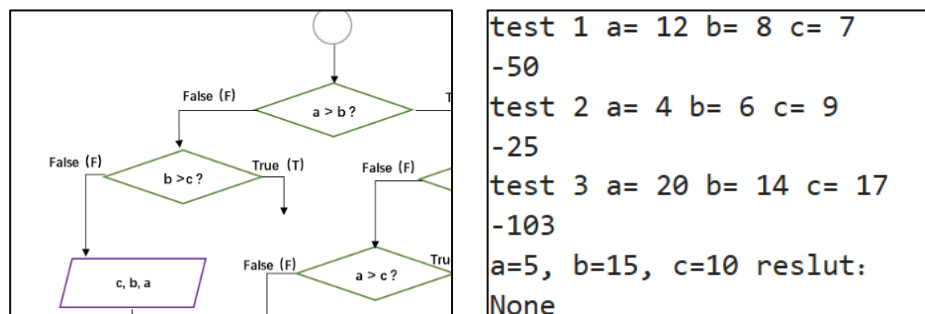


## 1. Flowchart

[10 points] Write a function `Print_values` with arguments `a`, `b`, and `c` to reflect the following flowchart. Here the purple parallelogram operator on a list `[x, y, z]` is to compute and print  $x+y-10z$ . Try your output with some random `a`, `b`, and `c` values. Report your output when `a = 5`, `b = 15`, `c = 10`.

**Result:**

当 `a = 5`, `b = 15`, `c = 10`, 不执行运算，输出：none，运行结果如图：



## 2. Continuous celing function

[10 points] Given a list with `N` positive integers. For every element `x` of the list, find the value of continuous ceiling function defined as  $F(x) = F(\text{ceil}(x/3)) + 2x$ , where  $F(1) = 1$ .

**Result:**

首先看到 `ceil` 考虑使用 `math` 模块，观察函数为递归，定义边界情况后实现即可（最简单思路，若提高效率可考虑使用哈希表等记忆函数值），运行结果如图：

1	1
2	5
3	7
4	13
5	15
6	17
7	21
8	23
9	25
10	33

## 3. Dice rolling

3.1 [15 points] Given 10 dice each with 6 faces, numbered from 1 to 6. Write a function `Find_number_of_ways` to find the number of ways to get sum `x`, defined as the sum of values

on each face when all the dice are thrown.

Result:

来自于 Leetcode 1155 题 (<https://leetcode.com/problems/number-of-dice-rolls-with-target-sum/description/>), 是原题的简版, 即设定了  $n=10$ 、 $k=6$ , 最合适算法无疑是 dynamic programming



这里尝试更容易理解的递归, 即从后往前推算最后一个骰子的值为 1-6 时, 前 9 个骰子应该为多少, 递归解决剩余问题。基准情况: 如果没有骰子( $n=0$ )且目标为 0, 说明成功找到一种方法, 返回 1, 如果没有骰子但目标不为 0, 或者目标变成负数, 说明这种方法不可行, 返回 0。(递归思路: 1. <https://www.deepseek.com/zh>)

```
定义递归函数: ways(n, target) 表示用 n 个骰子掷出总和 target 的方法数

基准情况:
1. 如果 n == 0 且 target == 0: 返回 1 (成功)
2. 如果 n == 0 且 target != 0: 返回 0 (失败)
3. 如果 target < 0: 返回 0 (不可能)

递归关系:
最后一个骰子可能掷出 1 到 6:
ways(n, target) = ways(n-1, target-1) // 最后一个骰子为1
                  + ways(n-1, target-2) // 最后一个骰子为2
                  + ways(n-1, target-3) // 最后一个骰子为3
                  + ways(n-1, target-4) // 最后一个骰子为4
                  + ways(n-1, target-5) // 最后一个骰子为5
                  + ways(n-1, target-6) // 最后一个骰子为6
```

传统递归会出现重复计算的问题, 大量情况被重复计算浪费大量时间, 考虑使用记忆化, 所有值只计算一次。(记忆存储思路: B 站 up 主 flyingchow123)

Python 中的 functools.lru\_cache 装饰器, functools.lru\_cache 是 Python 标准库中 functools 模块的一部分。lru\_cache 可以用来为一个函数添加一个缓存系统。它可以帮助我们优化递归函数, 避免重复计算已经计算过的值。(@lru\_cache 装饰器理解: <https://zhuanlan.zhihu.com/p/640954732>), 运行结果如图:

```
总和10的方法数: 1
总和15的方法数: 2002
总和30的方法数: 2930455
总和45的方法数: 831204
总和60的方法数: 1
```

**3.2 [5 points]** Count the number of ways for any x from 10 to 60, assign the number of ways to a list called Number\_of\_ways, so which x yields the maximum of Number\_of\_ways?

