

# Assignment 11, Infinitesimal Calculus

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# 1 Problems

## 1.1 Problem 1

1. Prove that for any natural number  $n$  it holds that:

$$4^n \geq \binom{2n}{n}.$$

2. Prove by induction, or in any other way, that for any natural number  $n$  it holds that:

$$\binom{2n}{n} \geq \frac{4^n}{2n+1}.$$

### 1.1.1 Answer 1

**Discussion:** The idea behind the proof is to show that  $\binom{2n}{n}$  is a term in the binomial expansion representing  $4^n$ . Since all other terms in this expansion will be non-negative, then  $4^n$  will be at least as big as  $\binom{2n}{n}$ .

**Proof:** Put  $4^n = (1+1)^{2n}$ , using binomial formula obtains:

$$\begin{aligned} (1+1)^{2n} &= \sum_{k=0}^{2n} \binom{2n}{k} 1^{2n} 1^{2n-k} \\ &= \binom{2n}{n} + \sum_{k=0, k \neq n}^{2n} \binom{2n}{k}. \end{aligned}$$

We know that  $\binom{2n}{n}$  is a term of binomial expansion because we know that  $k_i \leq 2n$ , which implies that since  $k_i, n$  are natural numbers, there exists  $k_i = n$ . Besides, there might exist other terms in binomial expansion which are guaranteed to be non-negative. Hence,  $4^n \geq \binom{2n}{n}$ .

### 1.1.2 Answer 2

**Discussion:** In order to make the proof a little less verbose, I will prove a stronger claim, viz.  $\binom{2n}{n} \geq \frac{4^n}{2n}$ . Since  $n$  is positive,  $2n < 2n + 1$ , hence  $\frac{4^n}{2n} > \frac{4^n}{2n+1}$ . The proof will proceed by induction on  $n$ . First I will find a factor s.t. multiplying it with  $S_{n-1}$  I will obtain  $S_n$ , and then multiply it with the  $\frac{4^n}{2n}$  to show that it will necessary be at least as large as  $\frac{4^{n+1}}{2(n+1)}$ .

**Proof:** Using mathematical induction, let's first prove the base step, where  $n = 1$ :

$$\begin{aligned} \binom{2 * 1}{1} &\geq \frac{4^1}{2 * 1} \iff \\ \frac{2!}{1!(2-1)!} &\geq \frac{4}{2} \iff \\ \frac{2}{1} &\geq 2 \iff \\ 2 &\geq 2. \end{aligned}$$

Now, to the inductive step (for  $n > 1$ ), some useful simplification first:

$$\begin{aligned} \binom{2n}{n} &= \frac{(2n)!}{n!(2n-n)!} \\ &= \frac{(2n)!}{n!n!}. \end{aligned} \tag{1}$$

Invoking inductive hypothesis  $\binom{2(n+1)}{(n+1)} \geq \frac{4^{n+1}}{2(n+1)}$ :

$$\begin{aligned} \binom{2(n+1)}{n+1} &= \frac{(2(n+1))!}{(n+1)!(2(n+1)-(n+1))!} \\ &= \frac{(2n!)(2n+1)(2n+2)}{n!(n+1)n!(n+1)} \\ &= \frac{(2n!)(2n+1)2(n+1)}{n!(n+1)n!(n+1)} \\ &= \frac{(2n!)(2n+1)2}{n!n!(n+1)}. \end{aligned} \tag{2}$$

Dividing **2** by **1** gives us the factor  $\frac{(2n+1)2}{(n+1)}$ . Thus:

$$\begin{aligned}\frac{(2n+1)2}{n+1} \times \frac{4^n}{2n} &\geq \frac{4^{n+1}}{2(n+1)} \\ \frac{(2n+1)2 * 4^n}{2n(n+1)} &\geq \frac{4^{n+1}}{2(n+1)} \\ \frac{(2n+1)2 * 4^n}{n} &\geq 4^{n+1} \\ \frac{(2n+1)2 * 4^n}{n} &\geq 4^n * 4 \\ \frac{(2n+1)2}{n} &\geq 4 \\ \frac{4n+2}{n} &\geq 4 \\ 4 + \frac{2}{n} &\geq 4\end{aligned}$$

Since  $n > 1$ ,  $\frac{2}{n}$  is positive, hence the inequality holds. This completes the inductive step. Hence, by using mathematical induction the proof is complete.

