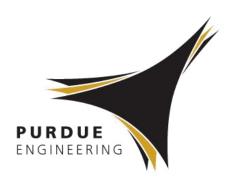
Al, app-defined OS, & the ossified Linux kernel

Felix Xiaozhu Lin 林小竹

http://xsel.rocks



Summary

Describe recent trends of OS research

Present three cases of our investigation

Share personal reflection on OS research

Self intro

- 2014 now. Asst. prof, Purdue ECE
- 2014 PhD in CS. Rice U (advisor: Lin Zhong)
 - Thesis: OS for mobile computing
- 2008 MS + BS. Tsinghua U

Crossroads Systems Exploration Lab



XSEL & Friends, 2017

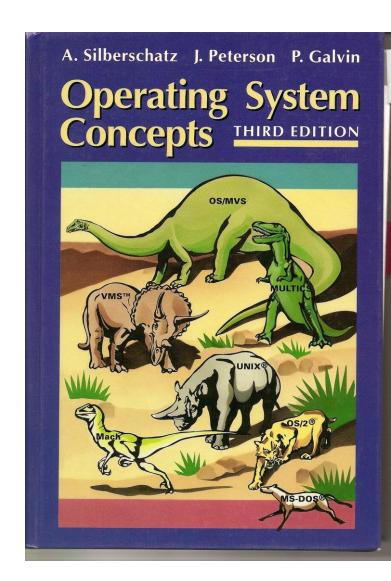
AI + big data: the workloads in our era

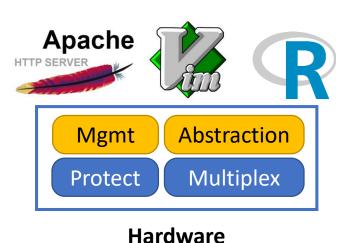


How does OS respond?

OS as in textbooks

- Resource management
- Abstraction
- Protection
- Multiplexing

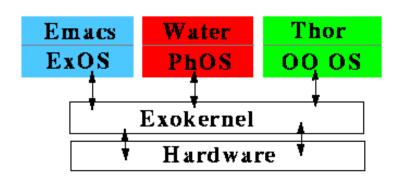




OS in 2000



Hardware
OS in 2018



Exokernel: An Operating System Architecture for Application-Level Resource Management

Dawson R. Engler, M. Frans Kaashoek, and James O'Toole Jr. M.I.T. Laboratory for Computer Science

SOSP 1995

Two premises of our research

- 1. Important apps define OSes
- 2. The Linux kernel is the new firmware

Like it or not ...

- Following the model: widely adopted
 - Various ML frameworks
 - Android
 - DPDK
 - RDMA
- Against the mode: less adopted
 - Library OS
 - Multikernel
 - Unikernels
 - (... and all kinds of research proposals)

This talk: our exploration in three cases

- App-defined OSes
 - Case 1: Memory mgmt
 - Case 2: Storage
- The ossified Linux Kernel
 - Case 3: Transkernel

Case 1: App-defined memory management

Exploiting hybrid memory for stream analytics

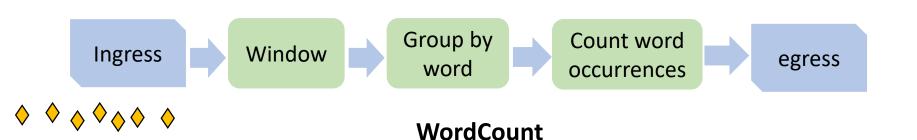
Background: Stream analytics in 1 minute



High input throughput: millions of events per sec Low delay processing: sub-seconds

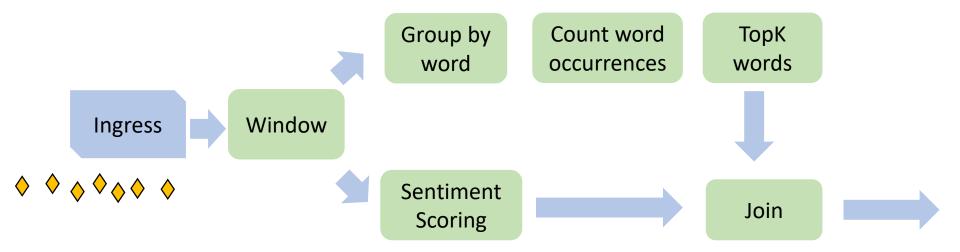
Background: streaming pipeline

Operators Computations consuming & producing streams **Pipeline** A graph of operators



