



Number Theory

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Notation

From here on:

- N denotes a positive integer.
- p denotes a prime.

Notation: $\mathbb{Z}_N = \{0, 1, 2, \dots, N - 1\}$

We can do addition and multiplication modulo N

Modular arithmetic

Examples: let $N = 12$

$$9 + 8 = 5 \quad \text{in } \mathbb{Z}_{12}$$

$$5 \times 7 = 11 \quad \text{in } \mathbb{Z}_{12}$$

$$5 - 7 = 10 \quad \text{in } \mathbb{Z}_{12}$$

Arithmetic in \mathbb{Z}_N works as you expect,
e.g., $x \cdot (y+z) = x \cdot y + x \cdot z$ in \mathbb{Z}_N

distributive law

Greatest common divisor

Def: For integers x, y , $\text{gcd}(x, y)$ is the greatest common divisor of x, y

Example: $\text{gcd}(12, 18) = 6$

Fact: For all integers x, y , there exist integers a, b such that

$$a \cdot x + b \cdot y = \text{gcd}(x, y) \quad 2 \cdot 12 + -1 \cdot 18 = 6$$

a, b can be found efficiently using the **extended Euclid alg.**

If $\text{gcd}(x, y) = 1$ we say that x and y are relatively prime

Modular inversion

Over the rational numbers, the inverse of 2 is $\frac{1}{2}$.

What about \mathbb{Z}_N ?

Def: The **inverse** of x in \mathbb{Z}_N is an element y in \mathbb{Z}_N s.t. $x \cdot y = 1$ in \mathbb{Z}_N
 y is denoted x^{-1} .

Example: Let N be an odd integer. The inverse of 2 in \mathbb{Z}_N is $(N+1)/2$

$$2 \cdot (N+1)/2 = N+1 = 1 \text{ in } \mathbb{Z}_N$$

Modular inversion

Which elements have an inverse in \mathbb{Z}_N ?

Lemma: x in \mathbb{Z}_N has an inverse if and only if $\gcd(x, N) = 1$

Proof:

$$\gcd(x, N) = 1 \Rightarrow \exists a, b: a \cdot x + b \cdot N = 1 \Rightarrow a \cdot x = 1 \text{ in } \mathbb{Z}_N$$

$$\Rightarrow x^{-1} = a \text{ in } \mathbb{Z}_N$$

$$\gcd(x, N) > 1 \Rightarrow \forall a: \gcd(a \cdot x, N) > 1 \Rightarrow a \cdot x \neq 1 \text{ in } \mathbb{Z}_N$$

More notation

Def: \mathbb{Z}_N^* = (set of invertible elements in \mathbb{Z}_N) =
= $\{ x \in \mathbb{Z}_N : \gcd(x, N) = 1 \}$

Examples:

1. for prime p , $\mathbb{Z}_p^* = \mathbb{Z}_p \setminus \{0\} = \{1, 2, \dots, p-1\}$
2. $\mathbb{Z}_{12}^* = \{1, 5, 7, 11\}$

For x in \mathbb{Z}_N^* , we can find x^{-1} using the extended Euclidean algorithm.

Euclid's rule for GCD

- How do you find the greatest common divisor of two integers?
 - Select the largest factor that can divide both
 - A brute force approach requires exponential time
- Ancient Greece mathematician Euclid discovered the rule:
 - If x and y are positive integers with $x \geq y$ then
$$\gcd(x, y) = \gcd(y, x \bmod y)$$
 - When $x \bmod y = 0$, then y is the gcd
- We can swap parameters each time to keep largest as 1st parameter

Euclid's algorithm

- This rule leads to the following algorithm

Function Euclid (a, b)

Input: Two integers a and b with $a \geq b \geq 0$ (n -bit integers)

Output: $\gcd(a, b)$

if $b=0$: return a

else: return Euclid($b, a \bmod b$)

Example of Euclid's algorithm

- As an example of the running of Euclid, consider the computation of $\text{gcd}(30, 21)$:

$$\begin{aligned}\text{Euclid}(30, 21) &= \text{Euclid}(21, 9) \\ &= \text{Euclid}(9, 3) \\ &= \text{Euclid}(3, 0) \\ &= 3\end{aligned}$$

