

Lemma 11 Let $\lambda \geq \aleph_0$.

(1) \Leftrightarrow (2)

- 1.) If $A \subseteq M$, $|A| \leq \lambda \Rightarrow |S_1(A)| \leq \lambda$
- 2.) If $A \subseteq M$, $n < \omega$, $|A| \leq \lambda \Rightarrow |S_n(A)| \leq \lambda$

Proof (2) \Rightarrow (1) trivial

Assume (1). Prove (2). Assume $|A| \leq \lambda$

By induction, $|S_{n-1}(A)| \leq \lambda$.

$$S_{n-1}(A) = \{ \text{tp}(\bar{b}_\alpha / A) : \alpha < \lambda \} \quad (\text{for some } \bar{b}_\alpha)$$

For each α , $|S_1(A \bar{b}_\alpha)| \leq \lambda$ by (1).

$$A \bar{b}_\alpha = A \cup \{ b_{\alpha,1}, \dots, b_{\alpha,n-1} \}$$

$$S_1(A \bar{b}_\alpha) = \{ \text{tp}(c_{\alpha,\beta} / A \bar{b}_\alpha) : \beta < \lambda \}.$$



Claim If $(\bar{d}, e) \in M^{n-1}$ then $\exists \alpha, \beta$ such that

$$\text{tp}(\bar{d}e / A) = \text{tp}(\bar{b}_\alpha, c_{\alpha,\beta} / A)$$

Proof $\text{tp}(\bar{d}/A) \in S_{n-1}(A)$

" $\text{tp}(\bar{b}_\alpha / A)$ for some α

$(\bar{d}e)$

$\exists \sigma \in \text{Aut}(M/A) \quad \sigma(\bar{d}) = \bar{b}_\alpha$

$\downarrow \sigma$

$\text{tp}(\sigma(e) / A \bar{b}_\alpha) \in S_1(A \bar{b}_\alpha)$

$(\bar{b}_\alpha, \sigma(e))$

" $\text{tp}(c_{\alpha,\beta} / A \bar{b}_\alpha)$ for some β .

$\downarrow \tau$

$\exists \tau \in \text{Aut}(M/A \bar{b}_\alpha) \quad \tau(\sigma(e)) = c_{\alpha,\beta}$

$(\bar{b}_\alpha, c_{\alpha,\beta})$

Then $\tau \circ \sigma \in \text{Aut}(M/A)$ and $(\tau \circ \sigma)(\bar{d}e) = (\bar{b}_\alpha, c_{\alpha,\beta})$

So $\text{tp}(\bar{d}e / A) = \text{tp}(\bar{b}_\alpha, c_{\alpha,\beta} / A)$. \square claim

Claim $\Rightarrow S_n(A) = \{ \text{tp}(\bar{b}_\alpha, c_{\alpha,\beta} / A) : \alpha, \beta < \lambda \}$

$$|S_n(A)| \leq \lambda^2 = \lambda.$$

\square

Definition T is λ -stable if $\forall A \subseteq M$, $|A| \leq \lambda$ then ~~$|S_n(A)| \leq \lambda$~~ .

If $\lambda \geq |\mathbb{L}|$, then (1) \Leftrightarrow (2) \Leftrightarrow (3) \Leftrightarrow (4)

1.) If $A \subseteq M$, $|A| \leq \lambda \Rightarrow |S_1(A)| \leq \lambda$

(1) \Leftrightarrow (2)

2.) If $A \subseteq M$, $n < \omega$, $|A| \leq \lambda \Rightarrow |S_n(A)| \leq \lambda$

\uparrow \downarrow

(*) 3.) If $M \preccurlyeq M$, $|M| \leq \lambda \Rightarrow |S_1(M)| \leq \lambda$

(3) (4)

4.) If $M \preccurlyeq M$, $|M| \leq \lambda \Rightarrow |S_n(M)| \leq \lambda$

Proof (1) \Leftrightarrow (2) above.

(3) \Rightarrow (1) Assume (3). Let $A \subseteq M$, $|A| \leq \lambda$.