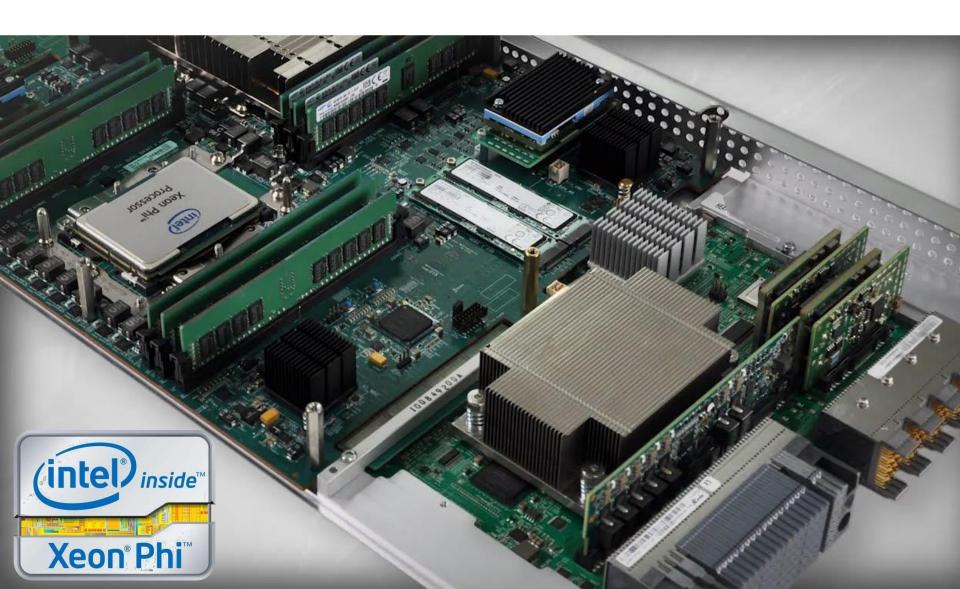
Background: streaming pipeline

Operators Computations consuming & producing streams **Pipeline** A graph of operators

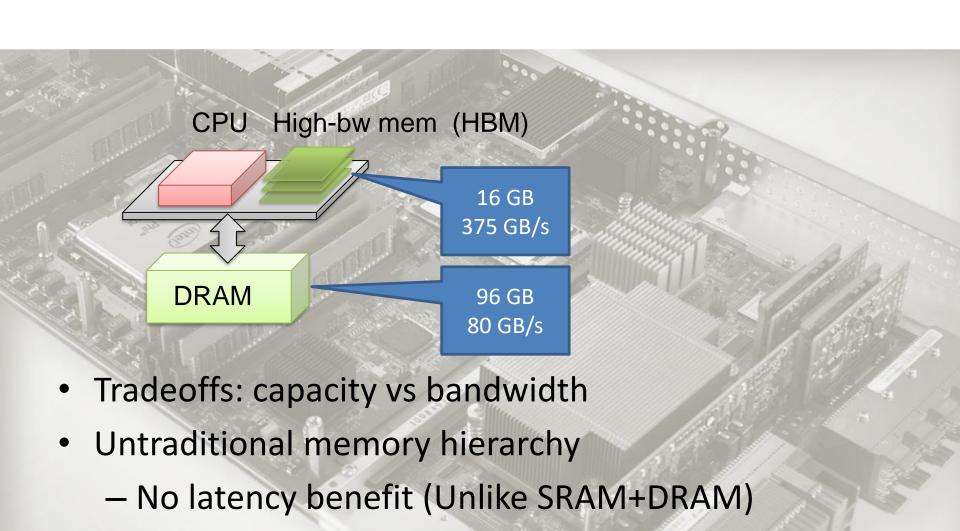


Sentiment Detection

Hybrid high-bandwidth memory



Hybrid high-bandwidth memory



- Configurable: sw-managed or hw-managed

Accelerate stream analytics with HBM?