C code

```
#include <stdio.h>
```

```
int gcd_algorithm(int x, int y)
{
    if (y == 0)
        return x;

    return gcd_algorithm(y, (x % y));
}
```

```
int main(void)
```

```
{
    int num1, num2, gcd;
    printf("\nEnter two numbers: ");
    scanf("%d%d", &num1, &num2);

    if (num1 < 0) num1 = -num1;
    if (num2 < 0) num2 = -num2;
    gcd = gcd_algorithm(num1, num2);

    if (gcd)
        printf("\nThe GCD of %d and %d is %d\n", num1, num2, gcd);
    else
        printf("\nInvalid input!!!\n");
    return 0;
}
```

Extended Euclid's Algorithm

function extended-Euclid (a, b)

Input: Two positive integers a and b with $a \geq b \geq 0$ (n -bits)

Output: Integers x, y, d such that $d = \gcd(a, b)$
and $ax + by = d$

if $b = 0$: return $(1, 0, a)$

$(x', y', d) = \text{extended-Euclid}(b, a \bmod b)$

return $(y', x' - \text{floor}(a/b)y', d)$

C code

```
int gcdExtended(int a, int b, int *x, int *y)
{
    // Base Case
    if (b == 0)
    { *x = 1;
      *y = 0;
      return a;
    }
    int x1, y1; // To store results of recursive call
    int gcd = gcdExtended(b, a % b, &x1, &y1);

    // Update x and y using results of recursive
    // call
    *x = y1;
    *y = x1 - (a/b) * y1;

    return gcd;
}
```

```
int main()
{
    int x, y;
    int a = 30, b = 21;
    int g = gcdExtended(a, b, &x, &y);
    printf("gcd(%d, %d) = %d, x = %d, y = %d",
           a, b, g, x, y);
    return 0;
}
```

```
gcd(30, 21) = 3, x = 3, y = -2
```

```
...Program finished with exit code 0
Press ENTER to exit console. █
```

Multiplicative inverses

- Two numbers a and b are relatively prime if $\gcd(a,b) = 1$
- If a and N are relatively prime, then we know the multiplicative inverse exists (e.g. 4 mod 7)
 - If a and N are relatively prime, then $ax + Ny = 1$
- Extended-Euclid(a,N) can be used:
 - Returns integers x, y, d such that $d = \gcd(a, b)$ and $ax + by = d$
 - First, it must return $d=1$ as the gcd to confirm that a and N are relatively prime
- $Ny = 0 \pmod{N}$ for all integers y
- Thus, $ax \equiv 1 \pmod{N}$ // ax is congruent to 1 modulo N
- Then x is the multiplicative inverse of a modulo N

Multiplicative inverse of 20 Mod 79

function extended-Euclid (a , b)

if $b = 0$: return (1 , 0 , a)

$(x', y', d) = \text{extended-Euclid}(b, a \bmod b)$

return (y' , $x' - \text{floor}(a/b)y'$, d)

a	b	x'	y'	d	ret 1	ret 2	ret 3
79	20						

Multiplicative inverse of 20 Mod 79

function extended-Euclid (a , b)

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79	20						
20	19						

Multiplicative inverse of 20 Mod 79

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79	20						
20	19						
19	1						

Multiplicative inverse of 20 Mod 79

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$(x', y', d) = \text{extended-Euclid}(b, a \bmod b)$

return (y' , $x' - \text{floor}(a/b)y'$, d)

a	b	x'	y'	d	ret 1	ret 2	ret 3
79	20						
20	19						
19	1						
1	0						

Multiplicative inverse of 20 Mod 79

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$(x', y', d) = \text{extended-Euclid}(b, a \bmod b)$

return (y' , $x' - \text{floor}(a/b)y'$, d)

a	b	x'	y'	d	ret 1	ret 2	ret 3
79	20						
20	19						
19	1						
1	0				1	0	1

Multiplicative inverse of 20 Mod 79

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a	b	x'	y'	d	ret 1	ret 2	ret 3
79	20						
20	19						
19	1	1	0	1			
1	0				1	0	1

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a	b	x'	y'	d	ret 1	ret 2	ret 3
79	20						
20	19	0	1	1			
19	1	1	0	1	0	1	1
1	0				1	0	1