Löwenheim-Skolem $\Rightarrow \exists M \preccurlyeq M \supseteq A$, $|M| \leq |A| + |\mathcal{L}| \leq \lambda$

$S_1(M) \xrightarrow{\text{restriction}} S_1(A)$ is surjective. $|S_1(A)| \leq |S_1(M)| \leq \lambda$ by (3).

(4) \Rightarrow (2) similar.

Lemma 1 Suppose: $\forall M \preccurlyeq M$, $\forall p \in S_1(M)$ is definable.

Then T is λ -stable for some λ .

Proof Take $\lambda = 2^{|\mathcal{L}|} > |\mathcal{L}|$. Suppose $M \preccurlyeq M$, $|M| \leq \lambda$.

Every type in $S_1(M)$ is definable.

$p \in S_1(M)$ is determined by

$$\varphi \mapsto d_p \varphi. \\ \mathcal{L}\text{-fmla} \rightarrow \mathcal{L}(M)\text{-fmla}$$

$$|L(M)| \leq \lambda$$

$$|S_1(M)| \leq |L(M)|^{|\mathcal{L}|} \leq \lambda^{|\mathcal{L}|} = (2^{|\mathcal{L}|})^{|\mathcal{L}|} = 2^{|\mathcal{L}|^2} = 2^{|\mathcal{L}|} = \lambda.$$

□.

Theorem 2 TFAE

- 1) $M \models T$, $p \in S_n(M) \Rightarrow p$ is definable.
- 2) $M \models T$, $p \in S_1(M) \Rightarrow p$ is definable.
- 3) No $\varphi(\bar{x}, \bar{y})$ has the dichotomy property
- 4) No $\varphi(\bar{x}, \bar{y})$ has the dichotomy property
- 5) T is λ -stable for some λ .

Proof

(3) \Rightarrow (4) trivial

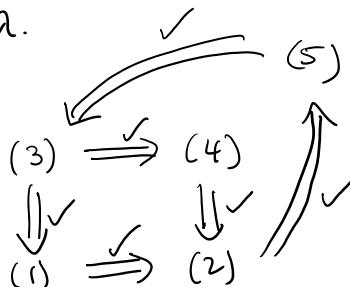
(1) \Rightarrow (2) trivial

(3) \Rightarrow (1) Prop 8 last week

(4) \Rightarrow (2) Prop 8

From (2) \Rightarrow (5) Lemma 2.

(5) \Rightarrow (3) Prop 10.



□.

T is STABLE if those conditions hold.

Fact If T is not stable, then for $\lambda > \lambda_0$, T has 2^λ -many models of size λ .

Cohesive

Def 3 If $M \preccurlyeq N$, if $p \in S_n(M)$, if $q \in S_n(N)$, then q is a coheir of p if $q \supseteq p$ and q is finitely supported for any $\varphi(\bar{x}) \in q(\bar{x})$, $\exists \bar{a} \in M$, $N \models \varphi(\bar{a})$.

($L(N)$ -formula)

If $\varphi_1, \dots, \varphi_m \in q$, then $\bigwedge_{i=1}^m \varphi_i(\bar{x}) \in q(\bar{x})$, so $\exists \bar{a} \in M : N \models \bigwedge_{i=1}^m \varphi_i(\bar{a})$, so $\{\varphi_1(\bar{x}), \dots, \varphi_m(\bar{x})\}$ is satisfiable in M .

q is a coher (of $q \upharpoonright M$) iff q is f.sat. in M .

Example $Q^{\text{alg}} \leq C$ $q = tp(\pi/C)$ $p = tp(\pi/(Q^{\text{alg}}))$
 $q \supseteq p$, but q isn't a coher
 $(x=\pi) \in q(x)$, not satisfiable in Q^{alg} .

Example If $M \leq N$ strongly minimal, $q(x) \in S_1(N)$, $p(x) \in S_1(M)$ transcendental
then q is a coher of p . $q \supseteq p$ ✓
 q is fsat in M : if $\varphi(x) \in q(x)$, then $\varphi(N)$ is cofinite
 M is infinite, so $M \cap \varphi(N) \neq \emptyset$. $\exists a \in M. N \models \varphi(a)$.