The most expensive stream operators: grouping



Grouping traditionally built as Hash Mismatch HBM!



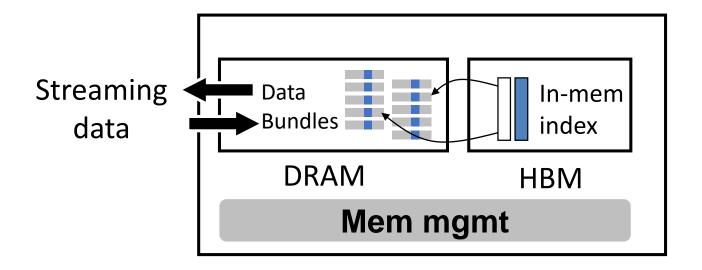
Grouping in Hash	High bandwidth memory
High capacity demand	Capacity limited
Extensive random access	Seq access + parallelism



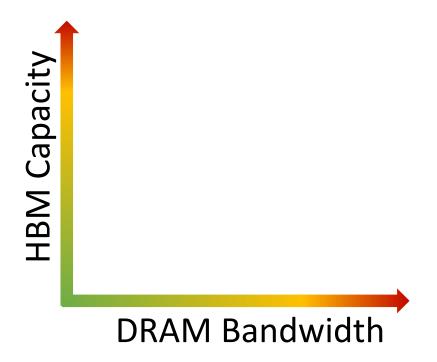
HBM can accelerate grouping!

- ... with an unconventional algorithm choice
- Grouping: hash → sort
- Sort is worse than Hash on algorithmic complexity
 - O(NlogN) vs. O(N)
- Yet, Sort beats Hash with ...
 - Abundant mem bw
 - High task parallelism
 - Wide SIMD (avx512)

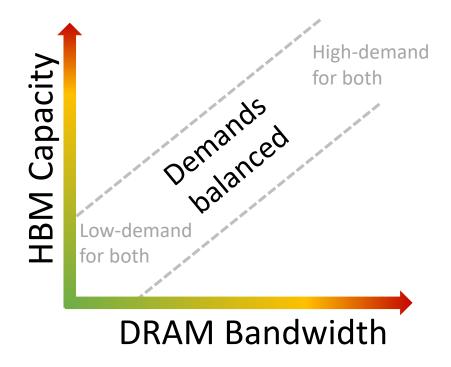
A bold design: only use HBM for in-mem index of parallel sorting



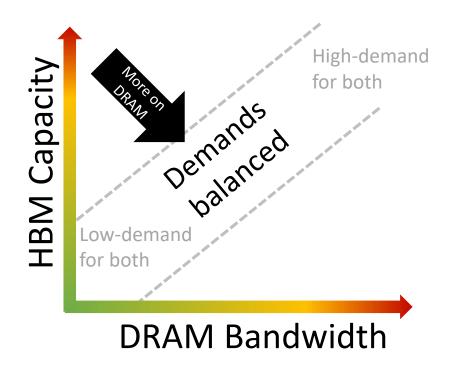
- Prevent either from becoming the bottleneck
- A single knob: where to allocate new index: HBM or DRAM?



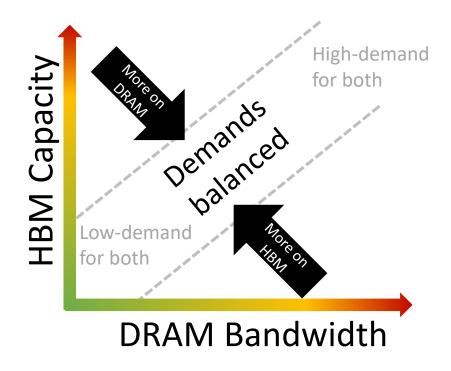
- Prevent either from becoming the bottleneck
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- Prevent either from becoming the bottleneck
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Case 1: First streaming engine optimized for hybrid memory

• Works on real hardware



Best stream analytics performance on a single machine

Generic memory allocators are inadequate!

