

Heirs and definable types

Advanced Model Theory

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Reference in the book: Sections 11.1–2 (beware typos!)

Motivation. Recall that a theory T is κ -categorical if T has a unique model of size κ , up to isomorphism. Our goal for the course is to prove *Morley's theorem*:

Let T be a theory in a countable language. If T is κ -categorical for one uncountable cardinal κ , then T is κ -categorical for all uncountable cardinals κ .

Theories that are κ -categorical for one (hence all) $\kappa > \aleph_0$ are called *uncountably categorical theories*.

The statement of Morley's theorem is not as interesting as the proof, which will be a tour of **stability theory**.

1 Definable types

Suppose $p(\bar{x}) \in S_n(M)$.

Definition 1. $p(\bar{x})$ is a *definable type* if for every formula $\varphi(\bar{x}; \bar{y})$, the set

$$\{\bar{b} \in M : \varphi(\bar{x}, \bar{b}) \in p(\bar{x})\}$$

is definable, defined by some $L(M)$ -formula $d\varphi(\bar{y})$.

Example (Constant types). If $p = \text{tp}(\bar{a}/M)$ for $\bar{a} \in M$, then

$$\varphi(\bar{x}, \bar{b}) \in p(\bar{x}) \iff M \models \varphi(\bar{a}, \bar{b}),$$

so p is definable with $d\varphi(\bar{y}) = \varphi(\bar{a}, \bar{y})$.

Definition 2. T is *strongly minimal* if for any model $M \models T$, M is infinite and for any definable $X \subseteq M^1$, either X or $M \setminus X$ is finite.

ACF, the theory of *algebraically closed fields* like $(\mathbb{C}, +, \cdot)$, is strongly minimal.

Proposition 3. *If T is strongly minimal and $M \models T$, there is a 1-type $p(x) \in S_1(M)$ such that*

$$\varphi(x, \bar{b}) \in p(x) \iff \exists^\infty a \in M : M \models \varphi(a, \bar{b}).$$

Moreover, $p = \text{tp}(c/M)$ for any $N \succeq M$ and $c \in N \setminus M$.

Proof. Take $N \succ M$ and $c \in N \setminus M$; let $p(x) = \text{tp}(c/M)$. We must show

$$N \models \varphi(c, \bar{b}) \iff \exists^\infty a \in M : M \models \varphi(a, \bar{b}).$$

Suppose (\Rightarrow) fails, so $N \models \varphi(c, \bar{b})$ but $\varphi(M, \bar{b})$ is a finite set $\{e_1, \dots, e_n\} \subseteq M$. Then $\forall x (\varphi(x, \bar{b}) \rightarrow \bigvee_{i=1}^n (x = e_i))$ holds in M and therefore N , and so $c \in \{e_1, \dots, e_n\} \subseteq M$, a contradiction. Thus (\Rightarrow) holds.

Conversely, suppose $N \models \neg \varphi(c, \bar{b})$. By (\Rightarrow) , the set $\neg \varphi(M, \bar{b})$ is infinite. By strong minimality, the complement $\varphi(M, \bar{b})$ is finite. This proves (\Leftarrow) . \square

$p(x)$ is called the *transcendental 1-type*

Proposition 4. *1. T eliminates the \exists^∞ quantifier: for any formula $\varphi(x, \bar{y})$ there is a formula $\psi(\bar{y})$ equivalent to $\exists^\infty x \varphi(x, \bar{y})$.*

2. If $M \models T$, the transcendental 1-type $p \in S_1(M)$ is definable.

Proof. 1. Fix $\varphi(x, \bar{y})$. We claim there is $n_\varphi < \omega$ such that for every $M \models T$ and $\bar{b} \in M$,

$$|\varphi(M, \bar{b})| < n_\varphi \text{ or } |\neg \varphi(M, \bar{b})| < n_\varphi.$$

Otherwise, by compactness we can find $M \models T$ and $\bar{b} \in M$ such that $\varphi(M, \bar{b})$ and its complement have size $> n$ for all finite n , contradicting strong minimality. Thus n_φ exists, and then

$$\underbrace{\exists^\infty x \varphi(x, \bar{b})}_{\text{not first-order}} \iff \underbrace{\exists^{\geq n_\varphi} x \varphi(x, \bar{b})}_{\text{first-order}}.$$

2. For each $\varphi(x, \bar{y})$, $d\varphi(\bar{y})$ is the formula (equivalent to) $\exists^\infty x \varphi(x, \bar{y})$. \square

Corollary 5. *If $p \in S_1(M)$ and M is strongly minimal, then p is definable.*

Definition 6. A theory T is *stable* if all n -types over models are definable.

