```
function TREE-SEARCH( problem, strategy) returns a solution, or failure
  initialize the search tree using the initial state of problem
  loop do
       if there are no candidates for expansion then return failure
       choose a leaf node for expansion according to strategy
       if the node contains a goal state then return the corresponding solution
       else expand the node and add the resulting nodes to the search tree
```

```
function GRAPH-SEARCH(problem, fringe) return a solution, or failure
          inge ← INSERT(MAKE-NODE(INITIAL-STATE[problem]), fringe)
           if fringe is empty then return failure
           node \leftarrow REMOVE\text{-FRONT}(fringe)
if GOAL-TEST(problem, STATE[node]) then return node
           if STATE[node] is not in closed then
                  \begin{array}{l} \operatorname{add} \operatorname{STATE}[node] \ \operatorname{to} \ closed \\ \operatorname{\mathbf{for}} \ child\text{-}node \ \operatorname{in} \ \operatorname{Expand}(\operatorname{STATE}[node], \ problem) \ \operatorname{\mathbf{do}} \end{array}
                          fringe \leftarrow INSERT(child-node, fringe)
                  end
```

DFS, m level, b children of one node,

Time O(b^m) Space O(bm)

Is it optinal? No, it finds the "leftmost" solution, regardless of depth or cost

BFS time O(bs) s is the depth of solution

Space O(b^s), ans it is complete, it is optimal if costs are all

UCS = uniform cost search

Is guaranteed to return an optimal path

Expand a cheapest node first, akes time $O(b^{C^*/\varepsilon})$, If that solution costs C^* and arcs cost at least ε , then the "effective depth" is roughly C^*/ε , space is the same, Processes all nodes with cost less than cheapest solution!

Search Heuristics

Admissible: $0 \le h(n) \le h^*(n)$

h*(n) is the true cost to a nearest goal

Consistency: heuristic "arc" cost ≤ actual cost for each

 $arc h(A) - h(C) \le cost(A to C)$

Greedy Search-best first search

Expand a mode that you think is closest to a goal state A* Search

Admissible heuristics never underestimate? False 不会 高估 cost

Uniform-cost =g(n), greedy-cost= h(n)

A* Search orders by the sum: f(n) = g(n) + h(n), stop: when we de-queue a goal

The different between a* and ucs is that a* add the heuristics which will intend to lead the agent to the goal with a purpose

Any complete search algorithm must have exponential (or worse) time complexity. True! (why?) BFS always returns the optimal cost path. False (in

Before expanding a node, check to make sure its state has never been expanded before

If not new, skip it, if new add to closed set

Important: store the closed set as a set, not a list

Constraint Search problem

Backtracking Search: Depth-first search with these two improvements is called backtracking search (not the best name) Idea 1: One variable at a time Idea 2: Check constraints as you go

Improvement: Arc consistency & forward checking

```
function Backtracking-Search(csp) returns solution/failure
    return Recursive-Backtracking({ }, csp)
function Recursive-Backtracking (assignment, csp) returns soln/failure
   \label{eq:constraint} \begin{split} & \text{if } assignment \text{ is complete then } \mathbf{return} \text{ } assignment \\ & var \leftarrow \mathbf{Select-Unassigned-Variable}(Variables[csp], assignment, csp) \end{split}
    for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do
         if \mathit{value} is consistent with \mathit{assignment} given \mathtt{Constraints}[\mathit{csp}] then
               add {var = value} to assignment
                result \leftarrow Recursive-Backtracking(assignment, csp)
               if result \neq failure then return result
               \mbox{remove } \{var = value\} \mbox{ from } assignment
    return failure
```

```
function AC-3( csp) returns the CSP, possibly with reduced domains
   \mathbf{inputs}:\ \mathit{csp},\ \mathsf{a}\ \mathsf{binary}\ \mathsf{CSP}\ \mathsf{with}\ \mathsf{variables}\ \{X_1,\ X_2,\ \dots,\ X_n\}
   {f local\ variables:}\ queue,\ {f a}\ {f queue},\ {f a}\ {f queue} and {f csp}
   while queue is not empty do
       (X_i, X_i) \leftarrow \text{Remove-First}(queue)
       if Remove-Inconsistent-Values(X_i, X_j) then
          for each X<sub>l</sub> in Neighbors[X<sub>l</sub>] do
             add (X_k, X_i) to queue
function Remove-Inconsistent-Values (X_i, X_j) returns true iff succeeds
   removed \leftarrow false
   for each x in Domain[X_i] do
       if no value y in \mathrm{DOMAIN}[X_j] allows (x,y) to satisfy the constraint X_i \leftrightarrow X_j
          then delete x from DOMAIN[X_i]; removed \leftarrow true
```

