

Advanced Data Bases

Course II - SQL 3

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part I

-

Introduction

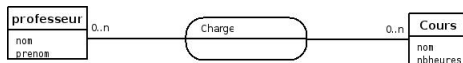
Design, Development, Use, Administration

- ① Conceptual step: Design and Database Modeling
- ② logical step: Establishment of a database
- ③ Physical Stage: Software (DBMS, interfaces, ...) & equipment

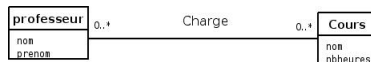
Database Modeling

- Analysis phase: definition of a conceptual schema
- Conceptual diagram of data (SCD): according to the formalism,
set of entities and Associations
or
set of Classes

- Formalism EA, ER



- UML formalism:



Database Modeling

Different modeling formalism of conceptual schemes of DB:

- Formalism EA, ER, EER
 - | Model Entity-Relationship (Entity-Relationship Model)
 - | Model Entity-Relationship Extended (Extended Entity-Relationship Model)
- Formalism UML (Unified Modeling Language)

Database Modeling

Both formalisms E / R and UML are very close / "equivalent"

Entity / Association

→

UML

Entity

Object

Type of entity

Class

Relationship

Object

Type of association

Class

Attribute / Property

Property

Role / Label

Role

Method

Field

domain constraint

Key

Key Constraint

Constraint

Constraint

cardinality

Multiplicity / cardinality

0.1 1.1 0, n 1, n, b, a

1 0..1 0 .. * 1 .. * a..ba

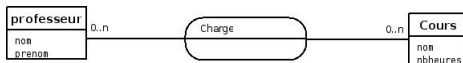
Diagram E / A

UML Class Diagram

Entity-Relationship Model

Recall

- Entity: Any concrete or abstract concept that can be individualized
- Class or type of entities: entities grouping similar (generic level)
- Association relation between multiple entities
- Class or type of association: group of associations with the same characteristics



Extended entity-relationship model

- Entity-relationship model: reduced concept but sufficient for modeling simple problems (or less complex)
- Extended Entity-relationship Model: more precise and expressive modeling of complex problems.

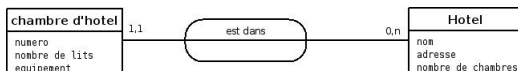
Introduction of abstraction mechanisms:

- | Weak types
- | classification
- | aggregation
- | inheritance
- | ...

Weak Types

Type of entities or weak associations:

- existence of a body subject to the existence of another type of entity or association



Here the identifier "number" of a single room is unique only for a given hotel!

Classification

- Grouping entities into classes based on common properties
- Possibility of classifying an object in several classes

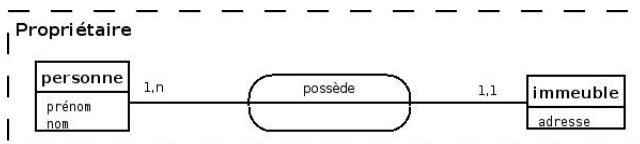
example:

- eBook: electronic file, and book
- Bus: Public transport vehicle, motor vehicle explosion

Aggregation

Description of complex entities' types.

One type of associations between entities' types is considered as a new type of entities

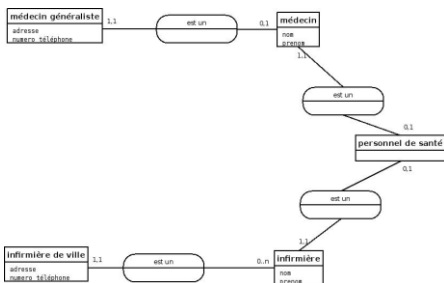


Inheritance

Specialization - Generalization

One type of entity A is a specialization of another type of entity B if

- each entity A is a B entity
- One entity (at most) of B is associated with a unit of A



Translation of functional constraints of integrity

Constraints: Score, Exclusion, Totality, Concurrency, Inclusion ...

All constraints can be set or programmed via:

- the declaration of constraints (constraints)
- programming
 - | functions (functions)
 - | procedures (procedures)
 - | packages (packages)
 - | triggers (triggers)

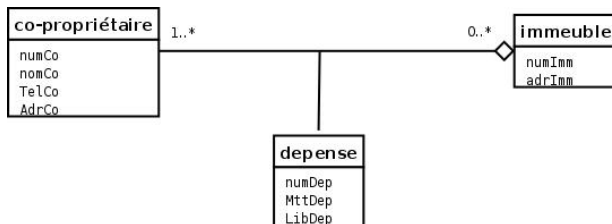
PL / SQL or with a host language such as C, C ++, Java

part II

-

Abstraction in SQL2

Translation aggregation relationships



Each building "includes" a number of owners.
Each building has at least 1 owner!

Translation aggregation relationships

REM ____ Un co-propriétaire peut posséder plusieurs immeubles

create table COPROPRIETAIRE

```
(  
    NUMCO number ( 7 ) ,  
    NOMCO varchar ( 10 ) ,  
    TELCO varchar ( 15 ) ,  
    ADRCO varchar ( 50 ) ,  
    constraint PK_COPROPRIETAIRE primary key (NUMCO)  
);
```

REM ____ Un immeuble doit être possédé par un ou

REM plusieurs copropriétaires

create table IMMEUBLE

```
(  
    NUMIMM number ( 7 ) ,  
    ADRIMM varchar ( 50 ) ,  
    constraint PK_IMMEUBLE primary key (NUMIMM)  
);
```


Translation aggregation associations

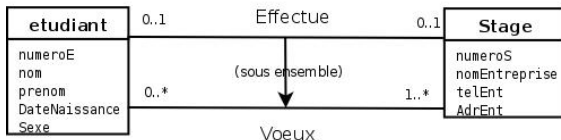
```
create table DEPENSE
(
    NUMCO number ( 7 ) ,
    NUMIMM number ( 7 ) ,
    DATEDEP date ,
    MTTDEP number ( 10 , 2 ) ,
    LIBDEP varchar ( 50 ) ,
    constraint PK_DEPENSE primary key (NUMCO,NUMIMM) ,
    constraint FK_DEPENSE_NUMCO_COPROPR foreign key (NUMCO)
        references COPROPRIETAIRE(NUMCO) on delete cascade ,
    constraint FK_DEPENSE_NUMIMM_IMMEUBLE foreign key (NUMIMM)
        references IMMEUBLE(NUMIMM ) on delete cascade
);
```

Translation aggregation associations

```
create table DEPENSE
(
    NUMCO number ( 7 ) ,
    NUMIMM number ( 7 ) ,
    DATEDEP date ,
    MTTDEP number ( 10 , 2 ) ,
    LIBDEP varchar ( 50 ) ,
    constraint PK_DEPENSE primary key (NUMCO,NUMIMM) ,
    constraint FK_DEPENSE_NUMCO_COPROPR foreign key (NUMCO)
        references COPROPRIETAIRE(NUMCO) on delete cascade ,
    constraint FK_DEPENSE_NUMIMM_IMMEUBLE foreign key (NUMIMM)
        references IMMEUBLE(NUMIMM) on delete cascade
);
```

The minimum cardinality of the association will be tested through a PL / SQL procedure

Inclusion constraint



```
create table STAGE
(
  NUMEROS number (7),
  NOMENTREPRISE varchar (40),
  TELENT varchar (15),
  ADRENT varchar (50),
  constraint PK_STAGE primary key (NUMEROS)
);
```

Inclusion constraint

Table per class

Create table ETUDIANT

```
(  
    NUMEROE number ( 7 ) ,  
    NOM varchar ( 10 ) ,  
    PRENOM varchar ( 10 ) ,  
    DATENAISSANCE date ,  
    SEXE char ( 1 ) ,  
    NUMEROS number ( 7 ) ,  
    constraint PK_ETUDIANT primary key (NUMEROE) ,  
    constraint t FK_ETUDIANT_NUMEROS_STAGE foreign key (NUMEROS)  
        references STAGE(NUMEROS),  
    constraint t CK_ETUDIANT_SEXE check (SEXE i n ( 'M' , 'F' ) )  
);
```

Inclusion constraint

```
create table VOEUX
(
  NUMEROE number (7),
  NUMEROS number (7),
  constraint PK_VOEUX primary key (NUMEROE,NUMEROS)
  constraint FK_VOEUX_NUMEROE_ETUDIANT foreign key (NUMEROE)
    references ETUDIANT (NUMEROE)
  constraint FK_WISHES_NUMEROS_STAGE foreign KEY (NUMEROS)
    references STAGE (NUMEROS)
);
```

inclusion Constraint?

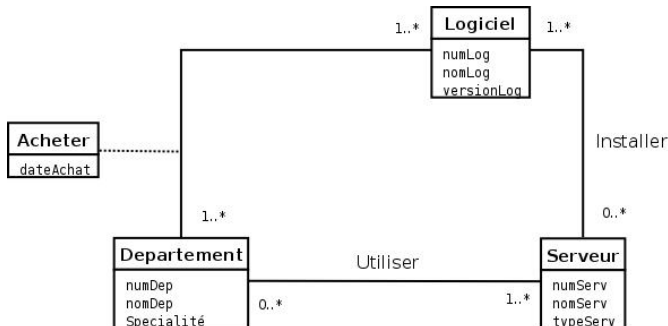
Inclusion constraint

```
create table VOEUX
(
  NUMEROE number (7),
  NUMEROS number (7),
  constraint PK_VOEUX primary key (NUMEROE,NUMEROS)
  constraint FK_VOEUX_NUMEROE_ETUDIANT foreign key (NUMEROE)
    references ETUDIANT (NUMEROE)
  constraint FK_WISHES_NUMEROS_STAGE foreign KEY (NUMEROS)
    references STAGE (NUMEROS)
);
```

```
Alter table ETUDIANT add
  constraint K_EFFECTUER_INCLUSION_VOEUX
  foreign key (NUMEROE,NUMEROS)
  references VOEUX (NUMEROE,NUMEROS);
```

Inclusion constraint

Inclusion constraint: The software must be installed on a server of the department that purchased the program.