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Case 2: App-defined storage

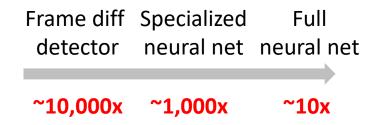
A data store for video analytics

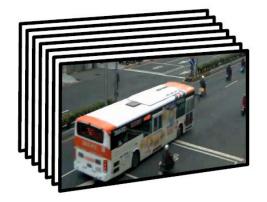




Retrospective video analytics in 30 seconds

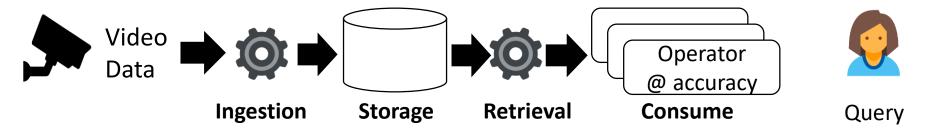
- A query: "find all white buses appeared yesterday"
- A cascade of operators





- Query picks operator accuracies
 - Lower accuracy → lower cost → faster execution

Need for a video store for retrospective analytics



"Aren't there many video databases already?" Yes.

But they are for **human** consumers. Not for **algorithmic** consumers.

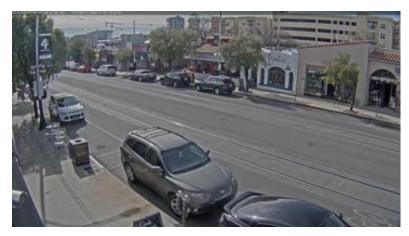
Central design issue: choose storage formats catering to operators

 Video format: resolution, frame rate, crop, compression level, color channel... Many knobs!

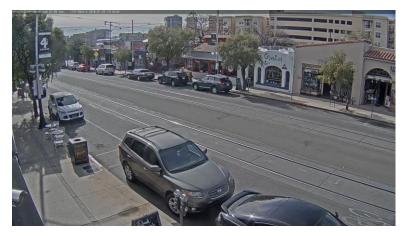
Central design issue: choose storage formats catering to operators

- Video format: resolution, frame rate, crop, compression level, color channel... Many knobs!
- Higher accuracy

 richer format & higher cost



200p 1FPS (~100KB/s)



720p 30FPS (~400KB/s)

Key problem: What video formats to store?

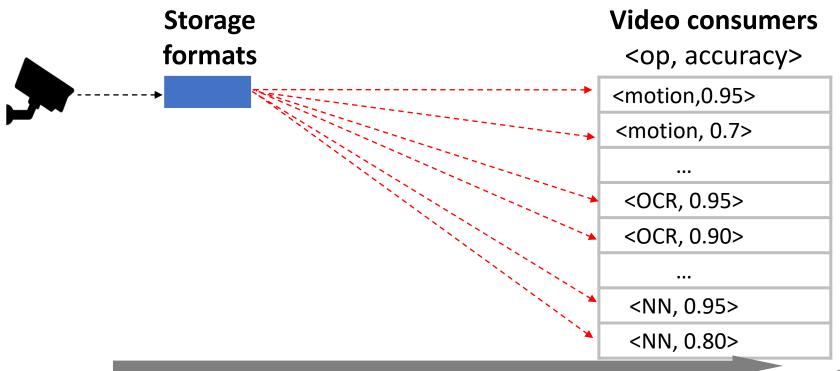
1. Retrieving formats must be fast enough to feed operators

2. Formats must be rich enough to satisfy all accuracy needs

Storing one unified format?

Retrieving & decoding slows down operators 😕

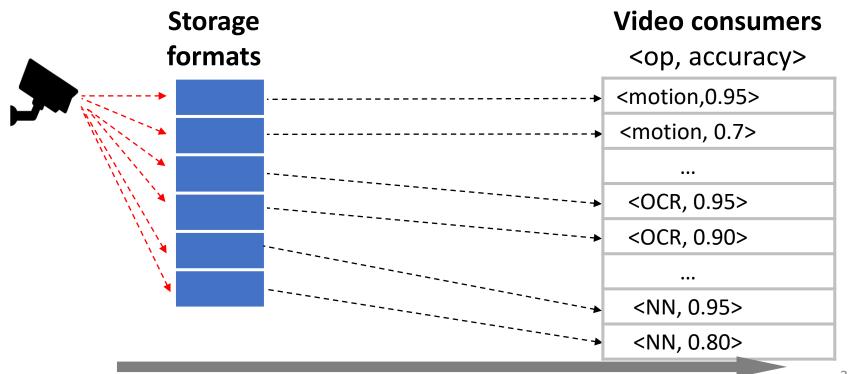




Storing all needed formats?

