?
- Compute $H(x_1), \dots, H(x_{2^{n/2}})$
 - What is the probability of a collision?
- Related to the so-called *birthday paradox*
 - How many people are needed to have a 50% chance that some two people share a birthday?

“Birthday attacks” analysis (1)

- How many possibilities that are all different?
 - $(K)_N = K(K-1)\dots(K-N+1)$: the number of all possible samples without replacement when we choose N different samples
- Probability of no repetition?

$$\frac{k * (k - 1) * (k - 2) * \dots * (k - n + 1)}{k * k * k \dots * k} =$$

$$\frac{k}{k} * \frac{k-1}{k} * \dots * \frac{k-n+1}{k} = 1 * \left(1 - \frac{1}{k}\right) * \left(1 - \frac{2}{k}\right) * \dots * \left(1 - \frac{n-1}{k}\right) \leq$$

$1 - \frac{a}{k} \leq e^{-\frac{a}{k}}$

$$e^{-\frac{1}{k}} * e^{-\frac{2}{k}} * e^{-\frac{3}{k}} * \dots * e^{-\frac{n-1}{k}} \cong e^{-\frac{n^2}{2k}}$$

“Birthday attacks” analysis (2)

$$e^{-\frac{n^2}{2k}} \leq \frac{1}{2} \Leftrightarrow \frac{n^2}{2k} \geq \ln 2 \Leftrightarrow$$

$$n^2 \geq 2 (\ln 2)k = 1.38k \Leftrightarrow$$

$$n \geq \sqrt{1.38k}$$

- Bottom line: For k=365, n=23 suffices
- In general $n = \Omega(\sqrt{k})$ suffices

Theorem for birthday attacks

- Thm. When the number of balls (in balls and N bins) is $O(N^{1/2})$, the probability of a collision is 50%. (see the CLR book)
- Need $2n$ -bit output length to get security against attackers running in 2^n time
 - Twice the length of block cipher keys
 - A block cipher with 128-bit key provides equivalent security to a hash function with 256-bit output.

Questions?