Lemma 4 If $M \leq N$, $\Sigma(\bar{x})$ partial type over N , $\Sigma(\bar{x})$ is f.sat. in M
THEN $\exists q(\bar{x}) \in S_n(N)$, $q(\bar{x})$ is fsat. in M .

Proof Let $\Psi(\bar{x}) = \{ \psi(\bar{x}) \in L(N) : \forall a \in M \quad N \models \psi(a) \}$

If $\bar{a} \in M$, then \bar{a} satisfies Ψ
 $\Sigma(\bar{x})$ fsat in $M \Rightarrow \Sigma(\bar{x}) \cup \Psi(\bar{x})$ is fsat. $\Rightarrow q \in S_n(N) \quad q \supseteq \Sigma \cup \Psi$.
If q isn't fsat. in M then $\varphi(\bar{x}) \in q(\bar{x})$, $\varphi(\bar{x})$ not sat. in M .
 $\neg \varphi(\bar{x}) \in \Psi(\bar{x}) \subseteq q(\bar{x})$, so No! II.

Theorem 5 If $p \in S_n(M)$, $N \geq M$, ~~then~~ then $\exists q \in S_n(N)$,
 q is a coher of p .

Proof Lemma 4 with $\Sigma = p$.

Theorem 6 Suppose $M_1 \leq M_2 \leq M_3$, $p_1 \in S_n(M_1)$, $p_2 \in S_n(M_2)$
 p_2 is a coher of p_1 . THEN $\exists p_3 \in S_n(M_3)$, p_3 is a coher of p_1 and p_2
Proof Apply lemma 4 ~~to~~ to $\Sigma = p_2$ $M = M_1$, $N = M_3$. p_3 is f.sat. in M_1 , $p_3 \supseteq p_2$ D.

Warning "coher" is not transitive.

Invariant types

M monster model If $A \subseteq M$ small

$$tp(\bar{a}/A) = tp(\bar{b}/A) \Leftrightarrow \exists \sigma \in \text{Aut}(M/A) [\sigma(\bar{a}) = \bar{b}]$$

$\bar{a} \equiv_A \bar{b}$ means this

$$\{ \sigma \in \text{Aut}(M) : \sigma \text{ fixes } A \text{ pointwise} \} \\ \forall x \in A \quad \sigma(x) = x$$

Lemma 8 If $X \subseteq M^n$, TFAE

- 1) $\sigma(X) = X$ if $\sigma \in \text{Aut}(M/A)$
- 2) If $\bar{a}, \bar{b} \in M^n$, $\bar{a} \equiv_A \bar{b} \Rightarrow (\bar{a} \in X \Leftrightarrow \bar{b} \in X)$

3) There is $f: S_n(A) \rightarrow \{0, 1\}$ such that $\bar{a} \in X \Leftrightarrow f(\text{tp}(\bar{a}/A)) = 1$.

Proof (2) \Leftrightarrow (3) easy.

(1) \Leftrightarrow (2): Rewrite (2) as

- If $\bar{a}, \bar{b} \in M$, $\sigma \in \text{Aut}(M/A)$, and $\sigma(\bar{a}) = \bar{b}$, then
 $\bar{a} \in X \Leftrightarrow \bar{b} \in X$.

i.e.

- If $\bar{a} \in M$, $\sigma \in \text{Aut}(M/A)$ then
 $\bar{a} \in X \Leftrightarrow \sigma(\bar{a}) \in X$

i.e.

- $X = \sigma^t(X)$

- $X = \sigma(X)$. □

Def 9 $X \subseteq M^n$ is $\text{Aut}(M/A)$ -invariant or A -invariant if

Lemma 8 (1) \rightarrow (3) hold.

Example If X is A -definable ($X = \varphi(M^n)$, $\varphi \in L(A)$)

then X is A -invariant ($\text{tp}(\bar{b}/A)$ determines whether $M \models \varphi(\bar{b})$.)

Lemma 10 If $D \subseteq M^n$ is definable and A -invariant,
then D is A -definable.

Proof: Compactness ...

