

# Competitive programming

wu

September 26, 2022

## 1 Dynamic Programming

### 1.1 Digit DP

*Problem 1.1.1* (LeetCode 788: Rotated Digits). An integer  $x$  is a **good** if after rotating each digit individually by 180 degrees, we get a valid number that is different from  $x$ . Each digit must be rotated - we cannot choose to leave it alone.

A number is valid if each digit remains a digit after rotation. For example:

- 0, 1, and 8 rotate to themselves,
- 2 and 5 rotate to each other (in this case they are rotated in a different direction, in other words, 2 or 5 gets mirrored)
- 6 and 9 rotate to each other, and
- the rest of the numbers do not rotate to any other number and become invalid.

Given an integer  $n$ , return the number of good integers in the range  $[1, n]$ .

*Solution.* Given  $n$ . Let  $f(pos, bound, diff)$  be the number of good numbers satisfying

1. Only consider  $pos$ th digit and  $pos$  starts from left, which means 0th digit is the highest digit. And we assume the first  $pos - 1$  digits are fixed
2. If digits in  $[0, pos - 1]$  are first  $pos$  digits of  $n$ , then  $bound$  is true

3. If digits in  $[0, pos - 1]$  has at least one 2/5/6/9, then  $diff$  is true

Therefore the answer is  $f(0, true, false)$ , and the transition formula is

$$f(pos, bound, diff) = \sum f(pos + 1, bound', diff')$$

- $bound'$  is true iff  $bound$  is true and the digit we choose is the  $pos$ th digit of  $n$
- $diff'$  is true iff  $diff$  is true or we chose 2/5/6/9

□

## 2 Trick and Bit

### 2.1 Bit operation

*Problem 2.1.1* (Leetcode: Missing Two LCCI). You are given an array with all the numbers from 1 to  $N$  appearing exactly once, except for two number that is missing. How can you find the missing number in  $O(N)$  time and  $O(1)$  space?

You can return the missing numbers in any order.

Input	Output
[1]	[2, 3]
[2, 3]	[1, 4]

`nums.length <= 30000`

*Solution.* Suppose the missing two numbers are  $x_1$  and  $x_2$ , and if we add  $1, \dots, N$  to the end of the array  $A$ , then  $x = \oplus A = x_1 \oplus x_2$ .

By  $x \& -x$  we can get the lowest bit of  $x$ , assume it's in  $l$ th bit. Then we can assume  $x_1$ 's  $l$ th bit is 0, and  $x_2$ 's  $l$ th bit is 1, and we can partition  $A$  into  $A_1$  and  $A_2$  by whether the elements'  $l$ th bit is 1, then  $\oplus A_1 = x_1$  and  $\oplus A_2 = x_2$  □