## 1.2 Problem 2

1. Given  $k, l \in \mathbb{N}$ , prove that  $a = k + l\sqrt{2}$  is irrational.
2. Prove that for every natural number  $n$  it holds that:

$$\sum_{i=0}^n \sqrt{2}^i$$

is irrational.

### 1.2.1 Answer 3

**Discussion:** One way to see that summation of rational with irrational cannot produce a rational number is through invoking field axioms: sum-

mation must send the sum to the field of rationals, which would imply that the inverses of summands must be rationals too. This would also require summands to be rationals, but that's not possible.

**Proof:** Suppose, for contradiction that  $k + l\sqrt{2} = a$  is rational, then

$$\begin{aligned}
 k + l\sqrt{2} &= a && \text{Given} \\
 k + (-k) + l\sqrt{2} &= a - k && k \text{ must have additive inverse in } \mathbb{Q} \\
 0 + l\sqrt{2} &= a - k \\
 l^{-1}l\sqrt{2} &= l^{-1}(a - k) && l \text{ must have multiplicative inverse in } \mathbb{Q} \\
 1\sqrt{2} &= l^{-1}(a - k) \\
 \sqrt{2} &= l^{-1}(a - k).
 \end{aligned}$$

$l^{-1}$ ,  $q$  and  $k$  are all rationals, rationals are closed under multiplication and addition, hence  $l^{-1}(q - k)$  must be rational, but  $\sqrt{2}$  is not. Contradiction. Hence  $a = k + l\sqrt{2}$  is irrational.

### 1.2.2 Answer 4

**Discussion:** The way to see that this statement is true is to divide the sequence into odd and even terms. All even terms will produce rationals (even poverse of square root of two will be rational). While all odd terms will produce irrational numbers (a product of even number of square roots of two will give a rational, but them multiplied with an irrational number will give an irrational). Since  $n$  must be at least one (and thus we are guaranteed to have at least one odd term in this sequence), the sum of the sequence will always be irrational.

**Proof:** Let's rewrite this sum as two sums of the form:

$$\sum_{i=0}^{\lfloor n/2 \rfloor} \sqrt{2}^{2i} + \sum_{i=0}^{\lfloor (n+1)/2 \rfloor} \sqrt{2}^{2i+1}.$$

It is easy to see that the first term is a sum of powers of 2, viz:  $\sqrt{2}^2 +$

$\sqrt{2}^4 + \dots + \sqrt{2}^{\lfloor n/2 \rfloor}$ , which is just  $2 + 4 + \dots + 2^{\lfloor n/2 - 1 \rfloor}$ . Similarly, the terms of the other sum can be expressed as  $\sqrt{2}^1 + \sqrt{2}^3 + \dots + \sqrt{2}^{\lfloor (n+1)/2 \rfloor}$ .

Let's give names to the sequences we outlined:  $S_1 = \sum_{i=0}^{\lfloor n/2 \rfloor} \sqrt{2}^{2i}$  and  $S_2 = \sum_{i=0}^{\lfloor (n+1)/2 \rfloor} \sqrt{2}^{2i+1}$ . Now, for contradiction, assume  $S_2$  to be rational. Then let's define  $\div$  to be the element-wise division of sequences, i.e. for sequences  $A$  and  $B$  of length  $n$ ,  $\div$  is defined to be:

$$\sum_{i=0}^n \frac{A_i}{B_i}.$$

It's easy to see that if all elements of  $A$  and  $B$  are rational, then  $A \div B$  is rational too (because we only used addition and multiplication, which are known to be closed over rationals).

Observe that  $S_2 \div S_3 = \sum_{i=0}^{\lfloor (n+1)/2 \rfloor} \sqrt{2}$ , i.e. a sum of  $n/2$  square roots of 2, which is the same as  $n/2 * \sqrt{2}$ , but we've just showed that the product of a rational and irrational cannot be rational (in previous question). Hence  $S_2$  must be irrational, contrary to assumed. Hence it must be that  $\sum_{i=0}^n \sqrt{2}^i$  is irrational. This completes the proof.