Step 1 If $\bar{x} \in D$ then $\text{tp}(\bar{x}/A) \vdash \bar{x} \in D$
by "compactness", $\exists \varphi(\bar{x}) \in \text{tp}(\bar{x}/A)$ $\varphi(\bar{x}) \vdash \bar{x} \in D$.
 $\varphi(M^n) \subseteq D$

Step 2 So then D is covered by A -definable sets-subsets of D .

by "compactness", D is covered by finitely many of them,

$\Rightarrow D$ is A -definable. □

Def 11 A global type is $p \in S_n(M)$.

Def 12 $p \in S_n(M)$ is A -invariant (for small $A \subseteq M$) if

$$\sigma(p) = p \quad \forall \sigma \in \text{Aut}(M/A).$$

Equivalently, $\forall \varphi(\bar{x}, \bar{y})$, the set

~~$\{\bar{b} \in M : \varphi(\bar{x}, \bar{b}) \in p(\bar{x})\}$~~ is A -invariant.

Equivalently, if $\bar{b} \equiv_{\bar{A}} \bar{c}$ then $\varphi(\bar{x}, \bar{b}) \in p(\bar{x}) \Leftrightarrow \varphi(\bar{x}, \bar{c}) \in p(\bar{x})$.

Def 13 p is A -definable if $\forall \varphi, \{b \in M : \varphi(\bar{x}, b) \in p(\bar{x})\}$ is A -definable.

Remark 14 1) p is A -definable $\Rightarrow p$ is A -invariant

2) If p is definable, then p is A -invar. $\Leftrightarrow p$ is A -definable.

3) If p is definable then p is A -definable for some small A
Each $d_p \varphi$ uses only fin. many parameters $|A| \leq |L|$.

Prop 15 Suppose $M \preccurlyeq M$ small.

1) If $p \in S_n(M)$ definable and p^M is its heir over M , then $p^M \in S_n(M)$ is M -definable.

2) $p \mapsto p^M$ is a bijection from definable types over M $\not\cong$ to M -definable types over M .

Proof 1) p^M has same def. as p , so it's M -definable

2) $q \mapsto q|M$ is an inverse to $p \mapsto p^M$

If q is M -definable, then $q|M$ is definable and $q = (q|M)^M$.

Warning 16 an M -invariant type p is not determined by $p|M$

If $A \subseteq M$, A -definable type p is not determined by $p|A$.

Theorem 17 Suppose $M \preccurlyeq M$, $p \in S_n(M)$

1) If $q \in S_n(M)$ and q is a coherr of p , then q is M -invariant.

2) $\exists q \in S_n(M)$ q is M -invariant, $q \not\cong p$.

Proof (1) \Rightarrow (2) by Theorem 5.

(1): If q a coherr of p , but q is not M -invariant, then $\exists \bar{b}, \bar{c}, \varphi$

$\bar{b} \equiv_M \bar{c}$ but $\varphi(\bar{x}, \bar{b}) \in q(\bar{x})$ $\varphi(\bar{x}, \bar{c}) \notin q(\bar{x})$.

$\varphi(\bar{x}, \bar{b}) \wedge \neg \varphi(\bar{x}, \bar{c}) \in q(\bar{x})$.

q is fsat. in M , $\exists \bar{a} \in M$ $M \models \varphi(\bar{a}, \bar{b}) \wedge \neg \varphi(\bar{a}, \bar{c})$.

$\Rightarrow \bar{b} \not\equiv_M \bar{c}$ since $\varphi(\bar{a}, \bar{y})$ is true for \bar{b} , false for \bar{c} . \square

In stable theories:

Lemma 19: If T is stable and p is A -invariant,
then p is A -definable.

\square .

Theorem 20 Suppose T stable, $M \preccurlyeq M$ small, $p \in S_n(M)$.

Let p^M be the global heir.

- 1) p^M is the only M -invariant global type extending p
- 2) p^M is the only global coher of p
- 3) If $M \preceq N \preceq M$ and q is the heir of p over N , then q is the unique coher of p over N .

Proof 1) M -invar \Leftrightarrow M -definable, use Prop 15
 p^M is the only M -definable type $\supseteq p$.