The software L purchased by D department is installed on a server S, for among other things, this department

Inclusion constraint

Table per Class

```
create      table  DEPARTMENT
(
  NUMDEP number (7),
  NOMDEP varchar (1 0),
  SPECIALTE varchar (2 0),
  constraint PK_DEPARTEMENT primary key          (NUMDEP)
);
```

```
create      table  LOGICIEL
(
  NUMLOG number (7),
  NOMLOG varchar (1 0),
  VERSIONLOG varchar (1 0)      ,
  constraint PK_LOGICIEL          primary key    (NUMLOG)
);
```


Inclusion constraint

Table per Class

```
create   table SERVEUR
(
  NUMSERV number (7),
  NOMSERV varchar (10),
  TYPESERV varchar (10),
  constraint PK_SERVEUR primary key (NUMSERV)
);
```

Inclusion constraint

A table by Association or Class-Association

```
create   table ACHETER (  
    NUMDEP number (7), NUMLOG number (7), DATEACHAT date ,  
    constraint PK_ACHETER primary key (NUMDEP, NUMLOG),  
    constraint FK_ACHETER_NUMDEP_DEPARTEMENT  
        Foreign KEY (NUMDEP) references DEPARTMENT (NUMDEP)  
    constraint FK_ACHETER_NUMLOG_LOGICIEL  
        foreign KEY          (NUMLOG)    references LOGICIEL (NUMLOG)  
);
```

```
create table UTILISER(  
    NUMDEP number (7), NUMSERV number (7),  
    constraint PK_UTILISER primary key (NUMDEP, NUMSERV),  
    constraint FK_UTILISER_NUMDEP_DEPARTEMENT  
        foreign key (NUMDEP) references DEPARTMENT (NUMDEP),  
    constraint FK_UTILISER_NUMSERV_SERVEUR  
        foreign NKEY  (NUMSERV) references SERVER (NUMSERV)  
);
```

Inclusion constraint

A table by Association or Class-Association

```
Create Table INSTALLER (  
  NUMLOG number (7), NUMSERV number (7),  
  constraint PK_INSTALLER primary key (NUMLOG, NUMSERV)  
  constraint FK_INSTALLER_NUMLOG_LOGICIEL  
    foreign KEY(NUMLOG)      references      LOGICIEL (NUMLOG)  
  constraint FK_INSTALLER_NUMSERV_SERVEUR  
    foreign KEY(NUMSERV) references SERVEUR (NUMSERV)      );
```

Inclusion Constraint

trigger

The software purchased by D department is installed on a server S dedicated to this department

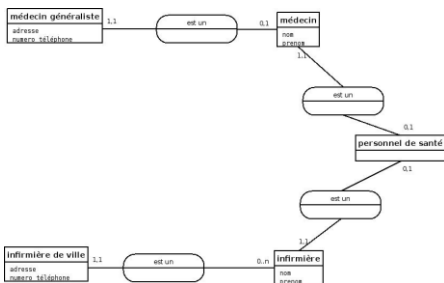
```
Create or replace trigger trig_contrainte_inclusion
before insert on INSTALLER
For each row
declared
LOGIC number (7);
SERV number (7);
begin
    select ACHETER.NUMLOG,    UTILISER .NUMSERV into LOGIC SERV
from ACHETER, UTILISER
Where ACHETER.NUMDEP = UTILISER .NUMDEP and
        ACHETER.NUMLOG = : new .NUMLOG and
        UTILISER .NUMSERV = : new .NUMSERV;
exception
When NO_DATA_FOUND    Then
    raise_application_error (-20,100,
        Le logiciel doit être installé sur
        un serveur du département acheteur');
end ;
/
```

Inheritance

Specialization - Generalization

One type of entity A is a specialization of another type of entity B if

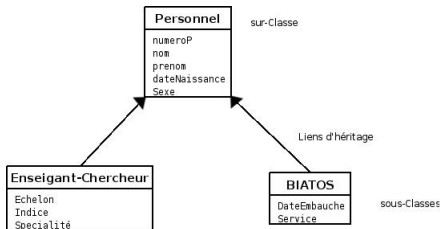
- each entity A is a B entity
- One entity (at most) of B is associated with a unit of A



Translation of inheritance associations

Translation of inheritance constraints

Personnel management in a university



Inheritance associations in UML (1)

- Presenting the different cases of inheritance based on instances
- Modeling of different inheritance in the UML formalism with constraints
 - | partition
 - | exclusion
 - | totality

Inheritance associations in UML (2)

Expression of inheritance cases using

- coverage
- disjunction

of instances in a given population

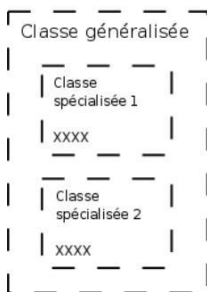
Four types of constraints are identified:

- partition
- totality
- exclusion
- the lack of constraints

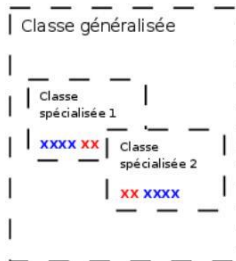
Inheritance constraints PARTITION and TOTALITY

- Disjunction & coverage → Partition
- Non-disjunction & Coverage → Totality

PARTITION



TOTALITE

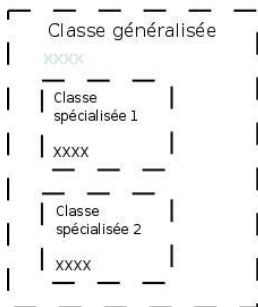


Inheritance Constraints EXCLUSION and NO CONSTRAINTS

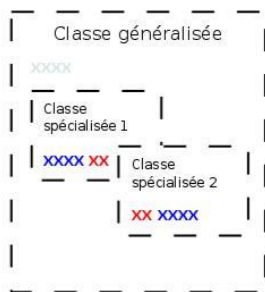
- Disjunction & Non-Coverage → Exclusion • Non-

disjunction & Non-Coverage → No Constraints

EXCLUSION



Absence de Contrainte



Example (1)

Personnel management in a university

- Cover + disjunction \rightarrow Partition

Personal (P) equals to the union of Teacher (EC) and BIATOS (B) and the EC and B Intersection is Empty

- Cover + Non-disjunction \rightarrow Totality

Personal (P) equals to the union of Teacher (EC) and BIATOS (B) and the EC and B Intersection is not Empty

Example (2)

Personnel management in a university

- Non-Coverage + disjunction \rightarrow Exclusion

The Union of Teacher (EC) and BIATOS (B) is included in P and EC and B Intersection is Empty

- Non-Coverage + non-disjunction \rightarrow No constraints

The Union of Teacher (EC) and BIATOS (B) is included in P and EC and B Intersection is not Empty

Transformation of Inheritance associations

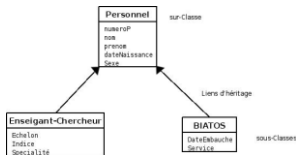
Translation of inheritance relationship depends on its constraints

→ 3 decomposition families:

- Decomposition by distinction
- Decomposition top-down (push-down)
- Decomposition bottom-up (push-up)

Decomposition by distinction

- Transformation of each subclass in a relationship
- Migration of the primary key of the super class to relationships or issues of the subclasses
- The primary key of the super class becomes both primary and foreign key



Distinction

PERSONNEL (Numéro , Nom, Prénom ,
DateNaissance , Sexe)

ENSEIGNANT(Numéro * ,
Echelon,
Indice , Spécialité)

BIATOS(Numéro _ , DateEmbauche ,
Service)

Top-down decomposition

Two possible cases according to the inheritance constraint:

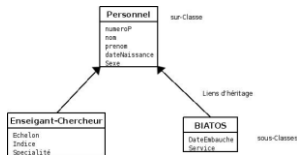
- Constraint of *partition* or *totality* on the association:
Possibility of no translation of the relationship resulting of the super class
→ Migration of all attributes on the relationships from the subclasses
- Otherwise: Migration of all attributes on relationships from the subclasses
→ Duplicate data

Top-down decomposition

Example

Partition Constraint:

- No staff can be both teacher and BIATOS
- It also no personal being neither teacher nor BIATOS.

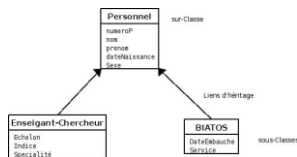


ENSEIGNANT(Numéro ,
Nom, Prénom , DateNaissance ,
Sexe , Echelon , Indice ,
Spécialité)
BIATOS(Numéro , Nom, Prénom ,
DateNaissance , Sexe
DateEmbauche , Service)

Bottom-up decomposition

- Removed all relationships resulting from or subclasses
- Migration of attributes on the relationship of the super class

Example: (No constraints)



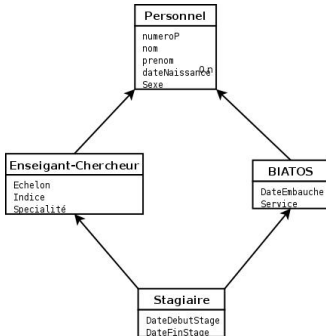
ascending

PERSONNEL(Numéro , Nom, Prénom ,
DateNaissance , Sexe ,
Echelon , Indice ,
Specialité,
DateEmbauche , Service)

Transformation of multiple inheritance associations

Same rules; several possibilities

Example: (Bottom up decomposition) teacher on exclusion Constraint on *Enseignant* and *BIATOS*



PERSONNEL(Numéro , Nom, Prénom ,
DateNaissance , Sexe)

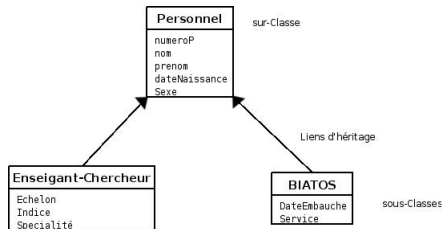
ENSEIGNANT (Numéro * ,
Echelon , Indice , Spécialité ,
DateDébutStage , DateFinStage)

BIATOS(Numéro* , DateEmbauche ,
Service ,
DateDébutStage , DateFinStage)

Transformation of inheritance associations on SQL 2

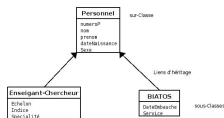
Example

Personnel management in a university



Transformation of inheritance associations

Decomposition by distinction



Distinction

PERSONNEL(Numéro , Nom, Prénom ,
DateNaissance , Sexe)

ENSEIGNANT(Nu ENSEIGNANT(Numéro_
méro_
Echelon , Echelon ,
Indice , Spécialité)

BIATOS(Numéro_ , DateEmbauche ,

REM*** A staff at University ed

create Table PERSONNEL

(NUMBER **number** (7),

NOM **varchar** (1 0),

PRENOM **varchar** (1 0),

DATENAISSANCE **date**,

SEXE char (1),

constraint PK_PERSONNEL primarykey (NUMERO)

constraint CK_SEXE_PERSONNEL check (In SEXE ('M', 'F'))

);

Transformation of inheritance associations

Decomposition by distinction

REM*** Personnel enseignant

```
create table ENSEIGNANT
(NUMERO number (7), ECHELON number (2),
 INDEX number (5) SPECIALTY varchar (20),
 constraint      PK_ENSEIGNANT primary key (NUMERO)
 constraint      FK_ENS_PERS foreign key (NUMERO)
                    References PERSONNEL
);
```

REM Personnel BIATOS (Ing, Adm, Tech, Ouv, Serv)

```
create table BIATOS
( NUMERO number (7), DATEEMBAUCHE dates
  SERVICE varchar (20),
  constraint PK_BIATOS primary key (NUMERO)
  constraint FK_BIATOS_PERS foreign key (NUMERO)
                    references PERSONNEL
);
```

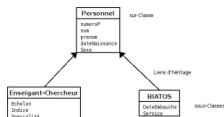
Transformation of inheritance associations

Down up decomposition

down

ENSEIGNANT(Numéro ,
Nom, Prénom , DateNaissance ,
Sexe , Echelon , Indice ,
Spécialité)

BIATOS(Numéro , Nom, Prénom ,
DateNaissance , Sexe ,
DateEmbauche , Service)



```
REM      **** Personnel  enseignant
create   table ENSEIGNANT
        (NUMERO number (7) NOM varchar (1 0),
        PRENOM varchar (1 0), DATENAissance
          SEX char (1), STEP number (2),
        INDEX number (5) SPECIALTY varchar (2 0),
        constraint CK_SEXE_ TEACHER check
          (SEX in
```

date ,

('M', 'F'))
key
(NUMERO));

```
constraint PK_ENSEIGNANT primary
```

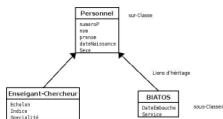
Transformation of inheritance associations

Down up decomposition

down

ENSEIGNANT(Numéro ,
Nom, Prénom , DateNaissance ,
Sexe , Echelon , Indice ,
Spécialité)

BIATOS(Numéro , Nom, Prénom ,
DateNaissance , Sexe ,
DateEmbauche , Service)



REM P ersonnel BIATOS

create table BIATOS

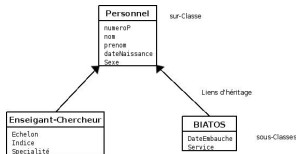
(NUMERO number (7) NAME varchar (1 0) NOM varchar (1 0),
BirthDate Date, SEX char (1), DATEEMBAUCHE date
SERVICE varchar (0 2),

constraint CK_SEXE_ BIATOS check (in SEX ('M', 'F'))

constraint PK_BIATOS primary key (NUMERO));