Runtime: $O(n^2d^3)$, can be reduced to $O(n^2d^2)$ but detecting all possible future problems is NP-hard ARC consistency- also called 2-consistency-a kind of binary consistency can affect any assigned variable in the graph that has a path to the assigned variable Limitation: 1, can have one solution left, 2, have multiply solution left. 3, none

Ordering: choose the MRV(Minimum(choose value) remaining values), Choose the variable with the fewest legal left values in its domain, also called most constrained variable, fall-fast ordering (which variable should be assigned next) Ordering: Least Constraining Value (choose node)

Given a choice of variable, choose the least constraining

(in what order should its values be tired)

Iterative Improvement & local search: try incorrect solution and modify it to be correct.

Algorithm: while not solved:

randomly select any conflicted variable Value selection: min-conflicts heuristic:

Choose a value that violates the fewest constraint

I.e., hill climb with h(n) = total number ofviolated constraints

local search: improve a single option until you can't make it better (no fringe!) 1, start whatever 2, repeat, move to the best neighboring state 3, if no better neighbors than current, return and quit

def max-value(state): def min-value(state):

initialize v = -∞

for each successor of state:

v = max(v, value(successor)) : v = min(v, value(successor))

value(successor))

return v def value(state):

if the state is a terminal state: return the state's utility if the next agent is MAX: return max-value(state)if the

next agent is MIN: return min-value(state) time,

space :DFS(without depth limited) Alpha-Beta Implementation

α: MAX's best option on path to root β: MIN's best option on path to root

def max-value(state, α , β):

initialize v = -∞

for each successor of state:

 $v = max(v, value(successor, \alpha, \beta))$

if $v \ge \beta$ return v

 $\alpha = \max(\alpha, v)$

return v

def min-value(state, α , β):

for each successor of state:

 $v = min(v, value(successor, \alpha, \beta))$

if $v < \alpha$ return v

 $\beta = \min(\beta, v)$

return v

Time complexity drops to O(b^{m/2})

Markov Decision Processes- MDPs are non-deterministic search problems One difference between MDP and expectimax: rewards are smeared throughout the tree(exit-max) instead of at the end.(MDP)

The value (utility) of a state s:

V (s) = expected utility starting in s and acting optimally

The value (utility) of a q-state (s,a):

Q (s,a) = expected utility starting out having taken action a from state s and (thereafter) acting optimally (bellman equation)

$$\begin{split} & V^*(s) = \max_{a} Q^*(s, a) \\ & Q^*(s, a) = \sum_{s'} T(s, a, s') \left[R(s, a, s') + \gamma V^*(s') \right] \\ & V^*(s) = \max_{a} \sum_{s'} T(s, a, s') \left[R(s, a, s') + \gamma V^*(s') \right] \end{split}$$

T 是概率 V like MAX and Q like EXP. But now incorporates rewards and discounts. transition function

Optimal action: action with highest Q value $V_{k+1}(s) \leftarrow \max_{a} \sum T(s, a, s') \left[R(s, a, s') + \gamma V_k(s') \right]$

Complexity of each iteration: O(S²A) How do we know the V_k vectors are going to converge?

Case 2: If the discount is less than 1 So as k increases, the values converge $V^{\pi}(s) = \sum T(s, \pi(s), s') [R(s, \pi(s), s') + \gamma V^{\pi}(s')]$

Utilities for a Fixed Policy calculate the V's for a fixed policy π (Policy Evaluation)

$$V_{k+1}^{\pi}(s) \leftarrow \sum_{s'} T(s, \pi(s), s') [R(s, \pi(s), s') + \gamma V_k^{\pi}(s')]$$

efficiency: O(S²) per iteration problem:1, slow 2, the max at each state rarely changes 3, the policy often converges long before the values policy iteration: Evaluation

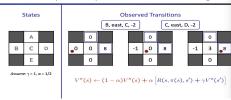
$$V_{k+1}^{\pi_i}(s) \leftarrow \sum_{s'} T(s, \pi_i(s), s') \left[R(s, \pi_i(s), s') + \gamma V_k^{\pi_i}(s') \right]$$

improvement: For fixed values, get a

better policy using policy extraction $\pi_{i+1}(s) = \arg\max_{a} \sum_{s} T(s, a, s') \left[R(s, a, s') + \gamma V^{\pi_i}(s') \right]$