2) \exists some coher of p . Any coher is M -invariant (Theorem 17)
so p^M is the only coher.

3) If q' is some coher over N , Theorem 6 gives $r \in S_n(M)$
 $r \geq q'$, r is a coher of p and q' .

By (2), $r = p^M$, then $q' = (p^M \upharpoonright N) = q$ ^{the heir over N .}

Cor 21 In a stable theory, cohers are unique and cohers = heirs.

Cor 22 In a stable theory, "coher" is transitive. (because "heir" is.)

Notation If $A \subseteq M$, $p \in S_n(A)$ then $\bar{b} \models p$ means
 \bar{b} satisfies p , ie $\text{tp}(\bar{b}/A) = p$.

Lemma 23 If p, q are A -invariant global types $p \in S_n(M)$, $q \in S_m(M)$
then there is $r \in S_{n+m}(A)$ such that $(\bar{b}, \bar{c}) \models r$ iff

$$\bar{b} \models p \upharpoonright A \quad \text{and} \quad \bar{c} \models q \upharpoonright (A\bar{b}) \quad (*)$$

Proof Let $X = \{(\bar{b}, \bar{c}) : \bar{b} \models p \upharpoonright A \text{ and } \bar{c} \models q \upharpoonright (A\bar{b})\}$

if $(\bar{b}, \bar{c}) \in X$ and $\sigma \in \text{Aut}(M/A)$, then

$$\not\exists \sigma(\bar{b}) \models p \upharpoonright A \quad \text{and} \quad \sigma(\bar{c}) \models q \upharpoonright (A\sigma(\bar{b})),$$

so $\sigma(\bar{b}, \bar{c}) \in X$.

X is A -invariant.

Fix $\not\exists \bar{b}_0 \models p \upharpoonright A$, $\bar{c}_0 \models q \upharpoonright (A\bar{b}_0)$ so $(\bar{b}_0, \bar{c}_0) \in X$

Let $r = \text{tp}(\bar{b}_0, \bar{c}_0 / A)$. If $(\bar{b}, \bar{c}) \models r$ then $(\bar{b}, \bar{c}) \in X$.

Conversely if $(\bar{b}, \bar{c}) \in X$, want $(\bar{b}, \bar{c}) \models r$ i.e. $(\bar{b}, \bar{c}) \equiv_A (\bar{b}_0, \bar{c}_0)$.

$$\bar{b} \models p \upharpoonright A = \text{tp}(\bar{b}_0 / A) \quad \text{so} \quad \bar{b} \equiv_A \bar{b}_0, \exists \sigma \in \text{Aut}(A)$$

$\sigma(\bar{b}) = \bar{b}_0$. Replace (\bar{b}, \bar{c}) with $(\sigma(\bar{b}), \sigma(\bar{c})) = (\bar{b}_0, \sigma(\bar{c}))$.

WMA $\bar{b} = \bar{b}_0$. Then \bar{c} and \bar{c}_0 both satisfy $q \upharpoonright (A\bar{b}_0)$

so $\bar{c} \equiv_{A\bar{b}_0} \bar{c}_0$ Move \bar{c} by $\tau \in \text{Aut}(M / A\bar{b}_0)$,

we may assume $\bar{c} = \bar{c}_0$. Then $(b, \bar{c}) = (b_0, \bar{c}_0)$

$$b\bar{c} = b_0\bar{c} \underset{A}{=} b_0\bar{c}_0.$$

□.

Prop 24 If $p \in S_n(M)$, $q \in S_m(M)$ and both are A -invariant, then $\exists A$ -invariant $p \otimes q \in S_{n+m}(M)$ such that for any small A' with $A \subseteq A' \subseteq M$,

$$(\ast \ast) \quad (b, \bar{c}) \models (p \otimes q) \upharpoonright A' \iff [b \models p \upharpoonright A' \text{ and } \bar{c} \models q \upharpoonright A' \bar{b}].$$

Proof ~~Lemma 23~~ Note p, q are A' -invariant for $A' \supseteq A$ so Lemma 23 gives $r_{A'} \in S_{n+m}(A')$ for each $A' \supseteq A$ such that $(b, \bar{c}) \models r_{A'} \iff [b \models p \upharpoonright A' \text{ and } \bar{c} \models q \upharpoonright A' \bar{b}]$.

