



Database Systems

Lecture04 – Introduction to SQL



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Nested Subqueries

- SQL provides a mechanism for the nesting of subqueries.
- A **subquery** is a **select-from-where** expression that is nested within another query.
- A common use of subqueries is to perform tests for set membership, set comparisons, and set cardinality.



Example Query

- Find courses offered in Fall 2009 and in Spring 2010

```
select distinct course_id
from section
where semester = 'Fall' and year= 2009 and
       course_id in (select course_id
                       from section
                       where semester = 'Spring' and year= 2010);
```

- Find courses offered in Fall 2009 but not in Spring 2010

```
select distinct course_id
from section
where semester = 'Fall' and year= 2009 and
       course_id not in (select course_id
                       from section
                       where semester = 'Spring' and year= 2010);
```



Example Query

- Find the total number of (distinct) students who have taken course sections taught by the instructor with *ID* 10101

```
select count (distinct ID)  
from takes  
where (course_id, sec_id, semester, year) in  
        (select course_id, sec_id, semester, year  
          from teaches  
          where teaches.ID= 10101);
```

- Note:** Above query can be written in a much simpler manner. The formulation above is simply to illustrate SQL features.



Set Comparison

- Find names of instructors with salary greater than that of some (at least one) instructor in the Biology department.

```
select distinct T.name  
from instructor as T, instructor as S  
where T.salary > S.salary and S.dept_name = 'Biology';
```

- Same query using > **some** clause

```
select name  
from instructor  
where salary > some (select salary  
                        from instructor  
                        where dept_name = 'Biology');
```

Definition of Some Clause

- $F <\text{comp}> \text{some } r \Leftrightarrow \exists t \in r \text{ such that } (F <\text{comp}> t)$
Where $<\text{comp}>$ can be: $<, \leq, >, =, \neq$

$(5 < \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{true}$ (read: 5 < some tuple in the relation)

$(5 < \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{false}$

$(5 = \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true}$

$(5 \neq \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true (since } 0 \neq 5)$

$(= \text{some}) \equiv \text{in}$

However, $(\neq \text{some}) \not\equiv \text{not in}$



Example Query

- Find the names of all instructors whose salary is greater than the salary of all instructors in the Biology department.

```
select name
from instructor
where salary > all (select salary
                        from instructor
                        where dept_name = 'Biology');
```

Definition of all Clause

- $F \text{ <comp> all } r \Leftrightarrow \forall t \in r (F \text{ <comp> } t)$

$$(5 \text{ < all } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{false}$$

$$(5 \text{ < all } \begin{array}{|c|} \hline 6 \\ \hline 10 \\ \hline \end{array}) = \text{true}$$

$$(5 = \text{all } \begin{array}{|c|} \hline 4 \\ \hline 5 \\ \hline \end{array}) = \text{false}$$

$$(5 \neq \text{all } \begin{array}{|c|} \hline 4 \\ \hline 6 \\ \hline \end{array}) = \text{true (since } 5 \neq 4 \text{ and } 5 \neq 6)$$

$(\neq \text{ all}) \equiv \text{not in}$

However, $(= \text{ all}) \not\equiv \text{in}$



Test for Empty Relations

- The **exists** construct returns the value **true** if the argument subquery is nonempty.
- **exists** $r \Leftrightarrow r \neq \emptyset$
- **not exists** $r \Leftrightarrow r = \emptyset$



Correlation Variables

- Yet another way of specifying the query “Find all courses taught in both the Fall 2009 semester and in the Spring 2010 semester”

```
select course_id
from section as S
where semester = 'Fall' and year= 2009 and
      exists (select *
              from section as T
              where semester = 'Spring' and year= 2010
                  and S.course_id= T.course_id);
```

- **Correlated subquery**
- **Correlation name** or **correlation variable**



Not Exists

- Find all students who have taken **all** courses offered in the Biology department.

```
select distinct S.ID, S.name  
from student as S  
where not exists ( (select course_id  
                    from course  
                    where dept_name = 'Biology')  
except  
                (select T.course_id  
                 from takes as T  
                 where S.ID = T.ID));
```

- Note that $X - Y = \emptyset \Leftrightarrow X \subseteq Y$
- Note: Cannot write this query using = **all** and its variants