Example: Temporal Difference **Learning** (model base learning)

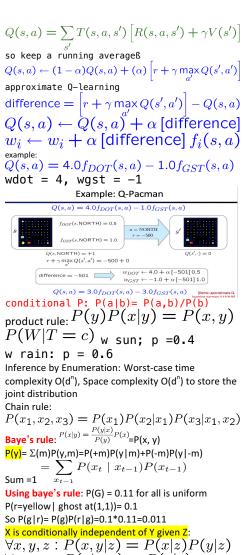
Example: Temporal Difference Learning



Value learning: watch some agent's performance under a policy and see how well it does Q-learning: watch yourself perform, possibly suboptimally, and still derive optimal

policies/values/etc. $V^{\pi}(s) \leftarrow (1 - \alpha)V^{\pi}(s) + \alpha |R(s, \pi(s), s') + \gamma V^{\pi}(s')|$

if we want to turn values into a (new) policy, we're sunk: $\pi(s) = \arg\max_{a} Q(s, a)$



$$\forall x, y, z : P(x, y|z) = P(x|z)P(y|z)$$

$$\forall x, y, z : P(x|z, y) = P(x|z)$$

Joint Distribution of a Markov Model now only related to past, future only related: ح

 $P(X_1, X_2, X_3, X_4) = P(X_1)P(X_2|X_1)P(X_3|X_2)P(X_4|X_3)$ $X_3 \perp \!\!\! \perp X_1 \mid X_2 \mid X_4 \perp \!\!\! \perp X_1, X_2 \mid X_3$ Stationary Distributions:最终 convegrent:

 $P_{\infty}(X) = P_{\infty+1}(X) = \sum P(X|x)P_{\infty}(x)$

HMM: evidence 're only conditionally independent given the hidden state

Quiz: does this mean that evidence variables are guaranteed to be independent?

No, they tend to correlated by the hidden state] Joint distribution of a HMM:

 $P(X_1, E_1, X_2, E_2, X_3, E_3) = P(X_1)P(E_1|X_1)P(X_2|X_1)P(E_2|X_2)P(X_3|X_2)P(E_3|X_3)$ $P(X_1, E_1, \dots, X_T, E_T) = P(X_1)P(E_1|X_1) \prod P(X_t|X_{t-1})P(E_t|X_t)$

inference: base case:

$$P(X_1|e_1)$$

$$P(x_1|e_1) = P(x_1, e_1)/P(e_1)$$

$$\propto_{X_1} P(x_1, e_1)$$

$$= P(x_1)P(e_1|x_1)$$

 $= \sum_{x_t}^{-\text{r-X}/-\text{r-}(X_t \mid e_{1:t}),} \text{we} \\ = \sum_{x_t}^{-\text{r-}(X_t \mid e_{1:t}),} P(X_{t+1}, x_t \mid e_{1:t})$ if we have e_1 to e_t and the $B(X_t)=P(X_t|e_{1:t})$, we reason

about : P(Xt+1|e1:t)=
$$x_t$$

= $\sum_{x_t} P(X_{t+1}|x_t)P(x_t|e_{1:t})$

or we can say:

$$B'(X_{t+1}) = \sum P(X'|x_t)B(x_t)$$

this is the guess part

the following is the observation part:

we get the believe for Xt+1: $B'(X_{t+1}) = P(X_{t+1}|e_{1:t})$

then the evidence comes in

 $\frac{P(X_{t+1}|e_{1:t+1})}{P(X_{t+1},e_{t+1}|e_{1:t})/P(e_{t+1}|e_{1:t})}$ $\propto_{X_{t+1}} P(X_{t+1}, e_{t+1}|e_{1:t})$ $= P(e_{t+1}|e_{1:t}, X_{t+1})P(X_{t+1}|e_{1:t})$ $= P(e_{t+1}|X_{t+1})P(X_{t+1}|e_{1:t})$

Unlike passage of time, we have to renormalize Basic idea: beliefs "reweighted" by likelihood of evidence, and compactly:

 $B(X_{t+1}) \propto_{X_{t+1}} P(e_{t+1}|X_{t+1})B'(X_{t+1})$

the forward algorithm

we want to know $B_t(X) = P(X_t|e_{1:t})$ we can derive the following updates:

 $P(x_t|e_{1:t}) \propto_X P(x_t,e_{1:t})$ $=\sum P(x_{t-1},x_t,e_{1:t})$ $= \sum_{t=0}^{x_{t-1}} P(x_{t-1}, e_{1:t-1}) P(x_{t}|x_{t-1}) P(e_{t}|x_{t})$ $= P(e_t|x_t) \sum P(x_t|x_{t-1})P(x_{t-1},e_{1:t-1})$