Transformation of inheritance associations

Bottom updecomposition



ascending

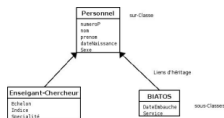
PERSONNEL (Numéro , Nom, Prénom ,
DateNaissance , Sexe ,
Echelon , Indice ,
Spécialité ,
DateEmbauche , Service)

Create table PERSONNEL

```
(NUMERO number (7) NOM varchar (1 0) PRENOM varchar (1 0),  
DateNaissance Date, SEXE char (1), Echelon number (2),  
INDICE number (5) SPECIALTE varchar (2 0),  
DATEEMBAUCHE dates SERVICE varchar (2 0),  
constraint CK_SEXE_ STAFF check (in SEXE ( 'M', 'F'))  
constraint PK_ PERSONNEL PRIMARY KEY (NUMERO)  
);
```


Translation of inheritance constraints

Decomposition distinction



• → score Constraint •
Stress screen • exclusion
constraint • Without
restraint

Inheritance Constraints:

- (Strain A) There is no staff both teaching and BIATOS
- (Strain B) There is no staff or teacher or BIATOS

Translation of inheritance constraints

Decomposition distinction

Implementing the constraint A: 2 triggers

```
REM          D e Trigger          TEACHER on

creat r eo          puts          trigger TRIG_ENSEIGNANT
before insert update to r          NUMBER of one TEACHER
for EACH rOW
declared
    number num;
begin
    select NUMBER INTO num
    from Where BIATOS NUMBER =: new .Reference;
    __ raise Application error (-20001, 'Lep erson the ||
        __ to char (num) || 'Es td e j BIATOS! !! ');
exception
    When          Then null ;
    NO_DATA_FOUND
end ;
/
```

Translation of inheritance constraints

Decomposition distinction

```
REM      D e Trigger      on BIATOS

creat r eo      puts      trigger TRIG_BIATOS      /
before inser update to r      NUMBER one of BIATOS
for EACH rOW
declared
    number num;
begin
    select NUMBER INTO num
    from TEACHER Where NUMBER =: new .Reference;      '||
    __ raise Application error (-20001, 'The      staff
        _ to char (num) || 'Es td e j? E nseignant! !! ');

exception
    When Then NO_DATA_FOUND null ;
end ;
```

Translation of inheritance constraints

Decomposition distinction

Implementing the constraint B:

- stored procedures (Insert, Delete)
- triggers (Change)

REM A joutd an E nseignant

```

create or replace procedure AJOUT_ENSEIGNANT
(NUM number, NAME varchar, varchar PREN, DNAIS date
SEX varchar, ECHEL number, IND number, varchar SPEC)
begin
insert into PERSONNEL values (NUM, NOUN, PREN, DNAIS, SEX);
insert into TEACHER into gains
(NUM, ECHEL, IND, SPEC);
end;
/

```


Translation of inheritance constraints

Decomposition distinction

```
REM      S uppression d an E nseignant
      gol
create   d  replaceprocedure      SUPPR_ENSEIGNANT
(NUM number) is
begin
  delete   from  TEACHER      Where NUMBER = num      ;
  delete   from  Where PERSONNEL NUMBER = num      ;
end;
/
```

```
REM      S uppression d a BIATOS
      gol
create   d  replaceprocedure      SUPPR_BIATOS
(NUM number) is
begin
  delete   from  Where BIATOS NUMBER = num;
  delete   from  Where PERSONNEL NUMBER = num      ;
end ;
/
```

Translation of inheritance constraints

Decomposition distinction

REM D e Trigger for I ar ed percussio has ndel

change of Scan

```
creat r eo            puts            trigger TRIG_ENSBIATOS
before            update NUMBER of it STAFF
for    EACH rOW
begin
    begin
        update TEACHER
        set NUMBER =: new .Reference
        Where NUMBER =: old .Reference;
    exception
        When    Then NO_DATA_FOUND null;
    end ;
    update    BIATOS
        set NUMBER =            : New .Reference
        Where NUMBER =            : Old .Reference;
    exception
        When Then NO_DATA_FOUND            null ;
    end ;
/
```

Translation of inheritance constraints

use

```

REM          Insertions      of the      ed given Essous      SQLPLUS

```

REM launching of the proc hard é

```
select 'Insertion of data é es' from dual;
```

PAUSE

executed	AJOUT_ENSEIGNANT	(1, 'TRAIFOR'	Clément '
----------	------------------	---------------	-----------

executed	AJOUT ENSEIGNANT	(2, 'TRAIFOR'	'C l e mentine'
----------	------------------	---------------	-----------------

'22-11-1969' 'F', 6 7 8 0, 'IA');

executed	AJOUT BIATOS (3,	'COOKING POT' , 'A lex'
----------	------------------	-------------------------

'16-10-1960', 'M' '01-01-2002' 'C ommercial');

PAUSE

```
select      from      STAFF :
```

```
select      from  TEACHER :
```

```
select      from      BIATOS;
```

PAUSE

Translation of inheritance constraints

Decomposition distinction

→ *Stress screen*

Inheritance Constraints:

- (Strain B) There is no staff or teacher or BIATOS
- (Strain C) There may be personal to both teacher and BIATOS

implementation:

- Constraint B: see above
- Constraint C: equivalent to not program the previous constraint A
⇒ No implementation of the triggers and tables TEACHER
BIATOS: TRIG_ENSEIGNANT, TRIG_BIATOS

Translation of inheritance constraints

Decomposition distinction

→ *exclusion constraint*

Inheritance Constraints:

- (Strain A) There is no staff both teaching and BIATOS
- (Strain D) There may be a teacher or staff or BIATOS

implementation:

- Constraint A: see above
- Constraint D: equivalent to not program the previous constraint B
⇒ not to implement the four procedures (add, delete) and trigger TRIG_ENSBIATOS

Translation of inheritance constraints

Decomposition distinction

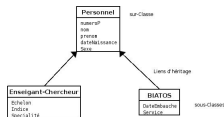
→ *Without restraint*

No compulsion is to be programmed!

Translation of inheritance constraints

down decomposition

down



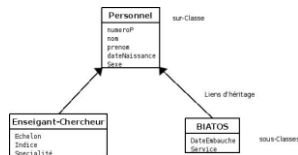
TEACHER (number,
Last name First Name ,
Gender, E chelon, Index ,
S p e c i a l i t e)
Date of birth ,

BIATOS (Number, Name,
Date of birth , gender,
HireDate , S service)

- Constraint score? → no staff can be both teacher and IATOS and it does not is a personal being neither teacher nor IATOS
- Forced to all?
- exclusion constraint? → should the staff table for non-teaching personnel and non BIATOS
- Without restraint !

Translation of inheritance constraints

decomposition upward



● → score Constraint ●

Stress screen ● exclusion
constraint ● Without
restraint

Inheritance Constraints:

- (Strain A) There is no staff both teaching and BIATOS
- (Strain B) There is no staff or teacher or BIATOS

Translation of inheritance constraints

decomposition upward

Implementation constraints A and B:

- staff at the table
- with the type of CHECK constraints
- Constraint A:
Check the ECHELON columns, INDEX, SPECIALTY,
DATEEMBAUCHE SERVICE and are not initialized all
- Constraint B:
Check the ECHELON columns, INDEX, SPECIALTY,
DATEEMBAUCHE SERVICE and are not all zero

Translation of inheritance constraints

decomposition upward

A REMCONTRAINTE

```
alter table STAFF
add constraint CK_CONTRAINTE_A
check (
    (ECHELON is null and INDEX is null and
      SPECIALTY is null )
    gold (HIREDATE is null and SERVICE is null )
);
```

REMCONTRAINTE B

```
alter table STAFF
add constraint CK_CONTRAINTE_B
check (
    (ECHELON is not null gold INDEX is not null gold
      SPECIALTY is not null )
    gold (HIREDATE is not null gold SERVICE is not null )
);
```

Translation of inheritance constraints

decomposition upward

→ *Stress screen*

Inheritance Constraints:

- (Strain B) There is no staff or teacher or BIATOS
- (Strain C) There may be personal to both teacher and BIATOS

Translation of inheritance constraints

decomposition upward

→ *Stress screen*

- Constraint B: see above
- Constraint C: Remove or disable the constraint A previous
(**DROP CONSTRAINT** or **DISABLE CONSTRAINT**)
 - | **DROP CONSTRAINT** : If reactivation of stress, you need to recreate (**ADD CONSTRAINT**)
 - | **DISABLE CONSTRAINT**: if reactivation of the constraint, simply reactivate the request with **ENABLE CONSTRAINT**

Translation of inheritance constraints

decomposition upward

→ *Stress screen*

REM The constraint C returns to fairela

REM Qu deactivation of At the STRESS

altertable STAFF

disableconstraint CK_CONTRAINTE_A;

Translation of inheritance constraints

decomposition upward

→ *exclusion constraint*

inheritance constraint:

- A Constraint → Reactivating the constraint A deleting tuples previously not responding to this constraint
- No-stress B → Disable strain B

Translation of inheritance constraints

decomposition upward

→ *exclusion constraint*

REM R e activation of At the STRESS
altertable STAFF
 enable **constraint** CK_CONTRAINTE_A;

REM Qu deactivation of the B STRAIN
altertable STAFF
 disable **constraint** CK_CONTRAINTE_B;

Translation of inheritance constraints

decomposition upward

→ *Without restraint*

No type of CHECK constraint is to program!

Conclusion / Summary

None of the solutions a panacea.

It is necessary to measure the performance of queries.

See also the kind of queries

part III

-

Inheritance on SQL 3

Transformation of Inheritance associations SQL 3

- Inheritance of types:

- | since version 9.1 of Oracle (November 2001)
- | No multiple inheritance:

- Inheritance of tables

- | Only on the Oracle's latest version



Inheritance of types

Definition of a staff at the University

```
---                               Creation of the super class  
                                type  
create type PERSONNEL_TYPE AS OBJECT  
(  
    NUMERO number ( 7 ),  
    NOM varchar ( 10 ),  
    PRENOM varchar ( 10 ),  
    DATENAISSANCE date ,  
    SEXE char ( 1 )  
)  
NOT FINAL    / *can include sub classes          */  
/
```

Inheritance of types

Definition of a teacher

--- *Creation of the sub class*

```
Create type ENSEIGNANT_TYPE UNDER PERSONNEL_TYPE  
(  
    ECHELON number ( 2 ) ,  
    INDICE number ( 5 ) ,  
    SPECIALITE v a r c h a r (20 )  
)  
FINAL  
/
```

Creating object tables and constraints

- Creating object tables depending in the types previously defined
- No guidelines specifying the inheritance; it is induced by the existing type hierarchy

```
— *** Personnel de l'université
create table PERSONNEL OF PERSONNEL_TYPE
(
  constraint PK_PERSONNEL primary key (NUMERO),
  constraint CK_SEXE_PERSONNEL check (SEXE in ('M', 'F'))
);
— *** Personnel enseignant
create table ENSEIGNANT OF ENSEIGNANT_TYPE ;
```

IMPORTANT: the constraints are defined only in the “Personnel” table

Creating object tables and constraints

Illustration

NB: the constraints are defined only in the “Personnel” table → It inherits a type

Inserting data into the “Personnel” table:

```
insert into personnel values (1, 'B', 'F', '17-09-2004', 'M');
insert into personnel values (1, 'B', 'F', '17-09-2004', 'M');
```

*
ERREUR à la ligne 1 :

ORA-00001: violation de contrainte **unique** (FB.PK_PERSONNEL)

```
select * from personnel;
NUMERO      NOM      PRENOM      DATENAISSA S
-----
1 B          F          17-09-2004 M
```