Subqueries in the From Clause

- SQL allows a subquery expression to be used in the **from** clause
- Find the average instructors' salaries of those departments where the average salary is greater than \$42,000.

```
select dept_name, avg_salary
from (select dept_name, avg (salary) as avg_salary
      from instructor
      group by dept_name)
where avg_salary > 42000;
```

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Another way to write above query

```
select dept_name, avg_salary
from (select dept_name, avg (salary)
      from instructor
      group by dept_name)
as dept_avg (dept_name, avg_salary)
where avg_salary > 42000;
```



With Clause

- The **with** clause provides a way of defining a temporary **view** whose definition is available only to the query in which the **with** clause occurs.
- Find all departments with the maximum budget

```
with max_budget (value) as  
    (select max(budget)  
     from department)  
select budget  
from department, max_budget  
where department.budget = max_budget.value;
```



Complex Queries using With Clause

- With clause is very useful for writing complex queries
- Supported by most database systems, with minor syntax variations
- Find all departments where the total salary is greater than the average of the total salary at all departments

```
with dept_total (dept_name, value) as
    (select dept_name, sum(salary)
     from instructor
     group by dept_name),
dept_total_avg(value) as
    (select avg(value)
     from dept_total)
select dept_name
from dept_total, dept_total_avg
where dept_total.value >= dept_total_avg.value;
```



Scalar Subquery

- Scalar subquery is one which is used where a single value is expected
- E.g.

```
select dept_name,  
       (select count(*)  
        from instructor  
        where department.dept_name = instructor.dept_name)  
       as num_instructors  
from department;
```
- E.g.

```
select name  
from instructor  
where salary * 10 >  
      (select budget from department  
       where department.dept_name = instructor.dept_name)
```
- Runtime error if subquery returns more than one result tuple



Modification of the Database

- Deletion of tuples from a given relation
- Insertion of new tuples into a given relation
- Updating values in some tuples in a given relation



Modification of the Database – Deletion

- Delete all instructors

delete from *instructor*

- Delete all instructors from the Finance department

delete from *instructor*
where *dept_name* = 'Finance';

- Delete all tuples in the *instructor* relation for those instructors associated with a department located in the Watson building.

```
delete from instructor  
where dept_name in (select dept_name  
                        from department  
                        where building = 'Watson');
```



Deletion (Cont.)

- Delete all instructors whose salary is less than the average salary of instructors

```
delete from instructor  
where salary < (select avg (salary) from instructor);
```

- Problem: as we delete tuples from deposit, the average salary changes
- Solution used in SQL:
 1. First, compute **avg** salary and find all tuples to delete
 2. Next, delete all tuples found above (without recomputing **avg** or retesting the tuples)



Modification of the Database – Insertion

- Add a new tuple to *course*

```
insert into course  
values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

- or equivalently

```
insert into course (course_id, title, dept_name, credits)  
values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

- Add a new tuple to *student* with *tot_creds* set to null

```
insert into student  
values ('3003', 'Green', 'Finance', null);
```



Insertion (Cont.)

- Add all instructors to the *student* relation with *tot_creds* set to 0

```
insert into student
    select ID, name, dept_name, 0
from instructor
```

- The **select from where** statement is evaluated fully before any of its results are inserted into the relation
 - otherwise queries like
insert into table1 select * from table1
would cause problems. (*table1* did not have any primary key defined)



Modification of the Database – Updates

- Increase salaries of instructors whose salary is over \$100,000 by 3%, and all others receive a 5% raise

- Write two **update** statements:

```
update instructor  
  set salary = salary * 1.03  
  where salary > 100000;
```

```
update instructor  
  set salary = salary * 1.05  
  where salary <= 100000;
```

- What if we change the order of the two statements?
 - The order is important
- Can be done better using the **case** statement (next slide)



Case Statement for Conditional Updates

- Same query as before but with case statement

```
update instructor  
  set salary = case  
    when salary <= 100000 then salary * 1.05  
    else salary * 1.03  
  end
```

Updates with Scalar Subqueries

- Recompute and update `tot_creds` value for all students

```
update student S  
  set tot_cred = ( select sum(credits)  
                    from takes natural join course  
                    where S.ID= takes.ID and  
                        takes.grade <> 'F' and  
                        takes.grade is not null);
```

- Sets `tot_creds` to null for students who have not taken any course
- Instead of `sum(credits)`, use:

```
  case  
  when sum(credits) is not null then sum(credits)  
  else 0  
end
```