Motivation. Later we will see the following:

- Strongly minimal theories are stable.
- Uncountably categorical theories are stable.
- Strongly minimal complete theories are uncountably categorical, and uncountably categorical theories are very close to being strongly minimal.

2 Heirs and strong heirs

Suppose $M \preceq N$ and $p \in S_n(M)$. An *extension* or *son* of p is $q \in S_n(N)$ with $q \supseteq p$, i.e., $p = q \upharpoonright M$.

Definition 7 (Heirs). $q \in S_n(N)$ is an *heir* of p , written $p \sqsubseteq q$, if for any $\varphi(\bar{x}, \bar{b}, \bar{c}) \in q(\bar{x})$ with $\bar{b} \in M$ and $\bar{c} \in N$, there is $\bar{c}' \in M$ with $\varphi(\bar{x}, \bar{b}, \bar{c}') \in p(\bar{x})$.

We will see soon that (1) definable types have unique heirs, (2) this property characterizes definable types, and (3) the unique heir is a definable type with the same definition as the original type.

The notation \sqsubseteq and \supseteq isn't very standard, except in the special case of stable theories.

Lemma 8. Suppose $M_1 \preceq M_2 \preceq M_3$ and $p_i \in S_n(M_i)$ for $i = 1, 2, 3$, with $p_1 \subseteq p_2 \subseteq p_3$.

1. If $p_1 \sqsubseteq p_2 \sqsubseteq p_3$, then $p_1 \sqsubseteq p_3$.

2. If $p_1 \sqsubseteq p_3$, then $p_1 \sqsubseteq p_2$.

Proof. 1. Suppose $\varphi(\bar{x}, \bar{b}, \bar{c}) \in p_3(\bar{x})$ with $\bar{b} \in M_1$ and $\bar{c} \in M_3$. Then there is $\bar{c}' \in M_2$ with $\varphi(\bar{x}, \bar{b}, \bar{c}') \in p_2(\bar{x})$. Then there is $\bar{c}'' \in M_1$ with $\varphi(\bar{x}, \bar{b}, \bar{c}'') \in p_1(\bar{x})$.

2. Suppose $\varphi(\bar{x}, \bar{b}, \bar{c}) \in p_2(\bar{x})$ with $\bar{b} \in M_1$ and $\bar{c} \in M_2$. Then $\varphi(\bar{x}, \bar{b}, \bar{c}) \in p_3(\bar{x})$, so there is $\bar{c}' \in M_1$ with $\varphi(\bar{x}, \bar{b}, \bar{c}') \in p_1(\bar{x})$. \square

Definition 9. If $p \in S_n(M)$, then (M, dp) is the expansion of M by relation symbols $d\varphi(\bar{y})$ for each $\varphi(\bar{x}, \bar{y})$, interpreted as follows:

$$(M, dp) \models d\varphi(\bar{b}) \iff \varphi(\bar{x}, \bar{b}) \in p(\bar{x}).$$

Remark. p is definable if and only if the new relations in (M, dp) are definable in the old structure M .

Remark 10. The class of structures of the form (M, dp) with $M \models T$ and $p \in S_n(M)$ is an elementary class, axiomatized by T plus the following:

$$\begin{aligned} \forall \bar{y}_1, \dots, \bar{y}_m \left(\bigwedge_{i=1}^m d\varphi_i(\bar{y}_i) \rightarrow \exists \bar{x} \bigwedge_{i=1}^m \varphi_i(\bar{x}, \bar{y}_i) \right) & \text{ for formulas } \varphi_1(\bar{x}, \bar{y}_1), \dots, \varphi_n(\bar{x}, \bar{y}_n) \\ \forall \bar{y} (d\varphi(\bar{y}) \vee d\neg\varphi(\bar{y})) & \text{ for each formula } \varphi(\bar{x}, \bar{y}) \end{aligned}$$

The first axiom schema expresses finite satisfiability of p , and the second expresses completeness of p .