### 1.3 Problem 3

1. Given real numbers  $a$  and  $b$  prove that if

$$\frac{|a|}{2} > \left| b - \frac{a}{2} \right|,$$

then

$$|b - a| < |a|.$$

#### 1.3.1 Answer 5

**Discussion:** The proof will be based on invariance of order under multiplication, in particular, it will rely on the fact that  $|x| < |y| \iff x^2 < y^2$ . This will allow us to solve the inequality without splitting it into several cases.

**Proof:** First, let's simplify the expression:

$$\begin{aligned}
\frac{|a|}{2} &> \left| b - \frac{a}{2} \right| && \text{Given} \\
|a| &> 2 \left| b - \frac{a}{2} \right| && \text{Invariance of order under multiplication} \\
|a| &> |2b - a| && \text{Distributivity of multiplication over addition} \\
|2b - a| &< |a| \\
|2b - a|^2 &< |a|^2 && \text{Invariance under exponentiation} \\
(2b - a)^2 &< a^2 && \text{Squares are always positive} \\
4b^2 - 4ba + a^2 &< a^2 \\
4b^2 - 4ba &< 0 \\
4b^2 &< 4ba \\
b^2 &< ba && \text{Invariance under multiplication}
\end{aligned}$$

Next, let's perform similar operations on  $|b - a| < |a|$

$$\begin{aligned}
|b - a| &< |a| && \text{Given} \\
|b - a|^2 &< |a|^2 && \text{Invariance under exponentiation} \\
(b - a)^2 &< a^2 && \text{Squares are always positive} \\
b^2 - 2ba + a^2 &< a^2 \\
b^2 - 2ba &< 0 \\
b^2 &< 2ba.
\end{aligned}$$

From transitivity of order it follows that  $b^2 < ba < 2ba$ , i.e.  $b^2 < 2ba$ , but this is exactly the condition we set out to prove in the very beginning. Thus, the proof is complete.

## 1.4 Problem 4

Given  $a, b, c \in \mathbb{R}$ ,

1. Prove that if  $a > 0$  and  $a + b > a + c$ , then  $b > c$ .



2. Prove that if  $a > 0$  and  $ab > ac$ , then  $b > c$ .
3. Prove that  $|a| > |b|$  iff  $a^2 > b^2$ .
4. Prove that if  $b > c$  and  $|a - b| > |a - c|$ , then  $b > a$ .
5. Show (my means of example) that from  $b > c$  and  $b > a$  it doesn't follow that  $|a - b| > |a - c|$ .

#### 1.4.1 Answer 6

**Proof:** The proof is simple algebra relying on invariance of order under addition:

$$\begin{array}{ll}
 a + b > a + c & \text{Given} \\
 -a + a + b > -a + a + c & \text{Invariance under addition} \\
 0 + b > 0 + c & \text{Devinition of inverse} \\
 b > c & \text{Devinition of inverse}
 \end{array}$$

#### 1.4.2 Answer 7

**Proof:** The proof is simple algebra relying on invariance of order under multiplication:

$$\begin{array}{ll}
 ab > ac & \text{Given} \\
 a^{-1}ab > a^{-1}ac & \text{Invariance under multiplication} \\
 1b > 1c & \text{Devinition of inverse} \\
 b > c & \text{Devinition of inverse}
 \end{array}$$

### 1.4.3 Answer 8

**Proof:** first I will prove the implies part, i.e.  $|a| > |b| \implies a^2 > b^2$ , and then the  $|a| > |b| \iff a^2 > b^2$ :

$$\begin{array}{ll}
 |a| > |b| & \text{Given} \\
 |a||a| > |a||b| & \text{Invariance under multiplication} \\
 a^2 > |ab| & \text{Simple algebra} \\
 |a| > |b| & \text{Reiteration of the given} \\
 |a||b| > |b||b| & \text{Invariance under multiplication} \\
 |ab| > b^2 & \text{Simple algebra} \\
 a^2 > |ab| > b^2 & \text{Reusing derivation from step 3} \\
 a^2 > b^2 & \text{By transitivity of order}
 \end{array}$$

The converse proof:

$$\begin{array}{ll}
 a^2 > b^2 |a||a| > |b||b| & \text{Simple algebra} \\
 |a||a| - |b||b| > 0 & \text{Simple algebra} \\
 |a|(|b| + c) - |b||b| > 0 & \text{Define } c = |a| - |b| \\
 |a||b| + |a|c - |b||b| > 0 & \\
 (|a| - |b|)|b| + |a|c > 0 & \\
 (|a| - |b|)|b| > -|a|c & \\
 (|a| - |b|)|b| > -|a|(|a| - |b|) & \text{Recall } c = |a| - |b| \\
 |b| > -|a| & \text{Invariance under multiplication} \\
 -|a| < |b| & \\
 |a| > |b| & \text{Multiplying by } -1
 \end{array}$$

Both the “if” part and its converse have been proved, thus the proof is complete.