Creating object tables and constraints

Illustration

NB: the constraints are defined only in the “Personnel” table → It inherits a type

Inserting data into the “Enseignant” table:

```
insert into enseignant values (7, 'B', 'F', '17-09-2004', 'M', 2, 780, 'BD');
1 ligne créée.
insert into enseignant values (7, 'B', 'F', '17-09-2004', 'M', 2, 780, 'BD');
1 ligne créée. !!!
insert into enseignant values (8, 'B', 'D', '17-10-2004', 'M', 2, 780, 'BD');
1 ligne créée.

select * from enseignant ;
```

NUMERO	NOM	PRENO	DATENAIS	SA	S	ECHOLON	INDICE	SPECIALITE
7	B	F	17-09-2004	M		2	780	BD
7	B	F	17-09-2004	M		2	780	BD
8	B	D	17-10-2004	M		2	780	BD

Creating object tables and constraints (2)

- Creating object tables depending on types previously defined
- Defining constraints on tables

```
-- *** Personnel de l'université
create table PERSONNEL OF PERSONNEL_TYPE
(
    constraint PK_PERSONNEL primary key (NUMERO),
    constraint CK_SEXE_PERSONNEL check (SEXE in ('M', 'F'))
);
```

Transformation of inheritance associations SQL 3

```
— *** Personnel enseignant  
create table ENSEIGNANT OF ENSEIGNANT_TYPE  
(  
  constraint PK_ENSEIGNANT primary key (NUMERO),  
  constraint CK_SEXE_ENSEIGNANT check (SEXE in ( 'M', 'F' ))  
);
```

WARNING :

- define constraints, in the “Enseignant” table
- Inheritance of a type

Creating tables and constraints object (2)

illustrations

The constraints should be defined in the “Enseignant” table too:

```
SQL> insert into enseignant values (7, 'B', 'F', '17-09-2004', 'M', 2, 780, 'BD');  
1 ligne créée.
```

```
SQL> insert into enseignant values (7, 'B', 'F', '17-09-2004', 'M', 2, 780, 'BD');
```

```
* ERREUR à la ligne 1 : ORA-00001: violation de contrainte unique (FB.PK_ENSEIGNANT)
```

```
SQL> insert into enseignant values (8, 'B', 'D', '17-10-2004', 'M', 2, 780, 'BD');  
1 ligne créée.
```

```
SQL> insert into enseignant values (9, 'B', 'D', '17-10-2004', 'K', 2, 780, 'BD');
```

```
* ERREUR à la ligne 1 : ORA-02290: violation de contraintes (FB.CK_SEXE_ENSEIGNANT)  
de vérification
```

```
SQL> select * from enseignant ;
```

NUMERO	NOM	PRENO	DATENAISSA	S	ECHELON	INDICE	SPECIALITE
7	B	F	17-09-2004	M	2	780	BD
8	B	D	17-10-2004	M	2	780	BD

Transformation of Inheritance associations SQL 3

Inheritance of tables

```
REM *** Un personnel à l'Université
create table PERSONNEL
(
  NUMERO number(7), NOM varchar(10),
  PRENOM varchar(10), DATENAISSANCE date,
  SEXE char(1),
  constraint PK_PERSONNEL primary key (NUMERO),
  constraint CK_SEXE_PERSONNEL check (SEXE in ('M', 'F'))
);
```

Transformation of Inheritance associations SQL 3

Inheritance of tables

```
REM *** Personnel enseignant
create table ENSEIGNANT under PERSONNEL
(
    ECHELON number(2),
    INDICE number(5),
    SPECIALITE varchar(20)
);
```



```
REM *** Personnel biatos
create table BIATOS under PERSONNEL
(
    DATEEMBAUCHE date,
    SERVICE varchar(20)
);
```

part IV

-

Objects in SQL3

Object Programming - SQL 3

- Object-Relational - Object
- Translation UML → Object / Relational Object

relational schema / SQL2

- Relational schema:

COURS (NUM_COURS, NOMC, NBHEURES, ANNEE)

PROFESSEURS (NUM_PROF, NOMP, SPECIALITE, DATE_ENTREE,
DER_PROM, SALAIRE_BASE, SALAIRE_ACTUEL
)

CHARGE(NUM_PROF*, NUM_COURS*)

relational schema / SQL2

SQL2:

```
create table COURS
( NUM_COURS      NUMBER(2)   NOT NULL,
  NOMC            VARCHAR2(20) NOT NULL,
  NBHEURES        NUMBER(2),
  ANNE            NUMBER(1),
  constraint PK_COURS primary key (NUM_COURS)
);

create table PROFESSEURS
( NUM_PROF        NUMBER(4)   NOT NULL,
  NOMP            VARCHAR2(25) NOT NULL,
  SPECIALITE      VARCHAR2(20),
  DATE_ENTREE     DATE,
  DER_PROM        DATE,
  SALAIRE_BASE    NUMBER,
  SALAIRE_ACTUEL  NUMBER,
  constraint PK_PROFESSEURS primary key (NUM_PROF)
);

create table CHARGE
( NUM_PROF        NUMBER(4)   NOT NULL,
  NUM_COURS       NUMBER(4)   NOT NULL,
  constraint PK_CHARGE primary key (NUM_COURS,
                                     NUM_PROF)
);
```

relational schema / SQL2

```
alter table CHARGE
  add constraint FK_CHARGE_COURS
    foreign key (NUM_COURS)
      references COURS (NUM_COURS);

alter table CHARGE
  add constraint FK_CHARGE_PROFESSEUR
    foreign key (NUM_PROF)
      references PROFESSEURS (NUM_PROF);
```

Schema object-relational / SQL3

- object-relational schema

COURS (NUM_COURS, NOMC, NBHEURES, ANNEE)

PROFESSEURS (
NUM_PROF, NOMP, SPECIALITE, DATE_ENTREE,
DER_PROM, SALAIRE_BASE, SALAIRE_ACTUEL,
EnsembleDe (COURS)
)

Schema object-relational / SQL3

SQL3:

```
create type cours_type as object
( num_cours number(2), nomc varchar2(20),
  nbheures number(2), annee number(1) )
/

create type lescours_type as table of cours_type
/

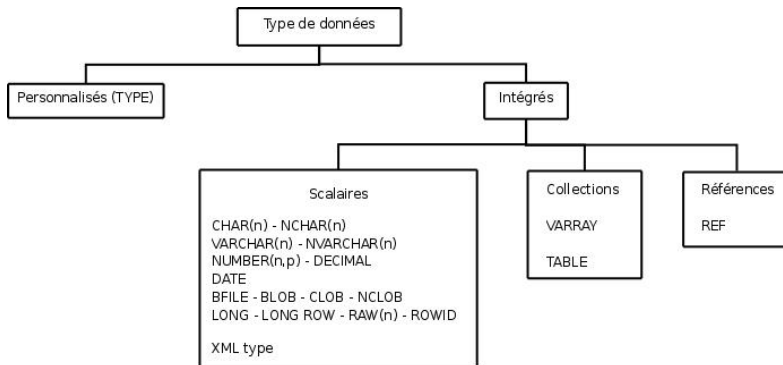
create type professeur_type as object
( num_prof number(4), nom varchar2(25),
  specialite varchar2(20), cours lescours_type ...)
/

create table professeur of professeur_type
( primary key (num_prof) )
nested table cours store as tabemp
/
```

Object Types

Data type

Main type of Oracle data:



Object Types

Persistence

Under Oracle, three categories of objects:

- Objects column (column objects) stored as structured column in a relational table;
- Online Items (row objects) stored as a line item table.
 - | possess a unique identifier known as OID (Object Identifier)
 - | can be indexed and partitioned
- non-persistent objects: not stored
 - | nor in a column of a relational table
 - | neither in a table object Online

These objects exist only during the execution of a PL / SQL program

Object Types

- Definition of each object from a type describing
 - | a data structure positioning in an inheritance hierarchy
 - | methods
- Using a Type:
 - | Build More types
 - | Define one or more object tables
 - | Define a column in a relational table
 - | Building object views

Object Types

Creating a type

• Creation

```
CREATE [OR REPLACE TYPE] schéma.nomType
      [AS OBJECT | UNDER schéma.nomSurType]
(
  REM *** définition de la structure
  colonne1 type1, colonne2 type2, ...,

  REM *** définition du comportement
  méthode1 (paramètres1), méthode2 (paramètres2) ...
)
[[NOT] INSTANTIABLE]

REM *** positionnement dans le graphe d'héritage
[[NOT] FINAL]
/
```

Creating a type

FINAL directive

- FINAL and NOT FINAL: positioning of a type in the inheritance graph
- NOT Final: to be applied to generic types
- By default, any type is FINAL
 - A FINAL type can be used to define subtypes

Creating a type

FINAL Directive - Examples

```
CREATE TYPE adresse_t AS OBJECT (  
  nrue NUMBER(3), rue VARCHAR(40), ville VARCHAR(30)  
/
```

```
CREATE TYPE Personnel_t AS OBJECT(  
  nom VARCHAR (10), prenom VARCHAR(10), adresse adresse_t))  
NOT FINAL  
/
```

```
CREATE TYPE Enseignant_t UNDER Personnel_t(  
  Echelon NUMBER, indice NUMBER)  
FINAL
```

Creating a type

INSTANTIABLE directive

- INSTANTIABLE and NOT INSTANTIABLE: to instantiate a type
All types are created by default INSTANTIABLE
- NOT INSTANTIABLE: similar to the concept of abstract class
- Each type has
 - | a constructor to create objects (persistent or not) with the NEW command or in an INSERT clause
 - | a constructor (default) and several in the case of overload
- One type NOT INSTANTIABLE can not be FINAL
- A subtype NOT INSTANTIABLE can inherit a type INSTANTIABLE

Creating a type

INSTANTIABLE Directive - Examples

```
CREATE TYPE Personnel_t AS OBJECT(  
    nom VARCHAR (10), prenom VARCHAR(10), adresse adresse_t))  
    NOT INSTANTIABLE NOT FINAL  
/
```

```
CREATE TYPE Enseignant_t UNDER Personnel_t(  
    Echelon NUMBER, indice NUMBER)  
    INSTANTIABLE FINAL  
/
```

Object Types

Deleting a type

DROP TYPE nomType [FORCE | VALIDATE] ;

- **FORCE**: remove a type even if there exist objects belong to this type in a database
Oracle : columns of this type, are labeled “UNUSED”, and they become inaccessible (not recommended)
- **VALIDATE**: Check if instances of the type to be deleted can be substituted by a superclass.

example:

DROP TYPE Personnel_t **FORCE**

Object Types

Creating a type

Specifying the Object

```
CREATE TYPE Bank_Account AS OBJECT (  
    acct_number INTEGER(5),  
    balance      REAL,  
    status       VARCHAR2(10),  
  
    MEMBER PROCEDURE open  
        (amount IN REAL),  
  
    MEMBER PROCEDURE verify_acct  
        (num IN INTEGER),  
  
    MEMBER PROCEDURE close  
        (num IN INTEGER, amount OUT REAL)  
);  
  
CREATE TYPE BODY Bank_Account AS  
    ...  
END;
```



Object Types

Creating a type

Definition of methods associated with the object

```
CREATE TYPE BODY Bank_Account AS
```

```
MEMBER PROCEDURE open (amount IN REAL) IS
BEGIN — open account with initial deposit
    IF NOT amount > 0 THEN
        RAISE_APPLICATION_ERROR(-20104, 'bad amount');
    END IF;
    — SELECT acct_sequence.NEXTVAL INTO acct_number FROM dual;
    status := 'open';
    balance := amount;
END open;
```

Object Types

Creating a type

```
MEMBER PROCEDURE verify_acct (num IN INTEGER) IS
BEGIN — check for wrong account number or closed account
  IF (num <> acct_number) THEN
    RAISE_APPLICATION_ERROR(-20105, 'wrong number');
  ELSIF (status = 'closed') THEN
    RAISE_APPLICATION_ERROR(-20106, 'account closed');
  END IF;
END verify_acct;

MEMBER PROCEDURE close (num IN INTEGER, amount OUT REAL) IS
BEGIN — close account and return balance
  verify_acct(num);
  status := 'closed';
  amount := balance;
END close;
END;
```

Object Types

Extraction of the description of a type

Defining new views to account types

example:

```
create type emp_type as object ( ninsee varchar2 ( 13 ) ,  
                                age number , nom varchar2 ( 30 ) ) har 2 (3 0))  
/
```

Description of the structure of the 1st level of a type:

```
SQL> DESC emp_type
```

Object Types

Extraction of the description of a type

Examples of views: (USER _..., DBA _..., ALL _...)