Lemma 11. If $(M, dp) \preceq (N, dq)$, then $M \preceq N$ and q is an heir of p .

Proof. The fact that $(N, dq) \succeq (M, dp)$ implies $N \succeq M$. Then:

- $q \supseteq p$: if $\varphi(\bar{x}, \bar{b}) \in p(\bar{x})$ (with $\bar{b} \in M$), then $(M, dp) \models d\varphi(\bar{b})$, so $(N, dq) \models d\varphi(\bar{b})$, and $\varphi(\bar{x}, \bar{b}) \in q(\bar{x})$.
- $q \sqsupseteq p$: suppose $\varphi(\bar{x}, \bar{b}, \bar{c}) \in q(\bar{x})$, with $\bar{b} \in M$ and $\bar{c} \in N$. Then $(N, dq) \models d\varphi(\bar{b}, \bar{c})$, and $(N, dq) \models \exists \bar{z} d\varphi(\bar{b}, \bar{z})$. Then $(M, dp) \models \exists \bar{z} d\varphi(\bar{b}, \bar{z})$, so there is $\bar{c}' \in M$ with $M' \models d\varphi(\bar{b}, \bar{c}')$ and $\varphi(\bar{x}, \bar{b}, \bar{c}') \in p(\bar{x})$. \square

Corollary 12. *If $p \in S_n(M)$, then there is $M_0 \preceq M$ with $|M_0| \leq |T|$, such that $p \sqsupseteq (p \upharpoonright M_0)$.*

Proof. Apply downward Löwenheim-Skolem to (M, dp) to find $(M_0, dq) \preceq (M, dp)$ with $|M_0| \leq |T|$. Then $q = p \upharpoonright M_0$ and $p \sqsupseteq q$. \square

Definition 13. If $M \preceq N$ and $p \in S_n(M)$ and $q \in S_n(N)$, then q is a *strong heir* of p if $(N, dq) \succeq (M, dp)$.

Proposition 14 (Types have heirs). *Suppose $M \preceq N$ and $p \in S_n(M)$.*

1. *There is $N' \succeq N$ and $q' \in S_n(N')$ a strong heir of p .*
2. *There is $q \in S_n(N)$ an heir of p .*

Proof. 1. Let \bar{c} be an infinite tuple enumerating N . Then $\text{tp}^L(\bar{c}/M)$ is finitely satisfiable in M , hence finitely satisfiable in the expansion (M, dp) . Therefore it is satisfied in some $(N', dq) \succeq (M, dp)$. So there is \bar{e} in N' with $\text{tp}^L(\bar{e}/M) = \text{tp}^L(\bar{c}/M)$. Then the map $f(c_i) = e_i$ is an L -elementary embedding of N into N' extending $\text{id}_M : M \rightarrow M$. Moving N' by an isomorphism, we may assume $N' \succeq N$.

2. Take $N' \succeq N$ and $q' \in S_n(N')$ a strong heir of p . Let $q = q' \upharpoonright N$. Then $q' \supseteq q \supseteq p$ and $q' \sqsupseteq p$, so $q \sqsupseteq p$. \square

Next week, we will see that every heir of p arises this way, as a restriction of a strong heir of p .

3 Heirs and definable types

Proposition 15. *Let $p \in S_n(M)$ be definable and $N \succeq M$.*

1. *p has a unique heir $q \in S_n(N)$.*
2. *For $\varphi(\bar{x}, \bar{y})$ and $\bar{b} \in N$,*

$$\varphi(\bar{x}, \bar{b}) \in q(\bar{x}) \iff N \models d_p \varphi(\bar{b}) \quad (*)$$

3. *In particular, q is definable with $d_q \varphi = d_p \varphi$ for all φ .*

Proof.