Note if $A'' \supseteq A' \supseteq A$, if $(b, \bar{c}) \models r_{A''} \Rightarrow (b, \bar{c}) \models r_{A'}$
So $r_{A'} = r_{A''} \upharpoonright A'$.

Let $p \otimes q = \bigcup_{A'} r_{A'}$. Then $p \otimes q \in S_{n+m}(M)$ and $r_{A'} = (p \otimes q) \upharpoonright A'$ for $A' \supseteq A$.

By def. of $r_{A'}$, $(\ast \ast)$ holds.

If $\sigma \in \text{Aut}(M/A)$, $\sigma(p \otimes q) = \sigma(p) \otimes \sigma(q) = p \otimes q$.

So $p \otimes q$ is A -invariant.

~~Fact~~ If $p \in S_n(M)$ p is "A-invariant" where M is $|A|^t$ -saturated (if $b, \bar{c} \in M$, $b \underset{A}{=} \bar{c} \Rightarrow \text{acl}(b) \in p \Leftrightarrow \text{acl}(\bar{c}) \in p$)

and $N \not\supseteq M$, then p has a unique "A-invariant" extension over N .

$p|N$ denotes this extension.

~~Fact~~ If $p, q \in S_{n+m}(M)$ A -invariant, Take $\bar{b} \models p \quad \bar{b} \in M, \bar{b} \not\in M$.

Take ~~such that~~ $\bar{c} \models q \upharpoonright M$, then $\text{tp}(\bar{b}, \bar{c}/M) = p \otimes q$.

Def $p \otimes q$ is The (Morley) product of invariant types p, q is $p \otimes q$ from Prop. 24.

If p, q are A -invar. then $(b, \bar{c}) \models (p \otimes q) \upharpoonright A \iff (b \models p \upharpoonright A \text{ and } \bar{c} \models q \upharpoonright A \bar{b})$

Example If T is strongly minimal and $p = q = \text{transcendental type}$,

what $p \otimes$ is $(p \otimes p)(x, y)$?

Def (In any theory) $\text{acl}(A) = \bigcup_{b \in \text{acl}(A)} \{\varphi(x) : \varphi(x) \in L(A), |\varphi(M)| < \infty\}$.

$b \in \text{acl}(A)$ if $\exists L(A)\text{-formula } \varphi(x), b \in \varphi(M) \subseteq_f M$.

example: in ACF, $\Sigma \in \text{acl}(\emptyset)$ because $\varphi(x) = (x^2 - 2 = 0)$.

Fact in ACF, if K subfield of M , then $\text{acl}(K)$ is K^{alg} field theoretic closure

Fact In any theory T , $\text{acl}(-)$ is a closure operation.

$$X \subseteq Y \Rightarrow \text{acl}(X) \subseteq \text{acl}(Y)$$

$$X \subseteq \text{acl}(X)$$

$$\text{acl}(\text{acl}(X)) = \text{acl}(X)$$

$$\text{acl}(\{X\}) = \bigcup_{X_0 \subseteq_p X} \text{acl}(X_0).$$

"transcendental" means "not algebraic".

If T is strongly minimal & $p \in S_1(M)$ transcendental

$$b \models p \upharpoonright A \Leftrightarrow b \notin \text{acl}(A).$$

$\varphi(x) \in p \Leftrightarrow \varphi(M)$ is cofinite ($\Leftrightarrow \varphi(M)$ not finite)

$$p \upharpoonright A = \{ \varphi(x) \in L(A) : \varphi(M) \text{ is cofinite} \}$$

Therefore $(b, c) \models (p \otimes p) \upharpoonright A$ iff $b \models p \upharpoonright A$ and $c \models p \upharpoonright Ab$
iff $b \notin \text{acl}(A)$ and $c \notin \text{acl}(Ab)$.

Idea: b, c are algebraically independent over A .

Idea (in stable theories), $(p \otimes q)(x, y)$ is the "most free" completion of $p(x) \vee q(y)$.