### 1.4.4 Answer 9

**Proof:** For the sake of diversity, this proof will rely on axioms rather than algebraic manipulations. Perusing the first axiom of order, we know that

either one of the three holds:  $c < a$ ,  $c = a$  or  $c > a$ . Consider  $c < a$ . We are given that the distance from  $a$  to  $b$  is greater than the distance from  $a$  to  $c$ . This means that  $b$  cannot lie between  $a$  and  $c$  (recall the triangle inequality), neither can it lie between  $c$  and  $a$  if  $a = c$ , hence  $b > a$ .

Consider now  $c > a$ , we are given that  $b > c$ , thus, by transitivity of order we conclude that  $b > c > a$  and in particular  $b > a$ . This completes the proof.

#### 1.4.5 Answer 10

Any combination of  $a, b$  and  $c$  s.t.  $c > a > b$  and  $a - c > b - a$  will satisfy the requirement, and in particular  $c = 1, a = 3, b = 4$  as is easy to see:  $|a - b| > |a - c| \implies |3 - 4| > |3 - 1| \implies 1 > 2$ . A way to see why this is true is to picture the number line, where the distance is measured from the central point  $a$  to two points on the left and on the right of it. We aren't constrained in any way as to which side should be bigger, thus we can certainly obtain the one that fits the requirement.

### 1.5 Problem 5

Solve the equation:

$$\lfloor |x + 1| - |x - 1| \rfloor = x.$$

#### 1.5.1 Answer 11

**Discussion:** First, observe that flooring on the left side guarantees that  $x$  is an integer. Having said that, we can drop the flooring operation altogether. The solution will examine three cases of  $x < 0$ ,  $x = 0$  and  $x > 0$  warranted by the first axiom of order.

**Case  $x = 0$ :**

$$\begin{aligned}|x + 1| - |x - 1| - x &= 0 \\0 + 1 - 0 - 1 - 0 &= 0 \\0 &= 0\end{aligned}$$

Substituting 0 back into original formula proves to give a correct identity.

**Case  $x > 0$ :**

$$\begin{aligned}|x + 1| - |x - 1| - x &= 0 \\x + 1 - |x - 1| - x &= 0 \\1 - |x - 1| &= 0 \\1 - x + 1 &= 0 \\x &= 2\end{aligned}$$

Substituting 2 back into equation gives  $|2 + 1| - |2 - 1| = 3 - 1 = 2$ , this we obtained additional solution.

**Case  $x < 0$ :**

$$\begin{aligned}|x + 1| - |x - 1| - x &= 0 \\x + 1 - |x - 1| - x &= 0 \\1 - |x - 1| &= 0 \\1 - x + 1 &= 0 \\x &= 2\end{aligned}$$

But we started with the assumption that  $x < 0$ , thus there are no solutions for  $x < 0$ . Hence the only solutions are 0 and 2.

## 1.6 Problem 6

**Definition:** set  $A$  of real numbers is called **dense in interval  $I$**  if for every  $x, y \in I$  s.t.  $x < y$  there exists  $a \in A$  such that  $x < a < y$ .

1. Let  $A$  be dense in interval  $[0, 1]$ , prove that set  $B = \{na | a \in A, n \in \mathbb{N}\}$  is dense in interval  $[0, \infty)$ .

2. Let  $A = \mathbb{R}$ , prove that  $A$  isn't dense in  $I$  iff exists an open interval  $(x, y)$  in  $I$ , such that  $A \cap (x, y) = \emptyset$ .
3. Let  $A$  be the real numbers in interval  $[0, 1]$ , prove that the set  $C = \{\frac{a+1}{n^2} | a \in A, n \in \mathbb{N}\}$  isn't dense in  $[0, 1]$ .

**1.6.1    Answer 12**

**1.6.2    Answer 13**

**1.6.3    Answer 14**