Claim. If $q \in S_n(N)$ and $q \sqsupseteq p$, then q satisfies $(*)$.

Proof. Take $\bar{a} \in N' \succeq N$ realizing q (and p). If $(*)$ fails then

$$\begin{aligned} (\varphi(\bar{x}, \bar{b}) \in q(\bar{x}) &\not\leftrightarrow N \models d_p \varphi(\bar{b})) \\ N' \models \neg(\varphi(\bar{a}, \bar{b}) &\leftrightarrow d_p \varphi(\bar{b})) \\ \neg(\varphi(\bar{x}, \bar{b}) &\leftrightarrow d_p \varphi(\bar{b})) \in q(\bar{x}). \end{aligned}$$

As $q \sqsupseteq p$, there is $b' \in M$ such that

$$\begin{aligned} \neg(\varphi(\bar{x}, \bar{b}') &\leftrightarrow d_p \varphi(\bar{b}')) \in p(\bar{x}) \\ N' \models \neg(\varphi(\bar{a}, \bar{b}') &\leftrightarrow d_p \varphi(\bar{b}')) \\ (\varphi(\bar{x}, \bar{b}') \in p(\bar{x}) &\not\leftrightarrow M \models d_p \varphi(\bar{b}')), \end{aligned}$$

contradicting the definition of $d_p \varphi$. □_{Claim}

There is at least one heir, and at most one $q \in S_n(N)$ satisfying $(*)$, so therefore there is a unique heir q and it satisfies $(*)$. Point (3) is clear. □

Example. Suppose T is strongly minimal and $M \preceq N$ are models of T . Let p and q be the transcendental 1-types over M and N . For any $\varphi(x, \bar{y})$,

$$d_p \varphi(\bar{y}) \equiv (\exists^\infty x \varphi(x, \bar{y})) \equiv d_q \varphi(\bar{y}),$$

so q is the unique heir of p .

Proposition 16. *The following are equivalent for $p \in S_n(M)$.*

1. p is definable.
2. For every $N \succeq M$, p has a unique heir over N .

Proof. (1) \implies (2) is above. Conversely, suppose p has unique heirs. Then for any $N \succeq M$, p has at most one strong heir over N . Therefore, there is at most one way to expand N to an elementary extension of (M, dp) . Then the elementary diagram (M, dp) explicitly defines the relations $d\varphi$. By Beth's implicit definability theorem, (M, dp) is an expansion of M by definable relations, meaning p is definable. □

Proposition 17. *Suppose $M_1 \preceq M_2 \preceq M_3$, and $p_i \in S_n(M_i)$ for $i = 1, 2, 3$, with $p_1 \subseteq p_2 \subseteq p_3$. Suppose p_1 is definable. Then $p_1 \sqsubseteq p_2 \sqsubseteq p_3$ if and only if $p_1 \sqsubseteq p_3$.*

Proof. We only need to show the implication $p_1 \sqsubseteq p_3 \implies p_2 \sqsubseteq p_3$. Suppose $p_1 \sqsubseteq p_3$. Take $p'_2 \sqsupseteq p_1$ and $p'_3 \sqsupseteq p'_2$. By the uniqueness of heirs of definable types, $p'_2 = p_2$, and p_2 is definable. Then $p'_3 = p_3$, and $p_3 = p'_3 \sqsupseteq p'_2 = p_2$. □

4 Types in ACF

Recall ACF is the theory of algebraically closed fields like \mathbb{C} or \mathbb{Q}^{alg} . If $K \models \text{ACF}$, then $K[x_1, \dots, x_n]$ denotes the set of polynomials $P(x_1, \dots, x_n)$ in n variables.

A *positive quantifier free formula* is a quantifier-free formula that doesn't use \neg .

Fix a model $M \models \text{ACF}$.

Definition 18. A set $V \subseteq M^n$ is an *algebraic set* if

$$V = \varphi(M^n; \bar{b}) = \{\bar{a} \in M^n : M \models \varphi(\bar{a}, \bar{b})\}$$

where φ is positive quantifier free. In other words, an algebraic set is a set defined by a positive quantifier free $L(M)$ -formula.