Example Suppose $M \models \text{ACF}$. Let p_V denote generic type of a variety $V \subseteq M^n$. $\{x \in V\} \cup \{x \notin W : W \subsetneq V, W \text{ algebraic}\}$

If $V \subseteq M^n, W \subseteq M^m$ varieties, then $V \times W$ is a variety and $p_V \otimes p_W$ is $p_{V \times W}$.

Proof $p_V \otimes p_W = p_z$ for some variety $z \subseteq M^{n+m}$. Goal: $z = V \times W$.

Take small $M \leq M$ s.t. V, W, z are M -definable.

Take $\bar{a} \models p_V \upharpoonright M$, take small $N \leq M$ $N \supseteq M\bar{a}$,

Take $\boxed{b \models p_W \upharpoonright N}$ ($\Rightarrow b \models p_W \upharpoonright M\bar{a}$, so $(\bar{a}, b) \models p_V \otimes p_W \upharpoonright M$
 $= p_z \upharpoonright M$).

$\bar{a} \in V$ ($\bar{x} \in V \in p_V \upharpoonright M$). $b \in W$, so $(\bar{a}, b) \in V \times W$.

Fact $p_z(\bar{x}) \vdash \bar{x} \in U \Leftrightarrow z \subseteq U$, for U algebraic

So $(\bar{a}, b) \in V \times W \Rightarrow z \subseteq V \times W$.

Suppose $z \subsetneq V \times W$. Take $(\bar{a}_0, \bar{b}_0) \in V \times W \setminus z$, $\bar{a}_0, \bar{b}_0 \in M$.

Let $z_{\bar{a}} = \{ \bar{y} \in M : (\bar{a}, \bar{y}) \in z \}$.

Then $z_{\bar{a}}$ is an algebraic set over $N \supseteq M\bar{a}$, $(\bar{a}, \bar{b}) \in z$
 $\bar{b} \models p_W \upharpoonright N \Rightarrow W \subseteq z_{\bar{a}}$. $\Rightarrow \bar{b} \in z_{\bar{a}}$.



Then $\bar{b}_0 \in w \subseteq Z_{\bar{a}}$ $(\bar{a}, \bar{b}_0) \in Z$.

Let $Z^{\bar{b}_0} = \{\bar{x} \in M; (\bar{x}, \bar{b}_0) \in Z\}$. $Z^{\bar{b}_0}$ is M -definable algebra.

$\bar{a} \in Z^{\bar{b}_0}$, $\bar{a} \models p_v \upharpoonright M$ so $V \subseteq Z^{\bar{b}_0}$.

$\bar{a}_0 \in V \subseteq Z^{\bar{b}_0}$ so $(\bar{a}_0, \bar{b}_0) \in Z \Rightarrow$.

$Z = V \times W$.

□.

Def \neq invariant types P, q "commute" if $(P \otimes q)(\bar{x}, \bar{y}) = (q \otimes P)(\bar{y}, \bar{x})$.

That is, for $C \subseteq M$ with p, q C -invariant,

$\bar{a} \models p \upharpoonright C$ and $\bar{b} \models q \upharpoonright C \bar{a} \Rightarrow (\bar{b} \models q \upharpoonright C \text{ and}) \bar{a} \models p \upharpoonright C \bar{b}$

Example In ACF, any two types commute:

$$(P_v \otimes P_w)(\bar{x}, \bar{y}) = P_{V \times W}(\bar{x}, \bar{y}) = P_{W \times V}(\bar{y}, \bar{x}) = (P_w \otimes P_v)(\bar{y}, \bar{x}).$$

$$(d_p x) \varphi(\bar{x}, \bar{y}) = \psi(\bar{y}) \quad d\varphi(\bar{y}) \text{ is } \exists^\infty x \varphi(x, \bar{y})$$

$$\exists^\infty x \quad (d_{p \otimes q} x, y) \varphi(x, y, z) \text{ is } (d_p x)(d_q y) \varphi(x, y, z).$$

$$\begin{matrix} \exists^\infty x \exists^\infty y \\ \exists^\infty y \exists^\infty x. \end{matrix}$$