## Math 578 HW#4

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**Problem 1.** The entirety of this code is contained in the included single file hw4.py. In the main executing loop (if \_\_name\_\_ == "\_\_main\_\_": ), various sections have been commented out. Uncomment as desired and run.

(a) Cholesky decomposition is performed by the function hw4.cholesky. The particular implementation handles sparse matrices as well, and most of the code is simply dedicated to choosing the right method based on the data object provided:

```
def cholesky(A):
2
            computes the cholesky decomposition for symmetric, positive definite
            matrices. returns a lower-triangular matril L with positive diagonal
            entries so that A=LL^T.
            also returns an integer nzl that gives the number of
            nonzero entries in the Cholesky factor L.
            INPUT:
10
            A - a positive definite matrix nxn
                   (may be np.array or scipy.sparse.spmatrix)
13
            OUTPUT:
14
            L - the cholesky factor L s.t. A = LL^T
16
            nzl - number of nonzero entries in L i.e. where |L {ij}| > 0
17
18
            if sparse.issparse(A):
                G = sparse.tril(A)
                #a sparse matrix that still allows elementwise access
23
                G = G.tocsc()
24
            else:
                G = np.tril(A)
26
            n = A.shape[0]
2.8
29
            for k in range(n):
30
                G[k:,k] -= G[k:,:k] @ G[k,:k].T
                G[k:,k] /= np.sqrt(G[k,k])
            if sparse.issparse(G):
                nzl = G.count nonzero()
            else:
36
                nzl = np.count nonzero(G)
37
38
            return G, nzl
39
```

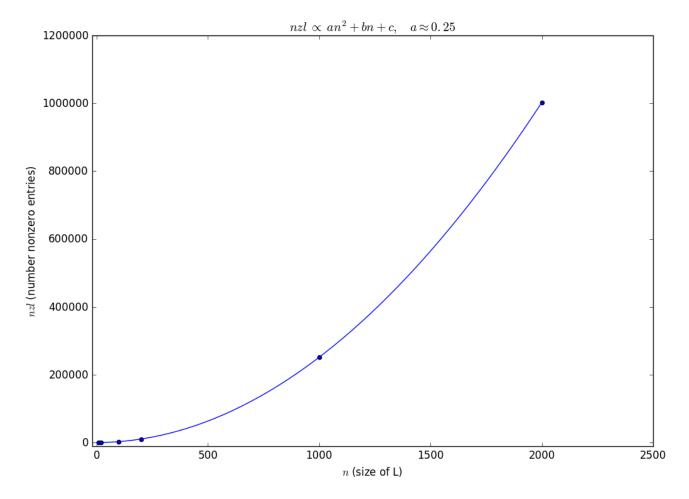


Figure 1: Relationship between size of system n and number of nonzero elements in Cholesky factor. The relationship was found to be quadratic with leading coefficient  $a \approx 0.25$ .

The following suggested sanity check was performed, which yielded the desired result:

- **(b)** Output of spy(L) for n = 2000:
- (c) To calculate the size n of system possible to solve with 8GB of memory, we simply use our coefficients we found in part (b) and solve for n, given that a 8GB of space can store 8GB  $\div$  8 bytes  $\rightarrow$  1,000,000 nonzero entries. Thus we calculate:

$$an^2 + bn + c = 10^9 \rightarrow \frac{1}{4}n^2 \approx 10^9 \rightarrow n \approx 63245$$

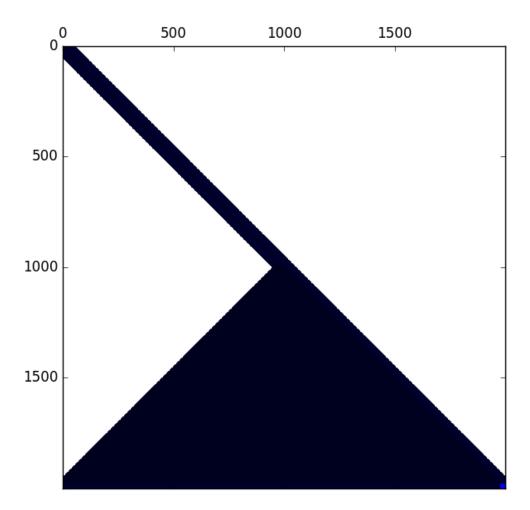


Figure 2: (for part (b)) A visualization of nonzero elements of the cholesky factor L of the matrix A when n=1000