Remark 19. V is an algebraic set iff V is defined by finitely many polynomial equations

$$V = \{\bar{a} \in M^n : P_1(\bar{a}) = \dots = P_m(\bar{a}) = 0\}.$$

Indeed, such a set is defined by the positive quantifier-free formula $\bigwedge_{i=1}^m (P_i(\bar{x}) = 0)$. Conversely, any positive quantifier-free formula can be put into conjunctive normal form

$$\bigwedge_{i=1}^m \bigvee_{j=1}^{\ell_i} (P_j(\bar{x}) = R_j(x))$$

which is equivalent to

$$\bigwedge_{i=1}^m \left(\underbrace{\prod_{j=1}^{\ell_i} (P_j(\bar{x}) - R_j(\bar{x}))}_{\text{a polynomial}} = 0 \right).$$

Example (Twisted cubic). $\{(t, t^2, t^3) : t \in M\}$ is an algebraic set in M^3 . It is defined by the system of equations

$$\begin{aligned} y &= x^2 \\ z &= x^3 \end{aligned}$$

i.e., by the two polynomials $x^2 - y$ and $x^3 - z$.

Lemma 20. 1. M^n and \emptyset are algebraic sets.

2. If $V, W \subseteq M^n$ are algebraic sets, then $V \cap W$ and $V \cup W$ are algebraic sets.

3. Any finite subset of M^n is an algebraic set.

Proof. (1) and (2) are easy, and then (3) reduces to showing that a point $\{\bar{a}\}$ is an algebraic set, which is easy, as it's defined by the system of polynomial equations $\{x_i = a_i : 1 \leq i \leq n\}$. \square

Fact (Quantifier elimination). *Every definable set $D \subseteq M^n$ is a finite boolean combination of algebraic sets.*

Fact (Consequence of Hilbert's basis theorem). *The class of algebraic sets has the descending chain condition (DCC): there is no infinite descending chain of algebraic sets $V_0 \supsetneq V_1 \supsetneq V_2 \supsetneq V_3 \supsetneq \dots$.*

Proof sketch (if you know ring theory). For V an algebraic set, let $I(V)$ be the set of polynomials in $K[\bar{x}]$ which vanish on V . Note that $I(V)$ is an ideal in $K[\bar{x}]$.

Claim. If V, W are algebraic sets and $V \supsetneq W$, then $I(V) \subsetneq I(W)$.

Proof. $I(V) \subseteq I(W)$ is easy. Take $\bar{a} \in V \setminus W$. Then $\bar{a} \notin W$ implies that $P(\bar{a}) \neq 0$ for one of the polynomials P defining W . Then P vanishes on W but not V , so $P \in I(W) \setminus I(V)$, and $I(W) \supsetneq I(V)$. □_{Claim}

Then a descending chain of algebraic sets would give an ascending chain of ideals in $K[\bar{x}]$, contradicting Hilbert's basis theorem, which says $K[\bar{x}]$ is Noetherian. □

Remark. In what follows, ACF can be replaced by any theory T with quantifier elimination, such that the class of sets defined by positive quantifier-free formulas satisfies the descending chain condition. One other theory with this property is the theory DCF of *differentially closed fields*, discussed in Section 6.2 of the textbook.

Corollary 21. *If \mathcal{S} is a non-empty collection of algebraic sets, then \mathcal{S} contains at least one minimal element.*

Corollary 22. *An infinite intersection $\bigcap_{i \in I} V_i$ of algebraic sets is an algebraic set.*

Proof. Apply Corollary 21 to the set of finite intersections $\{\bigcap_{i \in I_0} V_i : I_0 \subseteq_f I\}$. (Here, $A \subseteq_f B$ means A is a finite subset of B .) □

Corollary 23. *If $S \subseteq K[\bar{x}]$ is any set of polynomials, possibly infinite, then the subset of M^n defined by S is an algebraic set. All algebraic sets arise this way. (This is the usual definition of “algebraic set.”)*

Corollary 24 (Noetherian induction). *Let \mathcal{S} be a class of algebraic sets. Suppose the following holds:*

If X is an algebraic set, and every algebraic set $Y \subsetneq X$ is in \mathcal{S} , then $X \in \mathcal{S}$.