Description:

- collections: USER_COLL_TYPES
- indexes on the types: USER_INDEXTYPES
- types in general: USER_TYPES
- Attribute types: USER_TYPE_ATTRS
- Methods types: USER_TYPE_METHODS
- Revision types: USER_TYPE_VERSIONS

To the object

relational tables

— Table : MAGASINS2 SQL2

```
create table MAGASINS2
(
  NUMMAG          INTEGER          ,
  NOMMAG          CHAR(30)         ,
  TELMAG          CHAR(15)         ,
  ADRNUMMAG       VARCHAR2(10)    ,
  ADDRUEMAG       VARCHAR2(50)    ,
  ADRCPMAG        VARCHAR2(10)    ,
  ADRVILLEMAG    VARCHAR2(50)    ,
  ADRPAYSMAG      VARCHAR2(50)    ,
  constraint PK_MAGASINS2
    primary key (NUMMAG) );
```

```
insert into MAGASINS2 values (1, 'FB', '0145454545', '13', 'Avenue de la paix',
                              '75015', 'Paris', 'France');
```

NUMMAG	NOMMAG	TELMAG	ADRNU	ADDRUEMAG	ADRC	ADRVILLEMA	ADRPAYSMAG
1	FB	0145454545	13	Avenue de la paix	75015	Paris	France
2	FB	0155555555	20	Avenue de la liberté	06100	Nice	France
3	FB	0155555555	10	Avenue des Amis	6050	Bruxelles	Belgique
4	FB	71226002	10	Avenue du soleil	1001	Tunis	Tunisie

— Table : CLIENTS2 SQL2

```
create table CLIENTS2
(
  NUMCLI          INTEGER          ,
  NOMCLI          CHAR(20)         ,
  TELCLI          CHAR(15)         ,
  ADRNUMCLI       VARCHAR2(10)    ,
  ADDRUECLI       VARCHAR2(50)    ,
  ADRCPCLI        VARCHAR2(10)    ,
  ADRVILLECLI    VARCHAR2(50)    ,
  ADRPAYSCLI      VARCHAR2(50)    ,
  constraint PK_CLIENTS2
    primary key (NUMCLI));
```

NUMCLI	NOMCLI	TELCLI	ADRNU	ADDRUECLI	ADRC	ADRVILLECLI	ADRPAYSCLI
1	TRAIFOR	0645454545	13	Avenue de la paix	75015	Paris	France
2	CLEMENT	0607080910	17	Avenue de la paix	75015	Paris	France
3	SOUCY	98980307	77	Route de la corniche	4001	Sousse	Tunisie

Object Tables

Creating a type - TAD

First extension of the relational model: Abstract Data Types (ADT)
TAD (BD context):

- New attribute type defined by the user
- shared data structure
 - | Use of the type in one or more tables
 - | Participation in the composition of one or more other types

Remarks:

- A TAD includes methods that are procedures or functions
- They allow you to manipulate objects of the abstract type

Object Tables

Creating a type – example of TAD

```
create type ADRESSE_TYPE as object
(   ADRNUM      VARCHAR2(10),
    ADDRUE      VARCHAR2(50),
    ADRCP       VARCHAR2(10),
    ADRVILLE    VARCHAR2(50),
    ADRPAYS     VARCHAR2(50) )
/
```

```
create type MAG_TYPE as object
(   NUMMAG      INTEGER ,
    NOMMAG      CHAR(30),
    TELMAG      CHAR(15),
    ADRMAG      ADRESSE_TYPE )
/
```

```
create type CLI_TYPE as object
(   NUMCLI      INTEGER ,
    NOMCLI      CHAR(30),
    TELCLI      CHAR(15),
    ADRCLI      ADRESSE_TYPE )
/
```

Object Tables

Creating a table - Examples

```
create table MAGASINS3 OF MAG_TYPE  
( constraint PK_MAGASINS3 primary key (NUMMAG) );
```

```
create table CLIENTS3 OF CLI_TYPE  
( constraint PK_CLIENTS3 primary key (NUMCLI) );
```

Object Tables

Creating a type

Remarks:

- One type can not contain constraints (NOT NULL, CHECK, UNIQUE, DEFAULT, PRIMARY KEY, FOREIGN KEY, etc.).
- The constraints must be declared at the table object

Access to the description of the types from the Data Dictionary:

```
SQL > select table_name, object_id_type, table_type_owner,  
            table_type from user_object_tables;
```

Object Tables

Creation / description of a table

SQL> desc clients2

Nom	NULL ?	Type
NUMCLI	NOT NULL	NUMBER(38)
NOMCLI		CHAR(20)
TELCLI		CHAR(15)
ADRNUMCLI		VARCHAR2(10)
ADRRUECLI		VARCHAR2(50)
ADRCPCLI		VARCHAR2(10)
ADRVILLECLI		VARCHAR2(50)
ADRPAYSCLI		VARCHAR2(50)

SQL> desc clients3

Nom	NULL ?	Type
NUMCLI	NOT NULL	NUMBER(38)
NOMCLI		CHAR(30)
TELCLI		CHAR(15)
ADRCLI		ADRESSE_TYPE

Object Tables

Object identifier (OID)

- OID based on the primary key: Using the primary key option example:

```
create table CLIENTS3 OF CLI_TYPE
  ( constraint PK_CLIENTS3 primary key (NUMCLI)
    object identifier is primary key );
```

- OID Indexes:

```
create table CLIENTS3 OF CLI_TYPE
  ( constraint PK_CLIENTS3 primary key (NUMCLI) )
  object identifier is system generated OIDINDEX ndxclients3;
```

```
create table CLIENTS3 OF CLI_TYPE
  ( constraint PK_CLIENTS3 primary key (NUMCLI) )
  object identifier is system generated
  OIDINDEX ndxclients3 (storage (initial 100K next 50k
                               minextents 1 maxextents 50)
                        );
```

Object Tables

Instantiation - examples

Inserting a "line" (or rather an object):

```
insert into MAGASINS3 values (MAG_TYPE(1, 'FB', '0145454545',  
    ADRESSE_TYPE('13', 'Avenue de la paix', '75015', 'Paris', 'France')));  
insert into MAGASINS3 values (MAG_TYPE(2, 'FB', '0155555555',  
    ADRESSE_TYPE('20', 'Avenue de la liberté', '06100', 'Nice', 'France')));  
insert into MAGASINS3 values (MAG_TYPE(3, 'FB', '0155555555',  
    ADRESSE_TYPE('10', 'Avenue des Amis', '6050', 'Bruxelles', 'Belgique')));  
insert into MAGASINS3 values (MAG_TYPE(4, 'FB', '71226002',  
    ADRESSE_TYPE('10', 'Avenue du soleil', '1001', 'Tunis', 'Tunisie')));
```

```
SQL> select * from magasins3 ;  
NUMMAG  NOMMAG  TELMAG  ADRMAG(ADRRUE, ADRCP, ADRVILLE, ADRPAYS)  
1         FB    0145454545  ADRESSE_TYPE('13', 'Avenue de la paix',  
              '75015', 'Paris', 'France')  
2         FB    0155555555  ADRESSE_TYPE('20', 'Avenue de la liberté',  
              '06100', 'Nice', 'France')  
3         FB    0155555555  ADRESSE_TYPE('10', 'Avenue des Amis',  
              '6050', 'Bruxelles', 'Belgique')  
4         FB    71226002    ADRESSE_TYPE('10', 'Avenue du soleil',  
              '1001', 'Tunis', 'Tunisie')
```

Object Tables

Instantiation - examples

```
insert into CLIENTS3 values (CLI_TYPE(1, 'TRAIFOR', '0645454545',  
    ADRESSE_TYPE('13', 'Avenue de la paix', '75015', 'Paris', 'France')));  
insert into CLIENTS3 values (CLI_TYPE(2, 'CLEMENT', '0607080910',  
    ADRESSE_TYPE('17', 'Avenue de la paix', '75015', 'Paris', 'France')));  
insert into CLIENTS3 values (CLI_TYPE(3, 'SOUCY', '98980307',  
    ADRESSE_TYPE('77', 'Route de la corniche', '4001', 'Sousse', 'Tunisie')));
```

```
SQL> Select * from clients3 ;
```

NUMCLI	NOMCLI	TELCLI	ADRCLI(ADNRUM, ADDRUE, ADRCP, ADRVILLE, ADRPAYS)
1	TRAIFOR	0645454545	ADRESSE_TYPE('13', 'Avenue de la paix', '75015', 'Paris', 'France')
2	CLEMENT	0607080910	ADRESSE_TYPE('17', 'Avenue de la paix', '75015', 'Paris', 'France')
3	SOUCY	98980307	ADRESSE_TYPE('77', 'Route de la corniche', '4001', 'Sousse', 'Tunisie')

Object Tables

instantiation

object-relational table:

- dependent table of a type
- Records (rows) in this table seen as objects because they all have a single OID (Object Identifier)

```
SQL> SELECT * FROM clients3 ;
```

NUMCLI	NOMCLI	TELCLI	ADRCLI(ADRNUM, ADDRUE, ADRCP, ADRVILLE, ADRPAYS)
1	TRAIFOR	0645454545	ADRESSE_TYPE('13', 'Avenue de la paix', '75015', 'Paris', 'France')
2	CLEMENT	0607080910	ADRESSE_TYPE('17', 'Avenue de la paix', '75015', 'Paris', 'France')
3	SOUCY	98980307	ADRESSE_TYPE('77', 'Route de la corniche', '4001', 'Sousse', 'Tunisie')

Object Tables

instantiation

Reference OID objects of the table:

```
SQL> SELECT REF(c) FROM clients3 c ;
```

REF(C)

```
0000280209E9E229206EDF47DF9996946C4BBD571C4EB9AF259F2F42BC813E18E51603C0D4024001460000  
0000280209550141E8898C4859AF0F3D48FA3041944EB9AF259F2F42BC813E18E51603C0D4024001460001  
0000280209C2C96804847047F6856499690AAC9E254EB9AF259F2F42BC813E18E51603C0D4024001460002
```

Object Tables

Updates

Changes / Deletions of "lines" or objects

- Updating a standard column

```
update clients3  
set NOMCLI = 'CBON' where NUMCLI=2;
```

- Changing a column belonging to a nested type

```
update clients3 c  
set c.ADRCLI.ADRVILLE = 'MAVILLE' where c.NUMCLI=2;
```

- Deleting object

```
delete from clients3  
where numcli = 3;  
delete from clients3 c  
where upper ( c.ADRCLI.ADRPAYS) = 'FRANCE' ;
```

Object Tables

interrogations

- Use of standard columns

```
select numcli , nomcli from clients3 ;  
NUMCLI NOMCLI
```

```
-----  
1  TRAIFOR  
2  CLEMENT  
3  SOUCY
```

