

Implementing Distribution Testing Algorithms

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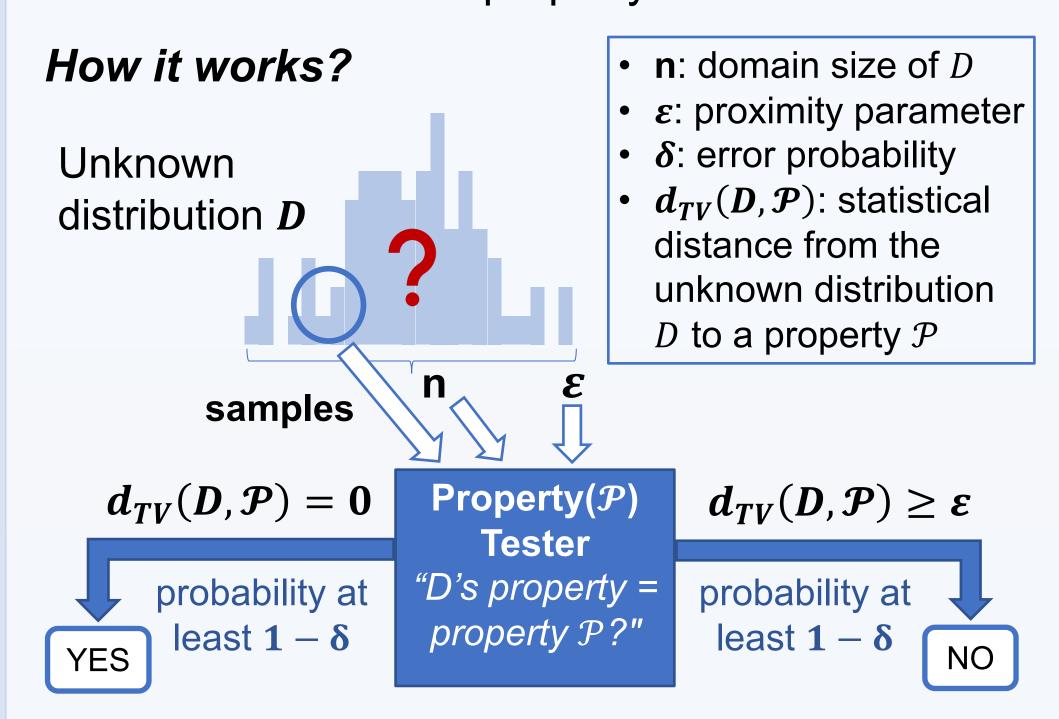
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MOTIVATION

- Sample-efficient distribution testing algorithms have yet to be implemented and utilized
- Specific number of samples required for these algorithms remain undetermined
- Theoretical claims of papers need to be verified

BACKGROUND

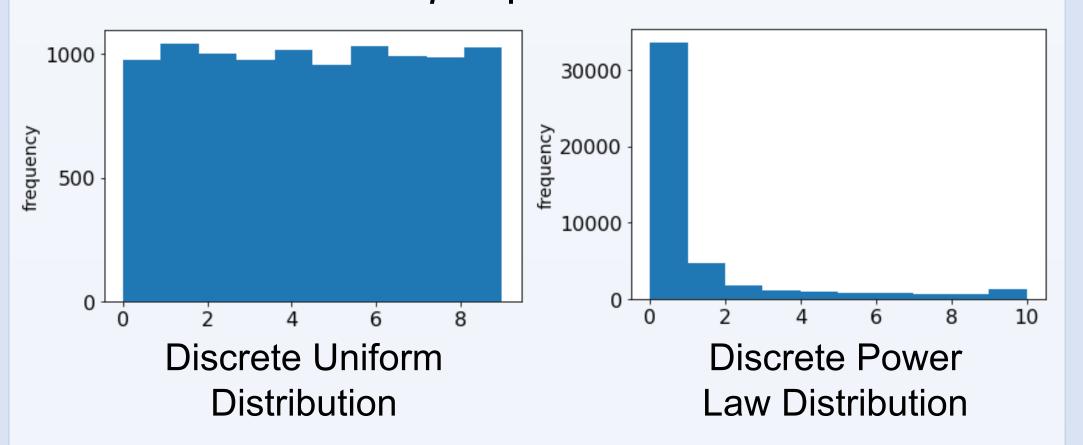
Distribution Testing Test if an unknown distribution has some property



Uniformity Testing Test if all discrete bins have the same weight

Identity Testing Test whether an unknown distribution matches a known distribution

EX. Is distribution p a power law distribution?



