# 论题 1-5 作业

姓名: 陈劭源 学号: 161240004

# 1 [DH] Problem 2.10

Let T be a vector of Booleans.

- (1) for I going from 1 to N do the following:
  - (1.1)  $T[I] \leftarrow \text{false}$ ;
- (2) for I going from 1 to N do the following:
  - (2.1) if P[I] < 1 or P[I] > N do the following:
    - (2.2.1) output 'NO';
    - (2.2.2) end;
  - (2.2) T[P[I]] = true;
- (3) for I going from 1 to N do the following:
  - (3.1) if T[I] = false do the following:
    - (3.1.1) output 'NO';
    - (3.1.2) end;
- (4) output 'YES'.

# **2** [DH] Problem **2.11**

Let *K* be a vector of Booleans, *L* be a vector of integers which stores a permutation.

#### subroutine **produce permutation** *I*

- (1) if I = N do the following:
  - (1.1) output *R*;
  - (1.2) return;
- (2) for *i* going from 1 to *N* do the following:
  - (2.1) if K[i] is false do the following:
    - $(2.1.1) R[I] \leftarrow i$
    - $(2.1.2) K[i] \leftarrow \text{true};$
    - (2.1.3) call **produce permutation** I + 1
    - $(2.1.4) K[i] \leftarrow \text{false}.$

- (1) i going from 1 to N do the following: (1.1)  $K[i] \leftarrow$  false;
- (2) call **produce permutation** 0.

- (a) i. read(X), push(X,S), read(X), push(X,S), read(X), print(X), pop(X,S), print(X)
  - ii. read(X), push(X,S), read(X), push(X,S), read(X), print(X), pop(X,S), print(X), pop(X,S), print(X)
  - iii. read(X), push(X,S), read(X), push(X,S), read(X),  $\mathbf{print}(X)$ , read(X), push(X,S), read(X), print(X), read(X), push(X,S), read(X), print(X), pop(X,S), print(X), read(X), print(X), pop(X,S), print(X), read(X), print(X), pop(X,S), print(X), read(X), print(X), pop(X,S), print(X)
- (b) i. 要生成 (3, 1, 2) 这个排列,由于 3 是最先输出的,1,2 依次在栈中,此时若要继续输出,必然是以 2,1 的形式输出,所以不可能用栈生成 (3, 1, 2) 这个排列。□
  - ii. 要生成 (4, 5, 3, 7, 2, 1, 6) 这个排列, 当输出 7 时, 栈中剩余的元素依次为 1, 2, 6, 下一个需要输出 2, 但输出 2 之前 6 必须输出,从而不可能用栈生成 (4, 5, 3, 7, 2, 1, 6) 这个排列。
- (c) 容易验证,以下排列可以用栈生成:
  - (1, 2, 3, 4) (1, 2, 4, 3) (1, 3, 2, 4) (1, 3, 4, 2) (1, 4, 3, 2) (2, 1, 3, 4) (2, 1, 4, 3) (2, 3, 1, 4) (2, 3, 4, 1) (2, 4, 3, 1) (3, 2, 1, 4) (3, 2, 4, 1) (3, 4, 2, 1) 以下排列不能用栈生成:
  - (1, 4, 2, 3) (2, 4, 1, 3) (3, 1, 2, 4) (3, 1, 4, 2) (3, 4, 1, 2) (4, 1, 2, 3) (4, 1, 3, 2) (4, 2, 3, 1) (4, 2, 1, 3) (4, 3, 1, 2)

所以共有10个排列不能用栈生成。

## 4 [DH] Problem 2.13

Let S be an empty stack. Assume that the length of the permutation is N.

#### function test

- $(1) I \leftarrow 1;$
- (2) for i going from 1 to N do the following:
  - $(2.1) \, read(X);$
  - (2.2) if  $I \leq X$  do the following:

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(2.2.1) for j going from I to X do the following:
                 (2.2.1.1) push(j,S);
         (2.2.2) pop(X, S);
         (2.2.3) I \leftarrow X + 1;
   (2.3) otherwise do the following:
         (2.3.1) pop(Y,S);
         (2.3.2) if X \neq Y then return false;
(3) return true.
subroutine print operations
(1) I \leftarrow 1;
(2) for i going from 1 to N do the following:
    (2.1) read(X);
   (2.2) if I \leq X do the following:
         (2.2.1) for j going from I to X do the following:
                 (2.2.1.1) print("read(X)");
                 (2.2.1.2) print("push(X, S)");
                 (2.2.1.3) push(j,S);
         (2.2.2) print("pop(X,S)");
         (2.2.3) print("print(X)");
         (2.2.4) pop(X, S);
         (2.2.5) I \leftarrow X + 1;
    (2.3) otherwise do the following:
         (2.3.1) pop(Y, S);
         (2.3.2) print("pop(X,S)");
         (2.3.3) print("print(X)").
(1) if test is true then do the following:
    (1.1) print("Yes");
    (1.2) call print operations;
(2) otherwise do the following:
    (2.1) print("No").
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(a) i. Obtain by a queue: read(X), add(X,Q), read(X), add(X,Q), read(X),
 print(X), remove(X,Q), print(X), remove(X,Q), print(X)
 Obtain by two stacks: read(X), push(X,S), read(X), push(X,S'), read(X),