- Use of a column belonging to a nested type

```
select numcli , nomcli , c. ADRCLI. ADRPAYS  
from clients3 c ;  
NUMCLI CUSTNAME      ADRCLI. ADRPAYS
```

```
-----  
1  TRAIFOR      La France  
2  CLEMENT      La France  
3  SOUCY        T unisia
```

Subject Tables

interrogations

- with formatting

A10 name collar size

He co oc A15 size

```
select   numcli   ace cli, CUSTNAME
          c. ADRCLI. ADRVILLE || " ||
          ADRPAYS
from     3 customers c;
```

```
ace name,
c. ADRCLI.      ac   loc
                e
```

1 TRAIFOR	P aris	La France
2 CLEMENT	P aris	La France
3 SOUCY	S ousse	T unisia

Subject Tables

interrogations

● constrained

```
SQL> COLC. ADRCLI. ADRPAYS      format  A10
SQL> COLC. ADRCLI. ADRVILLE     format   A10
SQL> select numcli,             CUSTNAME, c. ADRCLI. ADRPAYS,
      2  c. ADRCLI. ADRVILLE from customers 3      c
      3 WHERE upper (C. ADRCLI. ADRVILLE)      like 'P%';
```

NUMCLI	CUSTNAME	ADRCLI. ADRPAYS	ADRCLI. ADRVILLE
1	TRAIFOR	La France	P aris
2	CLEMENT	La France	P aris

part V

-

nested tables in SQL3

nested Tables

(NESTED TABLE)

nested table (NESTED TABLE): unordered collection rather limited of the same type elements

Example: table Department



One table containing a column (table):

Association of Type 1-N

NumDep	Budget	employees		
		NInsee	Name	Age



?? ? 1 or more tables NN type Association

Nested Tables (NESTED TABLE)

Creation

```
create type emp_type as object
(Ninsee varchar 2 (1 3), age          number , Name varchar 2 (3 0))
/

create type emps_type as table of emp_type
/

create type __Department type as object
(Varchar numdep 2 (1 1), the budget number,
employed emps_type)
/

CreateTable      departement of departem ent types __
( primary key (Numdep))
nested table     employee ss torus      ace tabemp
/
```

- NESTED TABLE clause: definition of a nested table
- STORE AS clause: Naming the internal structure that stores the "records" of this nested table

Nested Tables (NESTED TABLE)

Example

```
create type emp_type      object (Ninsee varchar 2 (1 3),
                                age number, name varchar 2 (3 0))
/
create type emps_type      table emp_type of
/
create type _ Department typeas      object      (Numdep varchar 2 (1 1),
                                budget number ,      employed emps_type)
/
CreateTable      departementofdepartem ent types _
                                (primary key (Numdep))
nested table      employee ss torus      ace tabemp
/
```

Nested Tables (NESTED TABLE)

Example

SQL> desc department

Name	NULL ?	Type

NUMDEP	NOT NULL	VARCHAR2(1 1) BUDGET NUMBER USED
EMPS_TYPE		

SQL> desc emps_type

emp_type	OF EMP_TYPE	NULL ?	Type
Name			

NINSEE			VARCHAR2(1 3)
AGE			NUMBER
NAME			VARCHAR2(30)

Nested Tables (NESTED TABLE)

Insertion

• Inserting data into a nested table

```
insert into department gains ( 'D1', 1 0 0 0 0 0, emps_type ());  
insert into department gains (D2 ' 2 0 0 0 0 0, emps_type ());
```

```
SQL> select from department;  
NUMDEP          BUDGET EMPLOYEES (NINSEE, age, name)  
-----  
D1              100000 EMPS_TYPE ()  
D2              200000 EMPS_TYPE ()
```


Nested Tables (NESTED TABLE)

Insertion

- NB: In the following example, the vacuum table is uninitialized

```
insert      into    department (numdep,          budget)
              gains  (D3 '          3 0 0 0 0 0);
```

```
SQL> select      from department          ;
NUMDEP          BUDGET EMPLOYEES (NINSEE, age, name)
```

```
-----
D1              100000 EMPS_TYPE ()
D2              200000 EMPS_TYPE ()
D3              300000
```


Nested Tables (NESTED TABLE)

Insertion

- Inserting data into a nested table

```
insert into department gains ( 'D4',      4000000,  
emp_type (emp_type ( 'N5'      2 5 'B ibi),  
            emp_type ( 'N6'      2 6 'C here),  
            emp_type (N7'      2 7 'D idi),  
            emp_type ( 'N8'      2 8, 'F ifi')));
```


Nested Tables (NESTED TABLE)

Insertion

```
SQL> select          from department;
NUMDEP              BUDGET EMPLOYEES (NINSEE, age, name)
-----
D1                100000  EMPS_TYPE ()
D2                200000  EMPS_TYPE ()
D3                300000
D4                400000  EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),
                                     EMP_TYPE ( 'N6'  2 6 'C here),
                                     EMP_TYPE ( 'N7'  2 7 'D idi),
                                     EMP_TYPE ( 'N8'  2 8, 'F ifi'))
```

Note: INSERT command with the builders of the types of NESTED TABLE

- stores an object in the table
- initializes the nested table associated with records

Nested Tables (NESTED TABLE)

Insertion

- Inserting data into a nested table

```
InsertInto departementvalues (D5 ', 4 0 0 0 0 0,  
                                emps_type (emp_type ( 'N5',          2 5 'B ibi),  
                                emp_type ( 'N8'          2 8, 'F ifi')));
```

```
SQL> select      from department;  
NUMDEP          BUDGET EMPLOYEES (NINSEE, age, name)  
-----  
D1              100000  EMPS_TYPE ()  
D2              200000  EMPS_TYPE ()  
D3              300000  
D4              400000   EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),  
                                EMP_TYPE ( 'N6'  2 6 'C here),  
                                EMP_TYPE ( 'N7'  2 7 'D idi),  
                                EMP_TYPE ( 'N8'  2 8, 'F ifi'))  
D5              400000   EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),  
                                EMP_TYPE ( 'N8'  2 8, 'F ifi'))
```


Nested Tables (NESTED TABLE)

Integration with operator TABLE

- TABLE insertion with the operator in a nested table
(D1 and D2 were initialized to empty)

```
insert into TABLE (Selectd. Employees from department d
                        Where d. numdep = 'D1')
                        gains ('N1' 2 1 'CLEMMENT');
insert into TABLE (Selectd. Employees from department d
                        Where d. numdep = 'D2')
                        gains (N2 ' 2 2, 'CLEMENTINE');
```

NB: the THE operator is obsolete and has been replaced by the TABLE operator

Nested Tables (NESTED TABLE)

Integration with operator TABLE

```
SQL> select      from department      ;
NUMDEP      BUDGET EMPLOYEES (NINSEE, age, name)
-----
D1          100000      EMPS_TYPE (EMP_TYPE ( 'N1'  2 1 'CLEMENT'))
D2          200000      EMPS_TYPE (EMP_TYPE (N2 '  2 2, 'CLEMENTINE'))
D3          300000
D4          400000 EMPS_TYPE (EMP_TYPE ( 'N5  '  2 5 'B ibi),
                                EMP_TYPE ( 'N6  '  2 6 'C here),
                                EMP_TYPE (N7  '  2 7 'D idi),
                                EMP_TYPE ( 'N8  '  2 8, 'F ifi'))
```

Remarks:

- INSERT INTO TABLE (SELECT ...): storage of a record in the nested table designated TABLE
- SELECT after TABLE: Returns a single object, which selects the associated nested table

Nested Tables (NESTED TABLE)

Integration with operator TABLE

- Integration with the operator in a nested table TABLE

(D3 was not initialized to empty)

Inserting an employee in the department D3
while it did not initialize

```
insert    into TABLE (Selectd. Employees from departementd
           Where d. numdep = 'D3')
           gains    ( 'N3', 2 3, 'DOES NOT');
```

Nested Tables (NESTED TABLE)

Integration with operator TABLE

```
SQL> insert      into      table      (Selectd. Employees from      departementd
      2 Where      d. numdep = 'D3') gains      ( 'N3', 2 3,  'DO NOT WORK ' ) ;
insert      into      table      (Selectd. Employees from departementd
                                   Where d. numdep = 'D3')
```

FAULT to the line 1 :

ORA-22908: reference at a value of table **NULL**

explanations:

- 1 The D3 department is an object of the table Department
- 2 but it does not have nested table
- 3 because it was not created during insertion. We

must destroy the D3 object and recreate it!

Nested Tables (NESTED TABLE)

change

• update the main table

```
update departementd  
  set d. Budget = d. budge t1. 5  
      Where d. Budget <=      200000;
```

```
SQL>      select tfrom department      ;  
NUMDEP      BUDGET EMPLOYEES (NINSEE, age, name)  
-----  
D1           150000  EMPS_TYPE (EMP_TYPE ( 'N1'  2 1  'CLEMENT'))  
D2           300000  EMPS_TYPE (EMP_TYPE (N2 '  2 2, 'CLEMENTINE'))  
D3           300000  
D4           400000  EMPS_TYPE (EMP_TYPE ( 'N5'  2 5  'B ibi),  
                                EMP_TYPE ( 'N6'  2 6  'C here),  
                                EMP_TYPE ( 'N7'  2 7  'D idi),  
                                EMP_TYPE ( 'N8'  2 8, 'F ifi'))  
D5           400000  EMPS_TYPE (EMP_TYPE ( 'N5'  2 5  'B ibi),  
                                EMP_TYPE ( 'N8'  2 8, 'F ifi'))
```


Nested Tables (NESTED TABLE)

Editing (continued)

- Update from the main table as a predicate in the nested table

```
update departementdsetd. Budget = d. Budget + 777
      wheree exists (select from
      table ( select dt. employees from department dt
              Where dt. numdep =
                  d. numdep) nt
      Where nt. age <25) ;
```

Description:

- Query that returns the employees in each department

```
select dt. employees from departementdt
      Where dt. numdep = d. numdep
      e
```

- Condition on an attribute of the nested table:

```
Where nt. age <25
```

- Alias of the nested table: nt

Nested Tables (NESTED TABLE)

Editing (continued)

```
SQL> select          from department          ;
NUMDEP  BUDGET EMPLOYEES (NINSEE, age, name)
-----
D1      150777  EMPS_TYPE (EMP_TYPE ( 'N1'  2 1 'CLEMENT'))
D2      300777  EMPS_TYPE (EMP_TYPE (N2 '  2 2, 'CLEMENTINE'))
D3      300000
D4      400000  EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),
                        EMP_TYPE ( 'N6'  2 6 'C here),
                        EMP_TYPE (N7 '   2 7 'D idi),
                        EMP_TYPE ( 'N8 '  2 8, 'F ifi'))
D5      400000 EMPS_TYPE (EMP_TYPE ( 'N5 '  2 5 'B ibi),
                        EMP_TYPE ( 'N8 '  2 8, 'F ifi'))
```

Nested Tables (NESTED TABLE)

change

- Update from the main table as a predicate in the nested table

```
update departementd set      d. budget = D. Budget + 999
  where exists
    (select      from table
     (select      dt. employees from departementdt
      where dt. numdep = d. numdep) nt
     where nt. age > 25) ;
```

Nested Tables (NESTED TABLE)

change

```
SQL> select          from department          ;
NUMDEP  BUDGET EMPLOYEES (NINSEE, age, name)
```

D1	150777	EMPS_TYPE (EMP_TYPE ('N1'	2 1	'CLEMENT'))
D2	300777	EMPS_TYPE (EMP_TYPE (N2 '	2 2,	'CLEMENTINE'))
D3	300000			
D4	400999	EMPS_TYPE (EMP_TYPE ('N5'	2 5	'B ibi),
		EMP_TYPE ('N6'	2 6	'C here),
		EMP_TYPE (N7 '	2 7	'D idi),
		EMP_TYPE ('N8 '	2 8,	'F ifi'))
D5	400999	EMPS_TYPE (EMP_TYPE ('N5 '	2 5	'B ibi),
		EMP_TYPE ('N8 '	2 8,	'F ifi'))

Note: the same employees are in two departments

Nested Tables (NESTED TABLE)

change

- Update in the nested table

update

```
table (select d. employees from departementd
       Where d. numdep = D2 ') nt
set
  Where nt. ninsee = 'N2' ;
```


Nested Tables (NESTED TABLE)

change

```
SQL> select      from department      ;
NUMDEP  BUDGET EMPLOYEES (NINSEE, age, name)
-----
D1      150777  EMPS_TYPE (EMP_TYPE ( 'N1'  2 1 'CLEMENT'))
D2      300777  EMPS_TYPE (EMP_TYPE (N2 '  4 4 'CLEMENTINE'))
D3      300000
D4      400999  EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),
                        EMP_TYPE ( 'N6'  2 6 'C here),
                        EMP_TYPE (N7 '   2 7 'D idi),
                        EMP_TYPE ( 'N8 '  2 8, 'F ifi'))
D5      400999 EMPS_TYPE (EMP_TYPE ( 'N5 '  2 5 'B ibi),
                        EMP_TYPE ( 'N8 '  2 8, 'F ifi'))
```

Note: It is impossible to change several different records nested tables with a single UPDATE!