**Result:**

```
3.2 计算所有可能总和的方式数:
方式数列表长度: 51
产生最大方式数的x值: 35
最大方式数: 4395456
```

#### 4. Dynamic programming

**4.1 [5 points]** Write a function Random\_integer to fill an array of N elements by randomly selecting integers from 0 to 10.

**Result:**

使用 random 模块的 random.randint 生成随机整数，使用 append 添加到列表中。

```
import random

# 4.1
def Random_integer(N):
    result = []
    for i in range(N):
        num = random.randint(0, 10)
        result.append(num)
    return result

arr = Random_integer(5)
print(arr)

[4, 5, 9, 5, 10]
```

**4.2 [15 points]** Write a function Sum\_averages to compute the sum of the average of all subsets of the array. For example, given an array of [1, 2, 3], you Sum\_averages function should compute the sum of: average of [1], average of [2], average of [3], average of [1, 2], average of [1, 3], average of [2, 3], and average of [1, 2, 3].

**Result:**

动态规划思路描述：（来源：1.<https://chat.deepseek.com/> 2.）

(1)状态定义：

$dp\_sum[k]$ : 存储所有长度为  $k$  的子集的和的总和

$dp\_count[k]$ : 存储所有长度为  $k$  的子集的个数

(2)初始化:

$dp\_count[0] = 1$  表示空集, 这是动态规划的起点

(3)核心动态规划过程:

遍历数组中的每个数字, 对于每个数字, 从后往前更新状态 (避免重复计算)

$dp\_sum[k] = dp\_sum[k] + dp\_sum[k-1] + num * dp\_count[k-1]$ :

$dp\_sum[k]$ : 保持已有的长度为  $k$  的子集和

$dp\_sum[k-1]$ : 之前长度为  $k-1$  的子集和

$num * dp\_count[k-1]$ : 当前数字添加到所有  $k-1$  长度子集中产生的新和

$dp\_count[k] = dp\_count[k] + dp\_count[k-1]$ : 更新子集数量

(4)计算最终结果:

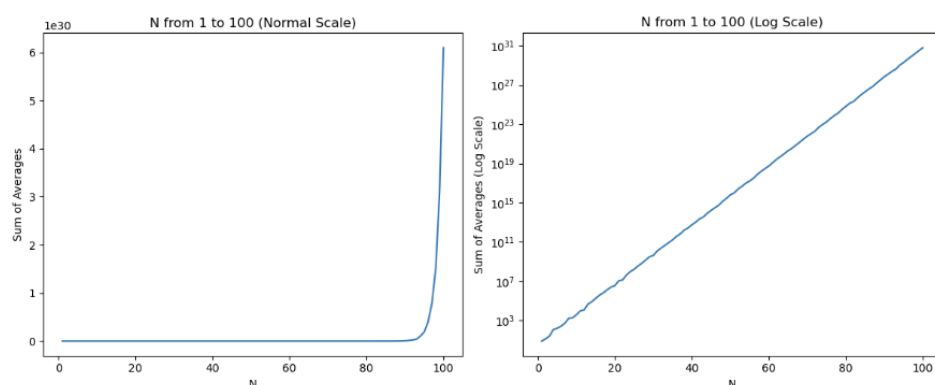
遍历所有可能的子集长度 (1 到  $n$ ) 对于每个长度  $k$ , 将  $dp\_sum[k] / k$  加到总和和  
这样就得到了所有子集平均值的总和。

数组  $[1, 2, 3]$  的结果: 14.0  
数组  $[1, 2, 3, 4, 5]$  的结果: 93.0

**4.3 [5 points]** Call `Sum_averages` with  $N$  increasing from 1 to 100, assign the output to a list called `Total_sum_averages`. Plot `Total_sum_averages`, describe what do you see.

**Result:**

从图中可以观察到, 随着数组长度  $N$  的增大, 所有子集平均值之和 (Sum of Averages) 呈现快速增长的指数趋势 (左图为正常图, 右图对  $y$  取了对数以便观察趋势), 子集数量的指数级增长主导了整个求和过程, 即使单个子集的平均值很小, 但数量庞大的子集使得总和呈现爆炸式增长。



## 5. Path counting

**5.1 [5 points]** Create a matrix with N rows and M columns, fill the right-bottom corner and top-left corner cells with 1, and randomly fill the rest of matrix with integer 0 or 1.