Potential Applications of Distribution Testing

- Check if the distribution of NY State Lotto's winning numbers is uniform
- Determine if the dataset to regress is linear by checking if residual is normally distributed

PROBLEM STATEMENT

 To provide working implementations of theoretically-efficient distribution testing algorithms, with determined sample-complexity and relative run-times

METHODS

Implemented Algorithms

Uniformity tester	Identity tester
I ₂ estimator	Chi-squared tester
Coincidence-based tester	instance-optimal tester
"New approach" uniformity	

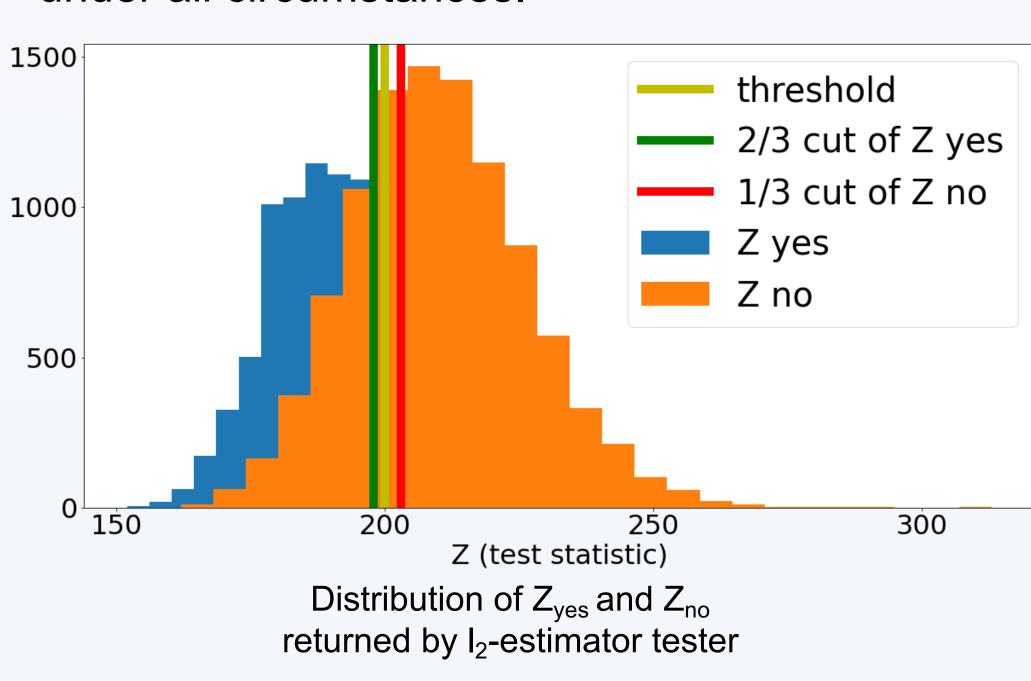
How to Pin Down Constants?

Algorithms compute test statistic Z on tested distribution and use Z to indicate if the distribution has the desired property.

STEP 1 Run the algorithms on both yes and no distributions for 10^6 times, to get z_{yes} and z_{no} .

STEP 2 Adjust number of samples to make the z_{yes} and z_{no} distribution overlap at percentile of $100 * \delta$ (ex. $\delta = \frac{1}{2}$, intersect at 33^{rd} percentile).