Nested Tables (NESTED TABLE)

suppression

Delete in the main table

```
delete      from department
           where numdep = 'D3';
e
```

```
SQL> select      from department      ;
```

```
NUMDEP      BUDGET EMPLOYEES (NINSEE, age, name)
```

```
-----
D1          150777      EMPS_TYPE (EMP_TYPE ( 'N1'  2 1 'CLEMENT'))
D2          300777      EMPS_TYPE (EMP_TYPE (N2 '  4 4 'CLEMANTINE'))
D4          400999      EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),
                        EMP_TYPE ( 'N6'  2 6 'C here),
                        EMP_TYPE (N7 '  2 7 'D idi),
                        EMP_TYPE ( 'N8 '  2 8, 'F ifi'))
D5          400999      EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),
                        EMP_TYPE ( 'N8 '  2 8, 'F ifi'))
```


Nested Tables (NESTED TABLE)

suppression

- Deleting from a value of the nested table Elimination of departments that employ a person whose name is FIFI

```
delete from departementd
  Where exists (select      from
table      (select      dt. employees from department      dt
              Where dt. numdep = d. numdep)      nt
  Where upper (Nt. Name)      like '%% FIFI');
```

Nested Tables (NESTED TABLE)

suppression

```
SQL> select          from department          ;
```

```
NUMDEP    BUDGET EMPLOYEES (NINSEE, age, name)
```

```
-----  
D1          150777 EMPS_TYPE (EMP_TYPE ( 'N1'          2 1 'CLEMENT'))  
D2          300777 EMPS_TYPE (EMP_TYPE (N2 '          4 4 'CLEMENTINE'))
```


Nested Tables (NESTED TABLE)

suppression

- Deleting a nested table

Elimination of departments that employ a person whose name CLEMENT

```
delet and reliable          (select      dt. employees
                             from departementdt where where      dt. numdep = D1 ')      nt
                             nt. name = 'CLEMENT';
```

```
SQL> select      from department      ;
NUMDEP      BUDGET EMPLOYEES (NINSEE, age, name)
```

```
-----
D1          150777
D2          300777 EMPS_TYPE (EMP_TYPE (N2 ', 4 4' CLEMENTINE '))
```

Nested Tables (NESTED TABLE)

interrogation

- *What are the numbers and employee names D4 department?*

```
SQL> select          from department          ;
```

```
NUMDEP    BUDGET EMPLOYEES (NINSEE, age, name)
```

```
-----  
D1          150777    EMPS_TYPE (EMP_TYPE ( 'N1'  2 1 'CLEMENT'))  
D2          300777    EMPS_TYPE (EMP_TYPE (N2 '  4 4 'CLEMENTINE'))  
D4          400999    EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),  
                                EMP_TYPE ( 'N6'  2 6 'C here),  
                                EMP_TYPE (N7 '  2 7 'D idi),  
                                EMP_TYPE ( 'N8 '  2 8, 'F ifi'))  
D5          400999    EMPS_TYPE (EMP_TYPE ( 'N5 '  2 5 'B ibi),  
                                EMP_TYPE ( 'N8 '  2 8, 'F ifi'))
```


Nested Tables (NESTED TABLE)

interrogation

```
select  nt. ninsee,      nt. name
         from table      (select  dt. employees from
                           Where dt. numdep = 'D4') nt;
```

department

dt

NINSEE

NAME

N5	B ibi
N6	C here
N7	D idi
N8	F ifi

Nested Tables (NESTED TABLE)

interrogation

- *What are the numbers and employee names D4 department with less than 26 years?*

```
SQL> select          from department      ;
```

```
NUMDEP  BUDGET EMPLOYEES (NINSEE, age, name)
```

```
-----  
D1      150777  EMPS_TYPE (EMP_TYPE ( 'N1'  2 1 'CLEMENT'))  
D2      300777  EMPS_TYPE (EMP_TYPE (N2 '  4 4 'CLEMENTINE'))  
D4      400999  EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),  
                                EMP_TYPE ( 'N6'  2 6 'C here),  
                                EMP_TYPE (N7 '  2 7 'D idi),  
                                EMP_TYPE ( 'N8 '  2 8, 'F ifi'))  
D5      400999 EMPS_TYPE (EMP_TYPE ( 'N5 '  2 5 'B ibi),  
                                EMP_TYPE ( 'N8 '  2 8, 'F ifi'))
```

Nested Tables (NESTED TABLE)

interrogation

```
select  nt. ninsee,      nt. name
        from table      (select      dt. employees from department      dt
        Where dt. numdep = 'D4') Where nt nt. age <2 6;
NINSEE      NAME
-----
N5           B ibi
```

Nested Tables (NESTED TABLE)

interrogation

- Query: What is the number of employees D4 department?

```
SQL> select          from department          ;
```

```
NUMDEP    BUDGET EMPLOYEES (NINSEE, age, name)
```

```
-----  
D1          150777    EMPS_TYPE (EMP_TYPE ( 'N1'  2 1 'CLEMENT'))  
D2          300777    EMPS_TYPE (EMP_TYPE (N2 '  4 4 'CLEMENTINE'))  
D4          400999    EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),  
                                EMP_TYPE ( 'N6'  2 6 'C here),  
                                EMP_TYPE (N7 '  2 7 'D idi),  
                                EMP_TYPE ( 'N8 '  2 8, 'F ifi'))  
D5          400999    EMPS_TYPE (EMP_TYPE ( 'N5 '  2 5 'B ibi),  
                                EMP_TYPE ( 'N8 '  2 8, 'F ifi'))
```


Nested Tables (NESTED TABLE)

interrogation

```
select COUNT( ) " Number of employees "  
      from table (select dt. employees  
      Where dt. numdep = 'D4') nt  
      from      department dt  
Number of employees  
-----  
4
```


Nested Tables (NESTED TABLE)

interrogation

- What are the numbers and names of employees of departments D1 and D2?

```
SQL> select          from department      ;
```

```
NUMDEP  BUDGET EMPLOYEES (NINSEE, age, name)
```

```
-----  
D1      150777  EMPS_TYPE (EMP_TYPE ( 'N1'  2 1 'CLEMENT'))  
D2      300777  EMPS_TYPE (EMP_TYPE (N2 '  4 4 'CLEMENTINE'))  
D4      400999  EMPS_TYPE (EMP_TYPE ( 'N5'  2 5 'B ibi),  
                                     EMP_TYPE ( 'N6'  2 6 'C here),  
                                     EMP_TYPE (N7 '  2 7 'D idi),  
                                     EMP_TYPE ( 'N8 '  2 8, 'F ifi'))  
D5      400999 EMPS_TYPE (EMP_TYPE ( 'N5 '  2 5 'B ibi),  
                                     EMP_TYPE ( 'N8 '  2 8, 'F ifi'))
```


Nested Tables (NESTED TABLE)

interrogation

```
select      select nt. ninsee,      nt. name from table
              (Selectdt. Employees from department      dt
              Where dt. numdep =      'D1')      nt

union
select      nt. ninsee, nt. name from table
              (Selectdt. Employees from department      dt
              Where dt. numdep =      'D2')      nt;
```

NINSEE	NAME
N1	CLEMENT
N2	CLEMENTINE

Several nested tables

Creation

Grouping nested tables and Teachers Training Course in the table

NUMC	title	professors		formations	
		Name	Specialty	spinner et	schedule

```
create      typeprof _ typeasobject
            (Name varchar 2 (3 0), specialty          varchar 2 ( 30 )
/
create      typeprofs __ kind typeastableofprof
/
create      typeformation _ typeasobject
            (Filierevarchar 2 (3 0), time              number (5))
/
create      typeformations typeastableofformatio _ n _ Type
/
```

Several nested tables

Creation

```
create      typeprof _ typeasobject
            (name      varchar 2 ( 30 ) , specialty          varchar 2 ( 30 ) )
/
create      typeprofs __ kind typeastableofprof
/
create      typeformation _ typeasobject
            (Filierevarchar 2 (3 0),          schedule      number (5))
/
create      typeformations typeastableofformatio _ n _ Type
/
create      typecours _ typeasobject
            (NUMC
              C          varchar 2 (5)  titrevarchar 2 (1 5)
              professeursprofs _ type formationsformations _ type)
/
CreateTable coursofcours _ kind
            (constraint      pk _ coursprimary      key (NUMC))
            nestedtableprofesseur sstoreastabprofs,
            nested      table formationsstoreastabf ormation;
```


Several nested tables

Insertion

Inserting an object in the Course table, without linking it to teachers or training

```
InsertInto coursvalues (BD 'Teacher Type ' ' Data base ' ,  
                        _ (), training _ the type ());
```

```
select      from course;
```

```
NUMC TITLE                FACULTY (NAME, SPECIALTY) FORMATIONS (DIE, SCHEDULE)
```

```
-----  
BD      B ases of PROFS_TYPE Data ()                FORMATIONS_TYPE ()
```


Several nested tables

Insertion

- Insert, with VALUES in 2 nested tables:

```
insert    into    course    gains    ( 'DW'    'Data    WareHouse '
```

```
Teacher Type _ ( _ prof template ( 'Clemence' BD '),
```

```
    prof _ deviation (Adam 'BD')),
```

```
    _-type formations (training _ template ( 'Master 1', 1 0 0),
```

```
        _ training template ( 'Master          SIA ' 2 0 0),
```

```
        training _ deviation (DEA          AIOC ' 2 0 0)    ));
```


Insertion

```
select      from      course;
```

NUMC TITLE	FACULTY (NAME, SPECIALTY)
------------	---------------------------

BD B ases of PROFS_TYPE Data ()

dwData WareHouse

```
FORMATIONS_TYPE (
```

```
PROF_TYPE ( 'Clemence' BD '),
1',
PROF_TYPE (Adam 'BD'))
2P'
```

FORMATIONS (DIE, SCHEDULE)

FORMATIONS_TYPE ()

PROFS TYPE (

FORMATION_TYPE ('M aster

```
FORMATION_TYPE ( 'Master
```

FORMATION_TYPE ('M aster 2R

Several nested tables

Insertion

Insert, with VALUES in 2 nested tables

```
insert into course gains 'B' ( 'BDA'  
                                ases of data A VANC ed es'
```

```
Teachers _ kind (Prof. _ template ( 'Mercy', 'BD'),  
                Prof. _ kind ( 'T raifor' BD '), Prof. _ the type ('  
                The Good ',' BD '))
```

```
training _ kind (  
    _ training template ( 'M aster 2P' 2 0 0),  
    _ training template ( 'M aster 2R', 2 0 0) ));
```


