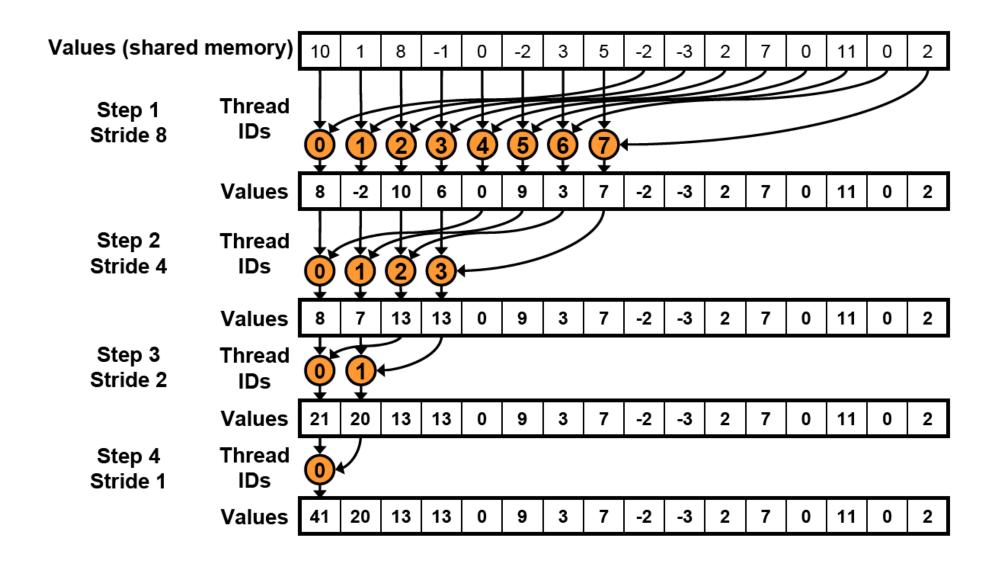


规约算法

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Reduction #3: Sequential Accesses



Eliminates bank conflicts

Reduction #3: Code Changes

Replace stride indexing in the inner loop:

```
// do reduction in shared mem
for (unsigned int s=1; s < blockDim.x; s *= 2) {
    int index = 2 * s * tid;

    if (index < blockDim.x == 0) {
        sdata[index] += sdata[index + s];
    }
    __syncthreads();
}</pre>
```

With reversed loop and threadID-based indexing:

```
// do reduction in shared mem
for (unsigned int s = blockDim.x/2; s > 0; s /= 2) {
      if (tid < s) {
            sdata[tid] += sdata[tid + s];
      }
      __syncthreads();
}</pre>
```

Performance for 4M element reduction

	Time (2 ²² ints)	Bandwidth	Step Speedup	Cumulative Speedup
Kernel 1: interleaved addressing with divergent branching	8.054 ms	2.083 GB/s		
Kernel 2: interleaved addressing non-divergent branching	3.456 ms	4.854 GB/s	2.33x	2.33x
Kernel 3: sequential addressing	1.722 ms	9.741 GB/s	2.01x	4.68x

Reduction #3: Bad resource utilization

- All threads read one element
- First step: half of the threads are idle
- Next step: another half becomes idle

```
// do reduction in shared mem
for (unsigned int s = blockDim.x/2; s > 0; s /= 2) {
      if (tid < s) {
            sdata[tid] += sdata[tid + s];
      }
      __syncthreads();
}</pre>
```



THANK YOU