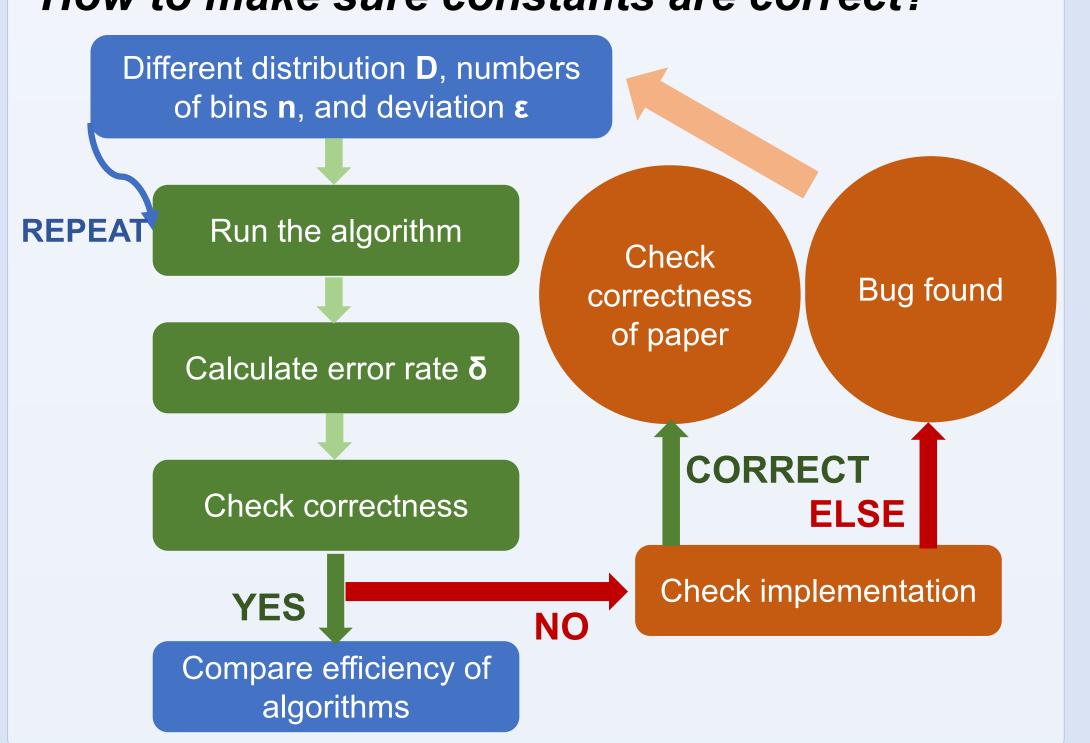
STEP 3 Adjust the constant of threshold, making the threshold match the "true" threshold under all circumstances.



Testing Distributions

Uniformity Yes Distribution	Uniform D. over N
Uniformity No Distribution	$(1\pm\epsilon)/n$ for all n bins
Identity Yes Distribution	$q_i = c_1(i+c_2)^{(-3/2)}$
Identity No Distribution	$p_i = q_i(1 \pm \varepsilon)$

How to make sure constants are correct?