SUMMAY:

We reason that $P(x_t|e_{1:t-1}) = \sum_{x_{t-1}} P(x_{t-1}|e_{1:t-1}) \cdot P(x_t|x_{t-1})$

Get evidence: we get: (update)

 $P(x_t|e_{1:t}) \propto_X P(x_t|e_{1:t-1}) \cdot P(e_t|x_t)$ Particle filter:

Step 1: base on the evidence-> guess what will happen x' = sample(P(X'|x))

Step 2: observe: w(x) = P(e|x)

 $B(X) \propto P(e|X) B'(X)$ As before, the probabilities don't sum to one, since all have been down weighted Step 3: resample: normalize

HMMS: most likely explanation: forward algorithm (Sum)

HMMOS most macry capacitans arg max $P(x_{1:t}|e_{1:t})$ $f_t[x_t] = P(x_t, e_{1:t})$ viterbi Algorithm(max): $m_t[x_t] = \max_{x_{1:t-1}} P(x_{1:t-1}, x_t, e_{1:t})$ $P(e_t|x_t) \max_{x_{t-1}} P(x_t|x_{t-1}) m_{t-1}[x_{t-1}]$

which state sequence $X_{1:T}$ is most likely given the evidence e 1:t?

 $x_{1:T}^* = \underset{x_1 \cdot x}{\arg \max} P(x_{1:T} | e_{1:T}) = \underset{x_1 \cdot x}{\arg \max} P(x_{1:T}, e_{1:T})$

Bayes Net: to see what probability a BN gives to a full assignment: 假设 x 只与 parent (x) 有关

 $P(x_1, x_2, \dots x_n) = \prod P(x_i | parents(X_i))$

 $P(x_i|x_1,\ldots x_{i-1}) = P(x_i|parents(X_i))$ How big is an N-node net if nodes have up to k parents? $O(N * 2^{k+1})$

General question: in a given BN, are two variables independent (given evidence)?

X->Y->Z, Guaranteed X independent of Z? No! Guaranteed X independent of Z given Y? yes! Y->X, Y->Z. Guaranteed X independent of Z? No! Guaranteed X independent of Z given Y? ves! Evidence along the chain "blocks" the influence X->Z, Y->Z. Are X and Y independent? Yes

Are X and Y independent given Z? NO: seeing traffic puts the rain and the ballgame in competition as explanation. Active means no independent (dependent)

A trail X1-...-Xn is active if : it has no vstructures Xi−1→Xi←Xi+1

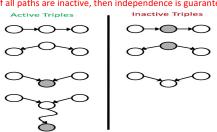
A trial X1-...-Xn is active given Z if:

1, for any v-structure Xi-1→Xi←Xi+1we have that Xi or one of its descendants ∈ Z 2, no other Xi is in Z

如果两个节点间没有 Active trial,则这两个节点是 dseparated 的 也就是独立的! Independence

 $X_i \perp \!\!\! \perp X_j | \{X_{k_1}, ..., X_{k_n}\}$? Given means Block (DARKED)

Check all (undirected!) paths between Xi and Xj If one or more active path, then independence not guaranteed If all paths are inactive, then independence is guaranteed



Decision Networks: choose the action which maximizes the expected utility given the evidence = EU(ACTION) = $EU(leave) = \sum P(w)U(leave, w)$

MEU(X)= max(EU for all)在当你获得新的 evidence 之 后,you update you believe (probability), but the value of Utility doesn't change! Evidence kinds of forecast: P(rain | forecast = bad) = 0.66, P(sun | forecast = bad) = 0.34, recalculate the EU, and MEU

Value of information kinds of evidence = MEU (after we get the information)- MEU(before we get the information)

$$\begin{aligned} & \text{VPI}(E'|e) = \left(\sum_{e'} P(e'|e) \text{MEU}(e,e')\right) - \text{MEU}(e) \\ & \forall E', e : \text{VPI}(E'|e) > 0 \end{aligned}$$

$$VPI(E_j, E_k|e) \neq VPI(E_j|e) + VPI(E_k|e)$$
 $VPI(E_j, E_k|e) = VPI(E_j|e) + VPI(E_k|e, E_j)$

$$= VPI(E_k|e) + VPI(E_j|e, E_k)$$

If Parents(U) $\prod Z$ | CurrentEvidence Then VPI(Z | CurrentEvidence) = 0