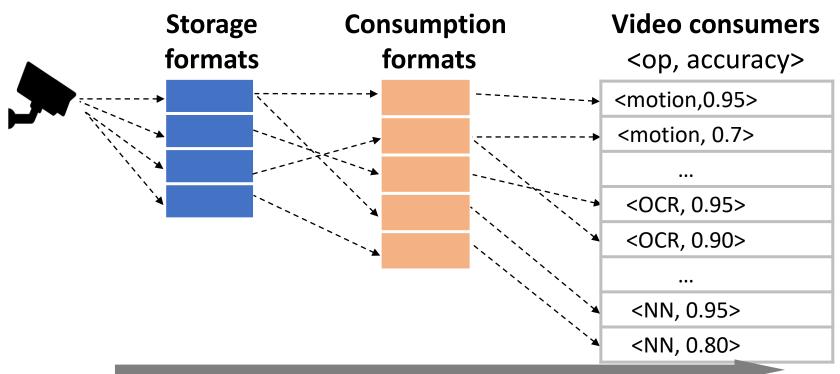
Ingestion & storage are expensive 🥦







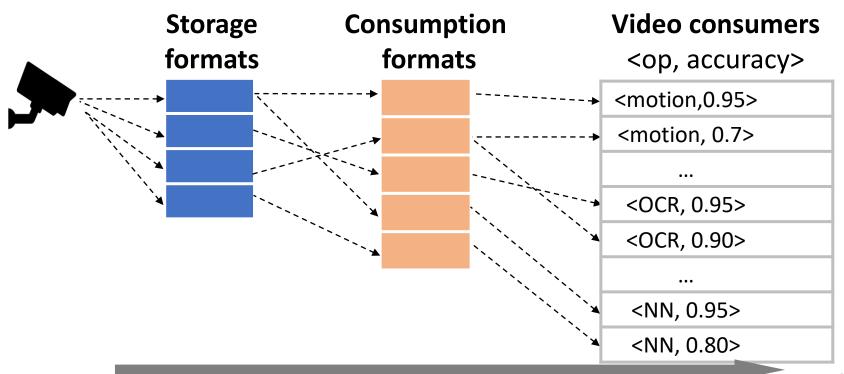
Proposal: Store minimum necessary formats





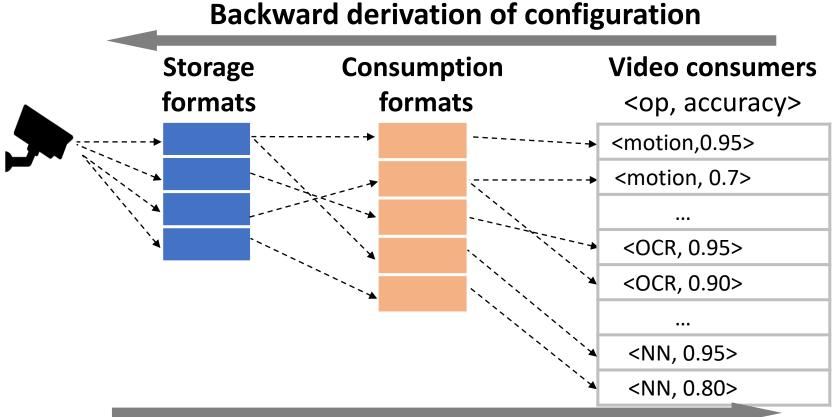
Proposal: Store minimum necessary formats

Key Challenge: Too many knobs to tune!