RESULTS & COMPARISONS

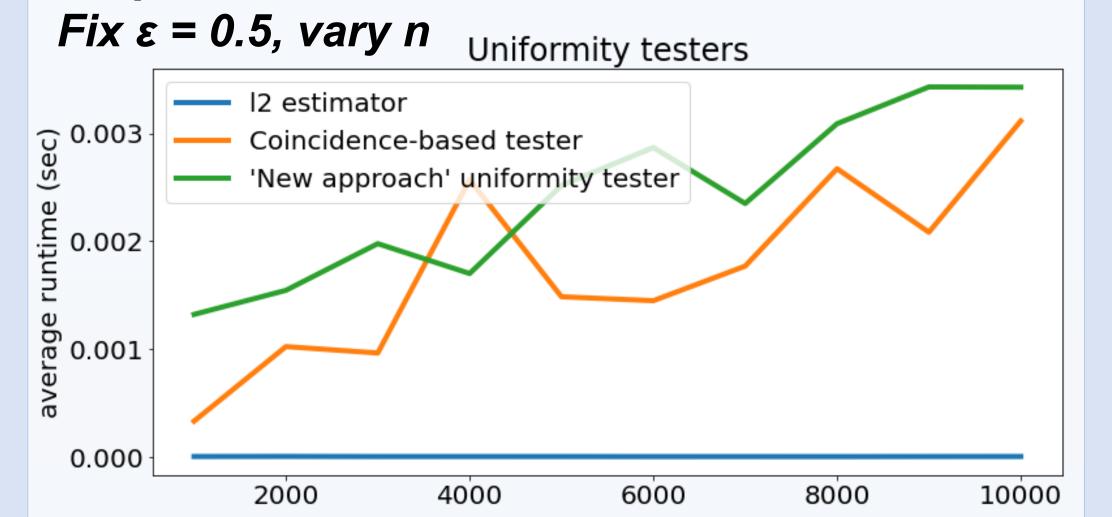
Worst-case Sample Size Pinned Down

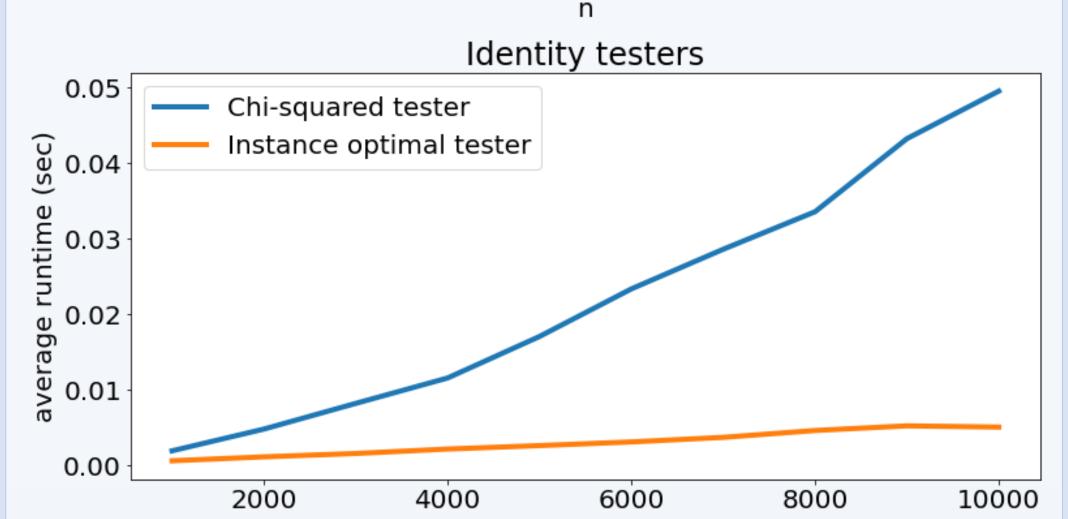
Uniformity tester	Sample size
I ₂ estimator	$1.76 * \sqrt{n}/\varepsilon^2$
Coincidence-based tester	$2.1 * \sqrt{n}/\varepsilon^2$
"New approach" uniformity tester	$5.9 * \sqrt{n}/\varepsilon^2$

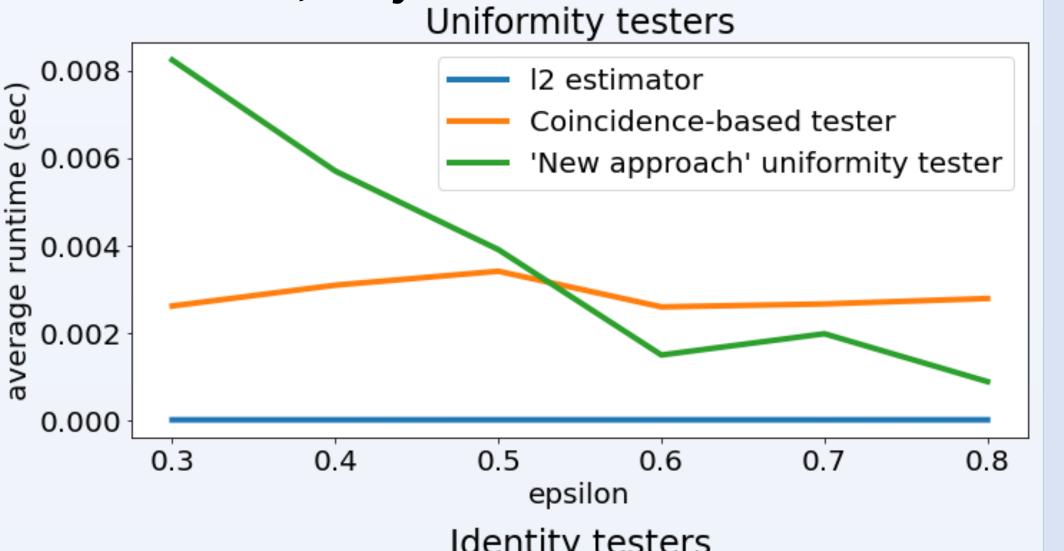
Identity tester	Sample size
Chi-squared tester	$1.45 * \sqrt{n}/\varepsilon^2$
Instance optimal tester	1.8 * max $(\frac{1}{\varepsilon}, p _{\frac{2}{\varepsilon}}/\varepsilon^2)$

