# Problem Solving:

## Combinatorics

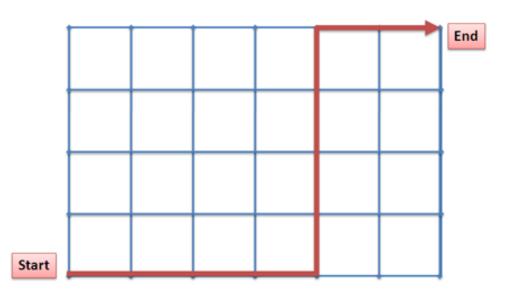
April 2019

Honguk Woo

## Q : Paths of grid

- The number of paths across a grid of (n rows X m columns)
  - e.g. start to end on (4 X 6) grid

### **How Many Paths?**



### Counting Techniques

- Combinatorics is a branch of mathematics concerning the study of finite or countable discrete structures, simply about "Counting"
- Suppose you have 5 shirts and 4 pants in your closet:
- Product rule
  - Different way to wear = Combinations : 5 shirts \* 4 pants = total 20
  - |A| X |B|
- Sum rule
  - if any one is missing from the closet, it is one of 9 clothing pieces
  - |A| + |B|
- Inclusion-Exclusion Formula
  - $|A \cup B| = |A| + |B| |A \cap B|$
  - $|A \cup B \cup C| = |A| + |B| + |C| |A \cap B| |B \cap C| |B \cap C| + |A \cap B \cap B|$
  - · eliminates double counting

## More counting techniques

#### (1) Permutation

- an arrangement of *n* items where every item appears exactly once  $n! = \prod_{i=1}^{n} i$
- arrange three letters word using a, b, c
  - abc, acb, bac, bca, cab, cba: 6 words (3!)
  - Very common for exhaustive search methods
  - What about 10! = 3,628,800 .... huge in complexity, approaching to the limits of exhaustive search
- what if letters are reused ((2) permutation with repetition)?
  - aaa, abb, ... on a, b, c: 3X3X3 = 27
  - e.g., r-length **strings** among n characters  $_{n}\prod_{r} = \mathbf{n}^{r}$

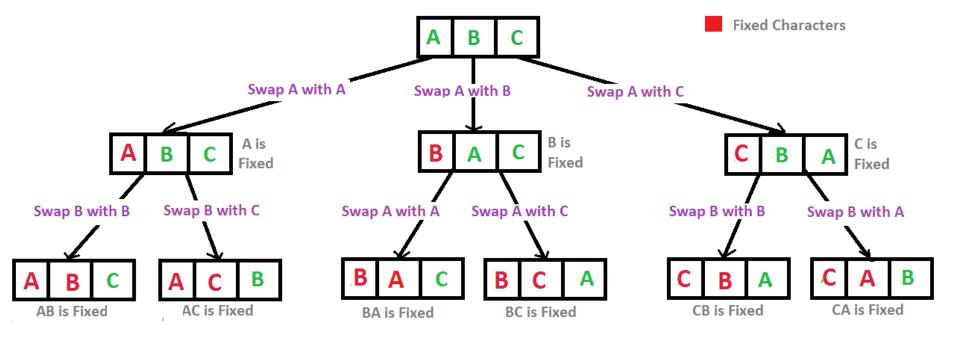
#### • (3) Subsets

- there are  $2^n$  subsets of *n* items
  - a, b, c, ab, bc, ac, abc, and an empty set: 8 subsets

## Print all permutations of a given string

- Given a string "abc", print all the permutations (rearrangements)
- "abc", 3! = 6
- "abcd", 4! ...
- Many loops ?
- Recursion
  - Base case: each string (the string length == the target length)
  - General case: "a" + permute("bc"), "b" + permute("ac") ...

If (base case) then print string else generate more permutation



#### **Recursion Tree for Permutations of String "ABC"**

```
void permute(char *a, int I, int r) {
  if (I == r) printf("%s\n", a);
  else {
    for (i = I; i <= r; i++) {
       swap((a+I), (a+i));
       permute(a, I+1, r);
       swap((a+I), (a+i));
    }
  }
}</pre>
```

char str[] = "abc";
Permute(str, 0, strlen(str) -1);