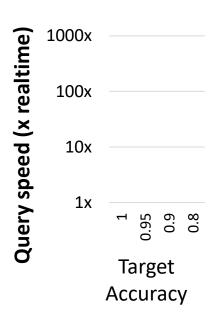




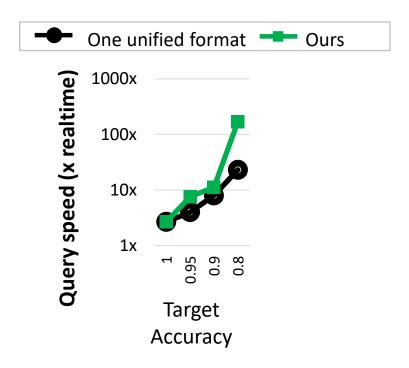
Proposal: Store minimum necessary formats & derive them automatically!



Benefit: faster analytics



Benefit: faster analytics

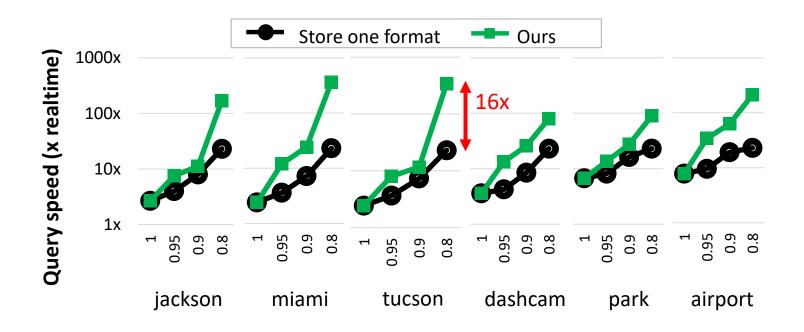


Why?

- Lower accuracy

 analytics tolerates cheaper video formats
- Video retrieval should be proportionally cheaper!

Benefit: faster analytics



Why?

- Lower accuracy

 analytics tolerates cheaper video formats
- Video retrieval should be proportionally cheaper!

Case 2: First data store for retro. video analytics

Baking analytics demands into storage decisions

 Showing clear advantage over generic file systems or storage layers

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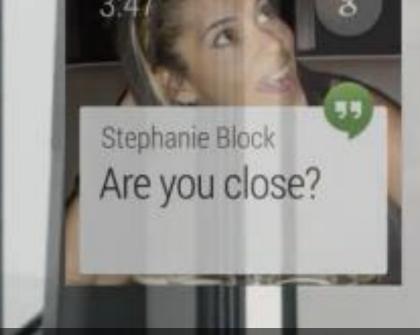
Case 3: enhancing the ossified Linux kernel

A new OS structure for device driver offloading

A commodity kernel still irreplaceable

- The major reason: device drivers
 - Linux: a common environment for all drivers
 - >70% Linux kernel source are device drivers
- All other OS functions have alternatives!

Problem: suspend/resume in ephemeral tasks



- Mobile/IoT run frequent, ephemeral tasks
 - Android smartwatch: each minute
 - Each task is short-lived
 - Smartwatch: 10 secs
 - Background task: < 1 sec

Under the hood...

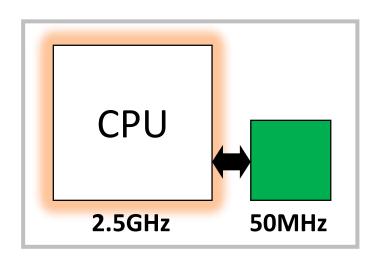
Wakeup Device Resume Thaw user **User Task** Freeze user&fs Device Suspend Sleep Time

Suspend/Resume Is Slow

	Nexus 5	Gear	Note 4	Panda
Suspend	119 ms	191 ms	231 ms	262 ms
Resume	88 ms	159 ms	316 ms	492 ms

Why? IO devices are slow

Proposal: run these driver paths on a low-power tiny core

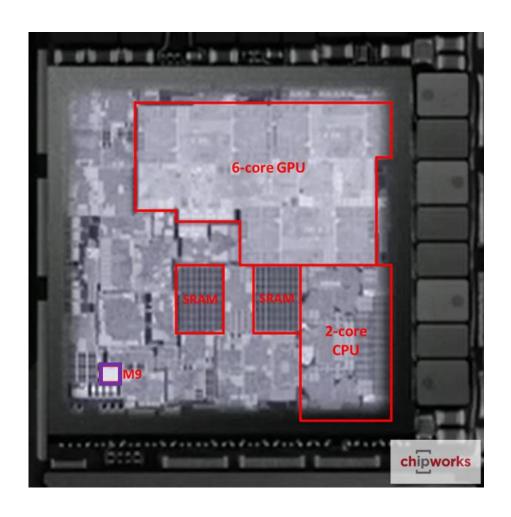


What are these tiny cores?