Then every algebraic set is in \mathcal{S} .

Proof. Otherwise, take a minimal algebraic set not in \mathcal{S} and get a contradiction. □

Definition 25. An algebraic set V is *reducible* if $V = W_1 \cup W_2$ for algebraic sets $W_1, W_2 \subsetneq V$. A *variety* is a non-empty irreducible algebraic set.

Remark 26. If V is an algebraic variety, then the set of algebraic proper subsets of V is closed under finite unions.

Proposition 27. *If V is an algebraic set, then V is a finite union of varieties.*

Proof. • $V = \emptyset$: V is a union of zero varieties.

- V is irreducible: V is a union of one variety (itself).
- V is reducible: $V = X \cup Y$ where $X, Y \subsetneq V$. By Noetherian induction, X and Y are finite unions of varieties, and so V is too. \square

Example. In M^1 , the algebraic sets are

- Finite sets
- M .

Thus, the varieties are

- Points
- M .

Fact. In M^2 , the varieties are exactly the following:

- *Points*
- M^2 .
- $\{(x, y) \in M^2 : P(x, y) = 0\}$ where $P(x, y)$ is a prime polynomial in $M[x, y]$.

In M^3 , one gets similar things, plus an additional class of “1-dimensional varieties,” one of which is the twisted cubic.

Let $V \subseteq M^n$ be a variety.

Definition 28. The *generic type* of V is the type generated by the following formulas:

- $x \in V$.
- $x \notin W$ for each algebraic proper subset $W \subsetneq V$.

We will write this type as $p_V(\bar{x})$.

Proposition 29. *Let V be a variety.*

1. $p_V(\bar{x})$ is a consistent complete type.
2. If W is an algebraic set, then $p_V(\bar{x}) \vdash \bar{x} \in W$ if and only if $W \supseteq V$.

Proof. Finite satisfiability: given finitely many proper algebraic subsets $W_1, \dots, W_m \subsetneq V$, we have $V \supsetneq \bigcup_{i=1}^m W_i$ by Remark 26, so there is $\bar{a} \in M^n$ with $\bar{a} \in V$ and $\bar{a} \notin W_i$ for $1 \leq i \leq m$.

(2): If $W \supseteq V$, then $p_V(\bar{x}) \vdash \bar{x} \in V \vdash \bar{x} \in W$. If $W \not\supseteq V$, then $(W \cap V) \subsetneq V$, so $p_V(\bar{x}) \vdash \bar{x} \notin W \cap V$. But $p_V(\bar{x}) \vdash \bar{x} \in V$, so $p_V(\bar{x}) \vdash \bar{x} \notin W$.

Completeness: by (2), for any positive quantifier-free formula $\varphi(\bar{x})$,

$$p_V(\bar{x}) \vdash \varphi(\bar{x}) \text{ or } p_V(\bar{x}) \vdash \neg\varphi(\bar{x}).$$

Then the same holds for quantifier-free formulas (which are boolean combinations of positive quantifier-free formulas), and then the same holds for all formulas (which are equivalent to quantifier-free formulas). \square

Example. The generic type of the variety M^1 is the transcendental type of Proposition 3.

Theorem 30. *The map $V \mapsto p_V$ is a bijection from the set of varieties $V \subseteq M^n$ to $S_n(M)$.*

Proof. Injectivity: suppose V, W are varieties and $V \neq W$. Without loss of generality, $V \not\subseteq W$. Then $p_W(\bar{x}) \vdash \bar{x} \in W$, but $p_V(\bar{x}) \not\vdash \bar{x} \in W$, so $p_W \neq p_V$.