Multiplicative inverse of 20 Mod 79

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return (y' , $x' - \text{floor}(a/b)y'$, d)

a	b	x'	y'	d	ret 1	ret 2	ret 3
79	20						
20	19	0	1	1	1	-1	1
19	1	1	0	1	0	1	1
1	0				1	0	1

Multiplicative inverse of 20 Mod 79

function extended-Euclid (a , b)

if $b = 0$: return (1 , 0 , a)

$(x', y', d) = \text{extended-Euclid}(b, a \bmod b)$

return (y' , $x' - \text{floor}(a/b)y'$, d)

a	b	x'	y'	d	ret 1	ret 2	ret 3
79	20	1	-1	1			
20	19	0	1	1	1	-1	1
19	1	1	0	1	0	1	1
1	0				1	0	1

Multiplicative inverse of 20 Mod 79

function extended-Euclid (a , b)

if $b = 0$: return (1 , 0 , a)

$(x', y', d) = \text{extended-Euclid}(b, a \bmod b)$

return (y' , $x' - \text{floor}(a/b)y'$, d)

a	b	x'	y'	d	ret 1	ret 2	ret 3
79	20	1	-1	1	-1	4	1
20	19	0	1	1	1	-1	1
19	1	1	0	1	0	1	1
1	0				1	0	1

Multiplicative inverse of 20 Mod 79

function extended-Euclid (a , b)

if $b = 0$: return (1 , 0 , a)

$(x', y', d) = \text{extended-Euclid}(b, a \bmod b)$

return (y' , $x' - \text{floor}(a/b)y'$, d)

a	b	x'	y'	d	ret 1	ret 2	ret 3
79	20	1	-1	1	-1	4	1
20	19	0	1	1	1	-1	1
19	1	1	0	1	0	1	1
1	0				1	0	1

$$ax + Ny = 1 = 20(4) + 79(-1)$$

$$\text{Thus } x = a^{-1} = 4$$

How can we choose a random prime?

- A number p is *prime* if
 - p is an integer and $p > 1$
 - The only (positive) factors of p are 1 and p .
- If $n > 1$ is not prime, it is *composite*
 - n has a positive factor other than 1 or n .

Q. Can you develop a practical algorithm?

Brute-force algorithm

- Input: integer $N > 1$
- Output: (x,y) such that $x > 1$, $y > 1$ and $xy = N$
or “Prime” if N is prime
- For $x = 2$ to $N-1$
 - If x evenly divides N
 - Return($x, N/x$)
- Return(“Prime”)

Slightly Better Algorithm

- Input: integer $N > 1$
- Output: (x,y) such that $x > 1$, $y > 1$ and $xy = N$
or “Prime” if N is prime

- For $x = 2$ to ~~$N-1$~~ \sqrt{N}
 - If x evenly divides N
 - Return($x, N/x$)
- Return(“Prime”)

If N has 200 digits

\sqrt{N} has 100 digits

If inner loop takes .1 ns

Algorithm take 3×10^{82} years

*Time proportional to the value
of N , not the number of digits*

PRIMES is in P

PRIMES is in P

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A. We can develop an efficient algorithm in $O(n^{12})$.

Best: $O(n^6)$.

Practical (randomized) algorithm

Suppose we want to generate a large random prime

say, prime p of length 1024 bits (i.e., $p \approx 2^{1024}$)

Step 1: Choose a random integer $p \in [2^{1024}, 2^{1025}-1]$

Step 2: Test if $2^{p-1} = 1$ in \mathbb{Z}_p^*

If so, output p and stop. If not, goto step 1 .

Simple algorithm (not the best)

Fermat's theorem (1640)

Thm: Let p be a prime

$$\forall x \in \mathbb{Z}_p^* : \quad x^{p-1} = 1 \text{ in } \mathbb{Z}_p$$

Example: $p=5$. $3^4 = 81 = 1$ in \mathbb{Z}_5

So: $x \in \mathbb{Z}_p^* \Rightarrow x \cdot x^{p-2} = 1 \Rightarrow x^{-1} = x^{p-2}$ in \mathbb{Z}_p

another way to compute inverses, but less efficient than Euclidian algorithm

Questions?

