알고리즘

01. Divide & Conquer

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미디어소프트웨어학과 민경하

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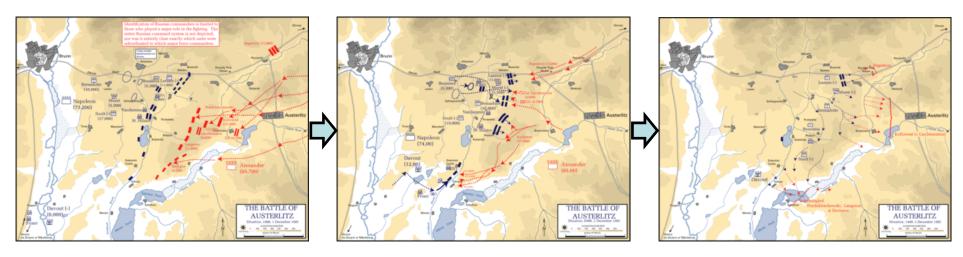
1. Divide & Conquer

- 1.0 Introduction
- 1.1 Multiplication
- 1.2 Recurrence relation
- 1.3 Sorting
- 1.4 Medians
- 1.5 Matrix multiplication
- 1.6 Finding closest pair

(1) Battle of Austerlitz



- (1) Battle of Austerlitz (1805)
 - Battle of three emperors
 - France (65,000) VS Austria + Russia (80,000)
 - The end of the 3rd anti-France aliance





- (2) League VS Tournament
 - Ex) Elite-8 of Worldcup 2014
 - BRA, COL, FRA, GER, NED, CRC, ARG, BEL
 - Champion by league
 - How many games do they play?
 - Champion by tournament
 - How many games do they play?

(2) League VS Tournament



(2) League VS Tournament

```
int Champion8 ( {BRA, COL, FRA, GER, NED, CRC, ARG, BEL} )
{
   Lwinner = Champion4 ( {BRA, COL, FRA, GER} );
   Rwinner = Champion4 ( {NED, CRC, ARG, BEL} );

   return Winner ( Lwinner, Rwinner );
}
```

(2) League VS Tournament

```
int Champion4 ( {BRA, COL, FRA, GER} )
{
   Lwinner = Champion2 ( {BRA, COL} );
   Rwinner = Champion2 ( {FRA, GER} );

   return Winner ( Lwinner, Rwinner );
}
```

```
int Champion4 ( {NED, CRC, ARG, BEL} )
{
   Lwinner = Champion2 ( {NED, CRC} );
   Rwinner = Champion2 ( {ARG, BEL} );

   return Winner ( Lwinner, Rwinner );
}
```

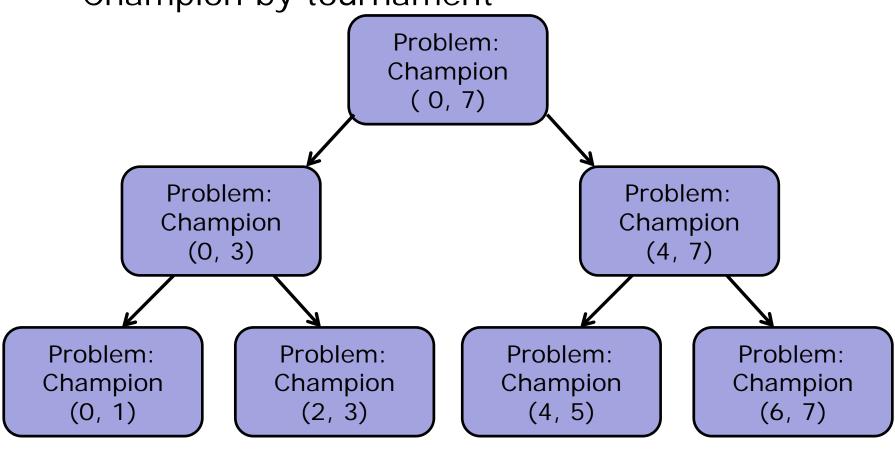
(2) League VS Tournament

```
int Champion2 ( {BRA, COL} )
{
   Lwinner = Champion1 ( {BRA} );  // Unnecessary
   Rwinner = Champion1 ( {COL} );  // Unnecessary
   return Winner ( BRA, COL);
}
```

```
int Champion2 ( {BRA, COL} )
{
    Lwinner = Champion1 ( {BRA} );  // Unnecessary
    Rwinner = Champion1 ( {COL} );  // Unnecessary
    return Winner ( BRA, COL);
}
```

(2) League VS Tournament

(2) League VS Tournament



(2) League VS Tournament

Champion by tournament

– Do we miss something?