**Result:**

使用 `matrix = [[0 for _ in range(M)] for _ in range(N)]` 的写法来代替传统写法:

`[0 for _ in range(M)]`: 创建一个长度为 M 的列表, 每个元素都是 0

`for _ in range(M)`: 循环 M 次, `_` 是约定表示不关心循环变量值的变量名

0: 每次循环都放入数字 0

`for _ in range(N)`: 外层循环 N 次

每次循环都创建一个新的长度为 M 的零列表

测试结果如下:

```
# 测试
matrix = create_matrix(3, 4)
for row in matrix:
    print(row)

[1, 0, 1, 0]
[0, 0, 1, 1]
[0, 1, 0, 1]
```

**5.2 [25 points]** Consider a cell marked with 0 as a blockage or dead-end, and a cell marked with 1 is good to go. Write a function `Count_path` to count total number of paths to reach the right-bottom corner cell from the top-left corner cell.

**Notice:** for a given cell, you are **only allowed** to move either rightward or downward.

来自于 Leetcode 63 题(<https://leetcode.com/problems/unique-paths-ii/>), 改动了有无障碍物的状态 (0 代表障碍物), 也是经典的动态规划问题。

**63. Unique Paths II**

Medium Topics Companies Hint

You are given an  $n \times m$  integer array `grid`. There is a robot initially located at the **top-left corner** (i.e., `grid[0][0]`). The robot tries to move to the **bottom-right corner** (i.e., `grid[n - 1][m - 1]`). The robot can only move either down or right at any point in time.

An obstacle and space are marked as 1 or 0 respectively in `grid`. A path that the robot takes cannot include **any** square that is an obstacle.

Return the number of possible unique paths that the robot can take to reach the bottom-right corner.

The testcases are generated so that the answer will be less than or equal to  $2 \times 10^5$ .

**Example 1:**

**Input:** `obstacleGrid = [[0,0,0],[0,1,0],[0,0,0]]`

**Output:** 2

**Explanation:** There is one obstacle in the middle of the 3x3 grid above. There are two ways to reach the bottom-right corner:

1. Right -> Right -> Down -> Down
2. Down -> Down -> Right -> Right

(算法思路: 1.B 站 up 代码随想录 2. <https://chat.deepseek.com/>)

(1) 状态定义

创建  $dp[i][j]$  表示从起点  $(0,0)$  到达  $(i,j)$  的路径数量

(2) 初始化

起点  $(0,0)$ : 如果可通过, 路径数为 1; 如果是障碍物, 路径数为 0

(3) 状态转移

对于每个单元格  $(i,j)$ :

如果是障碍物:  $dp[i][j] = 0$  (无法到达)

如果不是障碍物: 路径数 = 从上面来的路径数 + 从左边来的路径数

$dp[i][j] = dp[i-1][j] + dp[i][j-1]$

(4) 边界处理

第一行  $(i=0)$ : 只能从左边来 (不能从上面来)

第一列  $(j=0)$ : 只能从上面来 (不能从左边来)

时间复杂度:  $O(N \times M)$ , 遍历整个网格一次 空间复杂度:  $O(N \times M)$ , 存储 DP 表格

测试结果如下:

测试矩阵: [1, 1, 0, 0, 1, 1] [1, 0, 1, 0, 0, 0] [1, 1, 1, 1, 0, 0] [1, 1, 1, 0, 1, 1] [0, 1, 1, 1, 1, 1] [1, 0, 1, 1, 0, 1] 路径数量: 5	测试矩阵: [1, 1, 0] [1, 1, 1] [1, 1, 1] 路径数量: 5
--	---

**5.3 [5 points]** Let  $N = 10, M = 8$ , run `Count_path` for 1000 times, each time the matrix (except the right-bottom corner and top-left corner cells, which remain being 1) is re-filled with integer 0 or 1 randomly, report the mean of total number of paths from the 1000 runs.

Result:

<pre>[57]: #5. Path counting #5.3 def run_simulation():     N = 10     M = 8     total_paths = 0     runs = 1000      for i in range(runs):         matrix = create_matrix(N, M)         paths = Count_path(matrix)         total_paths += paths      mean_paths = total_paths / runs     return mean_paths  mean_result = run_simulation() print(f"1000次运行的平均路径数: {mean_result}")  1000次运行的平均路径数: 0.338</pre>	<pre>[58]: #5. Path counting #5.3 def run_simulation():     N = 10     M = 8     total_paths = 0     runs = 1000      for i in range(runs):         matrix = create_matrix(N, M)         paths = Count_path(matrix)         total_paths += paths      mean_paths = total_paths / runs     return mean_paths  mean_result = run_simulation() print(f"1000次运行的平均路径数: {mean_result}")  1000次运行的平均路径数: 0.426</pre>
--	--