What are these tiny cores?



What are these tiny cores?

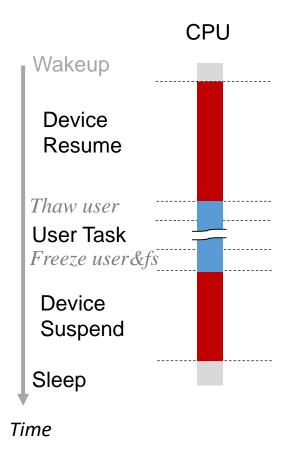


How can a tiny core help?

- Idling more efficiently
 - 3 mW vs 30 mW
- Executing kernel more efficiently
 - Small code working set
 - Less predictable control flow

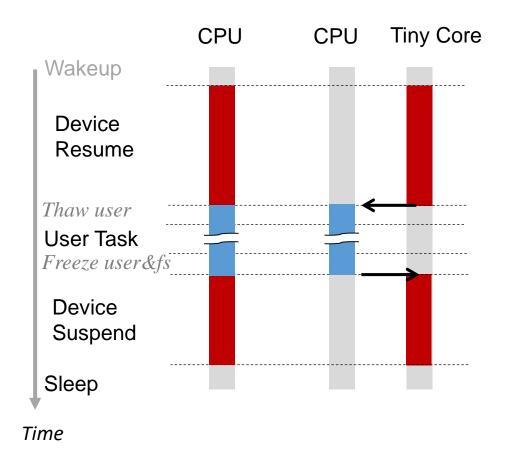
^{1.} J. Mogul, J. Mudigonda, N. Binkert, P. Ranganathan, and V. Talwar. Using asymmetric single-isa cmps to save energy on operating systems. Micro, IEEE, 2008

The workflows



Current

The workflows

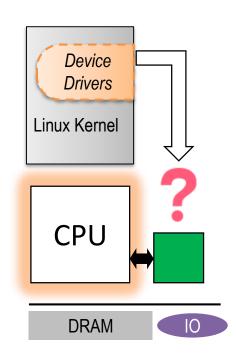


Current Proposed

How?

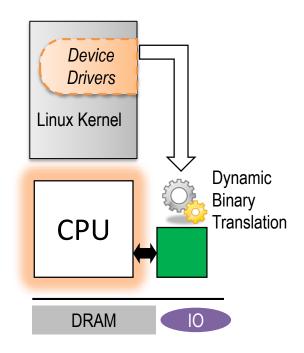
- The tiny core is very wimpy
 - Essentially a microcontroller
- Has a different ISA
 - More aggressive than big.LITTLE

- Re-engineering the kernel?
- "Linux kernel is the new firmware"



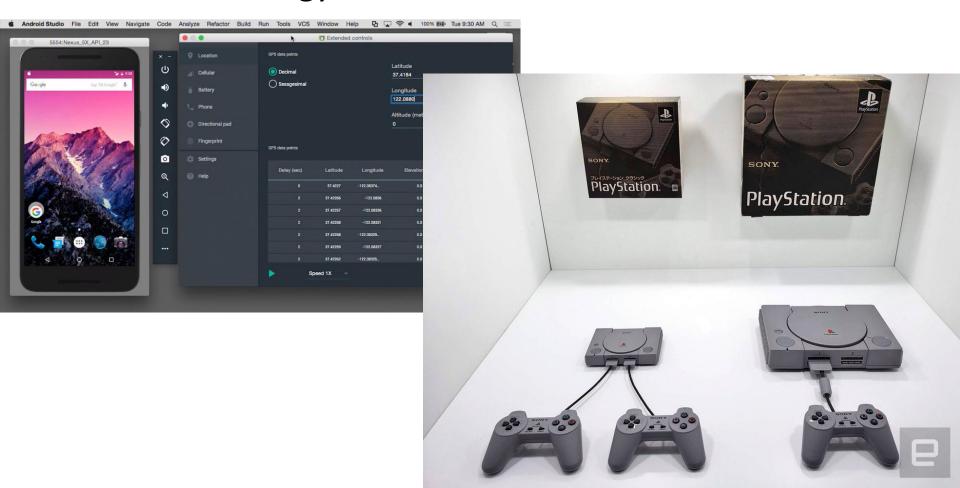
Proposal: use dynamic binary translation