#### Binomial Coefficient

The number of ways to choose k things out of n

$$_{n} C_{k} \equiv \binom{n}{k} \equiv \frac{n!}{(n-k)! \ k!}$$

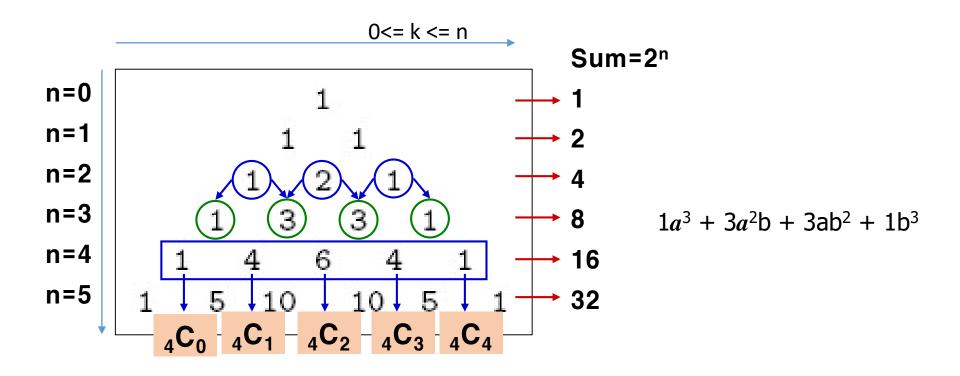
- e.g. how many ways to form a k-member committee from n people
- e.g., how many paths from a grid of k X n-k

## Binomial Coefficient (2)

- Coefficient of (a+b)<sup>n</sup>
  - $(a + b)^3 = 1a^3 + 3a^2b + 3ab^2 + 1b^3$
  - $a^2b \rightarrow choose 2 a (or 1 b) from all 3, (a+b) X (a+b) X (a+b)$

- what is the coefficient of  $a^kb^{n-k}$  ?
  - how many ways to choose k a-terms out of n
  - $(a + b)^3$  = aaa + 3 aab + 3 abb + bbb

## Pascal's Triangle



$$_{n}C_{k} = _{(n-1)}C_{(k-1)} + _{(n-1)}C_{k}$$

## Binomial Coefficient (3)

Consider: choose k from n items, and two cases about the below n



case1: if this (n) is contained, same as choosing k-1 from n-1 case2: if this (n) is not contained,

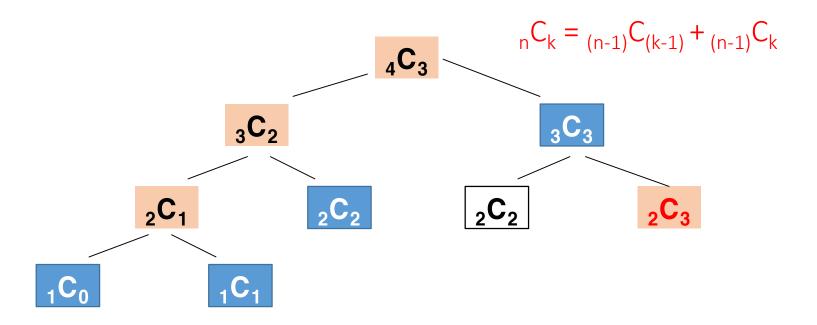
must choose k from n-1

$$_{n}C_{k} = {}_{(n-1)}C_{(k-1)} + {}_{(n-1)}C_{k}$$

Note that (n-k)!k! may cause overflow; more stable using recurrence relation.

$$\begin{bmatrix} n \\ k \end{bmatrix} = \frac{n!}{(n-k)!k!}$$

### Stop the recurrence



good ending points (base cases)

$$\cdot_{n}C_{0} = 1$$
  
 $\cdot_{n}C_{n} = 1$ 

### Pascal's Triangle in C

#### recursion

```
int pascal(int r, int c){
 if(c == 0 \mid\mid c == r) return 1;
 else
  return pascal(r-1, c-1) + pascal(r-1, c);
int main(){
 int n = 7;
 for(int i=0; i<n; i++) {
  for(int j=0; j< i+1; j++){
    printf("%d ", pascal(i, j));
  printf("\n");
 return 0;
```

```
1
11
121
1331
14641
15101051
1615201561
```

#### Pascal's Triangle in C

#### memorization

```
#define MAXN 100 /* largest n or m */
long binomial_coefficient(n,m) /* computer n choose m */
  int i,j; /* counters */
  long bc[MAXN][MAXN]; /* table of binomial coefficients */
  for (i=0; i \le n; i++) bc[i][0] = 1;
  for (j=0; j<=n; j++) bc[j][j] = 1;
   for (i=1; i<=n; i++)
      for (j=1; j<i; j++)
          bc[i][j] = bc[i-1][j-1] + bc[i-1][j];
   return( bc[n][m]);
```

#### Fibonacci Numbers

```
F_0 = 0
F_1 = 1
F_n = F_{n-1} + F_{n-2} \text{ for } n \ge 2
\text{int fb(int n) } \{ \text{ if } (n == 0 \mid\mid n == 1) \text{ return n; } // \text{ f}(0) = 0, \text{ f}(1) = 1 \text{ else return fb(n-1) + fb(n-2); } \}
```