Several nested tables

Insertion

COURSE

NUMC	title	professors		formations	
		Name	Specialty	spinneret	schedule
BD DW	Data base DataWareHouse				
		Clemency	BD	master 1	100
		Adam	BD	Master 2P	200
				Master 2R	200
BDA	Advanced Data Bases	Clemency	BD	Master 2P	200
		Traifor	BD	Master 2R	200
		Good	BD		

The SQL chage a ffi + is very bad ...

select from course;

```

NUMC  TITLE                FACULTY (NAME, SPECIALTY)    FORMATIONS (DIE, SCHEDULE)
BD     Data base           PROFS_TYPE ()               FORMATIONS_TYPE ()
DW     Data WareHouse      PROFS_TYPE (PROF_TYPE ( 'Clemence' BD ' ) PROF_TYPE (Adam' BD '))
                                     FORMATIONS_TYPE (FORMATION_TYPE ( 'M aster
FORMATION_TYPE ( 'M aster  2P ' , 2 0 0 ), FORMATION_TYPE ( 'M aster 2R ' , 2 0 0 ))
BDA    Data base          A VANC ed es PROFS_TYPE (PROF_TYPE ( 'Mercy', 'BD'), PROF_TYPE (T raifor '
PROF_TYPE ( 'The Good', 'BD'))
                                     FORMATIONS_TYPE (FORMATION_TYPE ( 'M aster
FORMATION_TYPE ( 'M aster  2R ' , 2 0 0 ))

```


Several nested tables

Insertion

Insert, and with TABLE VALUES in 2 nested tables Data recording:
the Traifor Parisi and teachers teach BD

```
insert into table (Selectc. Teachers from coursec
Where
NUMC = BD ') gains (T raifor ' ' IF ' ) ;

insert into table (Selectc. Teachers from coursec
Where
NUMC = BD ') gains ( 'P Arisi' 'DM');

select from course;
```

NUMC	title	professors		formations	
		Name	Specialty	spinneret	schedule
BD	Data base	Traifor	IF		
		Parisi	DM		
DW	DataWareHouse	Clemency	BD	master 1	100
		Adam	BD	Master 2P	200
				Master 2R	200
BDA	Advanced Data Bases	Clemency	BD	Master 2P	200
		Traifor	BD	Master 2R	200
		Good	BD		

Several nested tables

Insertion

Insert, with TABLE and SELECT in 2 nested tables

The BD course must now be taught in all sectors concerned by DW material provided that it had a volume of less than 150 hours

```
insert into table (select c. training from course
                    Where c. NUMC = 'BD')
select nestedf. filiere, nestedf. schedule
from table (select c. training from course c
              Where c. NUMC = 'DW') nestedf
Where nestedf. Zone <150;
```

Several nested tables

Insertion

select tfrom course;

NUMC	title	professors		formations	
		Name	Specialty	spinneret	schedule
BD	Data base	Traifor	IF	INFO1	70
DW	DataWareHouse	Parisi	DM	master 1	
		Clemency	BD	master 1	100
		Adam	BD	Master 2P	200
				Master 2R	200
BDA	Advanced Data Bases	Clemency	BD	Master 2P	200
		Traifor	BD	Master 2R	200
		Good	BD		

Several nested tables

Modification - Example

In the Data WareHouse subject, Professor Adam is replaced by the Saitout professor and schedules for 2P Master increase of 30%

update table (Selectc. Where teachers from course v. Title = 'Data WareHouse') nestedprfnestedprf. name = 'S aitou' where nestedprf. name = set Adam ';

updatet reliable (select c. training **from** course
Where c. title = 'Data WareHouse') nestedfrm
set nestedfrm. time = time 1.3
Where nestedfrm. filiere **like** 'M aster 2P%';

Several nested tables

change

NUMC	title	professors		formations	
		Name	Specialty	spinneret	schedule
BD	Data base	Traifor	IF	INFO1	70
		Parisi	DM	master 1	
DW	DataWareHouse	Clemency	BD	master 1	100
		Saitou	BD	Master 2P	260
BDA	Advanced Data Bases			Master 2R	200
		Clemency	BD	Master 2P	200
		Traifor	BD	Master 2R	200
		Good	BD		

Several nested tables

Modification - Example

explanations:

- Editing using the UPDATE of one or more attributes in one of two nested tables in the current table
- Changing a teachers and training as part of a given material: two distinct UPDATE requests
(Because the two nested tables are involved)
- Requires the use of an alias to identify the object in the nested table

Several nested tables

change

*For DW material, replacing the Master1 die through the die MASTER 2
and recording an hourly volume of 150 hours*

```
update table (Select  c. training      from coursec
                  Where c. Ti be = 'Data WareHouse') nestedfrm
set      nestedfrm. time = 1 5 0,      nestedfrm. filiere = 'MASTER 2'
Where nestedfrm. filiere = 'M aster 1';
```


Several nested tables

change

NUMC	title	professors		formations	
		Name	Specialty	spinneret	schedule
BD	Data base	Traifor	IF	INFO1	70
		Parisi	DM	master 1	
DW	DataWareHouse	Clemency	BD	MASTER 2	150
		Saitou	BD	Master 2P	260
BDA	Advanced Data Bases			Master 2R	200
		Clemency	BD	Master 2P	200
		Traifor	BD	Master 2R	200
		Good	BD		

Several nested tables

suppression

Deleting in 2 nested tables

Professor Parisi no longer teaches comics material. Recording this information

DeleteTable (Selectc. Teachers

c.
from coursc **Where** NUMC = 'BD') nt
Where nt.name = 'P Arisi';

NUMC	title	professors		formations	
		Name	Specialty	spinneret	schedule
BD	Data base	Traifor	IF	INFO1	70
DW	DataWareHouse			master 1	
		Clemency	BD	MASTER 2	100
		Saitou	BD	Master 2P	260
				Master 2R	200
BDA	Advanced Data Bases	Clemency	BD	Master 2P	200
		Traifor	BD	Master 2R	200
		Good	BD		

Several nested tables

suppression

Deleting in 2 nested tables

*The sector includes not Master1 comics material in its curriculum.
Recording this information*

DeleteTable (Selectc. Training
from Where course c. NUMC = BD ') nt
Where nt. filiere = ' Master 1 ' ;

NUMC	title	professors		formations	
		Name	Specialty	spinneret	schedule
BD	Data base	Traifor	IF	INFO1	70
DW	DataWareHouse	Clemency	BD	MASTER 2	100
		Saitou	BD	Master 2P	260
				Master 2R	200
BDA	Advanced Data Bases	Clemency	BD	Master 2P	200
		Traifor	BD	Master 2R	200
		Good	BD		

Several levels of nesting

NUMC	title	professors		formations		
		Name	Specialty	spinneret	schedule	dates
BD	Data base	Traifor	IF	INFO1	70	Day
DW	DataWareHouse	Clemency	BD	MASTER 2	100	
		Saitou	BD	Master 2P	260	
				Master 2R	200	
BDA	Advanced Data Bases	Clemency	BD	Master 2P	200	
		Traifor	BD	Master 2R	200	
		Good	BD			

Oracle 8 does not allow to install several nesting levels in an object-relational table

? ? in Oracle 9i and / or 10g? ?

Pre-sized arrays (VARRAY)

- VARRAY (varying ARRAY) ordered collection and limited items of its type
- If the maximum number of items contained in a nested table is known a priori possibility of using a VARRAY type of table instead of a nested table
- Example: storage of up to 3 telephone numbers per teacher

professors:

Nump	pname	Address				phones
		AdrNum	AdrRue	AdrVille	AdrCP	TEL NUM B R

Pre-sized size arrays (VARRAY)

Example

Storage of up to 3 telephone numbers per teacher

• Creation

```
create      AA stands _ typeasobject
            (AdrNum varchar 2 (1 0), AdrNom          varchar 2 ( 30 ) ,
            A V dr illev varchar 2 (2 0),          AdrCP          varchar 2 (5))
/
create      typetel _ typeasobject          (TEL NUMBR varchar 2 (2 0))
/

create      typetels typeasvarray _ (3) _ OFTEL Type
/

create      typeprofesseur _ typeasobject
(NUMP varchar 2 (5)          pname varchar 2 (2 0),
            A stands AA stands _ type _ elephonestels T type)
/
CreateTable professeursofprofesse kind heart _
            (constraint          pk _ professeursprimary key (NUMP));
```


Pre-sized size arrays (VARRAY)

- Insert: INSERT with VALUES

Storage Professor 3 type with no object, respectively, three and two telephone numbers (VARRAY of records)

```
insert into professors gains ('P1', 'Mercy'
AA stands _ kind (7, 'Avenue the Peace ', 'P aris' '75009'),
such kind _ ());

insert into professors gains ('P2', Adam '
'Stre th
AA stands _ kind (7 7 et of e freedom ', 'P aris' '75015'),
such _ type (such kind _ ('01 53 80 07 99 '),
such kind _ ('06 14 56 07 06 '),
such kind _ ('01 49 40 07 40 '));

insert into professeursvalues AA stands _ ('P3 ', Saitou '
Type (1, '_ Rue such type (such _ del al ibert é ', 'P aris', '75015'),
template (' _ 01 such deviation (
'06 14 56 14
53 80 53 80 '),
77 '), NULL));
```


Pre-sized size arrays (VARRAY)

Nump	pname	Address				phones
		AdrNum	AdrRue	AdrVille	AdrCP	TEL NUMBR
P1	Clemency	77	Avenue of peace	Paris	75009	NULL
						NULL
						NULL
P2	Adam	7	Liberty Street	Paris	75015	01 53 80 07 99
						06 14 56 07 06
						01 49 40 07 40
P3	Saitou	1	Liberty Street	Paris	75015	01 53 80 53 80
						06 14 56 14 77
						NULL

Pre-sized size arrays (VARRAY)

- Insert: INSERT into a VARRAY with PL / SQL
 - | With VARRAY tables, the operator is not operational TABLE (Version 8 Oracle - to check on the V9 and v10g)
 - | To manipulate the tables, it is necessary to use a program PL / SQL

DECLARED

```
new _ such such kind _ := _ such type (such kind _ ( '01 55 55 55 55'),  
                                         such kind _ ( '06 06 98 98 98'),  
                                         such kind _ ( '01 40 40 40 40'));
```

BEGIN

```
update professors
```

```
set Phones = new _ such
```

```
Where NUMP = 'P1';
```

END;

/

Pre-sized size arrays (VARRAY)

Nump	pname	Address				phones
		AdrNum	AdrRue	AdrVille	AdrCP	TEL NUMBR
P1	Clemency	77	Avenue of peace	Paris	75009	01 55 55 55 55
						06 06 98 98 98
						01 40 40 40 40
P2	Adam	7	Liberty Street	Paris	75015	01 53 80 07 99
						06 14 56 07 06
						01 49 40 07 40
P3	Saitou	1	Liberty Street	Paris	75015	01 53 80 53 80
						06 14 56 14 77
						NULL

Note :

- Insert a single telephone number for Professor P1 and place it in 2nd place in the table Phones

Writing following the trial of a ff assignment:

```
new _ such _ such kind: = _ such kind (NULL,
      such kind _ ( '06 06 98 98 98'), NULL);
```

Conclusion

Comparing NESTED TABLE and VARRAY

- A check by Oracle versions
- Ability to define an index in a NESTED TABLE
 - The number of elements is not limited in a nested table
- No ability to set index in a VARRAY
 - The number of elements is limited in a pre-sized table
- Ability to directly access records stored in both data structures
 - functions: EXISTS, FIRST, LAST, etc.
- Performance? : NestedTable > Varray