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print(X), pop(X,S), print(X), pop(X,S'), print(X)
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- ii. Obtain by a queue: read(X), add(X,Q), read(X), add(X,Q), read(X), print(X), read(X), add(X, O), read(X), print(X), remove(X,Q), add(X,Q), remove(X, Q), add(X, Q), remove(X, Q), print(X), read(X), add(X,Q), read(X), print(X), remove(X,Q), add(X,Q), remove(X,Q), print(X), remove(X, Q), add(X, Q), remove(X, Q), print(X),remove(X, Q), print(X)Obtain by two stacks: read(X), push(X,S), read(X), push(X,S), read(X), push(X,S), read(X), print(X), read(X), print(X), pop(X,S), print(X), read(X), push(X,S'), read(X), print(X), pop(X,S), print(X), pop(X,S), print(X), pop(X, S'), print(X)
- (b) 用以下方法可只用一个队列生成任一排列:对于某一个要输出的数字,如果不在队列中,则将该数字之前的数字全部入队,然后直接输出该数字;如果在队列中,反复将队首数字出队后再入队"翻找"整个队列,直到找到为止,输出这个数字。对于排列中的每一个数字,依次重复上述操作,即可用一个队列输出任意排列。□
- (c) 用以下方法可只用两个队列生成任一排列:对于某一个要输出的数字,如果不在任一栈中,则将该数字之前的数字全部压入任何一个栈,然后直接输出该数字;如果在某个栈中,将该数字之上的数字依次弹出并压入到另一个之中,然后输出这个数字。对于排列中的每一个数字,依次重复上述操作,即可用两个栈输出任意排列。

Let S, S' be two empty stacks. Assume that the length of the permutation is N.

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function top(S)
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- (1) **pop**(t, S);
- (2) **push**(t,S);
- (3) return *t*.
- $(1) I \leftarrow 1;$
- (2) for *i* going from 1 to *N* do the following:
  - (2.1)**read**(X);
  - (2.2) if  $I \leq X$  do the following:
    - (2.2.1) for *j* going from *I* to *X* do the following:
      - (2.2.1.1) **print**("**read**(X)");
      - (2.2.1.2) **print**("**push**(X, S)");
      - (2.2.1.3) **push**(j,S);
    - (2.2.2) print("pop(X,S)");
    - (2.2.3) print("print(X)");

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(2.2.4) pop(X, S);
     (2.2.5) I \leftarrow X + 1;
(2.3) otherwise do the following:
     (2.3.1) while is-empty(S') is false do the following:
             (2.3.1.1) print("pop(X, S')");
             (2.3.1.2) print("push(X, S)");
             (2.3.1.3) pop(X, S');
             (2.3.1.4) push(X,S);
     (2.3.2) while top(S) \neq I do the following:
             (2.3.2.1) print("pop(X, S)");
             (2.3.2.2) print("push(X, S')");
             (2.3.2.3) pop(X,S);
             (2.3.2.4) push(X, S');
     (2.3.3) pop(X, S);
     (2.3.4) print("pop(X,S)");
     (2.3.5) print("print(X)").
```

(a) Let T be an empty binary search tree, L be a list of integers, N be the number of integers in L.

#### subroutine add X to S

- (1) if *S* is empty then do the following:
  - (1.1)  $S \leftarrow X$ ;
  - (1.2) return;
- (2) if X < S then do the following:
  - (2.1) **add** *X* **to left**(*S*);
- (3) otherwise do the following:
  - (3.1) add X to right(S).
- (1) for i going from 1 to N do the following:
  - (1.1) add L[i] to T.
- (b) Let *T* be a binary search tree.
  - subroutine **visit** S
  - (1) if *S* is empty then return;
  - (2) **visit right**(*S*);

- (3) output *S*;
- (4) **visit left**(*S*).
- (1) **visit left**(T).