#### Closed form solution

$$F_n = \frac{1}{\sqrt{5}} \left( \left( \frac{1 + \sqrt{5}}{2} \right)^n - \left( \frac{1 - \sqrt{5}}{2} \right)^n \right)$$

#### **Eulerian Numbers**

- the number of permutations of length n with exactly k ascending sequences
  - E.g. Permutations of the numbers 0 to 9 (length 10) in which exactly 3 elements are greater than the previous element
    - 0, ..., 9, k=3
    - · 3 1 5 7 6 4 8 2 0
  - case 1: add 9 on solution of 0..8 k=3 w/o changing k
  - case 2: add 9 on solution of 0..8 k=2 increasing k
  - case 3: add 9 on solution of 0..8 k=1 impossible

$${n \choose k} = k {n-1 \choose k} + (n-k+1) {n-1 \choose k-1}$$

#### Recurrence Relations

- Recursive relation makes it easy to count "recursively defined structures"
  - Recursively defined structures: tree, list, ...,
    - divide & conquer algorithms : binary search, quick sort, merge sort, ...
- Recurrence relation
  - An equation defined in terms of itself

$$a_{n} = a_{n-1} + 1$$
 ,  $a_{1} = 1$   $\rightarrow$   $a_{n} = n$   $a_{n} = 2a_{n-1}$  ,  $a_{1} = 2$   $\rightarrow$   $a_{n} = 2^{n}$   $a_{n} = na_{n-1}$  ,  $a_{1} = 1$   $\rightarrow$   $a_{n} = n!$ 

### Mathematical Induction

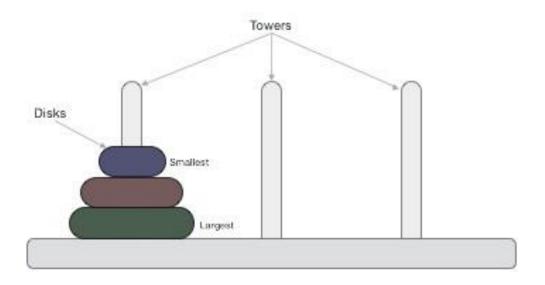
$$T_n = 2T_{n-1} + 1, T_0 = 0$$

n	0	1	2	3	4	5	6	7	
T <sub>n</sub>	0	1	3	7	15	31	63	127	

Prove that  $T_n = 2^n - 1$ 

#### Q: Tower of Hanoi

- The mission is to move all the disks to some another tower without violating the sequence of arrangement.
  - Only one disk can be moved among the towers at any given time
  - Only the "top" disk can be removed
  - No large disk can sit over a small disk



#### Tower of Hanoi

- First, we move the smaller (top) disk to aux
- Then, we move the larger (bottom) disk to destination
- And finally, we move the smaller disk from aux to destination

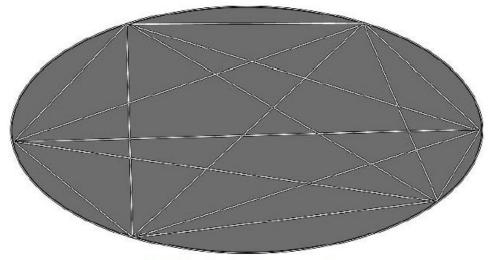
- In general
- Step 1 Move n-1 disks from source to aux
- Step 2 Move n<sup>th</sup> disk from source to dest
- Step 3 Move n-1 disks from aux to dest

```
#include <stdio.h>
void towerOfHanoi(int n, char from rod, char to rod, char aux rod) {
         if (n == 1)
                   printf("\n Move disk 1 from rod %c to rod %c", from rod, to rod);
                   return;
         towerOfHanoi(n-1, from rod, aux rod, to rod);
          printf("\n Move disk %d from rod %c to rod %c", n, from_rod, to_rod);
         towerOfHanoi(n-1, aux_rod, to_rod, from_rod);
int main()
         int n = 4; // Number of disks
         towerOfHanoi(n, 'A', 'C', 'B'); // A, B and C are names of rods
         return 0;
```

### Q: How Many Pieces of Land

- You are given a land and you are asked to choose n arbitrary points on its boundary. Then you connect each point with every other point using straight lines, forming n(n-1)/2 connections.
- What is the maximum number of pieces of land you will get by choosing the points on the boundary carefully?

N (input)	Maximum number of pieces (output)
1	1
2	2
3	4
4	8



Dividing the land when n = 6.