**(d)** The multiply function is as follows:

```
def multiply(x, diagonal=None):
      implicitly performs the matrix multiplication Ax
      for the matrix given in (i)
      # test multiply function
      n = 60
      x = np.expand dims(np.random.randn(n), -1)
      A = build p1(n)
10
      np.all(np.isclose(A@x,multiply(x))) -> True
13
      if diagonal is None:
14
          d = 3
16
      else:
17
          d = diagonal
18
19
      y = np.empty_like(x)
      n = x.size - 1
21
      n2 = x.size // 2 - 1
      # just overwrite the different entries later (ignore 0,n)
24
      for i in range(1,n):
          y[i] = d*x[i] - x[i+1] - x[i-1] - x[n-i]
      y[0] = d*x[0] - x[1] - x[-1] # first entry
      y[-1] = d*x[-1] - x[-2] - x[0] # nth (last) entry
29
      # then overwrite the following:
      y[n2] = d*x[n2] - x[n2+1] - x[n2-1]
      y[n2+1] = d*x[n2+1] - x[n2+2] - x[n2]
34
      return y
```

(e) Conjugate gradient was used to solve the system Ax = b as described. The following iteration numbers and relative errors were reported:

```
n iterations relative error

10 3 1.13492816813e-15

50 13 7.26595797191e-12

100 26 2.12606603652e-08

5000 1544 6.25257025384e-07

15000 4616 7.0254248905e-07

30000 9216 9.29915260829e-07
```

Clearly, CG does not fare well as the size of the system increases. The conjugate gradient method was implemented as follows:

```
def cg(b=None, n=None, mult=None, tol=1e-6):
      non-preconditioned CG
      initial estimate is b (defaults to normalized ones vector)
      size of system n (can be inferred from b or vice versa)
      using multiplication method mult
      assert mult is not None
      if b is None:
13
          try:
               # normalize vector of ones
14
               b = np.ones((n,1)) / np.sqrt(n)
15
          except NameError:
16
               raise Exception('must specify system size or initial quess')
17
      else:
18
          n = b.size
19
20
      # make sure initial guess is a column vector ala matlab
21
      if b.ndim == 1:
          b = np.expand dims(b, -1)
23
24
      # not sure if explicit copy is needed
      x = np.zeros like(b)
26
      r = b.copy()
      p = b.copy()
28
2.9
      d = mult(p)
30
      alpha = np.vdot(r,r) / np.vdot(p,d)
32
      for iterations in count(1):
34
35
          x += alpha*p
36
          r new = r - alpha*d
38
          beta = np.vdot(r new,r new) / np.vdot(r,r)
39
40
          p = r new + beta*p
42
          d = mult(p)
43
44
          alpha = np.vdot(r new, r new) / np.vdot(p,d)
45
46
          r = r \text{ new}
47
48
          err = norm(mult(x) - b) / norm(b)
49
50
          #print(iterations, err, sep='\t| ')
          if err <= tol:</pre>
               break
54
      return x, iterations, err
```

(f) The following iterations and errors were reported for CG when the diagonal elements of the system were perturbed slightly  $(3 \rightarrow 3.0001)$ 

```
n iterations relative error

10 3 1.88574575416e-15

50 13 3.00488927641e-12

100 26 9.88035379143e-07

5000 1519 8.48965754023e-07

15000 1835 9.97779149178e-07

30000 1797 9.90061636977e-07
```

Comparing to the unperturbed system, CG converges faster here. The small perturbation has the effect of lowering the condition number of the matrix.

(g) Another multiply function was created for a the system  $A = I + BB^T$ , and used for conjugate gradient. Implementation as shown:

```
def multiply2(x):
     multiply implicitly by A = I + BB.T
     where B = np.tril(np.ones((n,3))) and n = x.size
      i.e. B = array([[1, 0, 0],
                      [ 1, 1,
                                0],
                      [ 1, 1, 1],
                      [ . .
                                .],
                      [ .
                                .],
                      [ .
                                .],
10
                      [ 1, 1, 1]])
11
     this faster method is shown by first considering (B.T @ x) itself,
13
     which yields the 3-vector:
      [S, S-x[0], S-x[0]-x[1]], where S=x.sum()
16
17
     S = x.sum()
18
19
     # do BB.T mult first
20
     BBx = np.zeros like(x)
     BBx[0] = S
     BBx[1] = 2*S - x[0]
      BBx[2:] = 3*S - 2*x[0] - x[1]
24
      # now add I part and return
26
      return x + BBx
```

Size of system / Iterations / Relative error as shown:

```
10000 4 2.06562011245e-15
50000 4 1.82194210335e-14
100000 4 1.15615944448e-13
```

It's clear that A has at most 4 distinct eigenvalues. Clearly, the matrix  $BB^T$  has rank 3 and thus at most 3 distinct eigenvalues. Thus  $A = I + BB^T$  has these three plus 0 + 1, so 4.