Surjectivity: fix $p \in S_n(M)$. Take V a minimal algebraic set such that $p(\bar{x}) \vdash \bar{x} \in V$. (There's at least one such V , namely $V = M^n$.) V is non-empty because p is consistent. If V is reducible as $V = X \cup Y$ for smaller algebraic sets X, Y , then $p(\bar{x}) \vdash \bar{x} \in X$ or $p(\bar{x}) \vdash \bar{x} \in Y$ by completeness, contradicting the choice of V . Thus V is a variety. By choice of V , $p(\bar{x}) \vdash \bar{x} \in V$. If $W \subsetneq V$, then $p(\bar{x}) \not\vdash \bar{x} \in W$ (by choice of V), so $p(\bar{x}) \vdash \bar{x} \notin W$ (by completeness of p). Then $p(\bar{x}) \vdash p_V(\bar{x})$, so $p = p_V$, as both are complete types. \square

Proposition 31. *Let $V \subseteq M^n$ be a variety, defined by a formula φ .*

1. φ defines a variety $V_N \subseteq N^n$.
2. V_N depends only on V , not on the choice of φ .

Proof. Take ψ a positive quantifier-free formula defining V . Then $\forall \bar{x} (\varphi(\bar{x}) \leftrightarrow \psi(\bar{x}))$ is satisfied by M , and therefore by N . Let $V_N = \varphi(N^n) = \psi(N^n)$. This depends only on ψ , not on φ , so the choice of φ doesn't matter. As ψ is positive quantifier free, V_N is an algebraic set. As $M \models \exists \bar{x} \varphi(\bar{x})$, the same holds in N , so V_N is non-empty. Suppose for the sake of contradiction that V_N is reducible. Then $V_N = W_1 \cup W_2$, where for $i = 1, 2$, W_i is an algebraic proper subset of V_N defined by $\theta_i(\bar{x}, \bar{b}_i)$ for some positive quantifier-free L -formula θ_i and tuple of parameters $\bar{b}_i \in N$. Then

$$N \models \exists \bar{y}_1, \bar{y}_2 \left(\forall \bar{x} \left(\psi(\bar{x}) \leftrightarrow \bigvee_{i=1}^2 \theta_i(\bar{x}, \bar{y}_i) \right) \wedge \bigwedge_{i=1}^2 \exists \bar{x} (\psi(\bar{x}) \wedge \neg \theta_i(\bar{x}, \bar{y}_i)) \right)$$

namely $(\bar{y}_1, \bar{y}_2) = (\bar{b}_1, \bar{b}_2)$. The same holds in $M \preceq N$, which means that $V = \bigcup_{i=1}^2 \theta_i(M^n, \bar{b}_i)$ for some $\bar{b}_i \in N$, contradicting irreducibility of V . \square

Theorem 32. Let $M \preceq N$ be models of ACF. Let $V \subseteq M^n$ be a variety, and let $V_N \subseteq N^n$ be its extension. Then $p_{V_N} \in S_n(N)$ is the unique heir of $p_V \in S_n(M)$.

Proof. Let $q \in S_n(N)$ be an heir of p_V . Let φ be an $L(M)$ -formula defining V and V_N . Then $\varphi(\bar{x}) \in p_V(\bar{x}) \subseteq q(\bar{x})$, so $q(\bar{x}) \vdash \bar{x} \in V_N$. Suppose $q(\bar{x}) \not\vdash \bar{x} \notin W$ for some algebraic proper subset $W \subsetneq V_N$. By completeness of q , $q(\bar{x}) \vdash \bar{x} \in W$. Let $\psi(\bar{x}, \bar{b})$ be a positive quantifier-free formula defining W . Let $\theta(\bar{b})$ be the $L(M)$ -formula

$$\forall \bar{x} (\psi(\bar{x}, \bar{b}) \rightarrow \varphi(\bar{x})) \wedge \exists \bar{x} (\varphi(\bar{x}) \wedge \neg \psi(\bar{x}, \bar{b})).$$

