

1. (a) $\text{Rank}(A) = 3$

$$\text{Rank}(A|y) = 3$$

Since $\text{Rank}(A) = \text{Rank}(A|y)$

There is a solution $Ax=y$

(b) Pick 3 columns out of 5.

$$\{a_1, a_2, a_3\} : S_1$$

$$\{a_1, a_3, a_5\} : S_2$$

$$\{a_2, a_3, a_4\} : S_3$$

Results are:

$$\begin{pmatrix} 1 \\ -7 \\ 1 \\ 0 \\ 0 \end{pmatrix}_{S_1}, \begin{pmatrix} 8 \\ 0 \\ -1 \\ 0 \\ -7 \end{pmatrix}_{S_2}, \begin{pmatrix} 0 \\ -7 \\ \frac{4}{3} \\ \frac{1}{3} \\ 0 \end{pmatrix}_{S_3}$$

$$\lambda_1 + \lambda_2 + 1$$

$$\lambda_1 S_1 + \lambda_2 S_2 + (1 - (\lambda_1 + \lambda_2)) S_3$$

$$\begin{pmatrix} \lambda_1 + 8\lambda_2 \\ -7\lambda_2 - 7 \\ -\frac{1}{3}\lambda_1 - \frac{1}{3}\lambda_2 + \frac{4}{3} \\ -\frac{1}{3}\lambda_1 + \frac{1}{3}\lambda_2 + \frac{1}{3} \\ -7\lambda_2 \end{pmatrix}$$



1. (c) Output From Matlab:

```
>> IE531_hw1
```

```
A =
```

```
2  0  0  6  2
1 -1 -1  4  0
2 -2  2  4  0
2 -2  0  6  0
-4  4 -8 -4  0
4  0 -2 14  4
```

```
y =
```

```
2
7
18
16
-40
2
```

```
x =
```

```
8 - 4*b - 7*a
- 7*a - 7*b
1 - b
-b
7*a + 7*b - 7
```

```
A*x == y gives
```

```
ans =
```

```
2 == 2
7 == 7
18 == 18
16 == 16
-40 == -40
2 == 2
```

```
>>
```

1.(d) The solution (1) is verified by 1(c) already and the Matlab output of solution (2) is given below:

```
>> IE531_hw1
```

A =

```

2   0   0   6   2
1  -1  -1   4   0
2  -2   2   4   0
2  -2   0   6   0
-4   4  -8  -4   0
4   0  -2  14   4

```

y =

```

2
7
18
16
-40
2

```

x =

```

      a
    - 7*a - 7*b
11/3 - (7*b)/3 - (8*a)/3
 8/3 - (7*b)/3 - (8*a)/3
    7*a + 7*b - 7

```

A*x == y gives

ans =

```

2 == 2
7 == 7
18 == 18
16 == 16
-40 == -40
2 == 2

```

1.(e) Since the two solutions are linearly related. Knowing one indicates the other.

$$a = -7\lambda_1 - 4\lambda_2 + 8$$

$$b = 8\lambda_1 + 5\lambda_2 - 8$$

1.(d) After adding the two constraints, the rank of matrix of A is 5. Also, A has 5 columns. Hence, the matrix A has full column rank. It indicates that there will be no free variable and there can be only one solution to the combined set of equations.

$$2.(a) \quad f(x) = \begin{cases} x & 0 \leq x \leq 1 \\ 2-x & 1 < x \leq 2 \\ 0 & \text{otherwise} \end{cases}$$

Let $u_i \in \text{i.i.d. } [0, 1]$

Find CDF $F(x)$, inverse transform $x_i = F^{-1}(u_i)$

Then, x_i is i.i.d. r.v's

$$c_b). \quad F^{-1}(y) = \begin{cases} 1 + \sqrt{2y} & 0 \leq y \leq 0.5 \\ 1 - \sqrt{2-2y} & 0.5 < y \leq 1 \end{cases}$$

$$F(x) = \begin{cases} \frac{1}{2}(x^2 - 2x + 1) & 0 \leq x \leq 1 \\ \frac{1}{2}(-x^2 + 2x + 1) & 1 < x \leq 2 \\ 1 & \text{otherwise} \end{cases}$$

$$f(x) = \begin{cases} x-1 & 0 \leq x \leq 1 \\ -x+1 & 1 < x \leq 2 \\ 0 & \text{otherwise} \end{cases}$$

$$c_c). \quad F^{-1}(y) = \begin{cases} 2 - \sqrt{1-2y} & 0 \leq y \leq 0.5 \\ \sqrt{2y-1} & 0.5 < y < 1 \end{cases}$$

$$F(x) = \begin{cases} \frac{1}{2}(-x^2 + 4x - 3) & 1 \leq x \leq 2 \\ \frac{1}{2}(x^2 + 1) & 0 \leq x < 1 \\ 1 & \text{otherwise} \end{cases}$$

$$f(x) = \begin{cases} x & 0 \leq x < 1 \\ -x+2 & 1 \leq x \leq 2 \\ 0 & \text{otherwise} \end{cases}$$



3. (a) The program (I) outputs "kablooey"
since it prints the current character and
then iterates itself until the end.

The program (II) outputs "yeoolbak"
since it iterates itself until the end and
then prints the string in reverse order.

(b). Given two number (m, n) .

If $m < n$, the algorithm calls $\text{gcd}(m, n-m)$

If $m > n$, the algorithm calls $\text{gcd}(m-n, n)$

GCD of two numbers do not change if the
smaller number is subtracted from the large number.

Hence, $\text{gcd}(m, n-m) = \text{gcd}(m, n)$ and

$\text{gcd}(m-n, n) = \text{gcd}(m, n)$. After running the
recursive function gcd , it will finally return
 n , when $n=m$. It works for any pair of numbers.

(c) $\text{Putchar}(1234 \% 10 + 'A') = \text{Putchar}(4 + 'A')$

E is printed, $1234 / 10 = 123$

$\text{Putchar}(123 \% 10 + 'A') = \text{Putchar}(3 + 'A')$

D is printed, $123 / 10 = 12$

$\text{Putchar}(12 \% 10 + 'A') = \text{Putchar}(2 + 'A')$

C is printed, $12 / 10 = 1$

$\text{Putchar}(1 \% 10 + 'A') = \text{Putchar}(1 + 'A')$

B is printed.

Hence EDCB is the output.



(d) The algorithm is not correct and will cause illegal moves. When $n=5$,

The top disk goes from A to D, <1>

A to B, <2>

A B C D

D to B <3>

A to B <4>

The fourth move is illegal.

$n-k$ disks stay in B for destination from the second step while the third step let $k-1$ disk move to B. It causes illegal move.