(2) League VS Tournament

Champion by tournament

– Do we miss something? → degenerate case

(2) League VS Tournament

- (3) Key idea of divide & conquer
 - Solve a problem of n inputs by splitting the input into k subsets
 - Three steps of divide & conquer
 - Divide
 - Breaking a problem into subproblems
 - Conquer
 - Recursively solving these subproblems
 - Combine (optional)
 - Appropriately combining their answers

(4) Abstract algorithm for DnC (recursive)

```
global n, A(1:n);
DnC ( int p, int q )
{
   int m;
   if ( SMALL (p, q) )
      return G (p, q);
   else
      m ← DIVIDE (p, q);
      return COMBINE ( DnC (p, m), DnC (m+1, q) );
}
```

 Most divide & conquer algorithms are implemented using recursive call

- (5) Three check points
 - Same format
 - Reduced problem size
 - Degenerate case

Same format

```
global n, A(1:n);
DnC ( int p, int q )
{
    int m;
    if ( SMALL (p, q) )
        return G (p, q);
    else
        m 	DIVIDE (p, q);
        return COMBINE ( DnC (p, m), DnC (m+1, q) );
}
```

Reduced problem size

```
global n, A(1:n);
DnC ( int p, int q )
{
    int m;
    if ( SMALL (p, q) )
        return G (p, q);
    else
        m 	DIVIDE (p, q);
        return COMBINE ( DnC (p, m), DnC (m+1, q) );
}
```

Degenerate case

```
global n, A(1:n);
DnC ( int p, int q )
{
    int m;
    if ( SMALL (p, q) )
        return G (p, q);
    else
        m ← DIVIDE (p, q);
    return COMBINE ( DnC (p, m), DnC (m+1, q) );
}
```

(6) Performance analysis for DnC

- T (n): time complexity of DnC () for n inputs
- g (n): for small input
- $-f_1$ (n): for DIVIDE ()
- $-f_2$ (n): for COMBINE ()

$$T(n) = \begin{cases} g(n), & \text{for small } n \\ 2T(n/2) + f_1(n) + f_2(n), & \text{otherwise} \end{cases}$$

$$T(n) = \begin{cases} g(n), & \text{for small } n \\ aT(n/b) + O(n^d), & \text{otherwise} \end{cases}$$

(7) The simplest divide & conquer algorithm

- Binary search:
 - Let $A = \{a_1, ..., a_n\}$ be a list of elements which are sorted in nondecreasing order.
 - Determine whether x is in A or not. If x is in A, then find i such that $a_i = x$.

Straightforward search

```
int bsearch_straightforward ( int s, int e, int A[], int x )
{
    for ( i = s; i <= e; i++ ) {
        if ( A[i] == x )
            return i;
    }
    return NONE;
}</pre>
```

– Performance analysis?

Binary search (Divide & Conquer)

```
int bsearch (int s, int e, int A[], int x)
    if ( s == e )
        if (A[s] == x)
           return s;
        else
           return NONE;
    m \leftarrow (s+e)/2;
    if (A[m] == x)
       return m;
    else if (A[m] > x)
        return bsearch (s, m, A, x);
    else
       return bsearch ( m + 1, e, A, x );
```

Binary search -> check three key points

```
int bsearch (int s, int e, int A[], int x)
    if ( s == e )
        if (A[s] == x)
           return s;
        else
           return NONE;
    m \leftarrow (s+e)/2;
    if (A[m] == x)
        return m;
    else if (A[m] > x)
        return bsearch (s, m, A, x);
    else
        return bsearch ( m + 1, e, A, x );
```

- Binary search (Divide & Conquer)
 - Performance analysis

$$T(n) = \begin{cases} 1, & \text{for } n = 1 \\ T(n/2) + 1, & \text{otherwise} \end{cases}$$

$$T(n) = \log n$$