(This says $\psi(M^n, \bar{b}) \subsetneq \varphi(M^n)$.) In particular, $N \models \theta(\bar{b})$, because $W \subsetneq V$. Then $q(\bar{x}) \vdash \psi(\bar{x}, \bar{b}) \wedge \theta(\bar{b})$, because $\theta(\bar{b})$ is true in N . Because $q \supseteq p_V$, there is $\bar{b}' \in M$ such that

$$p_V(\bar{x}) \vdash \psi(\bar{x}, \bar{b}') \wedge \theta(\bar{b}').$$

In particular, $M \models \theta(\bar{b}')$ (or else $p_V(\bar{x})$ is inconsistent). If $W_0 = \psi(\bar{x}, \bar{b}')$, then W_0 is an algebraic proper subset of V , as $M \models \theta(\bar{b}')$. Then $p_V(\bar{x}) \vdash \psi(\bar{x}, \bar{b}')$ means $p_V(\bar{x}) \vdash \bar{x} \in W_0$, contradicting the definition of p_V .

So we see that $q(\bar{x}) \vdash \bar{x} \in V_N$, and $q(\bar{x}) \vdash \bar{x} \notin W$ for any proper algebraic subset $W \subsetneq V_N$. Thus $q = p_{V_N}$, and p_{V_N} is the only heir of p . \square

Corollary 33. In ACF, all types over models are definable. ACF is stable.

Motivation. On some level, all stable theories look like ACF, especially the strongly minimal and uncountably categorical theories. We will see examples of this throughout the course.

5 1-types in DLO

DLO is the theory of *dense linear orders* without endpoints, things like (\mathbb{R}, \leq) and (\mathbb{Q}, \leq) . It is complete and has quantifier elimination.

Suppose $M \models \text{DLO}$ and $p(x) \in S_1(M)$. By quantifier elimination, p is determined by the two sets:

$$\begin{aligned} X^- &= \{b \in M : p(x) \vdash b \leq x\} \\ X^+ &= \{b \in M : p(x) \vdash b \geq x\}. \end{aligned}$$

There are six possibilities:

X^-	X^+	Explanation	Example
$(-\infty, b]$	$[b, +\infty)$	Constant type at b	$\text{tp}(2/\mathbb{Q})$
$(-\infty, b]$	$(b, +\infty)$	Type at b^+	$\text{tp}((2 + \epsilon)/\mathbb{Q})$
$(-\infty, b)$	$[b, +\infty)$	Type at b^-	$\text{tp}((2 - \epsilon)/\mathbb{Q})$
M	\emptyset	Type at $+\infty$	$\text{tp}(10/(-1, 1))$
\emptyset	M	Type at $-\infty$	$\text{tp}(-10/(-1, 1))$
(has no maximum)	(has no minimum)	Irrational type	$\text{tp}(\pi/\mathbb{Q})$

Lemma 34. *Suppose $M \preceq N \models \text{DLO}$ and $p \in S_1(M)$ is the type at b^+ . If $q \in S_1(N)$ is an heir, then q is the type at b^+ .*

Proof. $p(x) \vdash x > b$, so $q(x) \vdash x > b$. If q isn't the type at b^+ , then there is $c > b$ in N such that $q(x) \vdash x > c$. Then $q(x) \vdash x > c \wedge c > b$. Because $q \sqsupseteq p$, there is $c' \in M$ such that $p(x) \vdash x > c' \wedge c' > b$. Then $c' > b$ (as p is consistent), and then c' shows p is *not* the type at b^+ , a contradiction. \square

So if $p \in S_1(M)$ is the type at b^+ , then p has a unique heir over any $N \succeq M$. Therefore p is definable.

Similarly...

1. The unique heir of the type at b^- is the type at b^- .
2. The unique heir of the type at $\pm\infty$ is the type at $\pm\infty$.
3. The unique heir (*or extension*) of the constant type at b is the constant type at b .

Then all these types are definable. On the other hand...

Fact (see pages 228–229 the textbook). *If $p \in S_1(M)$ is an irrational type and $N \succeq M$, then every non-constant extension of p to N is an heir of p . There is more than one such extension.*

Consequently, irrational types are not definable. Here is an example:

Example. $\text{tp}(\pi/\mathbb{Q})$ is not definable because $\{b \in \mathbb{Q} : p(x) \vdash x < b\} = (\pi, +\infty) \cap \mathbb{Q}$, which is not a definable subset of (\mathbb{Q}, \leq) .