 Empower the tiny core to execute unmodified kernel binaries



Dynamic binary translation in 30 sec

• The technology behind QEMU and a PS emulator



"This sounds impractical"

- Dynamic bin translation (DBT) is known expensive
 - > 10x overhead
- No one runs DBT on microcontrollers
 - Always weak over strong

Solution: the Transkernel model

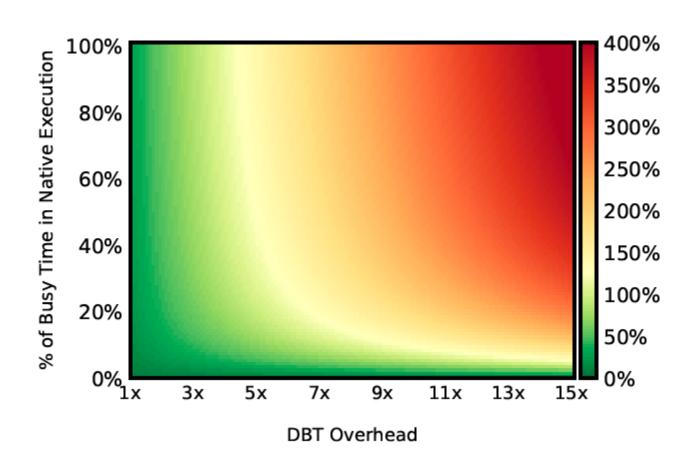
A transkernel Dynamic Device Binary **Drivers** Translation **Translated** code Linux Emu kernel Tiny core CPU DRAM

Key ideas:

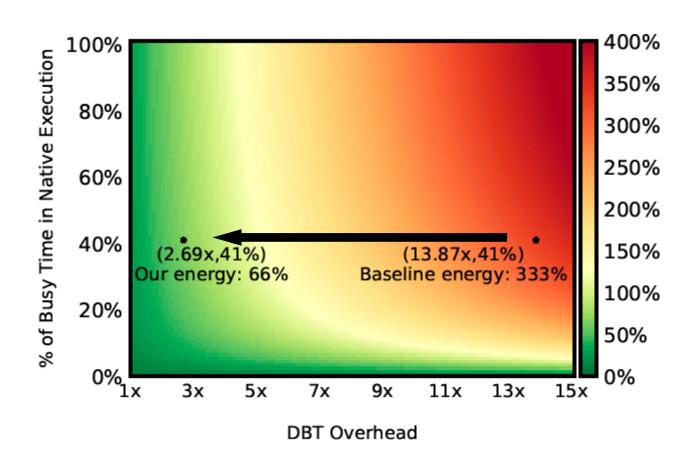
- 1. Translate most of driver code
- Emulate the remaining kernel infrastructure

Essentially a tiny VM for drivers!

Is this practical?



Only through careful optimization!



Case 3: Transkernel: tiny VMs for device driver paths

First DBT engine running on a microcontroller-like core!

Demonstrated:

- One can use DBT for efficiency gain
- The ossified kernel can be tamed!

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Reflection: a tale of two inquiries

- App-defined OS
 - Brave new world
 - Usefulness depends on the target apps
 - OS as a tool: serving AI and big data
- Kernel as the new firmware
 - Still lots of opportunities
 - Unique constraints. Require unconventional techniques
 - Like fixing an engine of an airplane in the air
 - OS as a craft. Fun!