Assignment 13, Discrete Mathematics

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1 Problems

1.1 Problem 1

Demonstrate five distinct sets: $A, B, A \cup B, A \oplus B, A \setminus B$, all of which have the same cardinality.

1.1.1 Answer 1

There's a popular programming puzzle called "FizzBuzz". The taks is to print natural numbers according to the rule: if a number is divizible by three, the program prints fizz, if the number is divisible by five the program prints buzz (when both condition holds, the program thus prints fizzbuzz, hence the name). Otherwise the program prints the number itself.

It is easy to see that the description above perfectly matches the requirement:

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A = \{x \in \mathbb{N} \mid x \bmod 3 = 0\}
B = \{x \in \mathbb{N} \mid x \bmod 5 = 0\}
A \cup B = \{x \in \mathbb{N} \mid x \bmod 3 = 0 \lor x \bmod 5 = 0\}
A \oplus B = \{x \in \mathbb{N} \mid (x \bmod 3 = 0 \lor x \bmod 5 = 0) \land \neg (x \bmod 15 = 0)\}
A \setminus B = \{x \in \mathbb{N} \mid x \bmod 3 = 0 \land \neg (x \bmod 5 = 0)\}.
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The cardinality of all sets is \aleph_0 and neither one of them is equal to any other set.

1.2 Problem 2

- 1. Let n be a natural number. Show that the set of all subsets of $\mathbb N$ of cardinality exactly n is a countable set.
- 2. Show that the set of all finite subsets of \mathbb{N} is countable.

- 3. Show that the set of all infinite subsets of \mathbb{N} is not countable.
- 4. Find the cardinality of the set given in (3).
- 5. $\aleph_0 = \big|\{X \in P(\mathbb{N})||X| = n\}\big|$ is the restatement of (1). Write a similar statement for (2) and (4).

1.2.1 Answer 2

We can construct the set of all subsets of \mathbb{N} by repeating cartesian product n times. Since $|\mathbb{N} \times \mathbb{N}| = \aleph_0$ so is the $|\mathbb{N}^n| = \aleph_0$ (because n is finite). Note that the requested sets in this case are given by the resulting ordered pairs. I.e. $(x_1, (x_2, \ldots (x_{n-1}, x_n) \ldots) \in \mathbb{N}^n$. In fact, we need only the unique tuples, of which there are even less.

1.2.2 Answer 3

This continues the argument from 1.2.1, but constructs the sets in question using recursive definition: $\bigcup_{i=1}^{n} \mathbb{N}^{i}$ since there are no infinite exponents involved in this construction, it must describe a countable set.

1.2.3 Answer 4

The set of all infinite sets can be described as $\mathbb{N}^{\mathbb{N}}$. We can substitute \mathbb{N} for any infinite subset because any infinite subset of \mathbb{N} is equivalent to \mathbb{N} , which follows from definition of infinite set). Because set exponentiation is non-decreasing in both arguments, this set has cardinality of at least $2^{\mathbb{N}}$, which is the cardinality of the powerset, which, by Cantor's theorem is infinitely larger than any countable set.

1.2.4 Answer 5

We just stated that the requested set is equivalent to $\mathbb{N}^{\mathbb{N}}$ in 1.2.3. Hence its cardinality is at least \mathfrak{c} and no greater than $2^{\mathfrak{c}}$. Assuming (here and through

the rest of the answers) continuity hypothesis is true, the cardinality of the set in question is \mathfrak{c} .

1.2.5 Answer 6

- 1. can be given by: $\aleph_0 = |\{X \in P(\mathbb{N}), n \in \mathbb{N} | |X| = n\}|.$
- 2. $\aleph_0 = |\{X \in P(\mathbb{N}) | |X| = |\mathbb{N}|\}|.$

1.3 Problem 3

Find the error in the following:

We shall define simmetric set difference as follows: let k and m be cardinalities, not necessarily distinct. Define $k \oplus m$ to mean A and B sets for which |A| = k and |B| = m, then the cardinality of $A \oplus B$ is the value of $k \oplus m$.

1.3.1 Answer 7

Since we are trying to define a binary operation, we must make sure that the result (or value) of this operation is uniquely defined. But it is easy to see it is not the case, since if $|A \cap B| \neq |A' \cap B'|$ while at the same time |A| = |A'| and |B| = |B'|, we will get that $k \oplus m \neq k \oplus m$. Thus our attempt at defining a ring-sum operation will fail.

1.4 Problem 4

- 1. Let k_1,k_2,m_1,m_2 be cardinalities. Prove that if $k_1 \leq k_2$ and $m_1 \leq m_2$ then $k_1 \times m_1 \leq k_2 \times m_2$.
- 2. Prove that $\mathfrak{c} \times \aleph_0 = \mathfrak{c}$.
- 3. Prove that $\mathfrak{c}^{\mathfrak{c}} = 2^{\mathfrak{c}}$.

1.4.1 Answer 8

Recall that the product of cardinalities is defined to be the cardinality of cartesian product. Also, recall that $|A| \leq |B|$, implies having an injective function from A to B. Now we can combine these two facts to build the proof. Suppose for contradiction that it was possible that $k_1 \times m_1 > k_2 \times m_2$. This would mean that for some sets A and B the cartesian product with cardinality $k_1 \times m_1$ defined by $A \times B$ contains a pair (a, b) with a property that for the chosen injective function from A to A' a is not the source of any element in A'. But we are given that there is an injective function from A' to A (because $k_1 \leq k_2$). The argument for b is identical. Recall that the function must be defined for each element in its domain, but this stands in contradiction to the earlier claim that a (or b) is not in the domain of this function. Since this is not possible, the proof is complete.

1.4.2 Answer 9

To prove $\mathfrak{c} \times \aleph_0 = \mathfrak{c}$ we can use the fact that there are no cardinal numbers between \mathfrak{c} and \aleph_2 . Also relying on the "sandwich" theorem, which says that if we can find a superset with cardinality n and a subset with the cardinality n, then the cardinality of the set "sandwiched" in between must be n. Consider $\mathfrak{c} \leq \mathfrak{c} \times \aleph_0 \leq \mathfrak{c} \times \mathfrak{c}$. We can show that $\mathfrak{c} \times \aleph_0$ is at least as big as \mathfrak{c} since \aleph_0 is defined to be the multiplicative identity. We can also show that $\mathfrak{c} \times \mathfrak{c} = \mathfrak{c}$ by perusing Cantor's theorem which proves that the cartesian product of an infinite set with itself has the same cardinality as the initial set. Thus $\mathfrak{c} \times \aleph_0 = \mathfrak{c}$.

1.4.3 Answer 10

Observe that $\mathfrak{c}^{\mathfrak{c}}$ is the cardinality of the set of all functions from a set of cardinality \mathfrak{c} to itself, while $2^{\mathfrak{c}}$ is the cardinality of the power-set of cardinality \mathfrak{c} . We know (from Cantor's diagonal argument) that |X| < |P(X)| for all sets, this means that in particular $2^{\mathfrak{c}} < 2^{2^{\mathfrak{c}}}$. Since we know that cardinality of exponentiation is non-decreasing in both arguments it follows that $\mathfrak{c}^{\mathfrak{c}} < 2^{2^{\mathfrak{c}}}$. And because there are no cardinal numbers between X and P(X) we can conclude that $\mathfrak{c}^{\mathfrak{c}} = 2^{\mathfrak{c}}$.