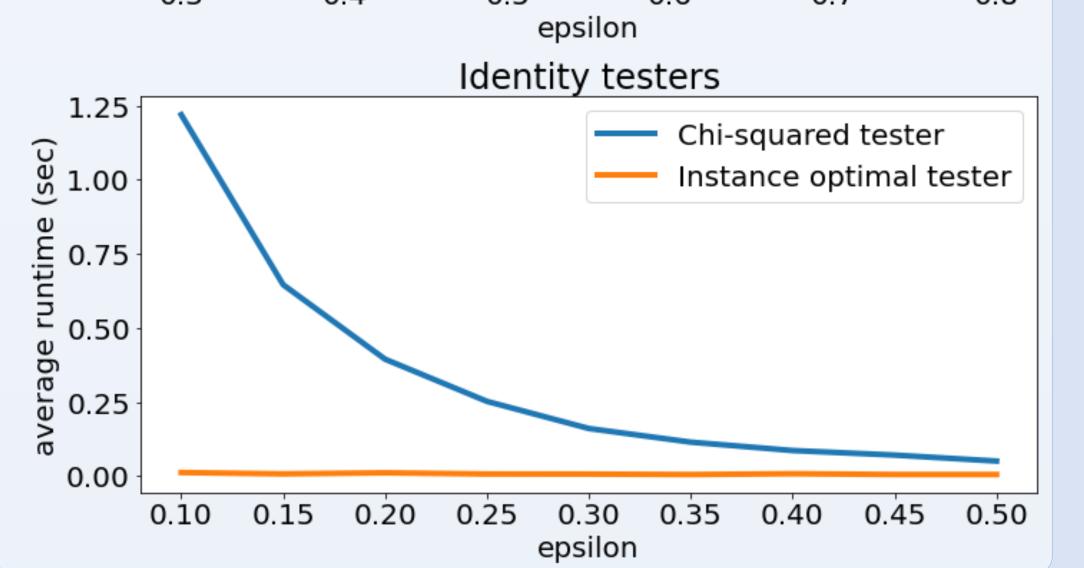
Graphs for Runtime

Fix n = 10000, vary ε









Sample Complexity Analysis

- The worst-case sample size required for each tester is proportional to \sqrt{n}/ε^2 (1.45 $\leq C \leq$ 5.9 for $C * \sqrt{n}/\varepsilon^2$)
- Our experimental results match the theoretical expectations.
- l₂ estimator is the most sample-efficient uniformity tester.
- Chi-squared tester is more sample-efficient when known distribution(q) is uniform; less efficient if q is far from uniform.

Runtime Analysis

Ran on MacBook Air
Processor: 1.8 GHz Dual-core Intel Core i5
Memory: 8 GB 1600 MHz DDR3

- All algorithms are fast enough to run on even general-purpose computers (~1 second).
- Outcome matches theoretical expectations.

Uniformity tester	
I ₂ estimator	Fastest
Coincidence-based tester	Constant with all $arepsilon$
"New approach" uniformity tester	Faster as ε increases

Identity tester	
Chi-squared tester	Fast with larger ε
Instance optimal tester	Faster

CONCLUSION

GitHub Repo for Implementations:



FUTURE WORK

- Implementation of closeness and independence testers
- Research paper to compile and share findings

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